

00:43

**QUESTION**

The distance between two stations M and N is  $L$  kilometres. All frames are  $K$  bits long. The propagation time per kilometer is  $t$  seconds. Let  $R$  bits/second be the channel capacity. Assuming the processing delay is negligible, the minimum number of bits for the sequence number field in a frame for maximum utilization, when the sliding window protocol is used, is:

[GATE CS 2007]

- A.  $\log_2 \left( \frac{2LtR+2K}{K} \right)$       C.  $\log_2 \left( \frac{2LtR+K}{K} \right)$   
B.  $\log_2 \left( \frac{2LtR}{K} \right)$       D.  $\log_2 \left( \frac{2LtR+K}{2K} \right)$

NESO ACADEMY

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**SOLUTION**

Let Propagation Delay =  $Lt$  sec

Round Trip Time =  $2 \times$  Propagation Delay

Round Trip Time =  $2 \times Lt$  sec

Round Trip Time =  $2 Lt$  sec

NESO ACADEMY

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**SOLUTION**

$$\text{No. of frames} = \left( \frac{2LtR}{K} \right)$$

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