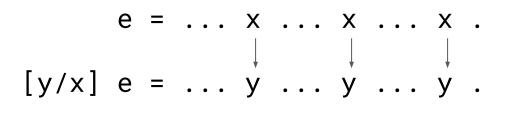
Rewriting Expressions As Computation

Substitution Revisited

- Substitute all free occurrences of x in e with y.
- 2. The substituted **y** remains free in **e**.

 Namely, it is not accidentally bound to an existing parameter in **e**.



$$e = \dots \times \dots \times \times \times$$

$$[y/x] e = \dots y \dots \times y$$
Not allowed

Alpha α-conversion

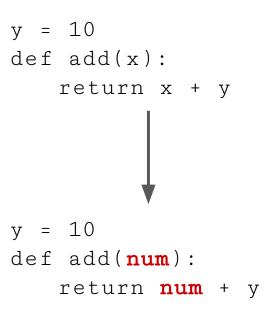
The alpha conversion is all about renaming parameters *correctly*.

$$\xspace x. e \rightarrow_{a} \yspace y. [y/x] e$$

Provided that we observe the following:

- We only replace free x in e with y
- 2. The name y must appear free in e

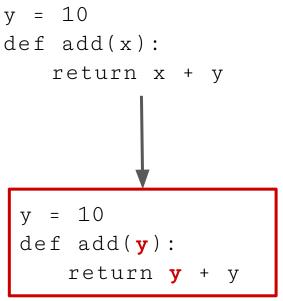
We can *practice* the alpha conversion using a more familiar programming language that supports function abstraction.



Observe that **y** is free in the function expression, while **x** is bounded to the parameter.

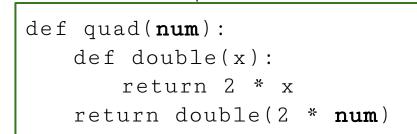
We can replace **x** with **num** without violating the alpha conversion requirements.

We can *practice* the alpha conversion using a more familiar programming language that supports function abstraction.



Observe that **y** is free in the function expression, while **x** is bounded to the parameter.

x with **y** because it already appears in expression as a free variable.



We can replace all *free* occurrences of **x** in the body.

But we cannot replace the bound occurrences of \mathbf{x} .

Try to evaluate quad(2).

Ken Q Pu, Faculty of Science, Ontario Tech University

```
def quad(num):
   def double(x):
      return 2 * num
   return double(2 * num)
```

If we naively substitutes all occurrences of **x** in the body of **quad**, we would get into trouble.

Try to evaluate quad(2).

Why alpha conversion is so important

Using alpha conversion, we can always rename the parameter of a function.

So, given two expressions, we can make sure that there are no name conflicts between the parameter names and free variable names.

$$Free(e_1) \cap Bound(e_2) = \emptyset$$

$$Bound(e_1) \cap Bound(e_2) = \emptyset$$

Bound(
$$e_1$$
) \cap Free(e_2) = \emptyset

Beta β-reduction

The beta reduction evaluates a function application expression.

It's called a reduction because it strictly reduces the length of the expression.

$$(x. e) e' \longrightarrow_{\beta} [e'/x] e$$

We must ensure that the substitution is sound.

e' does not create name shadowing.

We can ensure soundness by preprocessing e:

- Rename all parameters in e.
- Rename all the parameters in e'

so there is no overlap..

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Eta η-Reduction

Eta reduction can be seen as an optimization rule: it reduces a trivial function definition to a simpler form.

$$\xspace \xspace \xsp$$

