

```
CREATE TABLE InsuranceCo (
  name varchar(20) Primary key,
  phone INT,
);
CREATE TABLE Person (
  ssn INT Primary key,
  name varchar(20),
);
CREATE TABLE Driver (
  ssn INT Primary key references Person(ssn),
  driverID INT,
);
CREATE TABLE NonProfessionalDriver (
  ssn INT Primary key references Driver(ssn),
);
CREATE TABLE Professional Driver (
  ssn INT Primary key references Driver(ssn),
  medicalHistory varchar(200),
);
CREATE TABLE Vehicle (
  licensePlate varchar(20) Primary key,
  vear INT.
  name varchar(20) FOREIGN KEY references InsuranceCo(name),
  ssn INT FOREIGN KEY references Person(ssn),
  maxLiability REAL,
);
CREATE TABLE Car (
  licensePlate varchar(20) Primary key references Vehicle(licensePlate),
  make varchar(10),
);
CREATE TABLE Truck (
  licensePlate varchar(20) Primary key references Vehicle(licensePlate),
  capacity varchar(10),
  ssn INT FOREIGN KEY REFERENCES ProfessionalDriver(ssn),
);
CREATE TABLE Drives (
  licensePlate varchar(20) references Car(licensePlate),
  ssn INT references NonProfessionalDriver(ssn),
  PRIMARY KEY(ssn, licensePlate)
);
```

2.2 Which relation in your relational schema represents the relationship "insures" in the E/R diagram and why is that your representation?

Answer: my table vehicle represents the relationship "insures", because insures is a N to 1 relationship, so I can use a foreign key in relation vehicle to save one table.

2.3 Compare the representation of the relationships "drives" and "operates" in your schema, and explain why they are different.

Answer: relationship "drives" is a N to N relationship, which means a non professional driver can drive as many cars as he wants, and a car can be driven by one or many non professional drivers. However, Operates is a N to 1 relationship, which means a truck can only be operated by at most one professional driver, and one professional drive can drive one or many trucks.

- 3. Consider the following two relational schemas and sets of functional dependencies
- 3.1 R(A,B,C,D,E) with functional dependencies D \rightarrow B, CE \rightarrow A

Answer:

Using BCNF Decomposition Algorithm on R(A,B,C,D,E)

X = D, $X^+ = BD$, $X \neq X^+$ and $X^+ \neq [all attributes]$

Y = B, Z = ABCDE - BD = ACE

R1 = BD, R2 = ACDE

decompose R2

Using BCNF Decomposition Algorithm on R(A,C,D,E)

X = CE, $X^+ = ACE$, $X \neq X^+$ and $X^+ \neq [all attributes]$

Y = A. Z = ACDE - ACE = D

R1 = ACE, R2 = DCE

So we have R1(B,D), R2(A,C,E), R3(D,C,E)

3.2 S(A,B,C,D,E) with functional dependencies A \rightarrow E, BC \rightarrow A, DE \rightarrow B.

Answer:

Using BCNF Decomposition Algorithm on R(A,B,C,D,E)

X = A. $X^+ = AE$. $X \neq X^+$ and $X^+ \neq A$ [all attributes]

Y = E. Z = ABCDE - AE = BCD

R1 = AE, R2 = ABCD

decompose R2

Using BCNF Decomposition Algorithm on R(A,B,C,D)

X = BC, $X^+ = ABC$, $X \neq X^+$ and $X^+ \neq [all attributes]$

Y = A, Z = ABCD - ABC = D

R1 = ABC, R2 = BCD

So we have R1(AE), R2(ABC), R3(BCD)

- 4. A set of attributes X is called closed (with respect to a given set of functional dependencies) if X+=X. Consider a relation with schema R(A,B,C,D) and an unknown set of functional dependencies. For each closed attribute set below, give a set of functional dependencies that is consistent with it.
- 4.1 All sets of attributes are closed

```
A \rightarrow A

B \rightarrow B

C \rightarrow C

D \rightarrow D

All are trivial FDs

4.2 The only closed sets are {} and {A,B,C,D}.

A \rightarrow B

B \rightarrow C

C \rightarrow D

D \rightarrow A

4.3 The only closed sets are {}, {A,B}, and {A,B,C,D}.

A \rightarrow B

B \rightarrow A
```

C→ABC D→BCD