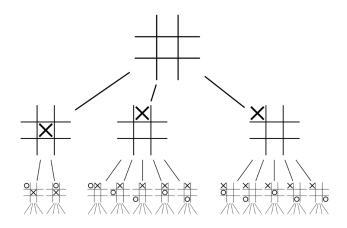
COMP250: Artificial Intelligence

#### 5: Game Tree Search

Minimax search

#### Game trees



#### **Minimax**

- ► Terminal game states have a value
  - ▶ E.g. +1 for a win, -1 for a loss, 0 for a draw
- ▶ I want to **maximise** the value
- ▶ My opponent wants to **minimise** the value
- Therefore I want to maximise the minimum value my opponent can achieve
- This is generally only true for two-player zero-sum games

#### Minimax search

- Recursively defines a value for non-terminal game states
- Consider each possible "next state", i.e. each possible move
- If it's my turn, the value is the maximum value over next states
- If it's my opponent's turn, the value is the minimum value over next states

# Minimax search – example

#### Minimax search pseudocode

```
procedure MINIMAX(state, currentPlayer)
   if state is terminal then
      return value of state
   else if currentPlayer = 1 then
      bestValue = -\infty
      for each possible nextState do
         v = MINIMAX(nextState, 3 - currentPlayer)
         bestValue = Max(bestValue, v)
      return bestValue
   else if currentPlayer = 2 then
      bestValue = +\infty
      for each possible nextState do
         v = MINIMAX(nextState, 3 - currentPlayer)
         bestValue = Min(bestValue, v)
      return bestValue
```

#### Stopping early

```
for each possible nextState do
  v = MINIMAX(nextState, 3- currentPlayer)
  bestValue = MAX(bestValue, v)
```

- $\blacktriangleright$  State values are always between -1 and +1
- ► So if we ever have bestValue = 1, we can stop early
- ► Similarly when minimising if bestValue = -1

#### Using minimax search

- To decide what move to play next...
- ► Calculate the minimax value for each move
- Choose the move with the maximum score
- If there are several with the same score, choose one at random

#### Minimax and game theory

- For a two-player zero-sum game with perfect information and sequential moves
- ► Minimax search will always find a Nash equilibrium
- I.e. a minimax player plays perfectly
- ▶ But...

#### Minimax for larger games

- ► The game tree for noughts and crosses has only a few thousand states
- Most games are too large to search fully
  - ► Connect 4 has  $\approx 10^{13}$  states
  - ► Chess has  $\approx 10^{47}$  states

### **Heuristics for search**

#### Depth limiting

- Standard minimax needs to search all the way to terminal (game over) states
- Depth limiting is a common technique to apply minimax to larger games
- $\blacktriangleright$  Still evaluate terminal states as +1/0/-1
- ► For nonterminal states at depth *d*, apply a heuristic evaluation instead of searching deeper
- ► Evaluation is a number between -1 and +1, estimating the probable outcome of the game

#### 1-ply search

- ► Case d = 1
- ► For each move, evaluate the state resulting from playing that move
- ► This is computationally fast
- Often easier to design a "which state is better" heuristic than to directly design a "which move to play" heuristic

#### Move ordering

- ► Minimax can stop early if it sees a value of +1 for maximising player or -1 for minimising player
- Modifications to minimax algorithm (e.g. alpha-beta pruning) lead to more of this
- Thus ordering moves from best to worst means faster search
- How do we know which moves are "best" and "worst"? Use a heuristic!

#### Designing heuristics

- The playing strength of depth limited minimax depends heavily on the design of the heuristic
- Good heuristic design requires in-depth knowledge of the tactics and strategy of the game
- Next time we will look at what we can do if we don't possess such knowledge

# Monte Carlo evaluation

#### From last time

- It is useful to have a heuristic evaluation function for nonterminal states
- ► Allows 1-ply search, depth-limited minimax, ...
- Designing a good heuristic requires in-depth knowledge of the game
- ➤ What if you don't have such knowledge?

#### Expected value

- ► Let X be a random variable
- ▶ Let p(x) be the probability that X has value x
- ► Then the **expected value** of X is

$$\sum_{x} x \cdot p(x)$$

#### Expected value — example

- ► A slot machine pays out:
  - ▶ £1 with probability 0.05
  - ▶ £5 with probability 0.03
  - £10 with probability 0.02
  - Nothing with probability 0.9
- ► The expected payout is

$$1 \times 0.05 + 5 \times 0.03 + 10 \times 0.02 + 0 \times 0.9 = 0.4$$

i.e. £0.40

▶ What this means: if you play the slot machine N times, on average you will win  $N \times £0.40$ 

### "Randomness" in computing

- Digital computers are deterministic, so there's no such thing as true randomness
  - Cryptographically secure systems use an external source of randomness e.g. atmospheric noise, radioactive decay
- What we actually have are pseudo-random number generators (PRNGs)
- A PRNG is an algorithm which gives an unpredictable sequence of numbers based on a seed
- Sequence is uniformly distributed, i.e. all numbers have equal probability
- Seed is generally based on some source of entropy,
   e.g. system clock, mouse input, electronic noise

#### Monte Carlo methods

- In computing, a Monte Carlo method is an algorithm based on averaging over random samples
- ► The average over a large number of samples is a good approximation of the expected value
- Used for quickly approximating quantities over large domains
- Generally designed to converge in the limit
  - An infinite number of samples would give an exact answer
  - As the number of samples increases, the accuracy of the answer improves
- Applications in physics, engineering, finance, weather forecasting, graphics, ...

#### Monte Carlo evaluation in games

Based on random rollouts

while s is not terminal do
let m be a random legal move from s
update s by playing m

- ► The value of a rollout is the value of the terminal state it reaches (i.e. 1 for a win, -1 for a loss, 0 for a draw)
- Averaging gives the expected value of the initial state
- ► Higher expected value = more chance of winning

#### Monte Carlo search

- ► Flat Monte Carlo search: 1-ply search with Monte Carlo evaluation
- ► How about minimax with d > 1 and Monte Carlo evaluation?
  - Minimax assumes the evaluation is deterministic, but Monte Carlo is not
  - Not commonly used, mainly because there's something better...

### **Monte Carlo Tree Search**

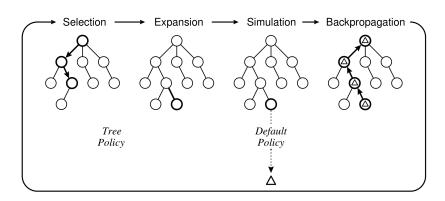
#### Monte Carlo Tree Search (MCTS)

- ► Like Monte Carlo evaluation, based on **rollouts**
- ► First few rollouts are random
- However, statistics from these rollouts are used to bias future rollouts
- Bias rollouts towards plausible lines of play, i.e. where each player is trying to play the best move

#### The MCTS algorithm

- MCTS builds a tree
- ▶ Initially, the tree consists of a single root node
- ► Each rollout has four stages:
  - Selection: Starting from the root, descend the tree by choosing moves. Continue until we reach a node which does not yet have children for all legal moves.
  - ► **Expansion**: Choose a random legal move for which the current node does not have a child node. Add this new node to the tree.
  - Simulation: Perform a Monte Carlo rollout, playing random moves until a terminal state is reached.
  - Backpropagation: For each node visited during selection and expansion, update the node's statistics based on the result of the simulation.
- Perform many rollouts, then use the statistics at the top level of the tree to choose the best move

### The MCTS algorithm



#### Selection policy

- Selection must balance:
  - Exploitation of moves that are known to be good
  - Exploration of moves that have not often been tried
- This can be modelled as a multi-armed bandit problem

#### Multi-armed bandits

- We have a row of one-armed bandits (slot machines)
- We do not know the payout probabilities of any of them, and they're all different
- ► How to maximise our winnings?
- Again must balance
  - Exploitation of machines that are known to have a high expected payout
  - Exploration of machines that have not been tried often, to get a better estimate of their expected payout

#### Upper Confidence Bound (UCB)

- ► For each machine *m*, record:
  - $\triangleright$   $n_m$ : the number of plays of this machine
  - $ightharpoonup V_m$ : the total winnings from playing this machine
  - $n = \sum_{m} n_m$ , total number of plays across all machines
- At each stage, play the machine for which

$$\frac{V_m}{n_m} + c\sqrt{\frac{\log n}{n_m}}$$

#### is largest

- ▶  $\frac{V_m}{n_m}$  is the **exploitation** part: average payout from this machine so far
- $\sqrt{\frac{\log n}{n_m}}$  is the **exploration** part: large if  $n_m$  is small
- c is a parameter for adjusting the balance between exploitation and exploration

#### UCB demo

# Upper Confidence Bound for Trees (UCT)

- Use UCB as the selection policy
- ▶ In each node x, record:
  - $\triangleright$   $n_x$ : the number of visits to this node
  - $\triangleright$   $V_x$ : the total value of rollouts through this node
- ▶ From node p, choose the child q such that

$$\frac{V_q}{n_q} + c\sqrt{\frac{\log n_p}{n_q}}$$

is largest

## UCT demo

#### Benefits of MCTS

- "Vanilla" MCTS is game independent
- But if game-specific heuristics are available, they can be used to enhance MCTS
- ► MCTS is anytime
  - Can stop it after any amount of computation (within reason) and get a reasonably good answer
  - ▶ Compare with minimax:  $O(e^d)$  for depth d
- ► Does not suffer from horizon effect
  - Minimax at depth d cannot "see" what happens d+1 moves in the future
  - MCTS can build the tree as deep as it likes
  - Selects which parts of the tree to expand more deeply

### Workshop

#### Workshop

- ► Clone https://github.com/ Falmouth-Games-Academy/comp250-workshop-5
- ► This is an implementation of Connect-4 in Python 3 with PyGame
- ► Edit ai\_player.py to implement MCTS
- ▶ Use http://mcts.ai/code/python.html as a guide