

Computation and Normal Forms in LC

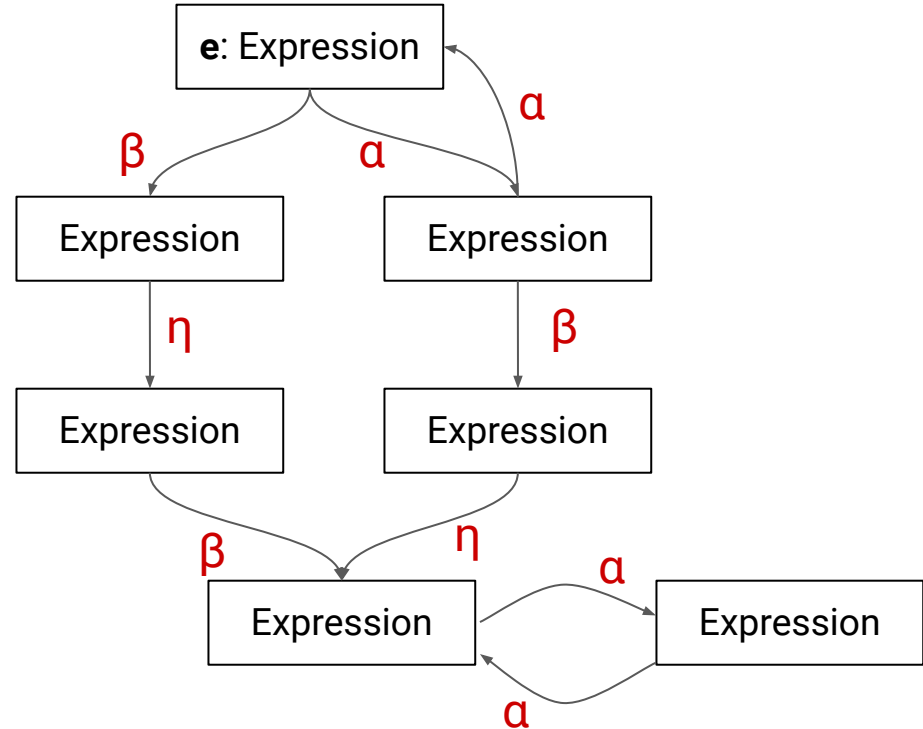
Lambda Calculus is a formal language

The expressions are strings in a formal language **LC**:

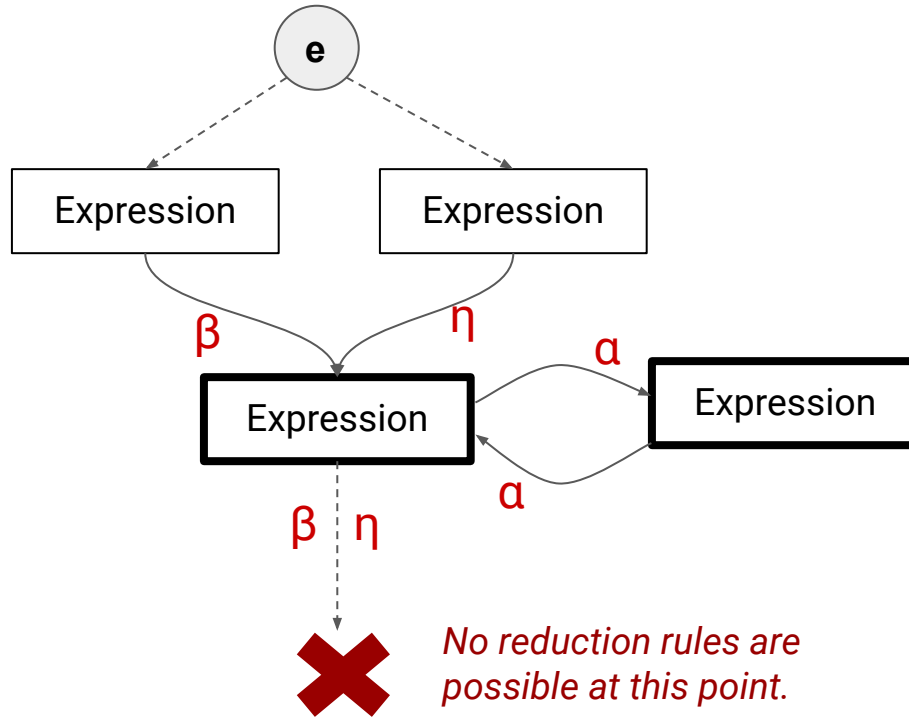
- "**<name>**" \in **LC**.
- If $e \in$ **LC**, then " **λ <name>. e**" \in **LC**
- If $e_1 \in$ **LC** and $e_2 \in$ **LC**, then "**e₁ e₂**" \in **LC**.
- If $e \in$ **LC**, then "**(e)**" \in **LC**

Rewrite rules are string operations

- Alpha conversion
- Beta reduction
- Eta reduction



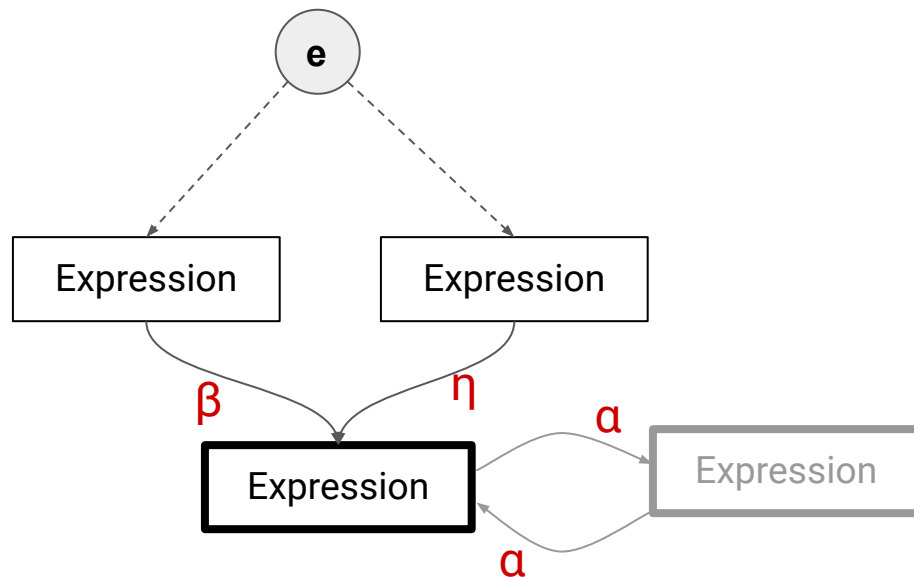
Normal forms



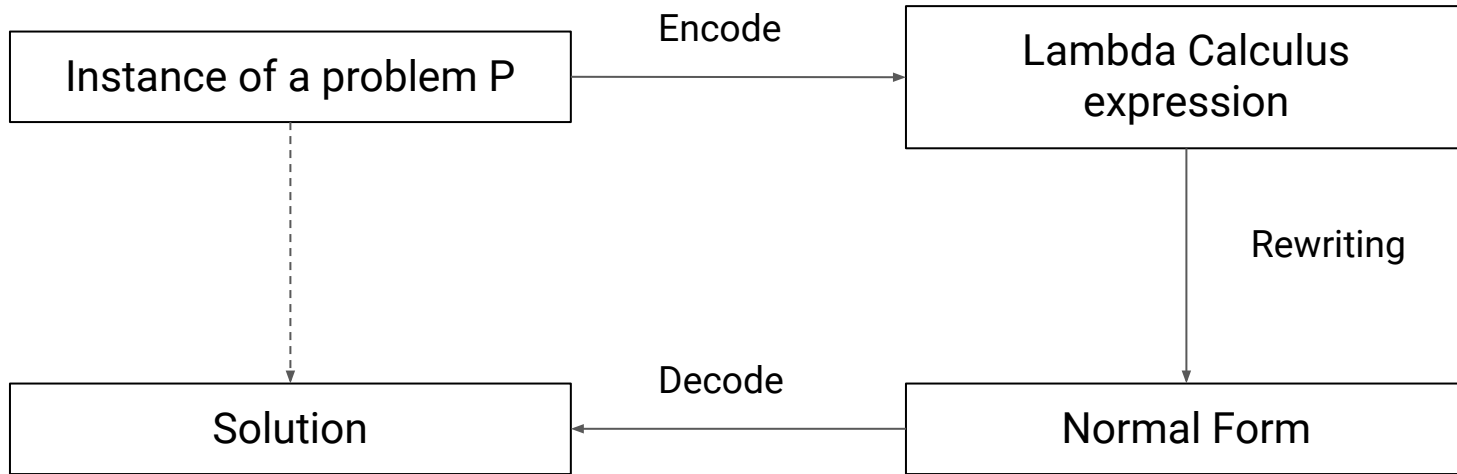
These expressions only differ in parameter names. They are equivalent.

Normal forms

The *normal form* of **e** is an expression that can be derived from **e** using alpha, beta, and eta rules, but no further beta reductions are possible.



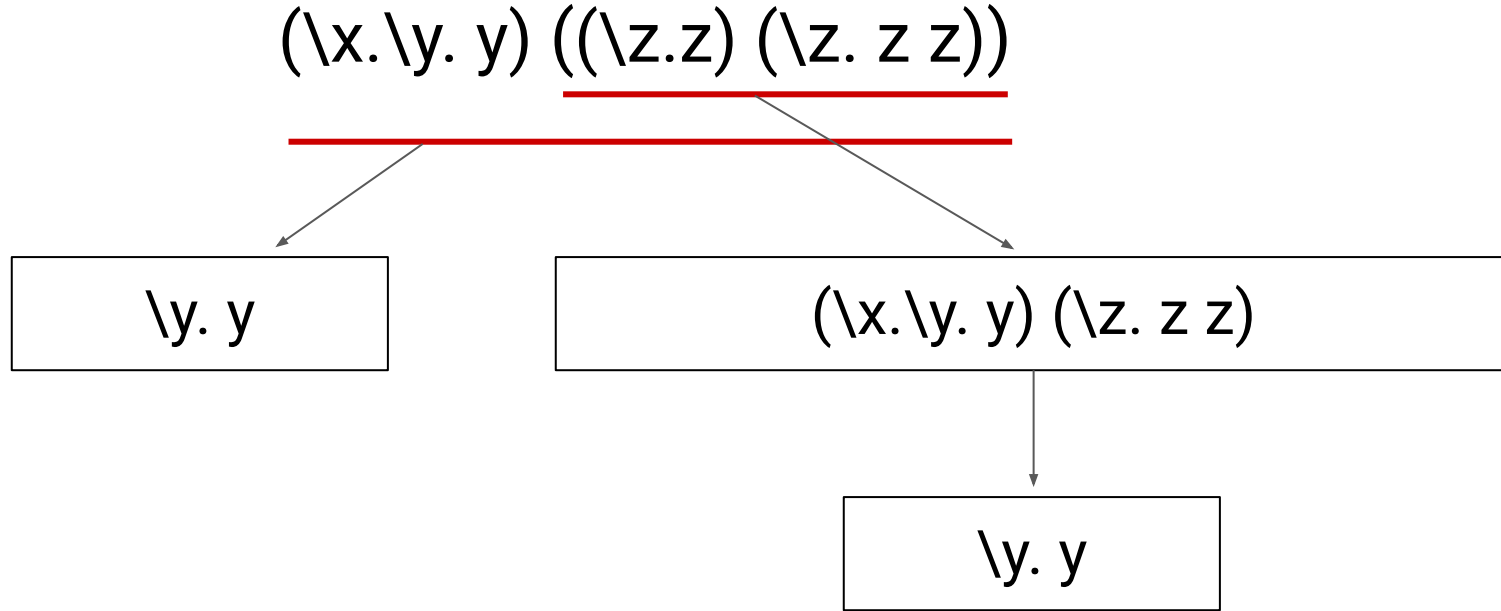
Computation in Lambda Calculus



Challenges with computation

- Ambiguity in selecting rewriting rules
- Termination

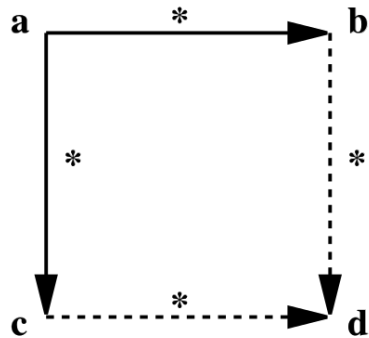
Ambiguity of reduction



Uniqueness of normal forms

Church-Rosser Theorem, 1936

If $a \rightarrow_* b$ and $a \rightarrow_* c$, then there exists some expression d such that $b \rightarrow_* d$ and $c \rightarrow_* d$.



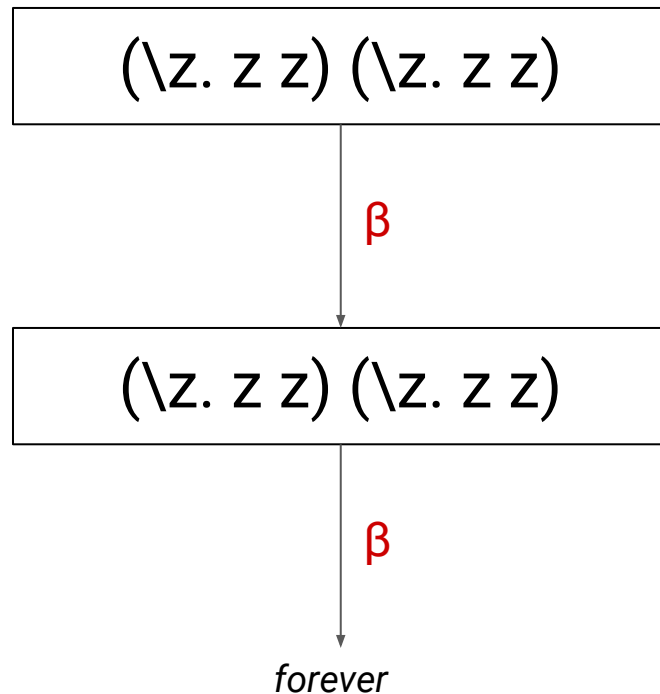
Consequence of Church-Rosser Theorem

Every expression has at most **one** normal form modulo alpha conversion.

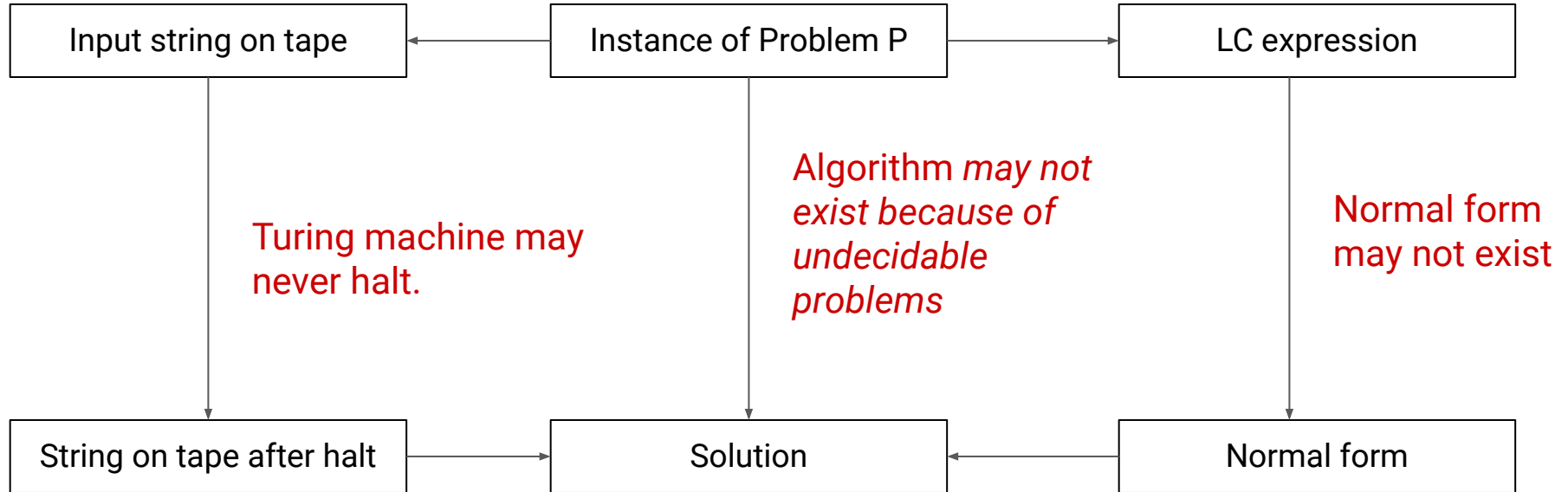
Termination

Consequence of Church-Rosser Theorem

Every expression has at most **one** normal form modulo alpha conversion.



Comparison with Turing Machine



Challenge

Consider the following decision problem:

Given a LC expression, does it have a normal form?

Can you prove that there is no general algorithm to solve this problem?

Namely, this problem is undecidable.