

COMP110: Principles of Computing

7: Data structures



Learning outcomes

- Explain the difference between pass-by-value and pass-by-reference
- Distinguish the basic data structures available in Python
- Determine the complexity of accessing and manipulating data in these data structures
- Choose the correct data structure for a given task

Worksheet D

- ▶ Data structures
- ► Due in class on **Monday 7th November** (next week)





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- Our picture of a variable: a labelled box containing a value
- ▶ For "plain old data" (e.g. numbers), this is accurate
- For objects (i.e. instances of classes), variables actually hold references (a.k.a. pointers)
- It is possible (indeed common) to have multiple references to the same underlying object

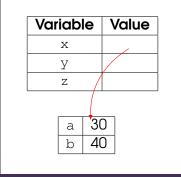
Variable	Value
X	
У	
Z	

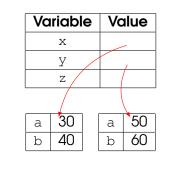
Variable	Vo	ilue
7.7	а	30
X	b	40
У		
Z		

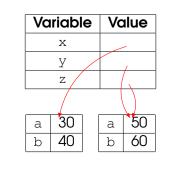
Variable	Value
Х	a 30
	b 40
У	a 50
	b 60
Z	

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```
a = 10
b = a
a = 20
print "a:", a
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class X:
    def __init__(self, value):
        self.value = value

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In **function parameters**, "plain old data" is passed by **value**

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def double(x):
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a = 7
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print a
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double does not actually do anything, as x is just a local copy of whatever is passed in!

However, instances are passed by reference

```
class Box:
    def __init__(self, v):
        self.value = v

def double(x):
        x.value *= 2

a = Box(7)
double(a)
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    def __init__(self, v):
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def double(x):
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double(a)
print a.value
```

double now has an effect, as x gets a reference to the Box instance

Lists are objects too

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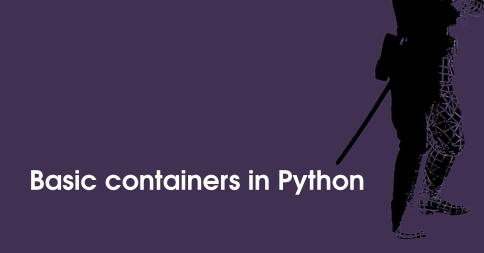
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a = ["Hello"]
b = a
b.append("world")
print a # ["Hello", "world"]
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... which means you should be careful when passing lists into functions, because the function might actually change the list!





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- Blocks can be allocated and deallocated at will, but can never grow or shrink



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- Memory management is hard and programmers are lazy
- Containers are an abstraction
 - Hide the details of memory allocation, and allow the programmer to write simpler code
- ► Containers are an **encapsulation**
 - Bundle together the data's representation in memory along with the algorithms for accessing it



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- ▶ E.g. if the array starts at address 1000 and each element is 4 bytes, the 3rd element is at address $1000 + 4 \times 3 = 1012$
- ▶ Accessing an array element is constant time O(1)





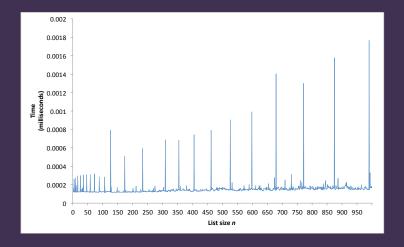
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- When the list needs to change size, it creates a new array, copies the contents of the old array, and deletes the old array
- ► Implementation details: http://www.laurentluce. com/posts/python-list-implementation/

Time taken to append an element to a list of size *n*



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 - Can't just insert new bytes into a memory block need to move all subsequent list elements to make room
- Similarly, deleting anything other than the last element is linear time

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- Create tuples with (), just as you create lists with []
 - Exception: a single element tuple is created as (foo,) because (foo) would be interpreted as a bracketed expression
- Can often omit the parentheses entirely, e.g.

```
my_tuple = 1,2,3
```

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```
a, b, c, d = foo
```

```
a = foo[0]
b = foo[1]
c = foo[2]
d = foo[3]
```

 Unpacking requires the number of elements to match exactly — if foo has more than 4 elements, the code on the left will give an error

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$$a, b = b, a$$

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- This isn't changing the string, it's creating a new one and throwing the old one away!
- ► Hence building a long string by appending can be slow (appending strings is O(n))

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- A dictionary is implemented as a hash table (see Session 5)

Using dictionaries

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Create them using {}:

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age = {"Alice": 23, "Bob": 36, "Charlie": 27}
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Access values using []:

```
print age["Alice"] # prints 23
age["Bob"] = 40  # overwriting an existing item
age["Denise"] = 21  # adding a new item
```

Iterating over dictionaries

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Use iteritems to get key, value pairs:

```
for key, value in age.iteritems():
    print key, "is", age, "years old"
```

Dictionaries are unordered

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What does this print?

```
square_root = {}
for i in xrange(30):
    square_root[i*i] = i

for key, value in square_root.iteritems():
    print "The square root of", key, "is", value
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Dictionaries are **unordered** — never rely on the order of their elements, because the order isn't guaranteed!

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Operation	List	Set
Add element	Append: <i>O</i> (1)	<i>O</i> (1)
	Insert: O(n)	
Delete element	O(n)	<i>O</i> (1)
Contains element?	O(n)	0(1)





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- ► E.g. w = 5, h = 4:

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 - Each element of the "outer" list represents a column of the array
- The element in column x row y is accessed by list[x][y], i.e. the yth element of the xth column

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- ► However, highly optimised for intensive calculations (e.g. "tinkering" with image pixel colours...?)

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There is no single "best" approach — it depends how you use it







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- Items can be dequeued from the front of the queue

Stacks and queues in Python

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Stacks can be implemented efficiently as lists

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- Stacks can be implemented efficiently as lists
- Queues can be implemented as lists, but not efficiently

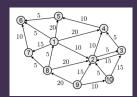


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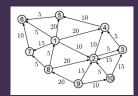
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- Stacks can be implemented efficiently as lists
- Queues can be implemented as lists, but not efficiently
- deque (from the collections module) implements an efficient double-ended queue
- ► Inserting and removing elements from the start and end of a deque is O(1)



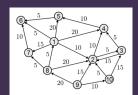




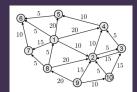




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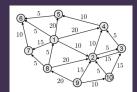


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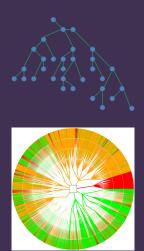
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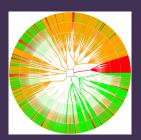




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- Often used to model networks

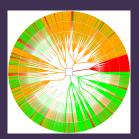
 (e.g. social networks, transport networks, game levels, finite state automata, ...)





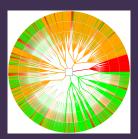
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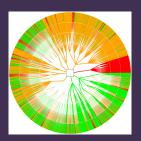
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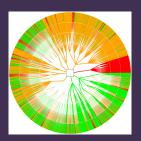
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 - Nodes can have 0, 1 or many children
- Used to model hierarchies (e.g. file systems, object inheritance, scene graphs, state-action trees, ...)

"Smart data structures and dumb code works a lot better than the other way around." — Eric S. Raymond