



COMP702: Classical Artificial Intelligence

7: Navigation

Paper Club

For next week's session:

Nathan R. Sturtevant, Devon Sigurdson, Bjorn Taylor, Tim Gibson. Pathfinding and Abstraction with Dynamic Terrain Costs. Proceedings of AllDE Conference, 2019. (PDF link on LearningSpace)





Pathfinding

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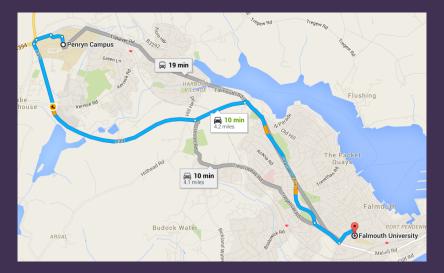
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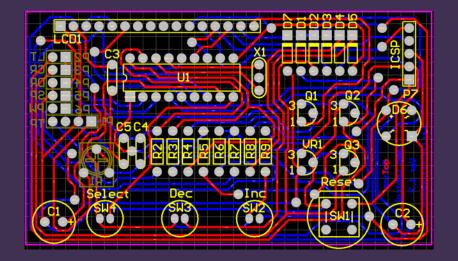
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 - "Shortest" in terms of edge lengths could be distance, time, fuel cost, ...























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 - Closed set: nodes which have been added to the tree, and shouldn't be revisited (otherwise we could get stuck in an infinite loop)

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 - "Best" according to some heuristic evaluation
 - Often implemented with the open set as a priority queue — a data structure optimised for finding the highest priority item

Greedy search

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- ▶ Not guaranteed to find the **shortest** path

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- Is guaranteed to find the shortest path
- ... but is not the most efficient algorithm for doing so

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- ▶ Choose a node that minimises g(x) + h(x)

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 - Heuristics are often used to prioritise search, i.e. explore the most promising options first



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- Different h(x) can lead to different paths (if there are multiple "shortest" paths)

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 - Repeat until there are no more points that can be removed







Pathfinding in videogames

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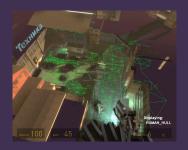
Pathfinding in videogames

- ► A* works on any graph
- But what if the game world is not a graph? E.g. complex 3D environments





 Manually place graph nodes in the world



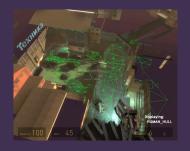
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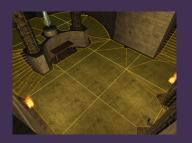
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 Automatically generate navigation graph from level geometry



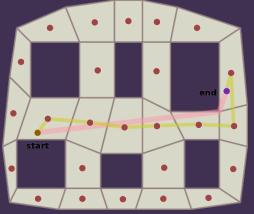
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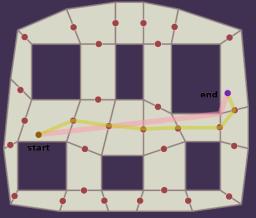
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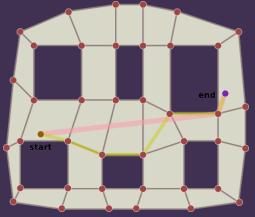
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 - Generate graph from polygons



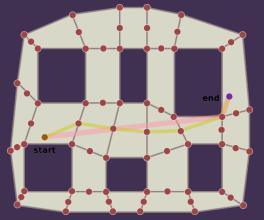
Centres of polygons



Centres of edges



Vertices of polygons



Hybrid approach: edges and vertices

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- ➤ **Dynamic environments**: may need to re-run pathfinder if environment changes (e.g. movable obstacles, destructible terrain)