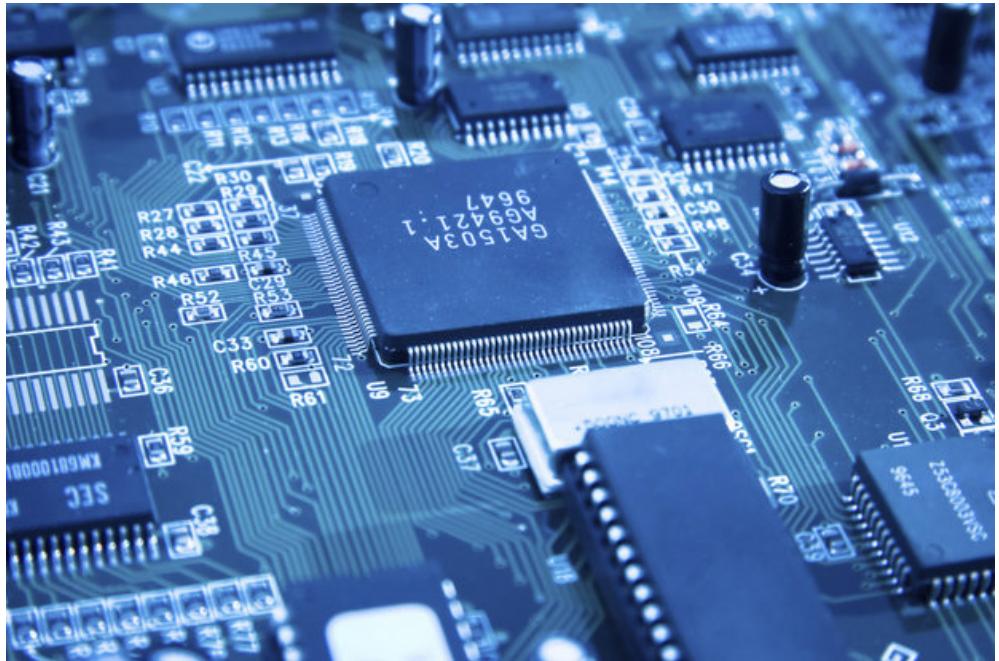


Lucas Fijen  
Falk Sinke [Chips and Circuits](#),  
Claartje Barkhof [Heuristieken, UvA 2017](#)

# **Oplossen van een constraint optimisation problem met een Heatmap gestuurde A\* implementatie**

# Chips and Circuits

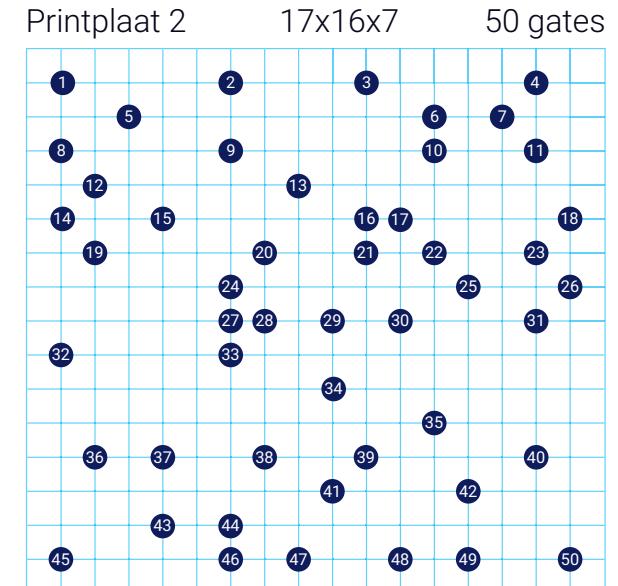
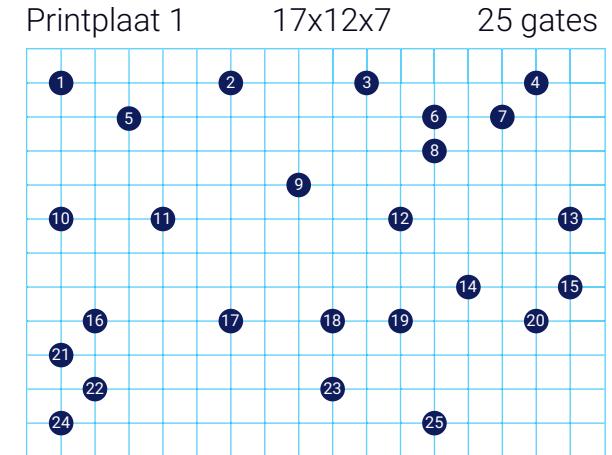
- Logische poorten (gates) verbinden voor werking
- Snijden = kortsluiting!
- Minder draad:
  - snellere chip
  - + lagere productiekosten per chip



# De casus

- Twee printplaten
- Netlist = voorschrift hoe gates verbonden moeten worden
- 3 netlists per printplaat (6 totaal)
- Manhattan grid + 7 lagen (3D)
- Niet snijden, niet buiten grid, niet door gates
- Constraint optimisation problem

Voorbeeld Netlist:  
(1, 11), (2, 8), (17, 6),  
(14, 20), (23, 25), (22, 13)



# Methode overview

1

Constraint Satisfaction

A\* & purpose-built Heatmap



2

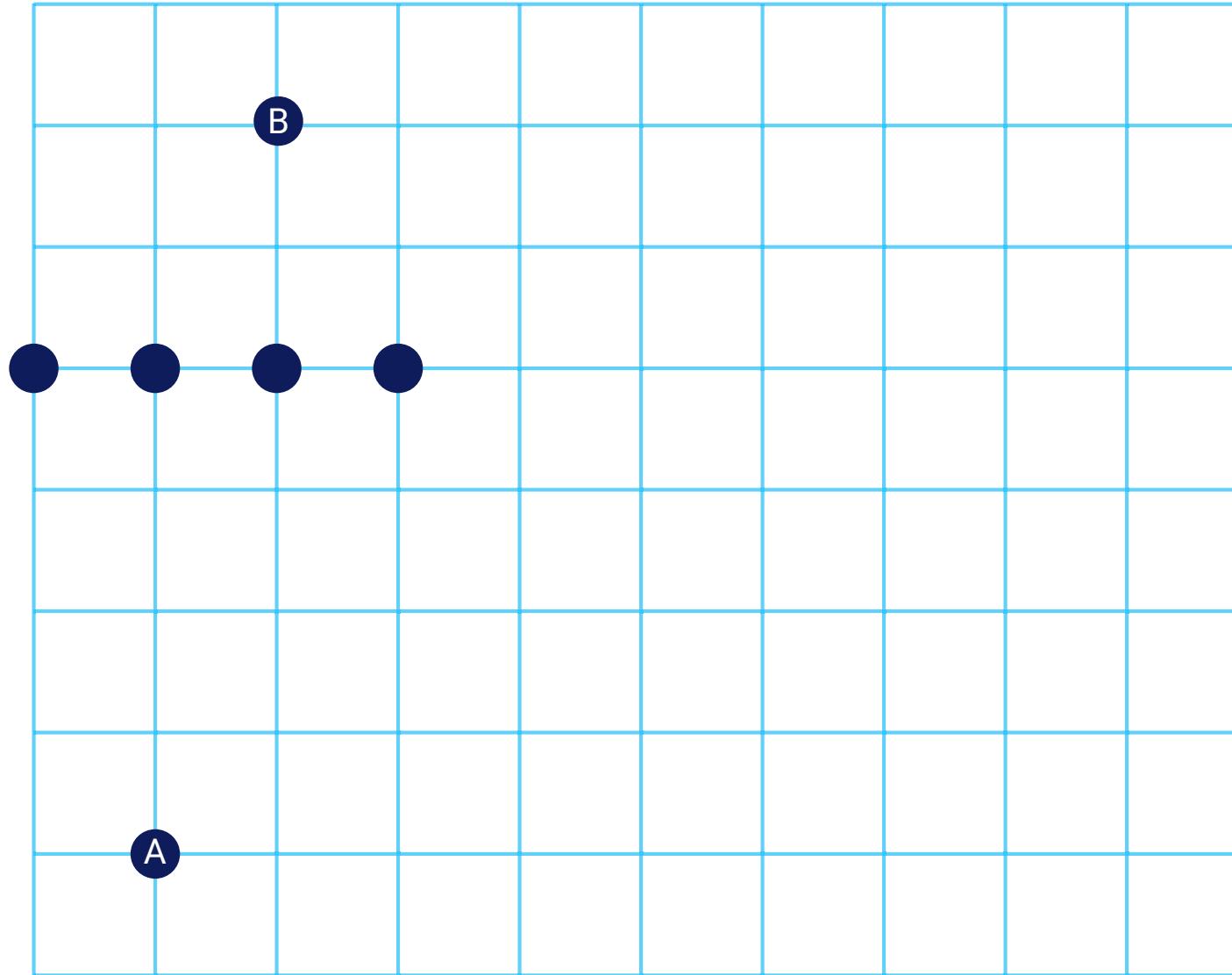
Optimisation

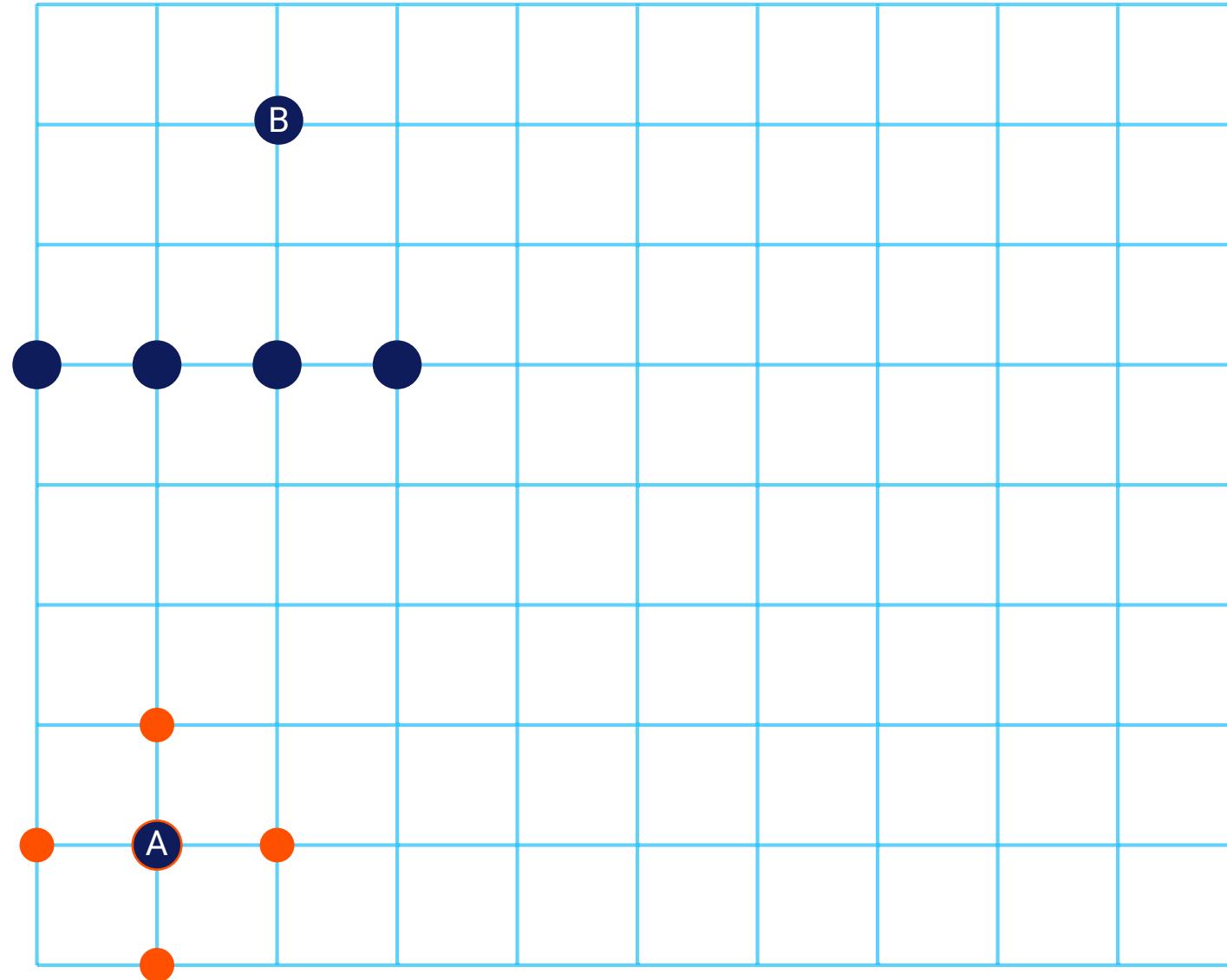
A\*

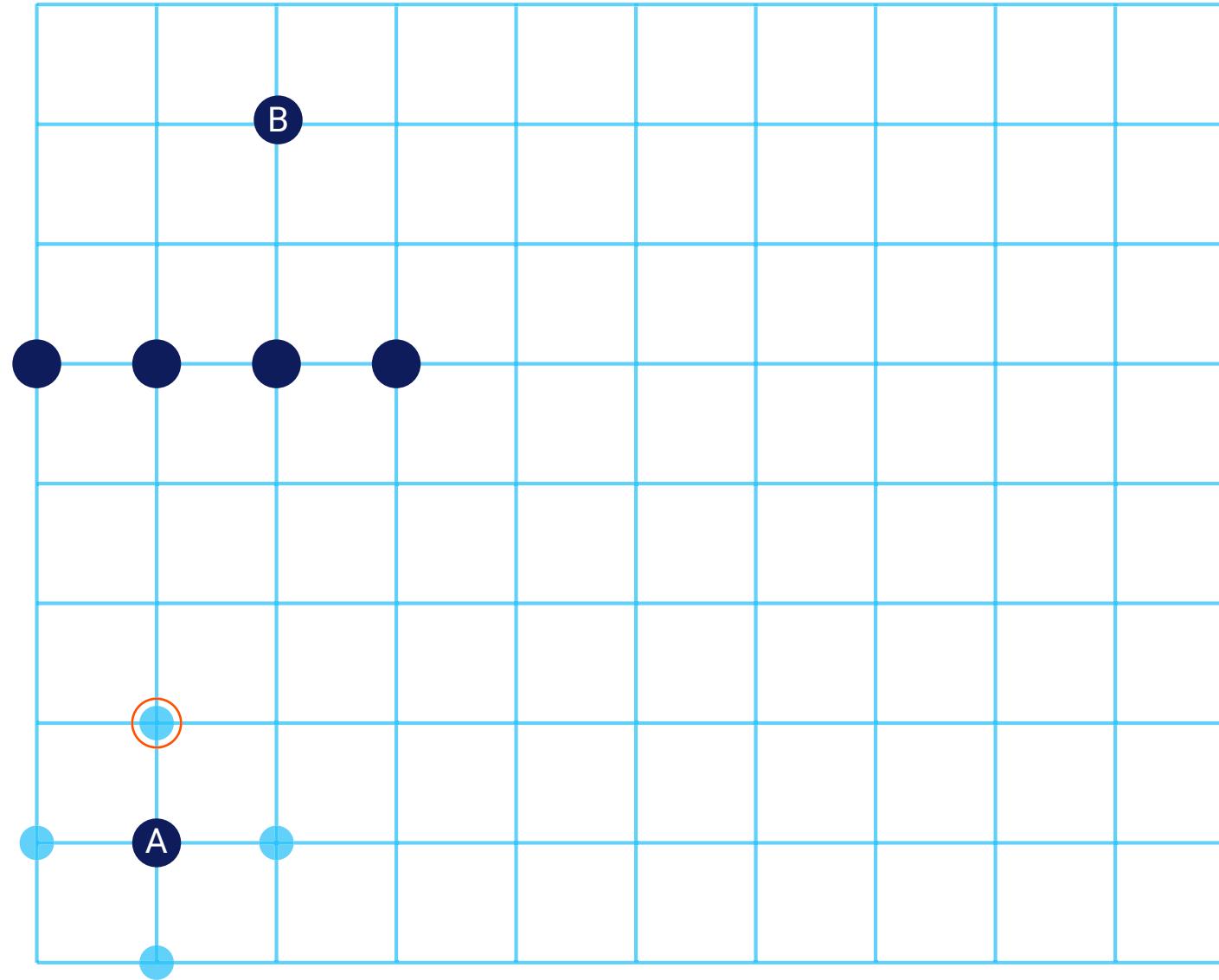
# A\* Zoek Algoritme

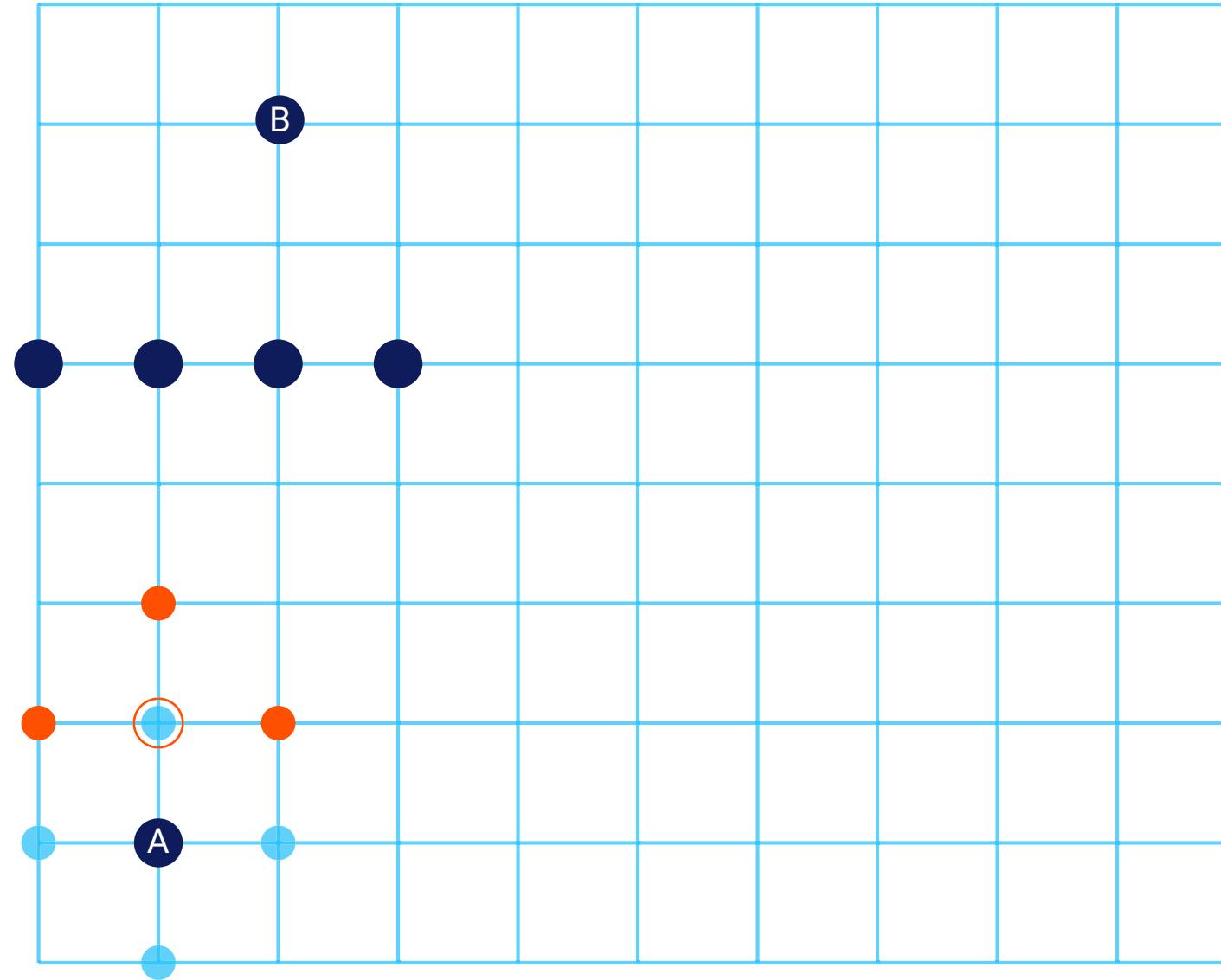
Zoek algoritme om het pad met de laagste kosten te vinden:

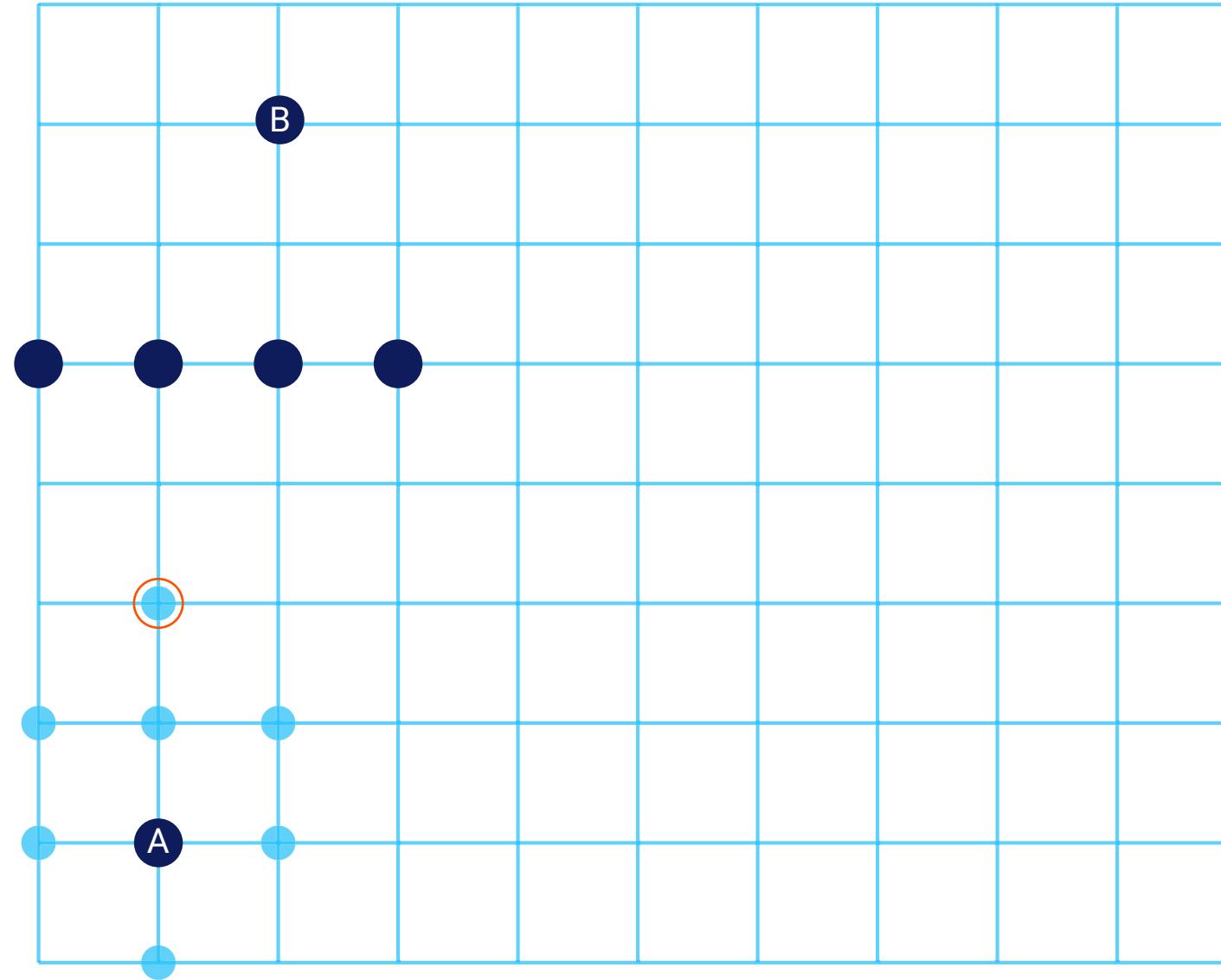
1. Huidige node -> evalueer vervolg states
2. Ken prijs toe aan vervolg states:  
*Afgelegd pad + Manhattan distance naar doel (heuristiek).*
3. Kies pad met laagste prijs uit lijst -> stap 1.

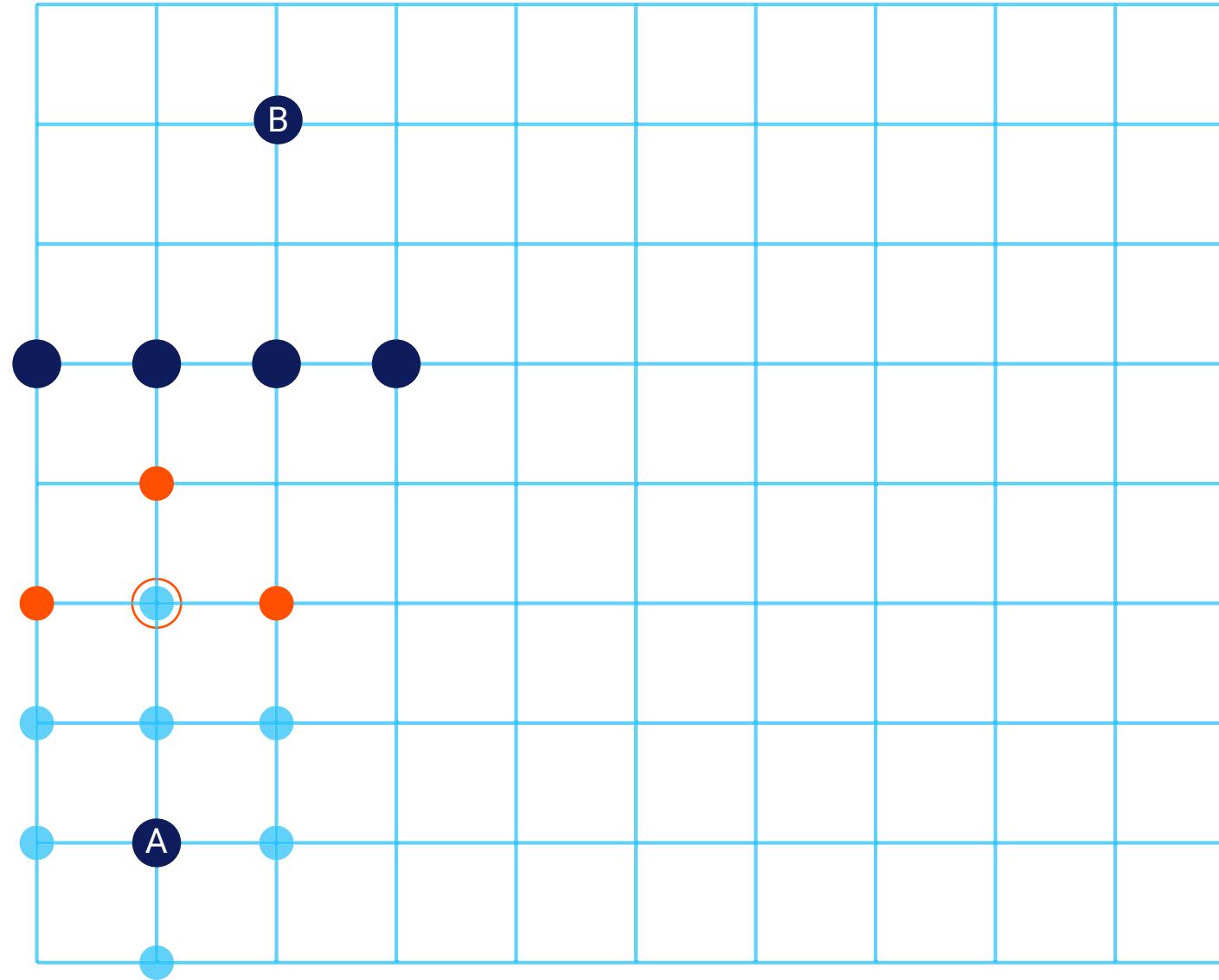


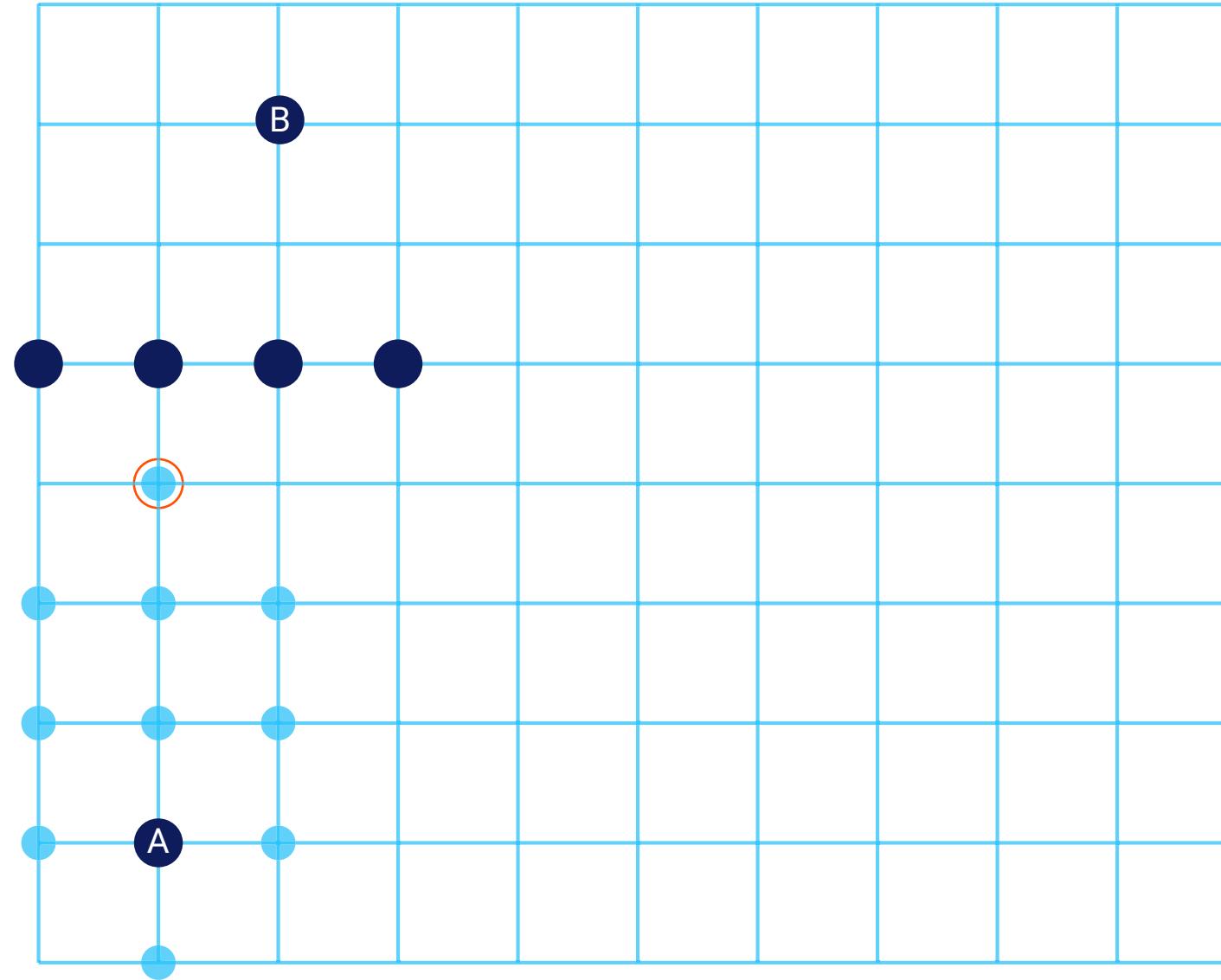


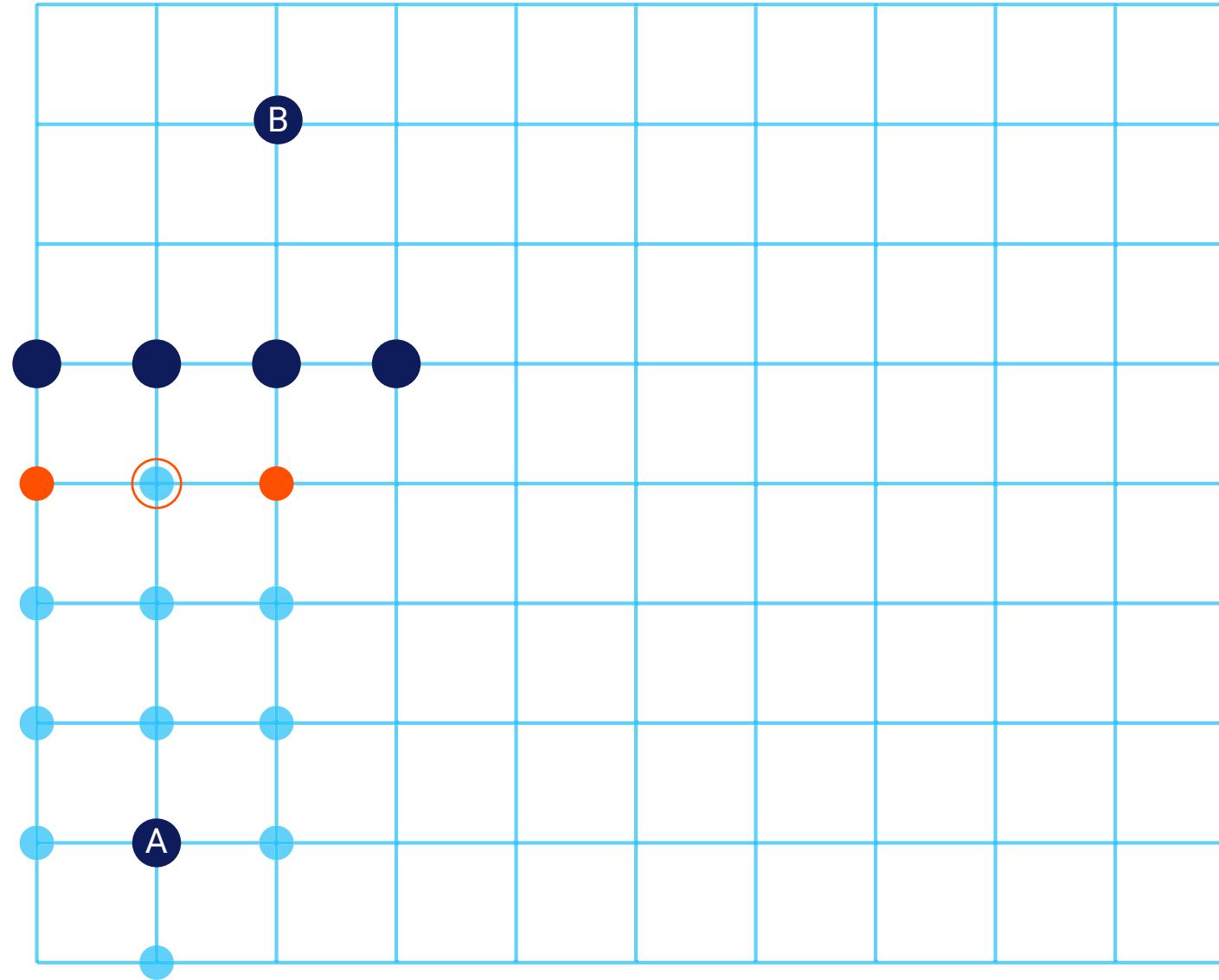


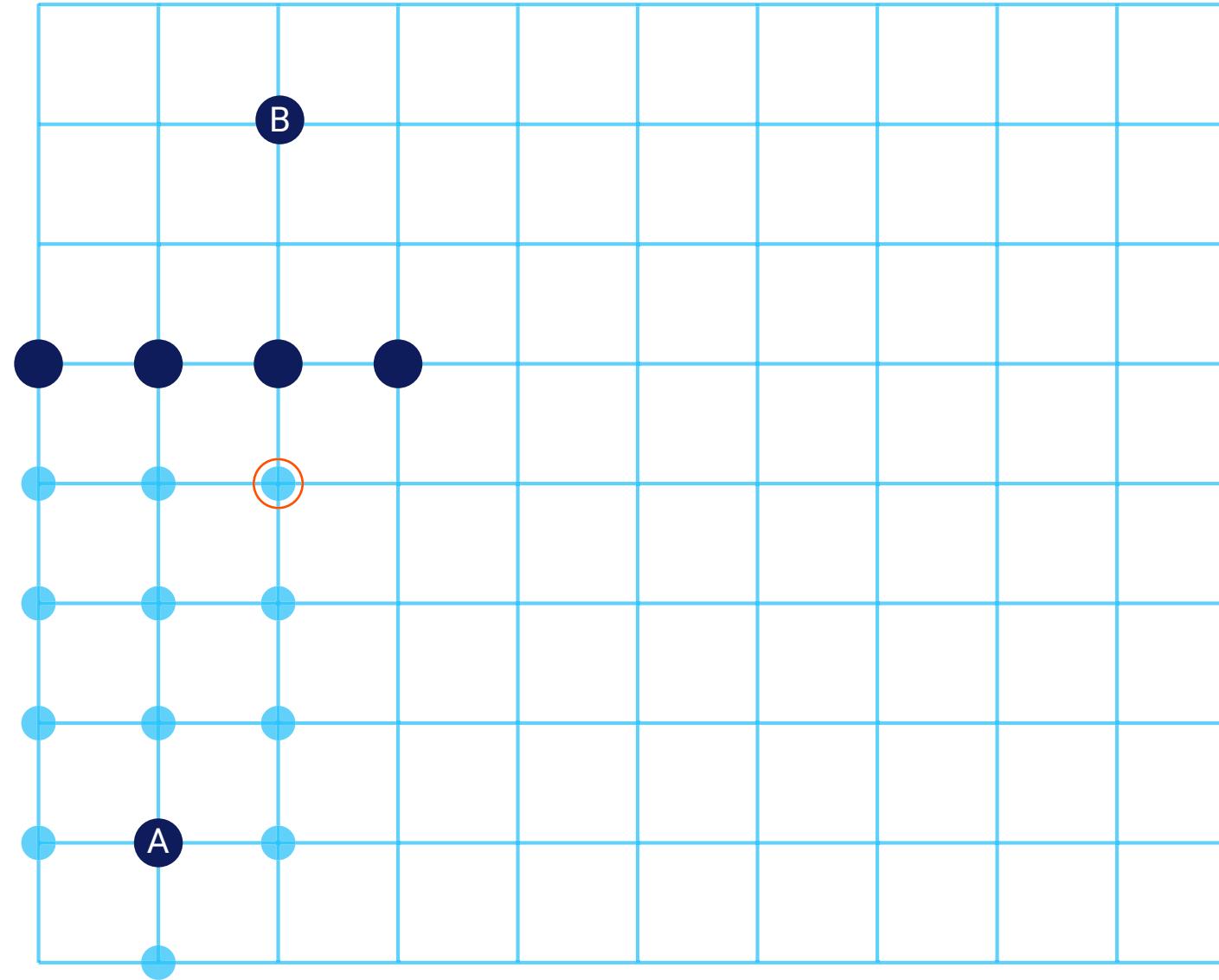


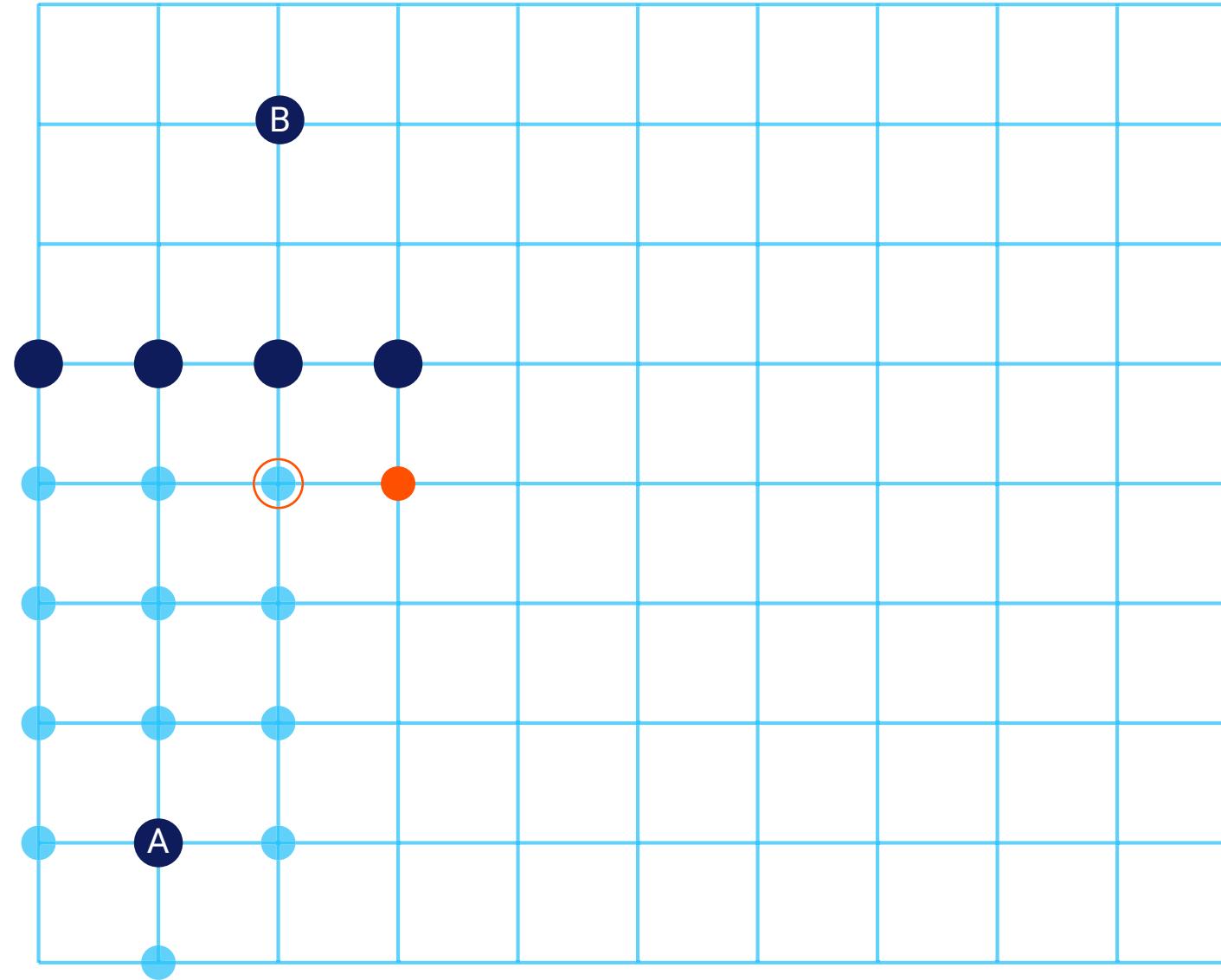


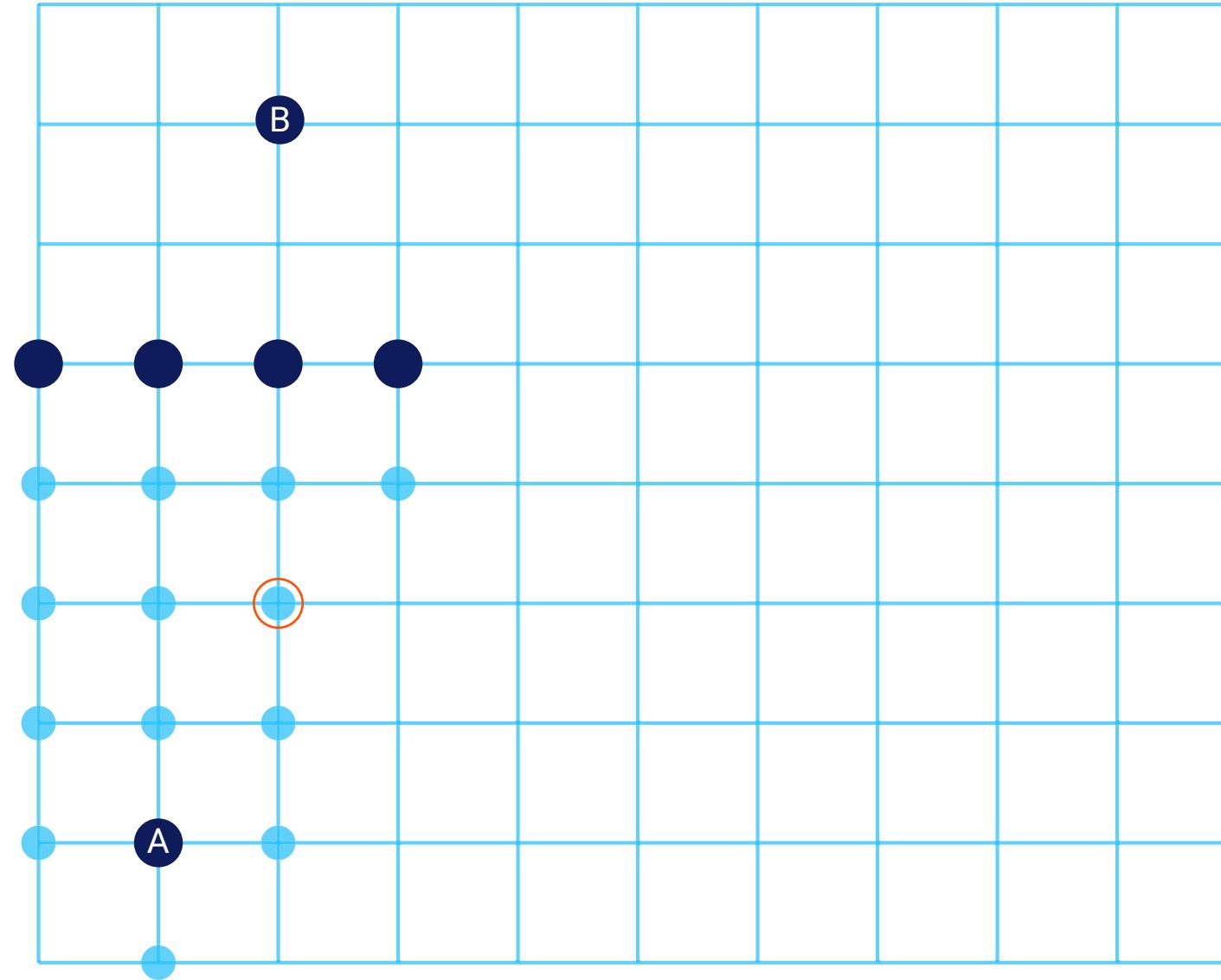


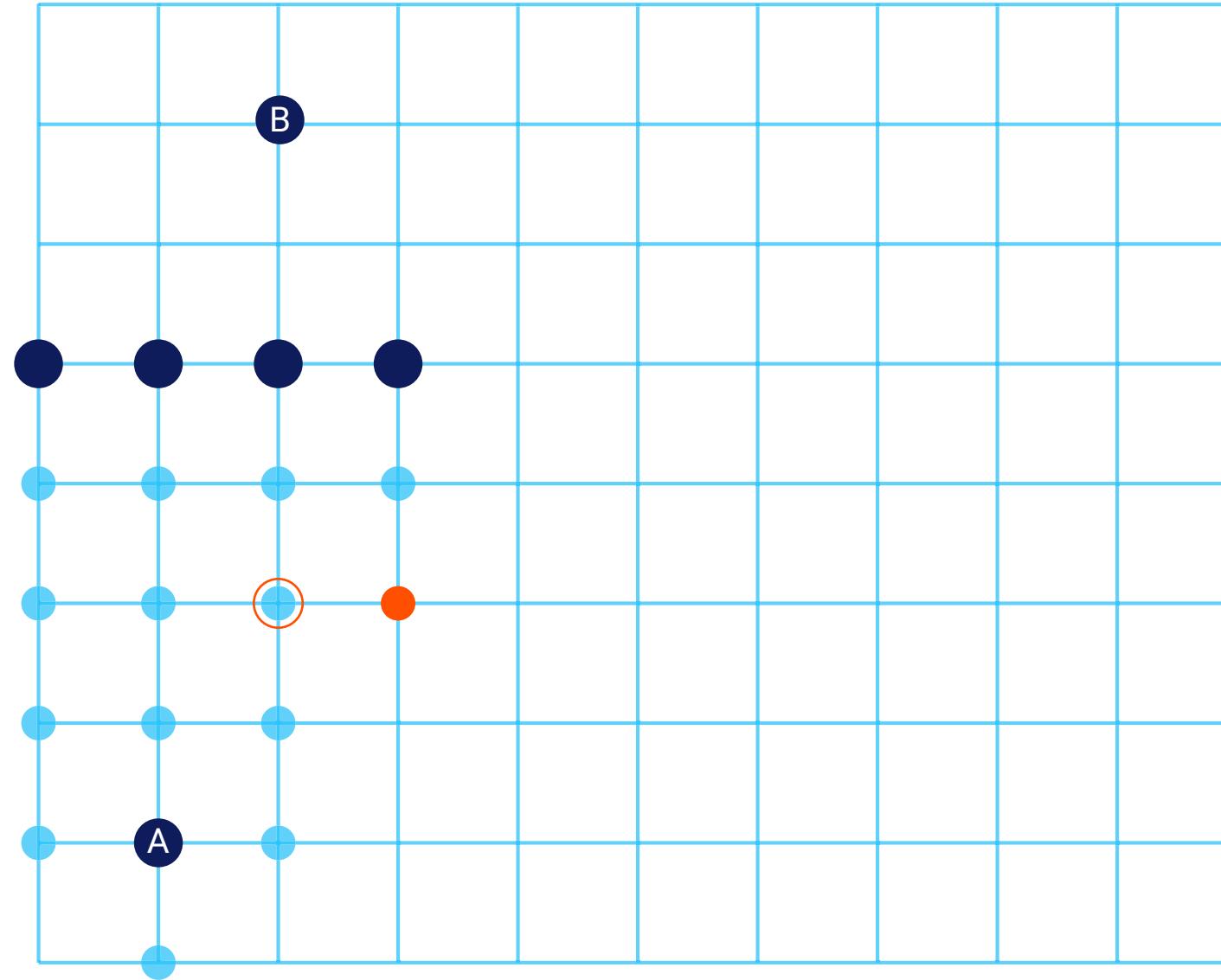


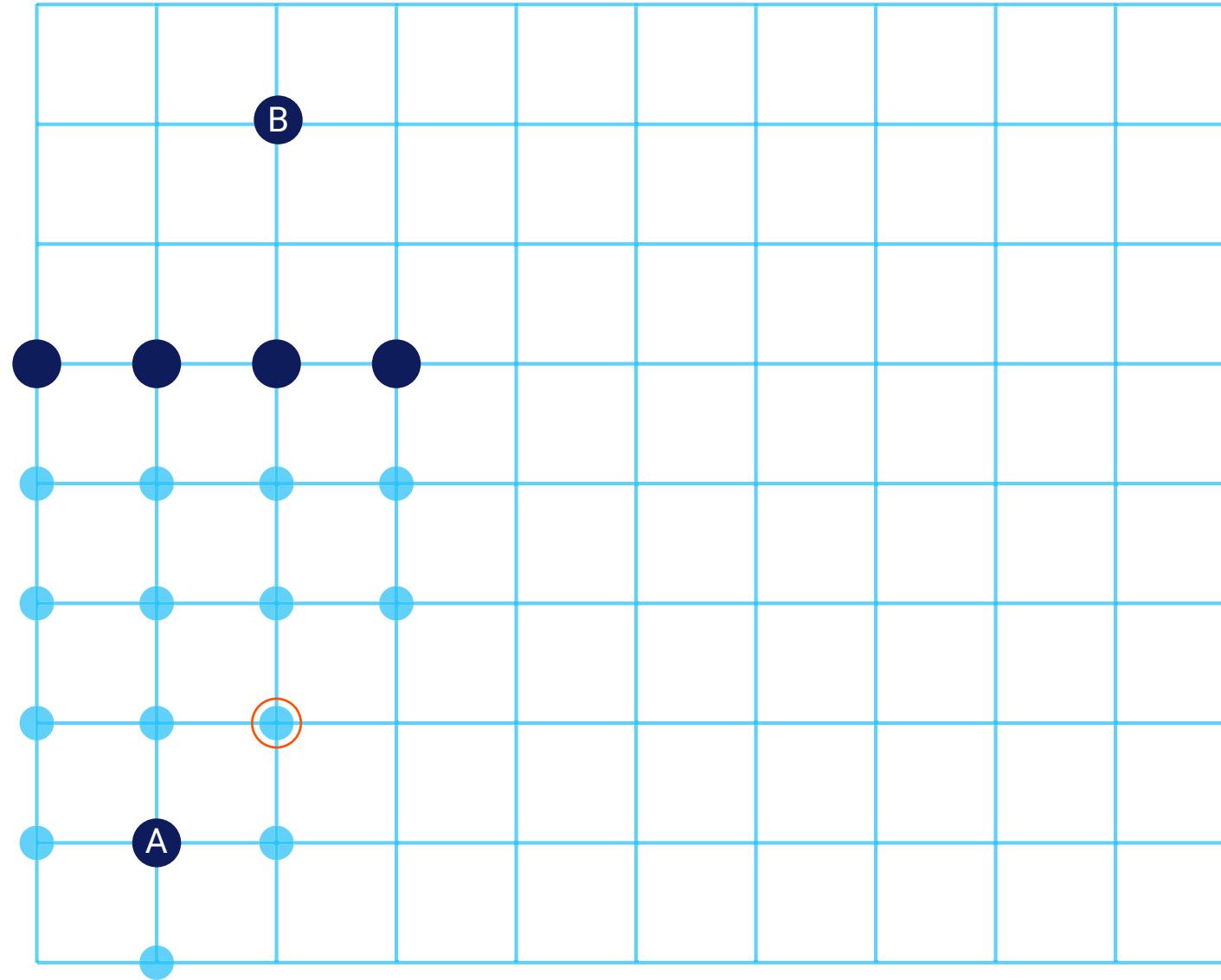


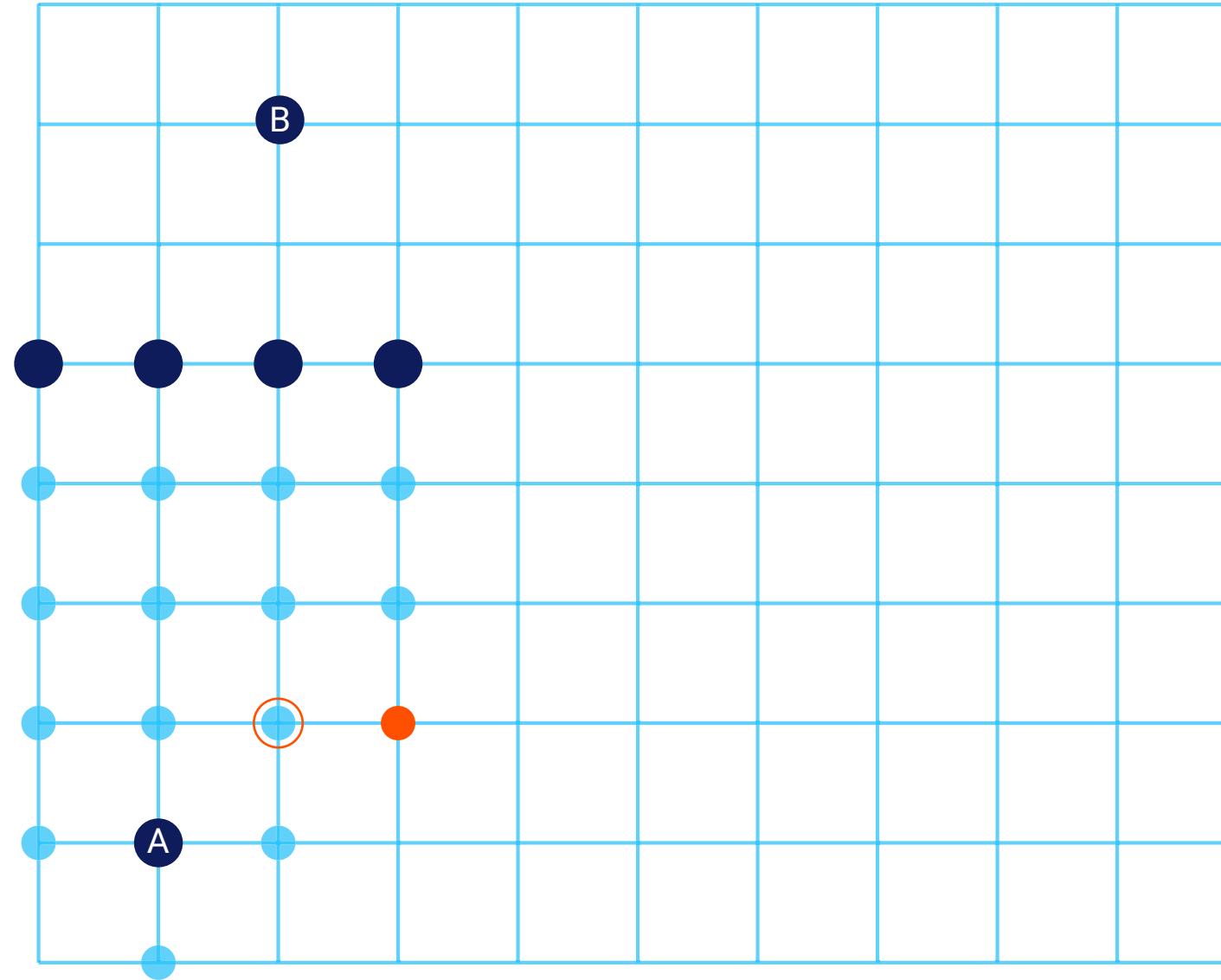


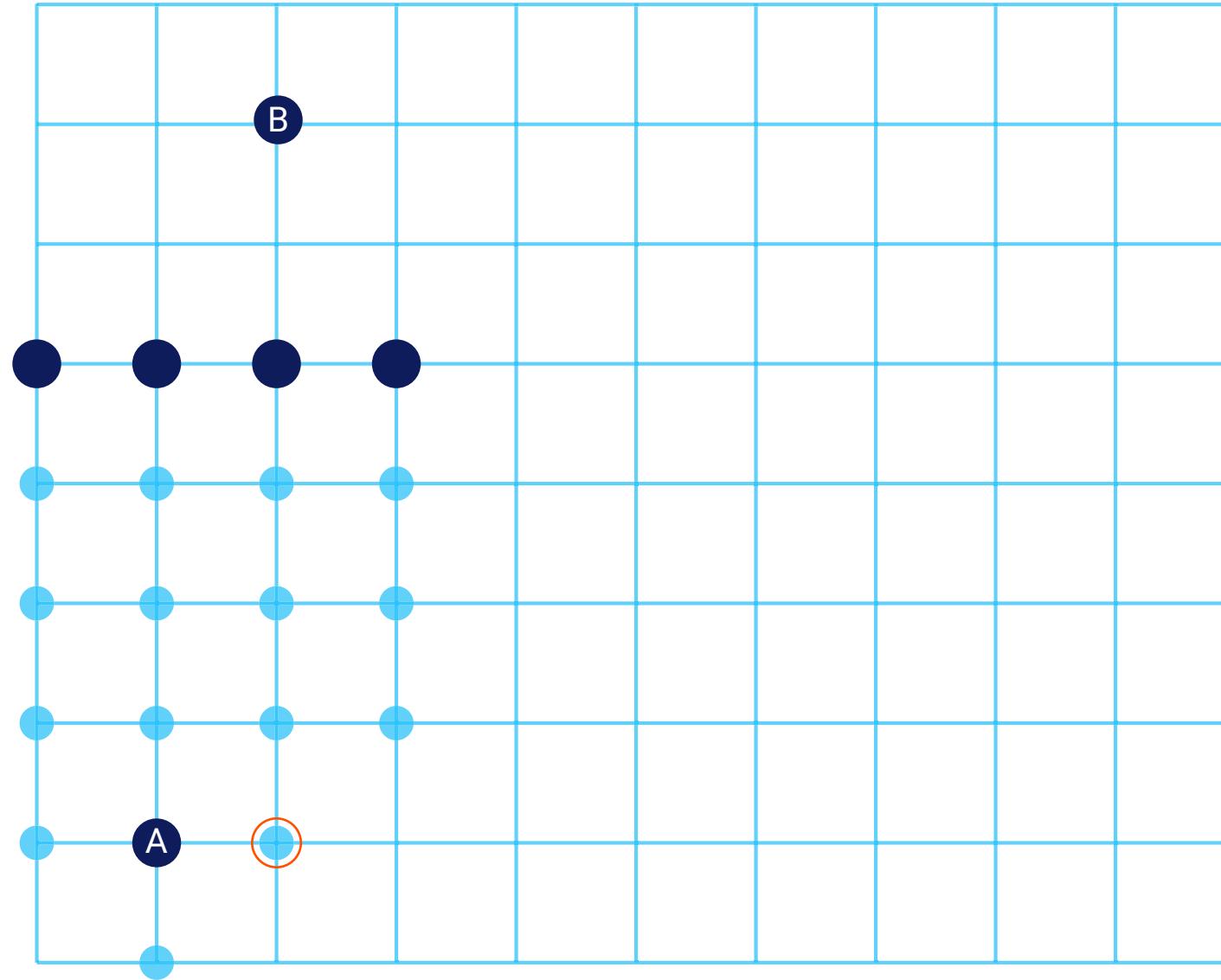


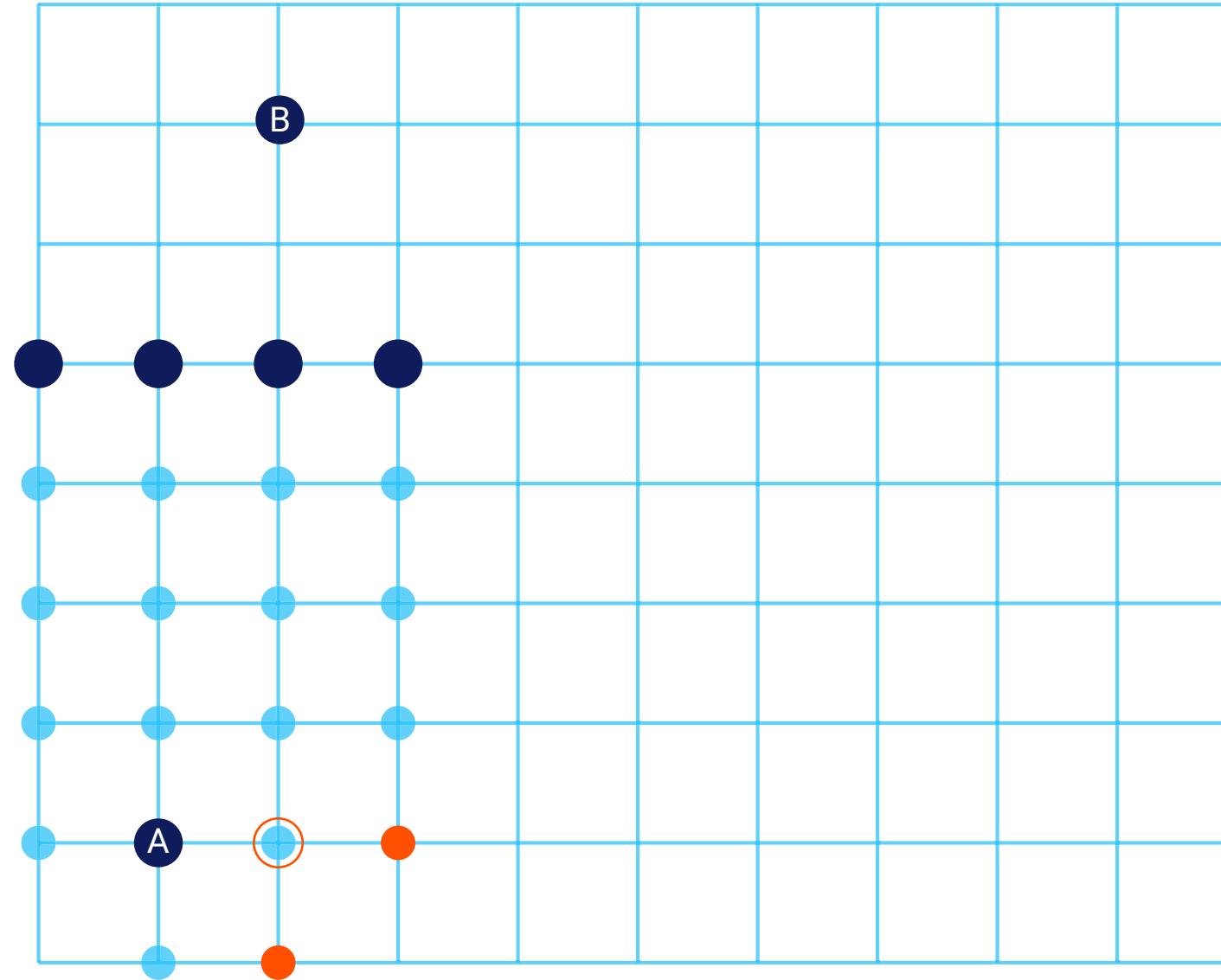


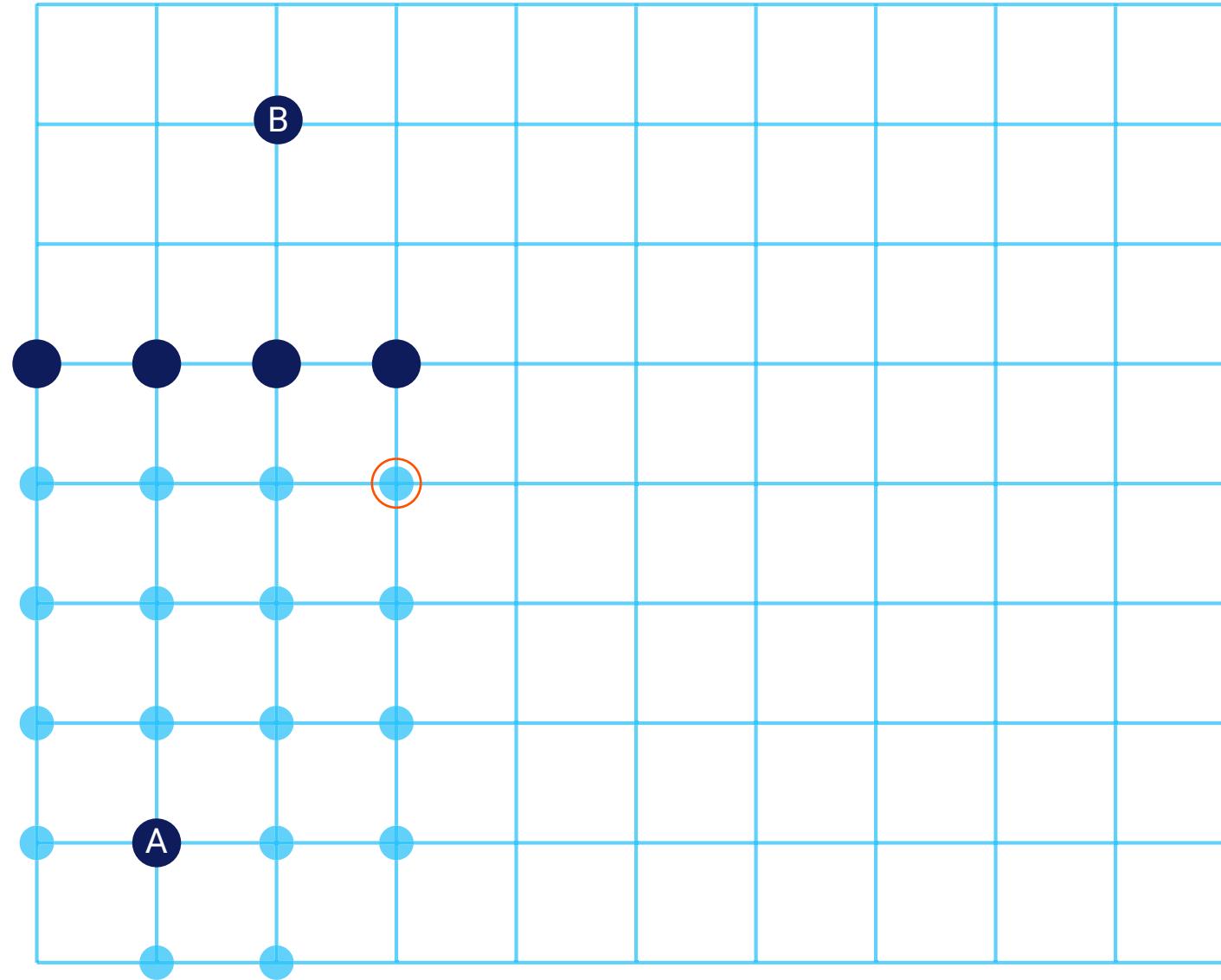


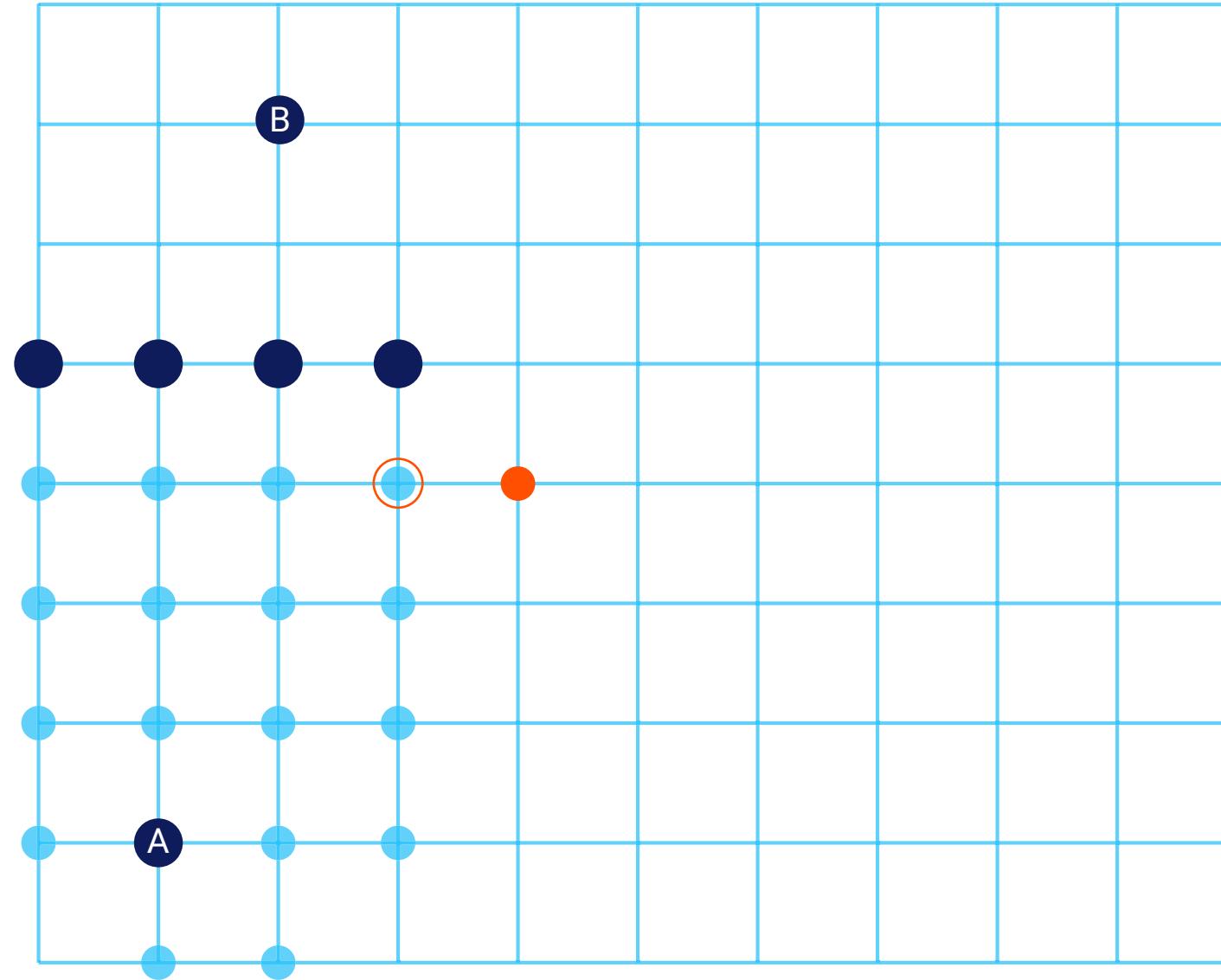


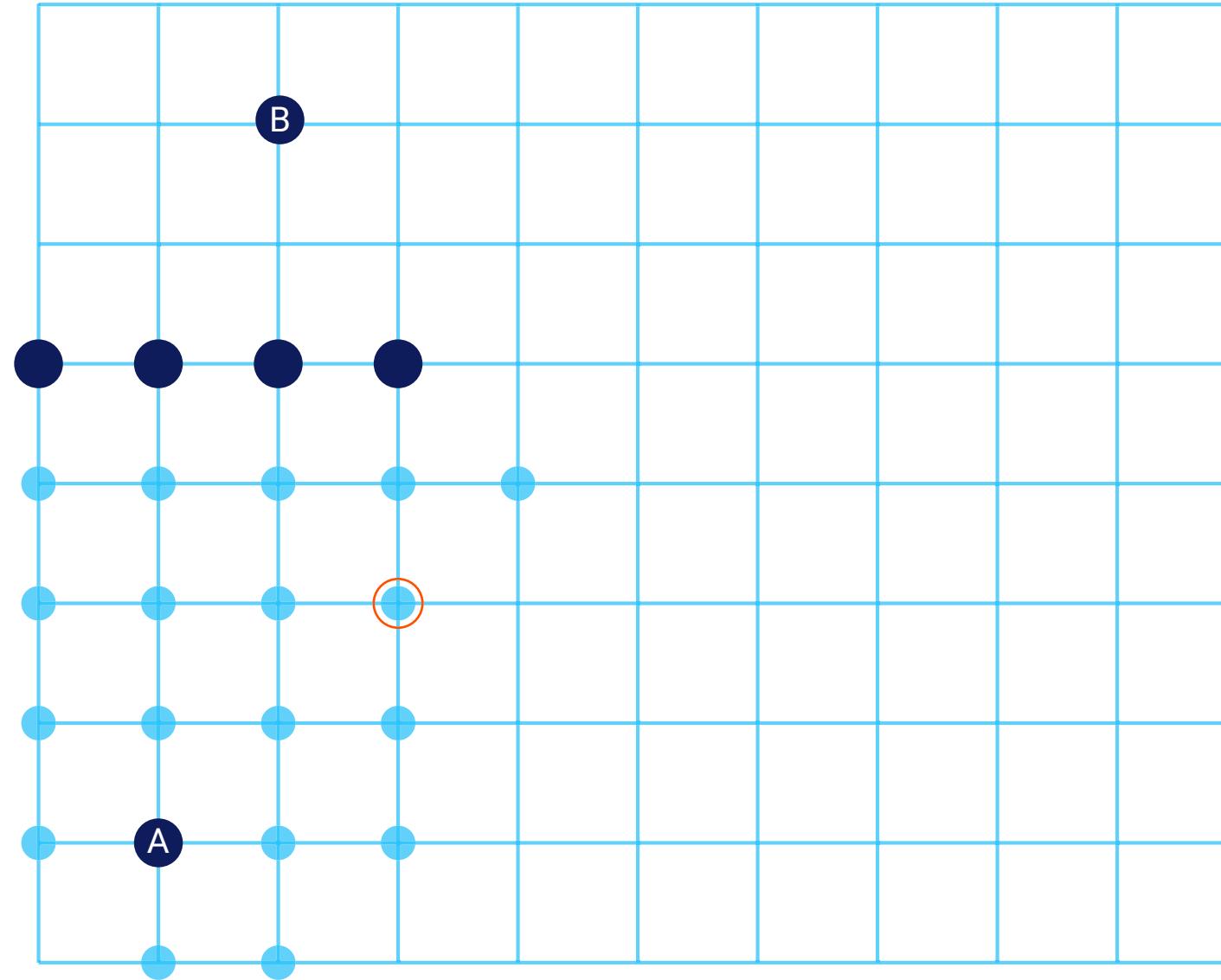


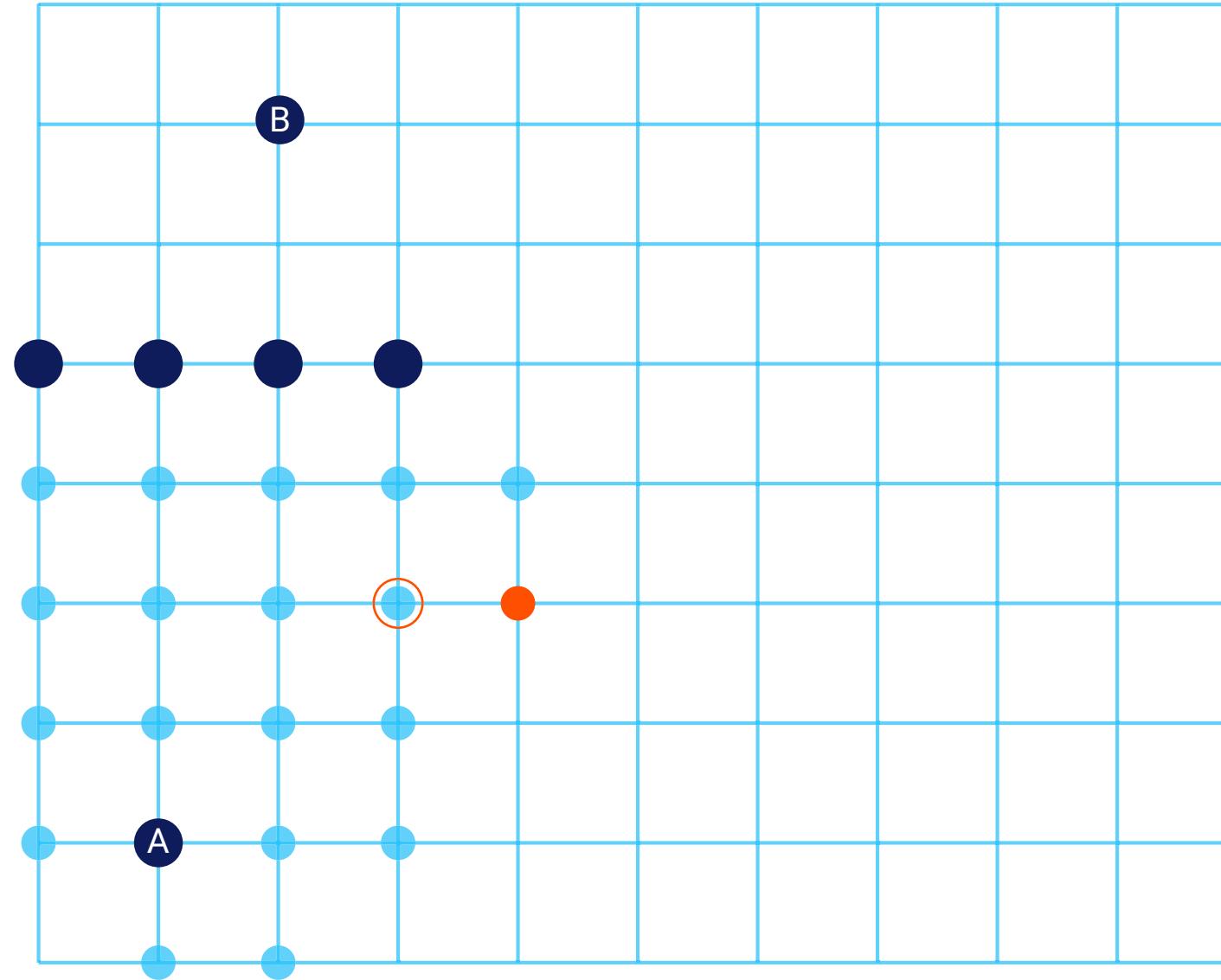


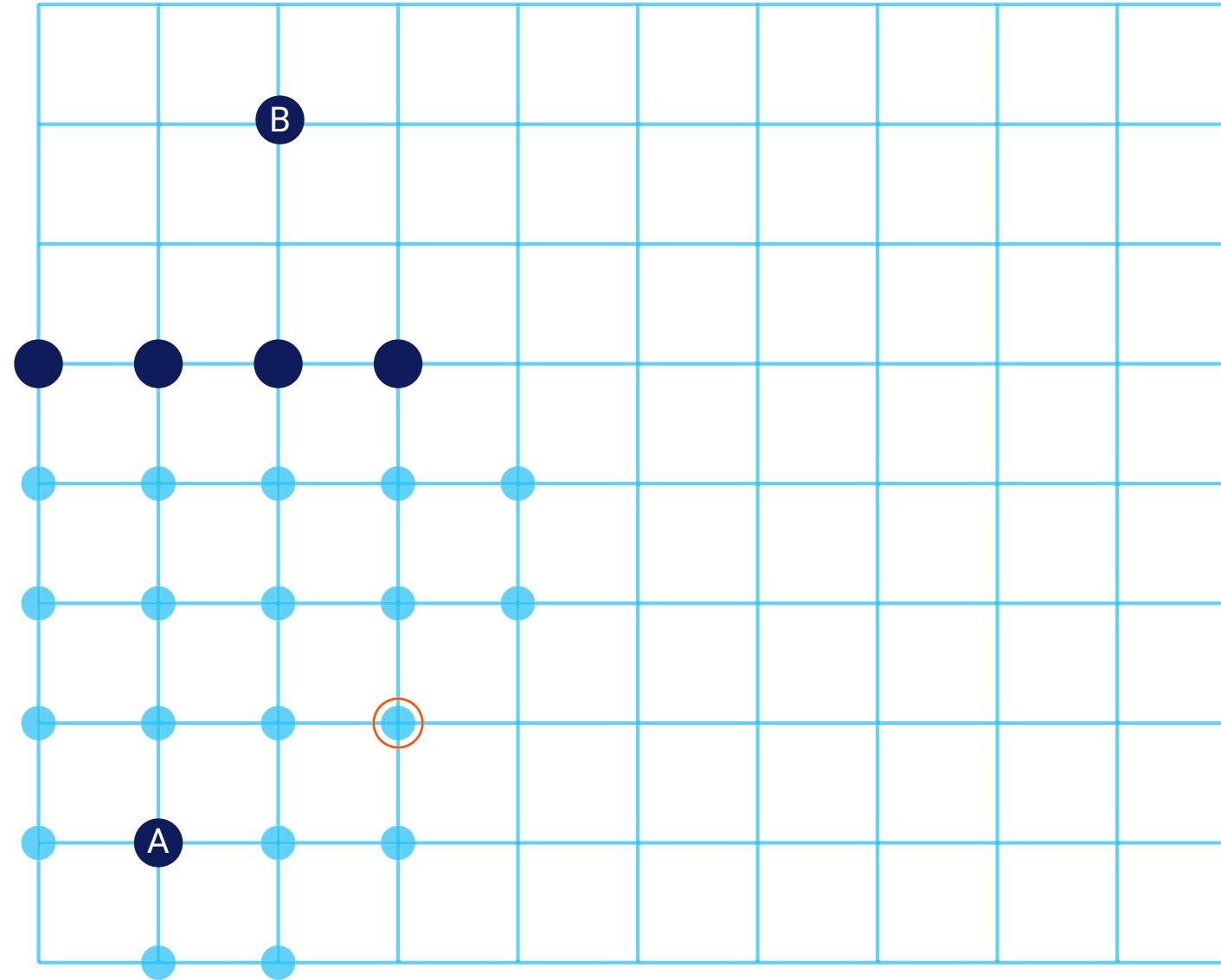


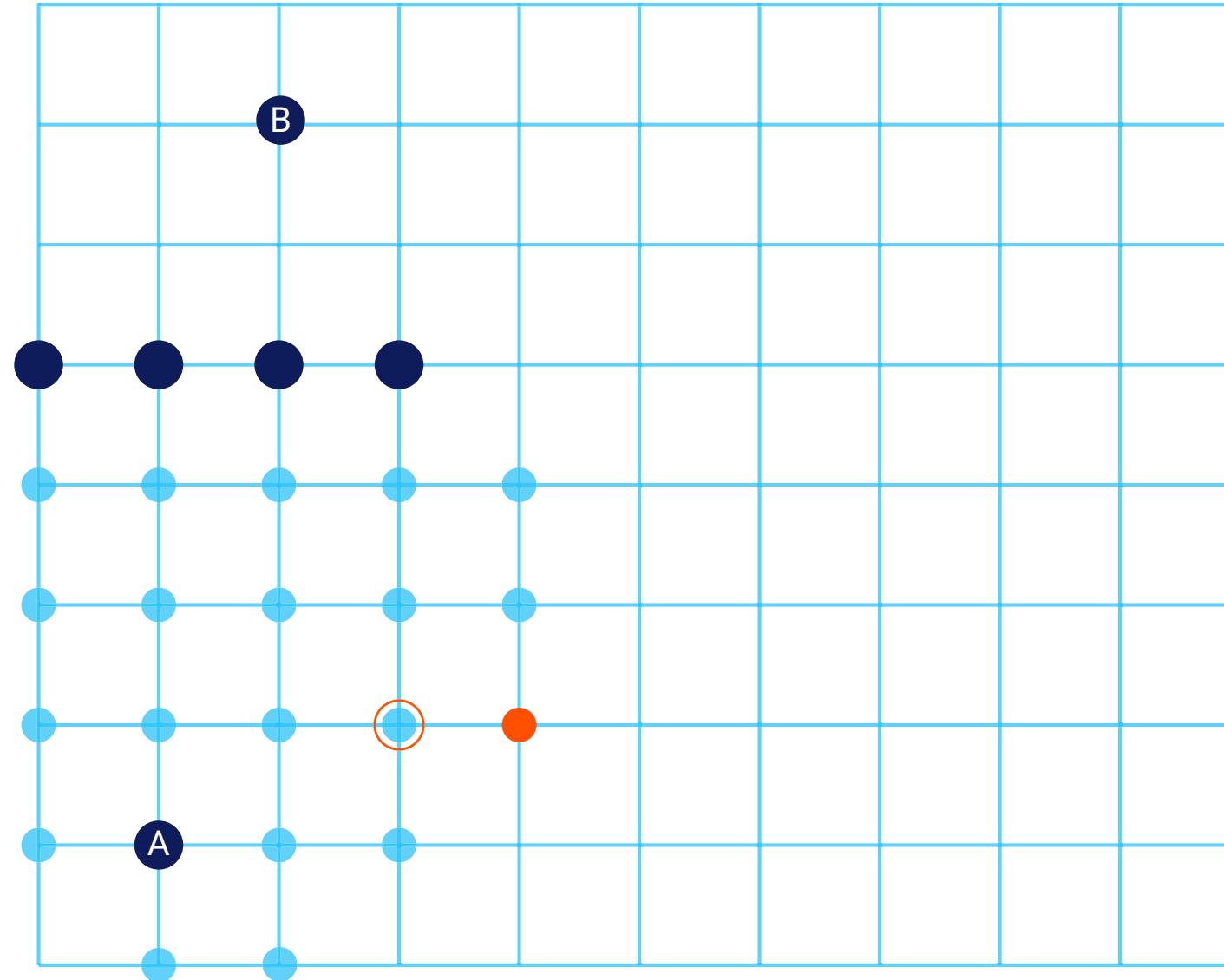


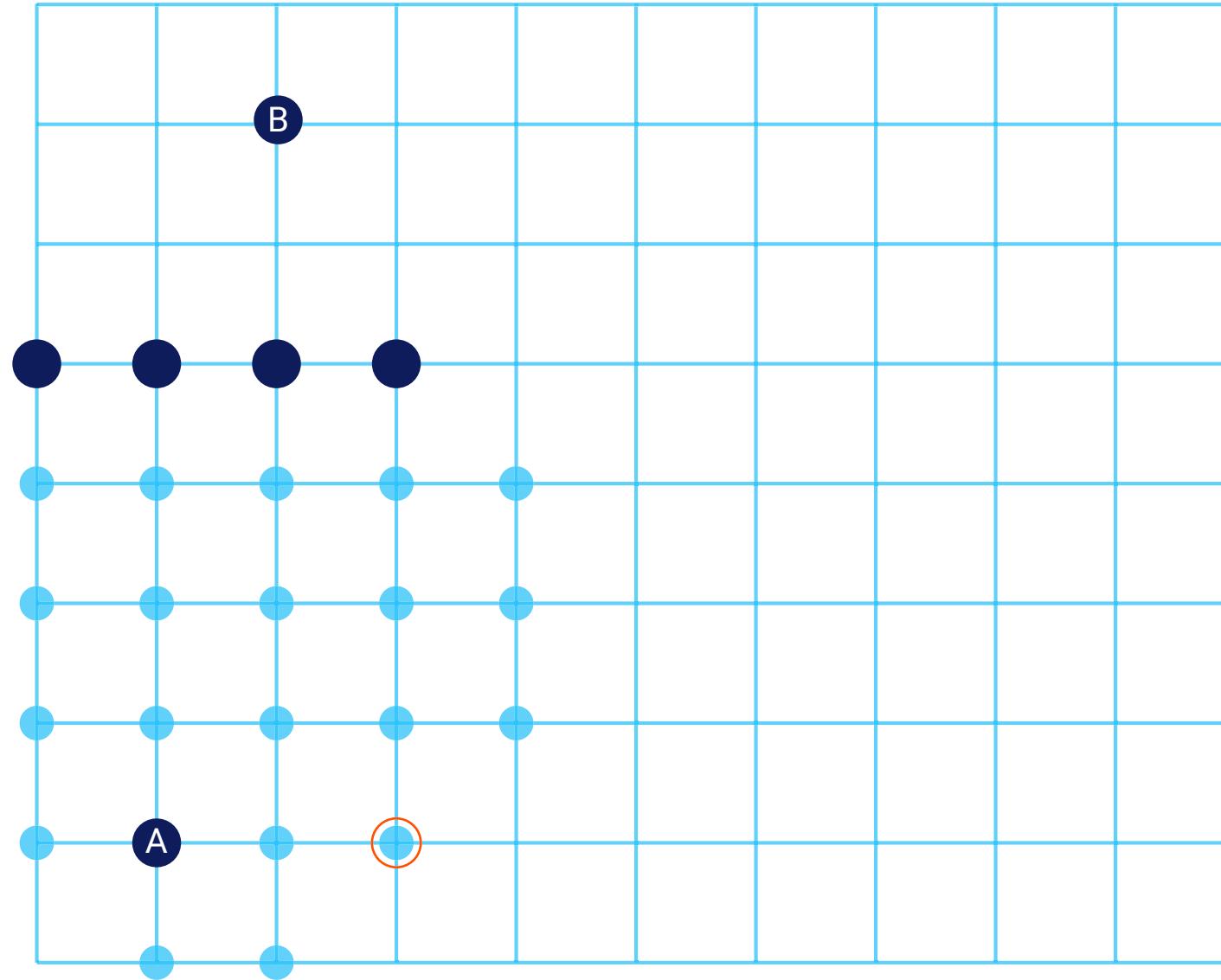


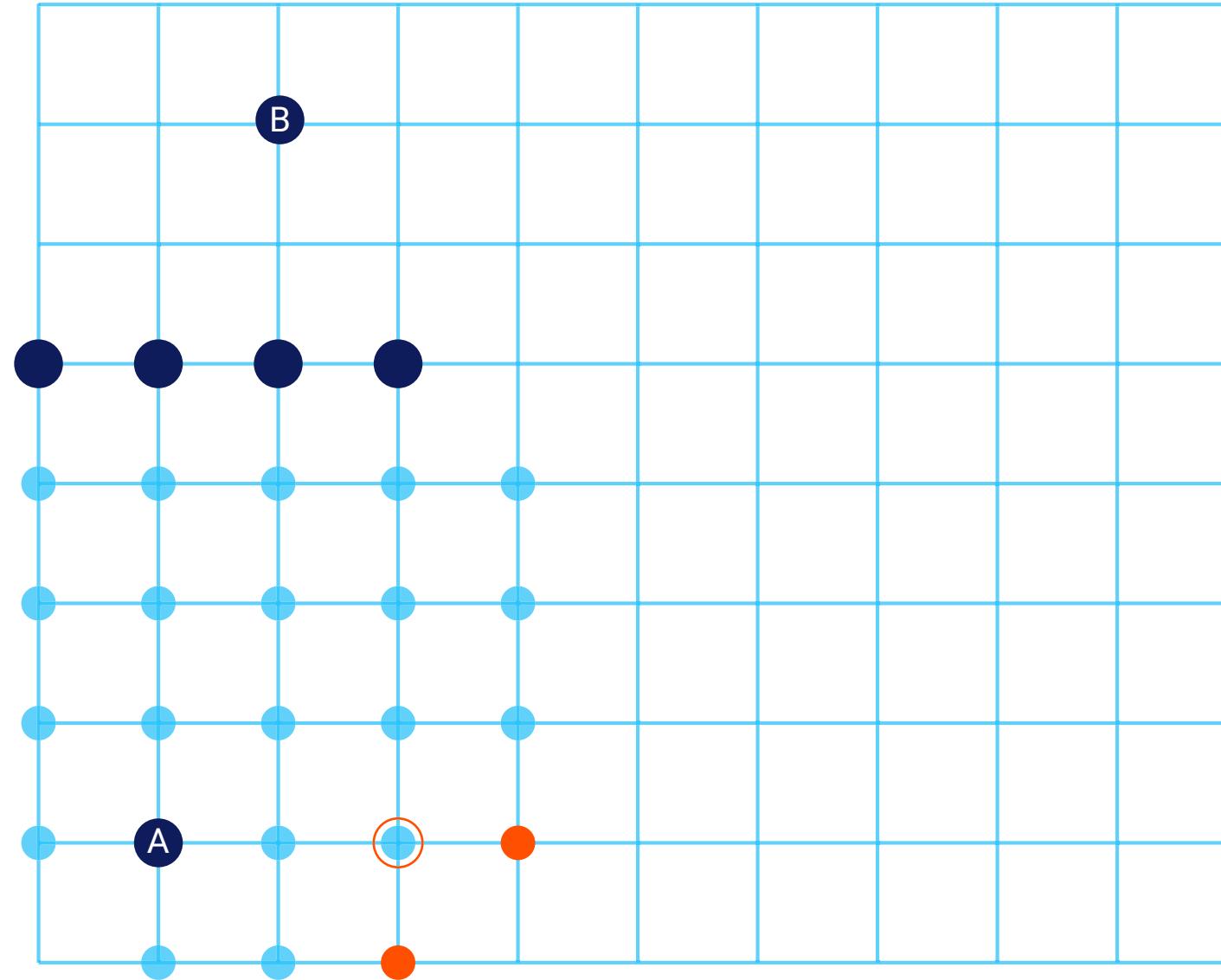


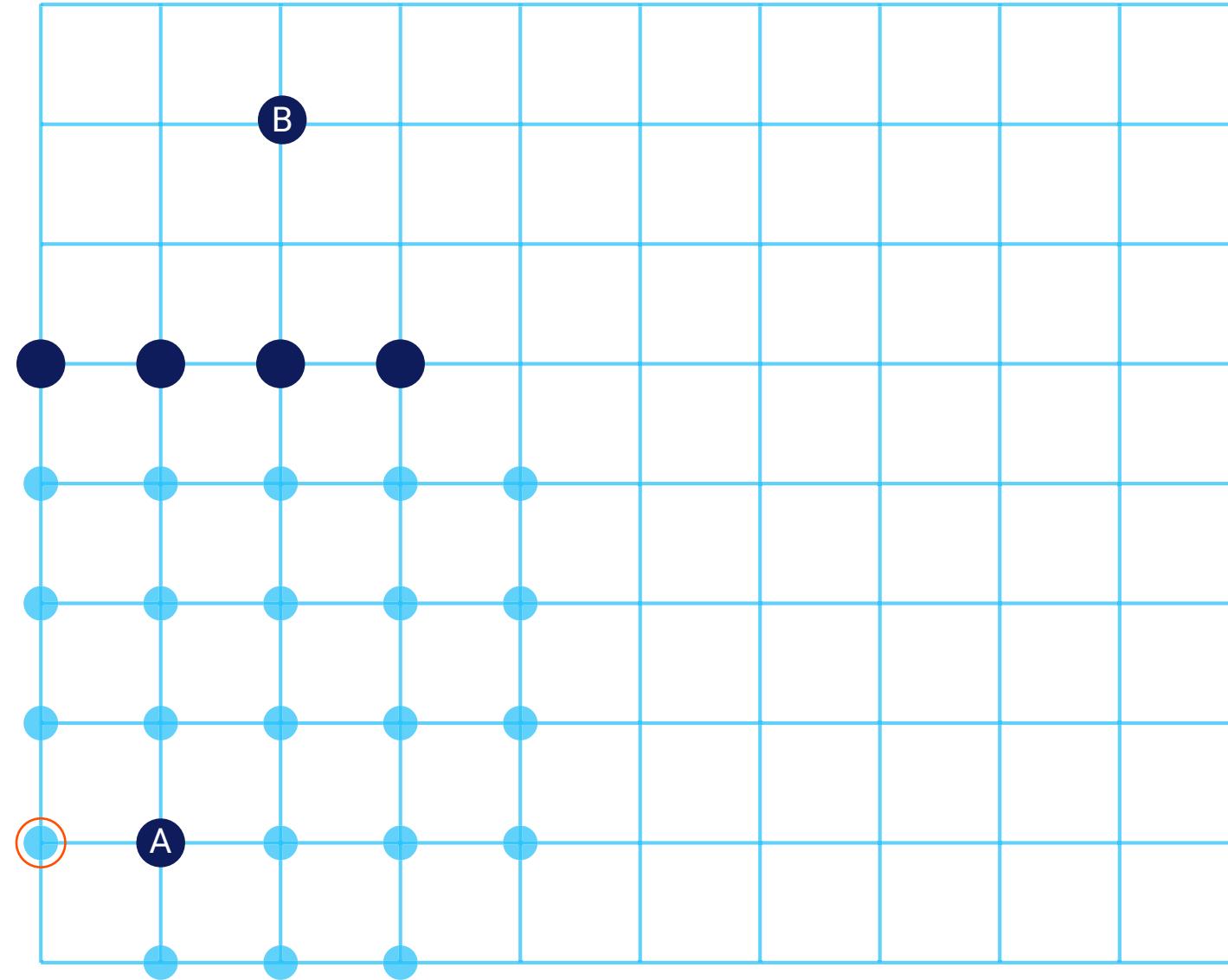


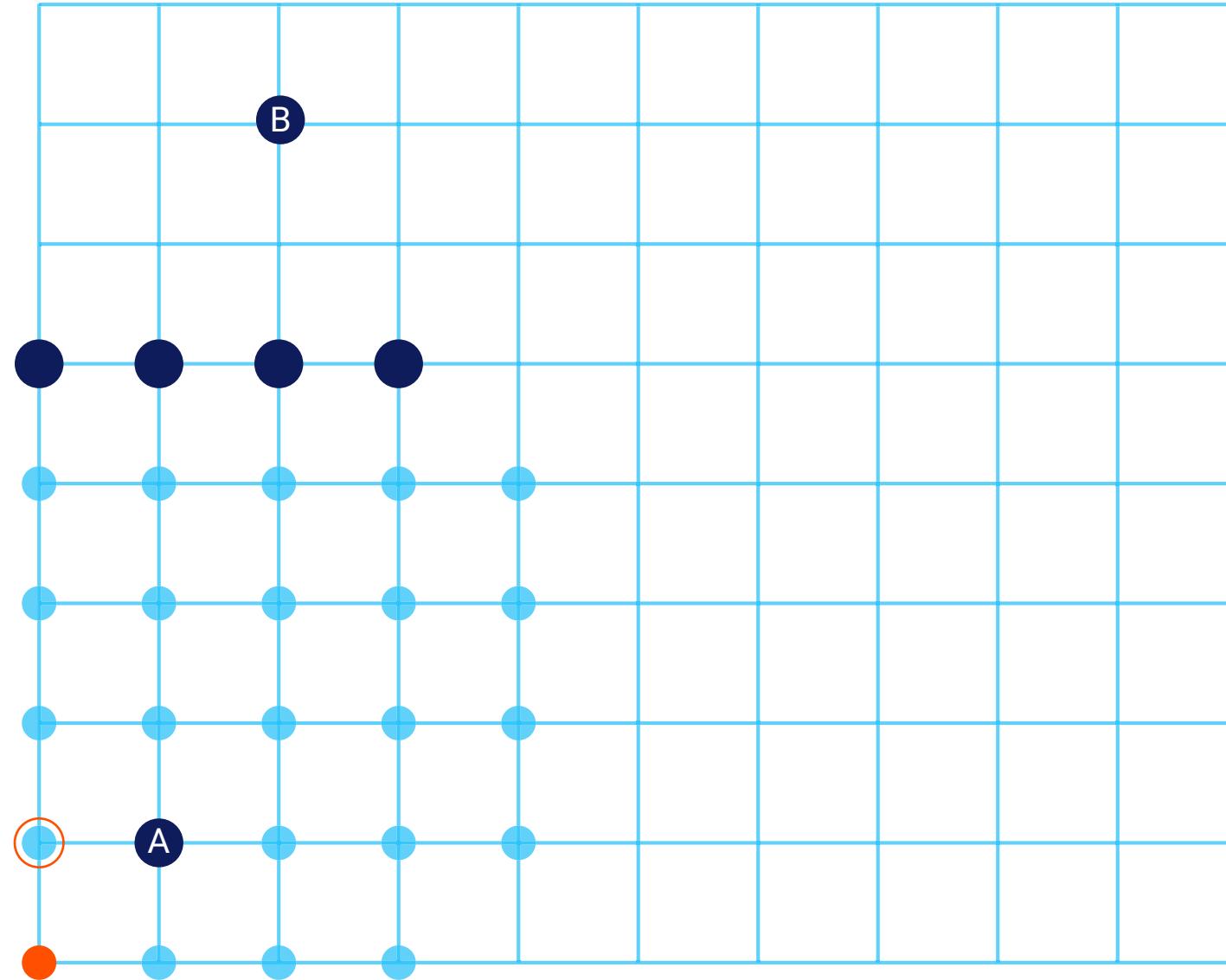


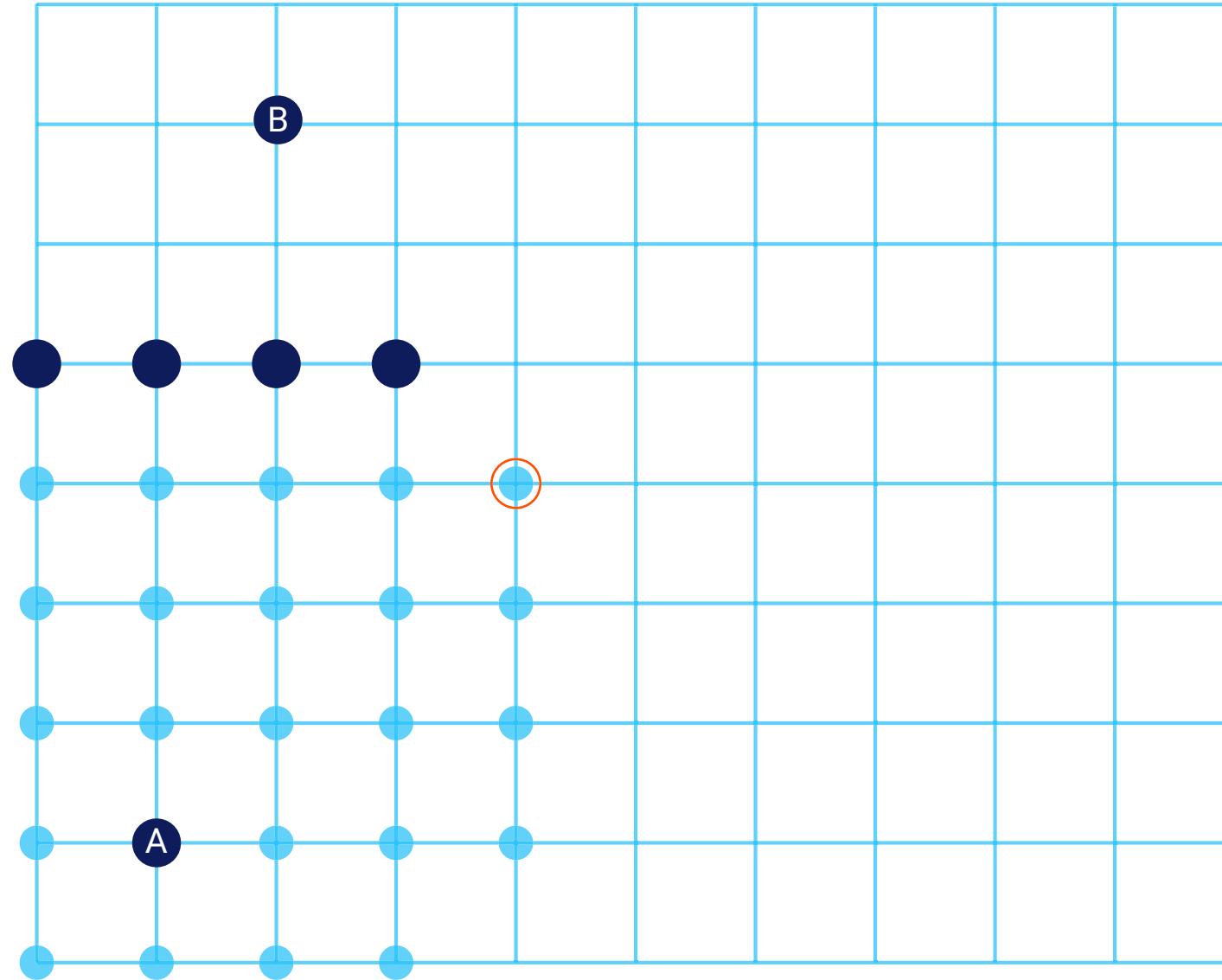


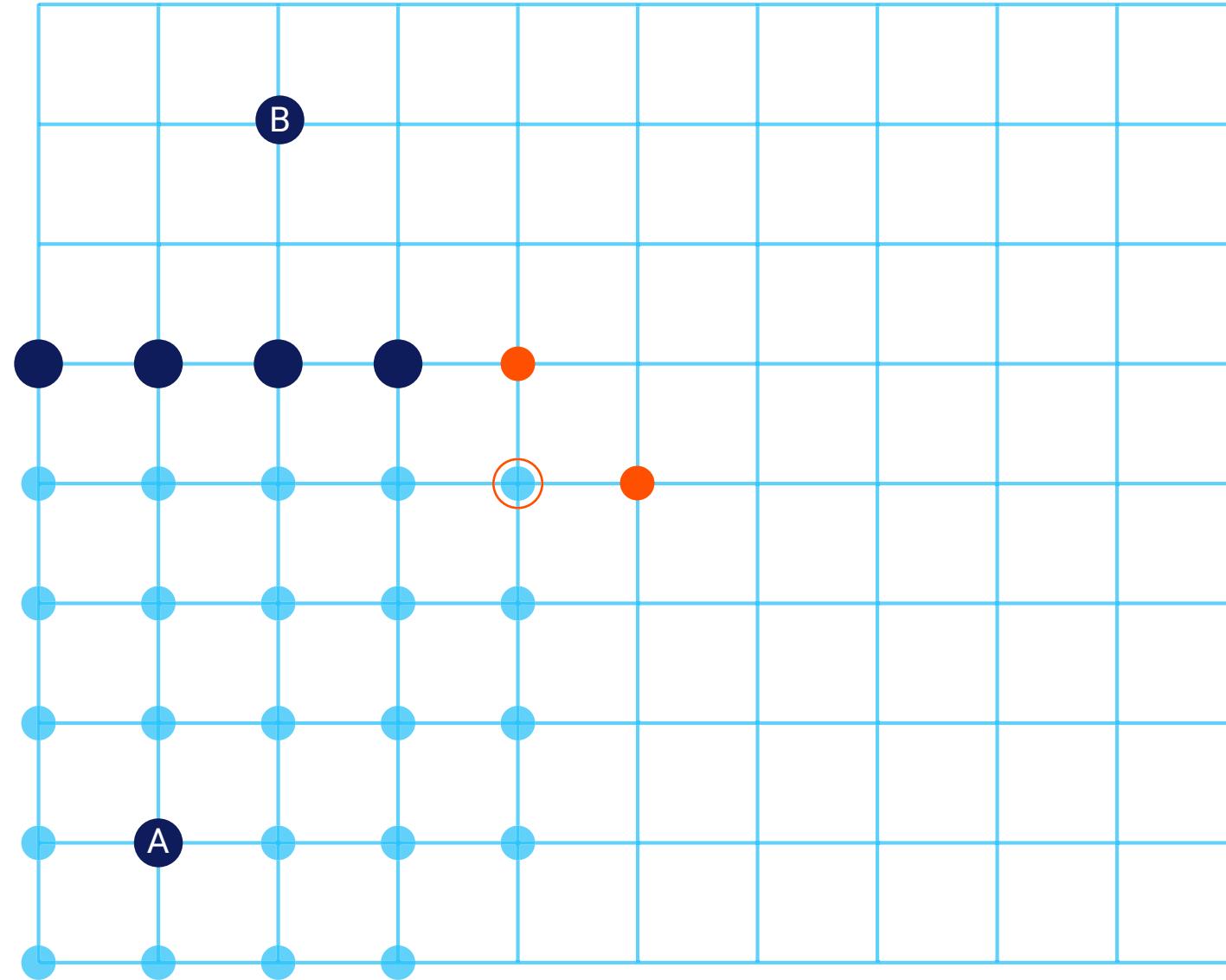


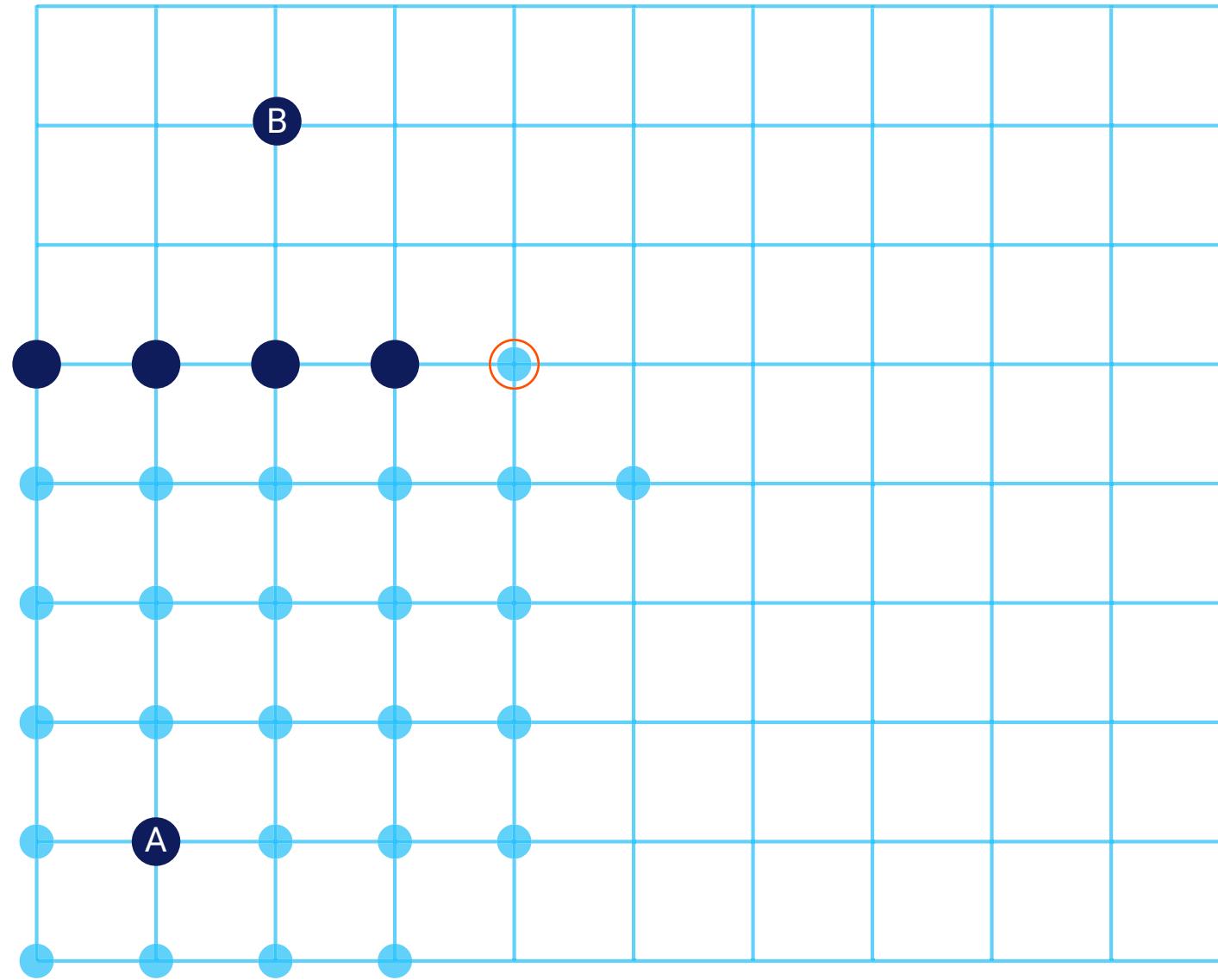


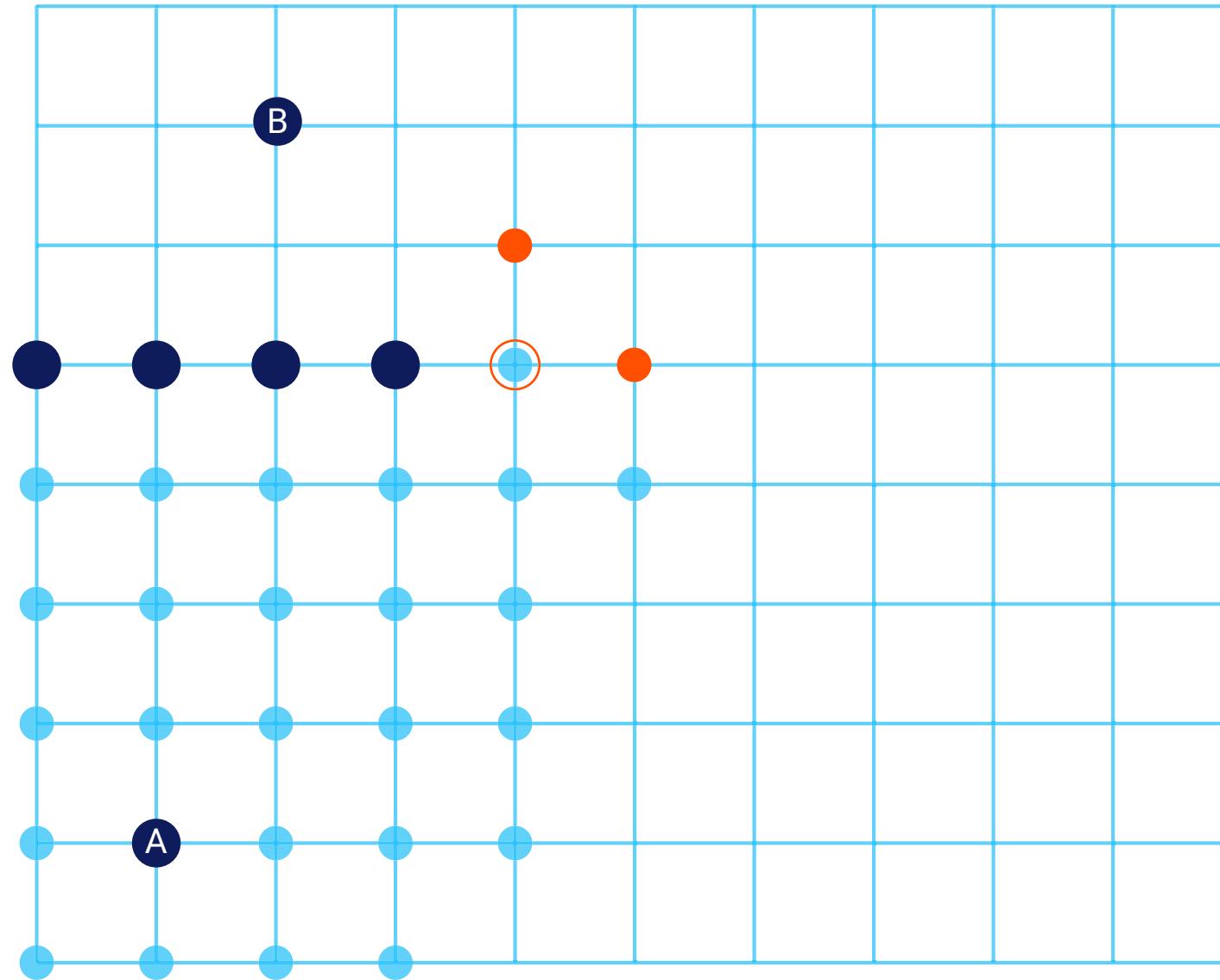


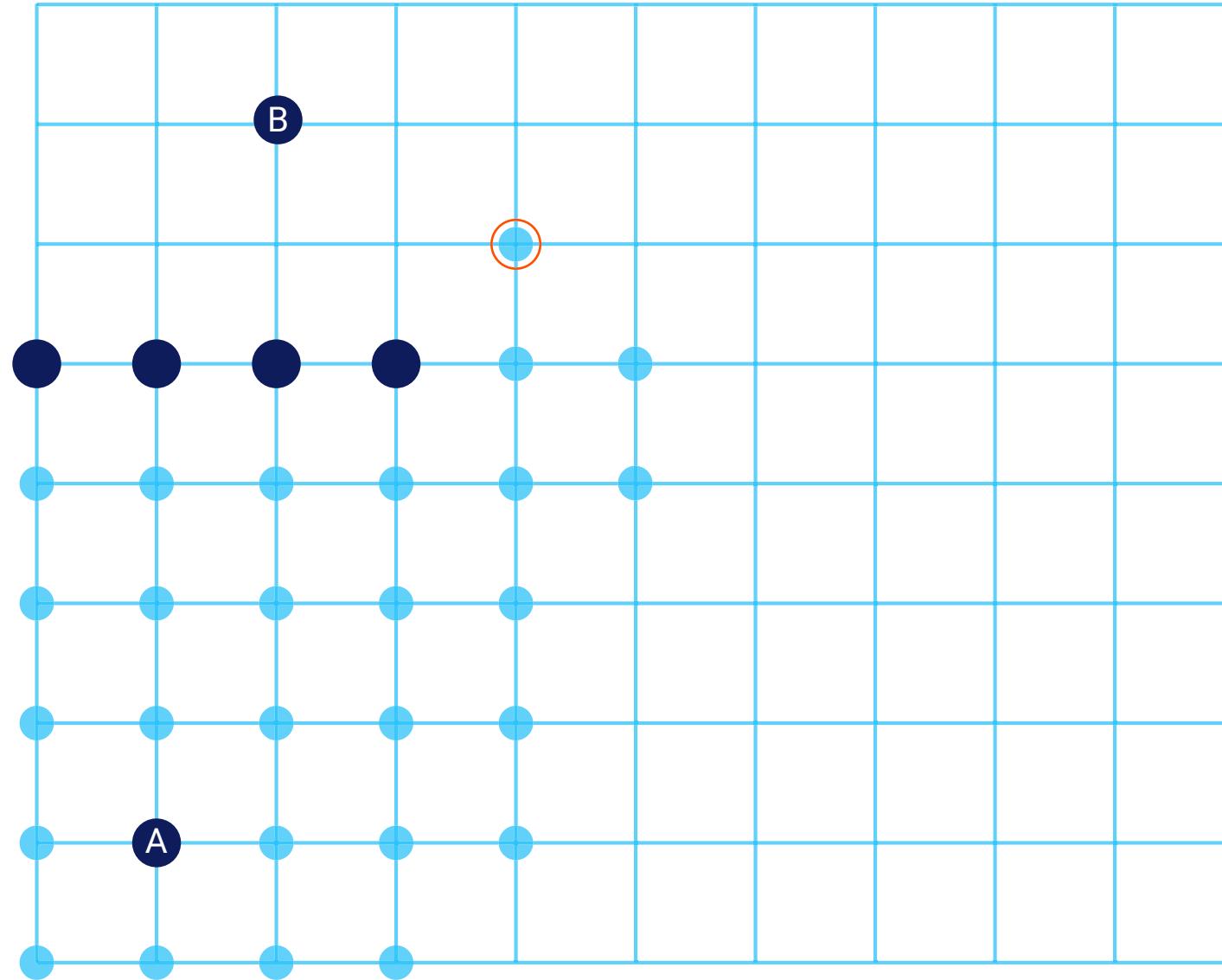


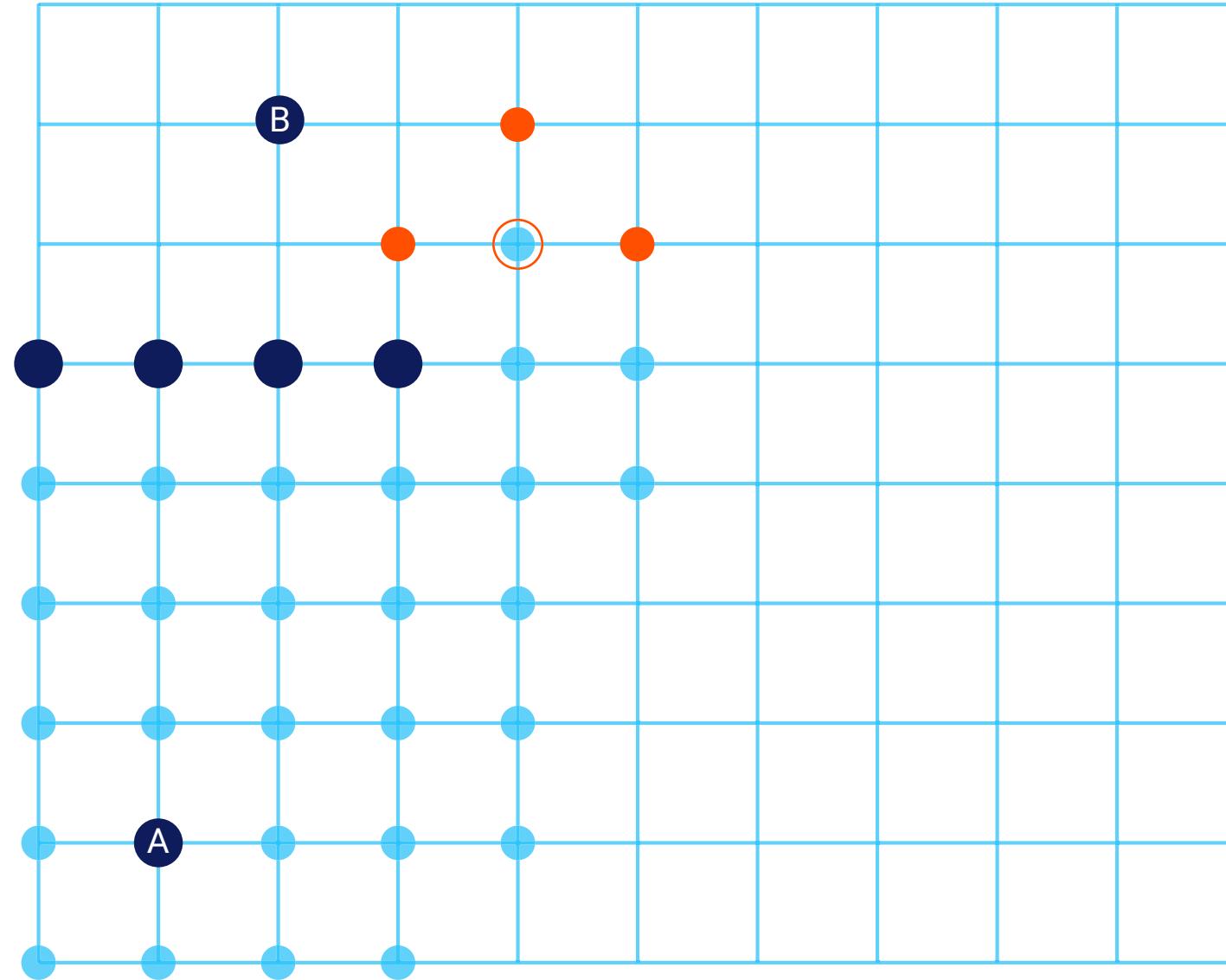


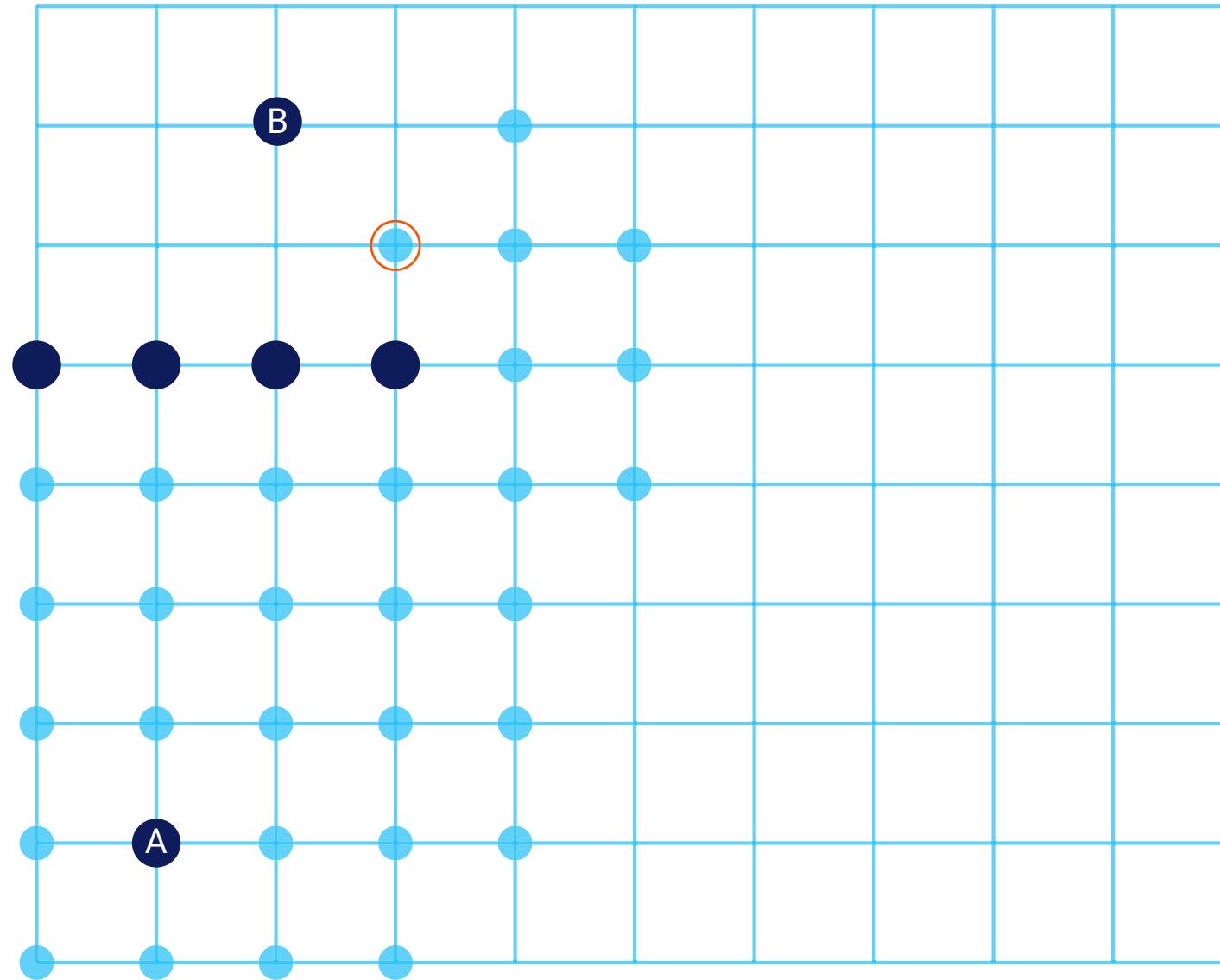


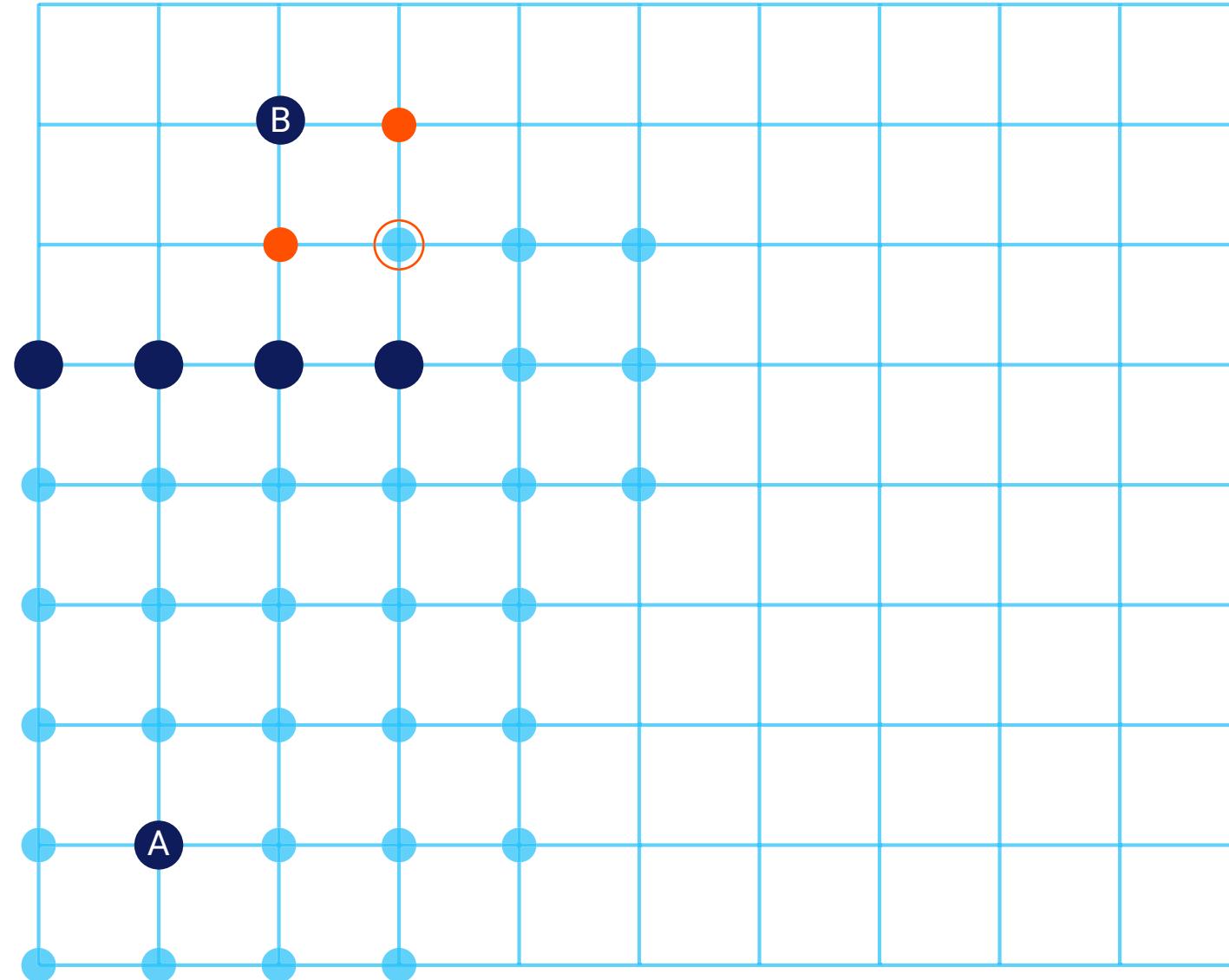


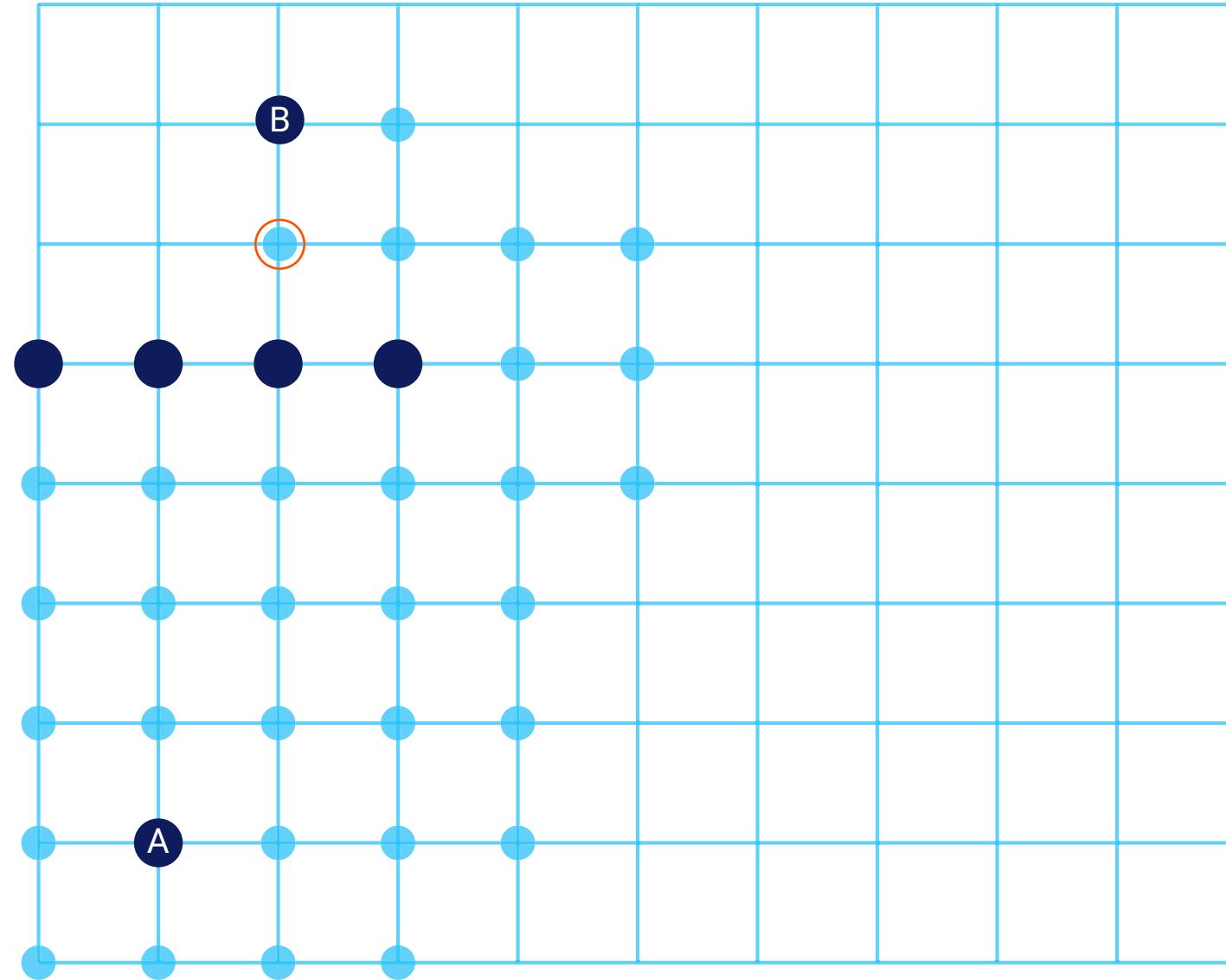


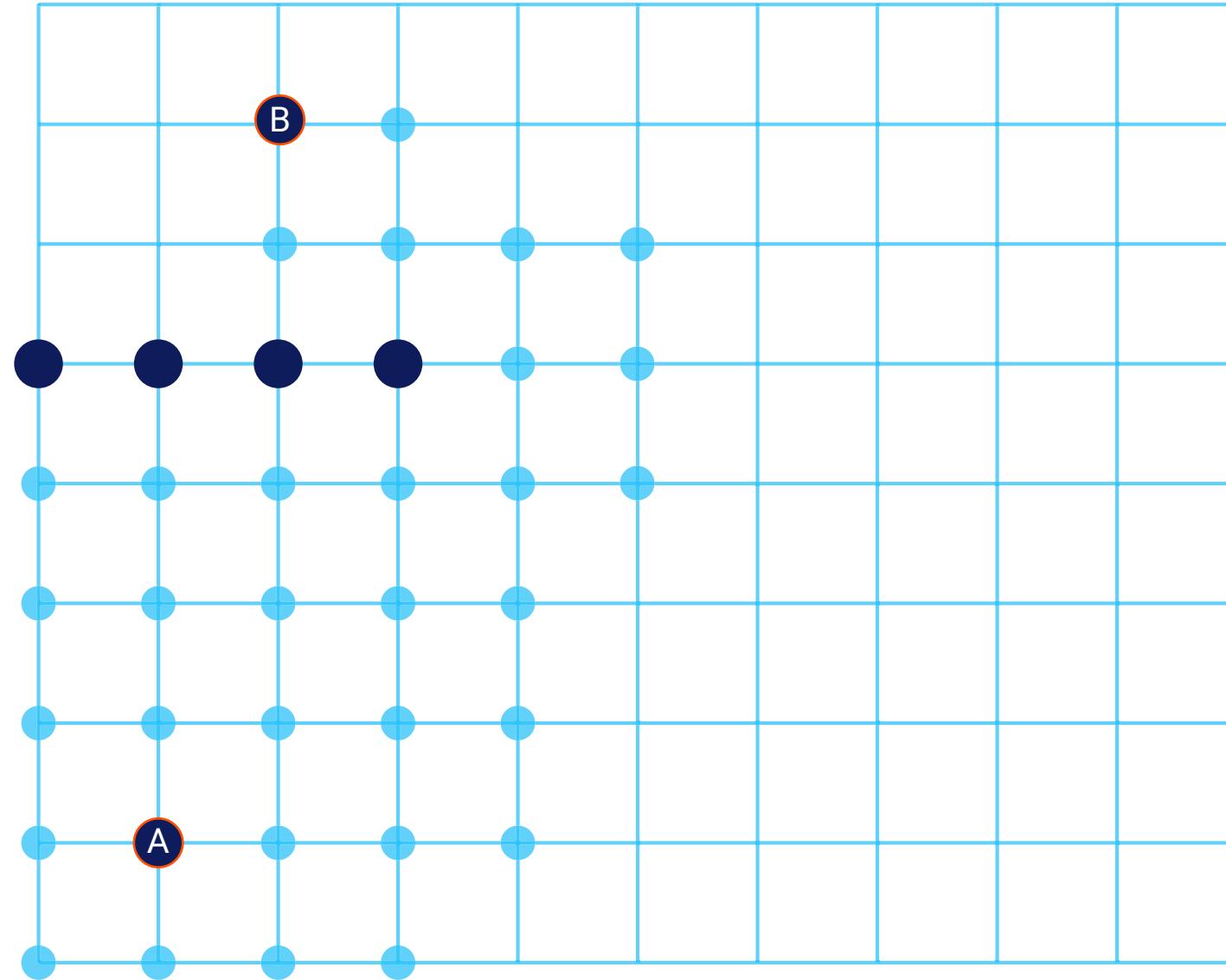


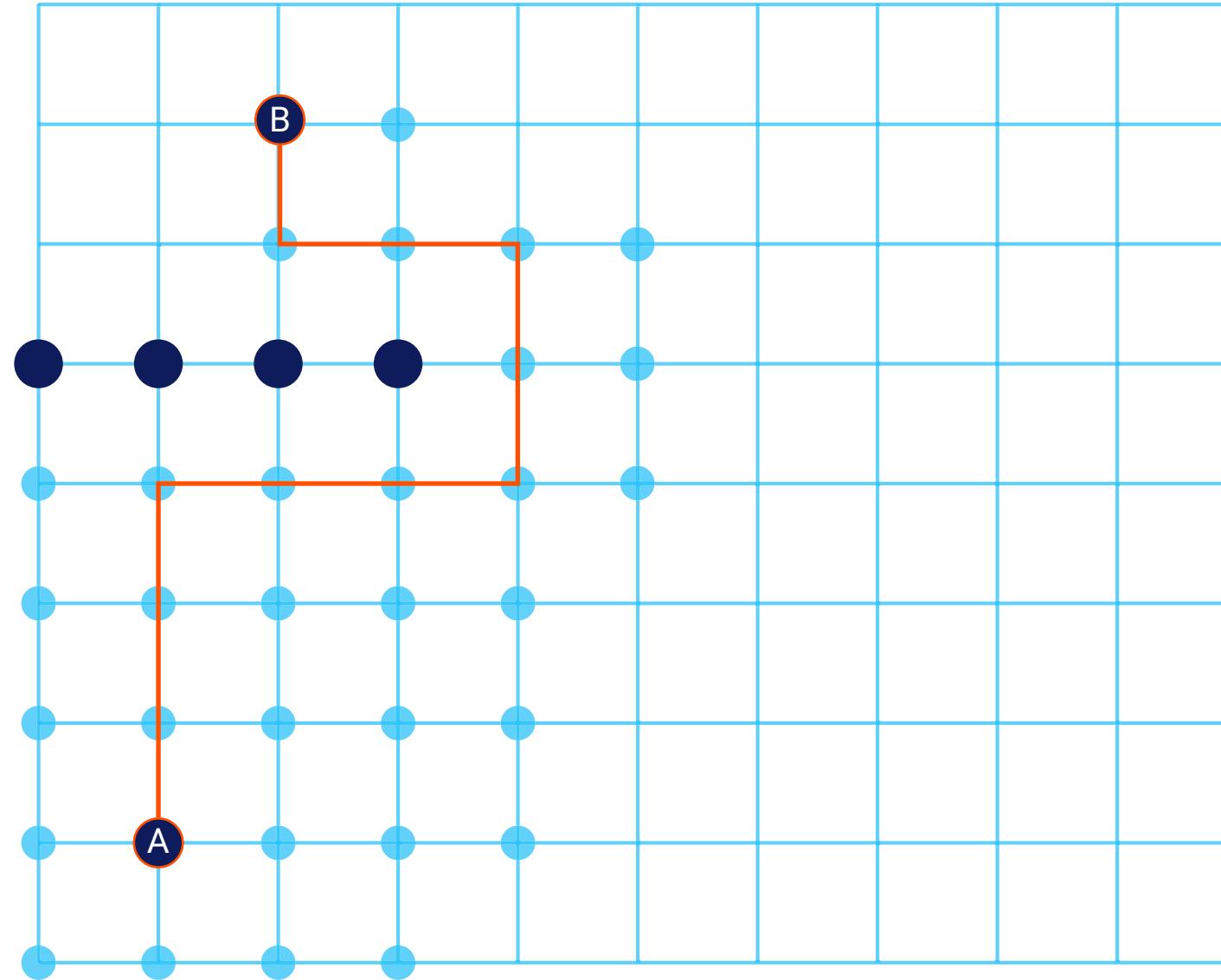










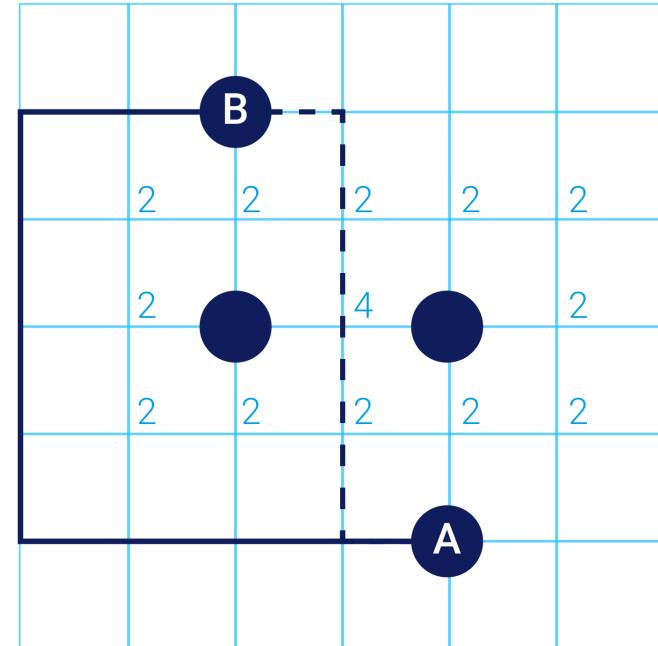


# Heatmap

Pad met laagste kosten  $\neq$  kortste pad

Heat penalty =

Heat constant - (2 · Manhattan distance to gate)

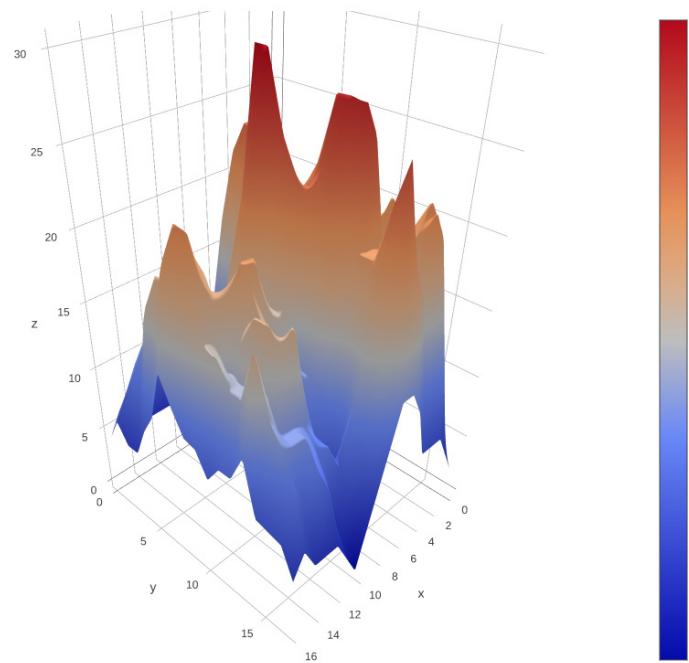


cost =  $6 + 8 = 14$

cost = 10

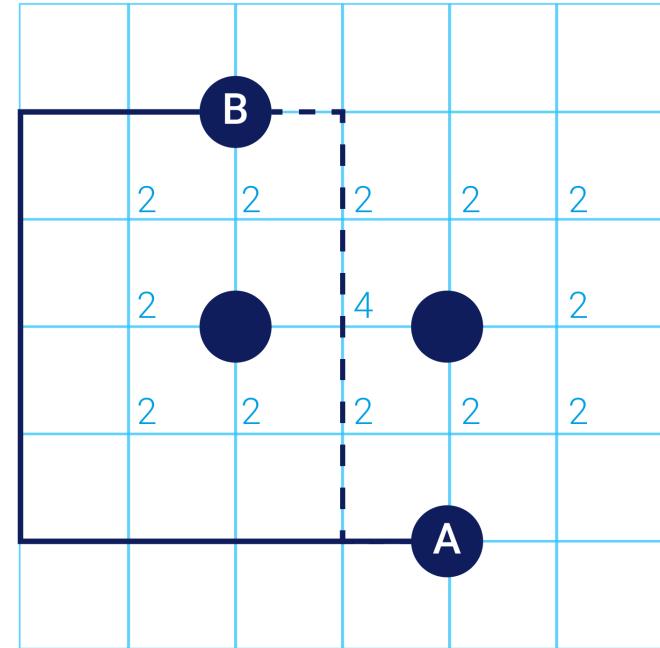
# Heatmap

Pad met laagste kosten  $\neq$  kortste pad



Heat penalty =

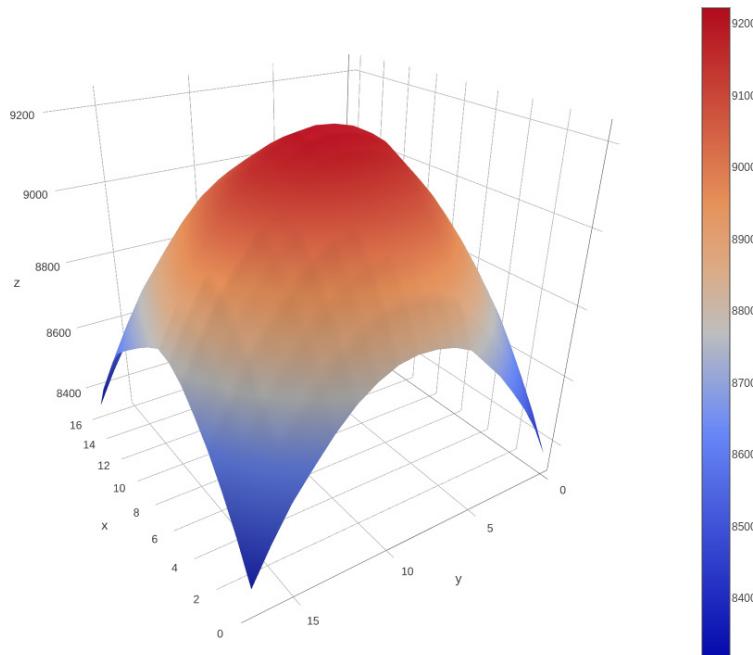
Heat constant - (2 · Manhattan distance to gate)



cost =  $6 + 8 = 14$   
cost = 10

# Heat range

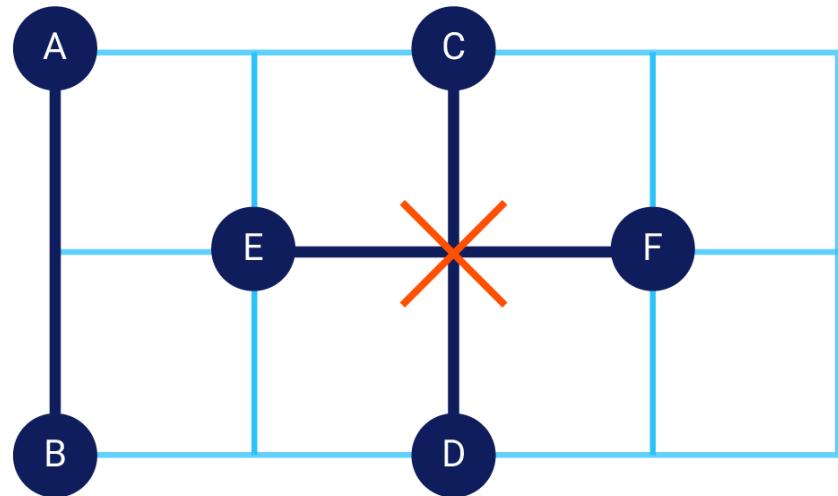
Pilot study: heat constant 5 - 30



Heat 2000, printplaat 2

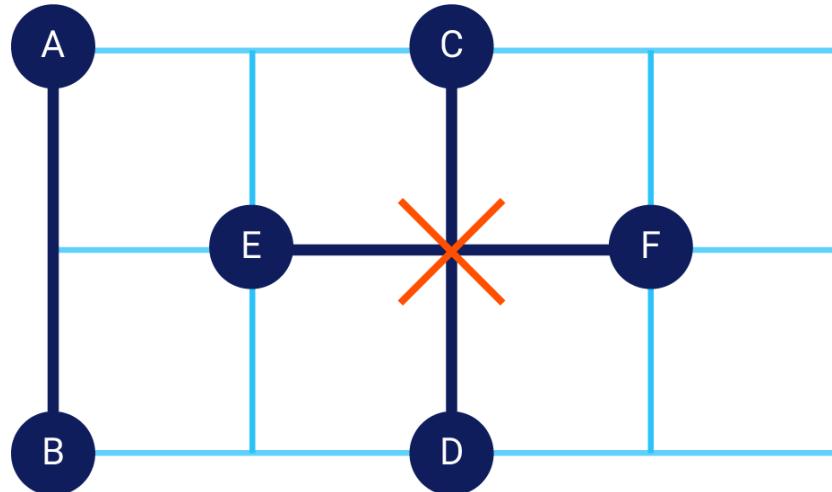
# Permutaties

Netlist: [[A, B], [C, D], [E,F]]

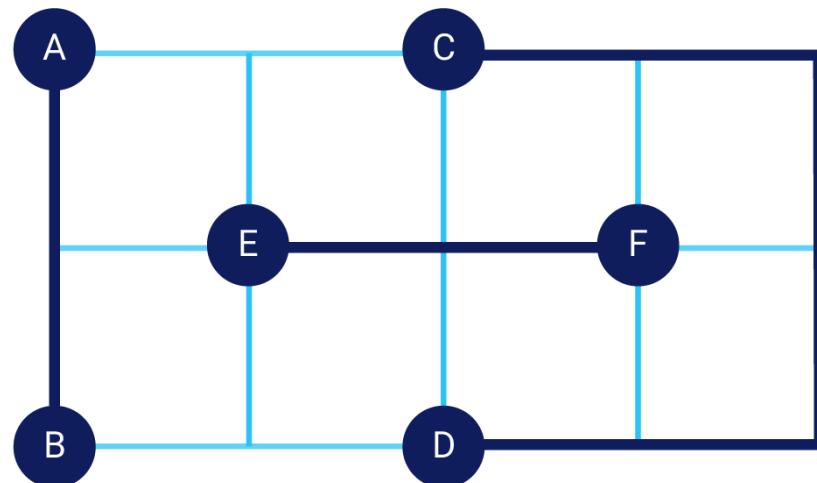


# Permutaties

Netlist: [[A, B], [C, D], [E,F]]

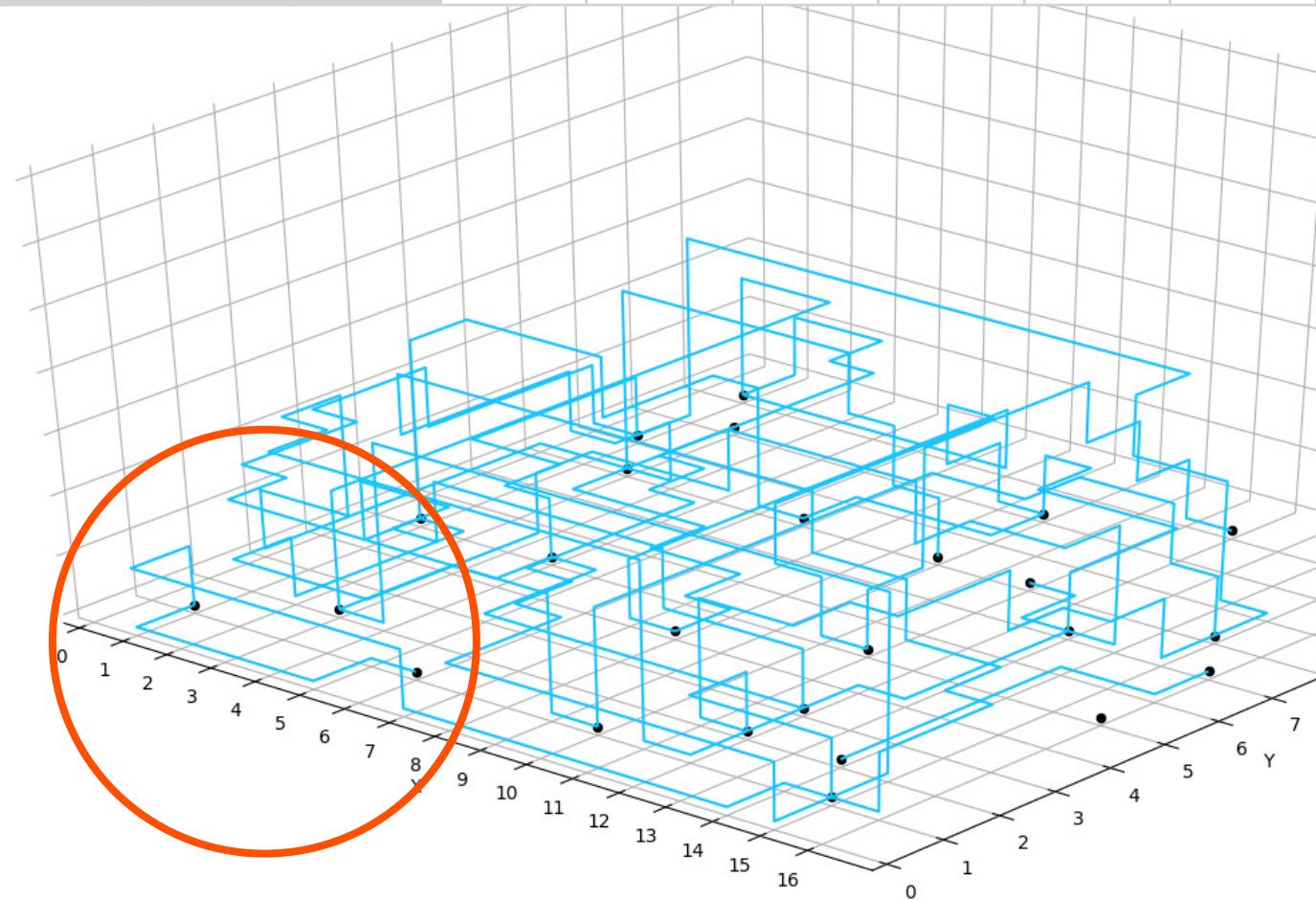


Netlist: [[A, B], [E, F], [C, D]]



# Optimalisatie

A\*



# Resultaten

Alleen A\*

500 permutations

Netlist	Circuit board 1			Circuit board 2		
	1	2	3	4	5	6
Number of nets laid down	27/30	32/40	37/50	40/50	45/60	47/70

# Resultaten

*A\* + Heatmap*

500 permutations

Heat constant:

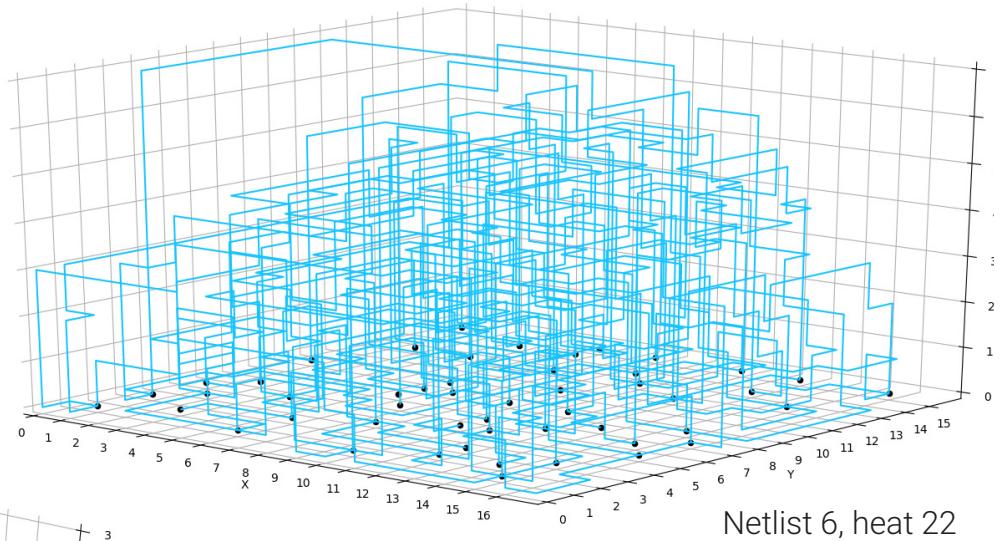
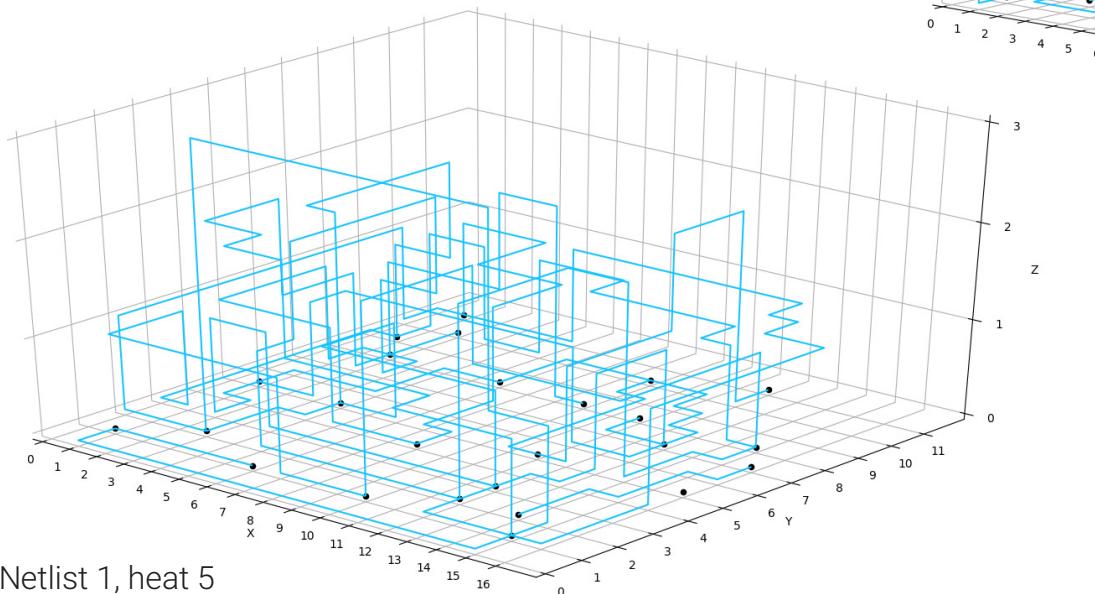
5-30

totaal: 12500 iteraties  
per netlist

Netlist	Circuit board 1			Circuit board 2		
	1	2	3	4	5	6
Number of nets laid down	30/30	40/40	50/50	50/50	60/60	70/70
Heat constant	5	19	14	29	18	22
Total length before optimisation	409	765	1003	1096	1416	1615
Total length after optimisation	<b>359</b>	<b>455</b>	<b>743</b>	<b>818</b>	<b>910</b>	<b>1341</b>
Decrease in total length after optimisation	12,2%	40,5%	25,9%	25,4%	35,7%	17,0%
Lower bound total length	291	341	475	600	578	761
Percentage of optimised total length exceeding the lower bound	23,4%	33,4%	56,4%	36,3%	57,4%	76,2%

# Resultaten

## 3D Plot



# Conclusie

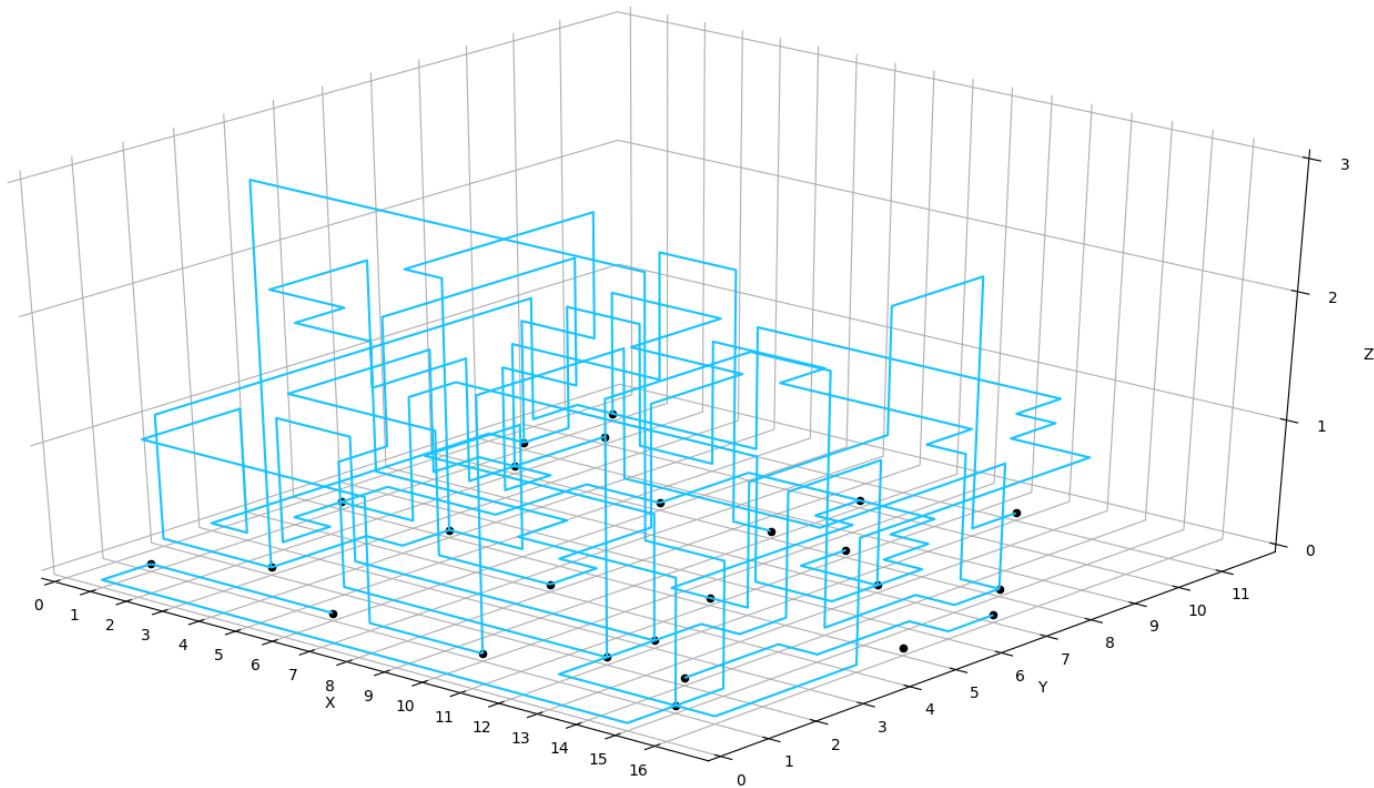
- Alle netlists zijn oplosbaar
- Heatmap was essentieel
- Geen universele optimale heat constant
- Volgorde van nets in de netlist cruciaal in deze methode

# Discussie

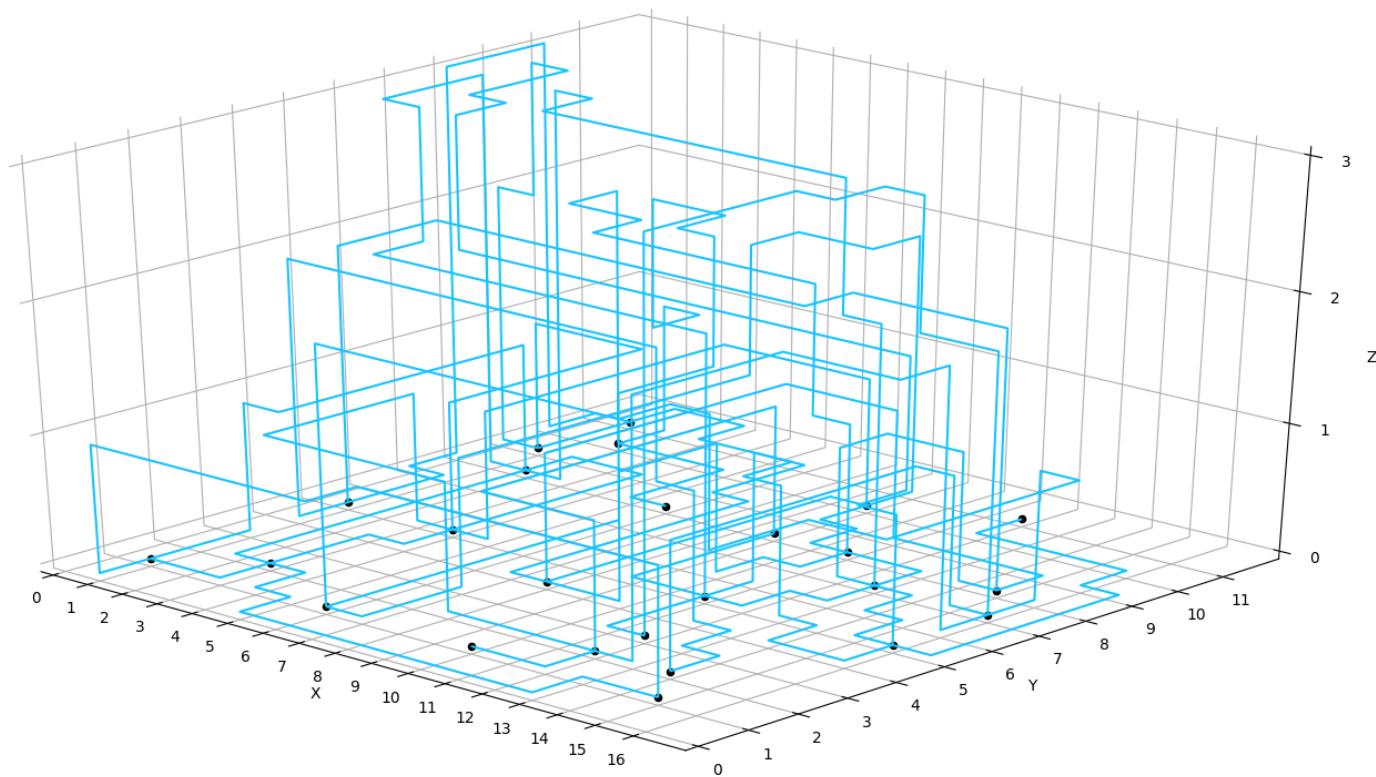
- Optimaal?  
*Lower bound,*  
*Random permutations: alle permutaties ( $n!$ )*
- 6<sup>e</sup> meest complex
- Vervolgonderzoek:  
*permutations + effectievere Heatmaps*  
*+ volgorde optimalisatie*

Vragen?

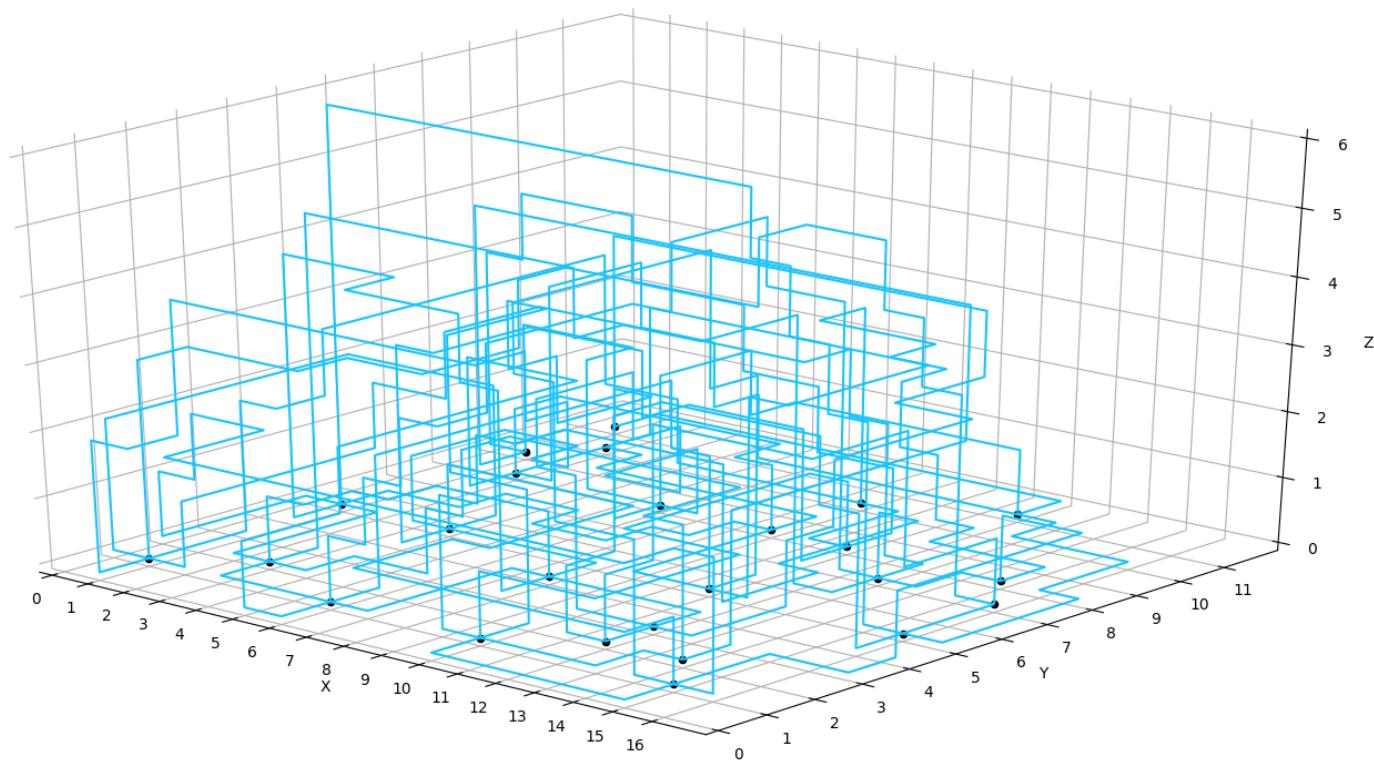
# *Netlist 1*



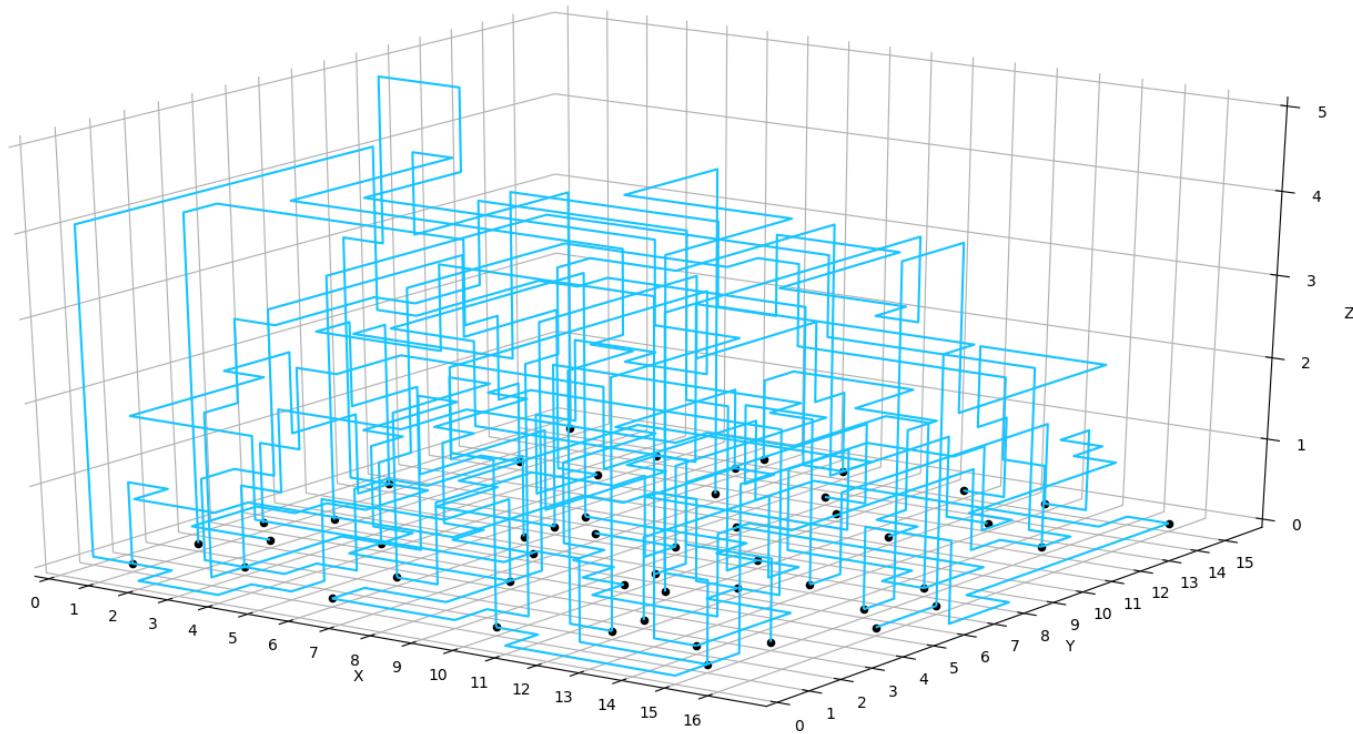
## Netlist 2



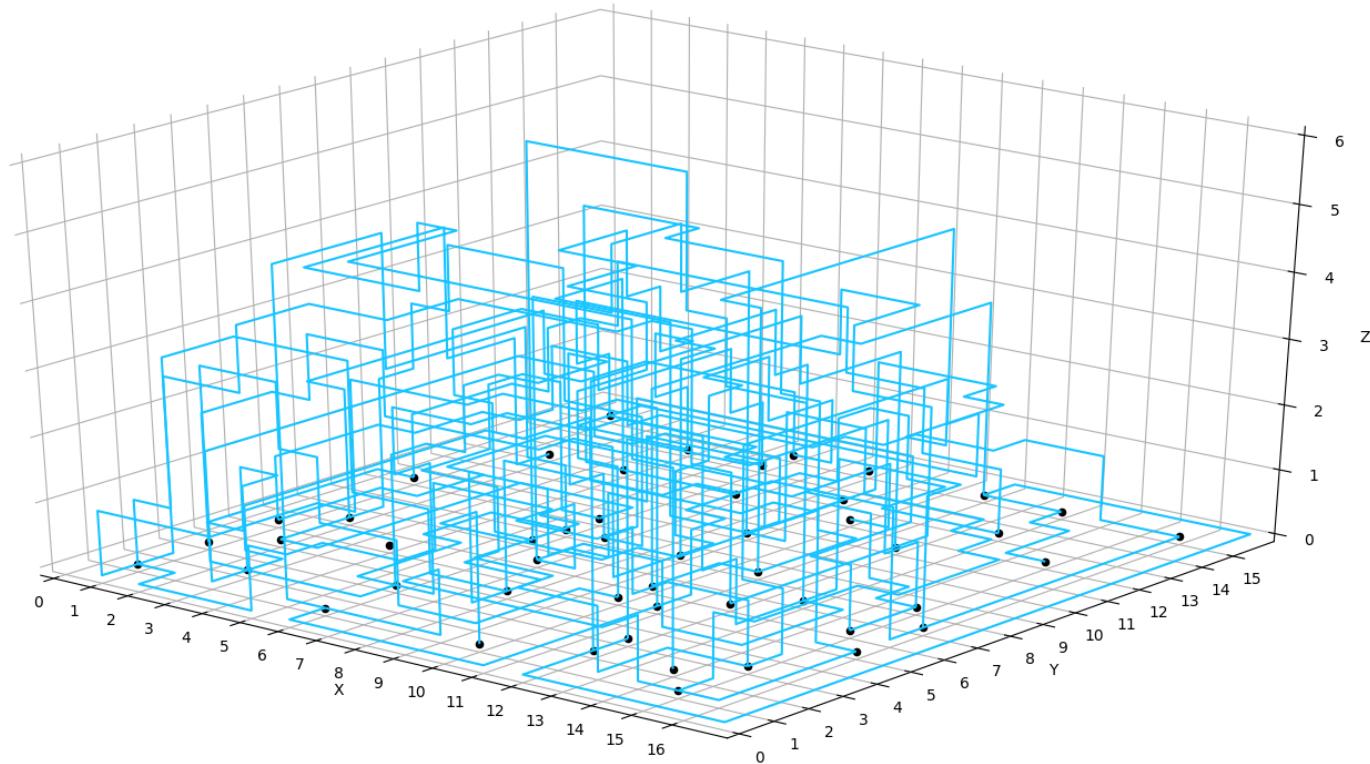
## *Netlist 3*



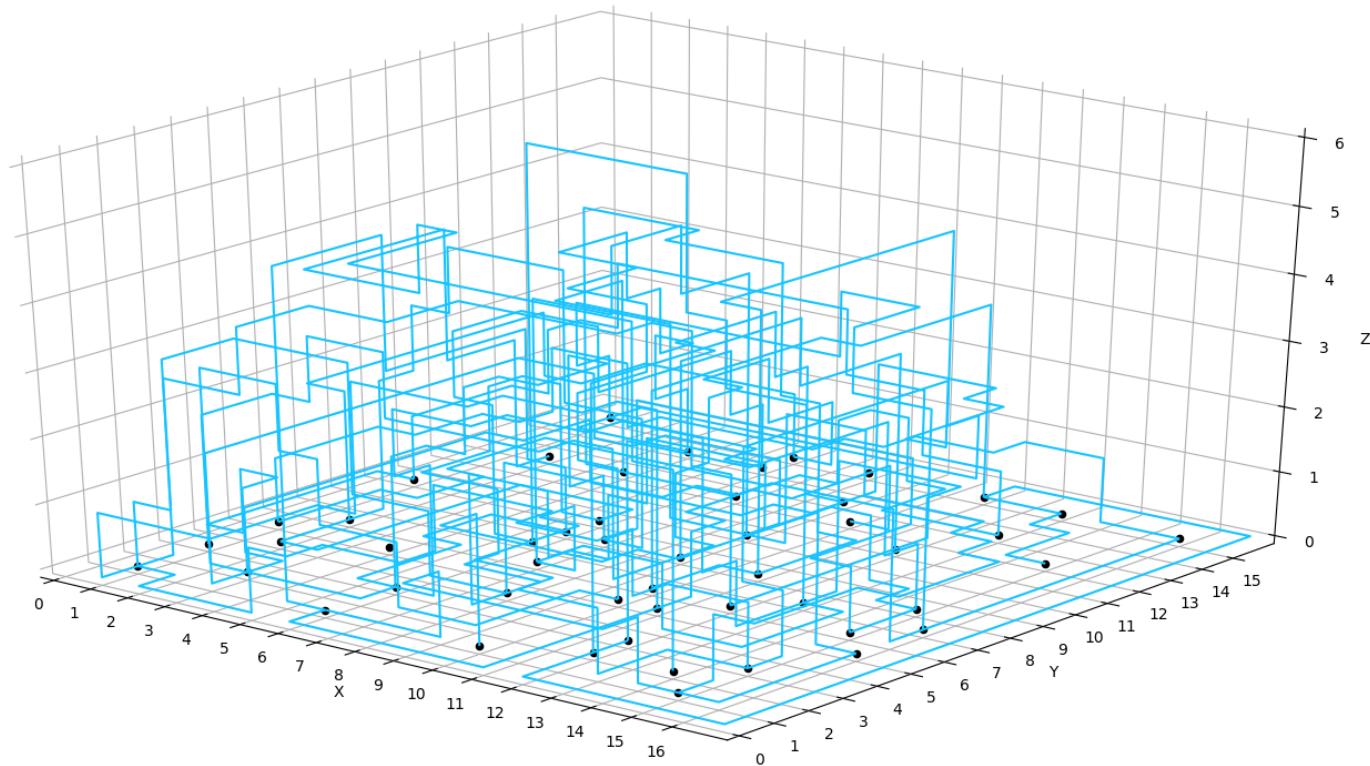
## Netlist 4



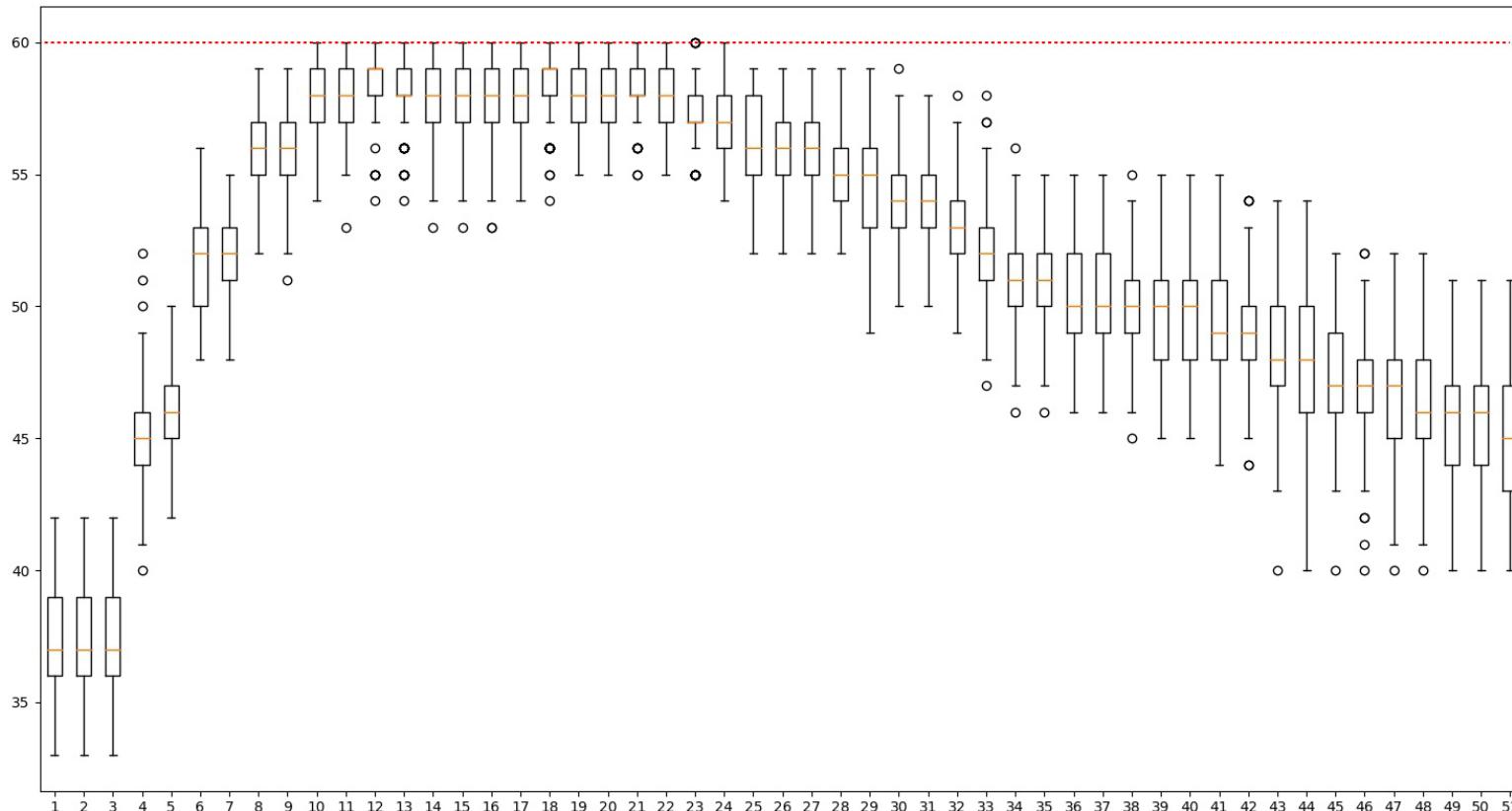
## Netlist 5



## Netlist 6



## *Netlist 5, 100 iteraties*



netlists/netlist\_1.txt

