# Artificial Intelligence Greedy and A\* Search

Portland Data Science Group Created by Andrew Ferlitsch Community Outreach Officer June, 2017

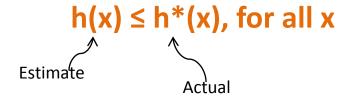


## **Greedy Algorithm**

- A search method of selecting the best local choice at each step in hopes of finding an optimal solution.
- Does not consider how optimal the current solution is.
- At each step, uses a heuristic to estimate the distance (cost) of each local choice from the goal.
- Steps:
  - 1. Define a heuristic function h(x) to estimate the distance to the goal from any state.
  - 2. From the current state, determine the search space (actions) for one step ahead.
  - 3. Select the action from the search space that minimizes the heuristic function.

# **Greedy - Completeness**

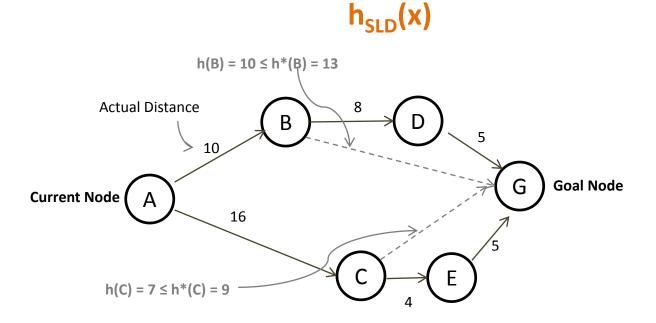
 The Greedy method will always find a solution if the heuristic function h(x), for all x, is always less than or equal the actual distance (cost) to the goal node, h\*(x).



 When the heuristic h(x) ≤ h\*(x), the method is said to be complete, since a solution will be found.

### Greedy – Straight Line Distance (SLD)

 One well-known heuristic that meets this criteria is the Straight Line Distance (SLD) in problems where the search space can be represented in a Euclidean Space.

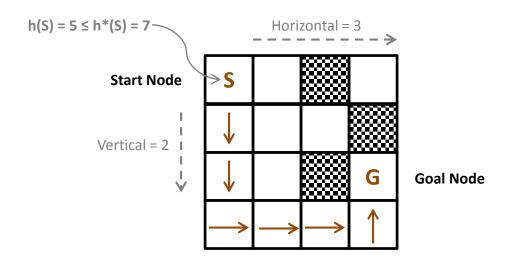


**Greedy will select node C, since h(C) < h(B)** 

## Greedy – Manhattan Distance

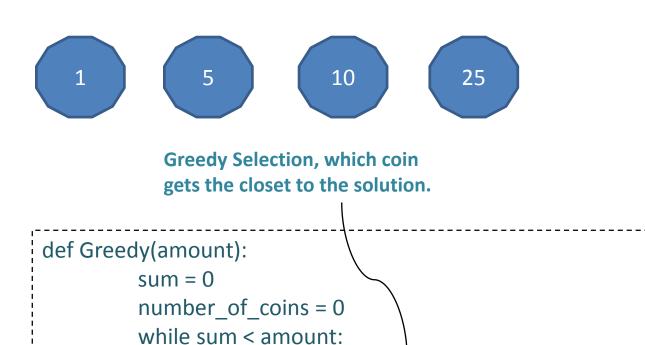
 In search spaces represented by a 2D grid (such as GridWorld), the Manhattan Distance h<sub>MAN</sub> is a heuristic which meets the criteria of always been less than or equal to the actual distance.

**h**<sub>MAN</sub> = the sum of the vertical and horizontal distance



# Least Coin Problem – Greedy Search

Problem: For any money amount, calculate the least number of coins to carry in your pocket.

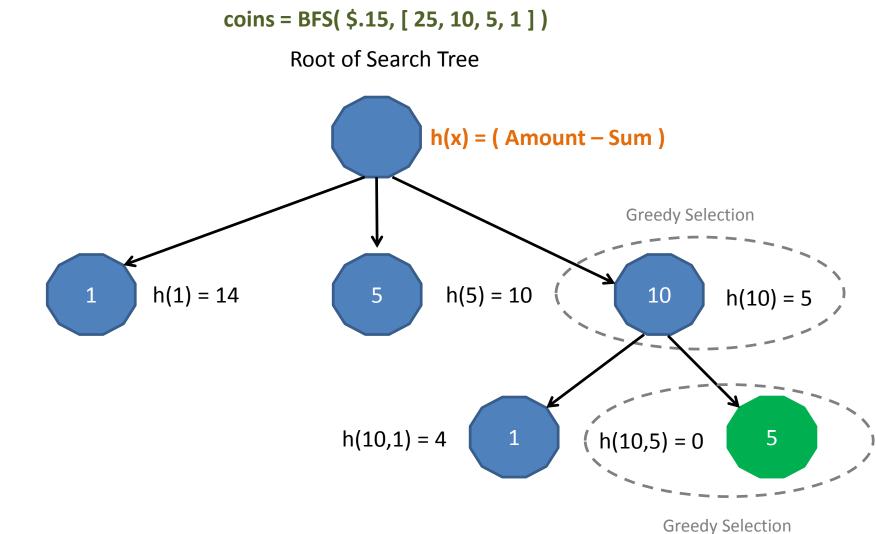


sum = sum + coin

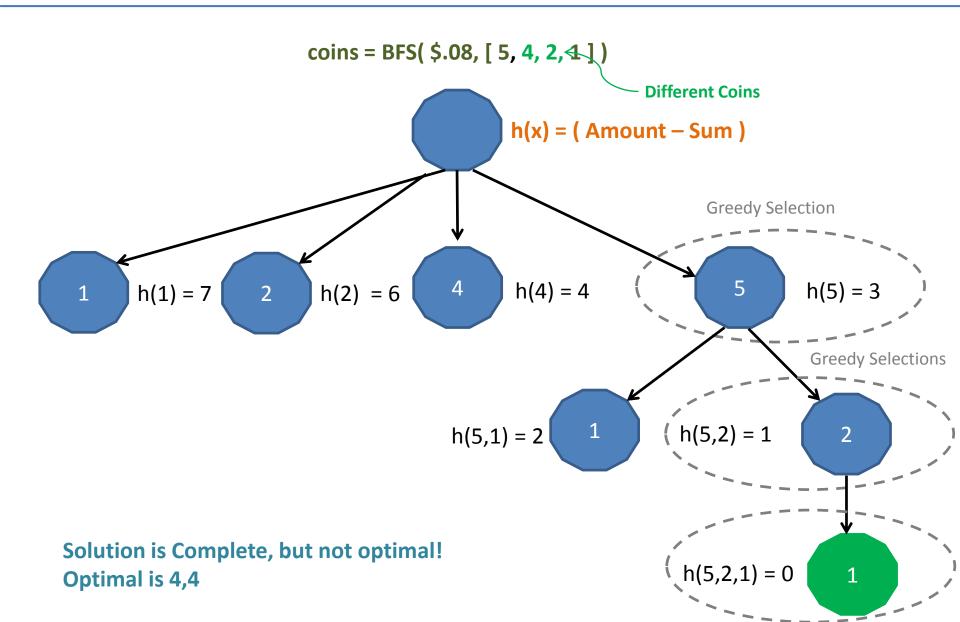
coin = largest coin( amount – sum )

number of coins = number of coins + 1

#### **Greedy Search Example**



### **Greedy Search Example**



# BFS Greedy – Algorithm

function Greedy(root, goal)

Heap = Priority Queue

initialize the <u>frontier</u> to the <u>root node</u> (Heap)



Difference from BFS

initialize the <u>visited</u> (explored) to the empty set (Heap)

while the **frontier** is not empty

remove (delete min) the next <u>node</u> from the <u>frontier</u> add (insert) the (node removed from <u>frontier</u>) <u>node</u> to the <u>visited</u>

if the <u>node</u> matches to <u>goal node</u> return <u>found goal</u>

if <u>child</u> not in <u>frontier</u> or <u>visited</u>

set key of <u>child</u> to h(child)

Difference from BFS

add (insert) child to the frontier

## A\* Algorithm

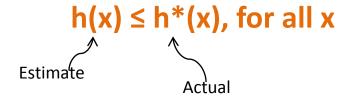
- A search method of selecting the best local choice at each step in hopes of finding an optimal solution.
- <u>Does</u> consider how optimal the current solution is.
- At each step, uses a heuristic h(x) to estimate the distance (cost)
  of each local choice from the goal, plus the accumulated distance
  (cost) g(x) to the current state.

#### Steps:

- 1. Define a heuristic function h(x) to estimate the distance to the goal from any state.
- 2. Define a function g(x) for measuring the actual distance from the start state to the current state.
- 2. From the current state, determine the search space (actions) for one step ahead.
- 3. Select the action from the search space that minimizes g(x) + h(x).

# A\* - Optimality

 The A\* method will always find the <u>optimal</u> solution if the heuristic function h(x), for all x, is always less than or equal the actual distance (cost) to the goal node, h\*(x).

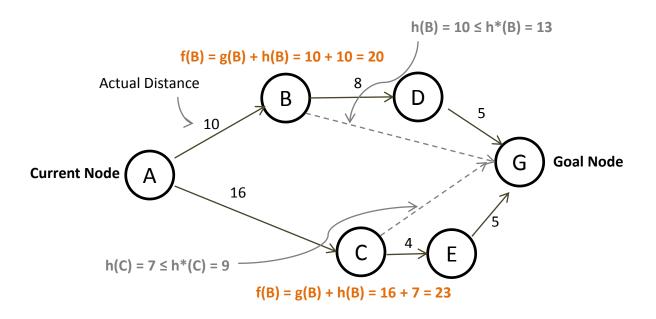


 When the heuristic h(x) ≤ h\*(x), and the evaluation function is f(x) = g(x) + h(x), the method is said to be optimal, since it will always find the optimal solution.

## A\* – SLD Example

 In the previous SLD example, Greedy picked a suboptimal solution (node B), while A\* will pick an optimal solution (node C).

$$f(x) = g(x) + h_{SLD}(x)$$



A\* will select node B, since g(B) + h(B) < g(C) + h(C)

# BFS A\* – Algorithm

```
function AStar(root, goal)
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Heap = Priority Queue

initialize the <u>frontier</u> to the <u>root node</u> (Heap)



Difference from BFS

initialize the <u>visited</u> (explored) to the empty set (Heap)

while the **frontier** is not empty

remove (delete min) the next <u>node</u> from the <u>frontier</u> add (insert) the (node removed from <u>frontier</u>) <u>node</u> to the <u>visited</u>

if the <u>node</u> matches to <u>goal node</u> return <u>found goal</u>

for each <u>child</u> (neighbor) of the <u>node</u>
if <u>child</u> not in <u>frontier</u> or <u>visited</u>

set key of child to g(node) + h(child)



Difference from BFS

add (insert) child to the frontier