

COMP250: Artificial Intelligence

7: Navigation







Research journal

Research wiki check-in

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Pathfinding

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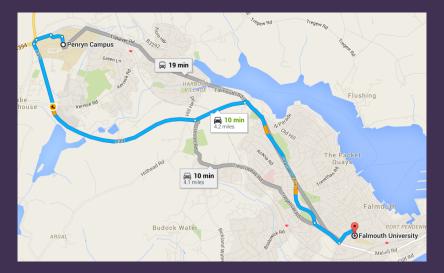
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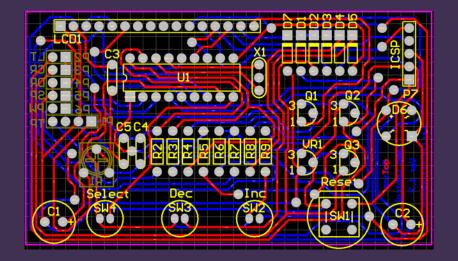
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 - "Shortest" in terms of edge lengths could be distance, time, fuel cost, ...









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Pathfinding example

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- ▶ Open COMP250/07_pathfinding in PyCharm

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 - Often implemented with a priority queue

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- ... but is not the most efficient algorithm for doing so

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 - Heuristics are often used to prioritise search, i.e. explore the most promising options first



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- Different h(x) can lead to different paths (if there are multiple "shortest" paths)

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 - Repeat until there are no more points that can be removed







Pathfinding in videogames

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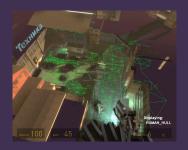
Pathfinding in videogames

- ► A* works on any graph
- But what if the game world is not a graph? E.g. complex 3D environments





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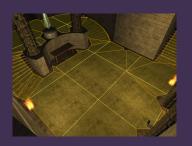
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 Automatically generate navigation graph from level geometry



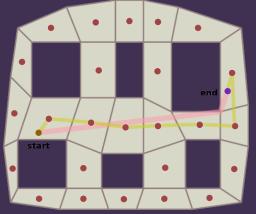
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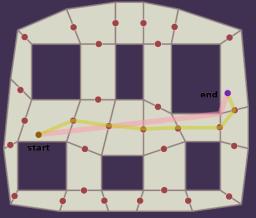
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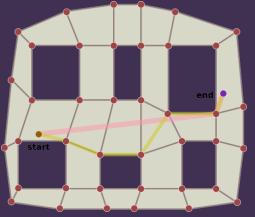
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 - Generate graph from polygons



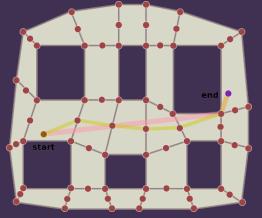
Centres of polygons



Centres of edges



Vertices of polygons



Hybrid approach: edges and vertices

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- Dynamic environments: may need to re-run pathfinder if environment changes (e.g. movable obstacles, destructible terrain)



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 - Pub crawls
 (http://www.math.uwaterloo.ca/tsp/pubs/)

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- Entire research field devoted to finding efficient search algorithms and heuristics