

ADVANCED

FUNCTIONAL PROGRAMMING

MEET YOUR INSTRUCTOR

- ▶ Ekaterina Verbitskaia
- ▶ Researcher @ JetBrains Research PLAN, Programming Languages and Program Analysis Lab
- ▶ Interested in:
 - ▶ Functional and Relational Programming
 - ▶ Partial Evaluation, Metacomputations
- ▶ @kajigor



ORGANIZATION

- ▶ 2 classes per week
- ▶ Written exam (50% of your mark)
- ▶ Work during the term (50% of your mark)
 - ▶ Homework assignments (40/100 points)
 - ▶ 75-minutes presentation on a research topic of your choice (40/100 points)
 - ▶ Active participation during classes (40/100 points)



WHAT YOU NEED TO START

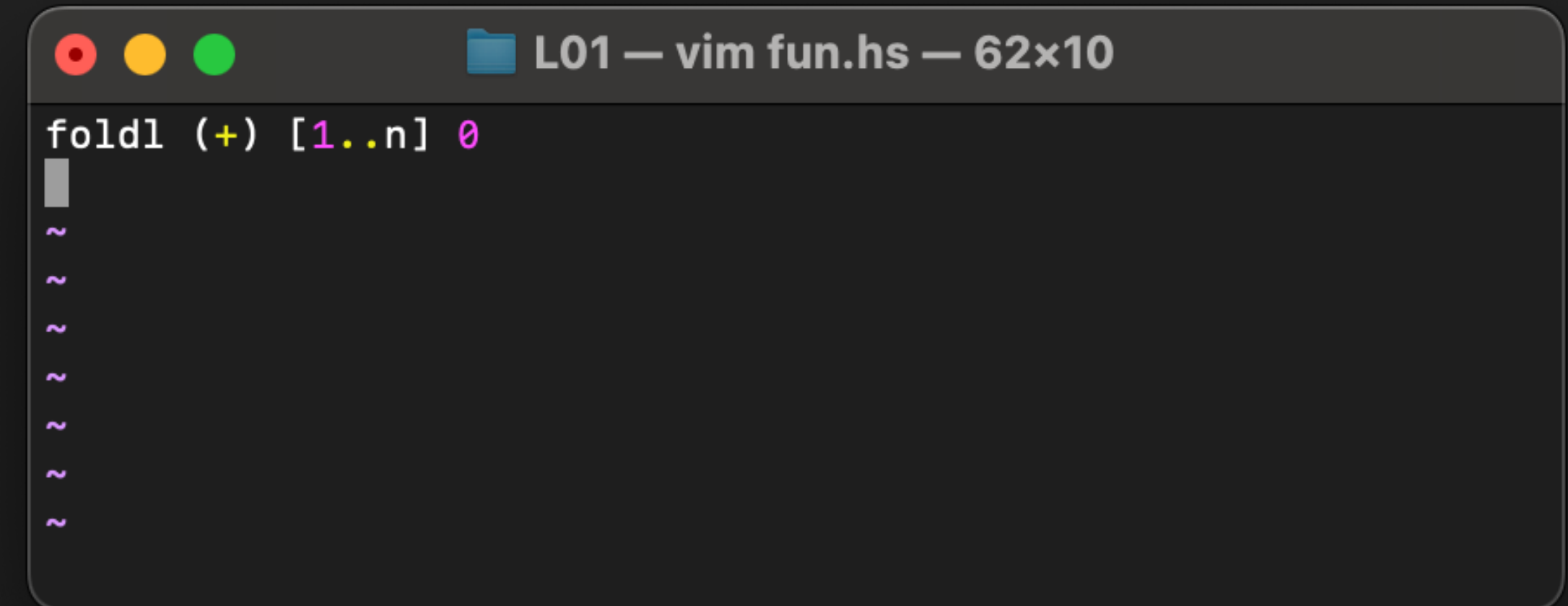
- ▶ Install with [GHCup](#):
 - ▶ [GHC](#) 9.6.6 – compiler
 - ▶ [Stack](#) 3.1.1 – build system
 - ▶ [HLS](#) 2.9.0.1 – language server
- ▶ [VSCode](#)
 - ▶ Haskell [extension](#)



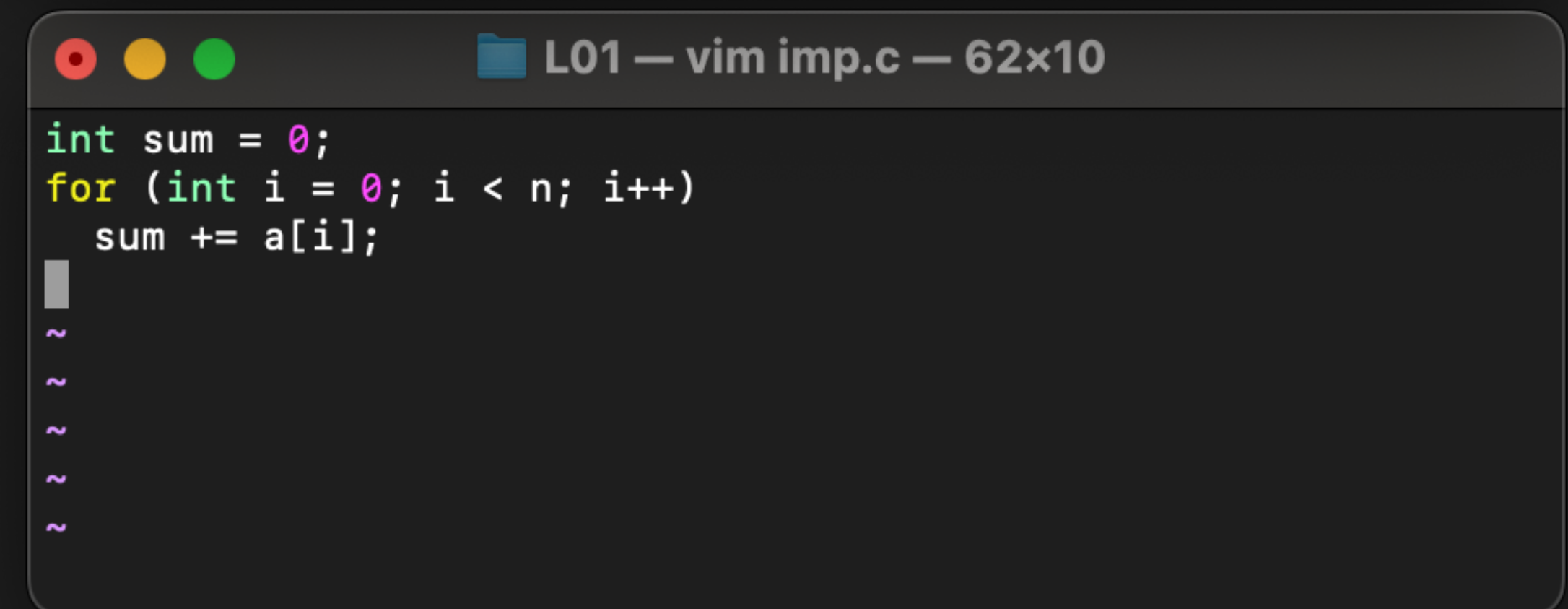
WHAT IS FUNCTIONAL PROGRAMMING?

WHAT IS FUNCTIONAL PROGRAMMING?

- ▶ Functional programming computes by means of evaluation of functions
- ▶ Imperative programming uses statements to change the program's state



A screenshot of a vim editor window titled "L01 — vim fun.hs — 62x10". The window displays a Haskell code snippet using the `foldl` function. The first line is `foldl (+) [1..n] 0`, where `(+)` is highlighted in pink and `0` is highlighted in purple. Below this line, there are several lines of tilde (~) characters, indicating that the code continues on multiple lines.

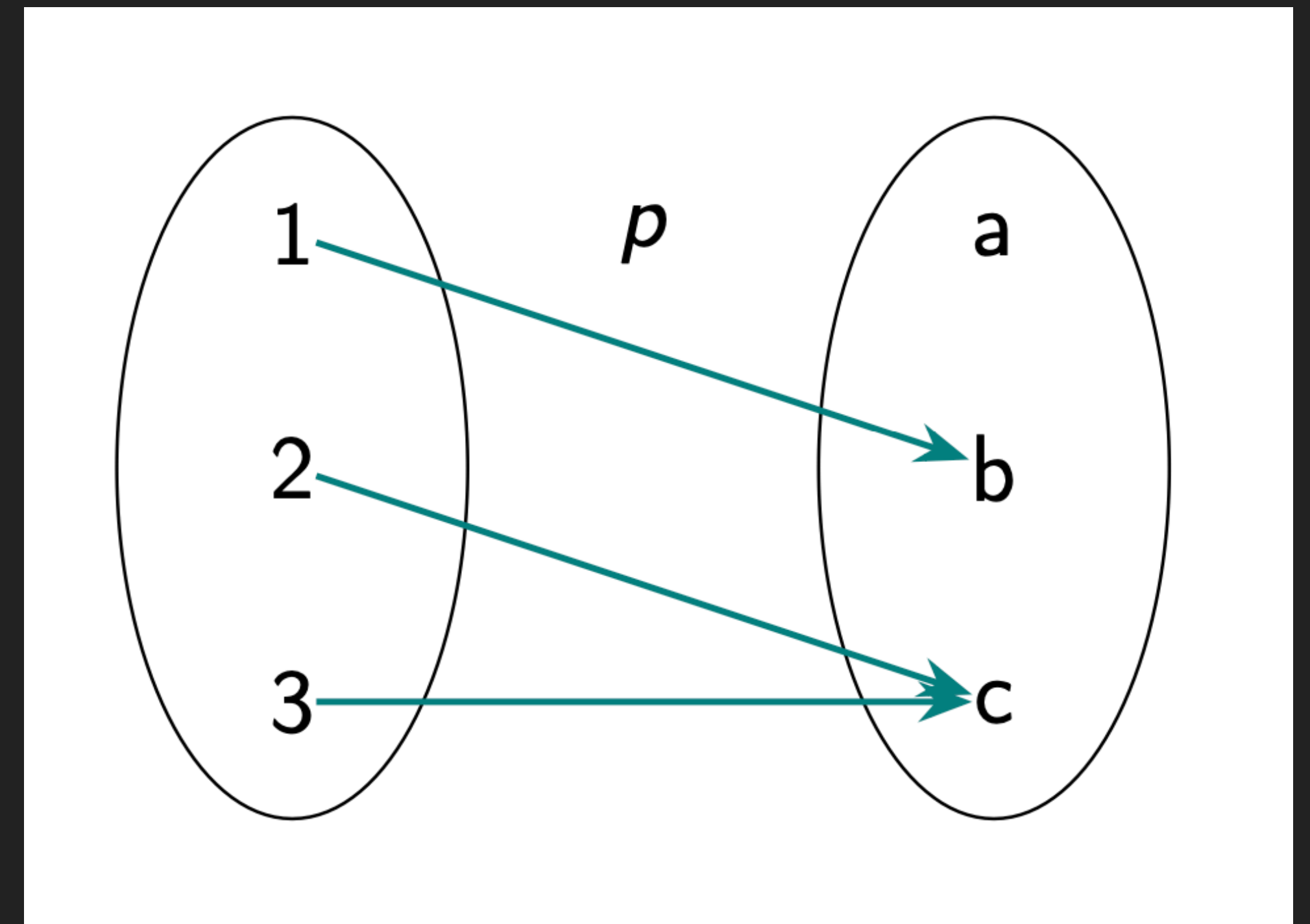


A screenshot of a vim editor window titled "L01 — vim imp.c — 62x10". The window displays a C code snippet using a `for` loop. The first line is `int sum = 0;`, where `int` is highlighted in green and `0` is highlighted in purple. The second line is `for (int i = 0; i < n; i++)`, where `for` is highlighted in yellow, `int` is highlighted in green, `i = 0` is highlighted in purple, and `i < n` is highlighted in green. The third line is `sum += a[i];`, where `sum` is highlighted in green and `a[i]` is highlighted in purple. Below these lines, there are several lines of tilde (~) characters, indicating that the code continues on multiple lines.

WHAT IS A FUNCTION?

WHAT IS A FUNCTION?

- ▶ Fundamental building block of a program
- ▶ Pure
 - ▶ No side-effects such as IO or mutable state
 - ▶ **Always** the same result for the same input
- ▶ First-class: functions can be passed as arguments and returned as results
- ▶ Composition to create new functions



IS MAIN A PURE FUNCTION?

IS MAIN A PURE FUNCTION?

- ▶ No
- ▶ `main :: IO ()`
 - ▶ Does not return any (useful) value
 - ▶ We are only interested in its effect

WHAT IS A MONAD?

**A MONAD IS JUST A
MONOID IN THE
CATEGORY OF
ENDOFUNCTORS.
WHAT'S THE PROBLEM ?**

WHAT IS A MONAD?

```
class Applicative m => Monad (m :: Type -> Type) where
  (>=>) :: forall a b. m a -> (a -> m b) -> m b
```

```
(>>) :: forall a b. m a -> m b -> m b
m >> k = m >=> \_ -> k
```

```
return :: a -> m a
return = pure
```

```
return a >> k = k a
m >=> return = m
m >=> (\x -> k x >=> h) = (m >=> k) >=> h
```

```
pure = return
m1 <*> m2 = m1 >=> \x1 -> m2 >=> \x2 -> return (x1 x2)
```

```
fmap f xs = xs >=> return . f
```

**A MONAD IS JUST A
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WHAT'S THE PROBLEM ?**

WHY DO WE USE MONADS?

- ▶ To represent impure computations
 - ▶ Errors
 - ▶ Non-determinism
 - ▶ State
 - ▶ IO
- ▶ To represent sequential computation
- ▶ **To hide boilerplate behind do-notation**

INTRO TO FP RECAP

ERRORS

ERRORS

- ▶ `Maybe a`
 - ▶ Use when there is a single way to fail
- ▶ `Either e a`
 - ▶ Use a custom error type to describe what went wrong

```
instance Monad Maybe where
  Just x >>= f = f x
  Nothing >>= _ = Nothing
```

```
safeDiv :: Int → Int → Maybe Int
safeDiv _ 0 = Nothing
safeDiv x y = Just $ x `div` y
```

```
chainDiv :: [Int] → Maybe Int
chainDiv [] = Nothing
chainDiv (h : t) = foldM safeDiv h t
```


NONDETERMINISM

NONDETERMINISM

- ▶ [a]
 - ▶ The default implementation features depth-first search

```
instance Monad [] where
  xs >>= f = concat (map f xs)
```

NONDETERMINISM

- ▶ [a]
- ▶ The default implementation features depth-first search

```
instance Monad [] where
  xs >>= f = concat (map f xs)

pairs =
  [ (x, y)
  | x <- ['a' .. 'z'] ← won't progress
  | y <- [0 ..]       ← until exhausted
  ]
```


NONDETERMINISM

- ▶ `[a]`
 - ▶ The default implementation features depth-first search
- ▶ `Logic a`
 - ▶ [Paper](#)
 - ▶ Interleaving search

```
instance Monad [] where
  xs >>= f = concat (map f xs)

pairs =
  [ (x, y)
  | x <- ['a' .. 'z'] ← won't progress
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  ]
```

INTRO TO FP RECAP

ENVIRONMENTS, LOGGING, STATE

ENVIRONMENTS, LOGGING, STATE

- ▶ `State s a`
- ▶ Functions as `Reader r a + Writer w a`
- ▶ Convenient interface to manipulate mutable state

```
newtype State s a
  = State { runState :: s → (s, a) }

instance Monad (State s) where
  m >>= k = State $ \s →
    let (s', x) = runState m s in
    runState (k x) s'

get :: State s s
get = State $ \s → (s, s)

put :: s → State s ()
put s = State $ \_ → (s, ())

modify :: (s → s) → State s ()
modify f = do
  s ← get
  put (f s)
```

IO MONAD

- ▶ Implementation is hidden
- ▶ You can think of it as a State monad over some RealWorld representation

```
newtype IO a
  = IO { runIO :: RealWorld
        → (RealWorld, a) }
```

EXERCISE

- Implement an interpreter for the formulas in Reverse Polish Notation (see RPN.hs)

```
ghci> eval "1 2 +"
```

```
3
```

```
ghci> eval "1 2 *"
```

```
2
```

```
ghci> eval "1 2 -"
```

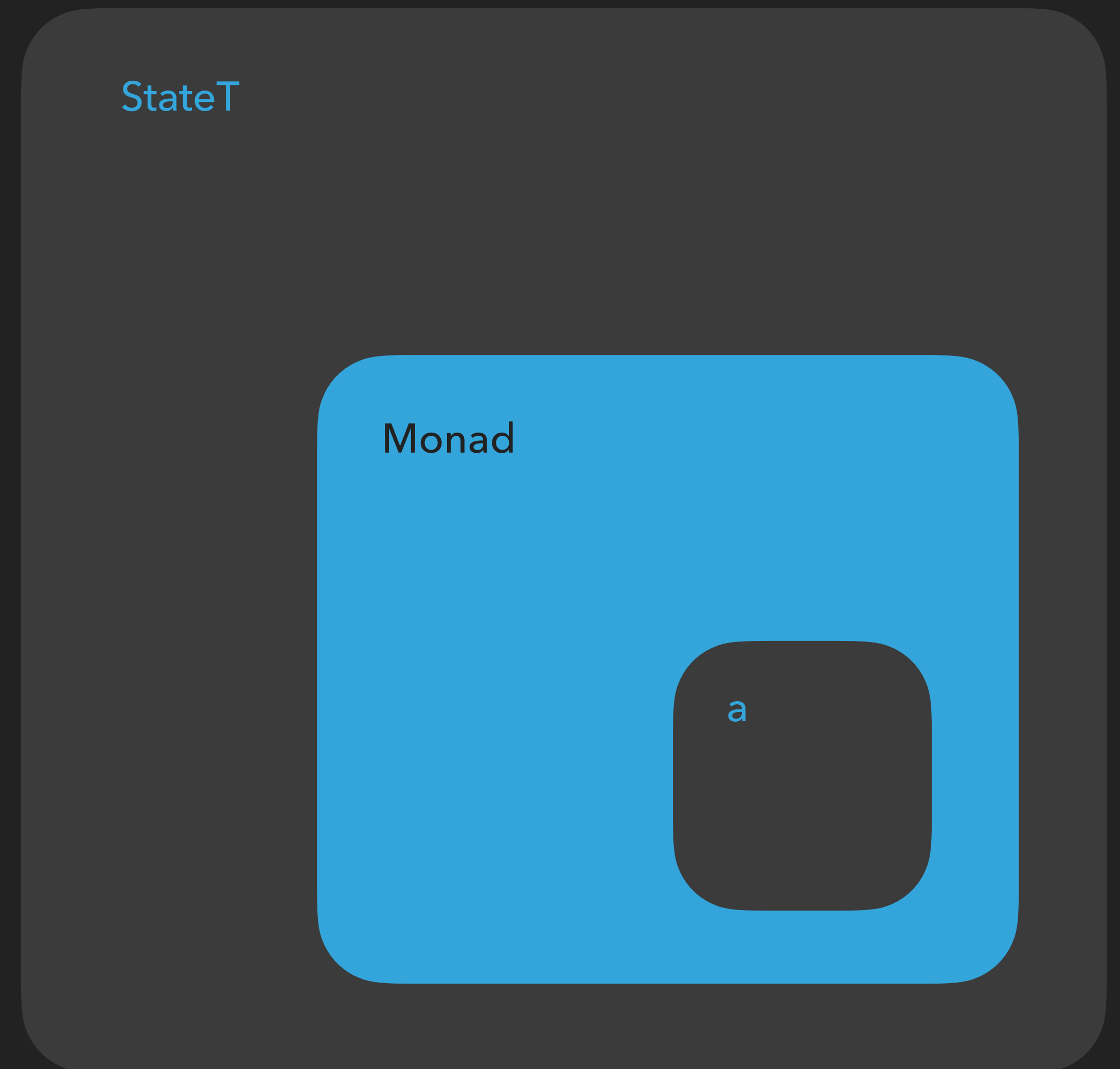
```
-1
```

```
ghci> eval "1 2 + 3 * 4 -"
```

```
5
```


HOW CAN WE COMBINE EFFECTS?

- ▶ Our `eval` function can sometimes fail, but it also needs to have access to mutable state
- ▶ Let's combine the two monads into a **stack**
- ▶ Which `>>=` do we use?



MONAD TRANSFORMER

- ▶ Should define a monad
- ▶ Should provide a way to use the features of both monads
- ▶ Should implement `lift`
 - ▶ Don't forget about the laws

```
class (forall m. Monad m => Monad (t m))  
  => MonadTrans t where  
  lift :: Monad m => m a -> t m a  
  
  -- lift . return = return  
  -- lift (m >=> f) = lift m >=> (lift . f)
```

THE ORDER OF MONADS IN A STACK

▶ StateT Maybe a

▶ $\sim s \rightarrow \text{Maybe } (a, s)$

▶ We lose both state and the result in case of an error

▶ MaybeT (State s)

▶ $\sim s \rightarrow (\text{Maybe } a, s)$

▶ Error only in the result; state survives

```
class (forall m. Monad m => Monad (t m))
  => MonadTrans t where
  lift :: Monad m => m a -> t m a

  -- lift . return = return
  -- lift (m >=> f) = lift m >=> (lift . f)
```

EXERCISE

- ▶ Implement `MyMaybeT` – monad transformer for `Maybe`
- ▶ Redefine the evaluator for RPN to use `MyMaybeT` as the outer monad

```
class (forall m. Monad m => Monad (t m))
  => MonadTrans t where
  lift :: Monad m => m a -> t m a

  -- lift . return = return
  -- lift (m >=> f) = lift m >=> (lift . f)

newtype MyMaybeT m a =
  MyMaybeT { runMyMaybeT :: m (Maybe a) }
```


LIFTING IS EXHAUSTING

- ▶ Using `lift` means we delegate some functionality to another monad
 - ▶ `lift (lift m)`
 - ▶ `lift (lift (lift m))`
- ▶ This is fragile if we change the monad stack
- ▶ Solution: use monad-specific interfaces
 - ▶ Any `MonadState` method within a monad stack build with `MaybeT` is silently lifted

```
class Monad m => MonadState s m | m -> s where
  get :: m s
  get = state (\s -> (s, s))

  put :: s -> m ()
  put s = state (\_ -> ((), s))

  state :: (s -> (a, s)) -> m a

instance MonadState s m =>
  MonadState s (MaybeT m) where

  state = lift . state
```

COMMON TRANSFORMERS

- ▶ Packages
 - ▶ [mtl](#)
 - ▶ [transformers](#)
- ▶ IdentityT
- ▶ MaybeT, ExceptT
- ▶ ReaderT, WriterT, StateT, RWST
- ▶ AccumT
- ▶ ContT

Modules

[\[Index\]](#) [\[Quick Jump\]](#)

Control

Monad

Control.Monad.Accum
Control.Monad.Cont
Control.Monad.Cont.Class

Error

Control.Monad.Error.Class
Control.Monad.Except
Control.Monad.Identity
Control.Monad.RWS
Control.Monad.RWS.CPS
Control.Monad.RWS.Class
Control.Monad.RWS.Lazy
Control.Monad.RWS.Strict

Control.Monad.Reader
Control.Monad.Reader.Class
Control.Monad.Select
Control.Monad.State
Control.Monad.State.Class
Control.Monad.State.Lazy
Control.Monad.State.Strict

Control.Monad.Trans

Control.Monad.Writer
Control.Monad.Writer.CPS
Control.Monad.Writer.Class
Control.Monad.Writer.Lazy
Control.Monad.Writer.Strict

Modules

[\[Index\]](#) [\[Quick Jump\]](#)

Control

Applicative

Control.Applicative.Backwards
Control.Applicative.Lift

Monad

IO

Control.Monad.IO.Class
Control.Monad.Signatures

Trans

Control.Monad.Trans.Accum
Control.Monad.Trans.Class
[Control.Monad.Trans.Cont](#)
Control.Monad.Trans.Except
Control.Monad.Trans.Identity
Control.Monad.Trans.Maybe
Control.Monad.Trans.RWS
Control.Monad.Trans.RWS.CPS
Control.Monad.Trans.RWS.Lazy
Control.Monad.Trans.RWS.Strict
Control.Monad.Trans.Reader
Control.Monad.Trans.Select
Control.Monad.Trans.State
Control.Monad.Trans.State.Lazy
Control.Monad.Trans.State.Strict
Control.Monad.Trans.Writer
Control.Monad.Trans.Writer.CPS
Control.Monad.Trans.Writer.Lazy

EXCEPTT

- ▶ Did you miss exceptions?
- ▶ Now you can have them in Haskell

```
newtype ExceptT e m a =  
    ExceptT (m (Either e a))
```

```
throwE :: Monad m => e -> ExceptT e m a
```

```
catchE ::  
    Monad m  
=> ExceptT e m a  
-> (e -> ExceptT e' m a)  
-> ExceptT e' m a
```

```
handleE ::  
    Monad m  
=> (e -> ExceptT e' m a)  
-> ExceptT e m a  
-> ExceptT e' m a
```

ACCUMT

- ▶ Limited version to `StateT`
 - ▶ Only uses `<>` to modify state
- ▶ Or `Writer` with the additional ability to see the result of previous `tells`
- ▶ Works faster than the default `StateT`

```
newtype AccumT w m a = AccumT (w → m (a, w))
```

```
look :: m w
```

```
add :: w → m ()
```

```
accum :: (w → (a, w)) → m a
```

