

Singular Value Decomposition

An application to Big Data

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Singular Value Decomposition

Definition of SVD

Theorem

Given a matrix $A \in \mathbb{R}^{m \times n}$, it can always be found a decomposition such that

$$A = U\Sigma V^T$$

where $U \in \mathbb{R}^{m \times m}$, $V \in \mathbb{R}^{n \times n}$ and $\Sigma \in \mathbb{R}^{m \times n}$.

U and V are two orthogonal matrices and Σ is a diagonal matrix, namely:

$$(\Sigma)_{ij} = \begin{cases} 0, & i \neq j \\ \sigma_i, & i = j \end{cases}$$

where $\sigma_1 \geq \sigma_2 \geq \dots \geq \sigma_p \geq 0$, $p = \min\{m, n\}$.

Definition of SVD

The non-zero entries of Σ , denoted by σ_i , are called *singular values*.

They are arranged in a nonincreasing order by convention.

The column vectors \mathbf{u}_i of U are called *left singular vectors* and those \mathbf{v}_i of V are called *right singular vectors*.

Since in general $m \neq n$, we have:

$$A = \sum_{i=1}^p \mathbf{u}_i \sigma_i \mathbf{v}_i^T$$

Definition of SVD

Theorem

If for some r such that $1 \leq r < p$ we have

$$\sigma_1 \geq \dots \geq \sigma_r > \sigma_{r+1} = \dots = \sigma_p = 0$$

then

- $\text{rank}(A) = r$
- $A = \sum_{i=1}^r \mathbf{u}_i \sigma_i \mathbf{v}_i^T$

This means that all other $p - r$ dimensions of matrix A are linear combinations of the first r .

Definition of SVD

Lower rank approximation

Let $A \in \mathbb{R}^{m \times n}$ be a matrix whose rank is $\text{rank}(A) = r$.

If for a fixed integer value $k < r$ we define

$$A_k = \sum_{i=1}^k \sigma_i \mathbf{u}_i \mathbf{v}_i^T \quad (1)$$

and

$$\mathcal{B} = \{B \in \mathbb{R}^{m \times n} : \text{rank}(B) = k\}$$

then

$$\min_{B \in \mathcal{B}} \|A - B\|_2 = \|A - A_k\|_2 = \sigma_{k+1}$$

This result tell us that A_k represents the best approximation (considering the *spectral norm*) of rank k of matrix A .

Singular values computation

To compute the singular values, consider the transpose of A given its decomposition:

$$A^T = (U\Sigma V^T)^T = V\Sigma^T U^T$$

The symmetric matrix $A^T A$ is equal to:

$$A^T A = (V\Sigma^T U^T)(U\Sigma V^T) = V\Sigma^T \Sigma V^T$$

Furthermore, this equation can be written as:

$$A^T A V = V \Sigma^T \Sigma$$

This means that the diagonal entries of the square matrix $\Sigma^T \Sigma$, which are the square of the singular values, are the eigenvalues of matrix $A^T A$ and V is the matrix of eigenvectors.

Singular values computation

Similarly, consider the product of AA^T . It is equal to:

$$AA^T = (U\Sigma V^T)(V\Sigma^T U^T) = U\Sigma\Sigma^T U^T$$

Which means that:

$$AA^T U = U\Sigma\Sigma^T$$

Hence U is the matrix of eigenvectors of AA^T .

Since $\text{rank}(A) = r$, only the first r eigenvalues of AA^T and $A^T A$ are non-zero.

Finding eigenvalues and eigenvectors

A possible method to find eigenvalues and eigenvectors of a matrix is based on QR decompositions and this theorem:

Theorem

Suppose $A \in \mathbb{R}^{n \times n}$ is a matrix having eigenvalues $\lambda_1, \lambda_2, \dots, \lambda_n$ satisfying

$$|\lambda_1| > |\lambda_2| > \dots > |\lambda_n| \quad (2)$$

then the following sequence for $A_1 = A$ and $k = 1, 2, \dots$

$$\begin{cases} A_k = Q_k R_k \\ A_{k+1} = R_k Q_k \end{cases} \quad (3)$$

converges to an upper triangular matrix where $(A_k)_{ii} = \lambda_i$, $i = 1, 2, \dots, n$.

In case (2) is not satisfied, this sequence converges to a triangular matrix with square blocks of order at most 2 along the diagonal.

If A is symmetric, then the sequence converges to a diagonal matrix.

So a basic implementation would be like this:

```
1  while err < toll
2      [Q, R] = qr(A);
3      A = R * Q;
4
5      err = max( max( tril(A, -1) ) );
6  end
```

This method could be speed up using a technique called *shifting*:

```
1  n = length( A );
2  while err < toll
3      % A(n, n) is an usual choice, it could be any real number
4      T = A(n, n) * eye(n);
5      [Q, R] = qr( A - T );
6      A = R * Q + T;
7
8      err = max( max( tril(A, -1) ) );
9  end
```

In our case, this method must be applied to AA^T and $A^T A$, so if (3) converges then we have a diagonal matrix.

Now, consider the diagonalization of B , where $B = AA^T$ or $B = A^T A$:

$$B = P\Lambda P^{-1} = P\Lambda P^T$$

As B can be factored using (3), the matrix containing the eigenvectors must be equal to:

$$P = \prod_i Q_i = Q_1 Q_2 Q_3 \cdots$$

Hence, for every iteration $B_k = Q_k R_k$ and $B_{k+1} = R_k Q_k$, requiring each step to have a computational cost equal to $O(\frac{2n^3}{3})$.

Hessenberg Reduction

To achieve a lower computational cost, one solution consists in transforming B in a similar tridiagonal matrix using Householder matrices.

A triangular matrix can be obtained because B is symmetric.

In general, the process of transforming a matrix in a similar matrix using Householder matrices allows us to obtain a matrix in which $B_{ij} = 0$, for all $i > j$. Such a matrix is called a *Hessenberg matrix*.

Consequently, after B has been transformed, each step of the previous code can be realized using *Givens rotation matrices*.

Givens rotation matrices

A *Givens rotation* can be represented with the following matrix:

$$G_{ij} = \begin{pmatrix} 1 & \cdots & 0 & \cdots & 0 & \cdots & 0 \\ \vdots & \ddots & \vdots & & \vdots & & \vdots \\ 0 & \cdots & c & \cdots & -s & \cdots & 0 \\ \vdots & & \vdots & \ddots & \vdots & & \vdots \\ 0 & \cdots & s & \cdots & c & \cdots & 0 \\ \vdots & & \vdots & & \vdots & \ddots & \vdots \\ 0 & \cdots & 0 & \cdots & 0 & \cdots & 1 \end{pmatrix} \begin{matrix} \leftarrow i \\ \\ \leftarrow j \\ \\ \end{matrix}$$

$\begin{matrix} \uparrow & \uparrow \\ i & j \end{matrix}$

where $c = \cos \theta$ and $s = \sin \theta$, for a particular value of $\theta \in [0, 2\pi]$.

A rotation occurs in the plane spanned by the two coordinates axes i and j .

Givens rotation matrices

Using the following values:

$$\cos \theta = \frac{|b_{ii}|}{\sqrt{b_{ii}^2 + b_{ji}^2}}$$
$$\sin \theta = \operatorname{sign}\left(\frac{b_{ji}}{b_{ii}}\right) \frac{|b_{ji}|}{\sqrt{b_{ii}^2 + b_{ji}^2}}$$

when B gets left multiplied by G_{ij} , the final matrix will have the element in position (j, i) ¹ equal to 0.

Therefore, instead of factoring B at each step by the QR method, if B is tridiagonal, we can simply iterate n times by constructing $G_{k\ k+1}$ and eliminating the elements in position $(k+1, k)$ and $(k, k+1)$.

¹Note that the indices are reversed

Pseudocode for finding eigenvalues and eigenvectors

A pseudocode of what to do should clarify what we have done so far (excluding shifting for brevity):

$B, P \leftarrow \text{hessemberg}(B)$

while error > toll **do**

for $k = 1, 2, \dots, n$ **do**

 construct $G_{k\ k+1}$

$B \leftarrow G_{k\ k+1} B G_{k\ k+1}^T$

$P \leftarrow P G_{k\ k+1}^T$

end for

 update error

end while

The function implementing the Hessemberg reduction must return the matrix used for the transformation, meaning that $B = PHP^T$.






P is updated in this way because $Q_i = G_{12}^i G_{23}^i \cdots G_{n-1n}^i$, where i denotes the generic iteration for error minimization.

Prova i

```
1 function [c, s] = givens(A, i, j)
2 %GIVENS Function to compute cos and sin of Gij
3 %   Given a matrix A, this function returns a vector containing the values
4 %   of givens matrix Gij such that Gij * A is equal to A except for the
5 %   element (i, j), which will be zero.
6     c = abs( A(i, i) ) / sqrt( A(i, i)^2 + A(j, i)^2 );
7     s = - sign( A(j, i) / A(i, i) ) * abs( A(j, i) ) / sqrt( A(i, i)^2 + A
      (j, i)^2 );
8 end
```

References

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