

Rewriting Expressions As Computation

Substitution Revisited

$[y/x] e$

1. Substitute all *free* occurrences of x in e with y .
2. The substituted y remains free in e . Namely, it is not accidentally bound to an existing parameter in e .

$$\begin{array}{c}
 e = \dots x \dots x \dots x \cdot \\
 \downarrow \quad \quad \downarrow \quad \quad \downarrow \\
 [y/x] e = \dots y \dots y \dots y \cdot
 \end{array}$$

$$\begin{array}{c}
 e = \dots x \dots \lambda x. x \\
 \downarrow \quad \quad \quad \downarrow \\
 [y/x] e = \dots y \dots \lambda x. y
 \end{array}$$

Not allowed

$$\begin{array}{c}
 e = \lambda y. \dots x \dots x \\
 \quad \quad \quad \downarrow \quad \quad \downarrow \\
 [y/x] e = \lambda y. \dots y \dots y
 \end{array}$$

Not allowed

Alpha α -conversion

The alpha conversion is all about renaming parameters *correctly*.

$$\lambda x. e \rightarrow_{\alpha} \lambda y. [y/x] e$$

Provided that we observe the following:

1. We only replace free **x** in **e** with **y**
2. The name **y** must appear free in **e**

Thinking about alpha conversion in Python

We can *practice* the alpha conversion using a more familiar programming language that supports function abstraction.

```
y = 10
def add(x):
    return x + y
```



```
y = 10
def add(num):
    return num + y
```

Observe that **y** is free in the function expression, while **x** is bounded to the parameter.

We can replace **x** with **num** without violating the alpha conversion requirements.

Thinking about alpha conversion in Python

We can *practice* the alpha conversion using a more familiar programming language that supports function abstraction.

```
y = 10
def add(x):
    return x + y
```

Observe that **y** is free in the function expression, while **x** is bounded to the parameter.



```
y = 10
def add(y):
    return y + y
```

But we cannot replace **x** with **y** because it already appears in expression as a free variable.

Thinking about alpha conversion in Python

```
def quad(x):  
    def double(x):  
        return 2 * x  
    return double(2 * x)
```

_____ This **x** are bound in the body of **quad**.

_____ This **x** is free in the body of **quad**.



```
def quad(num):  
    def double(x):  
        return 2 * x  
    return double(2 * num)
```

We can replace all *free* occurrences of **x** in the body.

But we cannot replace the bound occurrences of **x**.

Try to evaluate `quad(2)`.

Thinking about alpha conversion in Python

```
def quad(x):  
    def double(x):  
        return 2 * x  
    return double(2 * x)
```

_____ This **x** are bound in the body of **quad**.

_____ This **x** is free in the body of **quad**.



```
def quad(num):  
    def double(x):  
        return 2 * num  
    return double(2 * num)
```

If we naively substitutes all occurrences of **x** in the body of **quad**, we would get into trouble.

Try to evaluate `quad(2)`.

Why alpha conversion is so important

Using alpha conversion, we can always rename the parameter of a function.

So, given two expressions, we can make sure that there are no name conflicts between the parameter names and free variable names.

$$\text{Free}(e_1) \cap \text{Bound}(e_2) = \emptyset$$

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Beta β -reduction

The beta reduction evaluates a function application expression.

It's called a reduction because it strictly reduces the length of the expression.

$$(\lambda x. e) e' \rightarrow_{\beta} [e'/x] e$$

We must ensure that the substitution is sound.

e' does not create name shadowing.

We can ensure soundness by preprocessing e :

- Rename all parameters in e .
- Rename all the parameters in e'

so there is no overlap..

Eta η -Reduction

Eta reduction can be seen as an optimization rule: it reduces a trivial function definition to a simpler form.

$$\lambda x. (e x) \rightarrow_{\eta} e$$

