Operating Systems CSE 511

Assignment 3 Directions File Systems

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- The document will be released tonight
 - You can already get started with the code https://classroom.github.com/a/uBT6S-5T
 - Academic integrity is very important! (do not use any online implementations)
 - Use our slides and official manuals!
 - You need to implement lock-based and lock-free queues and stacks
 - Use pthread_mutex for lock-based
 - Use C11 atomics for lock-free

- Review pthread_mutex documentation
- Review C11 atomics (stdatomic.h)
 - atomic_compare_exchange_strong (CAS)
 - atomic_load (LOAD)
 - atomic_store (STORE)
 - You do not need atomic_compare_exchange_weak!

- For ABA queues/stacks
 - Create two stacks/queues (e.g., store two head and two tail pointers for queues)
 - e.g., can create a separate structure with head/tail and then create 'aq' and 'fq' instances inside the queue object
 - Allows to use the same code for aq and fq
 - Recycle elements through 'fq'

- [Hint] Stacks: need tagging only for 'top'
 - Not needed for 'next' (think why)
- [Hint] Queues: need tagging everywhere (head, tail, next)
- When you need tagging, create a separate structure like this

```
typedef struct aba_ptr {
    void *ptr; // or any other pointer type
    size_t tag;
} aba ptr;
```

- Just pass pointers/values of aba_ptr directly to atomic_compare_exchange_strong, atomic_load, atomic_store like you pass pointers/values to regular pointers
- Use _Atomic(...) to encapsulate custom atomic types, e.g., _Atomic(queue_node *) head or _Atomic(aba_ptr) aba_head.

Filesystem interface

- What is a file? Named storage
- File operations:
 - read(fd, buffer, length)
 - write(fd, buffer, length)
 - fd = open(pathname, flags)
 - close(fd)

syscalls (unbuffered)

- fread(buffer, size, num, file)
- fwrite(buffer, size, num, file)
- file = fopen(pathname, mode)
- fclose(file)

Buffered C library

File descriptor (fd)

- Purpose: identify file being accessed
- What actually is a file descriptor?
 - 0-indexed number index of file descriptor table

File descriptor	Referencing
0 (stdin)	Console
1 (stdout)	Console
2 (stderr)	Console
3	My_open_file.txt

- Important points:
 - Each process has its own table
 - There is a limit (default 1024)

Filesystems

- Purpose: making storage usable
- 1. Tracking file information (metadata)
- 2. Mapping files to blocks (i.e., LBN)
- 3. Tracking/allocating free blocks

File metadata

Metadata: data that's not the actual data

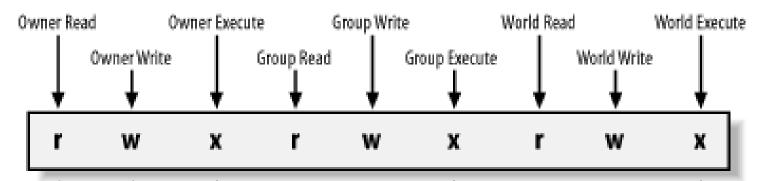
- Examples:
 - Filename
 - Size of file
 - Last access timestamp
 - Last modification timestamp
 - File owner
 - Permission
 - Location of data

inode

File permission

- Files have owners/groups
 - Use chown to modify owner/group
- File permission
 - Use chmod to modify permission

Unix Permissions



Mapping files to blocks (i.e., LBN)

- Simple: linked-list of data blocks (FAT)
- Better: indexed table (ext2/ext3)
- Problem: how big of a table?
- Solution: multi-level index

inode

owners (2)

timestamps (3)

size block count

data

ext4 uses an "extent tree" instead

Tracking/allocating free blocks

- Simple: free list (slow)
- Better: bitmap (easy to scan)
- Complex: list of extents (i.e., ranges of free blocks)