

# COMP250: Artificial Intelligence

## 7: Navigation

# Research journal



# Research wiki check-in

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# Pathfinding



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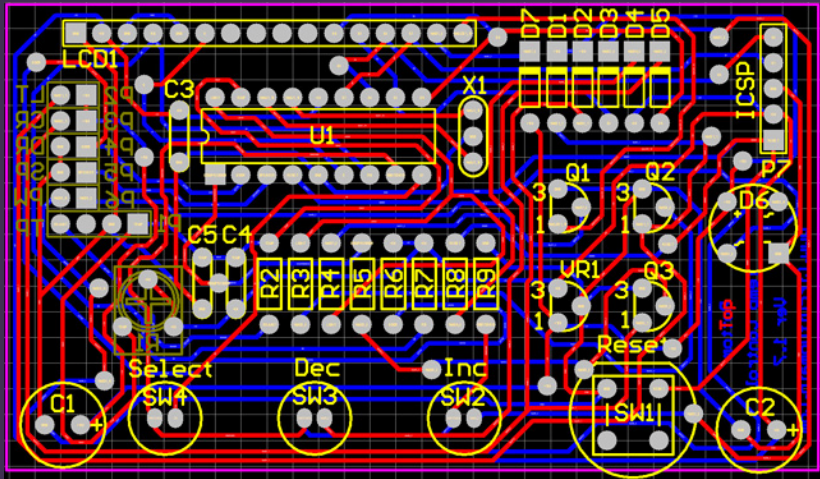
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  - ▶ “Shortest” in terms of edge lengths — could be distance, time, fuel cost, ...



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- ▶ Open 07\_pathfinding in PyCharm

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  - ▶ Often implemented with a **priority queue**



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- ▶ ... but is not the most efficient algorithm for doing so

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  - ▶ Heuristics are often used to prioritise search, i.e. explore the most promising options first

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- ▶ Different  $h(x)$  can lead to different paths (if there are multiple “shortest” paths)

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  - ▶ Repeat until there are no more points that can be removed

# Navigation meshes



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- ▶ A\* works on any **graph**
- ▶ But what if the game world is not a graph? E.g. complex 3D environments

# Waypoint navigation





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  - ▶ No good for dynamic environments

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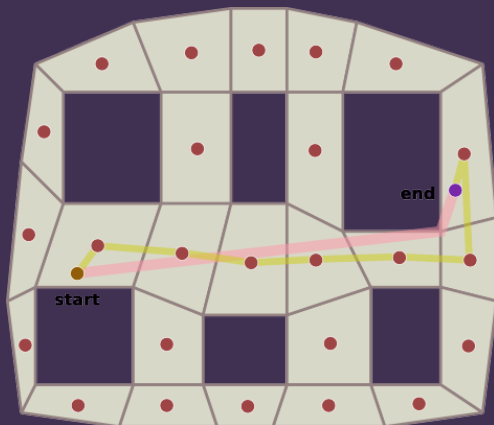
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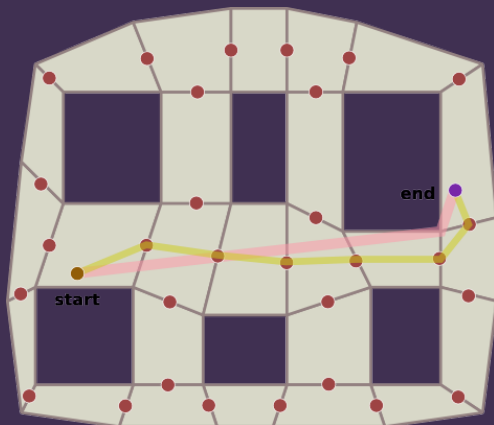
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# Meshes to graphs



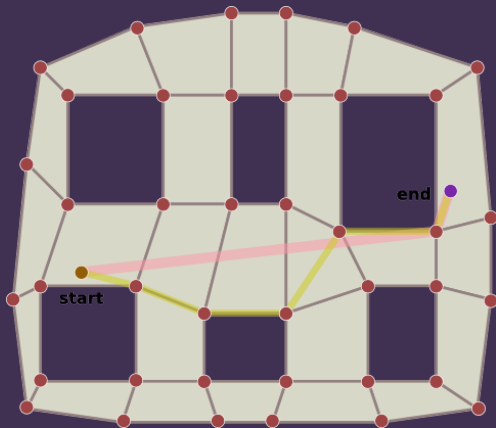
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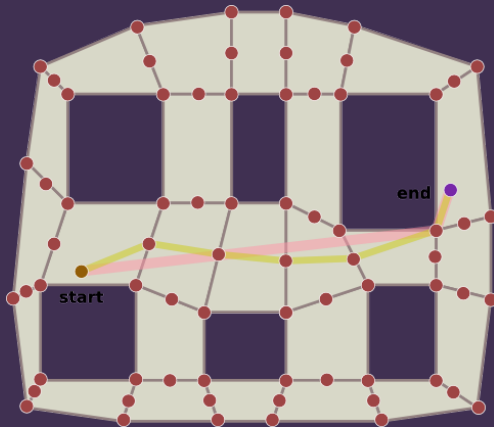
## Centres of edges

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Vertices of polygons

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Hybrid approach: edges and vertices

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- ▶ **Steering:** don't have your AI agent follow the path exactly, but instead try to stay close to it
- ▶ **Dynamic environments:** may need to re-run pathfinder if environment changes (e.g. movable obstacles, destructible terrain)

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    - ▶ Pub crawls
- (<http://www.math.uwaterloo.ca/tsp/pubs/>)

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- ▶ Entire research field devoted to finding efficient **search algorithms** and **heuristics**