中国科学技术大学计算机学院

计算机网络实验报告

实验三 利用 Wireshark 观察 TCP 报文

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一、实验目的

1、 通过捕获观察并分析 TCP 报文,理解 TCP 的细节,包括:为了 reliable 传输的 SEQ、ACK 序号使用; TCP 的拥塞控制算法-慢启动和拥塞避免; TCP 的流控制机制; TCP 连接的建立。

二、实验原理

Wireshark 是一个 packet 分析工具,可以抓取 packet,并分析出详细信息。Wireshark 使用 wincap 作为接口,直接与网卡进行 packet 交换,监听共享网络上传送的 packet。

三、实验条件

1、 硬件条件: 联想拯救者 Y7000:

i5-8300H 2.30GHz

16G 内存

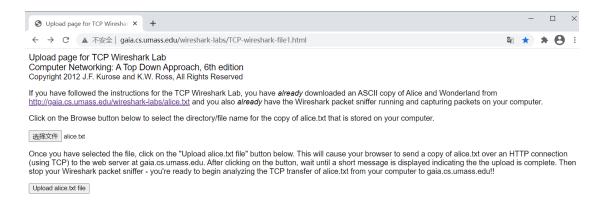
Intel UHD Graphics 630

2、 软件条件: Win10

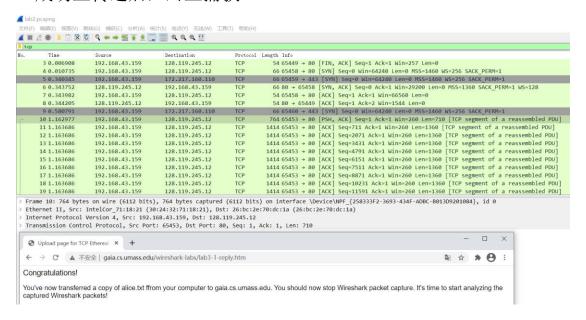
Wireshark3.4.0

四、实验过程

- 1、 向远程服务器发送一个 txt 文件,并捕获 TCP 报文。
- 首先下载 alice. txt
- 到指定页面选中要上传的文件。



- 打开 wireshark 开始捕获。
- 开始上传文件。
- 成功上传之后,终止捕获。



2、 开始回答问题。

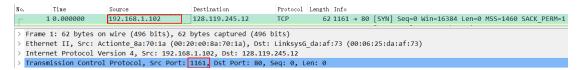
五、 结果分析

以下是 pdf 中 14 个问题对应的回答

(除 3、14 题,都是用的提供的 trace 文件。)

1. What is the IP address and TCP port number used by the client computer (source) that is transferring the file to gaia.cs.umass.edu? To answer this question, it's probably easiest to select an HTTP message and explore the details of the TCP packet used to carry this HTTP message, using the "details of the selected packet header window" (refer to Figure 2 in the "Getting Started with Wireshark" Lab if you're uncertain about the Wireshark windows).

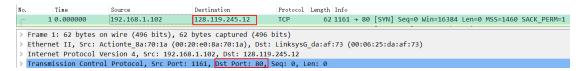
答: 客户端电脑的 ip 地址: 192.168.1.102; TCP 端口: 1161。如下图:



2. What is the IP address of gaia.cs.umass.edu? On what port number is it sending and receiving TCP segments for this connection?

答: IP地址为: 128.119.245.12; 端口号: 80。

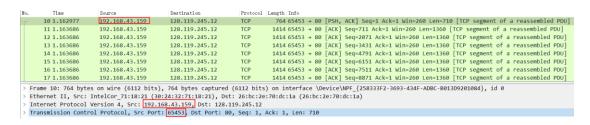
图:



3. What is the IP address and TCP port number used by your client computer (source) to transfer the file to gaia.cs.umass.edu?

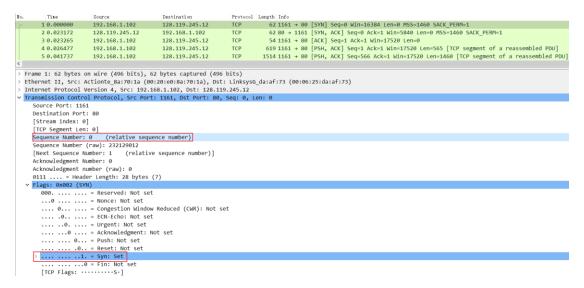
答: 自己捕获一遍之后,得到下图。我自己电脑的 IP 地址为:

192.168.43.159, 端口号: 65453

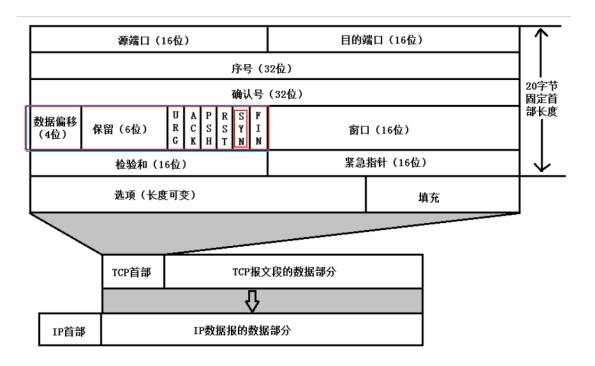


4. What is the sequence number of the TCP SYN segment that is used to initiate the TCP connection between the client computer and gaia.cs.umass.edu? What is it in the segment that identifies the segment as a SYN segment?

答:如下图,SYN的Seq序号为0。



确认报文为 SYN 报文的标志是报文的 TCP HEADER 中的 flag field 中被置为 1 的 SYN 标志。



5. What is the sequence number of the SYNACK segment sent by gaia.cs.umass.edu

to the client computer in reply to the SYN? What is the value of the ACKnowledgement field in the SYNACK segment? How did gaia.cs.umass.edu determine that value? What is it in the segment that identifies the segment as a SYNACK segment?

答: 如下图。

序号为0。

Acknowledgement field 为 1。

gaia.cs.umass.edu 将该值设置为所期望的下一个来自客户端的报文的 Sequence Number。

Flag field 中被置为 1 的 ACK 位和 SYN 位。

6. What is the sequence number of the TCP segment containing the HTTP POST command? Note that in order to find the POST command, you'll need to dig into the packet content field at the bottom of the Wireshark window, looking for a segment with a "POST" within its DATA field.

答: 如下图, 1。

7. Consider the TCP segment containing the HTTP POST as the first segment in the TCP connection. What are the sequence numbers of the first six segments in the TCP connection (including the segment containing the HTTP POST)? At what

time was each segment sent? When was the ACK for each segment received? Given the difference between when each TCP segment was sent, and when its acknowledgement was received, what is the RTT value for each of the six segments? What is the <code>EstimatedRTT</code> value (see page 249 in text) after the receipt of each ACK? Assume that the value of the <code>EstimatedRTT</code> is equal to the measured RTT for the first segment, and then is computed using the <code>EstimatedRTT</code> equation on page 249 for all subsequent segments.

(注: 在第七题时,已经更改时间格式为日期和时间)

答:根据下图,再根据等式:

EstimatedRTT = $(1 - \alpha)$ · EstimatedRTT + α · SampleRTT

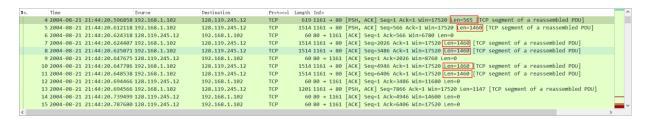
No.	Time Source	Destination	Protocol	Length Info	
	4 2004-08-21 21:44:20.596858 192.168.1.102	128.119.245.12	TCP	619 1161 → 80 [PSH, ACK] Seq=1 Ack=1 Win=17520 Len=565 [TCP segment of a reassembled PDU]	
	5 2004-08-21 21:44:20.612118 192.168.1.102	128.119.245.12	TCP	1514 1161 → 80 [PSH, ACK] Seq=566 Ack=1 Win=17520 Len=1460 [TCP segment of a reassembled PDU]	
	6 2004-08-21 21:44:20.624318 128.119.245.12	192.168.1.102	TCP	60 80 → 1161 [ACK] Seq=1 Ack=566 Win=6780 Len=0	
	7 2004-08-21 21:44:20.624407 192.168.1.102	128.119.245.12	TCP	1514 1161 → 80 [ACK] Seq=2026 Ack=1 Win=17520 Len=1460 [TCP segment of a reassembled PDU]	
	8 2004-08-21 21:44:20.625071 192.168.1.102	128.119.245.12	TCP	1514 1161 → 80 [ACK] Seq=3486 Ack=1 Win=17520 Len=1460 [TCP segment of a reassembled PDU]	
	9 2004-08-21 21:44:20.647675 128.119.245.12	192.168.1.102	TCP	60 80 → 1161 [ACK] Seq=1 Ack=2026 Win=8760 Len=0	
	10 2004-08-21 21:44:20.647786 192.168.1.102	128.119.245.12	TCP	1514 1161 → 80 [ACK] Seq=4946 Ack=1 Win=17520 Len=1460 [TCP segment of a reassembled PDU]	
	11 2004-08-21 21:44:20.648538 192.168.1.102	128.119.245.12	TCP	1514 1161 → 80 [ACK] Seq=6406 Ack=1 Win=17520 Len=1460 [TCP segment of a reassembled PDU]	
	12 2004-08-21 21:44:20.694466 128.119.245.12	192.168.1.102	TCP	60 80 → 1161 [ACK] Seq=1 Ack=3486 Win=11680 Len=0	
	13 2004-08-21 21:44:20.694566 192.168.1.102	128.119.245.12	TCP	1201 1161 → 80 [PSH, ACK] Seq=7866 Ack=1 Win=17520 Len=1147 [TCP segment of a reassembled PDU]	
	14 2004-08-21 21:44:20.739499 128.119.245.12	192.168.1.102	TCP	60 80 → 1161 [ACK] Seq=1 Ack=4946 Win=14600 Len=0	-
	15 2004-08-21 21:44:20.787680 128.119.245.12	192.168.1.102	TCP	60 80 → 1161 [ACK] Seq=1 Ack=6406 Win=17520 Len=0	

(令^α为 0.125,) 得到下表:

i	seq	发送时间	收到ACK时间	RTT(ms)	EstimatedRTT(ms)
1	1	21:44:20.596858	21:44:20.624318	27.5	27.5
2	566	21:44:20.612118	21:44:20.647675	35.6	28.5125
3	2026	21:44:20.624407	21:44:20.694466	70.1	33.7109375
4	3486	21:44:20.625071	21:44:20.737499	114.4	43.79707031
5	4946	21:44:20.647786	21:44:20.787680	139.9	55.80993652
6	6406	21:44:20.648538	21:44:20.838183	189.7	72.54619446

8. What is the length of each of the first six TCP segments?

答: 分别是565,1460,1460,1460,1460,1460。如下图:



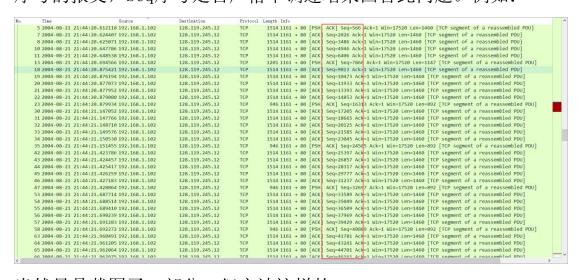
9. What is the minimum amount of available buffer space advertised at the received for the entire trace? Does the lack of receiver buffer space ever throttle the sender?

答:服务器最小的缓冲空间(Win)为5840字节。全程服务器端的Win在慢慢变大,最后变到62780字节,没有throttle发送端。



10. Are there any retransmitted segments in the trace file? What did you check for (in the trace) in order to answer this question?

答:没有重传的报文。通过检查发送端是否发送过两个具有相同Seq 序号的报文,Seq序号是否严格单调递增来回答此问题。例如:

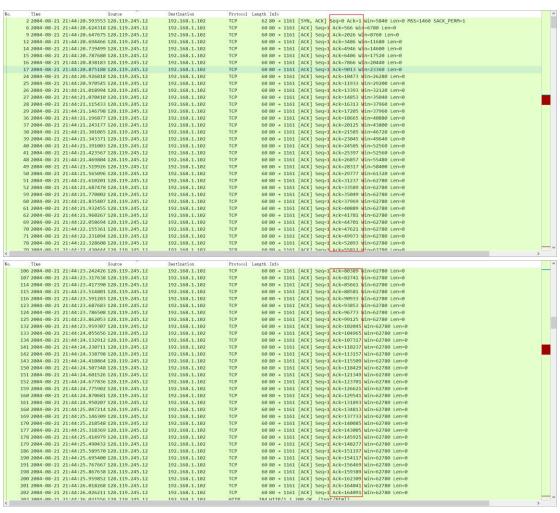


当然只是截图了一部分,但方法这样的。

11. How much data does the receiver typically acknowledge in an ACK? Can you identify cases where the receiver is ACKing every other received segment (see Table 3.2 on page 257 in the text).

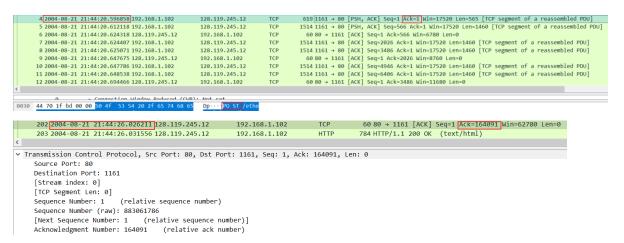
答: 典型的有1460(比如第24个报文和第26个报文以及第39个报文等等)和2920字节(比如第78个报文和第79个报文等等)。此外还

有50(比如第202个报文),566(比如第6个报文),1147(比如第17个报文),892(比如第41个报文),2352(892+1460)(比如第143个报文),2920(比如第60个报文)等等。一次ACK所acknowledge的字节数可以根据相邻Ack的差值算出。一次确认两个报文的有确认2920和2352字节的。比如trace文件中第52个报文,就是确认了2352字节;比如第60个报文,就是确认了2920字节。



12. What is the throughput (bytes transferred per unit time) for the TCP connection? Explain how you calculated this value.

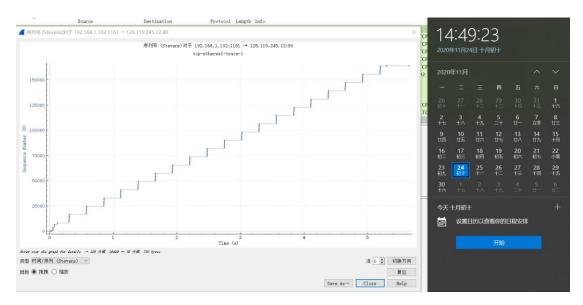
答: 考虑第一次发送post到发送端收到最后一条Ack这一过程。 总的时长t = 21:44:26.026211-21:44:20.596858=5.429353 (s)。 所传递的字节数w = 164091-1=164090 (Bytes)。 吞吐量为 w/t = 30222 Bps=30.222 KBps



13. Use the *Time-Sequence-Graph(Stevens)* plotting tool to view the sequence number versus time plot of segments being sent from the client to the gaia.cs.umass.edu server. Can you identify where TCP's slowstart phase begins and ends, and where congestion avoidance takes over? Comment on ways in which the measured data differs from the idealized behavior of TCP that we've studied in the text.

答: 慢启动在post发送开始时开始,但是看不出何时结束,也看不出拥塞避免何时开始。因为从这个trace可以看出没有出现3次冗余ACK或者超时,也就没有丢包,没有拥塞发生。不同点:

不同于课本上所说的先第一次慢启动,等到拥塞时再第二次进入 慢启动同时调整ssthresh,cwnd值,等到cwnd大于等于ssthresh时进 入拥塞避免这样的策略,我们的trace文件中没有冗余ACK,没有丢 包,没有拥塞,没有以上的状态转换。而是:数据发送的速率被严 格限制在了第一次慢启动结束前,并且到了后面一直重复这一过 程:连续发送1460*5+892*1字节=8192字节的包就会暂停发送,直到 这8192字节被全部ACK。这一数据传输过程,应该是被应用程序所 严格的控制着,而不仅仅依赖于TCP自身的拥塞控制机制。



14. Answer each of two questions above for the trace that you have gathered when you transferred a file from your computer to gaia.cs.umass.edu

答:同样,慢启动在post发送开始时开始,在0.425s左右结束,因为从0s附近到0.425s附近cwnd值翻了一番;同时,0.839s左右出现了三个冗余ACK,即慢启动结束,拥塞避免开始。2.052s左右出现了超时情况。

不同点:我这里的实际情况是,在慢启动开始时,cwnd不是从1MSS 开始,而是以大约9个MSS(MSS=1360Bytes)开始,之后再指数递增。慢启动遇到三个冗余ACK结束,并且进入拥塞避免状态,并不是像书上所说的每次只增加一个MSS,而是每次较上次的增加量增加一个MSS,比如19 22 26 31对应中间四段的MSS,后三次每次较上次的增加分别为3 4 5。遇到超时情况,也不是设置为1个MSS,而是从6个MSS开始。直到这里,便已发送完所有数据。这一数据传输过程,应该是应用程序和TCP自身的拥塞控制机制相互作用完成的。

