

# CN-Advanced L43

## Hierarchical Routing

Dr. Ram P Rustagi

[rprustagi@ksit.edu.in](mailto:rprustagi@ksit.edu.in)

<http://www.rprustagi.com>

<https://www.youtube.com/rprustagi>

# Chapter 4

## Wireless and Mobile Networks

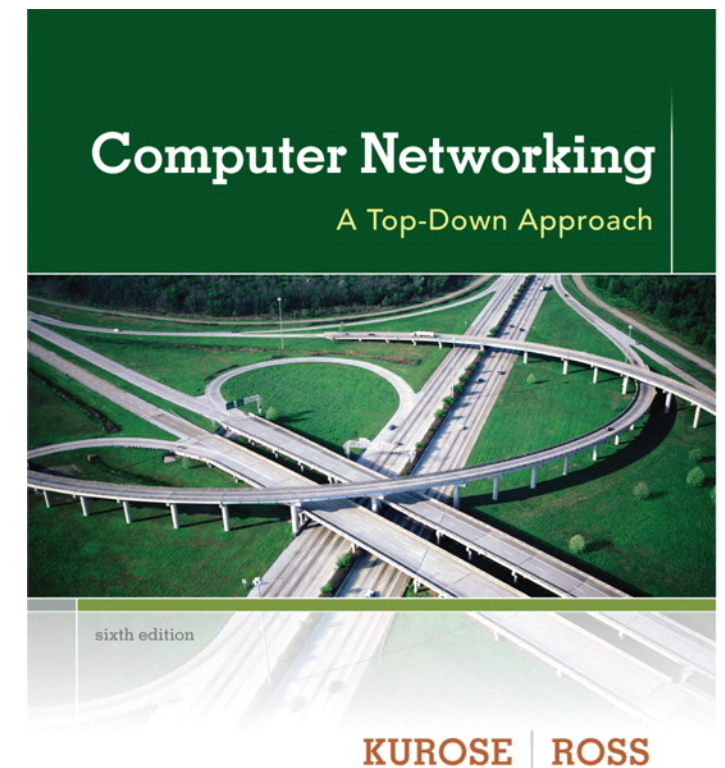
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*Computer  
Networking: A Top  
Down Approach*  
6<sup>th</sup> edition  
Jim Kurose, Keith Ross  
Addison-Wesley  
March 2012

# Hierarchical routing

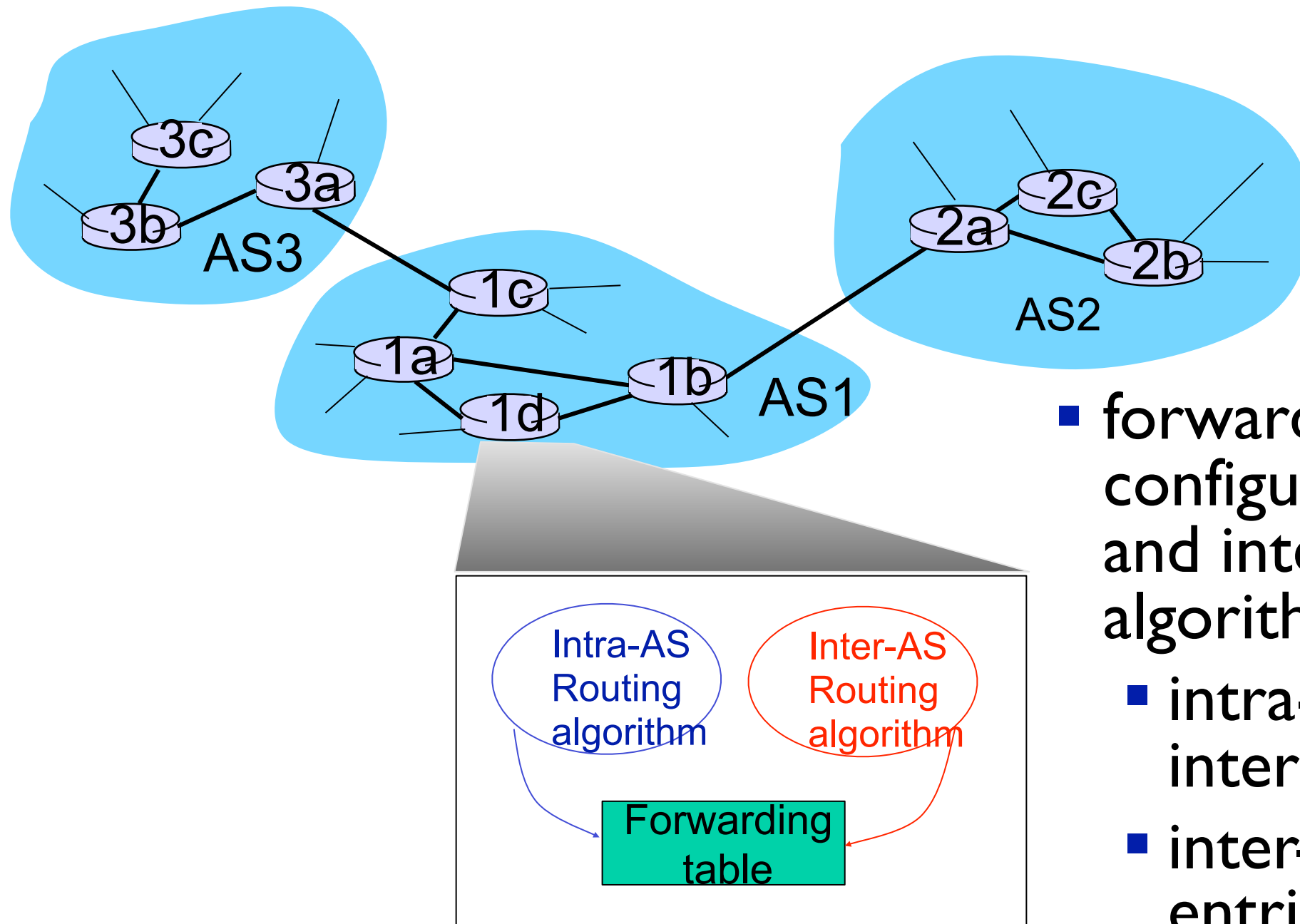
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- *scale*: with 600 million destinations:
  - can't store all dest's in routing tables!
  - routing table exchange would swamp links!
  - computing cost too high
- 
- our routing study thus far - idealization
  - all routers identical
  - network “flat”
  - ... *not* true in practice
- 
- *administrative autonomy*
  - internet = network of networks
  - each network admin may want to control routing in its own network

# Hierarchical routing

- aggregate routers into regions, “autonomous systems” (AS)
- routers in same AS run same routing protocol
  - “intra-AS” routing protocol
  - routers in different AS can run different intra-AS routing protocol
- *gateway router:*
- at “edge” of its own AS
- has link to router in another AS

# Interconnected ASes



- forwarding table configured by both intra- and inter-AS routing algorithm
  - intra-AS sets entries for internal dests
  - inter-AS & intra-AS sets entries for external dests

How is AS1 forwarding table different from those of AS2 or AS3?

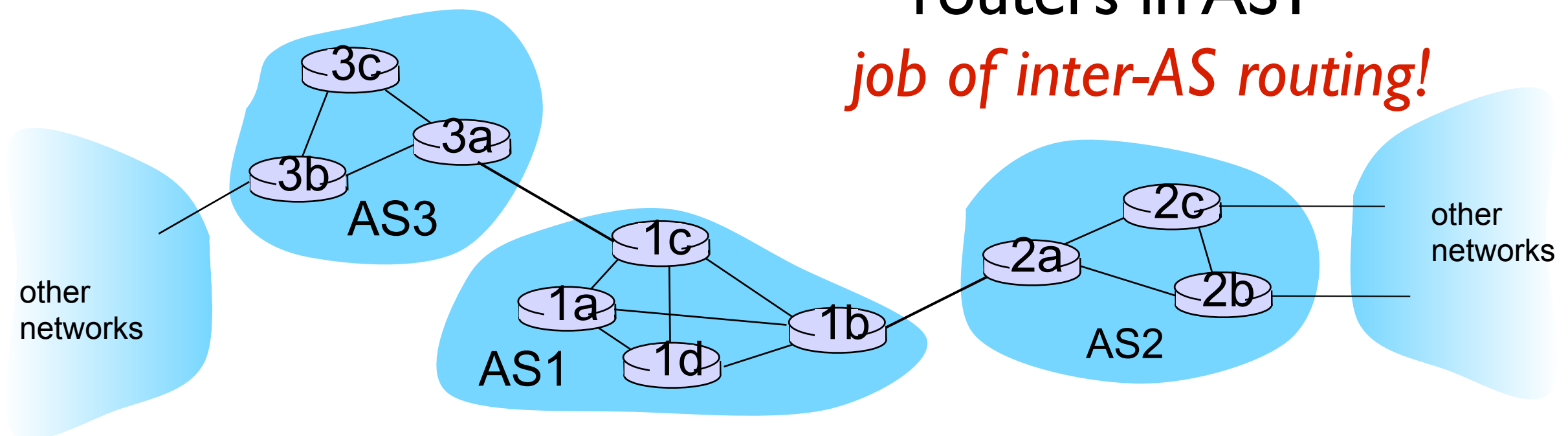
# Inter-AS tasks

- suppose router in AS1 receives datagram destined outside of AS1:
  - router should forward packet to gateway router, but which one?

*AS1 must:*

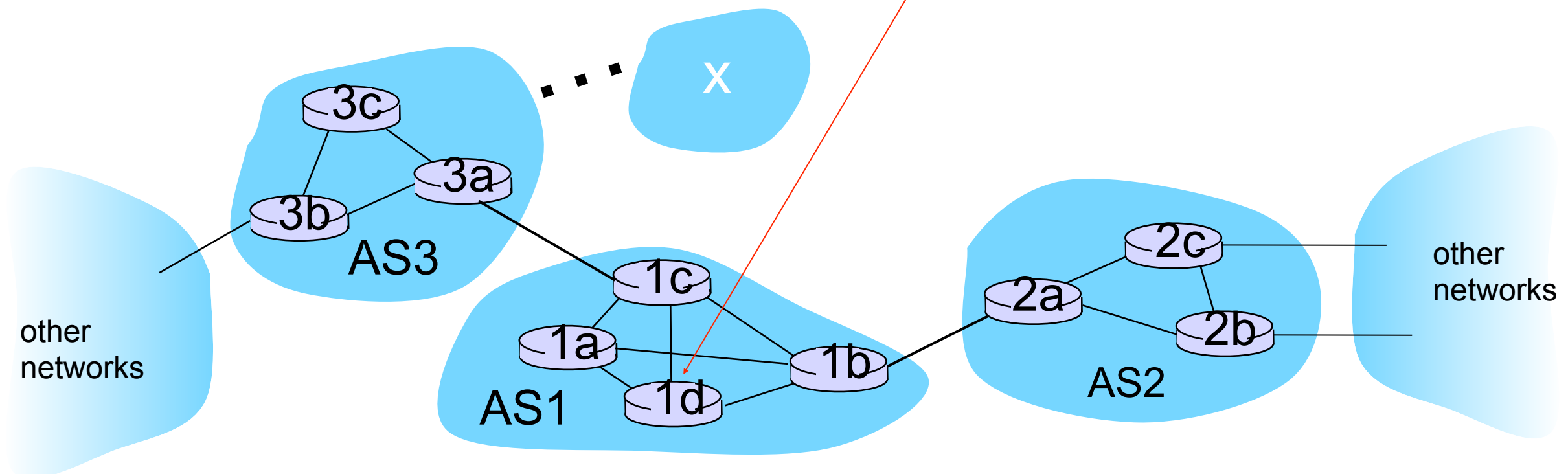
1. learn which destds are reachable through AS2, which through AS3
2. propagate this reachability info to all routers in AS1

*job of inter-AS routing!*



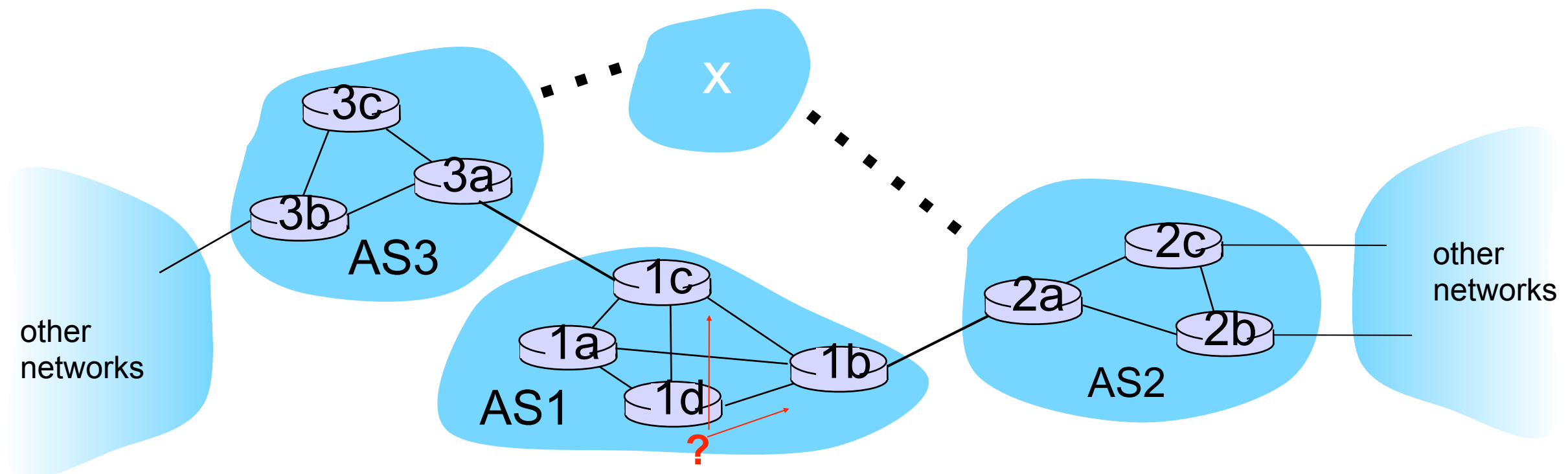
## Example: setting forwarding table in router 1d

- suppose AS1 learns (via inter-AS protocol) that subnet **x** reachable via AS3 (gateway 1c), but not via AS2
  - inter-AS protocol propagates reachability info to all internal routers
- router 1d determines from intra-AS routing info that its interface **I** is on the least cost path to 1c
  - installs forwarding table entry **(x, I)**



# Example: choosing among multiple ASes

- now suppose AS1 learns from inter-AS protocol that subnet **x** is reachable from AS3 *and* from AS2.
- to configure forwarding table, router 1d must determine which gateway it should forward packets towards for dest **x**
  - this is also job of inter-AS routing protocol!





## Example: choosing among multiple ASes

- now suppose AS1 learns from inter-AS protocol that subnet **x** is reachable from AS3 *and* from AS2.
- to configure forwarding table, router 1d must determine towards which gateway it should forward packets for dest **x**
  - this is also job of inter-AS routing protocol!
- **hot potato routing: send** packet towards closest of two routers.

# Adding Outside AS info to Forwarding Table

learn from inter-AS protocol that subnet  $x$  is reachable via multiple gateways



use routing info from intra-AS protocol to determine costs of least-cost paths to each of the gateways



determine from forwarding table the interface  $I$  that leads to least-cost gateway.



hot potato routing: choose the gateway that has the smallest least cost

Enter  $(x, I)$  in forwarding table

# ISP and AS

- Does one ISP imply one AS?
  - Some ISP break their network into multiple ASs
- Summary -
  - What AS achieves
    - Problem of scale
    - Problem of autonomy
  - Routers within AS run same Intra-AS routing protocol
  - Two ASs run same Inter-AS routing protocol
  - Inter AS protocol provides policy support

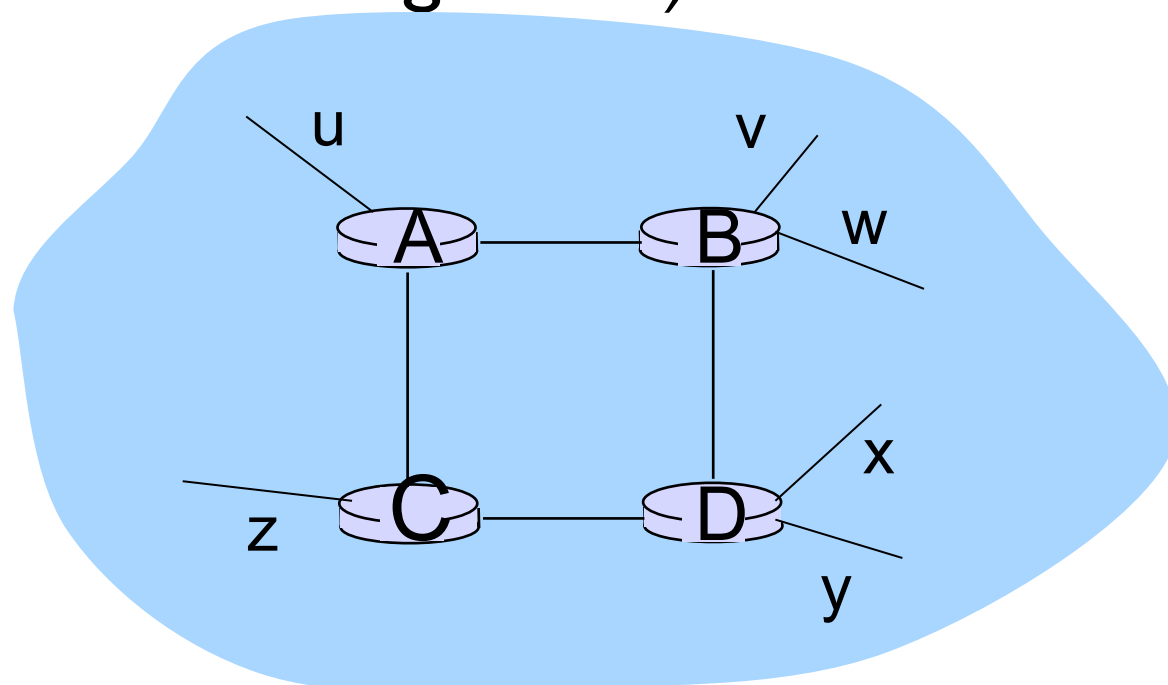
# Intra-AS Routing

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- also known as *interior gateway protocols (IGP)*
- most common intra-AS routing protocols:
  - RIP: Routing Information Protocol
    - RFC 1058, RIPv1
    - RFC 2453 RIPv2
  - OSPF: Open Shortest Path First
    - RFC 2328 OSPFv2
  - IGRP: Interior Gateway Routing Protocol (Cisco proprietary)

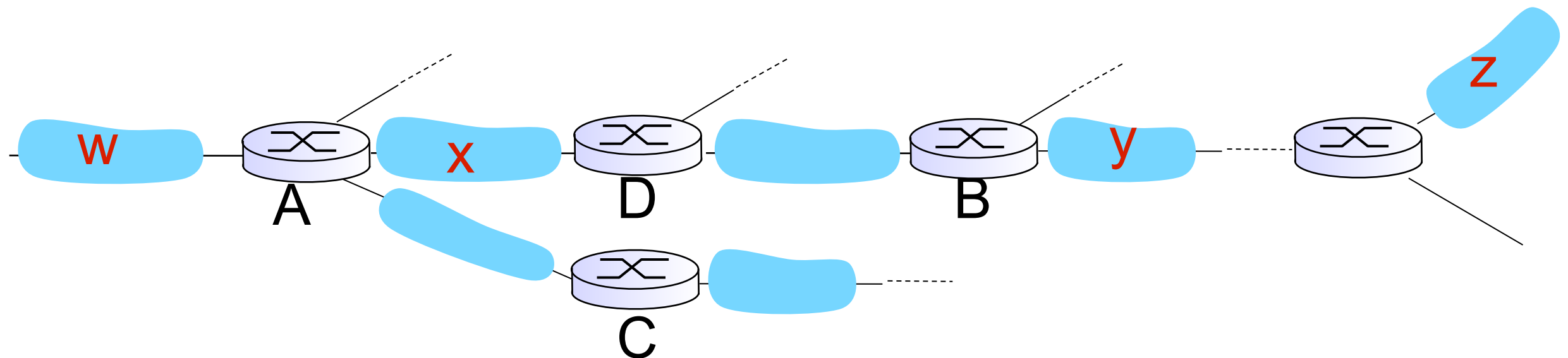
# RIP ( Routing Information Protocol)

- Included in BSD-UNIX supporting TCP/IP in 1982
  - Reason for widespread distribution
- distance vector algorithm
  - distance metric: # hops (max = 15 hops), each link has cost 1
  - DVs exchanged with neighbors every 30 sec in response message (aka **advertisement**)
  - each advertisement: list of up to 25 destination **subnets** (in IP addressing sense)  
from router A to destination **subnets**:



<u>subnet</u>	<u>hops</u>
u	1
v	2
w	2
x	3
y	3
z	2

# RIP: example

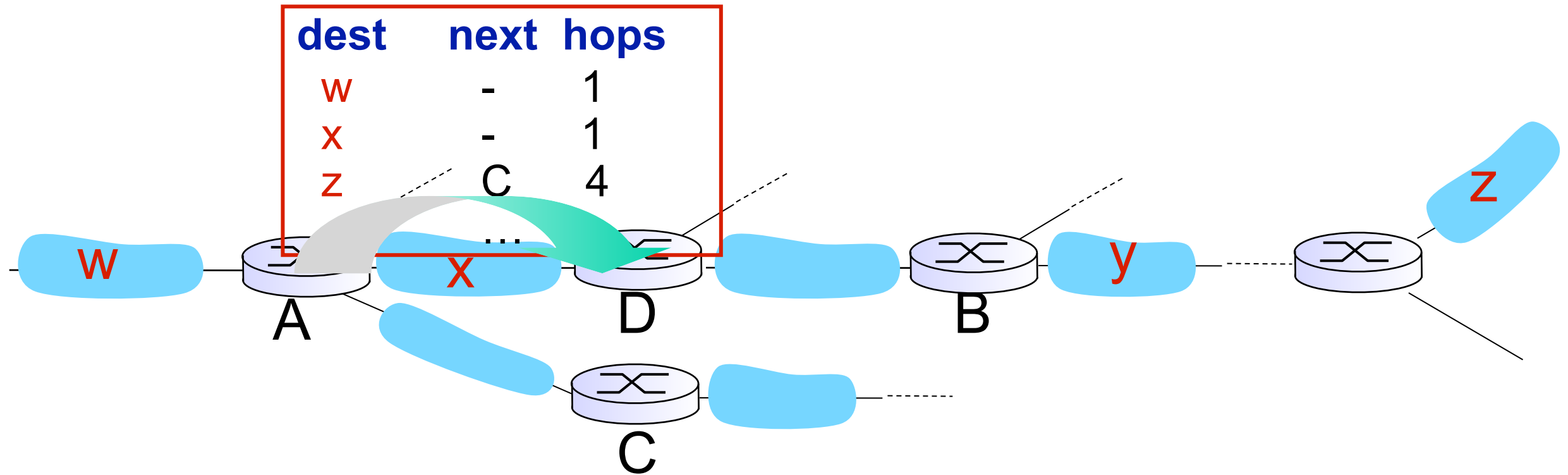


routing table in router D

destination subnet	next router	# hops to dest
<i>W</i>	A	2
<i>y</i>	B	2
<i>Z</i>	B	7
<i>X</i>	--	1
....	....	....

# RIP: example

A-to-D advertisement



routing table in router D

destination subnet	next router	# hops to dest
W	A	2
y	B	2
Z	B → A	7 → 5
X	--	1
....	....	....

# RIP: link failure, recovery

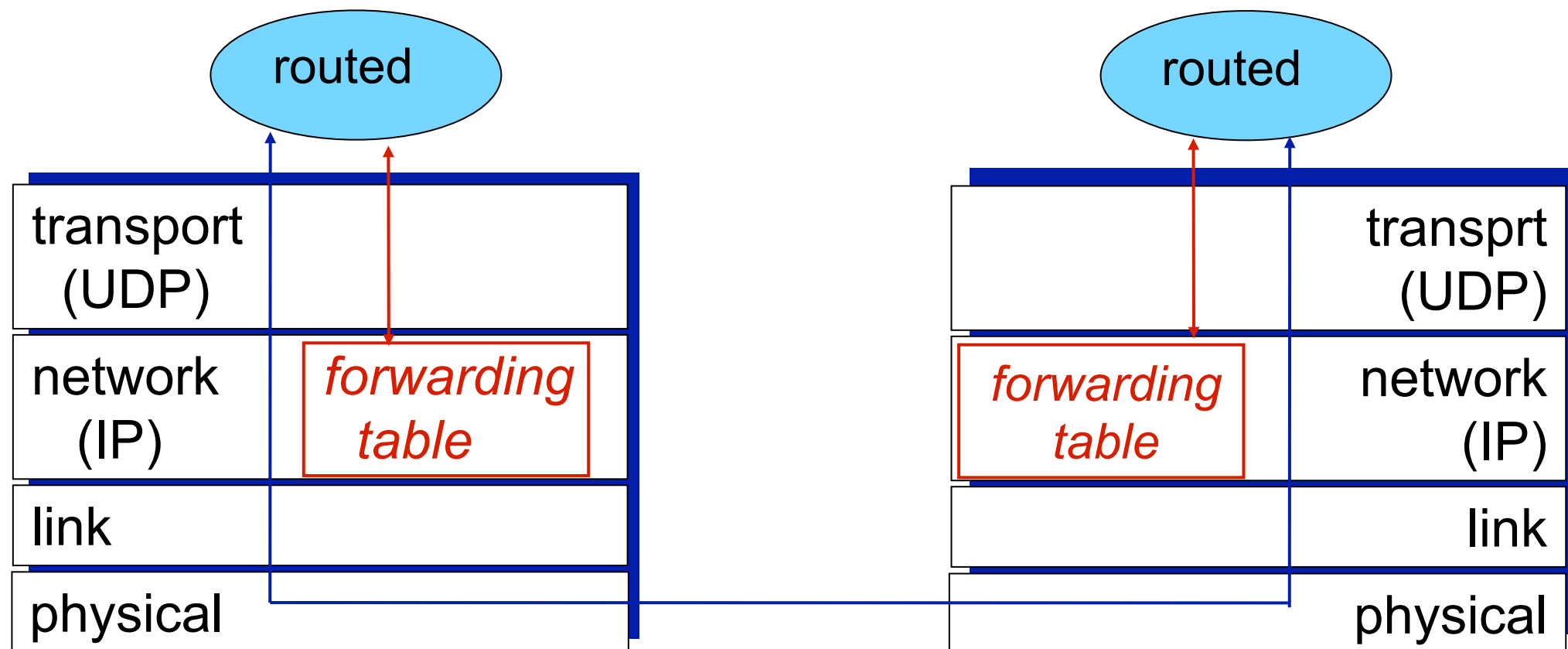
if no advertisement heard after 180 sec --> neighbor/  
link declared dead

- routes via neighbor invalidated
- new advertisements sent to neighbors
- neighbors in turn send out new advertisements
  - (if tables changed)
- link failure info quickly (?) propagates to entire net
- *poison reverse* used to prevent ping-pong loops  
(infinite distance = 16 hops)



# RIP table processing

- RIP routing tables managed by *application-level* process called route-d (daemon)
- advertisements sent in UDP packets, periodically repeated



# RIP

- **RIPv1**
  - works with classful networks
- **RIPv2**
  - works with CIDR
  - communicates network mask along with network number
  - supports aggregation and hierarchical routing

# Routing protocols on Linux

- Quagga protocol suite
  - <http://www.quagga.net>

# Summary

- Hierarchical Routing
- Interconnected ASes
- Inter AS tasks
- RIP