

Max-SMT

Albert Rubio

Programación con restricciones,
Facultad de Informática, UCM

The “maximum satisfiability problem” is defined as:

Given a conjunction of clauses (disjunction) over boolean literals

$$(p \vee q) \wedge (p \vee \neg q) \wedge (\neg p \vee q) \wedge (\neg p \vee \neg q)$$

Find a solution (model) that

- **maximizes** the number of satisfied clauses

Partial Max-SAT

Given a conjunction of clauses (disjunction) over boolean literals **hard**

$$(p \vee q) \wedge (p \vee \neg q)$$

and a conjunction of clauses (disjunction) over boolean literals **soft**

$$(\neg p \vee q) \wedge (\neg p \vee \neg q)$$

Find a solución that

- satisfies all **hard** clauses
- **maximizes** the number of satisfied **soft** clauses

(Generalized) Partial Max-SAT

Given a conjunction of formulas (constraints) over boolean literals **hard**

$$(p \vee q) \wedge (p \vee \neg q)$$

and a conjunction of formulas (constraints) over boolean literals **soft**

$$(\neg p \vee q) \wedge (\neg p \vee \neg q)$$

Find a solución that

- satisfies all **hard** constraints
- **maximizes** the number of satisfied **soft** constraints

Weighted (partial) Max-SAT

Given a conjunction of formulas (constraints) over boolean literals **hard**

$$(p \vee q) \wedge (p \vee \neg q)$$

and a conjunction of formulas (constraints) over boolean literals **soft** with a **weight**

$$(\neg p \vee q : w(3)) \wedge (\neg p \vee \neg q : w(2))$$

Find a solución that

- satisfies all **hard** constraints
- **maximize** the sum of the weights of satisfied **soft** constraints

Weighted (partial) Max-SAT

Given a conjunction of formulas (constraints) over boolean literals **hard**

$$(p \vee q) \wedge (p \vee \neg q)$$

and a conjunction of formulas (constraints) over boolean literals **soft** with a **weight**

$$(\neg p \vee q : w(3)) \wedge (\neg p \vee \neg q : w(2))$$

Find a solución that

- satisfies all **hard** constraints
- **maximize** the sum of the weights of satisfied **soft** constraints
- **minimize** the sum of the weights of not satisfied **soft** constraints

(Weighted partial) Max-SMT

Given a conjunction of **hard** constraints over some theory T (e.g. LIA)

$$(x \leq y \vee x \leq z) \wedge (x \leq y \vee x > z)$$

and a conjunction of **weighted** constraints over some theory T

$$(x > y \vee x \leq z : w(3)) \wedge (x > y \vee x > z : w(2))$$

Find a solución that

- satisfies all **hard** constraints
- **maximize** the sum of the weights of satisfied **soft** constraints
- **minimize** the sum of the weights of not satisfied **soft** constraints

This is not fully standard, but ...

```
(assert-soft (<= t 8) :weight 2)
```

The weight must be a positive natural number, but is optional.

```
(assert-soft (<= t 8))
```

If omitted, the weight is set to 1.

We can have more than one kind of weights

```
(assert-soft (<= t 8) :weight 2 :id A)
```

```
(assert-soft (<= 0 t) :weight 2 :id B)
```

Finds the best solution lexicographically (in order of appearance of tags):

- **maximize** the sum of the weights of satisfied **A** **soft** constraints
- If same weight for **A**
 maximize the sum of the weights of satisfied **B** **soft** constraints
- ...

Choose the names you like for the tags!!

We should use `add_soft` instead of `add`:

```
add_soft (self, arg, weight="1", id=None)
```

Given no optimizer `s` we write:

```
s = Optimize()  
...  
s.add_soft (t <= 8)  
s.add_soft (t <= 8, 2)  
s.add_soft (t <= 8, 2, "A")
```

We can add an objective function.

“still not fully competitive”, but ...

```
(minimize t)
```

```
(maximize t)
```

where t is an expression of type integer or real.

Must be added before `(check-sat)`

With z3py use:

```
s.minimize(t)
```

```
s.maximize(t)
```

We can have multiple objective functions.

Finds the best solution lexicographically (in order of appearance of objectives).

Other combination options can be considered: pareto, independent, ... (check the solver)

Available tools:

- Z3 <https://github.com/Z3Prover/z3>
- (Opti)MathSat <http://optimathsat.disi.unitn.it>

Example

(Borrowed from Z3 tutorial)

Queremos crear tres máquinas virtuales que requieren 100, 50 y 15 GB de disco respectivamente.

Disponemos de tres servidores con 100, 75 and 200 GB de disco respectivamente.

Además estos tres servidores tienen un coste diario de 10, 5 y 20 euros respectivamente.

Encontrar una solución (usando Max-SMT) que decida en que servidor se crea cada máquina virtual tal que:

- Minimize el número de servidores usados
- Minimize el coste diario

Example

Usad variables booleanas

- $x_{i,j}$ para expresar que la MV i se crea en el servidor j
- y_j para expresar que el servidor j se está usando

```
(declare-fun x_1_1 () Bool)
```

```
(declare-fun x_1_2 () Bool)
```

```
(declare-fun x_1_3 () Bool)
```

```
(declare-fun x_2_1 () Bool)
```

```
(declare-fun x_2_2 () Bool)
```

```
(declare-fun x_2_3 () Bool)
```

```
(declare-fun x_3_1 () Bool)
```

```
(declare-fun x_3_2 () Bool)
```

```
(declare-fun x_3_3 () Bool)
```

```
(declare-fun y_1 () Bool)
```

```
(declare-fun y_2 () Bool)
```

```
(declare-fun y_3 () Bool)
```

Example

Otra codificación:

Usad variables

- x_i Int, para expresar que la MV i se crea en el servidor j
- y_j Bool, para expresar que el servidor j se está usando

```
(declare-fun x_1 () Int)
```

```
(declare-fun x_2 () Int)
```

```
(declare-fun x_3 () Int)
```

```
(declare-fun y_1 () Bool)
```

```
(declare-fun y_2 () Bool)
```

```
(declare-fun y_3 () Bool)
```