

# S-expressions

# S-expressions: a formal syntax from Lambda Calculus

Lambda Calculus inspired language design:

1. Data representation format
  - a. S-expression
  - b. Extended Data Notation (EDN)
2. Programming languages based on s-expression and EDN

# S-expression

## Atoms:

- Name, aka symbol  
E.g. first-name, last-name, age, x, y, z,  
Even function names are symbols:  
+, -, \*, /, ...
- Constant: number, string, character  
E.g. 3.1415, 42, "Hello world", \a

## Lists:

- Empty list: (), nil
- Parenthesis list: ( ... )  
E.g.  
(+ 2 3)  
(+ 1 2 (\* 3 4) "hello world")  
(((1) (1 2)))

# From (extended) LC to S-expressions: a first look

Extended LC:

- Multi-arity in function declaration
- Symbol binding of names to expressions

## Lambda Calculus

$\lambda$

$\lambda x.y.z \text{ <expr>}$

$\text{<name> = <expr>}$

$\text{<func> <arg1> <arg2> ...}$

## S-expression

fn

(fn (x y z) <expr>)

(def <name> <expr>)

(<func> <arg1> <arg2> ...)

*We will be using Clojure, which has an extended version of S-expression with better readability.*

# S-expressions: data representation vs functional programs

S-expressions can be describe:

1. Data
2. Program

Any well-formed S-expression is a data value.

However, only some well-formed S-expressions are valid programs in Lisp.

## Data s-expressions

```
((first-name "Ken")  
 (last-name "Pu")  
 (courses "CSCI 3055U" "CSCI 1000U"))
```

*They are not valid programs.*

```
(123 "blah" () () (((123))))
```

## Program s-expressions

```
(print "Hello world")  
  
(def f +)  
(f 1 2 3)
```

*These s-expressions are also data representations.*

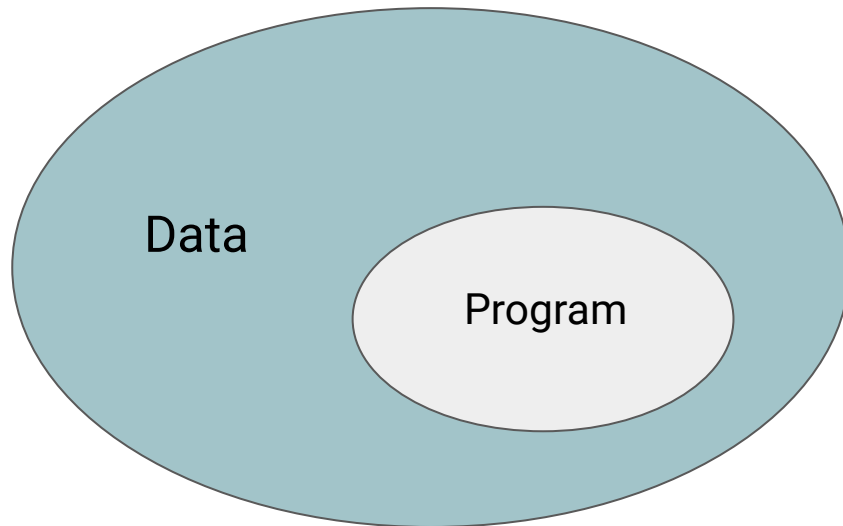
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Extended Data Notation

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A first look at Clojure's EDN

# EDN: Extensible Data Notation

EDN is used to represent:

1. Data structures
2. Clojure programs

Why?

- EDN can describe data structures efficiently. It goes beyond lists.
- EDN uses more than parenthesis to improve readability.

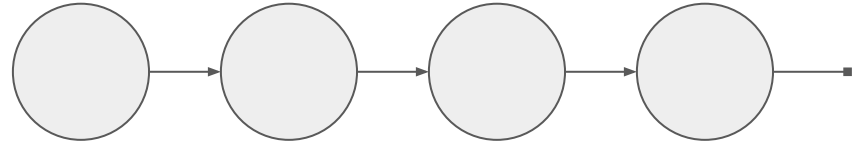
## Features in EDN:

- Vectors
- Dictionaries
- Keywords
- Comments



# S-expressions are computationally inefficient

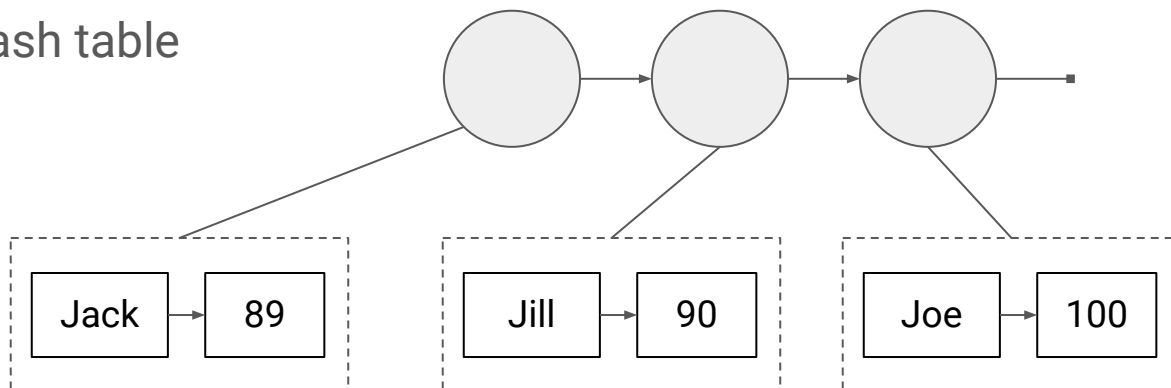
S-expressions can only lists. So classical Lisp implementations use *linked lists* as the underlying data structure to express s-expressions.



# S-expressions are computationally inefficient

So we have to *simulate* other data structures, such as a hash table using nested lists:

```
(( "Jack" 89)  
  ("Jill" 90)  
  ("Joe" 100))
```



How do we retrieve by key?

*What's the grade of "Jill"?*

Sequential scan:  **$O(n)$**

But if we use hash table, we have  **$O(1)$** . If we use balanced tree, we have  **$O(\log n)$** .

# Going beyond lists

There is **no reason** to limit the data model to lists.

- Vectors: dynamic arrays that support random access by index
- Hash-map: dynamic key-value pairs that support fast key-based lookup

Vectors in EDN:

[ ... ]

Hashmap in EDN:

{ key1 val1 key2 val2 ... }

# Going beyond basic scalars

- Consider the difference between `java.lang.String` and [Java enum](#)?

**:busy**

- Python uses strings to represent symbolic constants like:

**:age**

**if status == "busy": ...**

This is inefficient because string comparison is costly.

**:csci-3055u**

- Clojure introduces **keywords**, which are constants with efficient comparison:  $O(1)$

# Examples

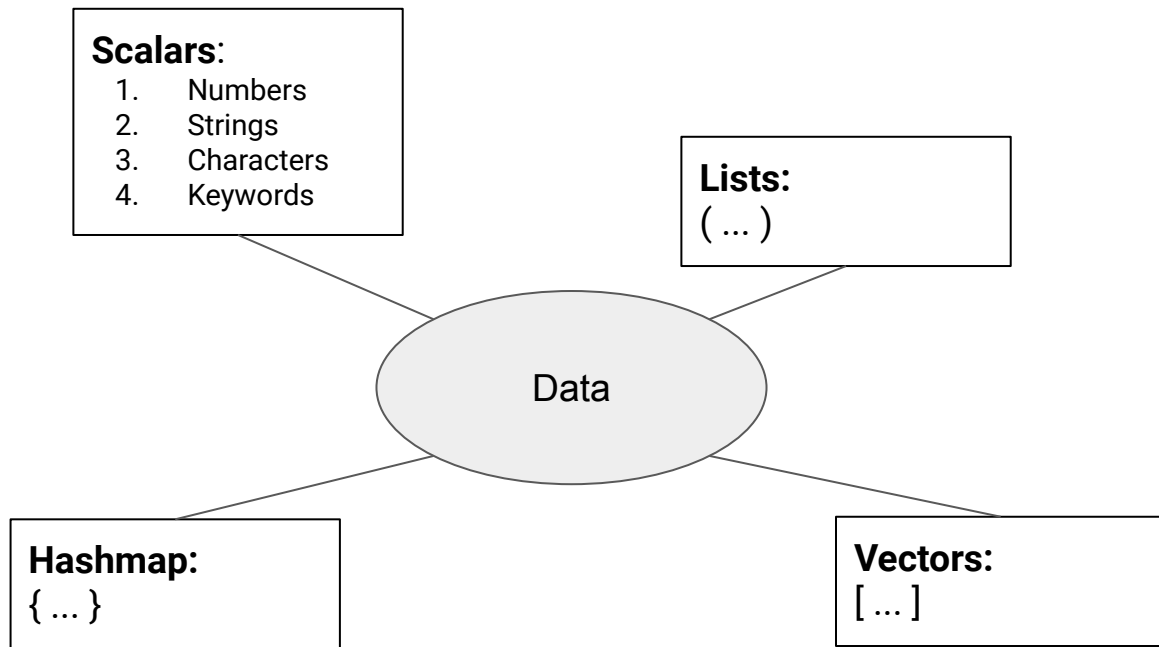
## A class of students

```
(( "Jack" ("CS" 89)
    ("Math" 76))
  ("Jill" ("CS" 90)
    ("Math" 80))
  ("Joe"  ("CS" 78)
    ("Math" 95)))
```

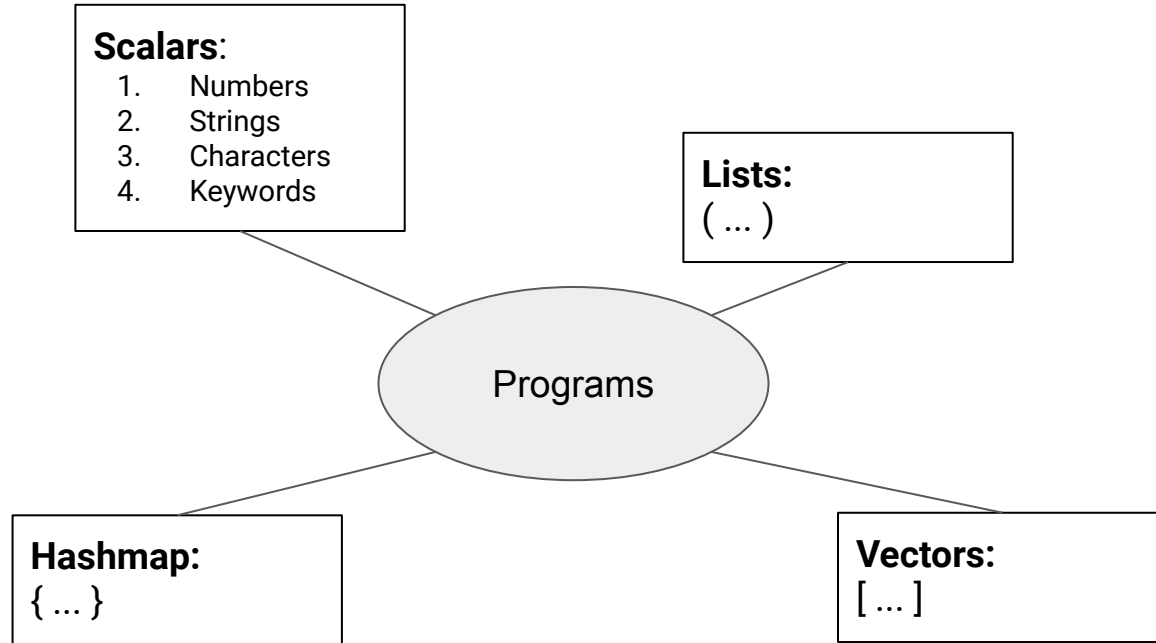
## A class of students in EDN

```
[ { :name "Jack"
    :grades { :cs 89 :math 76 } }
  { :name "Jill"
    :grades { :cs 90 :math 80 } }
  { :name "Joe"
    :grades { :cs 78 :math 95 } }
]
```

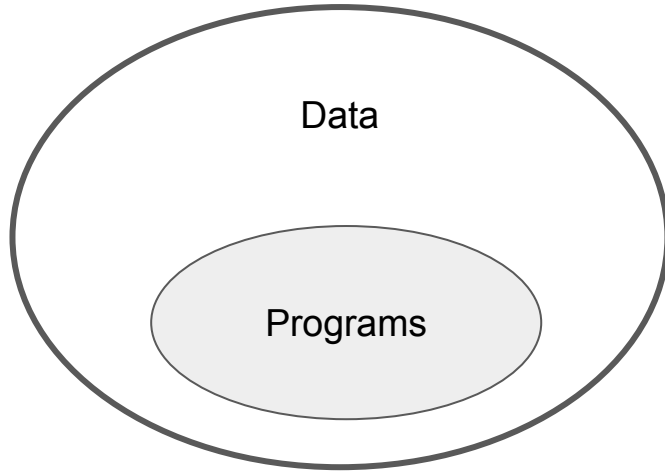
# Working with EDN



# Working with EDN



# Working with EDN



## Homoiconicity