

LECTURE 5: COMMUNITY STRUCTURE IN NETWORKS

Prof. Pan Hui

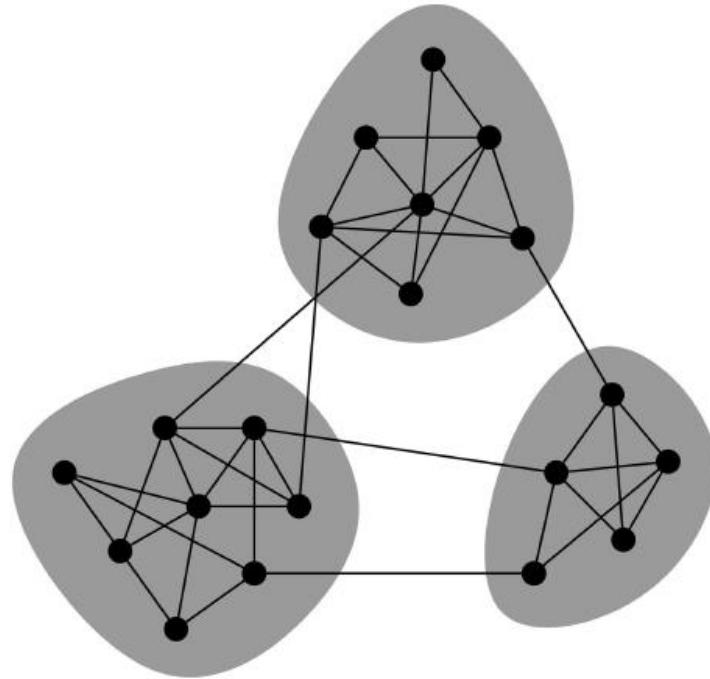
CSIT 6000K: Social Networks and Social Computing: A Data Science Perspective

Thursdays 07:30 PM - 10:20 PM

Networks & Communities

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- We often think of networks “looking” like this:



- What lead to such conceptual picture?

Networks: Flow of Information

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- **How information flows through the network?**
 - ▣ What structurally distinct roles do nodes play?
 - ▣ What roles do different **links** (**short** vs. **long**) play?
- **How people find out about new jobs?**
 - ▣ Mark Granovetter, part of his PhD in 1960s
 - ▣ People find the information through personal contacts
- **But:** Contacts were often **acquaintances** rather than close friends
 - ▣ **This is surprising:** One would expect your friends to help you out more than casual acquaintances
- **Why is it that acquaintances are most helpful?**

Granovetter's Answer

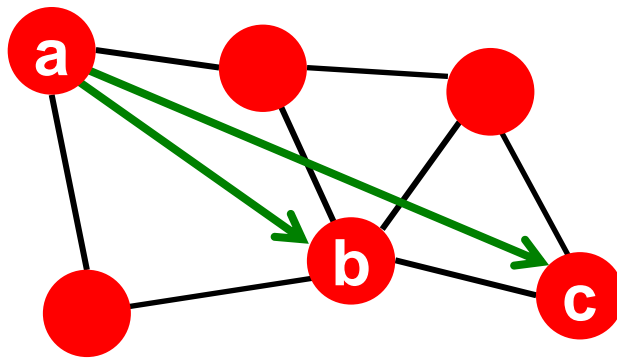
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□ Two perspectives on **friendships**:

- **Structural**: Friendships span different parts of the network

- **Interpersonal**: Friendship between two people is either **strong** or **weak**

□ **Structural role: Triadic Closure**



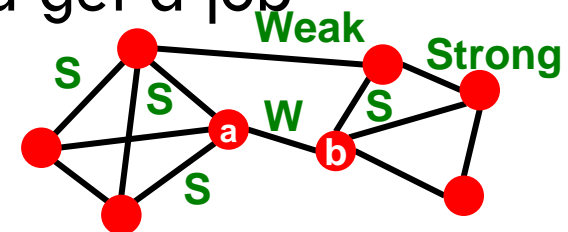
Which edge is more likely a-b or a-c?

If two people in a network have a friend in common there is an increased likelihood they will become friends themselves

Granovetter's Explanation

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- **Granovetter makes a connection between social and structural role of an edge**
- **First point:**
 - ▣ Structurally embedded edges are also socially strong
 - ▣ Edges spanning different parts of the network are socially weak
- **Second point:**
 - ▣ The long range edges allow you to gather information from different parts of the network and get a job
 - ▣ Structurally embedded edges are heavily redundant in terms of information access



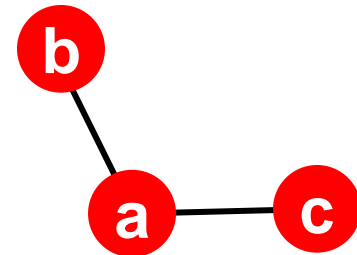
Triadic Closure

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- **Triadic closure == High clustering coefficient**

Reasons for triadic closure:

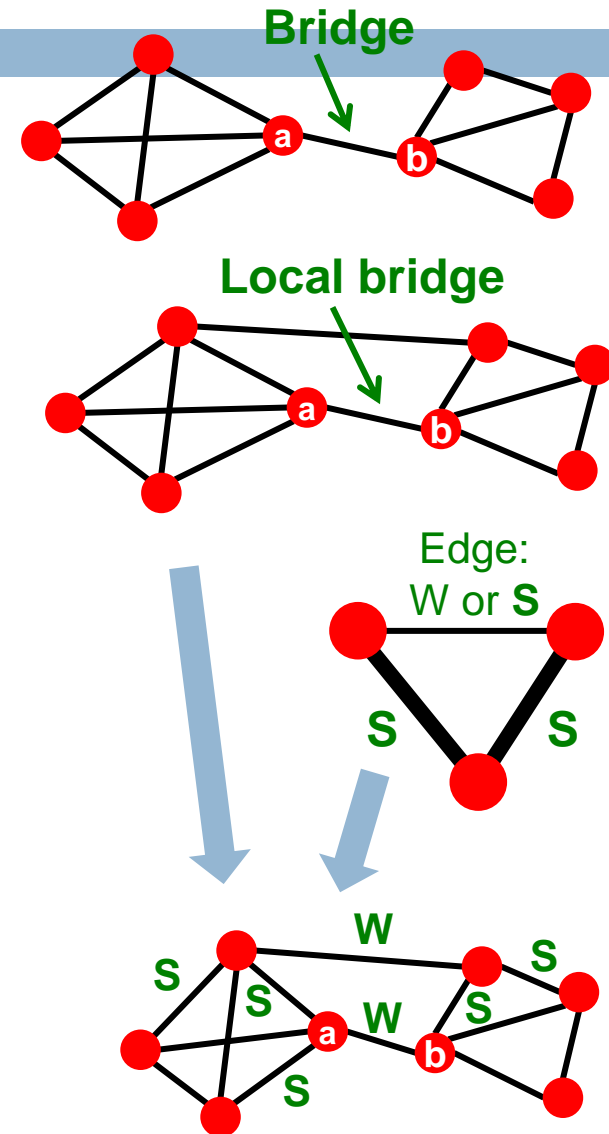
- If B and C have a friend A in common, then:
 - B is **more likely to meet** C
 - (since they both spend time with A)
 - B and C **trust** each other
 - (since they have a friend in common)
 - A has **incentive** to bring B and C together
 - (as it is hard for A to maintain two disjoint relationships)
- **Empirical study by Bearman and Moody:**
 - Teenage girls with low clustering coefficient are more likely to contemplate suicide



Granovetter's Explanation

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- Define: **Bridge edge**
 - If removed, it disconnects the graph
- Define: **Local bridge**
 - Edge of $\text{Span} > 2$
(**Span** of an edge is the distance of the edge endpoints if the edge is deleted. **Local bridges with long span are like real bridges**)
- Define: Two types of edges:
 - **Strong** (friend), **Weak** (acquaintance)
- Define: **Strong triadic closure:**
 - Two strong ties imply a third edge
- **Fact:** If strong triadic closure is satisfied then **local bridges are weak ties!**



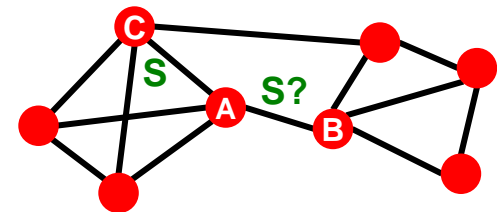
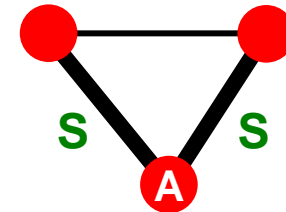
Local Bridges and Weak ties

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- **Claim:** If node satisfies **Strong Triadic Closure** and is involved in at least **two strong ties**, then any **local bridge** adjacent to must be a **weak tie**.

- **Proof by contradiction:**

- satisfies **Strong Triadic Closure**
- Let be local bridge and a **strong tie**
- Then must exist because of **Strong Triadic Closure**
- But then is **not a bridge!**



Tie strength in real data

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- **For many years the Granovetter's theory was not tested**
- But, today we have large who-talks-to-whom graphs:
 - ▣ Email, Messenger, Cell phones, Facebook
- **Onnela et al. 2007:**
 - ▣ Cell-phone network of 20% of country's population
 - ▣ **Edge strength:** # phone calls

Neighborhood Overlap

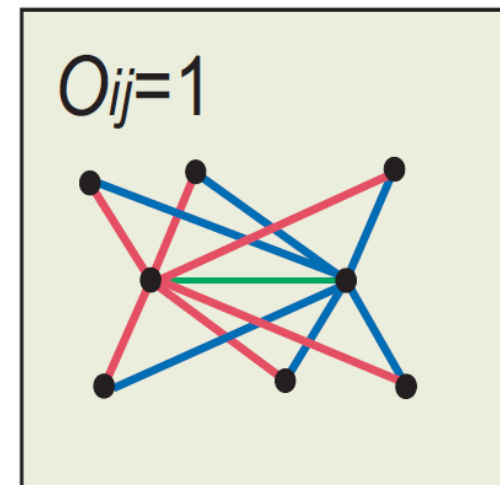
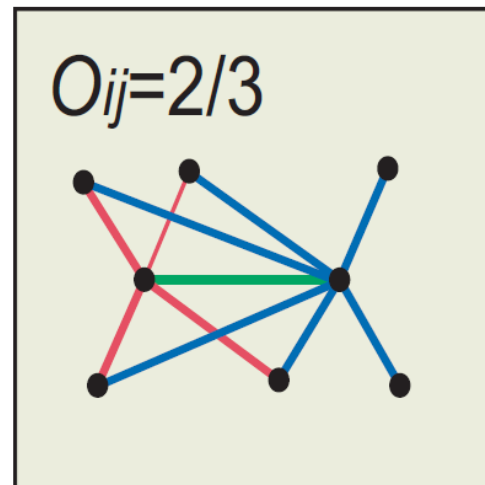
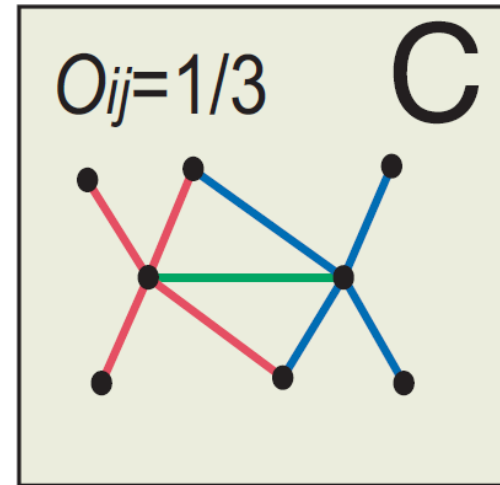
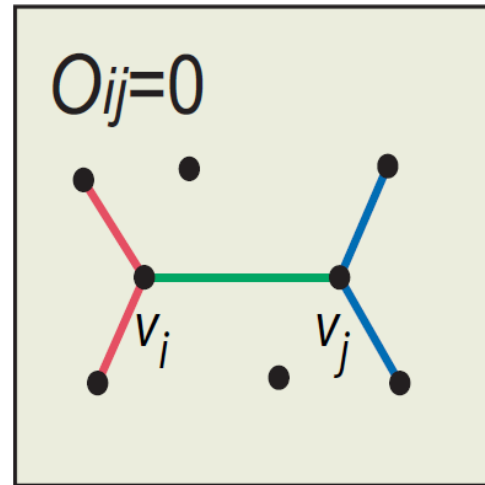
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□ Edge overlap:

$$O_{ij} = \frac{|\hat{N}(i) \cap \hat{N}(j)|}{|\hat{N}(i) \cup \hat{N}(j)|}$$

□ $N(i)$... a set of neighbors of node i

□ Overlap = 0 when an edge is a **local bridge**



Phones: Edge Overlap vs. Strength

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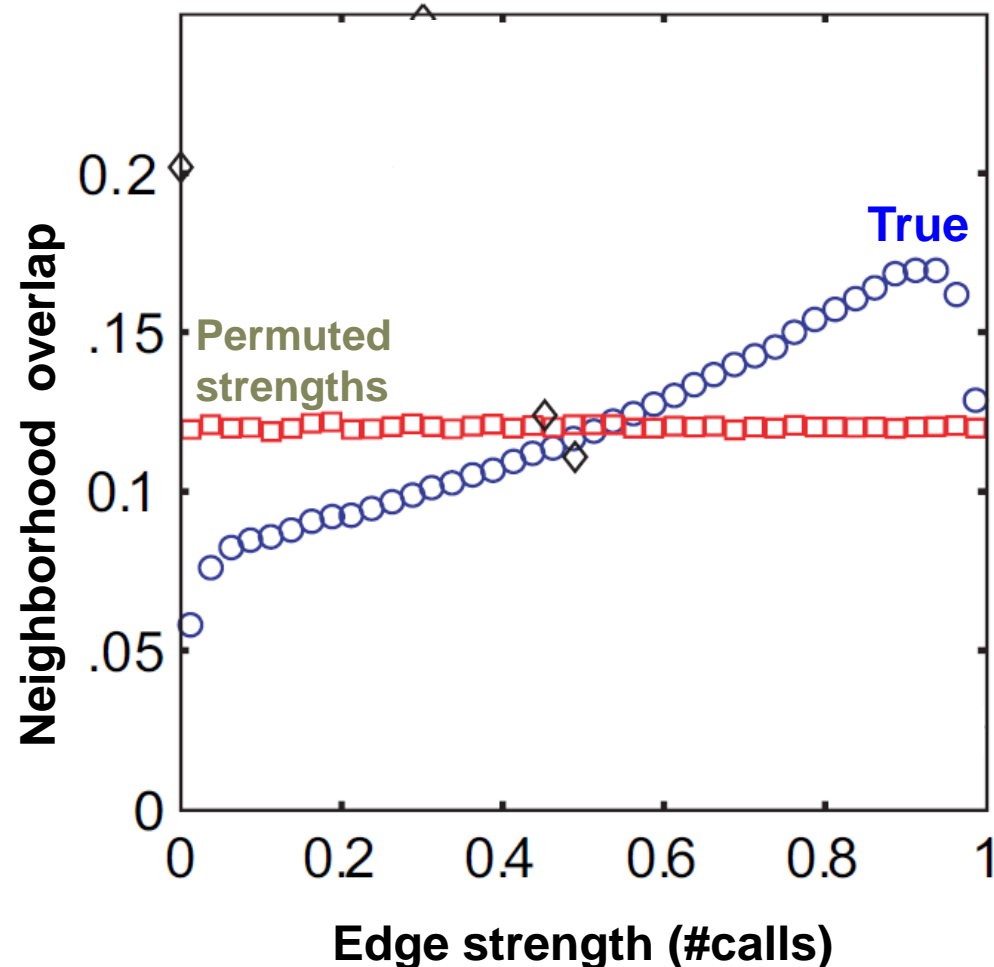
Cell-phone network

Observation:

- Highly used links have high overlap!

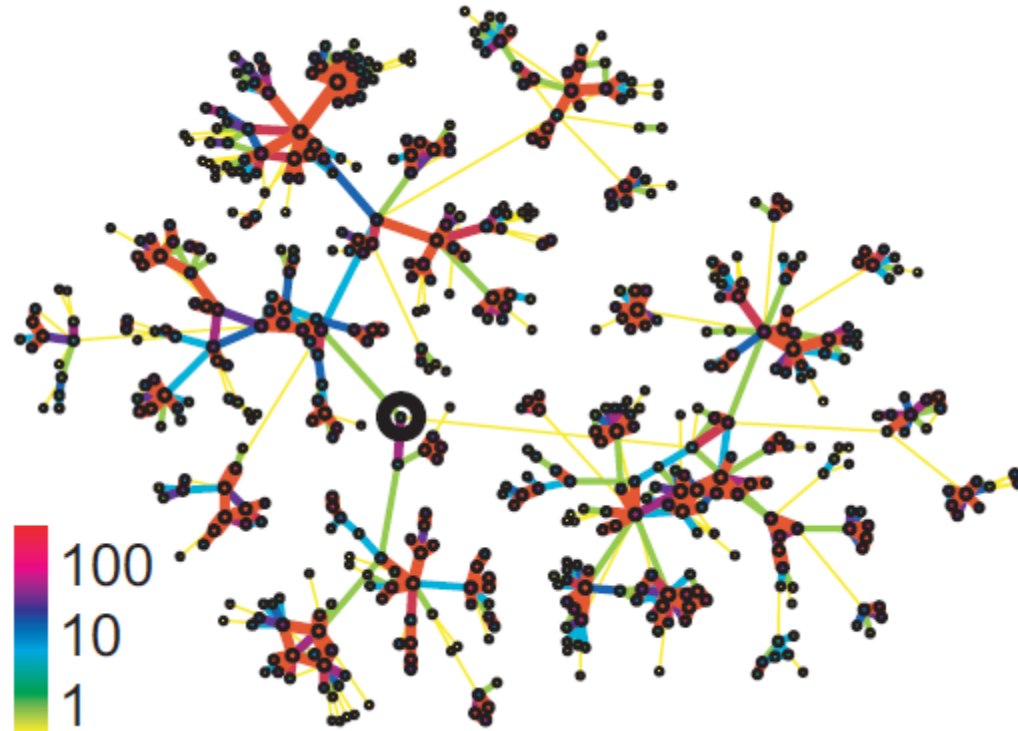
Legend:

- True:** The data
- Permuted strengths:** Keep the network structure but randomly reassign edge strengths



Real Network, Real Tie Strengths

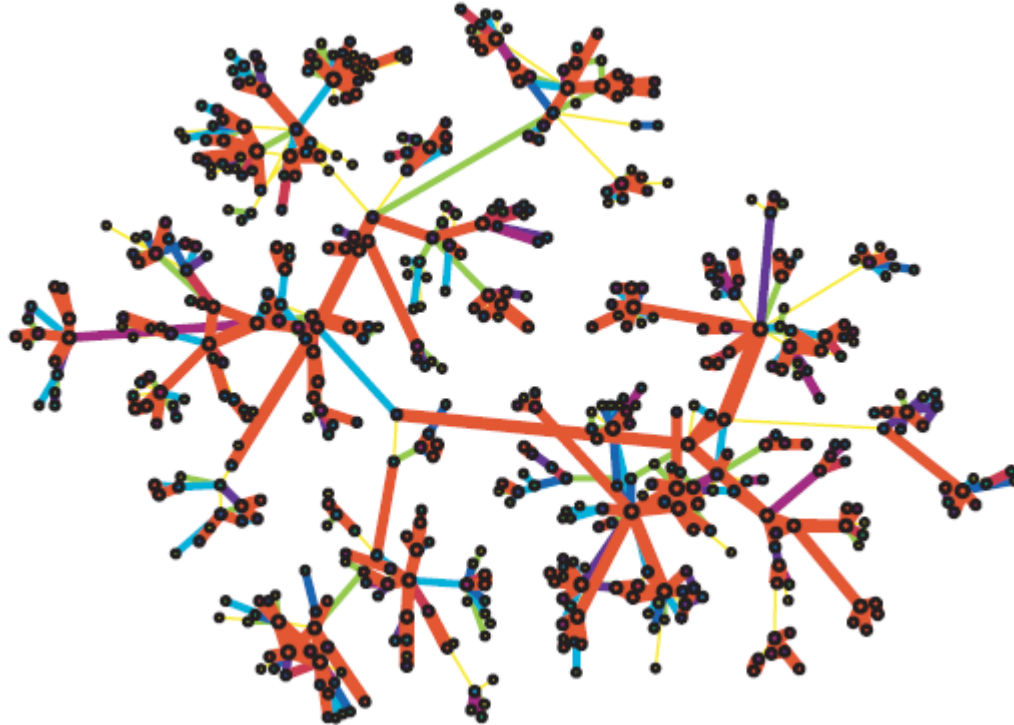
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- **Real edge strengths in mobile call graph**
 - Strong ties are more embedded (have higher overlap)

Real Net, Permuted Tie Strengths

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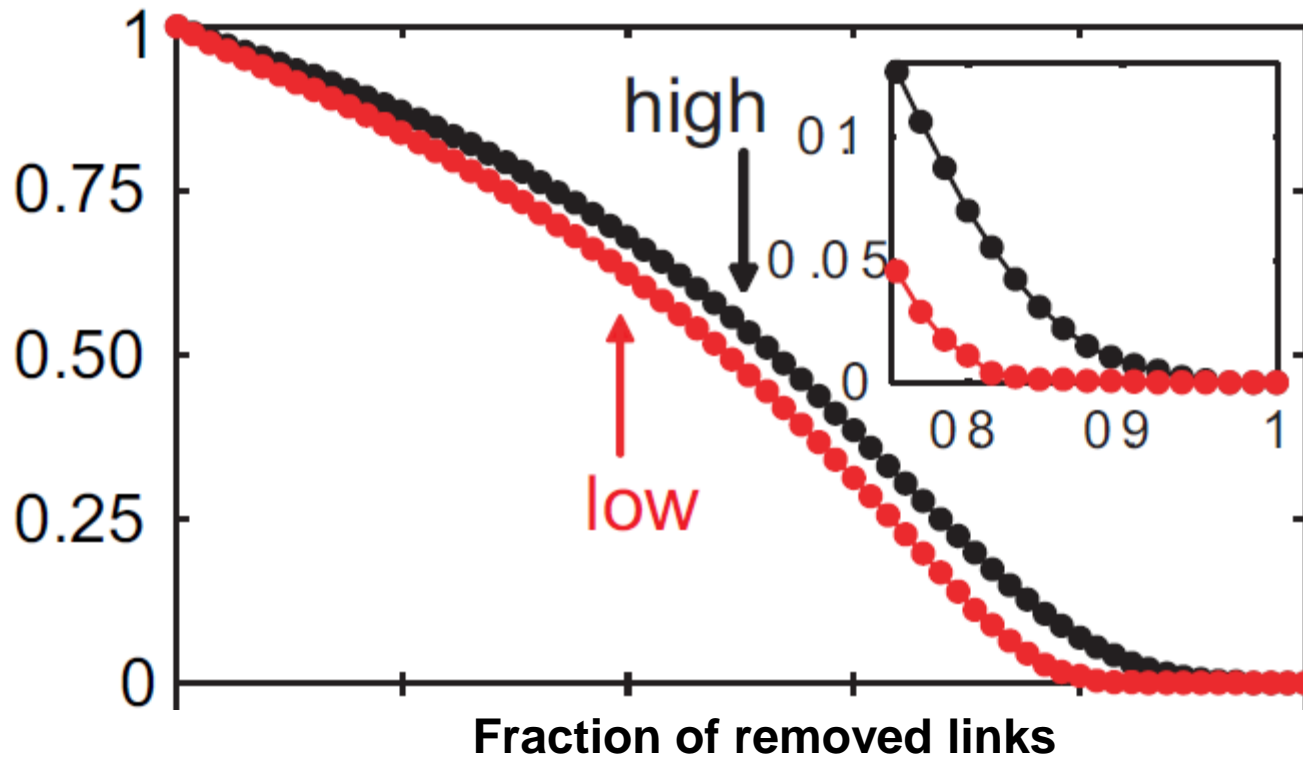


- Same network, same set of edge strengths but now **strengths are randomly shuffled**

Link Removal by Strength

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Size of largest component

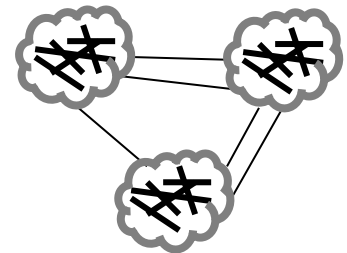


Low
disconnects
the network
sooner

□ Removing links by **strength (#calls)**

□ Low to high

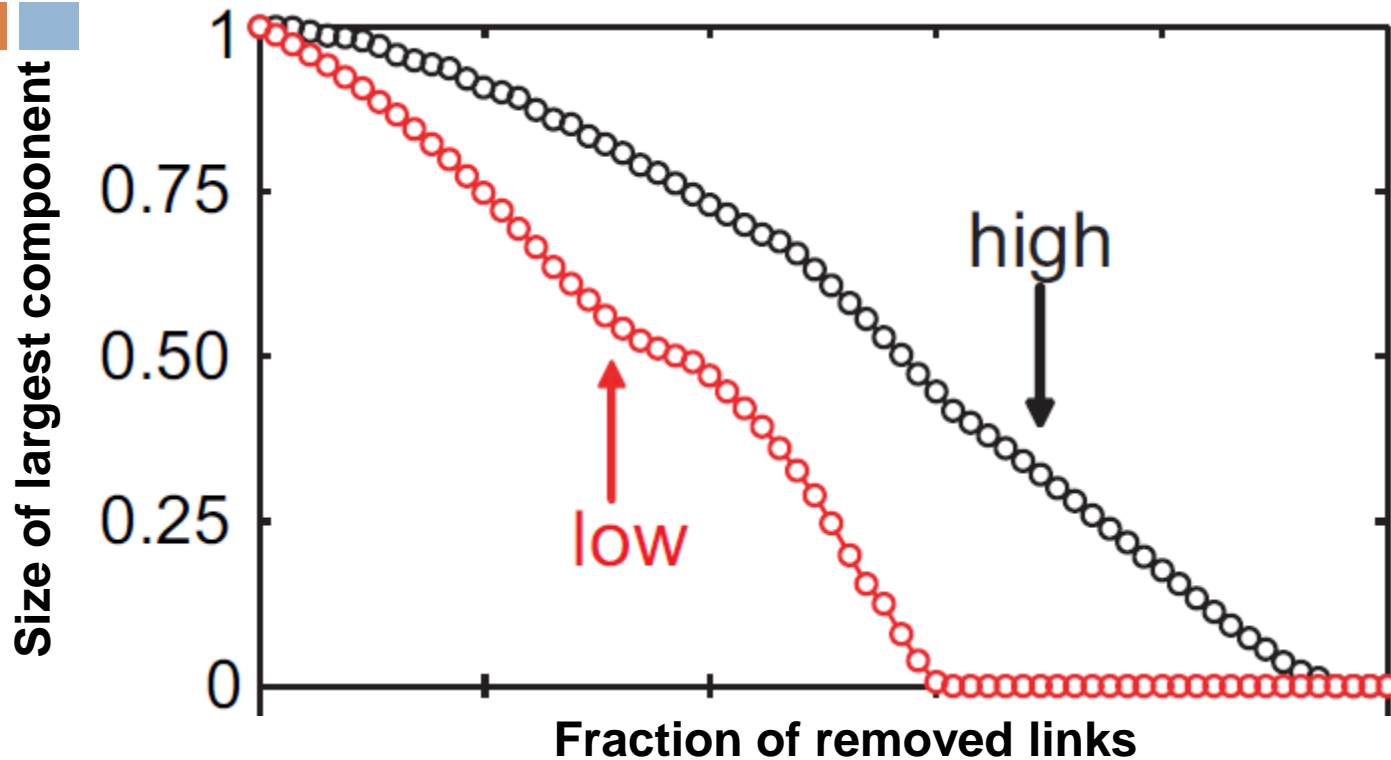
□ High to low



Conceptual picture
of network structure

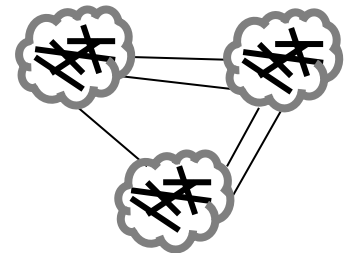
Link Removal by Overlap

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Low
disconnects
the network
sooner

- Removing links based on **overlap**
 - Low to high
 - High to low

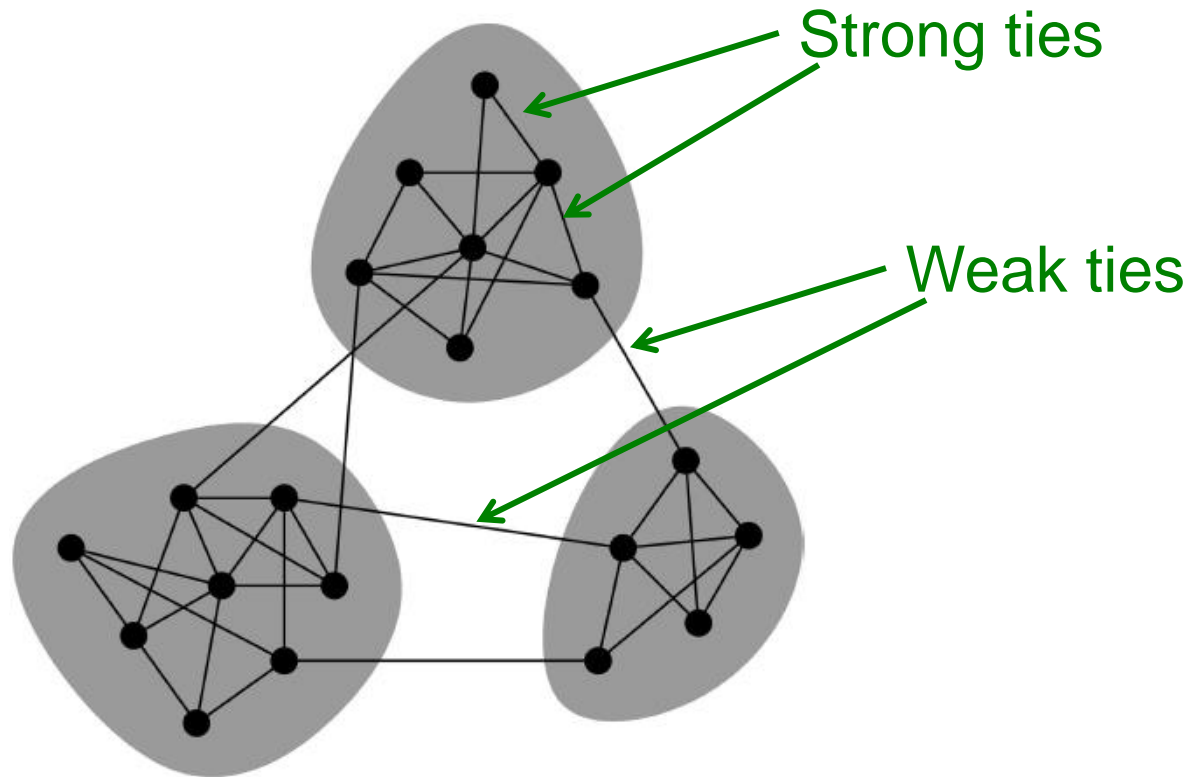


Conceptual picture
of network structure

Conceptual Picture of Networks

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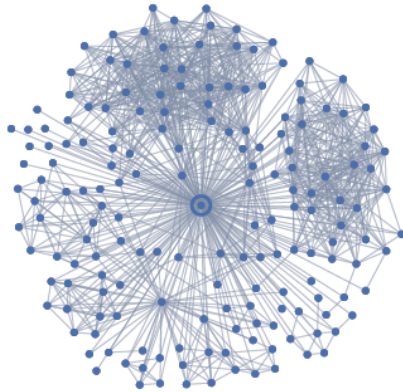
- Granovetter's theory leads to the following conceptual picture of networks



Facebook User's Tie Strength

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All Friends



Maintained Relationships



One-way Communication



Mutual Communication

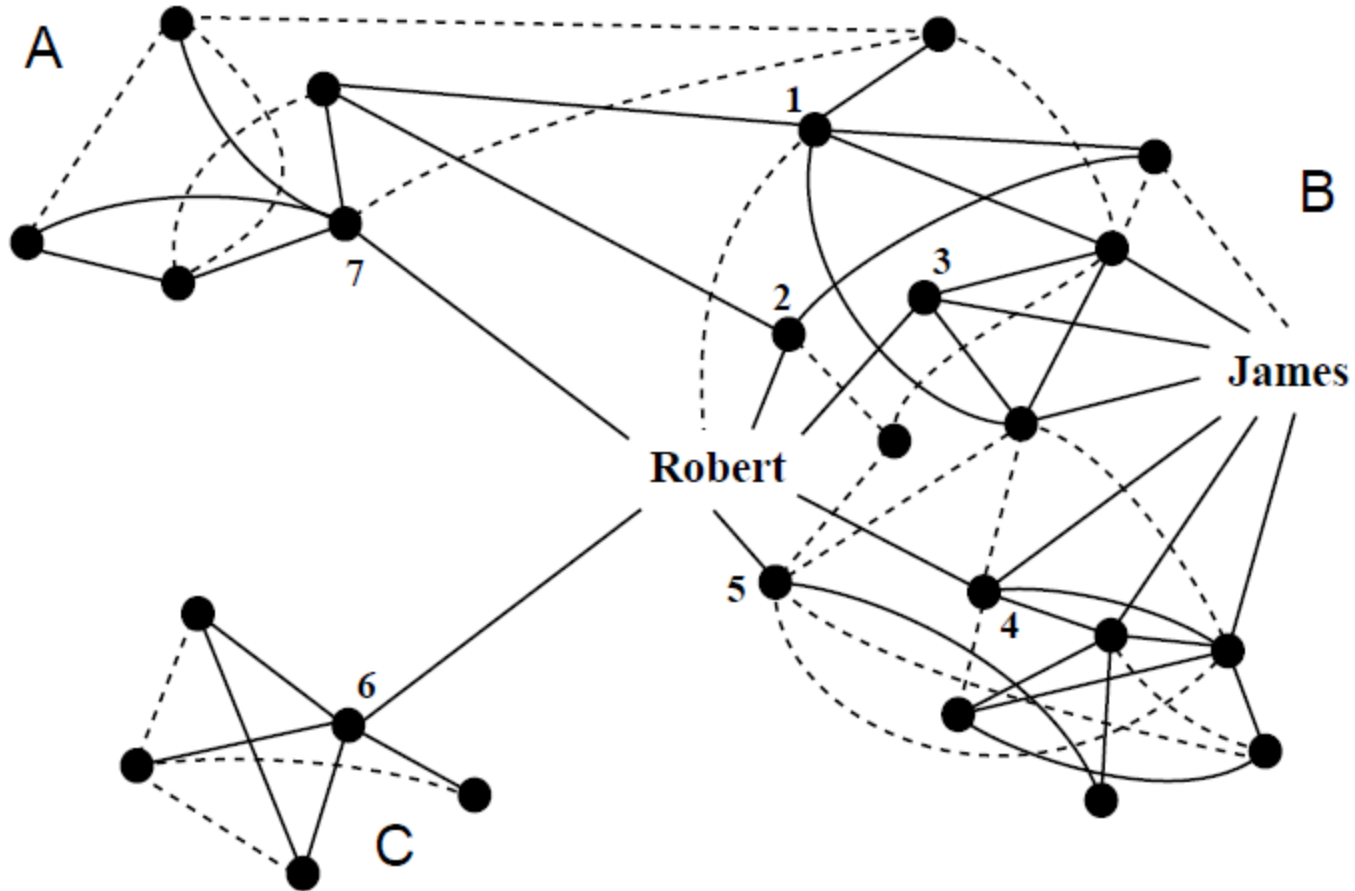


Four different views of a Facebook User's network neighborhood, show the structure of links corresponding respectively to all declared friendships, maintained relationships, one-way communication, and reciprocal communication

SMALL DETOUR: STRUCTURAL HOLES

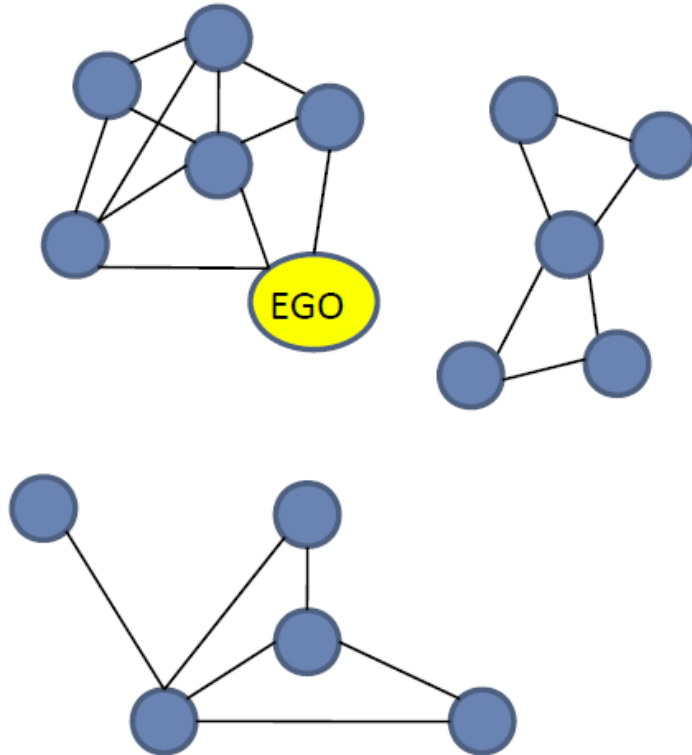
Small Detour: Structural Holes

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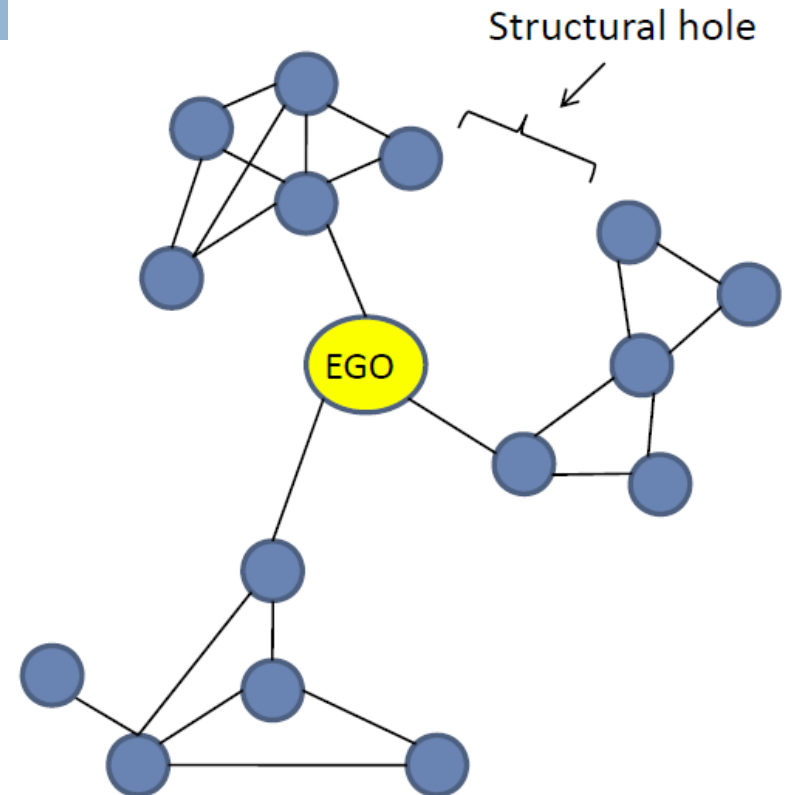


Structural Holes

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Few structural holes



Many structural holes

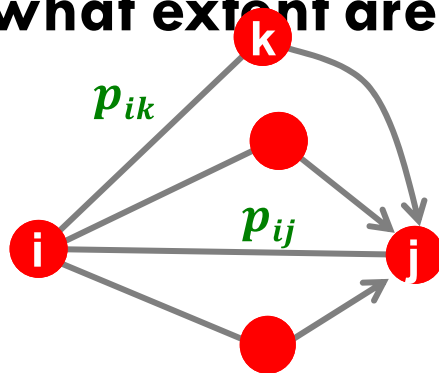
**Structural Holes provide ego with access
to novel information, power, freedom**

Structural Holes: Network Constraint

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□ The “network constraint” measure [Burt]:

▣ To what extent are person’s contacts redundant



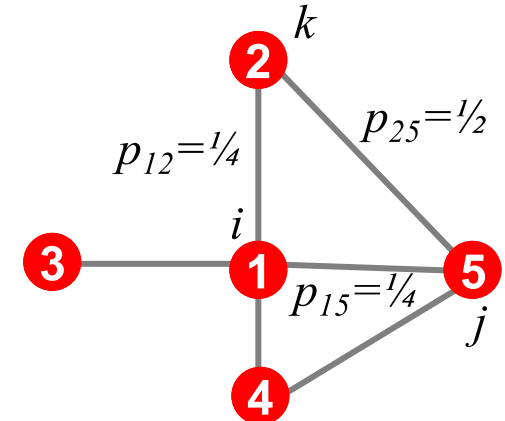
$$p_{uv} = 1/du$$

■ **Low:** disconnected contacts

■ **High:** contacts that are close or strongly tied

$$c_i = \sum_j c_{ij} = \sum_j \left[p_{ij} + \sum_k (p_{ik} p_{kj}) \right]^2$$

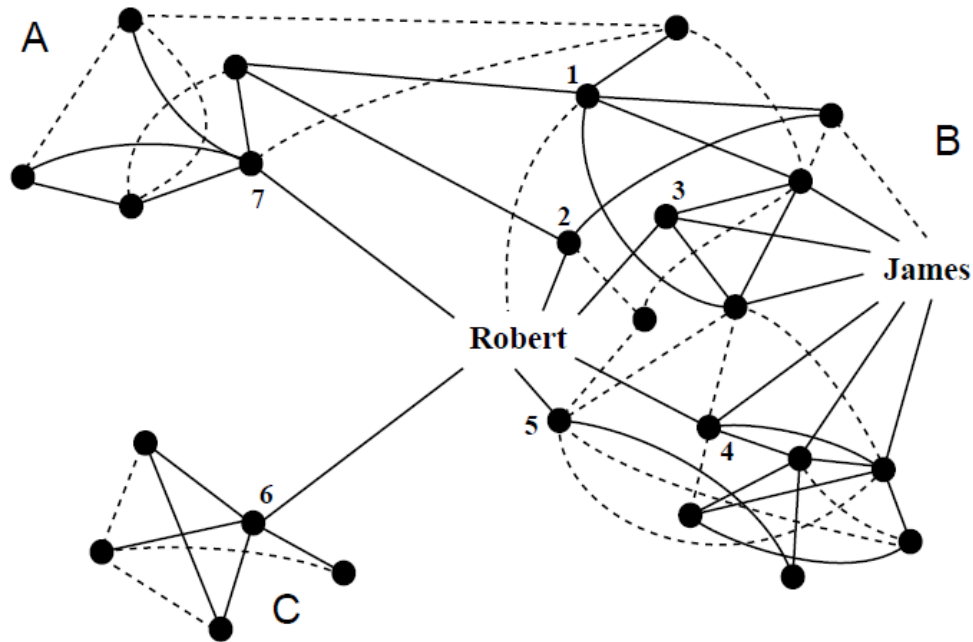
p_{uv} ... prop. of u 's “energy” invested in relationship with v



	p_{uv}				
	1	2	3	4	5
1	.00	.25	.25	.25	.25
2	.50	.00	.00	.00	.50
3	1.0	.00	.00	.00	.00
4	.50	.00	.00	.00	.50
5	.33	.33	.00	.33	.00

Example: Robert vs. James

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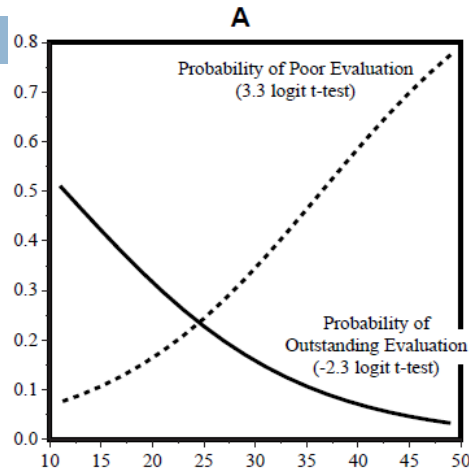
- **Constraint:** To what extent are person's contacts redundant
 - **Low:** disconnected contacts
 - **High:** contacts that are close or strongly tied

□ **Network constraint:**

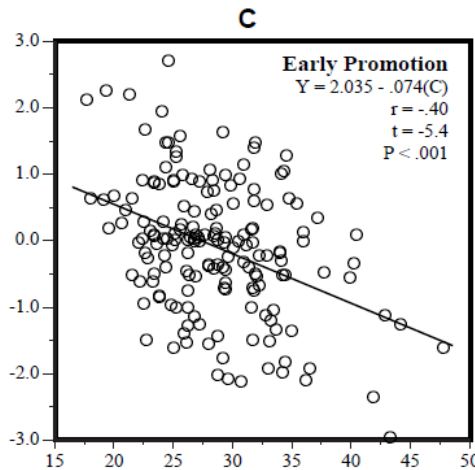
- James:
- Robert:

Spanning the Holes Matters

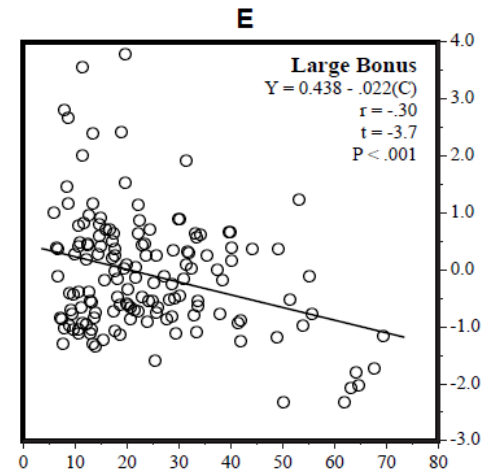
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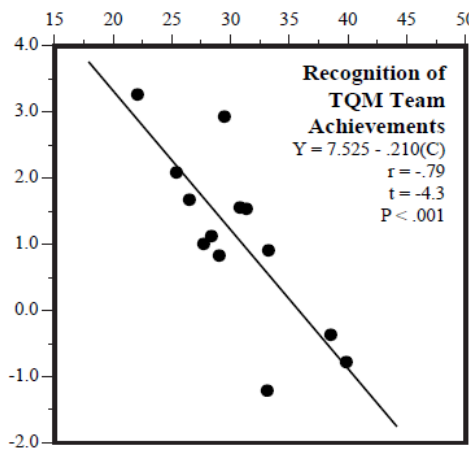
Network Constraint
many ——— Structural Holes ——— few
(manager C above, mean C in team below)



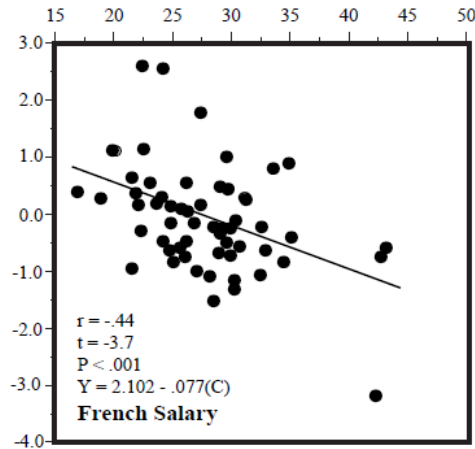
Network Constraint
many ——— Structural Holes ——— few
(C for manager's network)



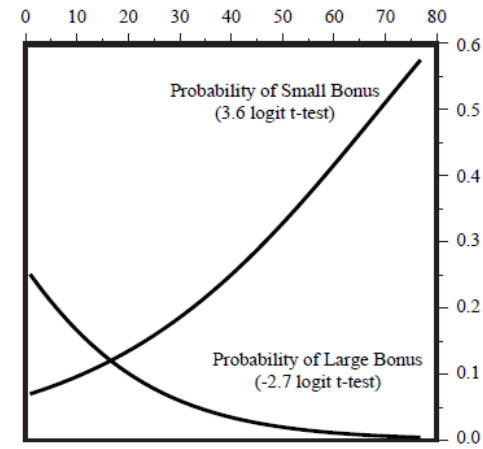
Network Constraint
many ——— Structural Holes ——— few
(C for officer's network)



B



D



F

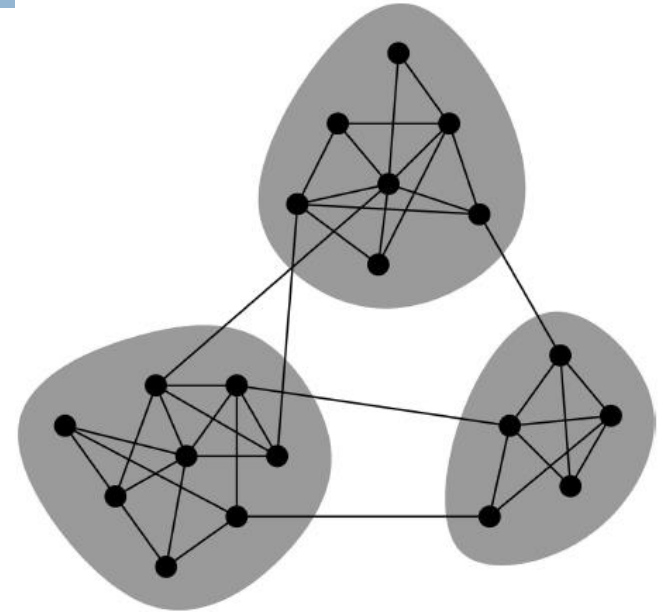
NETWORK COMMUNITIES

2/23/2022

Network Communities

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- Granovetter's theory (and common sense) suggest that networks are composed of **tightly connected sets of nodes**



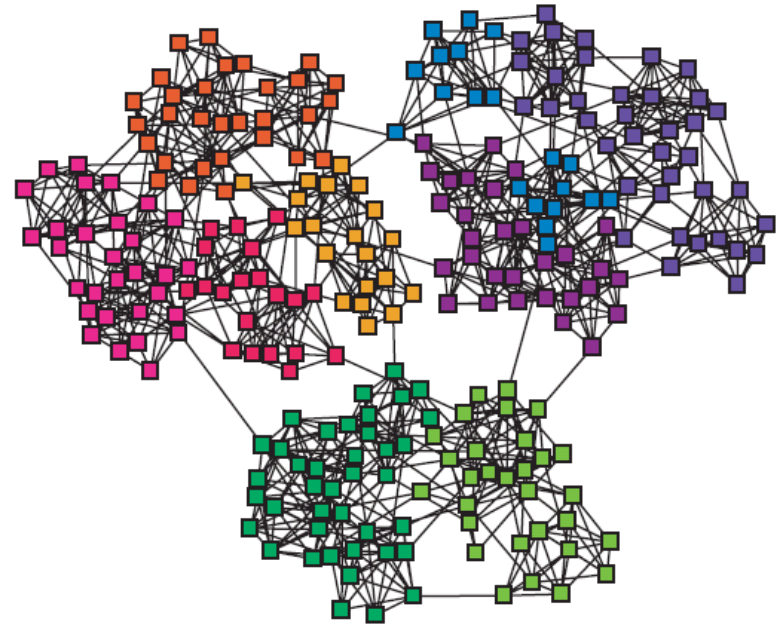
Communities, clusters, groups, modules

- **Network communities:**
 - Sets of nodes with **lots** of connections **inside** and **few** to **outside** (the rest of the network)

Finding Network Communities

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- **How to automatically find such densely connected groups of nodes?**
- Ideally such automatically detected clusters would then correspond to real groups
- **For example:**

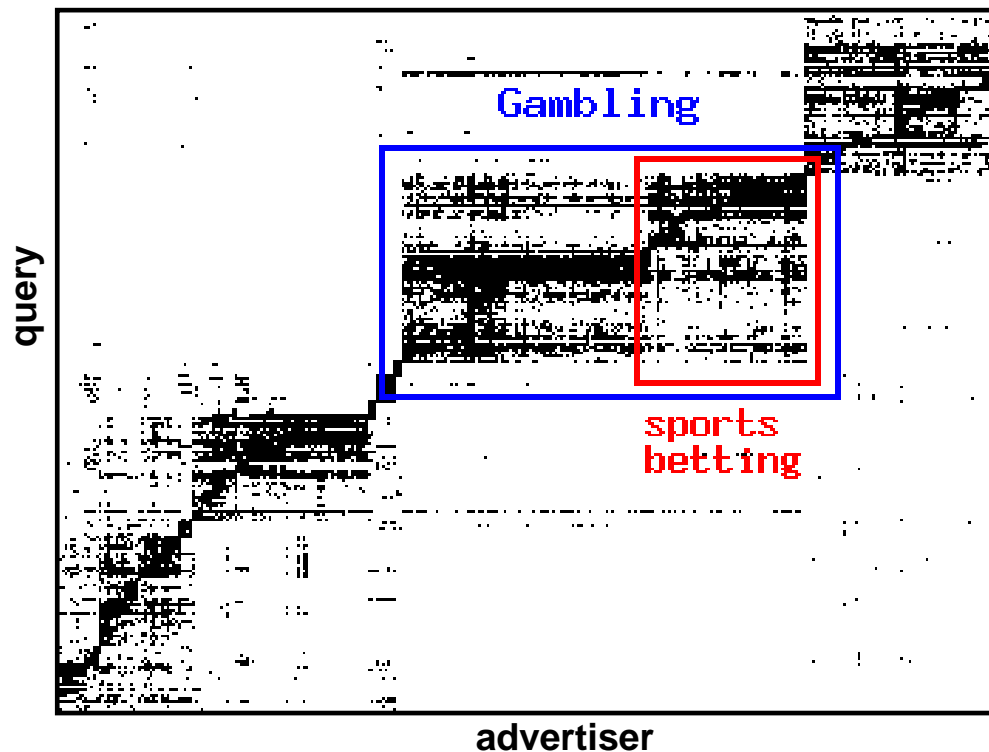


Communities, clusters,
groups, modules

Micro-Markets in Sponsored Search

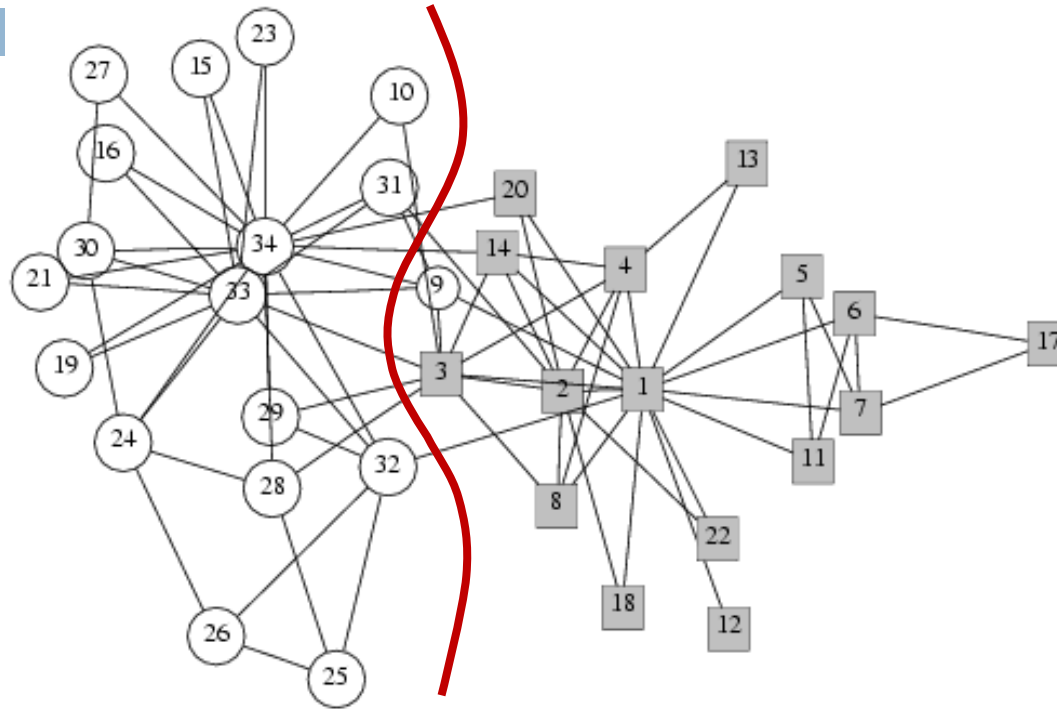
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Find micro-markets by partitioning the “query x advertiser” graph:



Social Network Data

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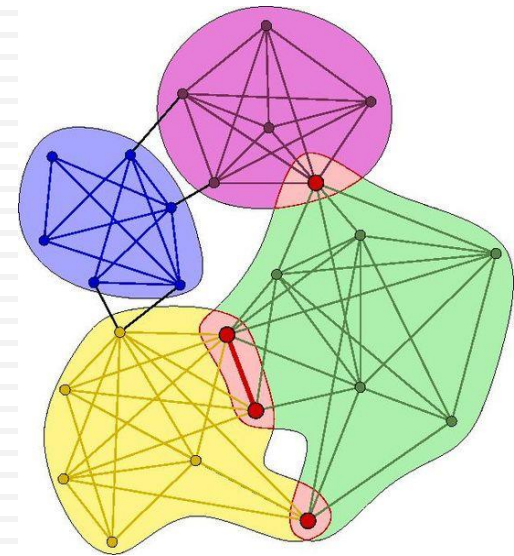
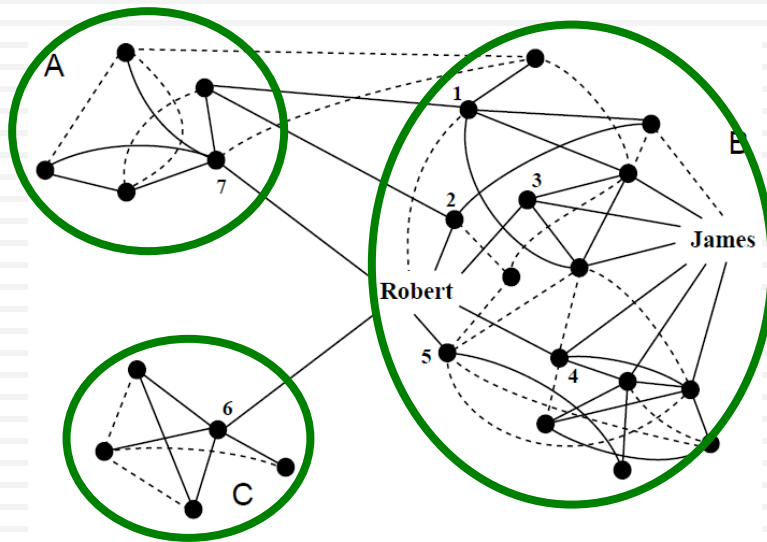


□ Zachary's Karate club network:

- Observe social ties and rivalries in a university karate club
- During his observation, conflicts led the group to split
- Split could be explained by a minimum cut in the network

Community Detection

How to find communities?



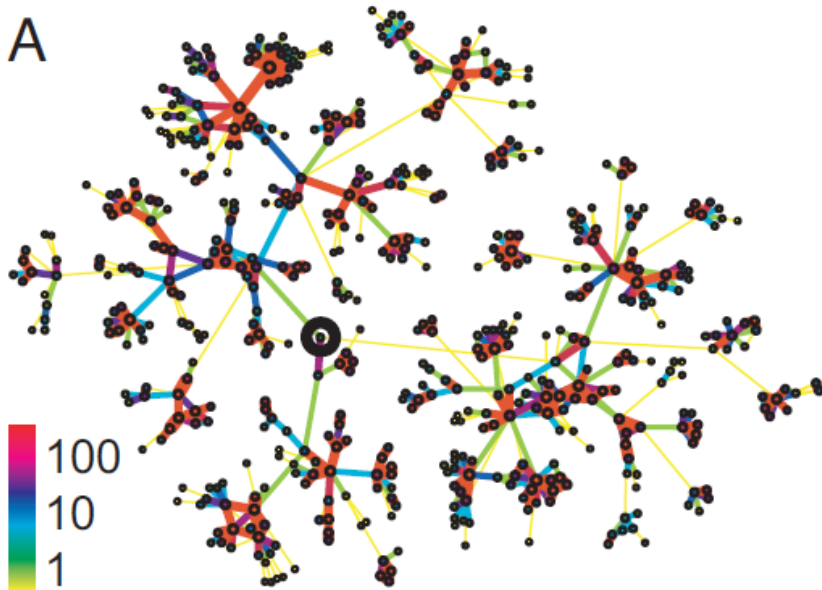
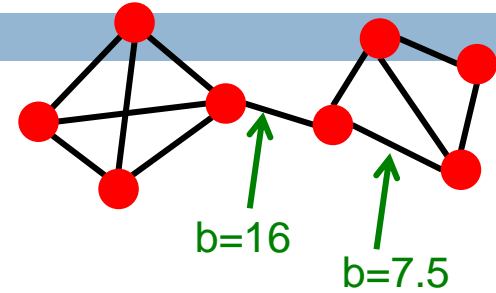
We will work with **undirected** (unweighted) networks

Method 1: Strength of Weak Ties

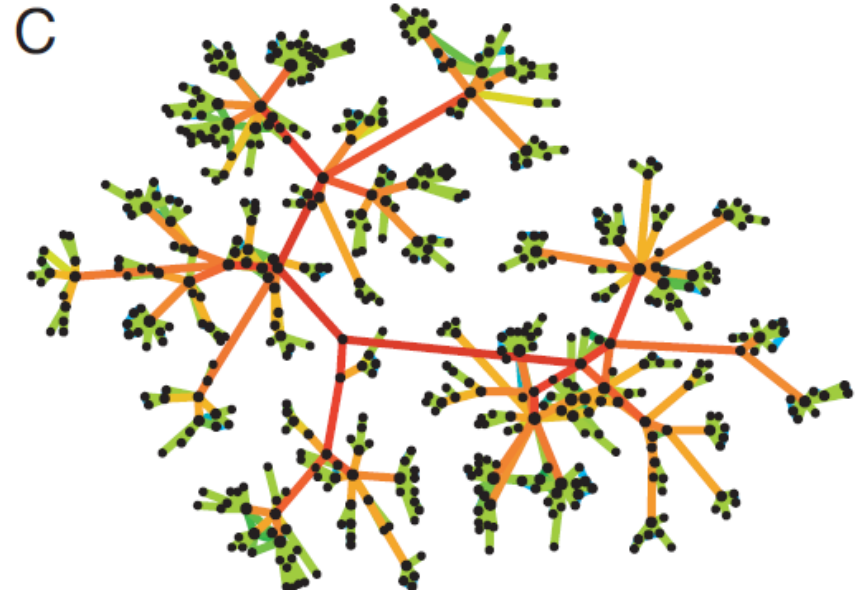
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- **Edge betweenness:** Number of shortest paths passing over the edge

- **Intuition:**



Edge strengths (call volume)
in real network



Edge betweenness
in real network

Method 1: Girvan-Newman

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- Divisive hierarchical clustering based on the notion of edge **betweenness**:

Number of shortest paths passing through the edge

- **Girvan-Newman Algorithm:**

- **Undirected unweighted networks**

- **Repeat until no edges are left:**

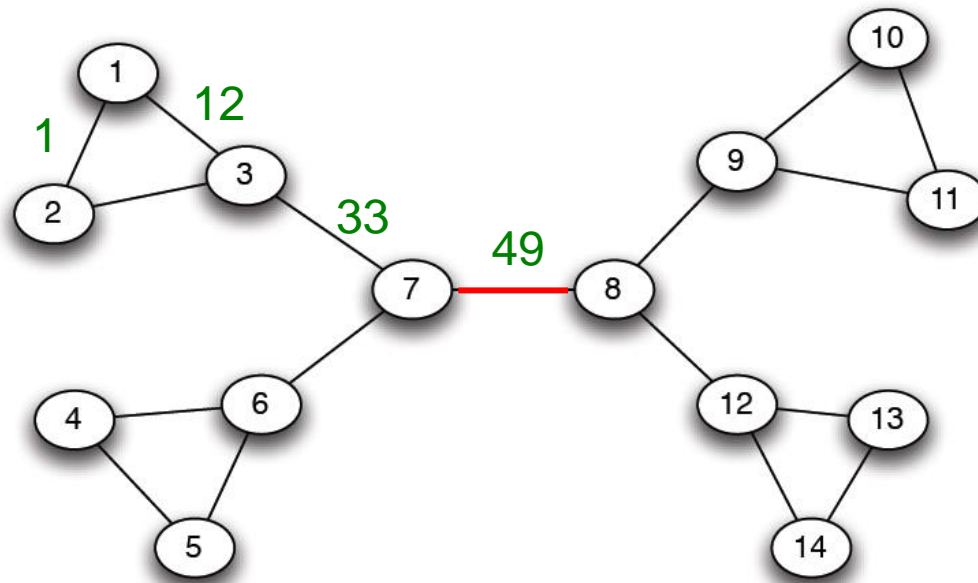
- Calculate betweenness of edges
 - Remove edges with highest betweenness

- Connected components are communities

- Gives a hierarchical decomposition of the network

Girvan-Newman: Example

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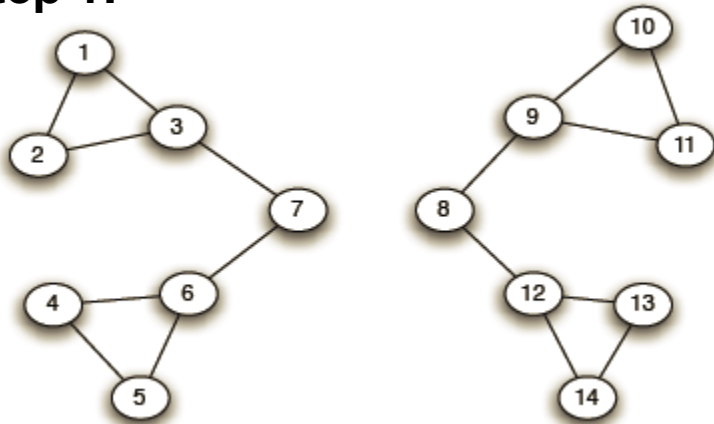


Need to re-compute
betweenness at
every step

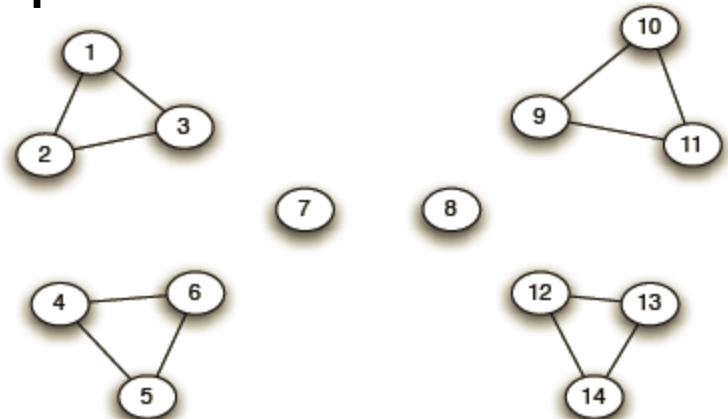
Girvan-Newman: Example

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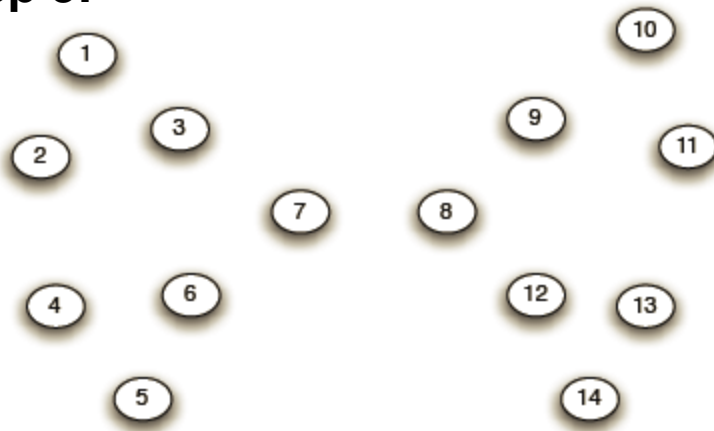
Step 1:



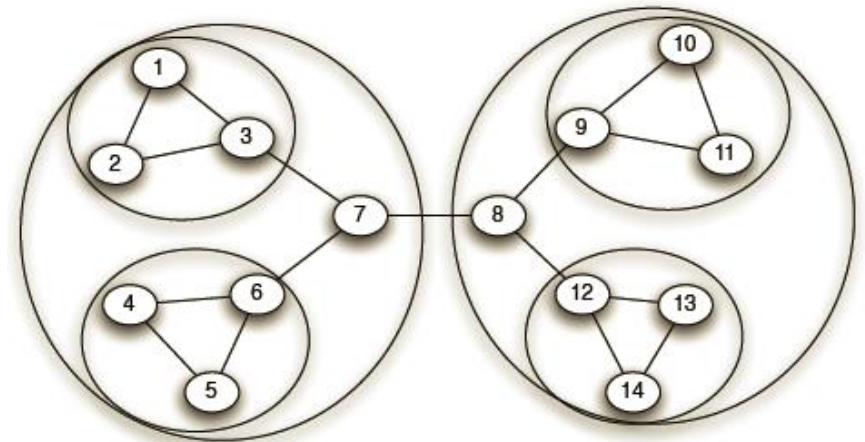
Step 2:



Step 3:

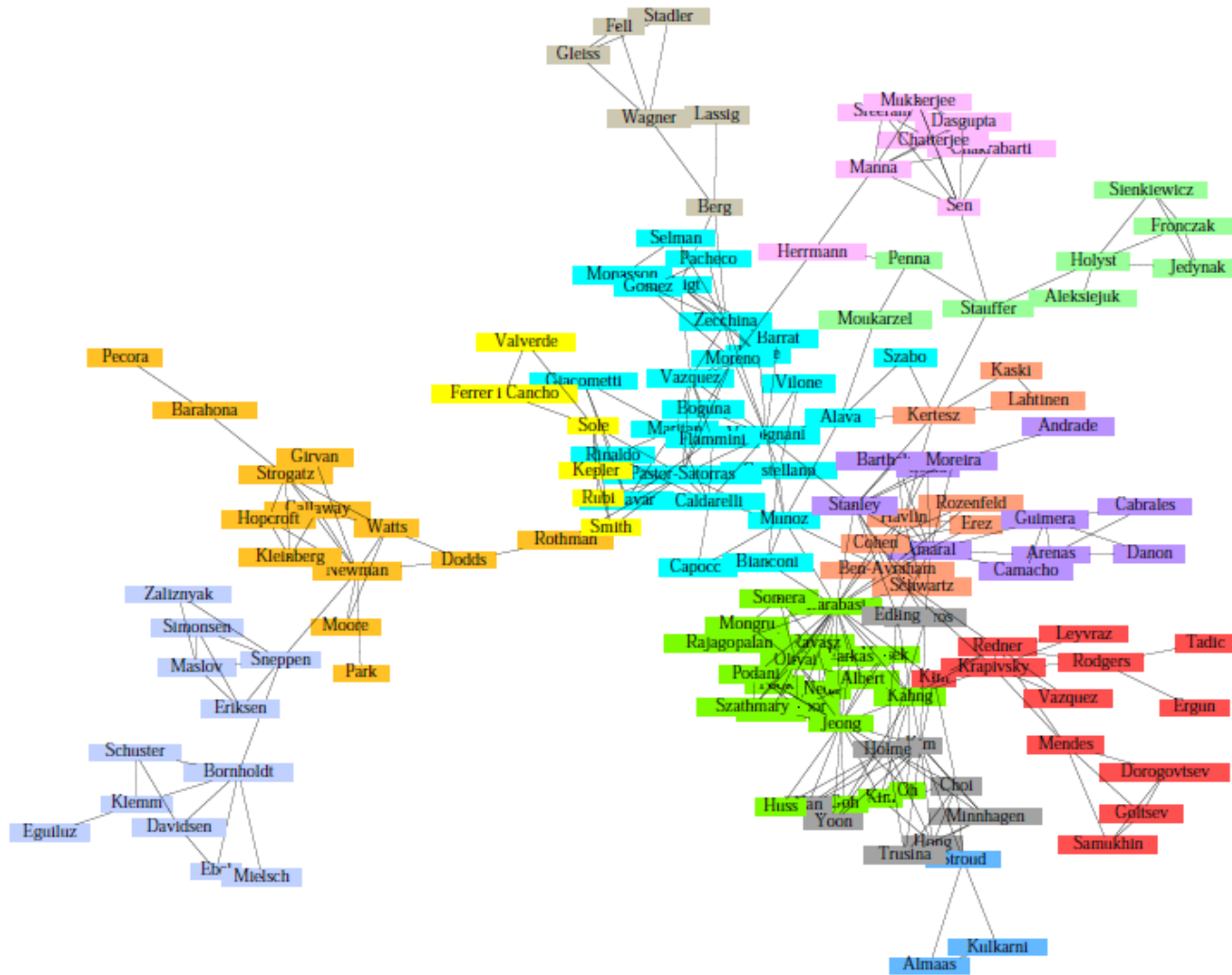


Hierarchical network decomposition:



Girvan-Newman: Results

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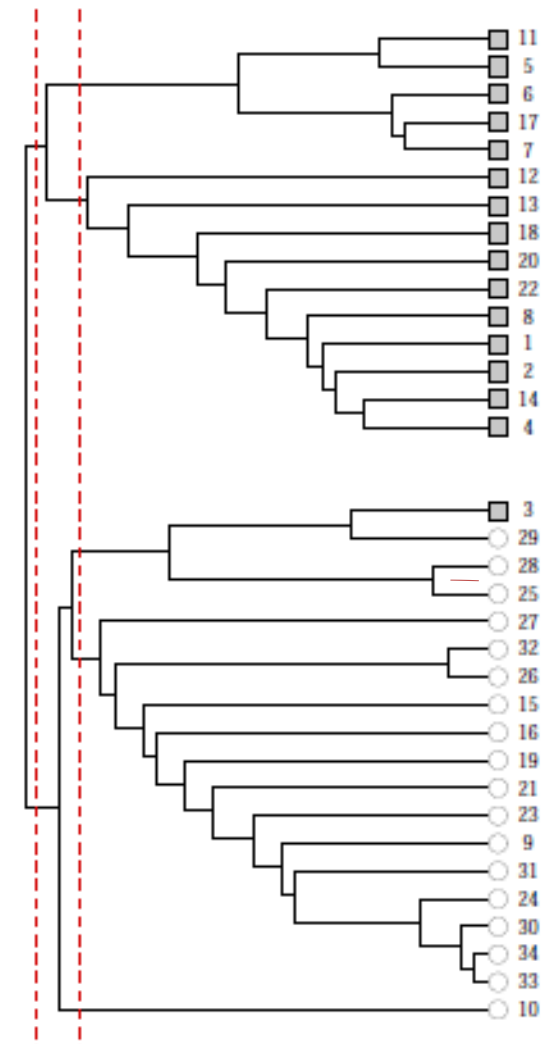
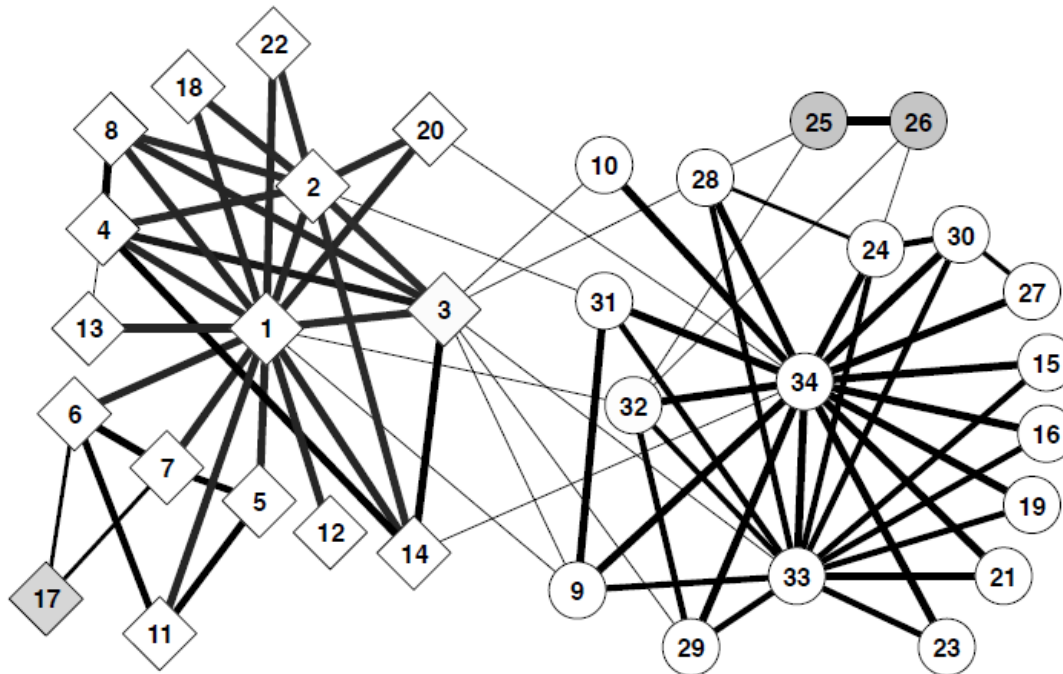


Communities in physics collaborations

Girvan-Newman: Results

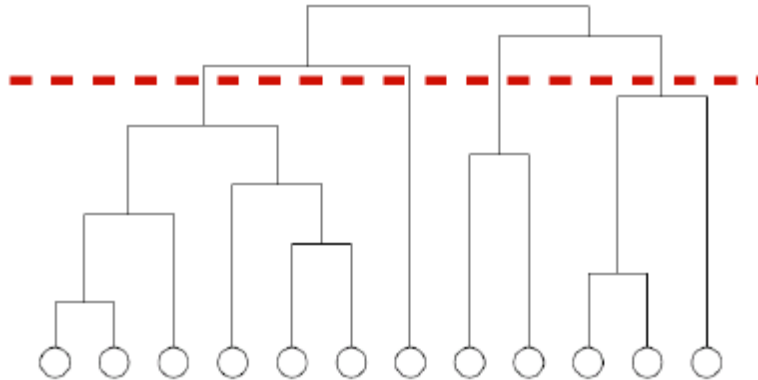
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□ Zachary's Karate club: Hierarchical decomposition



We need to resolve 2 questions

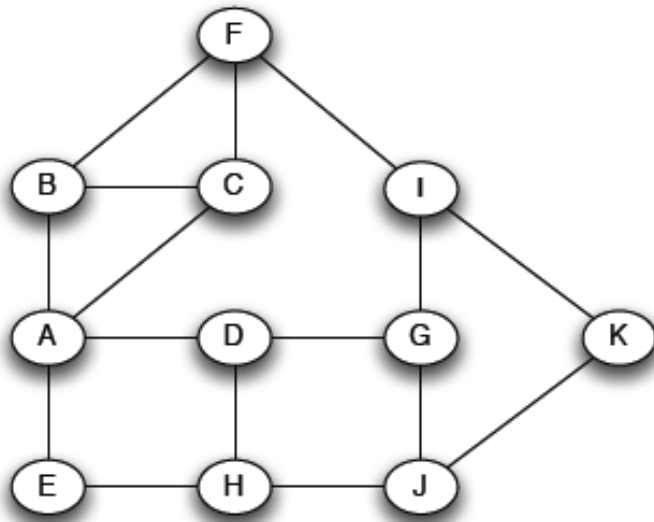
1. **How to compute betweenness?**
2. **How to select the number of clusters?**



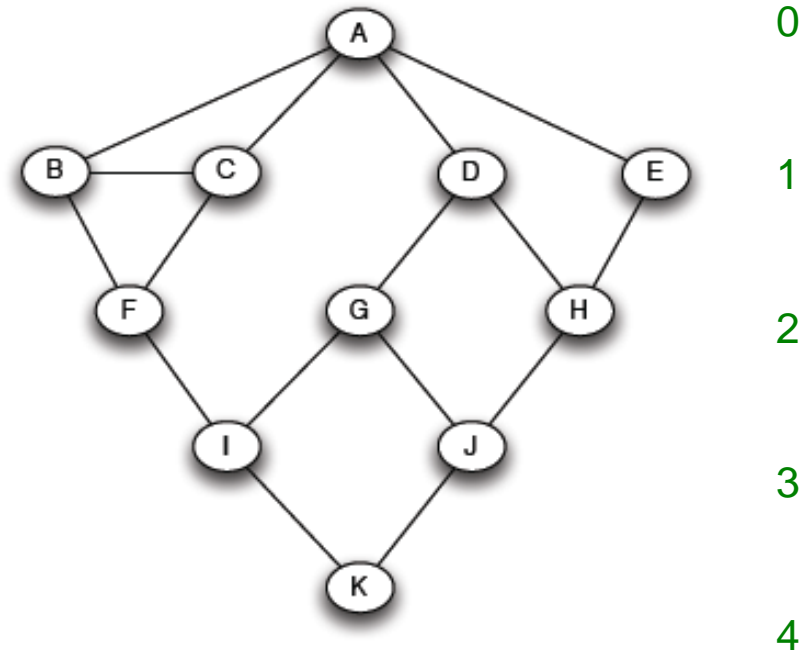
How to Compute Betweenness?

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- Want to compute betweenness of paths starting at node *A*



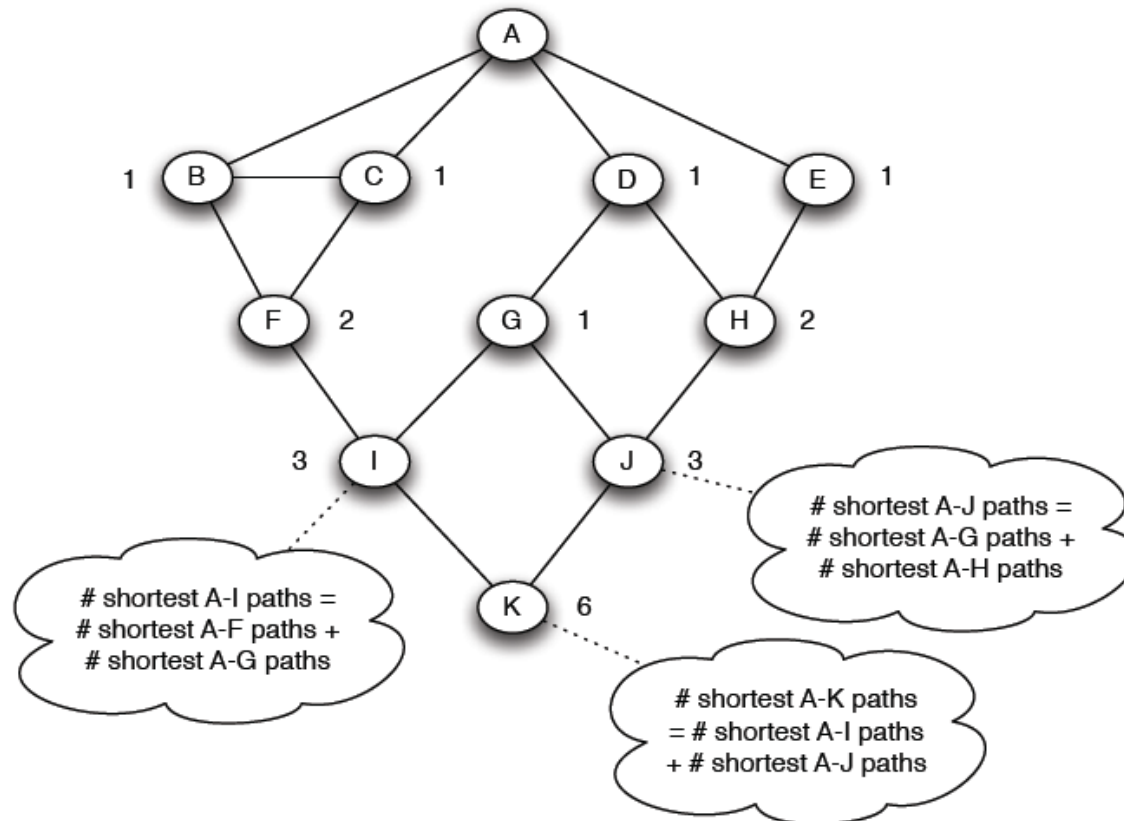
- Breath first search starting from *A*:



How to Compute Betweenness?

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- Count the number of shortest paths from *A* to all other nodes of the network:



How to Compute Betweenness?

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- **Compute betweenness by working up the tree:** If there are multiple paths count them fractionally

The algorithm:

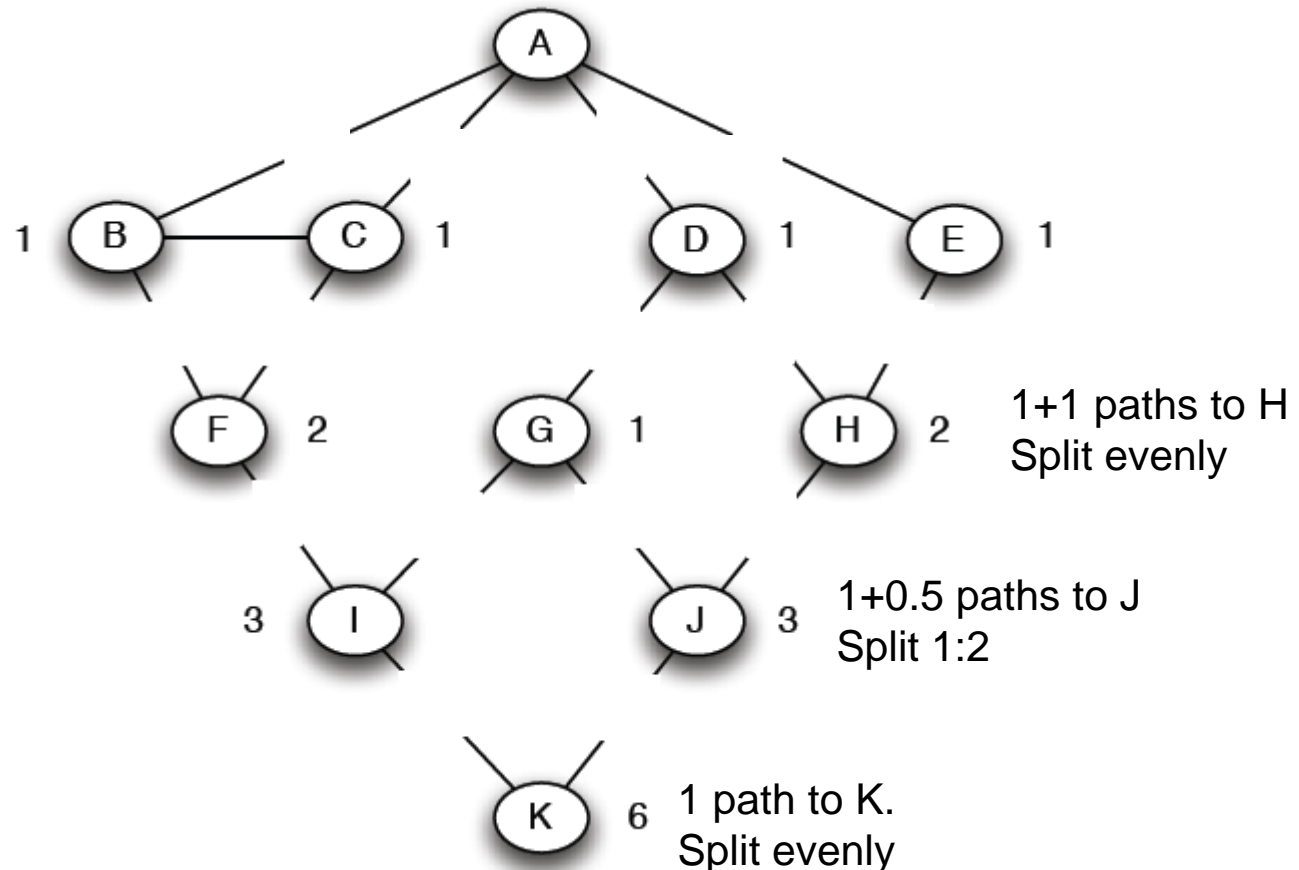
• Add edge flows:

-- node flow =

$$1 + \sum \text{child edges}$$

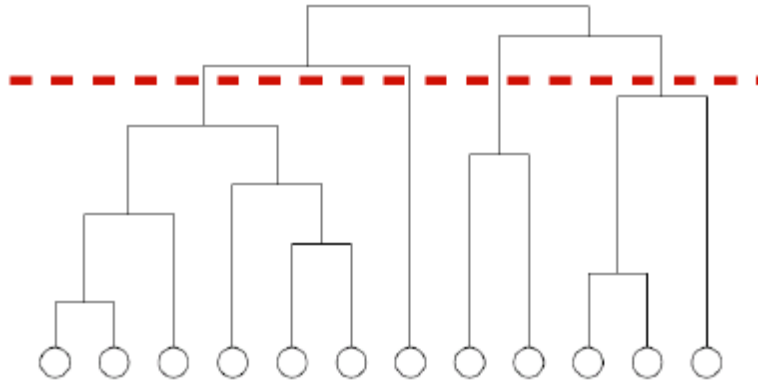
-- split the flow up based on the parent value

• Repeat the BFS procedure for each starting node U



We need to resolve 2 questions

1. **How to compute betweenness?**
2. **How to select the number of clusters?**



Network Communities

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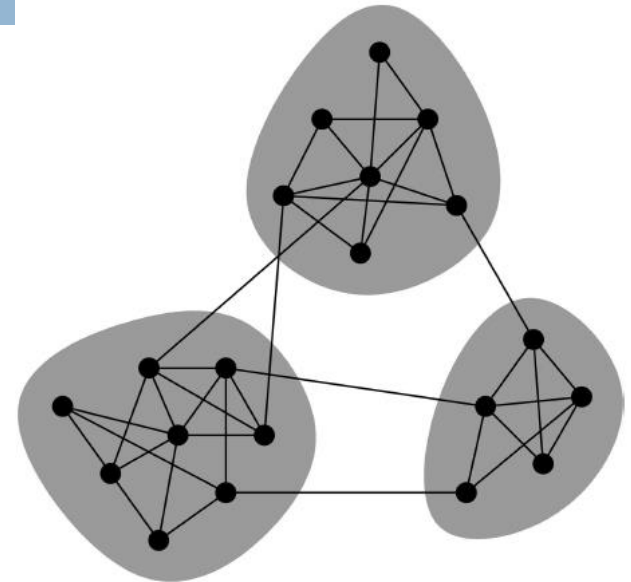
□ **Communities:** sets of tightly connected nodes

□ Define: **Modularity Q**

- A measure of how well a network is partitioned into communities
- Given a partitioning of the network into groups $s \in S$:

$$Q \propto \sum_{s \in S} [(\# \text{ edges within group } s) - (\text{expected } \# \text{ edges within group } s)]$$

Need a null model!



Null Model: Configuration Model

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□ Given real on nodes and edges, construct rewired network

- Same degree distribution but random connections

- Consider as a **multigraph**

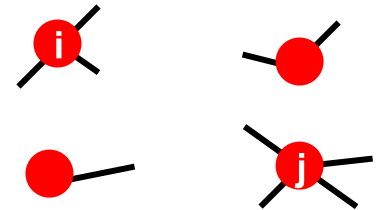
- The expected number of edge between nodes

i and *j* of degrees k_i and k_j equals to: $k_i \cdot \frac{k_j}{2m} = \frac{k_i k_j}{2m}$

- The expected number of edges in (multigraph) G' :

$$\blacksquare = \frac{1}{2} \sum_{i \in N} \sum_{j \in N} \frac{k_i k_j}{2m} = \frac{1}{2} \cdot \frac{1}{2m} \sum_{i \in N} k_i \left(\sum_{j \in N} k_j \right) =$$

$$\blacksquare = \frac{1}{4m} 2m \cdot 2m = m$$



Note:

$$\sum_{u \in N} k_u = 2m$$

Modularity

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□ Modularity of partitioning S of graph G :

□ $Q \propto \sum_{s \in S} [(\# \text{ edges within group } s) - (\text{expected } \# \text{ edges within group } s)]$

□
$$Q(G, S) = \underbrace{\frac{1}{2m} \sum_{s \in S} \sum_{i \in s} \sum_{j \in s} \left(A_{ij} - \frac{k_i k_j}{2m} \right)}_{\text{Normalizing cost.}}$$

Normalizing cost.: $-1 < Q < 1$

$A_{ij} = 1$ if $i \rightarrow j$,
0 else

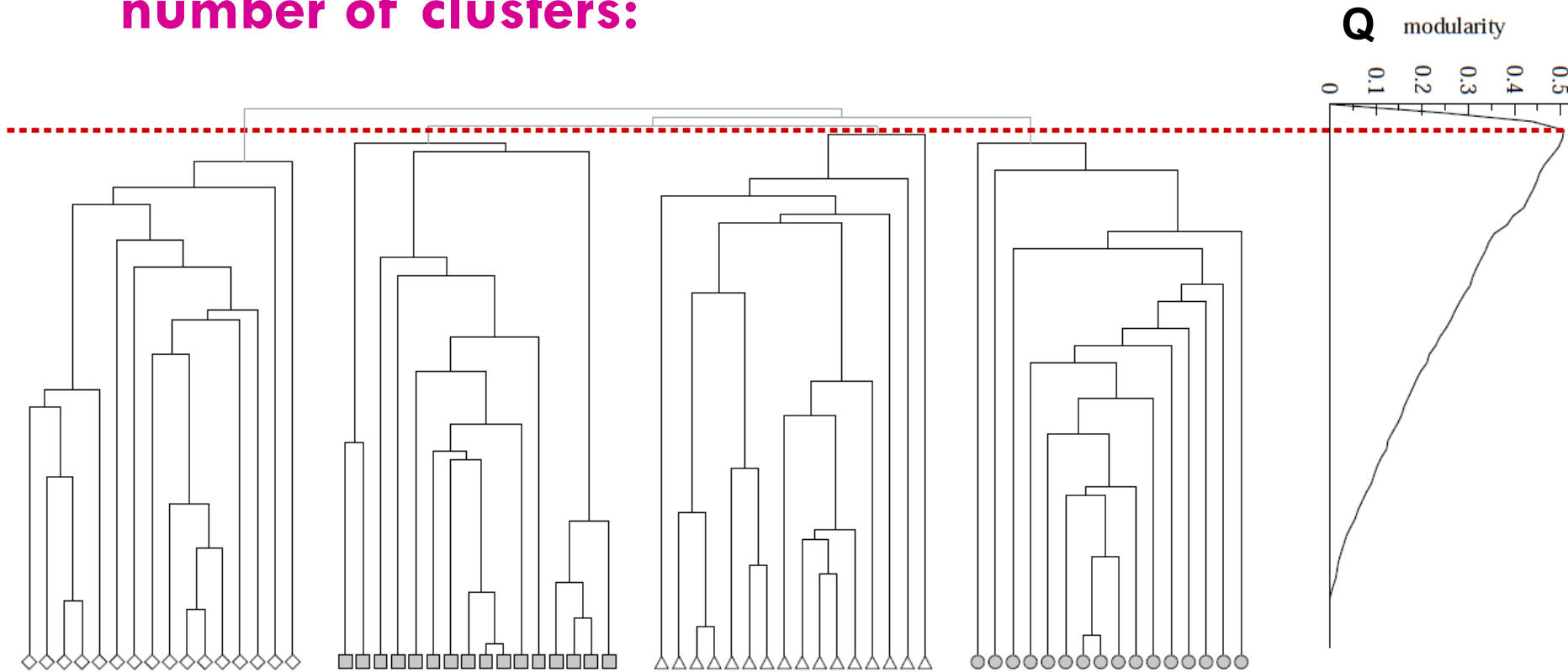
□ Modularity values take range $[-1, 1]$

- It is positive if the number of edges within groups exceeds the expected number
- $0.3 < Q < 0.7$ means significant community structure

Modularity: Number of clusters

45

- **Modularity is useful for selecting the number of clusters:**



Why not optimize modularity directly?