

CSE 511: Operating Systems Design

Lectures 22

TAS and CAS

What Should you do if you can't get a lock?

Keep trying

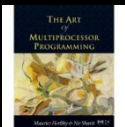
“spin” or “busy-wait”

Good if delays are short

Yield the processor

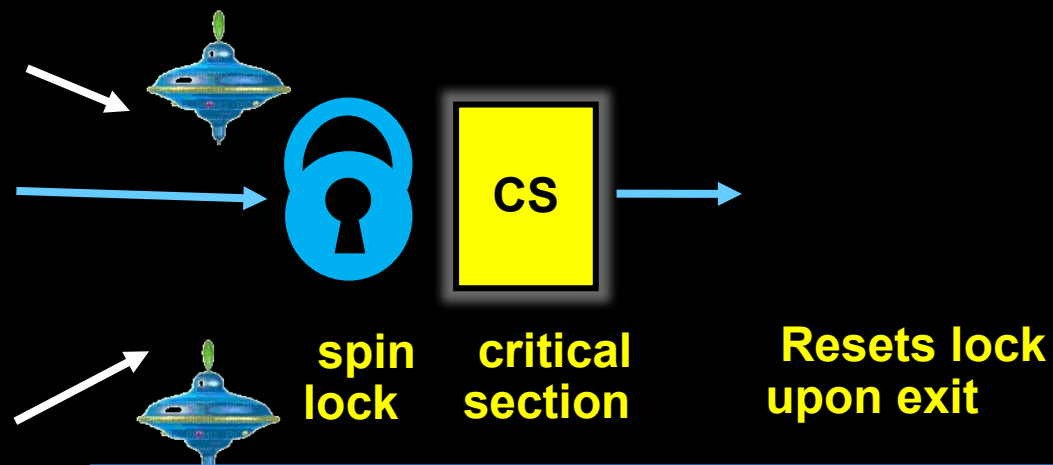
Good if delays are long

Always good on uniprocessor

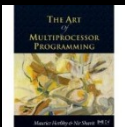


Basic Spin-Lock

lock means sequential bottleneck



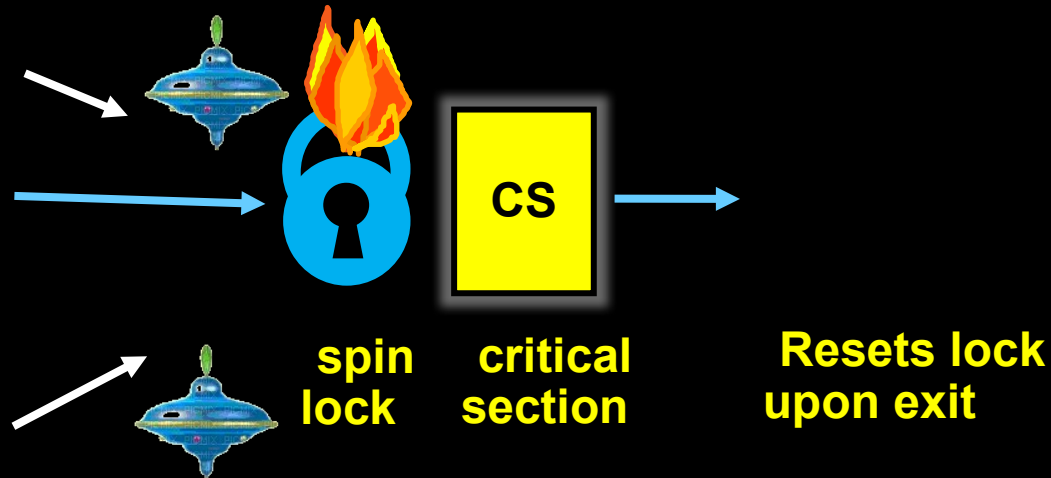
Bottleneck means no parallelism



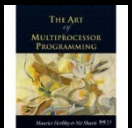
Basic Spin-Lock

Lock suffers from contention

Honk!



Honk!



Test-and-Set Lock

Boolean state

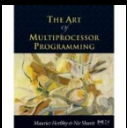
Lock: Test-and-set (TAS)

Swap *true* with current value

Returns prior value

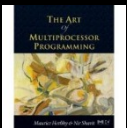
Unlock: write *false*

TAS sometimes called “getAndSet”



Test-and-Set

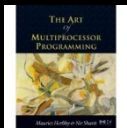
```
bool getAndSet(bool *globalVal, bool newValue) {  
    /* Done atomically (as a single instruction)! */  
    bool prior = *globalValue;  
    *globalValue = newValue;  
    return prior;  
}
```



Test-and-Set

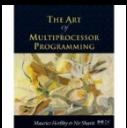
```
bool getAndSet(bool *globalVal, bool newValue) {  
    /* Done atomically (as a single instruction)! */  
    bool prior = *globalVal;  
    *globalVal = newValue;  
    return prior;  
}
```

Swapping in *true* is
called “*test-and-set*” or
TAS
(usually a HW machine
instruction)



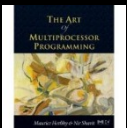
TAS-based Lock

```
struct spinlock {  
    bool value;  
};  
  
struct spinlock lock = { false };  
  
void spinlock_lock(struct spinlock *lock) {  
    while (getAndSet(&lock->value, true)  
           == true) {  
        /* Retry */  
    }  
    /* The previous value was false! */  
}
```



TAS-based Lock

```
struct spinlock {  
    bool value;  
};  
  
struct spinlock lock = { false };  
  
void spinlock_unlock(struct spinlock *lock) {  
    lock->value = false;  
}
```



Test-and-Set Locks

Lock is free: value is *false*

Lock is taken: value is *true*

Spin: repeatedly call TAS

When result is *false*, stop

While result is *true*, try again

Release lock by writing *false*



Space Complexity

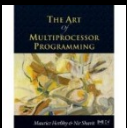
TAS spin-lock has small “footprint”

n thread spin-lock uses $O(1)$ space

As opposed to $O(n)$ Peterson/Bakery*

* Please review these algorithms for classical RW-approaches to implement locks

TAS is RMW, not RW ...



Compare-and-Set

```
bool compareAndSet(long *globalVal, long oldValue,
                  long newValue)
{
    /* Done atomically (as a single instruction)! */
    long prior = *globalVal;
    if (prior == oldValue) {
        *globalVal = newValue;
        return true;
    } else {
        return false;
    }
}
```

