



A T M E
College of Engineering

COLLEGE OF ENGINEERING
13th KM Stone, Bannur Road, Mysore-560028

DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING – AI & ML
(ACADEMIC YEAR 2024-25)

LABORATORY MANUAL

SUBJECT: ANALYSIS & DESIGN OF ALGORITHMS

SUB CODE: BCSL404

SEMESTER: IV SCHEME: 2022

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INSTITUTIONAL MISSION AND VISION

Objectives

- To provide quality education and groom top-notch professionals, entrepreneurs and leaders for different fields of engineering, technology and management.
- To open a Training-R&D-Design-Consultancy cell in each department, gradually introduce doctoral and postdoctoral programs, encourage basic & applied research in areas of social relevance, and develop the institute as a center of excellence.
- To develop academic, professional and financial alliances with the industry as well as the academia at national and transnational levels.
- To cultivate strong community relationships and involve the students and the staff in Local community service.
- To constantly enhance the value of the educational inputs with the participation of students, faculty, parents and industry.

Vision

- Development of academically excellent, culturally vibrant, socially responsible and globally competent human resources.

Mission

- To keep pace with advancements in knowledge and make the students competitive and capable at the global level.
- To create an environment for the students to acquire the right physical, intellectual, emotional and moral foundations and shine as torchbearers of tomorrow's society.
- To strive to attain ever-higher benchmarks of educational excellence.

Department of Computer Science & Engineering

Vision of the Department

- To develop highly talented individuals in Computer Science and Engineering to deal with realworld challenges in industry, education, research and society.

Mission of the Department

- To inculcate professional behavior, strong ethical values, innovative research capabilities and leadership abilities in the young minds & to provide a teaching environment that emphasizes depth, originality and critical thinking.
- Motivate students to put their thoughts and ideas adoptable by industry or to pursue higher studies leading to research.

Program outcomes (POs)

Engineering Graduates will be able to:

PO1. Engineering knowledge: Apply the knowledge of mathematics, science, engineering fundamentals, and an engineering specialization to the solution of complex engineering problems.

PO2. Problem analysis: Identify, formulate, review research literature, and analyze complex engineering problems reaching substantiated conclusions using first principles of mathematics, natural sciences, and engineering sciences.

PO3. Design/development of solutions: Design solutions for complex engineering problems and design system components or processes that meet the specified needs with appropriate consideration for the public health and safety, and the cultural, societal, and environmental considerations.

PO4. Conduct investigation of complex problems: Use research-based knowledge and research methods including design of experiments, analysis and interpretation of data, and synthesis of the information to provide valid conclusions.

PO5. Modern tool usage: Create, select, and apply appropriate techniques, resources, and modern engineering and IT tools including prediction and modeling to complex engineering activities with an understanding of the limitations.

PO6. The engineer and society: Apply reasoning informed by the contextual knowledge to assess societal, health, safety, legal and cultural issues and the consequent responsibilities relevant to the professional engineering practice.

PO7. Environment and sustainability: Understand the impact of the professional engineering solutions in societal and environmental contexts, and demonstrate the knowledge of, and need for sustainable development.

PO8. Ethics: Apply ethical principles and commit to professional ethics and responsibilities and norms of the engineering practice.

PO9. Individual and teamwork: Function effectively as an individual, and as a member or leader in diverse teams, and in multidisciplinary settings.

PO10. Communication: Communicate effectively on complex engineering activities with the engineering community and with society at large, such as, being able to comprehend and write effective reports and design documentation, make effective presentations, and give and receive clear instructions.

PO11. Project management and finance: Demonstrate knowledge and understanding of the engineering and management principles and apply these to one's own work, as a member and leader in a team, to manage projects and in multidisciplinary environments.

PO12. Life-long learning: Recognize the need for, and have the preparation and ability to engage in independent and life-long learning in the broadest context of technological change.

Program Specific Outcomes

PSO1: Ability to apply skills in the field of algorithms, database design, web design, cloud computing and data analytics.

PSO2: Apply knowledge in the field of computer networks for building network and internet based applications

Program Educational Objectives (PEOs):

1. Empower students with a strong basis in the mathematical, scientific and engineering fundamentals to solve computational problems and to prepare them for employment, higher learning and R&D.
2. Gain technical knowledge, skills and awareness of current technologies of computer science engineering and to develop an ability to design and provide novel engineering solutions for software/hardware problems through entrepreneurial skills.
3. Exposure to emerging technologies and work in teams on interdisciplinary projects with effective communication skills and leadership qualities.
4. Ability to function ethically and responsibly in a rapidly changing environment by applying innovative ideas in the latest technology, to become effective professionals in Computer Science to bear a life-long career in related areas.

Analysis&DesignofAlgorithmsLab		Semester	4
CourseCode	BCSL40 4	CIEMarks	50
TeachingHours/Week(L:T:P:S)	0:0:2:0	SEEMarks	50
Credits	01	ExamHours	2
Examinationtype(SEE)	Practical		
Sl.No	Experiments		
1	DesignandimplementC/C++ProgramtofindMinimumCostSpanningTreeofagivenconnected undirectedgraphusingKruskal'salgorithm.		
2	DesignandimplementC/C++ProgramtofindMinimumCostSpanningTreeofagivenconnected undirectedgraphusingPrim'salgorithm.		
3	a. DesignandimplementC/C++ProgramtosolveAll-PairsShortestPathsproblemusingFloyd's algorithm. b. DesignandimplementC/C++ProgramtofindthetransitiveclosureusingWarshal's algorithm.		
4	DesignandimplementC/C++Programtofindshortest pathsfromagivenvertexinaweighted connectedgraphtootherverticesusingDijkstra'salgorithm.		
5	Designand implementC/C++ProgramtoobtaintheTopologicalorderingofverticesinagiven digraph.		
6	DesignandimplementC/C++Programtosolve0/1KnapsackproblemusingDynamic Programmingmethod.		
7	DesignandimplementC/C++ProgramtosolvediscreteKnapsackandcontinuousKnapsack problemsusing greedyapproximationmethod.		
8	DesignandimplementC/C++ProgramtofindasubsetofagivensetS={s ₁ ,s ₂ , ,s _n }ofn positiveintegerswhosesumisequaltoagivenpositiveintegerd.		
9	Design and implement C/C++ Program to sort a given set of n integer elements using Selection Sortmethod andcomputeits timecomplexity. Runthe programforvaried values of n>5000 and record the time taken to sort. Plot a graph of the time taken versus n. The elementscan be read fromafileorcanbegeneratedusingtherandomnumbergenerator.		
10	Design and implement C/C++ Program to sort a given set of n integer elements using Quick Sort methodandcomputeitstimecomplexity.Runtheprogramforvariedvaluesofn>5000andrecord the time taken to sort. Plota graph of the time taken versus n. The elementscanbe read fromafileorcanbegeneratedusingtherandomnumbergenerator.		
11	Design and implement C/C++ Program to sort a given set of n integer elements using Merge Sort method and compute its time complexity. Run the program for varied values of n> 5000, and record the time taken to sort. Plota graph of the time taken versus n. The elementscan be read fromafileorcanbegeneratedusingtherandomnumbergenerator.		
12	DesignandimplementC/C++ProgramforNQueen'sproblemusingBacktracking.		

Course outcomes (Course Skill Set):

At the end of the course the student will be able to:

1. Develop programs to solve computational problems using suitable algorithm design strategy.
2. Compare algorithm design strategies by developing equivalent programs and observing running times for analysis (Empirical).
3. Make use of suitable integrated development tools to develop programs
4. Choose appropriate algorithm design techniques to develop solution to the computational and complex problems.
5. Demonstrate and present the development of program, its execution and running time(s) and record the results/inferences.

Assessment Details (both CIE and SEE)

The weightage of Continuous Internal Evaluation (CIE) is 50% and for Semester End Exam (SEE) is 50%. The minimum passing mark for the CIE is 40% of the maximum marks (20 marks out of 50) and for the SEE minimum passing mark is 35% of the maximum marks (18 out of 50 marks). A student shall be deemed to have satisfied the academic requirements and earned the credits allotted to each subject/course if the student secures a minimum of 40% (40 marks out of 100) in the sum total of the CIE (Continuous Internal Evaluation) and SEE (Semester End Examination) taken together

Continuous Internal Evaluation (CIE):

CIE marks for the practical course are **50 Marks**.

The split-up of CIE marks for record/journal and test are in the ratio **60:40**.

- Each experiment is to be evaluated for conduction with an observation sheet and record write-up. Rubrics for the evaluation of the journal/write-up for hardware/software experiments are designed by the faculty who is handling the laboratory session and are made known to students at the beginning of the practical session.
- Records should contain all the specified experiments in the syllabus and each experiment write-up will be evaluated for 10 marks.
- Total marks scored by the students are scaled down to **30 marks** (60% of maximum marks).
- Weightage to be given for neatness and submission of record/write-up on time.
- Departments shall conduct a test of 100 marks after the completion of all the experiments listed in the syllabus.
- In a test, test write-up, conduction of experiment, acceptable result, and procedural knowledge will carry a weightage of 60% and the rest 40% for viva-voce.
- The suitable rubrics can be designed to evaluate each student's performance and learning ability.
- The marks scored shall be scaled down to **20 marks** (40% of the maximum marks).

The sum of scaled-down marks scored in the report write-up/journal and marks of a test is the total CIE marks scored by the student.

Semester End Evaluation (SEE):

- SEE marks for the practical course are 50 Marks.

- SEE shall be conducted jointly by the two examiners of the same institute, examiners are appointed by the Head of the Institute.
- The examination schedule and names of examiners are informed to the university before the conduction of the examination. These practical examinations are to be conducted between the schedule mentioned in the academic calendar of the University.
- All laboratory experiments are to be included for practical examination.
- (Rubrics) Breakup of marks and the instructions printed on the cover page of the answerscript to be strictly adhered to by the examiners. **OR** based on the course requirement evaluation rubrics shall be decided jointly by examiners.
- Students can pick one question (experiment) from the questions slot prepared by the examiners jointly.
- Evaluation of test write-up/ conduction procedure and result/viva will be conducted jointly by examiners.

General rubrics suggested for SEE are mentioned here, write up-20%, Conduction procedure and result in -60%, Viva-voce 20% of maximum marks. SEE for practical shall be evaluated for 100 marks and scored marks shall be scaled down to 50 marks (however, based on course type, rubrics shall be decided by the examiners)

Change of experiment is allowed only once and 15% of Marks allotted to the procedure part to be made zero.

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INTRODUCTION

Needforstudyingalgorithms Theoretical importance

- The study of algorithms is the concept of computer science.
- It is a standard set of important algorithms, they further our analytical skills & help us in developing new algorithms for required applications

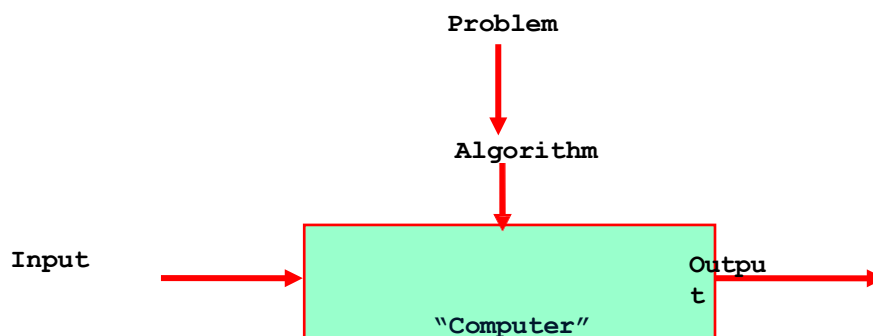
Algorithm

An algorithm is a sequence of unambiguous instructions for solving a problem, i.e., for obtaining a required output for any legitimate input in a finite amount of time.

In addition, all algorithms must satisfy the following criteria:

1. **Input:** Each algorithm should have zero or more inputs. The range of input for which algorithms work should be specified carefully.
2. **Output:** The algorithm should produce correct results. At least one quantity has to be produced.
3. **Definiteness:** Each instruction should be clear and unambiguous.
4. **Effectiveness:** Every instruction should be simple and should transform the given input to the desired output. so that it can be carried out, in
5. **Finiteness:** If we trace out the instruction of an algorithm, then for all cases, the algorithm must terminate after a finite number of steps.

Notion of algorithm



Fundamentals of Algorithmic problem solving

- Understanding the problem
- Ascertain the capabilities of the computational device
- Exact/approximate solution.
- Decide on the appropriate data structure
- Algorithm design techniques
- Methods of specifying an algorithm
- Proving an algorithm's correctness
- Analyzing an algorithm

Important Problem Types of different categories**• Sorting**

It refers to the problem of re-arranging the items of a given list in ascending or descending order. The various algorithms that can be used for sorting are bubble sort, selection sort, insertion sort, quick sort, merge sort, heap sort etc.

• Searching

* This problem deals with finding a value, called a search key, in a given set.

* The various algorithms that can be used for searching are binary search, linear search, hashing, interpolation search etc.

• String processing

* This problem deals with manipulation of characters or strings, string matching, search and replace, deleting a string in a text etc.

* String processing algorithms such as string matching algorithms are used in the design of assemblers and compilers.

• Graph problems

* The graph is a collection of vertices and edges.

* Ex: graph traversal problems, topological sorting etc

• Combinatorial problems

These problems are used to find combinatorial objects such as permutation and combinations.

Ex: travelling salesman problem, graph coloring problem etc

• Geometric problems

This problem deals with geometric objects such as points, curves, lines, polygon etc.

Ex: closest-pair problem, convex hull problem

• Numerical problems

These problems involving mathematical manipulation solving equations, computing differentiations and integrations, evaluating various types of functions etc.

Ex: Gauss elimination method, Newton-Raphson method

Fundamentals of data Structures

Since most of the algorithms operate on the data, particular ways of arranging the data play a critical role in the design & analysis of algorithms.

A data structure can be defined as a particular way of arrangement of data.

The commonly used data structures are:

1. Linear data structures
2. Graphs
3. Trees.
4. Sets and dictionaries

1. Linear data structures

The most common linear data structures are arrays and linked lists.

Arrays: Is a collection of homogeneous items. An item's place within the collection is called an index.

Linked list: finite sequence of data items i.e. it is a collection of data items in a certain order.

2. Graphs

A data structure that consists of a set of nodes and a set of edges that relate the nodes to each other is called a graph.

Undirected graph: A graph in which the edges have no direction

➤ **Directed graph (Digraph):** A graph in which each edge is directed from one vertex to another (or the same) vertex.

3. Tree

A tree is a connected acyclic graph that has no cycle.

4. Sets and dictionaries

* A set is defined as an unordered collection of distinct items called an element of the set

* Dictionary is a data structure that implements searching, adding of objects

Analysis of algorithms

Analysis of algorithms means to investigate an algorithm's efficiency with respect to resources:

Running time (time efficiency)

It indicates how fast an algorithm in question runs

• Memory space (space efficiency)

It deals with the extra space the algorithm requires

Theoretical analysis of time efficiency

Algorithm efficiency depends on the **input size**. And for some algorithm efficiency depends on **type of input**.

We have best, worst & average case efficiencies

Worst-case efficiency:

Efficiency (number of times the basic operation will be executed) **for the worst case**

input of size n . *i.e.* The algorithm runs the longest among all possible inputs of size n .

Best-case efficiency:

Efficiency (number of times the basic operation will be executed) **for the best case input of size n .** *i.e.* The algorithm runs the fastest among all possible inputs of size n

Average-case efficiency:

Average time taken (number of times the basic operation will be executed) **to solve all the possible instances (random) of the input.**

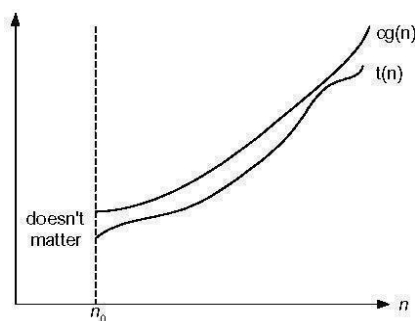
Asymptotic Notations

- Asymptotic notation is a way of comparing functions that ignores constant factors and small input sizes.
- Three notations used to compare orders of growth of an algorithm's basic operation are:

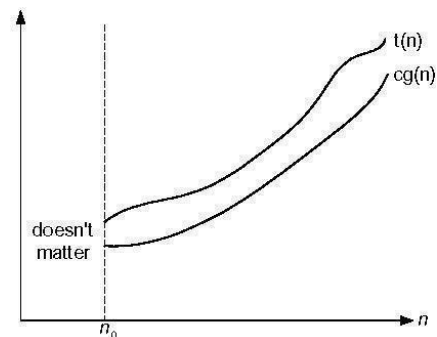
O, Ω , θ notations.

Big-oh notation:

A function $t(n)$ is said to be in $O(g(n))$, denoted $t(n) \in O(g(n))$, if $t(n)$ is bounded above by some constant multiple of $g(n)$ for all large n . *i.e.*, if there exist some positive constant c and some nonnegative integer n_0 such that $t(n) \leq cg(n)$ for all $n \geq n_0$



Big-oh notation $t(n) \in O(g(n))$



Big-omega notation $t(n) \in \Omega(g(n))$

Big-omega notation:

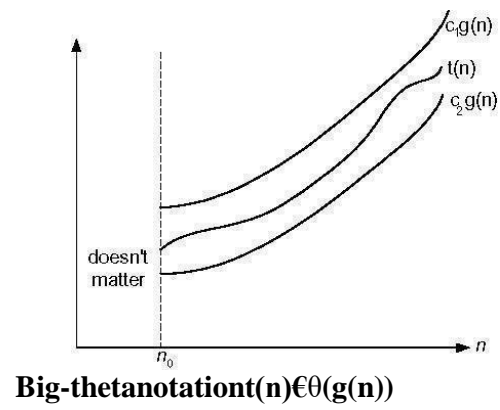
A function $t(n)$ is said to be in $\Omega(g(n))$, denoted $t(n) \in \Omega(g(n))$, if $t(n)$ is bounded below by some constant multiple of $g(n)$ for all large n . *i.e.*, if there exist some positive constant c and some nonnegative integer n_0 such that

$$t(n) \geq cg(n) \text{ for all } n \geq n_0$$

theta notation:

A function $t(n)$ is said to be in $\theta(g(n))$, denoted $t(n) \in \theta(g(n))$, if $t(n)$ is bounded both above and below by some positive constant multiple of $g(n)$ for all large n . *i.e.*, if there exist some positive constant c

$$c_2 g(n) \leq t(n) \leq c_1 g(n) \text{ for all } n \geq n_0$$



Basic Asymptotic Efficiency classes:

1	constant
$\log n$	logarithmic
N	linear
$n \log n$	$n \log n$
n^2	quadratic
n^3	cubic
2^n	exponential
$n!$	factorial

1. Design and implement C/C++ Program to find Minimum Cost Spanning Tree of a given connected undirected graph using Kruskal's algorithm.

```
#include<stdio.h>int

e=1,min_cost=0;

voidmain()
{
    intn,i,j,min,a,u,b,v,cost[20][20],parent[20];
    printf("Enter the no. of vertices:");
    scanf("%d",&n);
    printf("\nEnterthecostmatrix:\n");
    for(i=1;i<=n;i++)for(j=1;j<=n;j++)
        scanf("%d",&cost[i][j]);
    for(i=1;i<=n;i++)
        parent[i]=0;
    printf("\nTheedgesofspanningtreeare\n"); while(ne<n)
    {
        min=999;
        for(i=1;i<=n;i++)
        {
            for(j=1;j<=n;j++)
            {
                if(cost[i][j]<min)
                {
                    min=cost[i][j];
                    a=u=i;
                    b=v=j;
                }
            }
            while(parent[u])
                u=parent[u];
            while(parent[v])
                v=parent[v];
            if(u!=v)
            {
                printf("Edge%d\t(%d->%d)=%d\n",ne++,a,b,min);
                min_cost=min_cost+min;
                parent[v]=u;
            }
            cost[a][b]=cost[a][b]=999;
        }
        printf("\nMinimumcost=%d\n",min_cost);
    }
}
```


Output:

Enter the no. of vertices: 6

Enter the cost matrix:

```
9993      999  999  65
3  999  1  999  9994
9991      999  6  9994
999999  6  999  8  5
6  9999998      999  2
5  4  4  5  2  999
```

The edges of spanning tree are

```
Edge 1    (2->3)=1
Edge2     (5->6)=2
Edge3     (1->2)=3
Edge4     (2->6)=4
Edge5     (4->6)=5
```

Minimum cost=1

2. Design and implement C/C++ Program to find Minimum Cost Spanning Tree of a given connected undirected graph using Prim's algorithm.

```
#include<stdio.h>
inta,b,u,v,n,i,j,ne=1;
intvisited[10]={0},min,mincost=0,cost[10][10];
void main()
{
printf("\nEnter the number of nodes:");
scanf("%d",&n);
printf("\nEnter the adjacency matrix:\n");
for(i=1;i<=n;i++)
for(j=1;j<=n;j++)
{
scanf("%d",&cost[i][j]);
if(cost[i][j]==0)
cost[i][j]=999;
}
visited[1]=1;
printf("\n");
while(ne<n)
{
for(i=1,min=999;i<=n;i++)
for(j=1;j<=n;j++)
if(cost[i][j]<min)
if(visited[i]!=0)
{
min=cost[i][j];
a=u=i;
b=v=j;
}
if(visited[u]==0|| visited[v]==0)
{
printf("\nEdge %d:(%d%d)cost:%d",ne++,a,b,min);
mincost+=min;
visited[b]=1;
}
cost[a][b]=cost[b][a]=999;
}
printf("\nMinimuncost=%d",mincost);
}
```

Output

Enterthenumberofnodes:4

Entertheadjacencymatrix:

999 1 5 2

1 999999999

59999993

29993 999

Edge1:(12)cost:1

Edge2:(14)cost:2

Edge3:(43)cost:3

Minimun cost=6

3a)DesignandimplementC/C++ProgramtosolveAll-PairsShortestPathsproblem using Floyd's algorithm.

```
#include
<stdio.h>
#include<limits.h>
#define V 4
voidfloydWarshall(intgraph[V][V])
{
intdist[V][V];
for(inti=0;i<V;i++)for(int j=0; j<V;j++) dist[i][j]
= graph[i][j];
for(int k=0;k<V;k++) for (inti=0;i<V; i++)
for(intj=0;j<V; j++)
if(dist[i][k]!=INT_MAX&&dist[k][j]!=INT_MAX&& dist[i][k]+dist[k][j]
<dist[i][j])dist[i][j]=dist[i][k]+ dist[k][j]

("Shortestdistancesbetweeneverypairofvertices:\n"); for
(int i = 0; i < V; i++)

{
for(intj=0; j<V; j++)
{
if(dist[i][j]==INT_MAX)
printf("INF\t");
else
printf("%d\t",dist[i][j]);
}
printf("\n");
}
}

intmain(){
int graph[V][V]={ {0,INT_MAX,3,INT_MAX},
{2,0,INT_MAX,INT_MAX},
{INT_MAX,7,0,1},{6,INT_MAX,INT_MAX,0}};
floydWarshall(graph);
return 0;
}
```

Output:

Shortestdistancesbetweeneverypairofvertices:

0	10	3	4
2	0	5	6
7	7	0	1
6	16	9	0

3BDesignandimplementC/C++ProgramtofindthetransitiveclosureusingWarshal's algorithm.

```
#include<stdio.h>
intn,a[10][10],p[10][10];
voidpath()
{
int i,j,k;
for(i=0;i<n;i++)
for(j=0;j<n;j++)
p[i][j]=a[i][j];
for(k=0;k<n;k+)
for(i=0;i<n;i++)
for(j=0;j<n;j++)
if(p[i][k]==1&& p[k][j]==1)
p[i][j]=1;
}
voidmain()
{
inti,j;
printf("Enterthenumberofnodes:");
scanf("%d",&n);printf("\nEntertheadjacencymatrix:\n");for(i=0;i<n;i++)
for(j=0;j<n;j++)
scanf("%d",&a[i][j]);path();
printf("\nThepathmatrixisshownbelow\n");
for(i=0;i<n;i++)
{
for(j=0;j<n;j++)
printf("%d",p[i][j]);
printf("\n");
}
}
```

Output:

Enterthefnumberofnodes:4

Enter the adjacency matrix:

0100

0010

0001

1000

Thepathmatrixisshownbelow1111

11 11

11 11

11 11

4. Design and implement C/C++ Program to find shortest paths from a given vertex in a weighted connected graph to other vertices using Dijkstra's algorithm.

```
#include<stdio.h>
void dij(int,int[20][20],int[20],int[20],int); void
main()
{
int i,j,n,visited[20],source,cost[20][20],d[20]; printf("Enter
no. of vertices: ");
scanf("%d",&n);
printf("Enter the cost adjacency matrix\n");
for (i= 1; i <= n; i++)
{
for(j=1;j<=n;j++)
{
scanf("%d",&cost[i][j]);
}
}
printf("\nEnter the source node:");
scanf("%d", &source);
dij(source,cost,visited,d,n); for
(i = 1; i <= n; i++)
{
if(i!=source)
printf("\nShortest path from %d to %d is %d",source,i,d[i]);
}
}
void dij(int source,int cost[20][20],int visited[20],int d[20],int n)
{
int i,j,min,u,w;
for(i=1; i<=n;i++)
{
visited[i]=0;
d[i]=cost[source][i];
}
visited[source]=1;
d[source]=0;
for(j=2;j<=n;j++)
{
min =999;
for(i=1;i<=n;i++)
{
if(!visited[i]) { if(d[i]<min)
{
min=d[i];u = i;
}
}
```



```
}  
}  
visited[u]= 1;  
for(w=1;w<=n;w++)  
{  
if(cost[u][w]!=999&&visited[w] ==0)  
{  
    if(d[w] >cost[u][w]+d[u])  
d[w] =cost[u][w]+d[u];  
}  
}  
}  
}
```

Output:

Enterno.ofvertices:6
Enterthecostadjacencymatrix
9993 999 9996 5
3 999 1 9999994
9991 999 6 999 4
9999996 999 8 5
6 999999 8 9992
5 4 4 5 2 999

Enter the source node: 1

Shortestpathfrom1to2is3
Shortestpathfrom1to3is4
Shortestpathfrom1to4is10
Shortestpathfrom1to5is6
Shortestpathfrom1to6is5

5. Design and implement C/C++ Program to obtain the Topological ordering of vertices in a given digraph.

```
#include<stdio.h>
void findIndegree(int [10][10],int [10],int);
void topological(int,int [10][10]);
void main()
{
    int a[10][10],i,j,n;
    printf("Enter the number of nodes:");scanf("%d",&n);
    printf("\nEnter the adjacency matrix\n");
    for(i=1;i<=n;i++)
    for(j=1;j<=n;j++)
    scanf("%d",&a[i][j]);
    printf("\nThe adjacency matrix is:\n");for(i=1;i<=n;i++)
    {
        for(j=1;j<=n;j++)
        {
            printf("%d\t",a[i][j]);
        }
        printf("\n");
    }
    topological(n,a);
}
void findIndegree(int a[10][10],int indegree[10],int n)
{
    int i,j,sum;for(j=1;j<=n;j++)
    {
        sum=0;for(i=1;i<=n;i++)
        {
            sum=sum+a[i][j];
        }
        indegree[j]=sum;
    }
}
void topological(int n,int a[10][10])
{
    int k,top,t[100],i,stack[20],u,v,indegree[20];k=1;
    top=-1; findIndegree(a,indegree,n);
    for(i=1;i<=n;i++)
    {
        if(indegree[i]==0)
        {
            stack[++top]=i;
        }
    }
    while(top!=-1)
```

```
{
u=stack[top--];t[k++]=u;
for(v=1;v<=n;v++)
{
if(a[u][v]==1)
{
indegree[v]--;
if(indegree[v]==0)
{
stack[++top]=v;
}
}
}
}
printf("\nTopological sequence is\n");for(i=1;i<=n;i++)
printf("%d\t",t[i]);
}
```

Output:

Enterthenumberofnodes:5 Enter

the adjacency matrix

```
0 0 10 0
0 0 10 0
0 0 01 1
0 0 00 1
0 0 0 0 0
```

Theadjacencymatirxis:

```
0 0 1 0 0
0 0 1 0 0
0 0 0 1 1
0 0 0 0 1
0 0 0 0 0
```

Topologicalsequenceis

```
2 1 3 4 5
```

6. Design and implement C/C++ Program to solve 0/1 Knapsack problem using Dynamic Programming method.

```
#include<stdio.h>
#define MAX 50
int p[MAX],w[MAX],n;
int knapsack(int,int);
int max(int,int);
void main()
{
    int m,i,optsoln;

    printf("Enter no. of objects:");scanf("%d",&n);printf("\nEnter the weights:\n");for(i=1;i<=n;i++)
    scanf("%d",&w[i]); printf("\nEnter the profits:\n"); for(i=1;i<=n;i++) scanf("%d",&p[i]);
    printf("\nEnter the knapsack capacity:");
    scanf("%d",&m); optsoln=knapsack(1,m);
    printf("\nThe optimal solution is:%d",optsoln);
}

int knapsack(int i,int m)
{
    if(i==n)
        return(w[n]>m)?0:p[n];
    if(w[i]>m)
        return knapsack(i+1,m);
    return max(knapsack(i+1,m),knapsack(i+1,m-w[i])+p[i]);
}

int max(int a,int b)
{
    if(a>b) return a; else return b;
}
```

Output:

Enter no. of objects: 3

Enter the weights:

100 10 14

Enter the profits:

20 18 15

Enter the knapsack capacity: 116

The optimal solution is: 38

7. Design and implement C/C++ Program to solve discrete Knapsack and continuous Knapsack problems using greedy approximation method.

```
#include<iostream>
#include
<vector>#include<algorithm>
using namespace std;
//Structure to represent an item
struct Item
{
    int weight; int value;
};
//Function to solve discrete knapsack using greedy approach
int discreteKnapsack(vector<Item>& items, int capacity)
{
    //Sort items based on their value per unit weight
    sort(items.begin(), items.end(), [](const Item& a, const Item& b)
    {
        return (double)a.value/a.weight > (double)b.value/b.weight;
    });
    int totalValue=0;
    int currentWeight=0;
    //Fill the knapsack with items
    for (const Item& item : items) {
        if (currentWeight+item.weight<=capacity)
        {
            currentWeight+=item.weight;
            totalValue += item.value;
        }
    }
    return totalValue;
}
// Function to solve continuous knapsack using greedy approach
double continuousKnapsack(vector<Item>& items, int capacity) {
    // Sort items based on their value per unit weight
    sort(items.begin(), items.end(), [](const Item& a, const Item& b)
    {
        return (double)a.value/a.weight > (double)b.value/b.weight;
    });
    double totalValue=0.0;
    int currentWeight = 0;
    //Fill the knapsack with items fractionally
    for (const Item& item: items)
    {
        if (currentWeight+item.weight<=capacity)
        {
            currentWeight+=item.weight;
            totalValue += item.value;
        }
    }
}
```

```
    }
    else
    {
        intremainingCapacity=capacity-currentWeight;
        totalValue+=(double)item.value/item.weight*remainingCapacity;
        break;

    }
}
returntotalValue;
}
intmain()
{vector<Item>items;intn,capacity;
//Inputnumber ofitemsandcapacityofknapsackcout
<<"Enterthenumberofitems:";cin >>n;
cout<<"Enterthecapacityofknapsack:";cin>>capacity;
//Inputtheweightandvalueofeach item
cout<<"Entertheweightandvalueofeachitem:"<<endl;for(int i =0;i <n; i++)
{
    Itemitem;
    cout<<"Item"<<i+1<<":";cin>>
    item.weight >> item.value;
    items.push_back(item);
}
//Solvediscreteknapsackproblem
intdiscreteResult=discreteKnapsack(items,capacity);
cout<<"Maximumvalue fordiscreteknapsack:"<<discreteResult<<endl;
//Solvecontinuousknapsackproblem
doublecontinuousResult=continuousKnapsack(items,capacity);
cout<<"Maximumvalueforcontinuousknapsack:"<<continuousResult<<endl;return0;
}
```


Output:

Enterthe number of items:4

Enterthecapacityofknapsack:10

Entertheweightandvalueofeachitem:

Item1:2 10

Item2:3 5

Item3:5 15

Item 4:7 7

Maximum value for discrete knapsack :30

Maximumvalueforcontinuousknapsack:30

8. Design and implement C/C++ Program to find a subset of a given set $S = \{s_1, s_2, \dots, s_n\}$ of n positive integers whose sum is equal to a given positive integer d .

```
#include<stdio.h>void
subset(int,int,int);
int x[10],w[10],d,count=0;
void main()
{
    int i,n,sum=0;

    printf("Enter the no. of elements:");
    scanf("%d",&n);
    printf("\nEnter the elements in ascending order:\n");
    for(i=0;i<n;i++)
        scanf("%d",&w[i]);
    printf("\nEnter the sum:");
    scanf("%d",&d);
    for(i=0;i<n;i++)
        sum=sum+w[i];
    if(sum<d)
    {
        printf("No solution\n");
        return;
    }
    subset(0,0,sum);
    if(count==0)
    {
        printf("No solution\n");
        return;
    }
}

void subset(int cs,int k,int r)
{
    int i; x[k]=1;
    if(cs+w[k]==d)
    {
        printf("\n\nSubset %d\n",++count);
        for(i=0;i<=k;i++)
            if(x[i]==1)
                printf("%d\t",w[i]);
    }
    elseif(cs+w[k]+w[k+1]<=d)
        subset(cs+w[k],k+1,r-w[k]);
    if(cs+r-w[k]>=d&&cs+w[k]<=d)
    {
        x[k]=0;
```

```
subset(cs,k+1,r-w[k]);  
}  
}
```

Output:

Enterthno.ofelements:5

Entertheelementsinascendingorder: 1

2

5

6

8

Enterthesum:9

Subset1

1 2 6

Subset2

1 8

9. Design and implement C/C++ Program to sort a given set of n integer elements using Selection Sort method and compute its time complexity. Run the program for varied values of n>5000 and record the time taken to sort. Plot a graph of the time taken versus n. The elements can be read from a file or can be generated using the random number generator.

```
#include <stdio.h>
#include<stdlib.h>
#include <time.h>
//Function to perform selection sort
void selectionSort(int arr[], int n)
{
    int i, j, minIndex, temp;
    for (i = 0; i < n - 1; i++) {
        minIndex = i;
        for (j = i + 1; j < n; j++) {
            if (arr[j] < arr[minIndex])
                minIndex = j;
        }
        //Swap the found minimum element with the first element
        temp = arr[minIndex];
        arr[minIndex] = arr[i];
        arr[i] = temp;
    }
}
//Function to generate random numbers between 0 and 999
int generateRandomNumber() {
    return rand() % 1000;
}
int main() {
    //Set n value
    int n = 6000;
    //Allocate memory for the array
    int* arr = (int*) malloc(n * sizeof(int));
    //Generate random elements for the array
    srand(time(NULL));
    printf("Random numbers for n = %d:\n", n);
    for (int i = 0; i < n; i++) {
        arr[i] = generateRandomNumber();
        printf("%d ", arr[i]);
    }
    printf("\n");
    //Record the start time
    clock_t start = clock();
    //Perform selection sort
```

```

selectionSort(arr,n);
//Recordtheendtime
clock_tend=clock();
//Calculatethetimetakenfor sorting
doubletime_taken=((double)(end -start))/CLOCKS_PER_SEC;
//Outputthetime takentosortforthevalueofn
printf("\nTimetaketosort forn=%d:%lfseconds\n\n",n,time_taken);
//Displaysortednumbers
printf("Sortednumbersforn=%d:\n",n); for
(int i = 0; i < n; i++)
{
printf("%d",arr[i]);
}
printf("\n\n");
//Freethedynamicallyallocatedmemoryfree(arr);return 0;
}

```

Output:

Randomnumbersforn=6000:

243112599677912413721547640822...(morenumbers).....394.....

Timetaketosprtforn=6000:1.058000seconds

Sorted numbers for n=6000:

001223334 5.....(morenumbers).....995996 996 997998 999999

- 10. Design and implement C/C++ Program to sort a given set of n integer elements using Quick Sort method and compute its time complexity. Run the program for varied values of n > 5000 and record the time taken to sort. Plot a graph of the time taken versus n. The elements can be read from a file or can be generated using the random number generator.**

```
#include<stdio.h>
#include<stdlib.h>
#include<time.h>
//Function to swap two elements void
swap(int* a, int* b)
{
    int temp = *a;
    *a = *b;
    *b = temp;
}
//Function to partition the array and return the pivot index int
partition(int arr[], int low, int high) {
    int pivot = arr[high]; int i
    = (low - 1);
    for (int j = low; j <= high - 1; j++)
    {
        if (arr[j] < pivot) { i++;
        swap(&arr[i], &arr[j]);
        }
    }
    swap(&arr[i + 1], &arr[high]);
    return (i + 1);
}

//Function to perform Quick Sort
void quickSort(int arr[], int low, int high)
{
    if (low < high) {
        int pi = partition(arr, low, high);

        quickSort(arr, low, pi - 1);

        quickSort(arr, pi + 1, high);
    }
}
//Function to generate random numbers between 0 and 999 int generateRandomNumber()
{
    return rand() % 1000;
}
int main() {
    // Set n value int n = 6000;
    // Allocate memory for the array
    int* arr = (int*) malloc(n * sizeof(int));
```

```
//Generaterandomelementsforthearray srand(time(NULL));
printf("Randomnumbersfor n=%d:\n",n);for(int i =
0; i < n; i++)
{
arr[i]=generateRandomNumber();printf("%d",arr[i]);
}
printf("\n");
//Recordthestarttimeclock_t
start = clock();
// Perform quick sort
quickSort(arr,0,n-1);
//Recordtheendtimeclock_t end
= clock();
//Calculatethetimetakenfor
//sorting
double time_taken =
((double)(end-start))/
CLOCKS_PER_SEC;
//Outputthetimetakentosort forthecurrentvalueofn
```

```

printf("\nTimetaken to sort for n=%d: %f seconds\n\n", n, time_taken);
//Displays sorted numbers
printf("Sorted numbers for n=%d:\n", n); for
(int i = 0; i < n; i++)
{
printf("%d", arr[i]);
}
printf("\n\n");
//Free the dynamically allocated memory free(arr);
return 0;
}

```

Output:

Random numbers for n=6000:

243 112 5996 77912 413 721 547 640 822...(more numbers).....394.....

Time taken to sort for n=6000: 1.058000 seconds

Sorted numbers for n=6000:

00 12 23 3 3 4 5(more numbers).....995 996 996 997 998 999
999

11. Design and implement C/C++ Program to sort a given set of n integer elements using Merge Sort method and compute its time complexity. Run the program for varied values of n > 5000, and record the time taken to sort. Plot a graph of the time taken versus n. The elements can be read from a file or can be generated using the random number generator.

```
#include <stdio.h>
#include<stdlib.h>
#include <time.h>
//Merge two subarrays of arr[]
//First subarray is arr[l..m]
//Second subarray is arr[m+1..r]
void merge(int arr[], int l, int m, int r)
{
    int i, j, k;
    int n1 = m - l + 1; int n2 = r - m;
    //Create temporary arrays L[n1], R[n2];
    //Copy data to temporary arrays
    L[] and R[] for (i = 0; i < n1; i++)
    L[i] = arr[l + i]; for (j = 0; j < n2; j++) R[j] = arr[m + 1 + j]; // Merge the
    temporary arrays back into arr[l..r]
    i = 0; //Initial index of first subarray
    j = 0; //Initial index of second subarray k = l;
    //Initial index of merged subarray
    while (i < n1 && j < n2)
    {
        if (L[i] <= R[j])
        {
            arr[k] = L[i]; i++;
        }
        else
        {
            arr[k] = R[j]; j++;
        }
        k++;
    }
    //Copy the remaining elements of L[], if there are any while (i
    < n1)
    {
        arr[k] = L[i]; i++;
        k++;
    }
    //Copy the remaining elements of R[], if there are any while (j < n2)
    {
        arr[k] = R[j]; j++; k++;
    }
}
```

```

}
//Mergesortfunction
voidmergeSort(intarr[],intl,int r)
{
if(l<r)
{
//Sameas(l+r)/2,butavoidoverflowforlargelandr int m =
l + (r - l) / 2;
//Sortfirstandsecondhalves mergeSort(arr,
l,m);mergeSort(arr,m+1,r);
//Mergethesortedhalves
merge(arr, l, m, r);
}
}
//Functiontogeneraterandomnumbersbetween0and999 int
generateRandomNumber()
{
returnrand()% 1000;
}
intmain()
{
//Setnvalueint n
= 6000;
//Allocatememoryforthe array
int*arr=(int*)malloc(n*sizeof(int));
//Generaterandomelementsforthearray
srand(time(NULL));
printf("Randomnumbersforn=%d:\n",n); for
(int i = 0; i < n; i++)
{
arr[i]=generateRandomNumber();printf("%d",arr[i]);
}
printf("\n");//Recordthestarttimeclock_tstart=clock();
// Perform merge sort
mergeSort(arr,0,n-1);
//Recordtheendtimeclock_tend=clock();
//Calculatethetimetakenfor sorting
doubletime_taken=((double)(end -start))/CLOCKS_PER_SEC;
//Outputthetimetakentosort forthecurrentvalueofn
printf("\nTimetakentosort forn=%d:%lfseconds\n\n",n,time_taken);
//Displaysortednumbers
printf("Sortednumbersforn=%d:\n",n); for
(int i = 0; i < n; i++)
{
printf("%d",arr[i]);
}
}

```

```
printf("\n\n");  
//Freethe dynamically allocated memory  
free(arr);  
return 0;  
}
```

Output:

Random numbers for n=6000:

243112 599677912 413 721 547 640 822...(more numbers).....394.....

Time taken to sort for n=6000: 1.058000 seconds

Sorted numbers for n=6000:

00 12 23 3 3 4 5(more numbers).....995 996 996 997998 999
999

12. Design and implement C/C++ Program for NQueen's problem using Backtracking.

```

#include <iostream>
#include<vector>using
namespace std;
//Function to print the solution
void printSolution(const vector<vector<char>>& board)
{
    for(const auto& row : board)
    {
        for(char cell: row)
            cout<<" "<<cell<<" ";cout<<endl;
    }
}
//Function to check if a queen can be placed on board[row][col] bool isSafe(const
vector<vector<char>>& board, int row, int col)
{
    int i, j;
    int n=board.size();
    //Check the row on the left side for
    (i = 0; i < col; i++)
        if(board[row][i]=='Q')return false;
    //Check upper diagonal on the left side
    for(i=row, j=col; i>=0&&j>=0; i--, j--) if
        (board[i][j] == 'Q')
            return false;
    //Check lower diagonal on the left side
    for(i=row, j=col; j>=0&&i<n; i++, j--) if (board[i][j] ==
        'Q')
            return false;return
        true;
}
//Recursive function to solve NQueens problem
bool solveNQUtil(vector<vector<char>>& board, int col)
{
    int n=board.size();
    //If all queens are placed, return true if (col>=n)
    return true;
    //Consider this column and try placing this queen in all rows one by one for (int i
        = 0; i < n; i++)
    {
        //Check if the queen can be placed on board[i][col] if
        (isSafe(board, i, col))
        {
            //Place this queen in board[i][col] board[i][col]='Q';
            //Recur to place rest of the queens

```

```

    if(solveNQUtil(board,col+1))
returntrue;
//If placingqueeninboard[i][col]doesn'tleadtoasolution,then removequeenfrom board[i][col]
board[i][col]='-';
}
}
//Ifthequeencannotbeplacedinanyrowinthiscolumn,thenreturnfalsereturn false;
}
//FunctiontosolveNQueensproblemfor4queensvoid
solve4Queens()
{
intn=4;
vector<vector<char>>board(n,vector<char>(n,'-'));if
(solveNQUtil(board, 0) == false)
{
cout<<"Solutiondoesnotexist"<<endl;return;
}
printSolution(board);
}
//Driverfunction
int main() {
cout<<"Solutionfor4Queensproblem:"<<endl;solve4Queens();return 0;
}

```

Output:

Solutionfor4Queensproblem:

```

-   -Q   -
-   Q   -   -   -
-   -   -   Q
-   Q-   -

```

VIVA QUESTIONS

1. What is an Algorithm?

Algorithm is a Step by step procedure to Solve a given problem for a finite number of input producing finite number of output with desired output.

2. What is a Flow Chart?

Flowchart is a Graphical Representation of a solution to the Problem.

3. What is the difference between Algorithm, Flow Chart, Program?

Algorithm specifies the different things to be followed for solving a Problem.

Flow Chart is a Graphical Representation of a Solution to the Problem. Both

Algorithm and Flow Chart are Machine Independent.

Program is a Set of Instructions which is used as a tool to communicate to the machine to get our work done, Program is Machine Dependent for particular Machine.

4. What is the Aim of DAA and why we need to study DAA Lab

DAA is a discipline, where we are dealing with designing or writing the algorithm keeping in Consideration of Space and Time Complexity, Such that Our Algorithms should execute in a very minimum amount of time by Minimum Space or RAM.

5. Define Space and Time Complexity? Among this which one is more prioritized? Space Complexity is a measure of Amount

of Space taken by a Program to finish its Execution. Time

Complexity is a measure of amount of time taken by a program to complete its

Execution. Depending Upon Application it is considered, EX: For Mobile or Handheld Devices, We give Preference for both Space and time.

For a Huge and Interactive Systems like Web Applications we give more Preference to time Complexity.

6. What is Design and what is Analysis of a Program?

Design is a Process of Writing an algorithm to a given Problem so that it should accept finite number of input and finite number of output with a definite output and Should Exit appropriately.

Analysis: Analysis is a next Phase of Writing an Algorithm, in this phase we calculate the Efficiency of an Algorithm i.e time and space needed by an algorithm.

7. Write the general plan for analyzing the recursive algorithms.

- Identify the inputs.
- Identify the output is depended only on number of inputs.
- Identify the Basic Operation in Algorithm.
- Form or write the Recursive Relation to the Algorithm.

8. What are the various notations used to write an algorithm?

(i) Pseudocode (ii) Natural Language and etc..

9. What is a Pseudocode?

It's an notation which is having the combination of Programming Constructs and English like Statements.

10. What is the Time Complexity of Bubble Sort, Selection Sort, Merge Sort, Quick Sort? (L3)

Bubble Sort - n^2 , Selection Sort - n^2 , Merge Sort - $n \log n$, Quick Sort - $n \log n$, Worst case for Quick Sort - n^2

11. Which sorting algorithm is more Efficient and why?

Quick Sorting is More Efficient, because this algorithm is an in-place algorithm.

12. What do you mean by the term Instable Algorithms?

The Instable Algorithms are one, which divide the array as certainly depending upon pivot or key element and hence i index precedes index j

13. Which algorithms are faster?

Instable Algorithms are much Faster compared to Stable Algorithms.

14. For what type of instance Merge Sort does better than Quick Sort?

For a Larger input and a sorted input values.

15. For what type of instance Quick Sort does better than Merge Sort?

For a Smaller Set of input numbers.

16. What are Inplace Algorithms?

Inplace Algorithms are the one which doesn't occupy Extra Space.

17. Write the order of growth terms as per the time Execution in Ascending Order.

$\log n, n, n \log n, n^2, n^3, \dots, n^n, 2^n, n!$

18. What is Brute Force Technique? When We Should Use?

Brute Force is a straight Forward Technique to solve a problem, We used to solve a Problem through this approach when we don't have sufficient data to solve a problem in Efficient Way.

19. What is the difference between Divide and Conquer, Decrease and Conquer?

Divide and Conquer can be solved to solve a problem with a larger data set and when there is no dependency between any of the data sets.

❖ Divide and Solve as Small as Small sets.

Conquer or Merge it to get one final resultant data set.

Decrease and Conquer is almost similar to Divide and Conquer but we are finding a solution to the problem in a different variations, EX: Decrease by Constant (Usually by One), Decrease by Constant factor which is almost similar to Divide and Conquer Technique (Usually by two), Decrease by Variable (The Dividing Criteria changes for each iteration depends upon the data set).

20. Define Greedy Technique.

Greedy Technique is always applied for the problem of the type optimization type, which reduces loss and increases profit.

21. Define Optimal and Feasible Solution.

Optimal Solution is a solution which is best among N Feasible Solution. Feasible Solution is a solution which Satisfies a Problem Constraints/conditions.

22. Can A Problem solved by all the algorithmic Techniques.

Yes, but some problems will give better results with some Algorithmic Technique and it may give worst result when it is applied with other technique.

23. State and Explain Knapsack Problem.

Filling the Maximum number of items to the Knapsack (Container) Which Increases the profit and decreases the Loss.

24. Which one is Most Admired algorithmic Technique?

Dynamic Programming.

25. What is Spanning tree and Minimum Spanning tree?

A tree Without Cycles are called as Spanning tree. A Minimum Spanning Tree is a spanning tree which yields the very less Cost when all the edges cost summed up.

26. How Many Spanning Tree can a Tree can have?

A tree can have 1 to many number of Possible ways of Spanning Tree.

27. Differentiate between Prim's and Kruskal's Algorithm for finding MST.

In Prim's We consider any one vertex in the graph as Source and We compute the distance from that source to other vertices, after computing the vertices which has minimum value among $(n-1)$ vertices is added to tree vertices and that respective edges added to tree Edges Set. The above mentioned Process continues till we reach $(n-1)$ vertices.

In Kruskal's we first arrange the edges in Ascending Order and then we start to form the tree which won't form cycles, if adding that edges forms cycles then that edges is dropped from adding to tree edges. The above said process is continues till we reach the count of $(n-1)$ Vertices.

28. What is the Application of Prim's and Kruskal's Algorithm?

In Network to remove the Cyclicity of the Network.

29. Explain Job Sequencing With Deadlines?

Placing or scheduling the maximum number of Jobs to a machine without violating the deadlines constraint of any of the Jobs in Sequence.

30. Why the Name Bubble Sort named?

Because in first Pass the first highest data will bubble up, so since the largest element bubbles up in the first and second largest element bubbles up in the Second pass and so on, so hence the name bubble sort.

31. Why the Name Selection Sort? (L3)

The Selection sort is named because we initially first select an array's first element as minimum and will compare with other elements, so in pass one first least element goes to the first position and so on so forth for 2nd, 3rd and so on. Selecting

32. What is the difference between Brute force string matching to Horspool String Matching Method? (L2)

In brute force we compare each and every element of the text to the pattern by shifting the text position by one and in Horspool method we shift it by number of shift positions recorded in the shift table.

33. Explain Merge Sort?

In Merge Sort will divide the entire input set by 2 until we reach low < high and later will find a solution to each item by comparing half of the array dataset to the other half array dataset and finally we merge it to form a single array (conquer)

34. What are the basic operations in Merge sort and Quick sort?

In Merge Sort the basic operation is comparison and in Quick sort basic operation is Partitioning and hence also known as partitioning sort.

35. Why the Insertion Sort?

We are inserting an element to its suitable place by comparing elements for each pass.

36. What is the use of DFS and BFS?

DFS and BFS both are used to check the connectivity of a graph, cyclicity in a graph, spanning tree of a graph.

37. Differentiate between DFS and BFS.

DFS and BFS are both the graph traversing techniques, in which DFS traverses the graph in a depthwise (vertical) and BFS traverses the graph from left to right (horizontal)

38. Which data structures are used in BFS and DFS.

BFS uses queue as its data structure and DFS uses stack as its data structure.

39. What are back edges in DFS and cross edges in BFS.

Back edges and cross edges are the edges which are already visited by an ancestor node.

40. What is Topological Sorting?

Topological Sorting is a sorting technique used for sorting vertices in the graph.

41. What are the conditions necessary for a Topological Sorting?

For a Topological Sorting the graph should be DAG (Directed Acyclic Graph)

42. What are the different methods used to solve a topological Sorting?

1. Source Removal Method
2. Using DFS based Scheme.

43. What is the use of Topological Sorting?

Use of Topological Ordering is in the field of Operating System for scheduling and in Networks, Automation and Robotics.

44. What is Dijkstra's Algorithm?

Dijkstra's Algorithm is used to find the single shortest path from source to the other vertex.

45. What is a graph?

Graph is a component which is having a set of Edges and vertices $G = \{V, E\}$

46. What are the different ways that can represent a graph?

Adjacency Matrix and Adjacency List.

47. What is Adjacency Matrix?

Is a Matrix which illustrates the Graph in the form of Matrix, if it is a weighted Graph then we initialize the value of the cost in that position (i,j) or else simply we write 1 to mention there exists an edge between (i,j) or else we use 0 or 9999 to mention non connectivity of a graph.

48. What are the limitations of Algorithms?

Algorithm can't find the better solutions when we come across the tight lower bound, So we can find the better solutions which is not possible with Algorithmic way. To find the tight lower bound we use Decision Trees.

49. What is Tight lower Bound?

It is a Lower bound which is a best lower bound for a problem, beyond that no algorithm will produce better results.

50. What are Decision Trees?

Decision trees are also known as Comparison Trees used to find the tight lower bound for a particular Problem EX: Tight Lower Bound For Sorting is $n \log n$ and tight lower bound for Searching is $\log n$ which is not possible to get the better result.

51. What is a polynomial problem (P-type)?

P-type problems are decision problems in which we can find the solutions in a polynomial time and is of type deterministic.

52. What is NP-problem?

NP-Problem belongs to decision problems and these problems are Non-Deterministic Polynomial, i.e. for which the problem doesn't have deterministic solutions and can be solved in Polynomial time. There are 2 phases in solving a problem.

(i) Guessing (Non-Deterministic stage) Producing N number of Candidate Outputs.

(ii) Verification (Deterministic Stage) Verifying the correctness of N Number of Candidate Outputs.

53. What are NP-Complete Problems?

NP-Complete Problems belong to Decision problems and NP type Problems. These problems can be found the solutions by converting or reducing to the problem which we know the Solutions.

54. What is a trivial lower bound?

Trivial bound can be derived by formulating the number of inputs that have to be given and number

of output that has to be generated.

55. Explain Backtracking W.r.t

(i) Subset Problem (ii) N-Queens Problem

56. Explain Subset Problem.

In a given Set S , find the Subset, in which the sum of all subset elements is equal to the sum d which is predefined in a problem.

57. Explain N-Queens Problem.

N-Queens Problem is of Placing a N-Queens in a $N \times N$ Chess board such that No 2- Queens Should be placed in the same Row, Column and same diagonal (N=Should consider both principal diagonal elements)

58. What is Hamiltonian Circuit? (L2)

Hamiltonian circuit is a problem in which that circuit starts from a source vertex and has another vertex in any order without repeating and Should end with the Source vertex only i.e source and Destination vertex should be same.

59. Explain the Problem of TSP. (L2)

Travelling Sales Person Problem is a problem in which he should Visit N number of cities with a minimum number of Cost by visiting every city Exactly one and the city what he is started should end with same city.

60. What is the Concept of Dynamic Programming? (L3)

Deriving a Solution to the basic Condition and Extracting the solutions for the rest of the other data sets by Previously drawn Solution. Computer to other algorithmic Technique Dynamic Programming because it avoids lot of reworking on the same Solution what we have solved in the earlier phases of deriving the solution to the problem.

61. What is the goal of Warshall's Algorithm? (L3)

Warshall's algorithm is to find the shortest distance between a node to all the other nodes in the graph. It Uses the Property of Transitive Closure i.e if there exist a path between (i,k) and (k,j) then there surely exists a path between (i,j)

$(i,k) \& (k,j) \rightarrow (i,j)$ What is the use of Floyd's algorithm? (L3) It is used to find

the All pairs shortest Path of an Graph.