

10.4.4.2.1

EE24BTECH11049 - Patnam Shariq Faraz Muhammed

Question:

Find the values of k for which the quadratic equation $2x^2 + kx + 5 = 0$. So they've two equal roots.

Solution:

Given Equation:

$$2x^2 + kx + 5 = 0 \quad (0.1)$$

Numerical Solution:

- If the roots are equal then the value of Discriminant is to 0

$$b^2 - 4ac = 0 \quad (0.2)$$

$$k^2 - 4 \times 2 \times 5 = 0 \quad (0.3)$$

$$k = \pm 2\sqrt{10} \quad (0.4)$$

We get two equations since there are two values for k .

$$2x^2 + 2\sqrt{10}x + 5 = 0 \quad (0.5)$$

$$2x^2 - 2\sqrt{10}x + 5 = 0 \quad (0.6)$$

- They've a single root that is $-\frac{b}{2a}$.

$$- \text{root of (0.5)} = -\sqrt{\frac{5}{2}}$$

$$- \text{root of (0.6)} = +\sqrt{\frac{5}{2}}$$

Computational Solutions:

1) Fixed Point iteration

- The fixed point iteration method in numerical analysis is used to find an approximate solution to algebraic and transcendental equations.
- **How it works?**
 - Rewrite into $X = g(x)$
 - Idea the root is $r = g(r)$, r = fixed Point
 - initial guess = x_0 , compute $g(x_0)$
 - Hopefully $x_1 = g(x_0)$ is closer to r (Occurs when it's convergence)
 - Do iteration until stop criteria

$$x_{n+1} = g(x_n) \quad (1.1)$$

– end

- **Stop Criteria:**

$$|x_{n+1} - x_n| \leq \text{tolerance} \quad (1.2)$$

$$|g(x_n) - x_n| \leq \text{tolerance} \quad (1.3)$$

- This algorithm succeeds with proper choice of $g(x)$ and must be convergent

- **Error Analysis:**

- Let r be the root, that is, $r = g(r)$
- Iteration:

$$x_{n+1} = g(x_n) \quad (1.4)$$

– Error e :

$$e_n = |x_n - r| \quad (1.5)$$

$$e_{n+1} = |x_{n+1} - r| \quad (1.6)$$

$$= |g(x_n) - g(r)| \quad (1.7)$$

$$= |g'(c)| |x_n - r| \quad \text{where, } x_n < c < r \quad (1.8)$$

$$= |g'(c)| e_n \quad (1.9)$$

– Observation:

* If $|g'(c)| < 1 \implies e_{n+1} < e_n$: convergence.

* If $|g'(c)| > 1 \implies e_{n+1} > e_n$: Divergence.

- **Convergence Condition:** There exist an interval $I = [r - c, r + c]$ for some $c > 0$ such that $|g'(x)| < 1$ on I and $x_0 \in I$

Quadratic 1: $2x^2 + 2\sqrt{10}x + 5 = 0$	Quadratic 2: $2x^2 - 2\sqrt{10}x + 5 = 0$
$g_1(x) = -\left(\sqrt{10}x + \frac{5}{2x}\right)$	$g_2(x) = -\left(-\sqrt{10}x + \frac{5}{2x}\right)$
$g_3(x) = -\frac{(2x^2+5)}{2\sqrt{10}}$	$g_4(x) = \frac{(2x^2+5)}{2\sqrt{10}}$

TABLE 1: Fixed point iteration

- For the given Question there are 2 choices for each Quadratic
- For all $g(x)$, fixed point algo diverges.
- Initial guess
Quad1: $x_0 = -5, x_1 = -3, x_2 = 0, x_3 = 2$
Quad2: $x_0 = -2, x_1 = 0, x_2 = 3, x_3 = 5$
- **Program Output:**

Quad1:

g_1

Guess 1 failed for root

Guess 2 failed for root

Failed to detect root.

Guess 1 failed for root

Guess 2 failed for root

Failed to detect root.

Fixed-point iteration failed to converge.

g_3

Guess 1 failed for root

Guess 2 failed for root

Failed to detect root.

Guess 1 failed for root

Guess 2 failed for root

Failed to detect root.

Fixed-point iteration failed to converge.

Quad2:

g_2

Guess 1 failed for root

Guess 2 failed for root

Failed to detect root.

Guess 1 failed for root

Guess 2 failed for root

Failed to detect root.

Fixed-point iteration failed to converge.

g_4

Guess 1 failed for root

Guess 2 failed for root

Failed to detect root.

Guess 1 failed for root

Guess 2 failed for root

Failed to detect root.

Fixed-point iteration failed to converge.

2) Newton's Method

- Newton's Method is an iterative numerical technique used to approximate the roots of a real-valued function. It's particularly effective when you have a good initial guess for the root.
- Newton's Method uses the idea of approximating a function by its tangent line at a given point:

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)} \quad (2.1)$$

- It is iterated until convergence.
- We can view newton's method as "optimal" fixed point because $g'(r) = 0$

$$x_{n+1} = g(x_n) = x_n - \frac{f(x_n)}{f'(x_n)} \quad (2.2)$$

$$g'(x) = \frac{f(x)f''(x)}{(f'(x))^2} \quad (2.3)$$

$$f(r) = 0 \implies g'(r) = 0 \quad (2.4)$$

- **Error Analysis:**

- Error $e_k = |x_n - r|$ From Taylor's expansion, Final result

$$e_{n+1} \leq \frac{1}{2} \max |g''(c)| e_n^2 \quad (2.5)$$

- It exhibits quadratic convergence near the root, provided the initial guess is close enough, the function is sufficiently smooth, and the derivative at the root is not zero.
- This means the error roughly squares with each iteration, making Newton's method much faster once it's near the solution.

• **Result:**

- By taking the initial guesses -3 and 0 (assuming it has two roots) for the first quadratic (0.5) we get the result in 24 iterations.

$$x = -1.5811389 \quad (2.6)$$

- By taking the initial guesses 3 and 0 (assuming it has two roots) for the second quadratic (0.6) we get the result in 24 iterations.

$$x = 1.5811389 \quad (2.7)$$

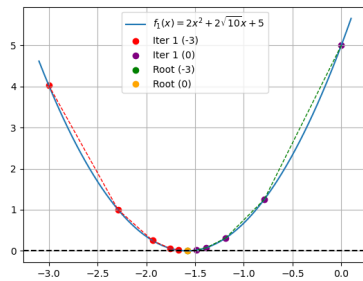


Fig. 2.1: Root of the function1

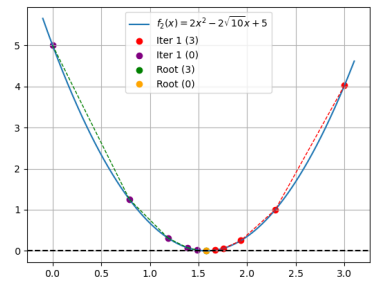


Fig. 2.2: Root of the function2

3) Secant Method

- It's a Hybrid idea(variant) of newton's and more robust.
- We need to obtain a good initial Guess x_0 And apply newton's to it.
- In this method we can avoid computing $f'(x)$
- In newton's we use an tangent line, here we use secant for approximation
- secant's iteration:

$$x_{n+1} = x_n - f(x_n) \frac{x_n - x_{n-1}}{f(x_n) - f(x_{n-1})} \quad (3.1)$$

• **Advantages:**

- No computation of $f'(x)$
- One function per iteration is calculated
- Very rapid convergence

- **Error**

$$e_{k+1} \leq M.e_k^\alpha \text{ Where } \alpha = \frac{1}{2}(1 + \sqrt{5}) \quad (3.2)$$

It is super linear convergence since $1 < \alpha < 2$

- **Failure**

- For the both the quadratic equations the secant's algo fails to find the root.
Output: Difference of the functions is too small.
- This shows us that both the Quadratics have roots as their minimum values.
- The secant method fails at minima or maxima of that function
 - * The slope between two points is almost zero.
 - * The secant line becomes nearly horizontal, and the intersection with the x-axis is poorly defined.

4) **Matrix Method**

If we consider the polynomial equation as the characteristic equation of a matrix, then by finding the eigen-values of that matrix, we can find the roots of the equation. The matrix whose eigen-values are the roots of polynomial equation is called the companion matrix of the said equation. If the given polynomial is,

$$P(x) = c_0 + c_1 x + c_2 x^2 + \dots + c_{n-1} x^{n-1} + c_n x^n \quad (4.1)$$

The companion matrix of this polynomial can be written as

$$c = \begin{pmatrix} 0 & 0 & \dots & 0 & -\frac{c_0}{c_n} \\ 1 & 0 & \dots & 0 & -\frac{c_1}{c_n} \\ 0 & 1 & \dots & 0 & -\frac{c_2}{c_n} \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & \dots & 1 & -\frac{c_{n-1}}{c_n} \end{pmatrix} \quad (4.2)$$

Power Iteration

- The Power iteration finds the largest eigen value λ_{max} of a companion matrix
- now divide the polynomial with $x - \lambda_{max}$ by performing synthetic division. This results in a new polynomial of degree one less.
- Repeat this process the degree becomes one.

The power Iteration follows the following steps iteratively:

$$\tilde{\mathbf{v}}_n = C\mathbf{v}_{n-1} \quad (4.3)$$

$$\mathbf{v}_n = \frac{\tilde{\mathbf{v}}_n}{\|\tilde{\mathbf{v}}_n\|} \quad (4.4)$$

This iteration stops when

$$|\lambda_{max}^{(n)} - \lambda_{max}^{(n-1)}| < \epsilon \quad (4.5)$$

After ϵ is tolerance, and

$$\lambda_n = \frac{\mathbf{x}_n^\top C \mathbf{x}_n}{\mathbf{x}_n^\top \mathbf{x}_n} \quad (4.6)$$

Once λ_{\max} is found, synthetic division is performed to reduce the polynomial:

$$P(x) = P(x) \div (x - \lambda_{\max}), \quad (4.7)$$

$$P(x) = c_{1,0} + c_{1,1}x + c_{1,2}x^2 + \cdots + c_{1,(n-1)}x^{n-1}. \quad (4.8)$$

Now, the new companion matrix will be evaluated.

This process is repeated iteratively until the polynomial is reduced to a degree-1 equation:

$$P_1(x) = c_{n-1,0} + c_{n-1,1}x. \quad (4.9)$$

Thus, the effective update equation would be:

$$P_{n-1}(x) = P_n(x) \div (x - \lambda_n), \quad (4.10)$$

5) Time complexity

- The construction of the companion matrix: $O(n^2)$.
- The power iteration process: $O(k \cdot n^2)$, where k is the number of iterations required for convergence.
- Synthetic division: $O(n)$.