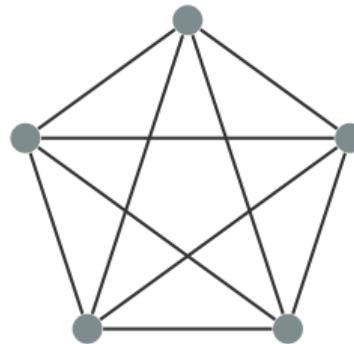
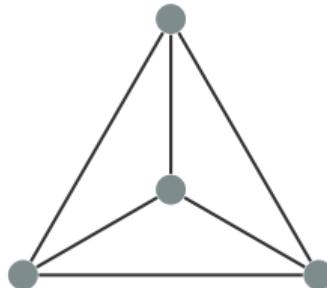


On Planar Graphs of Uniform Polynomial Growth

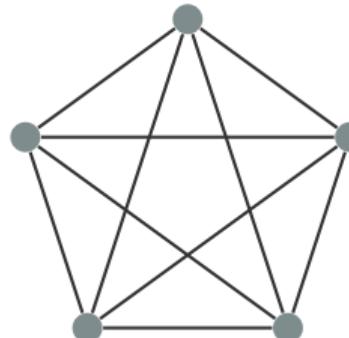
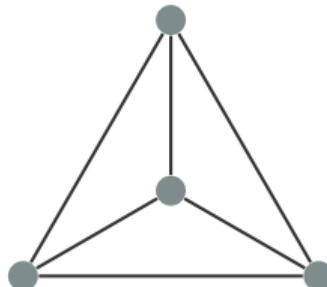
Farzam Ebrahimnejad

joint work with James R. Lee

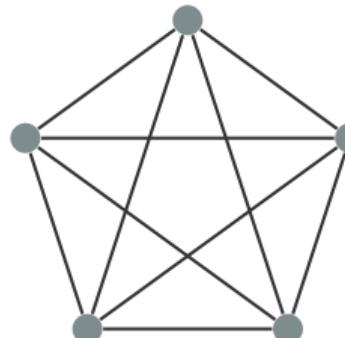
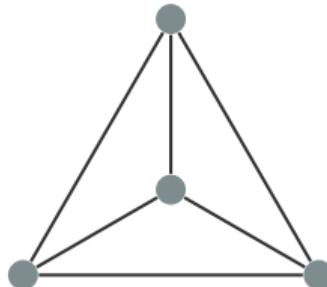
Planarity



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Volume Growth

Definition (Graph distance ball)

For a graph G , $v \in V(G)$ and $R \geq 0$ we define

$$B_G(v, R) = \{u : d_G(v, u) \leq R\},$$

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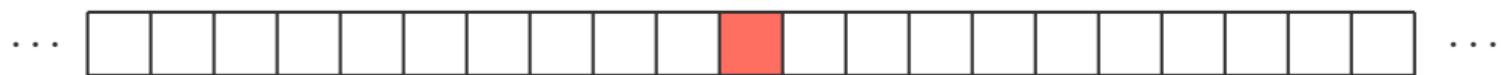
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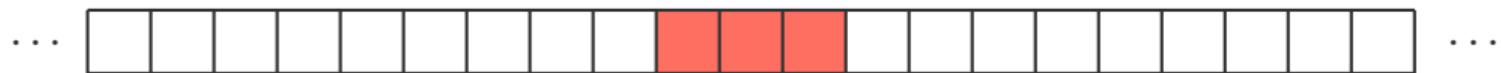
- ◎ We are interested in how $|B(v, R)|$ grows as a function of R .
 - Let's see some examples!

Volume Growth in \mathbb{Z}



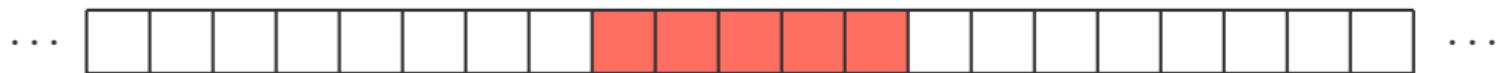
$$|B(v, 0)| = 1$$

Volume Growth in \mathbb{Z}



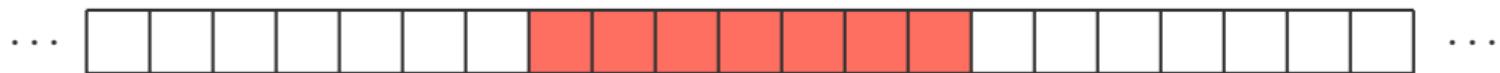
$$|B(v, 1)| = 3$$

Volume Growth in \mathbb{Z}



$$|B(v, R)| = 2R + 1$$

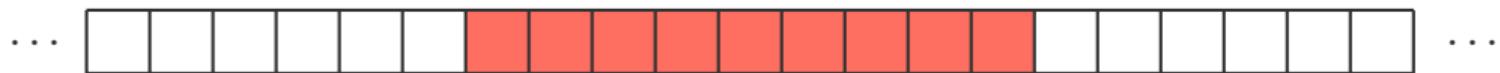
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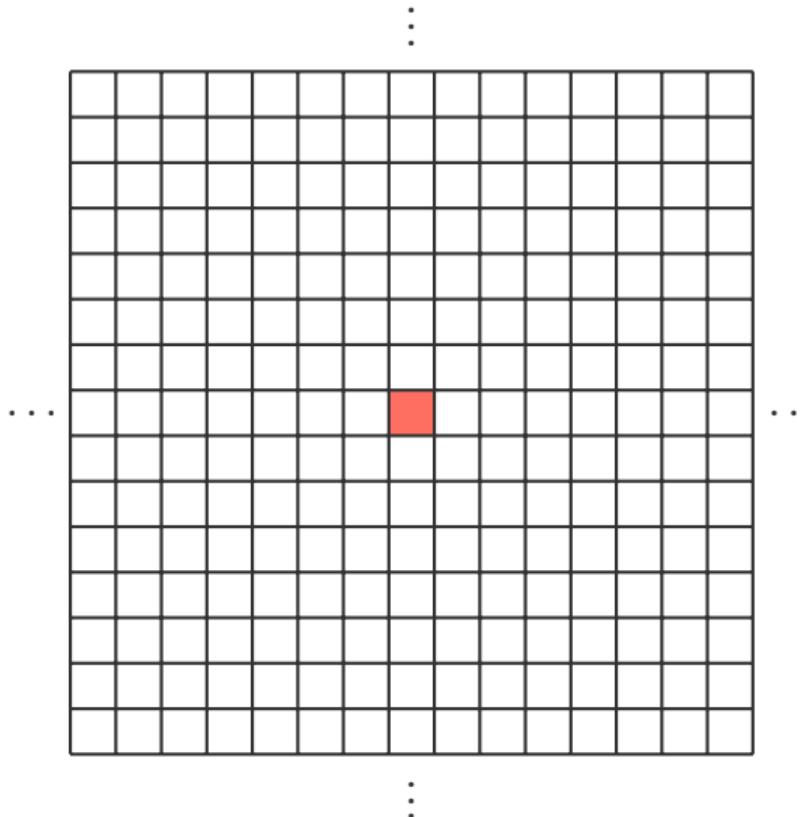
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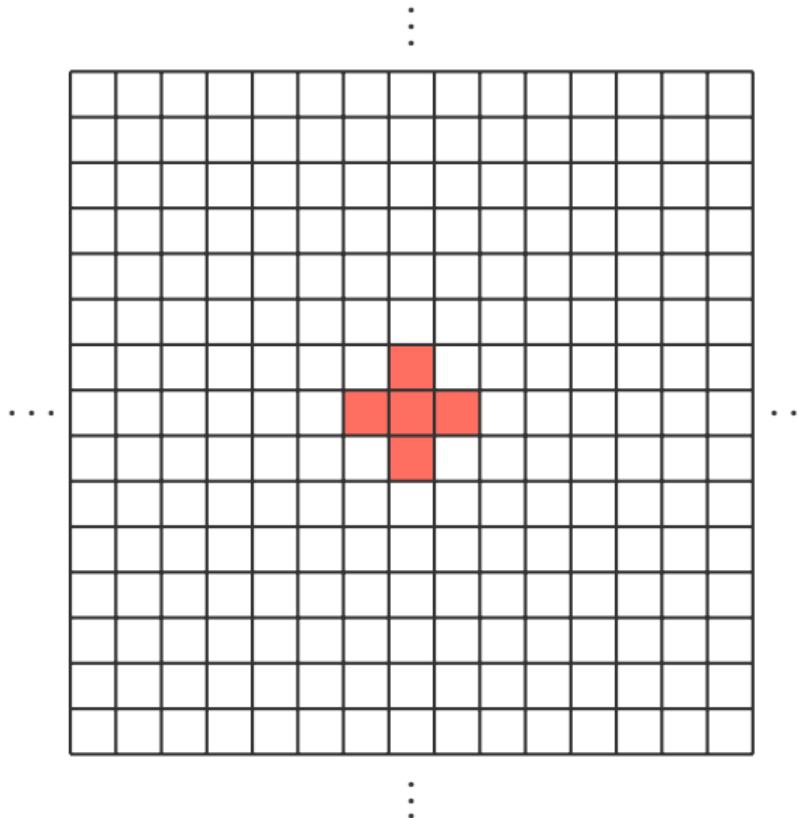
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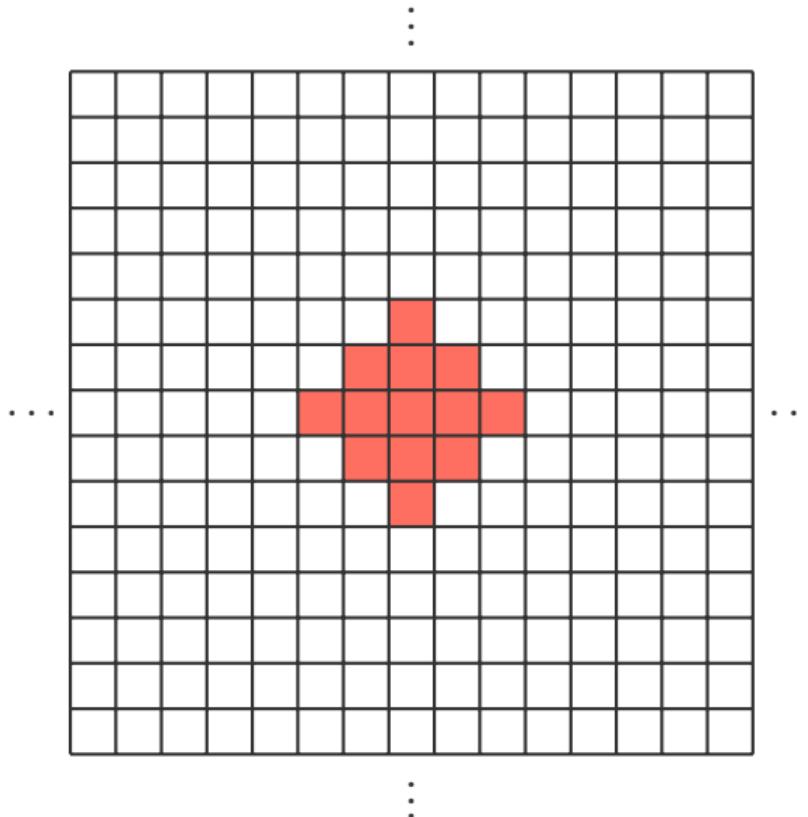
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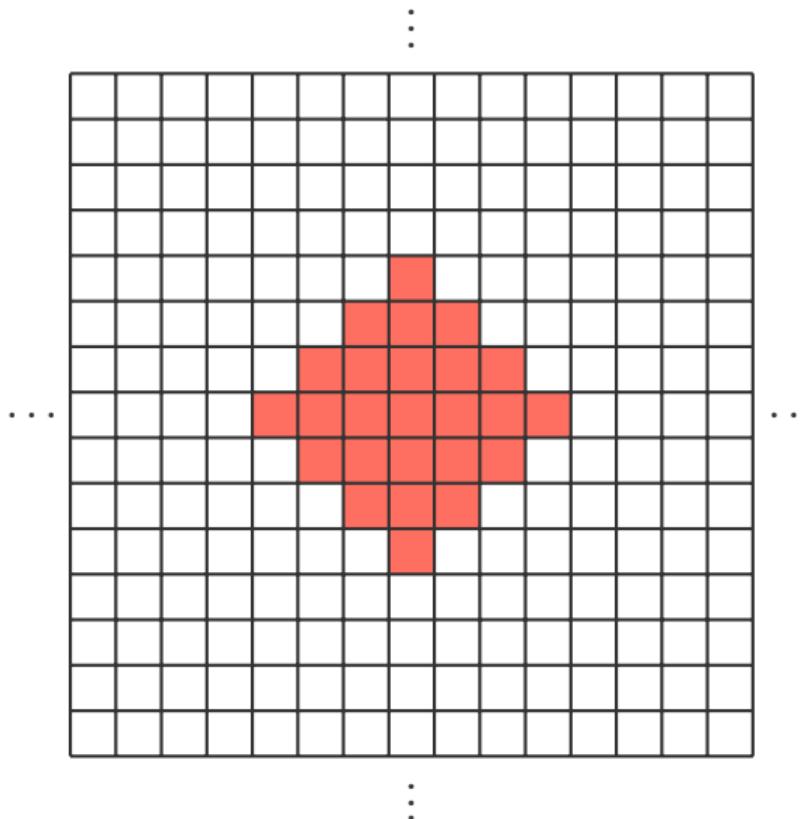
$$|B(v, 1)| = 5$$

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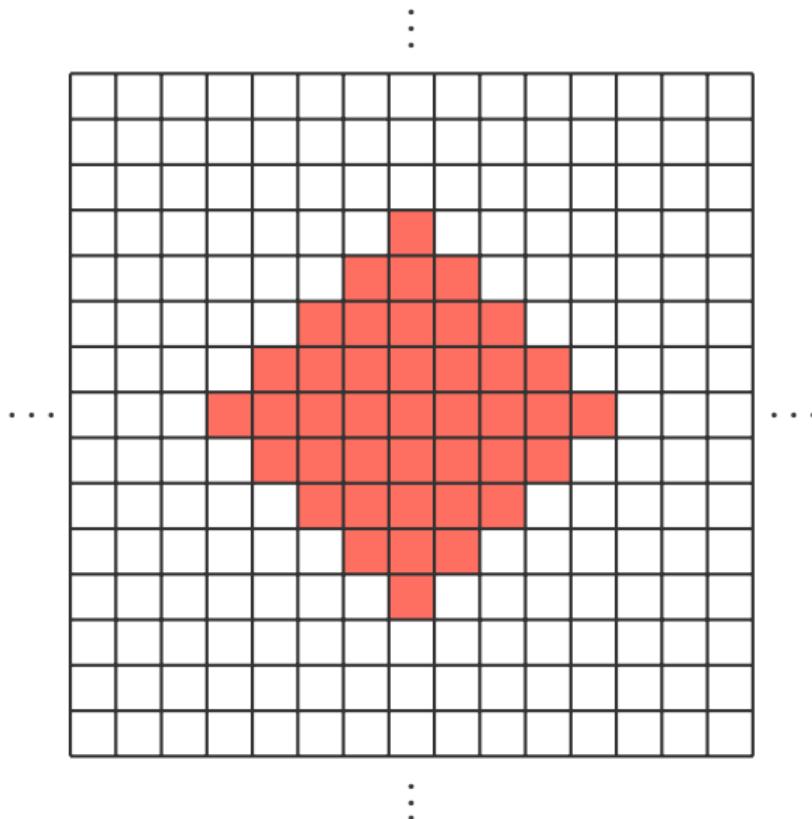
$$|B(v, 2)| = 13$$

Volume Growth in \mathbb{Z}^2



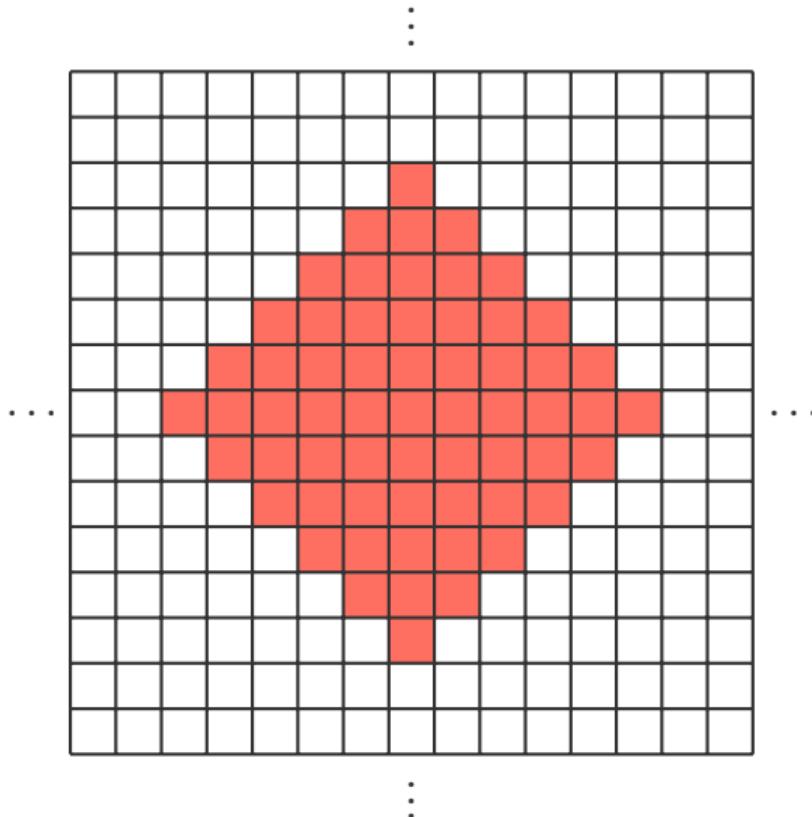
$$|B(v, R)| = 2R^2 - 2R + 1$$

Volume Growth in \mathbb{Z}^2



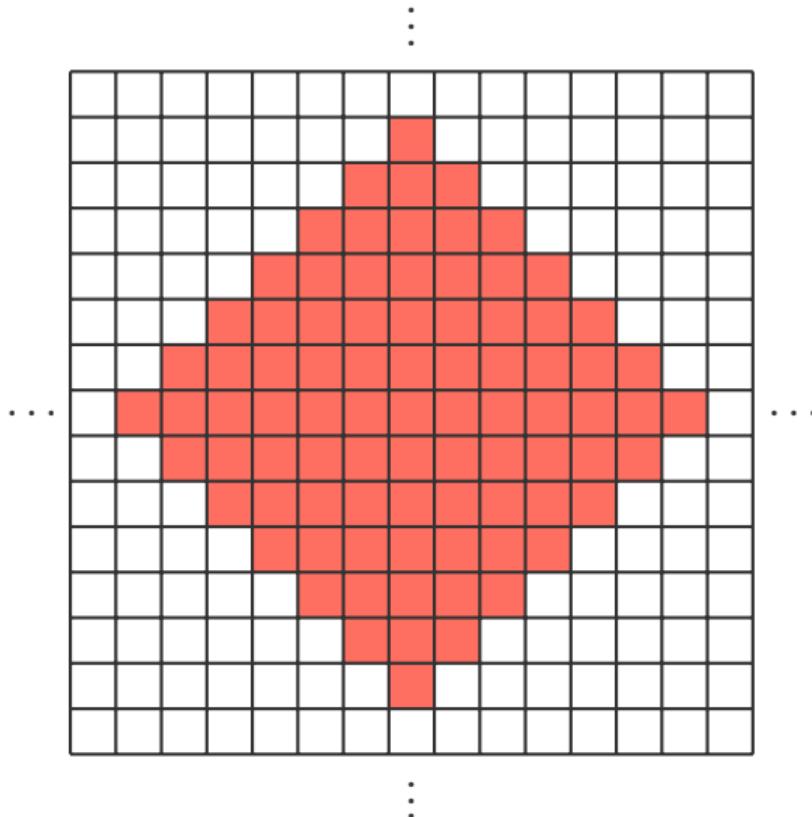
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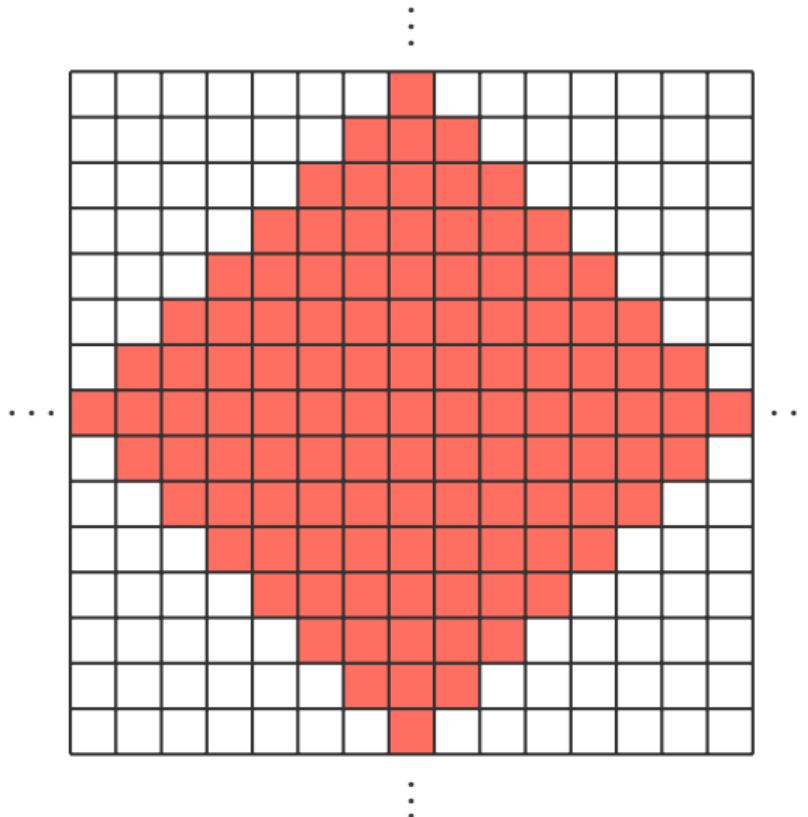
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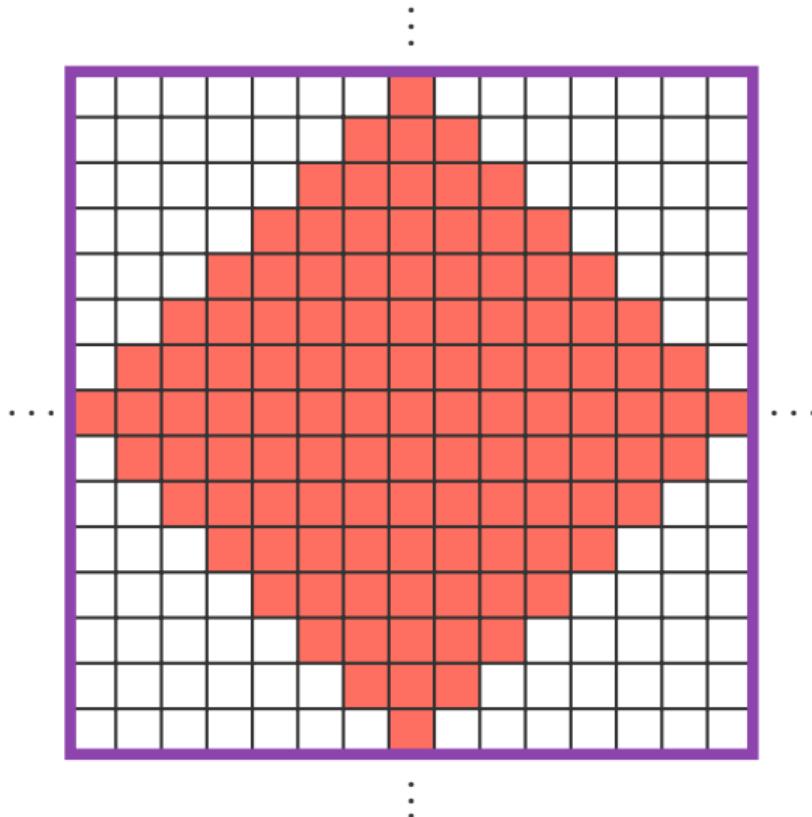
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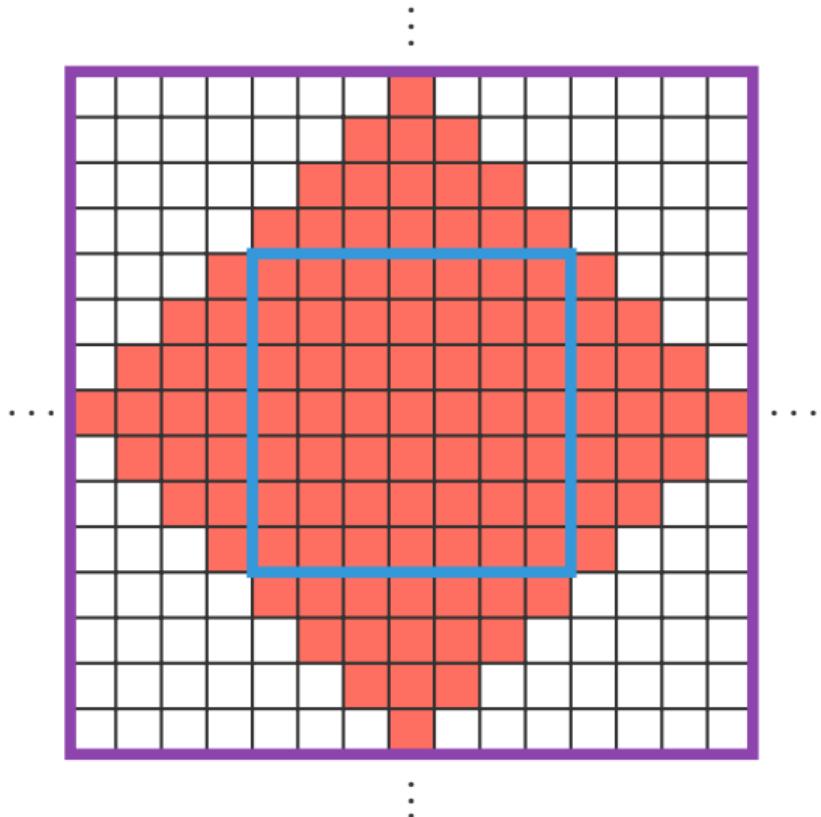
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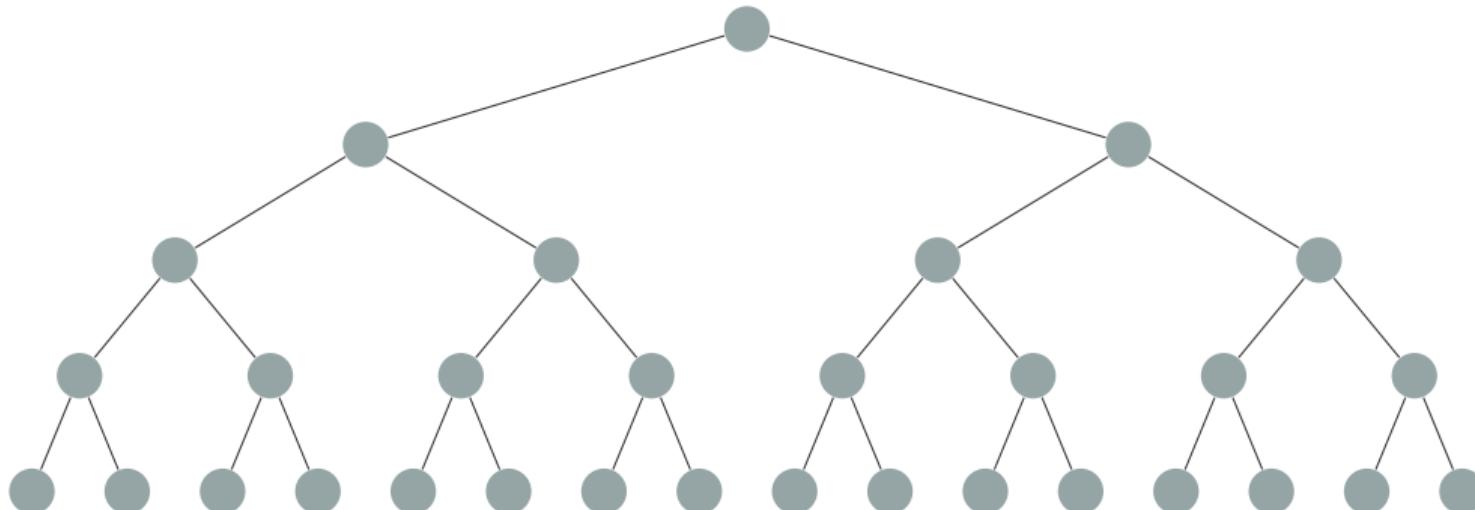
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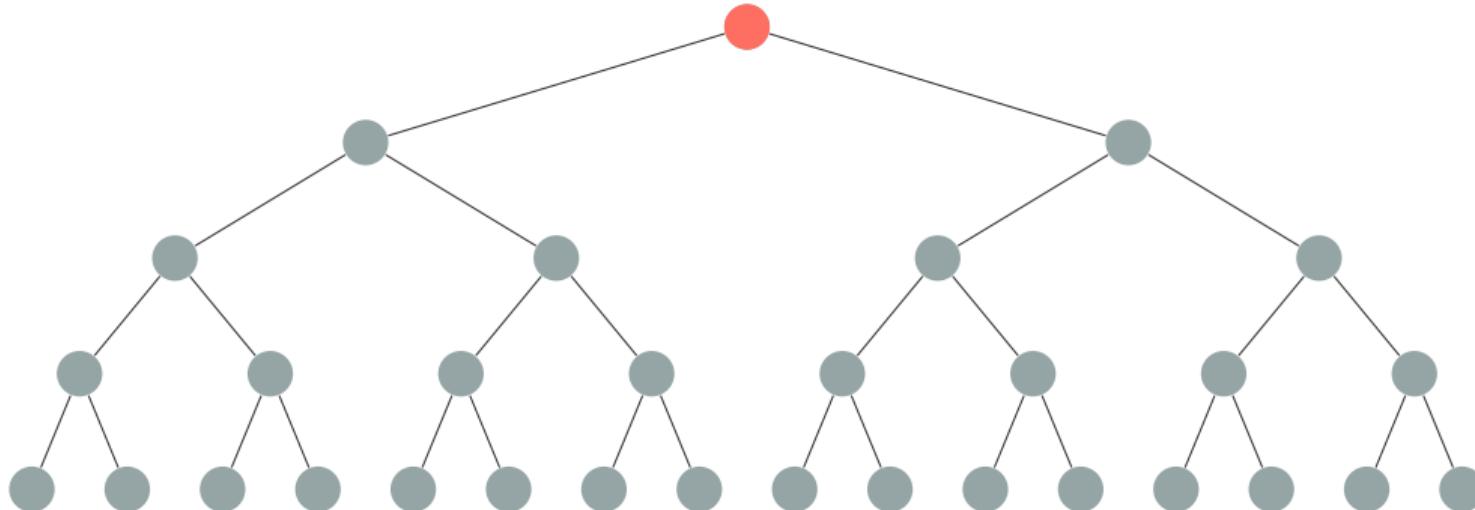


$$(R+1)^2 \leq |B(v, R)| \leq (2R+1)^2$$

Volume Growth in the Binary Tree

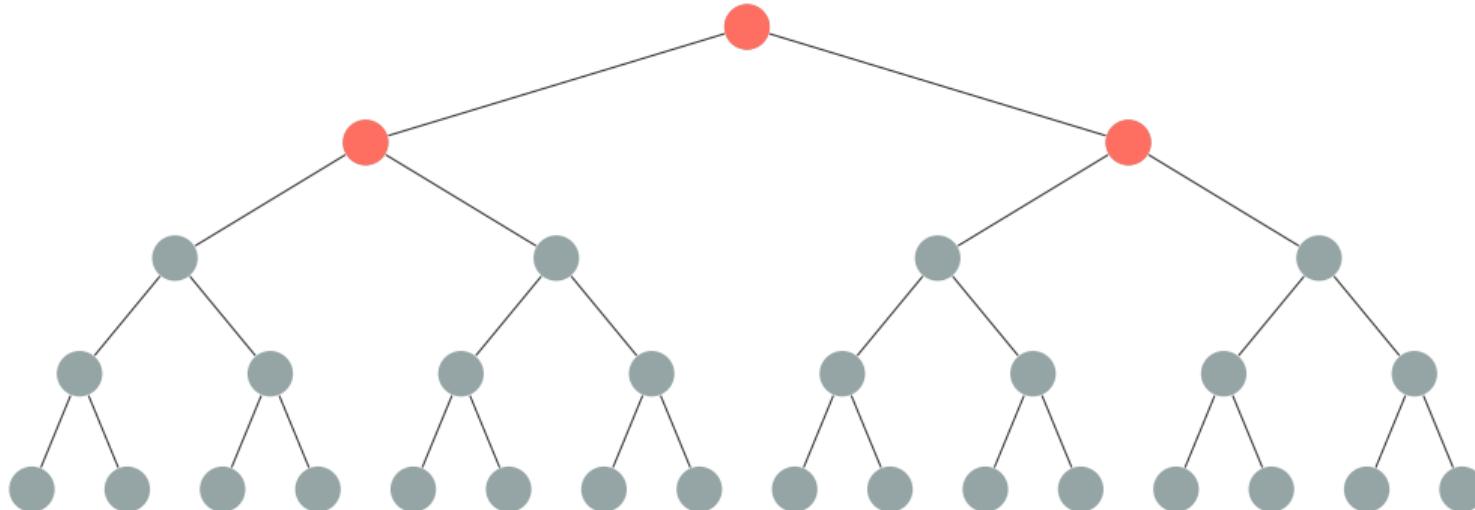


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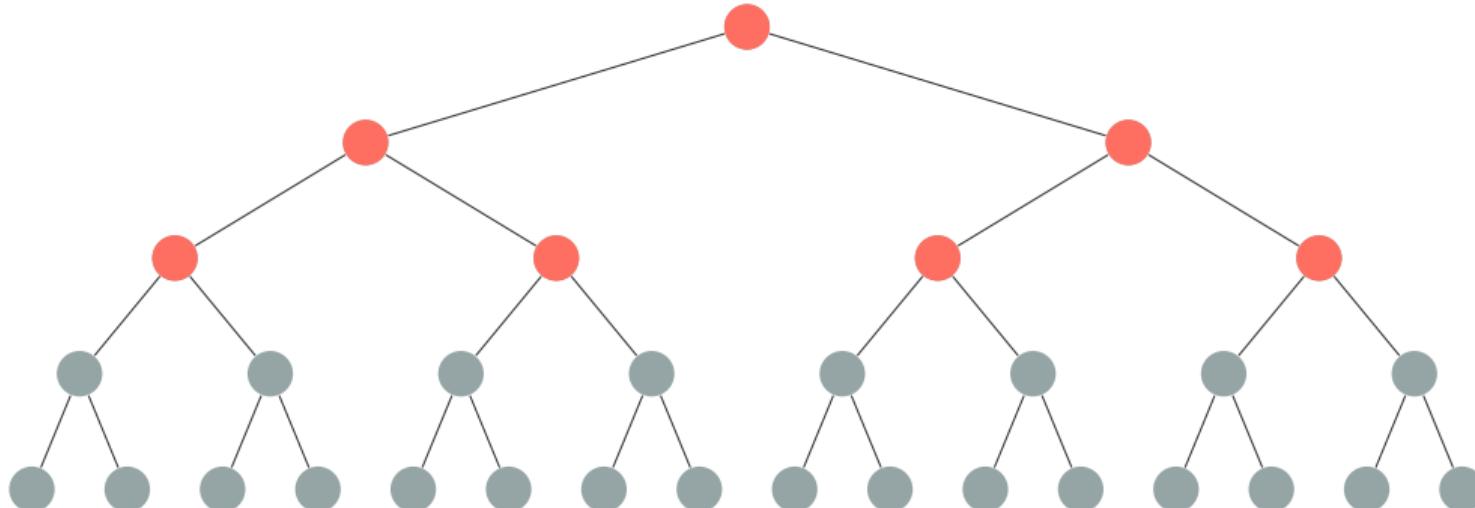
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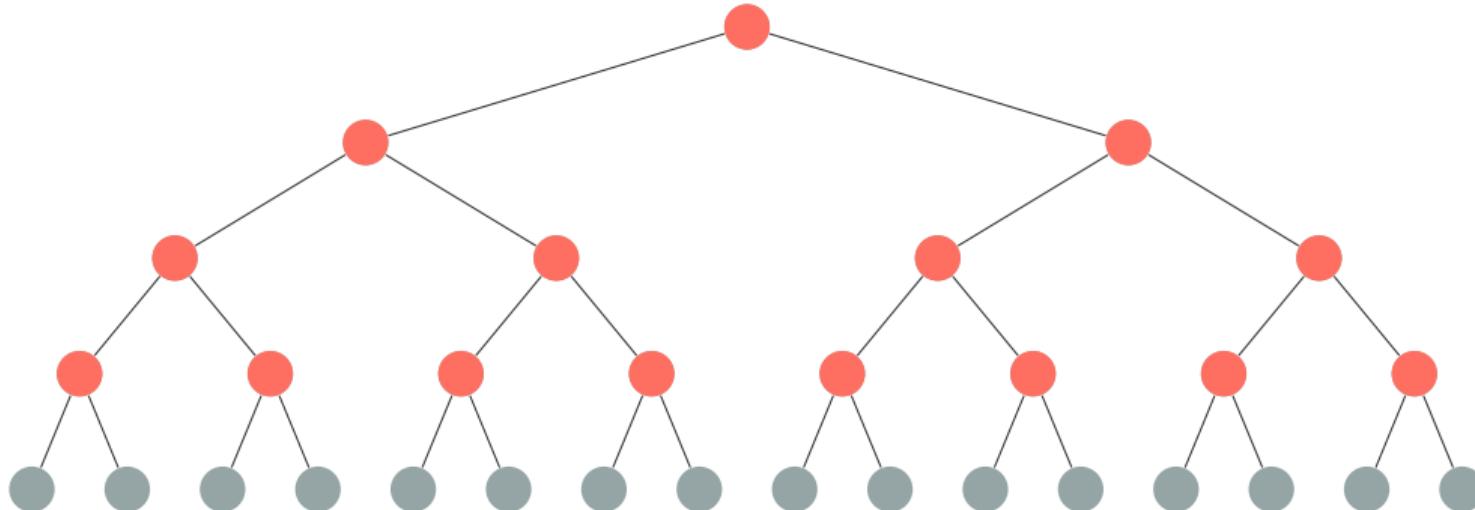
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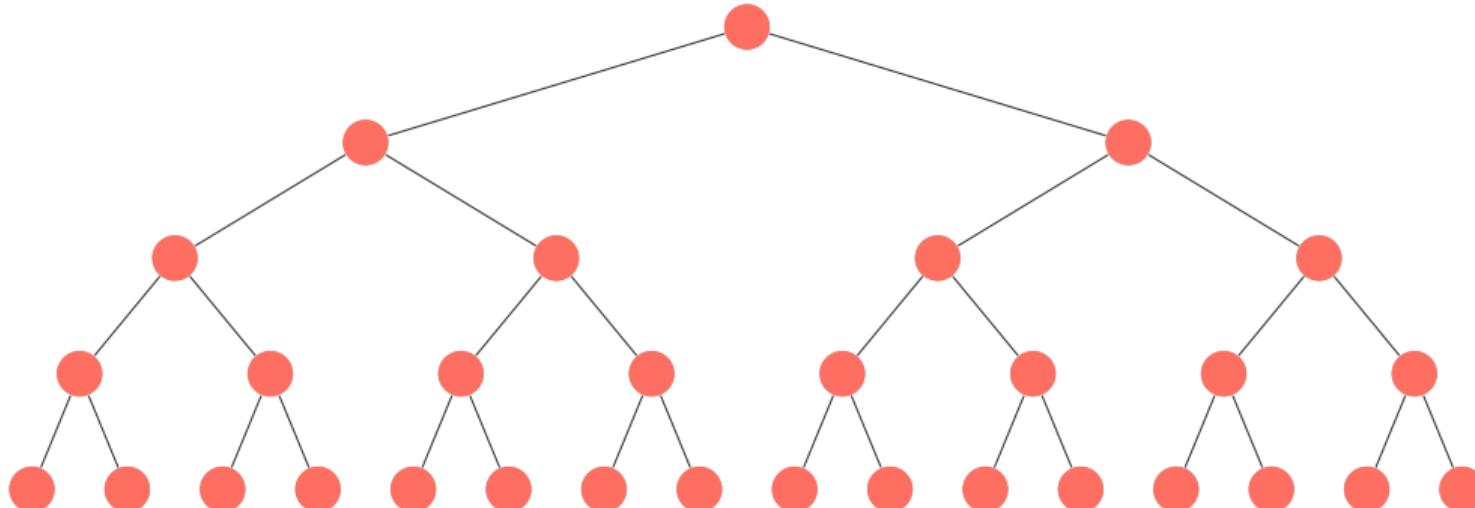
$$|B(v, 2)| = 7$$

Volume Growth in the Binary Tree



$$|B(v, R)| = 2^{R+1} - 1$$

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We say a graph G has *uniform volume growth* $f(R)$ if for all $v \in V(G)$ and $R \geq 0$ we have

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- ◎ Not all graphs have uniform growth rate.

On Planar Graphs of Uniform Polynomial Growth

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joint work with James R. Lee

Planar Graphs with Uniform Polynomial Growth > 2

Thm. [Babai '97]

If G is a *vertex-transitive* planar graph with uniform growth, then the growth is either **linear**, **quadratic**, or **exponential**.

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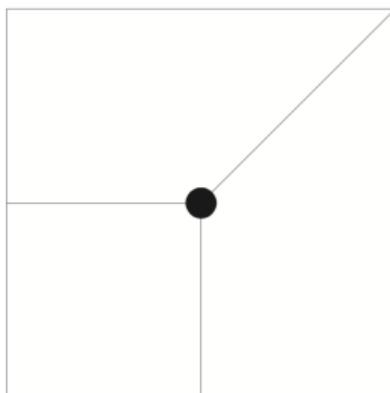
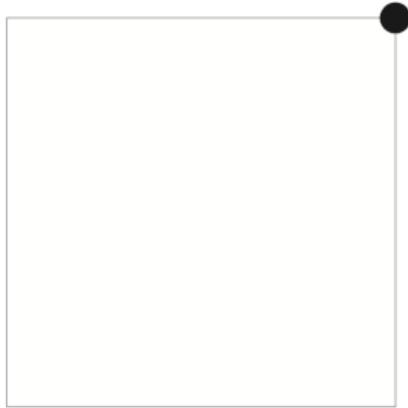
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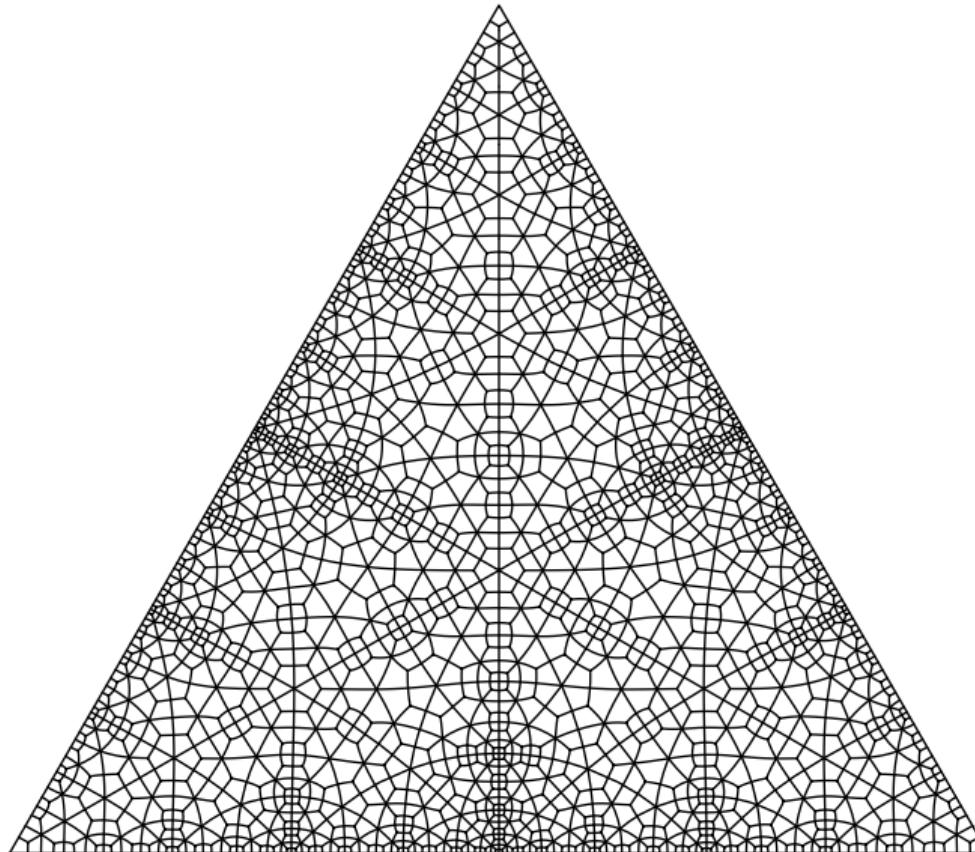
Thm. [Benjamini - Schramm '01]

For any $\alpha \geq 1$, there exists a planar graph with uniform growth R^α .

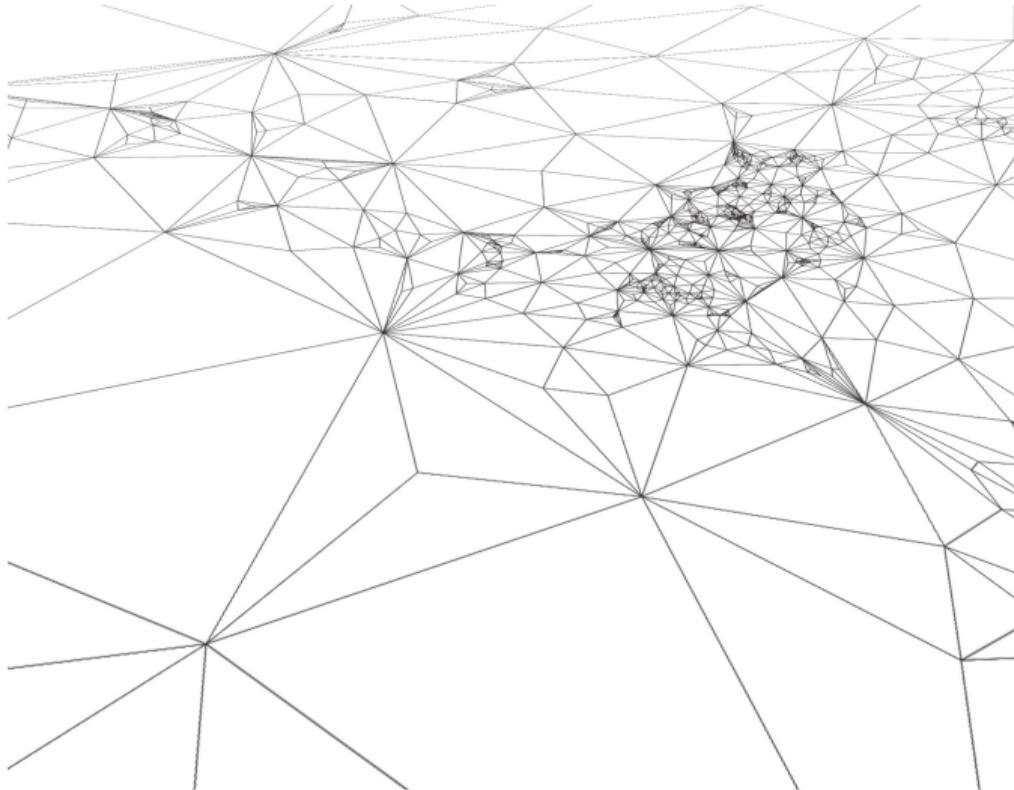
A Construction by Benjamini-Schramm



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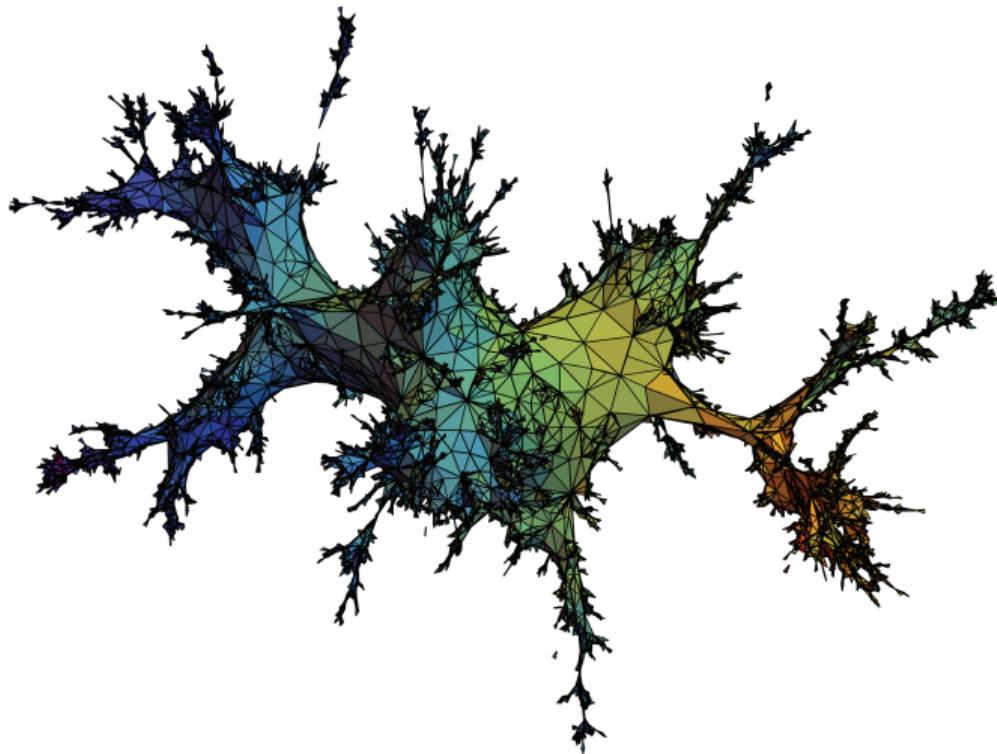


Uniform Infinite Planar Triangulation (UIPT)



Drawn by Igor Kortchemski

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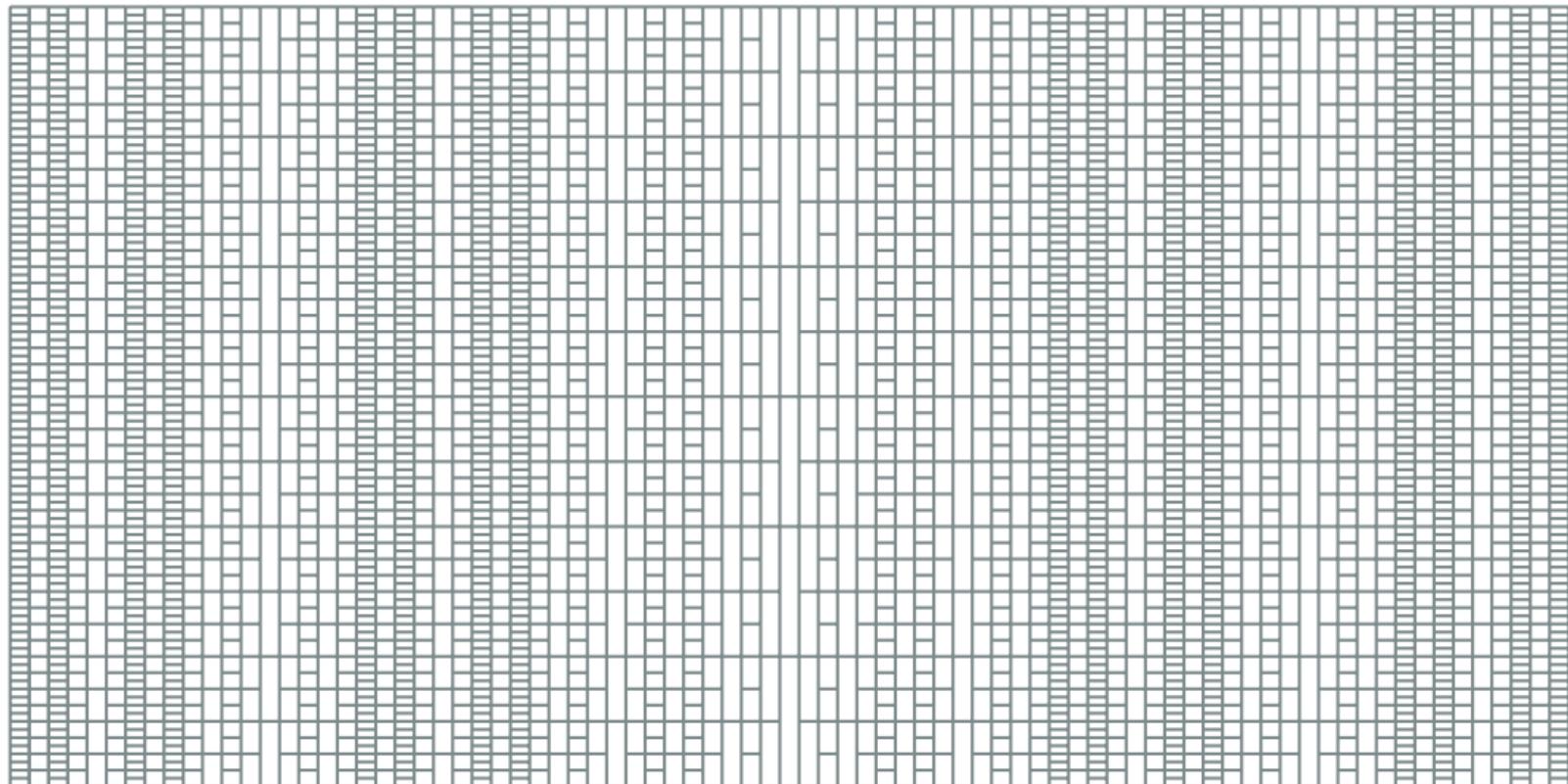
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Thm. [E - Lee '20]

For **any $\alpha > 2$** , there exists a planar graph with uniform growth R^α in which the complements of balls are all connected.

Sneak Peek



So far I have covered ...

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Up next ...

- *Random walks* and *effective resistances* in these graphs.
- Cannon's conjecture.
- Our construction.

Random Walks

Random Walk on Graphs

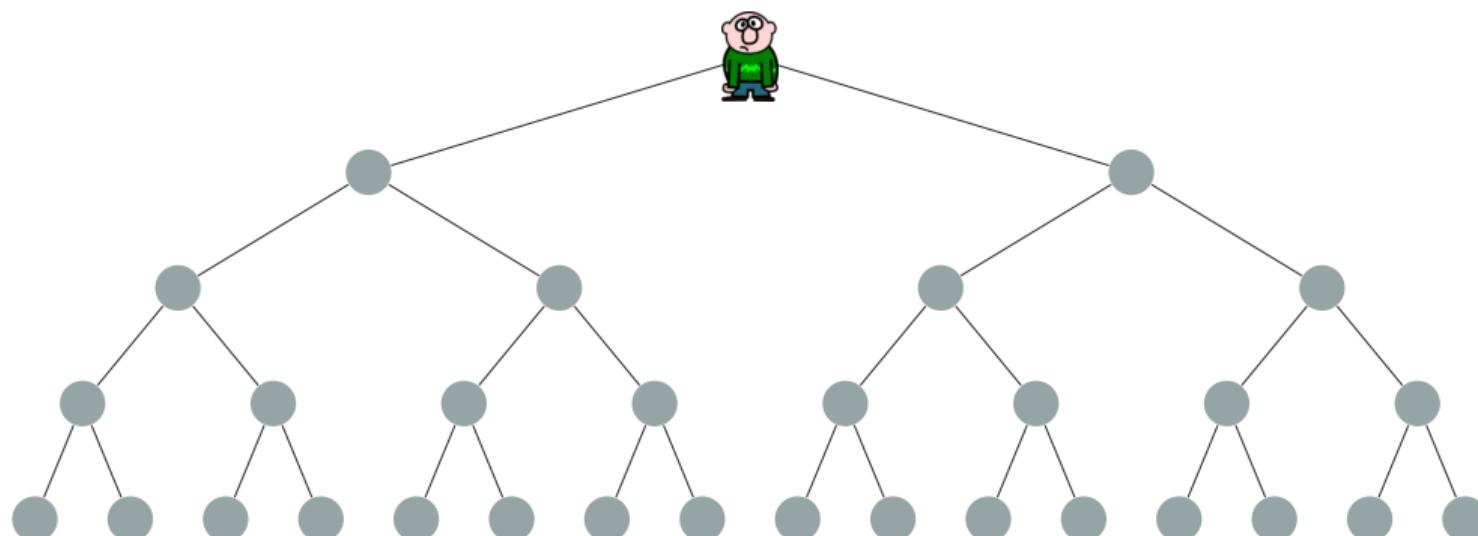
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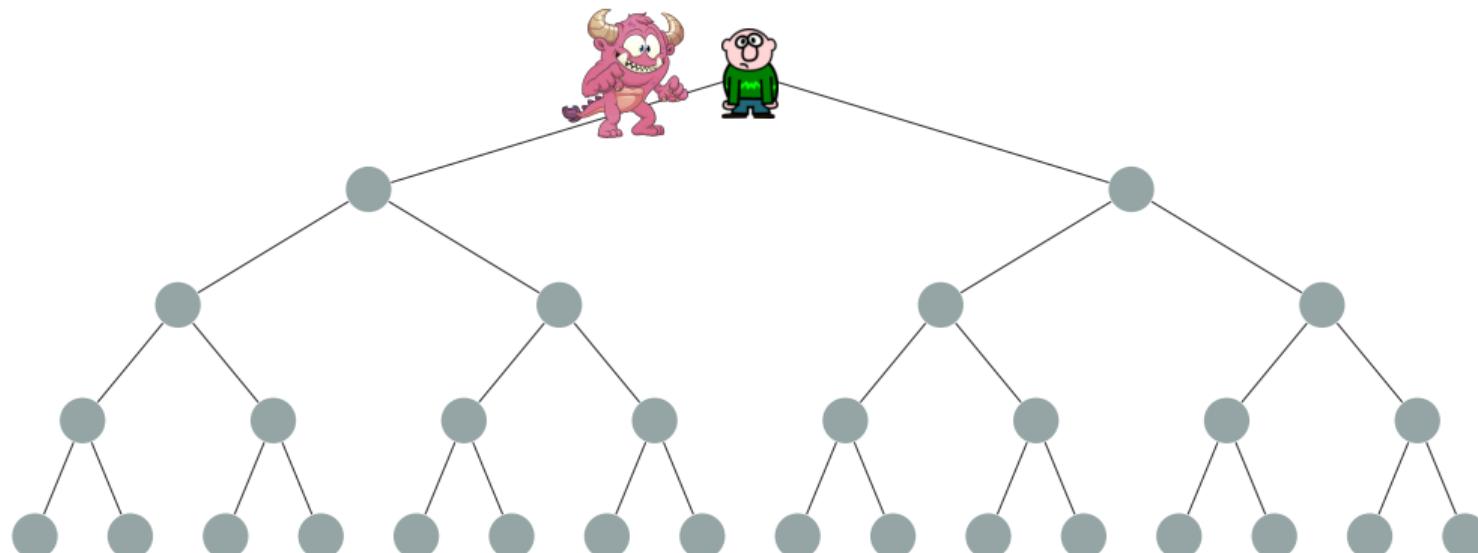
Random Walk on the Infinite Binary Tree

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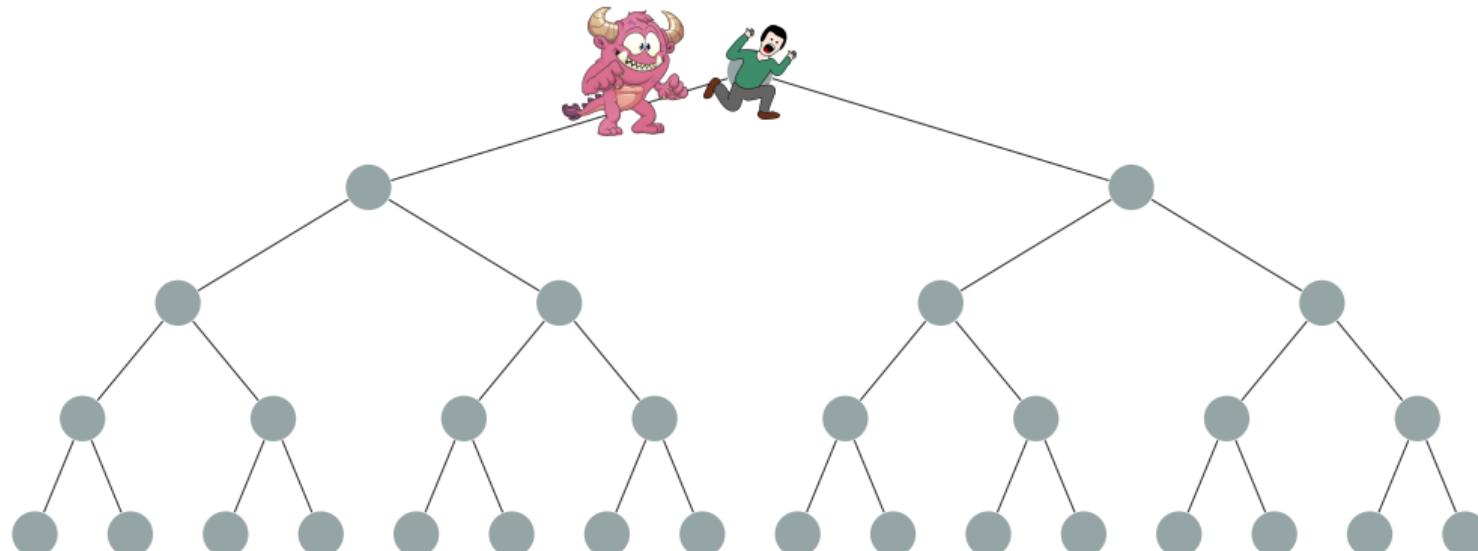
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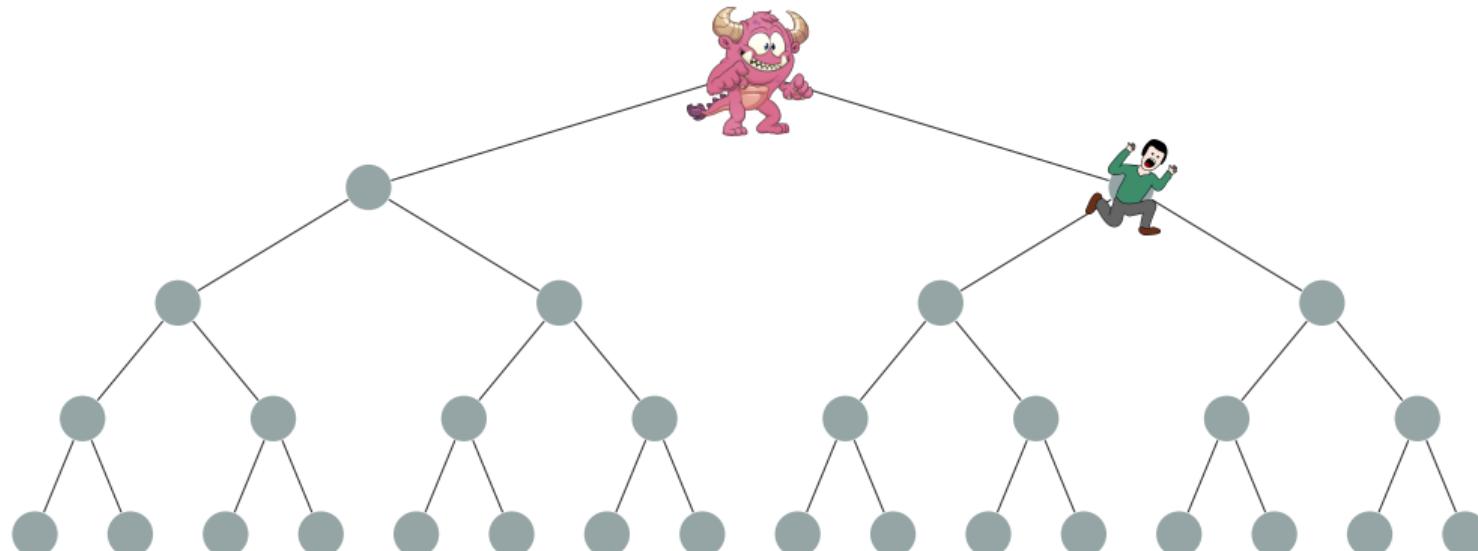
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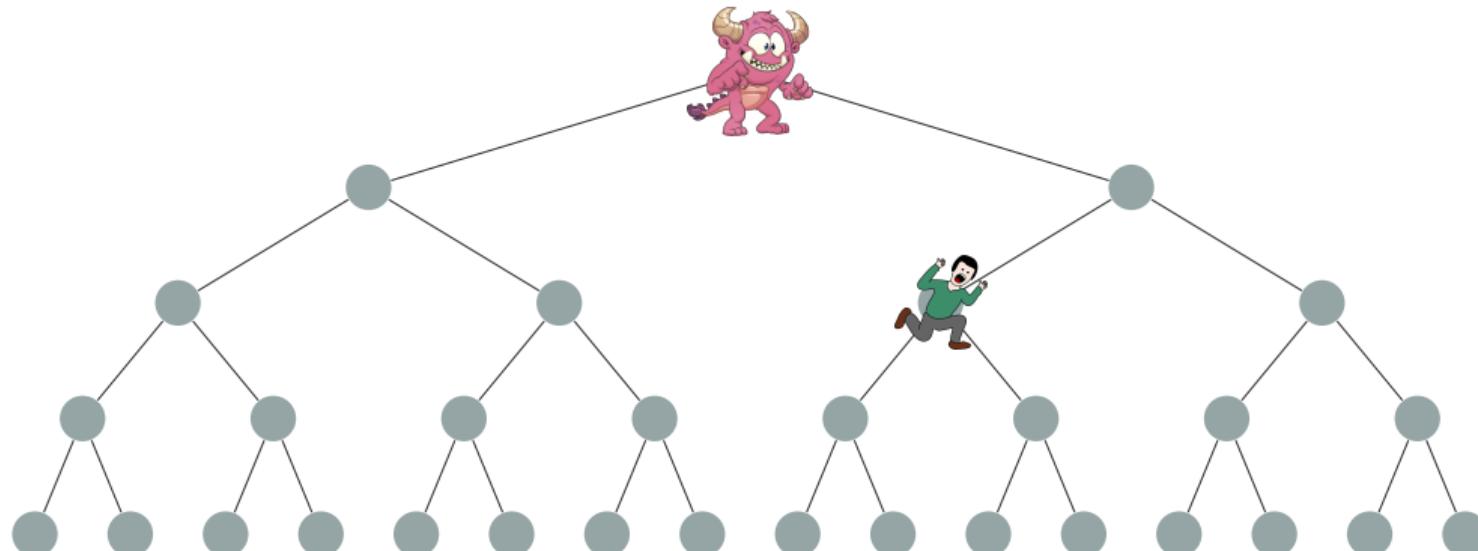
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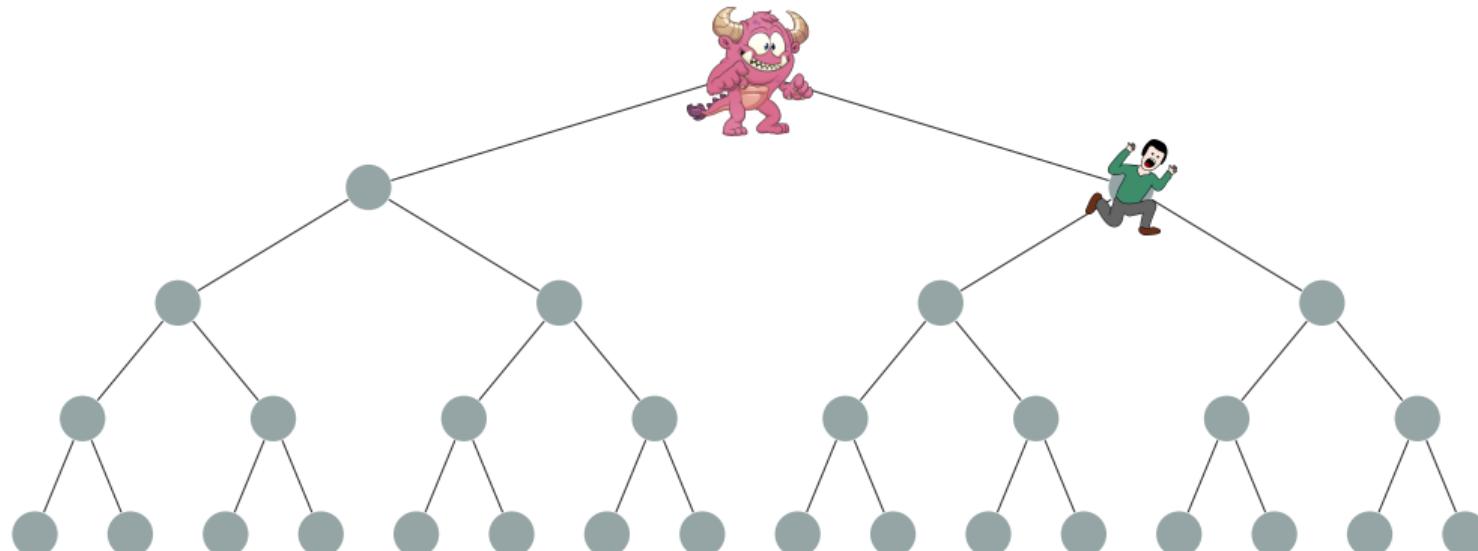
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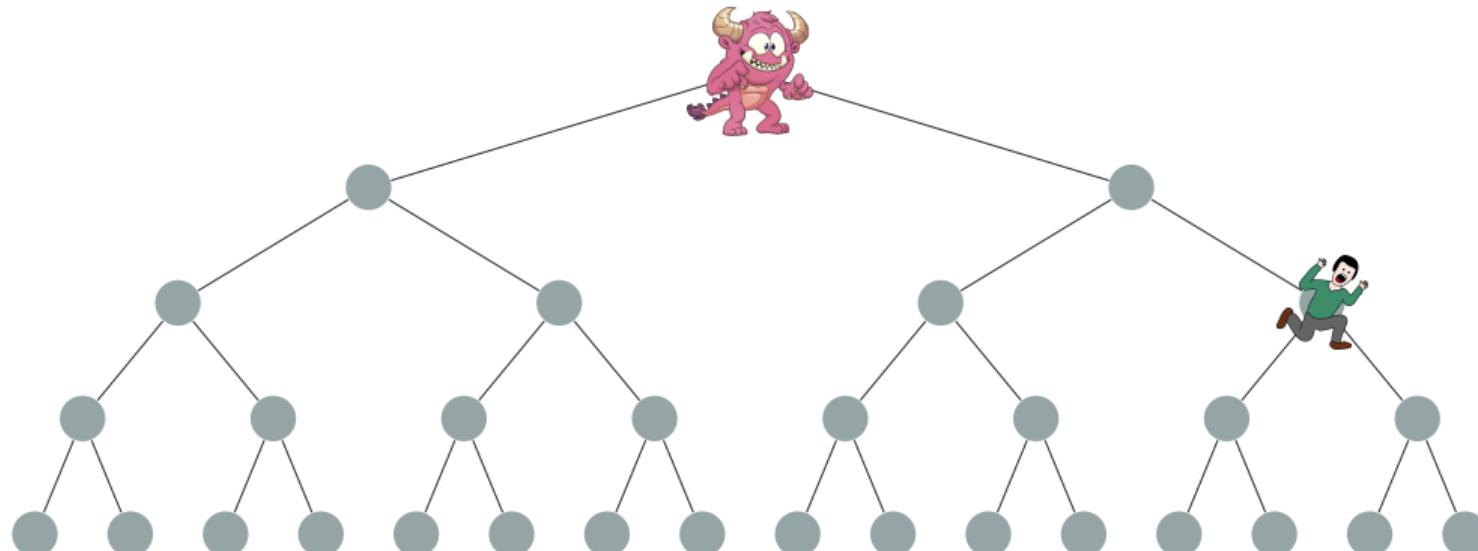
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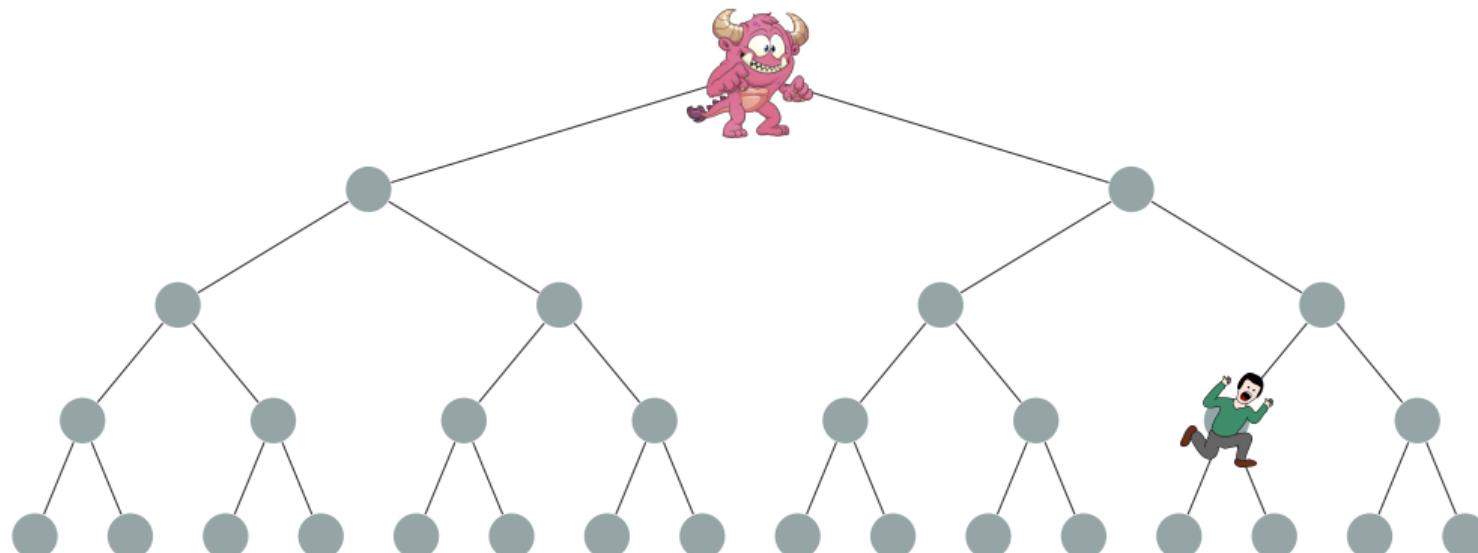
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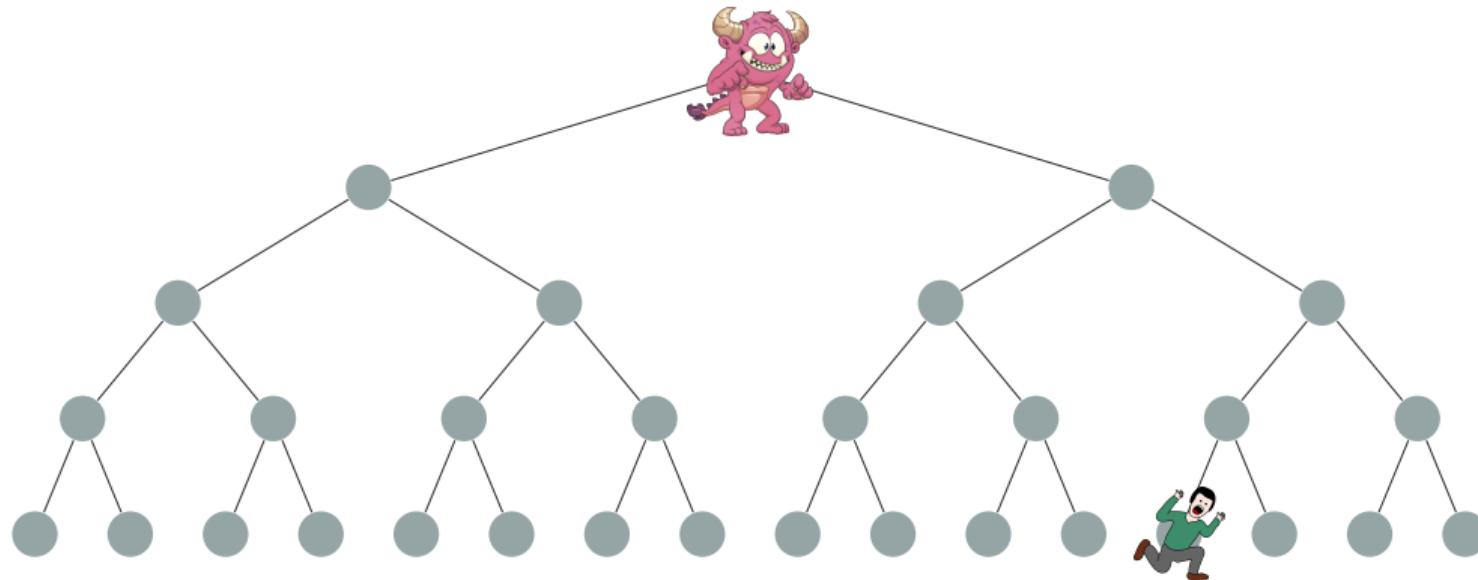
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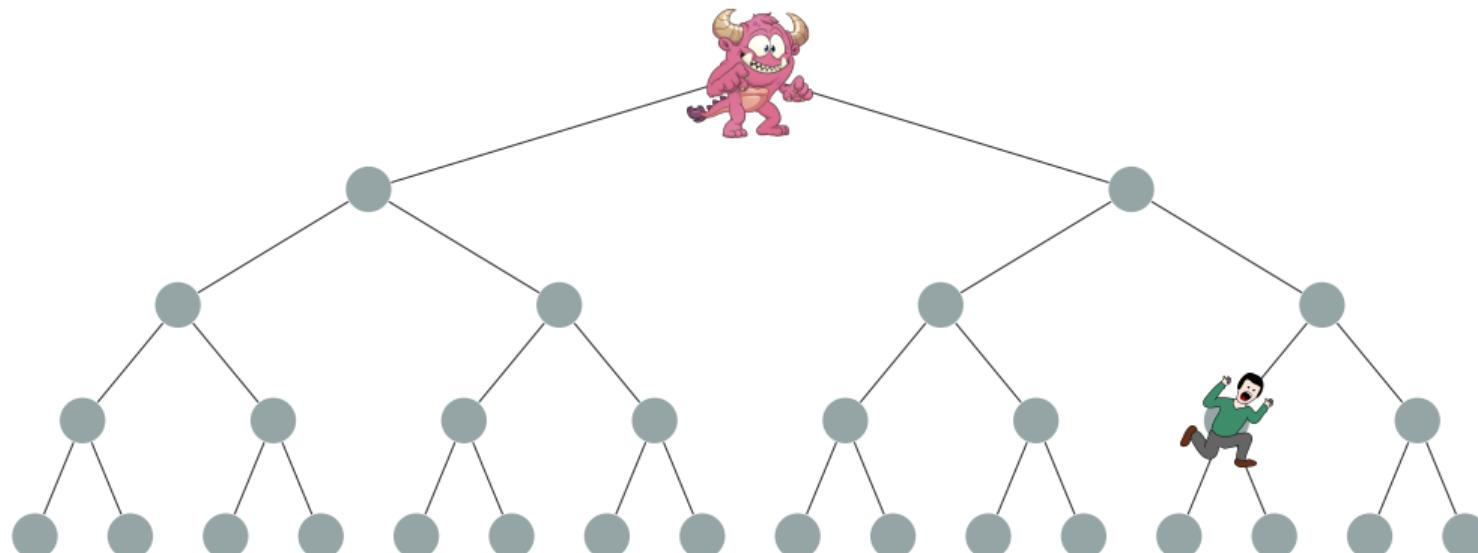
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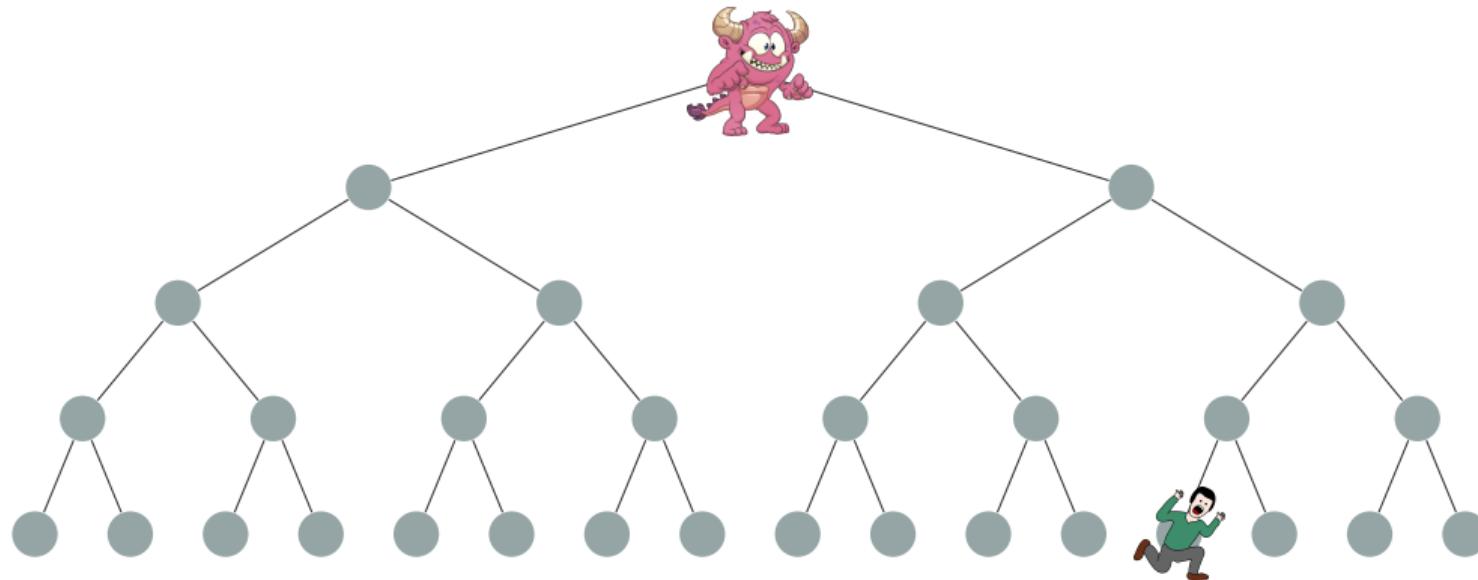
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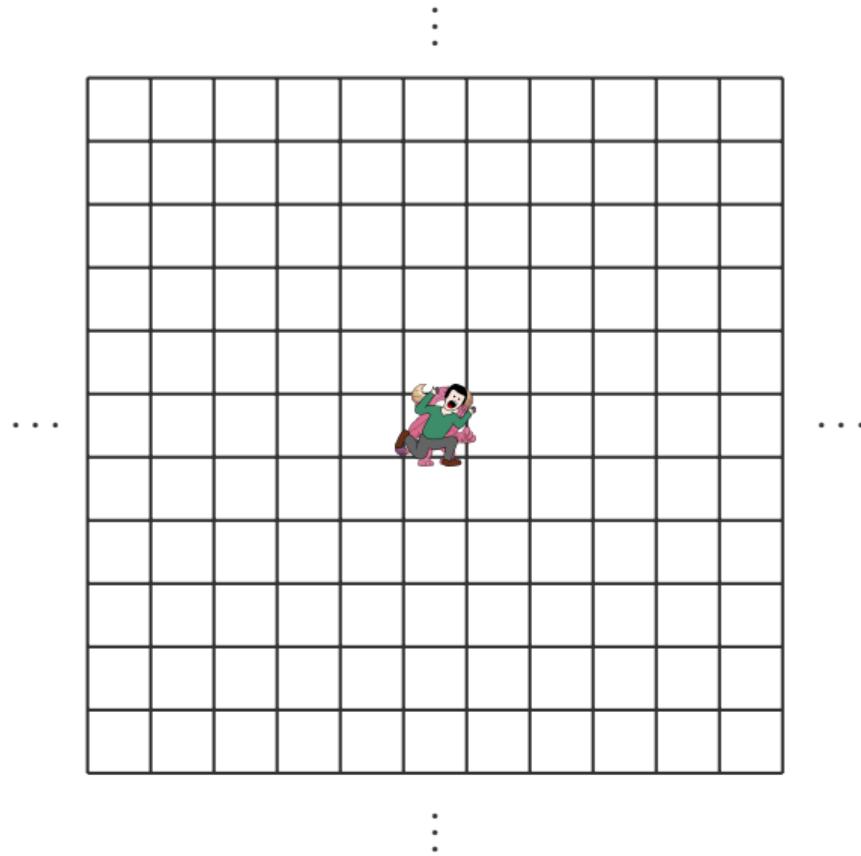
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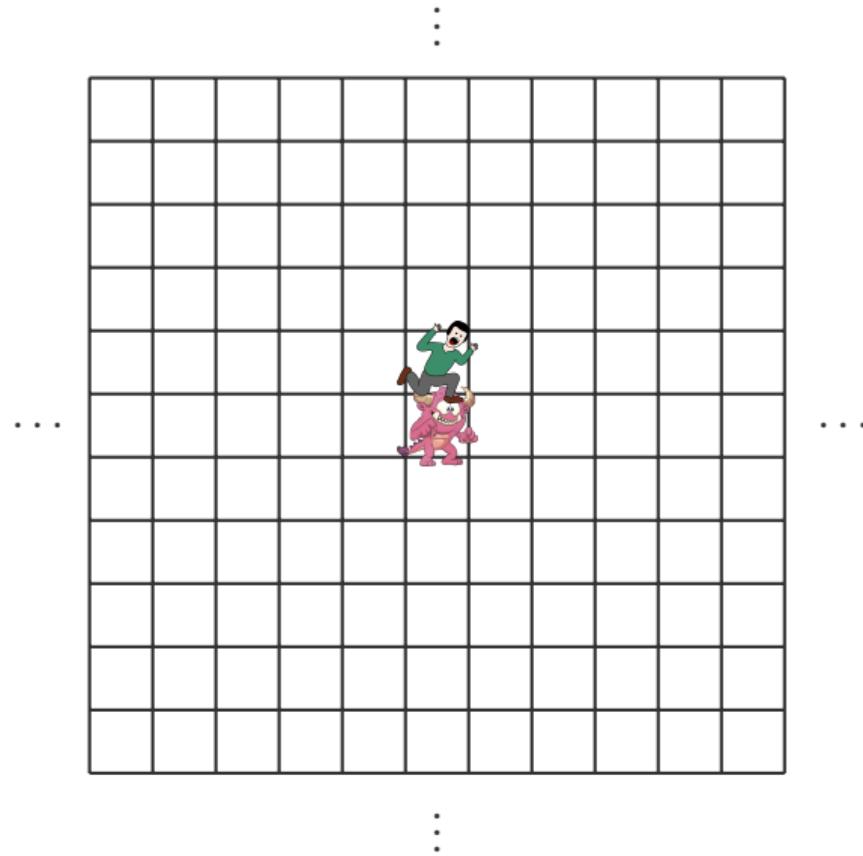
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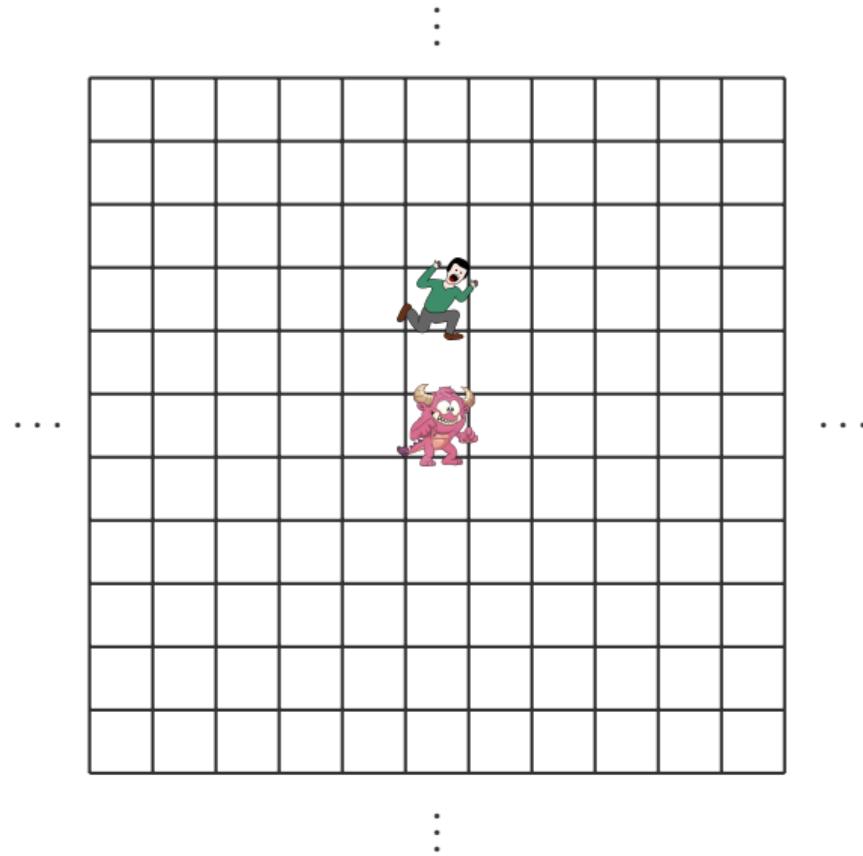
Random Walk on \mathbb{Z}^2



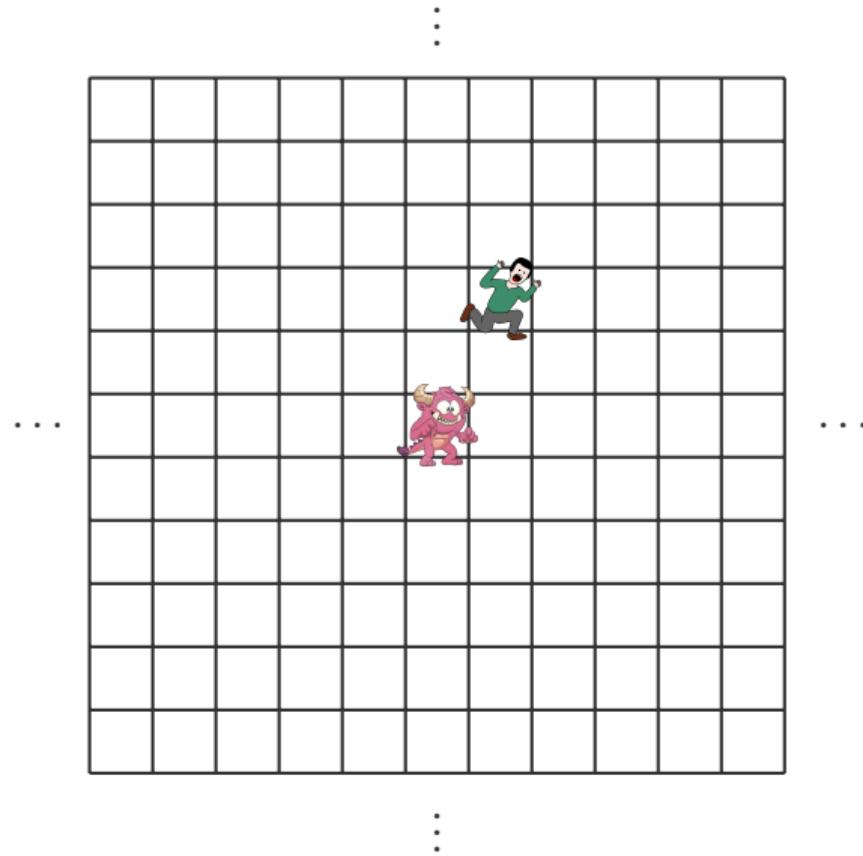
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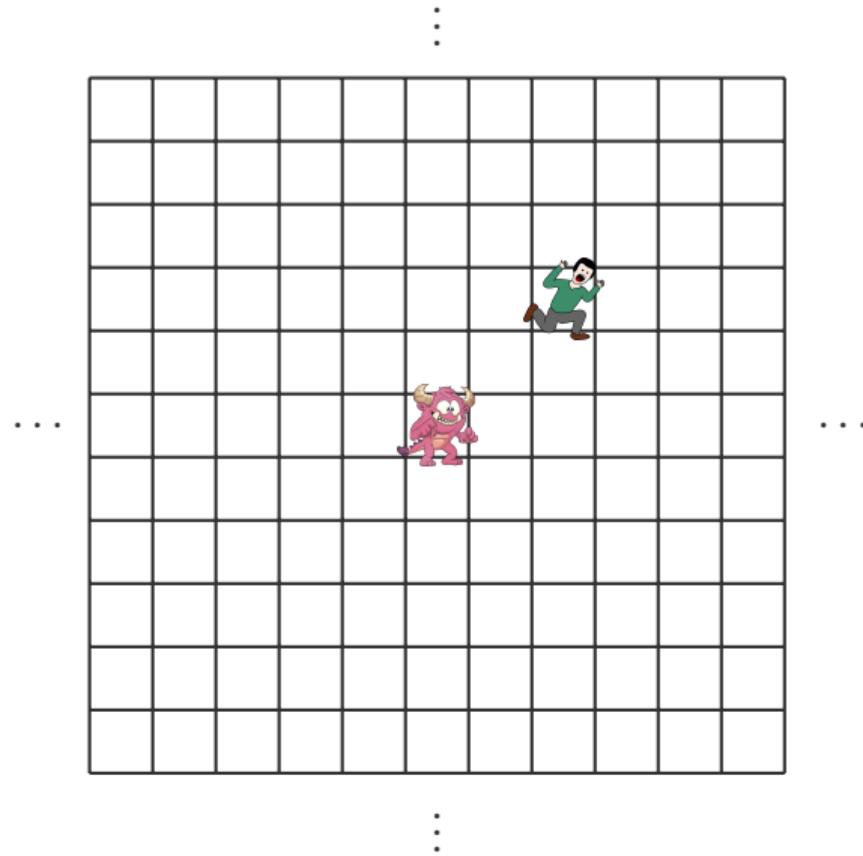
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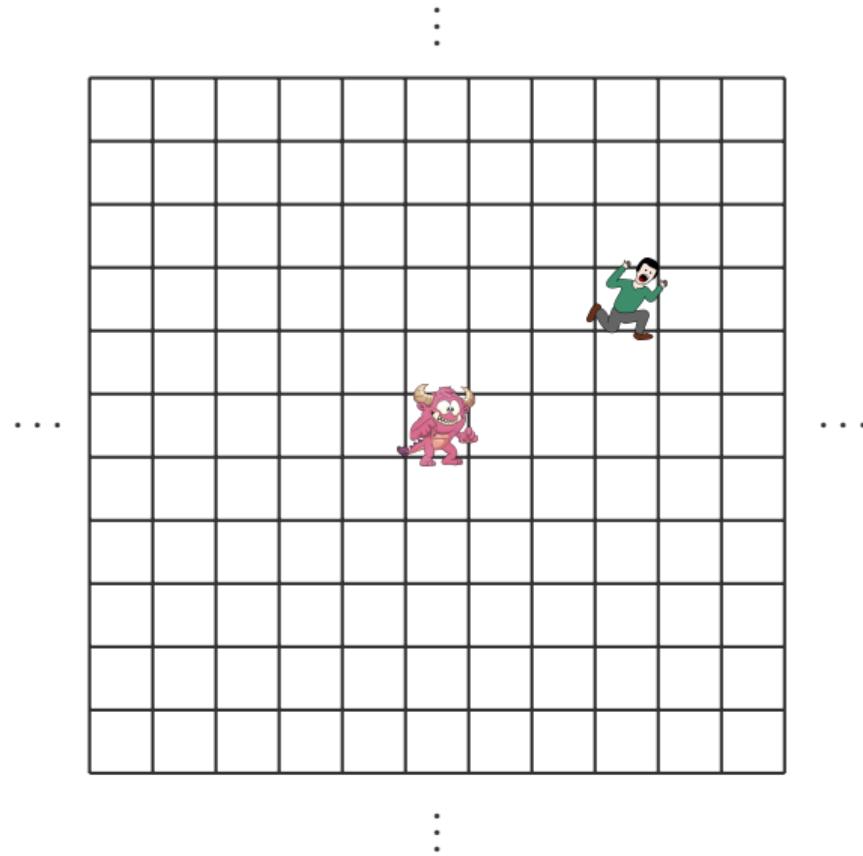
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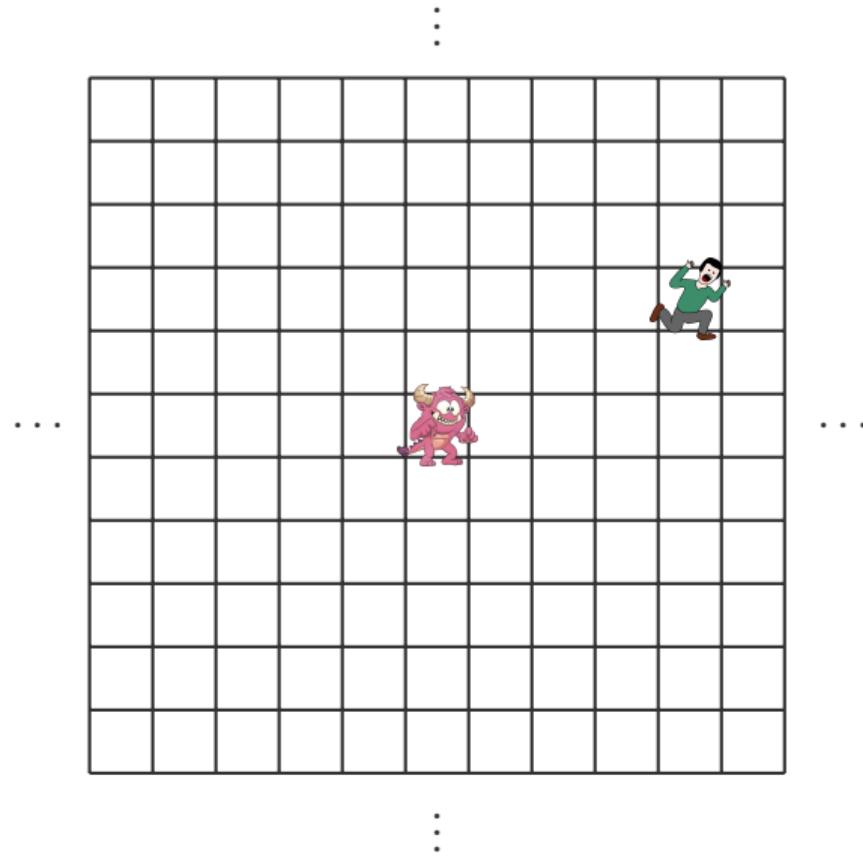
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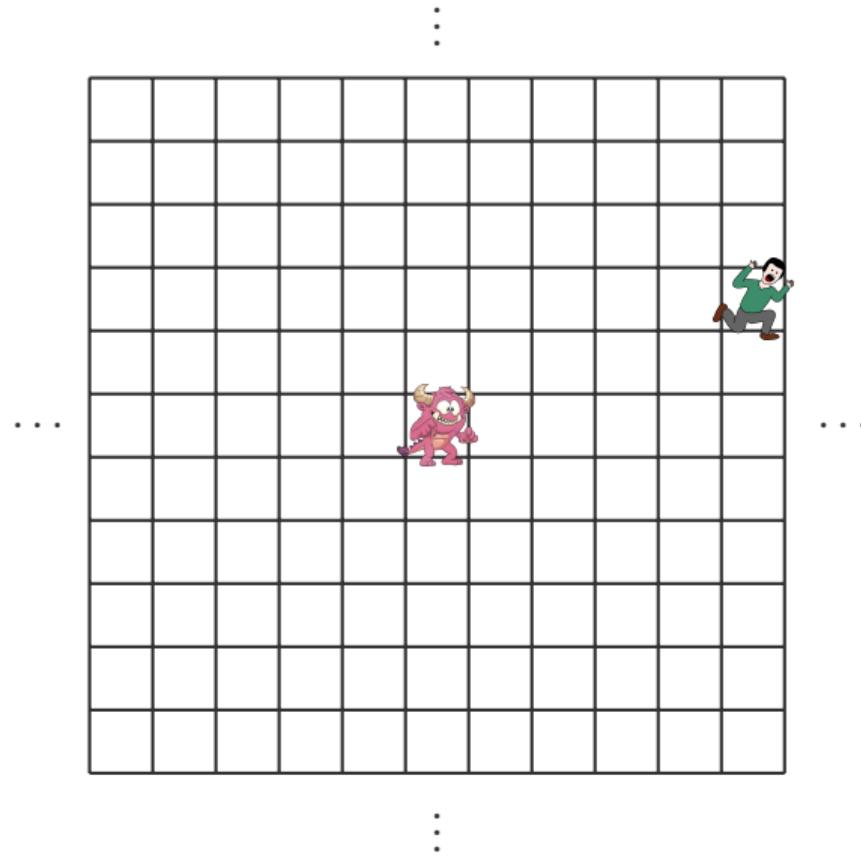
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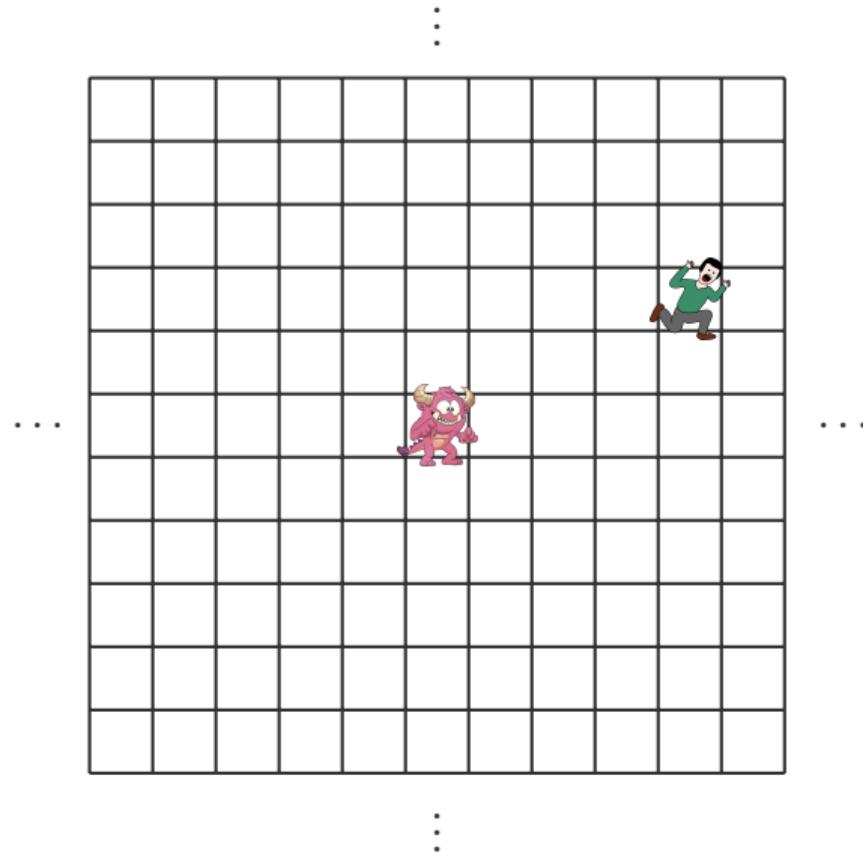
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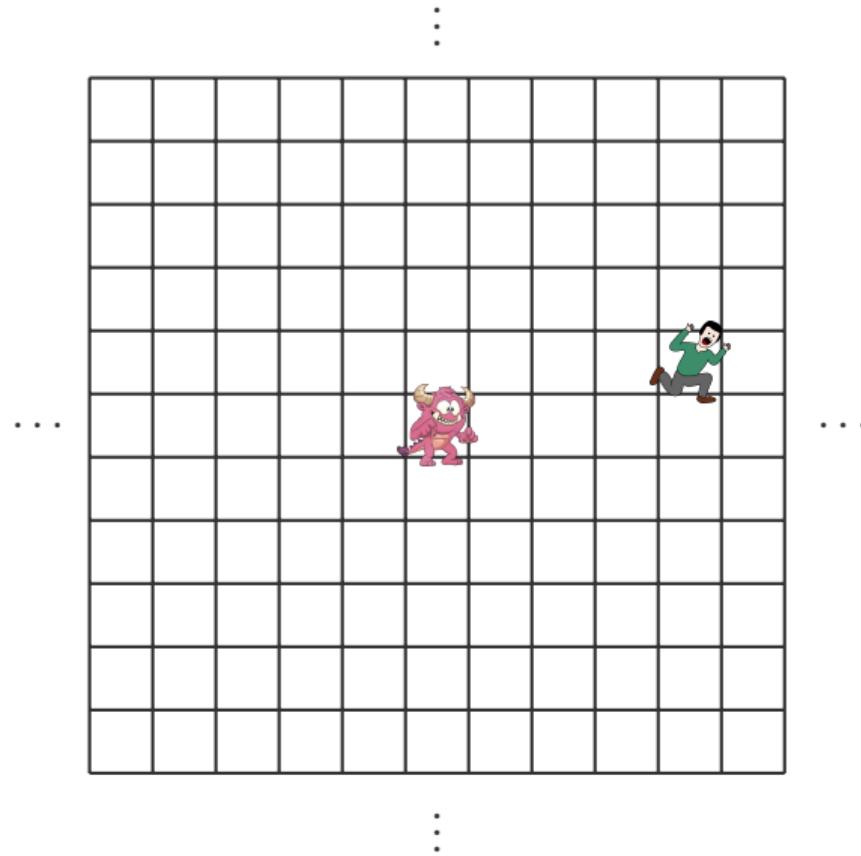
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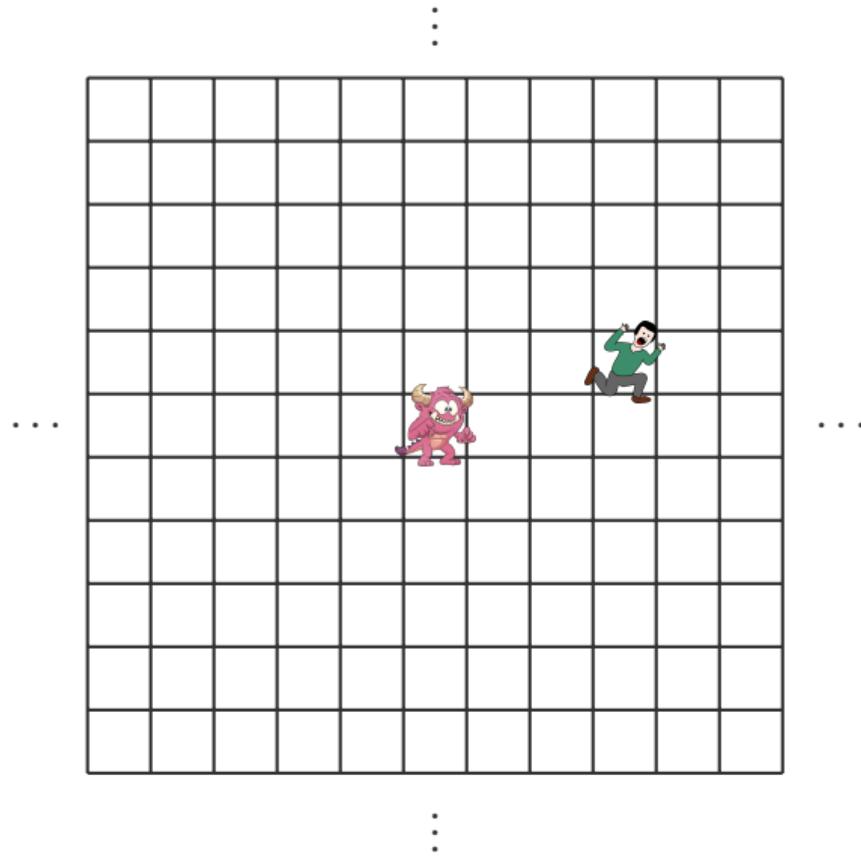
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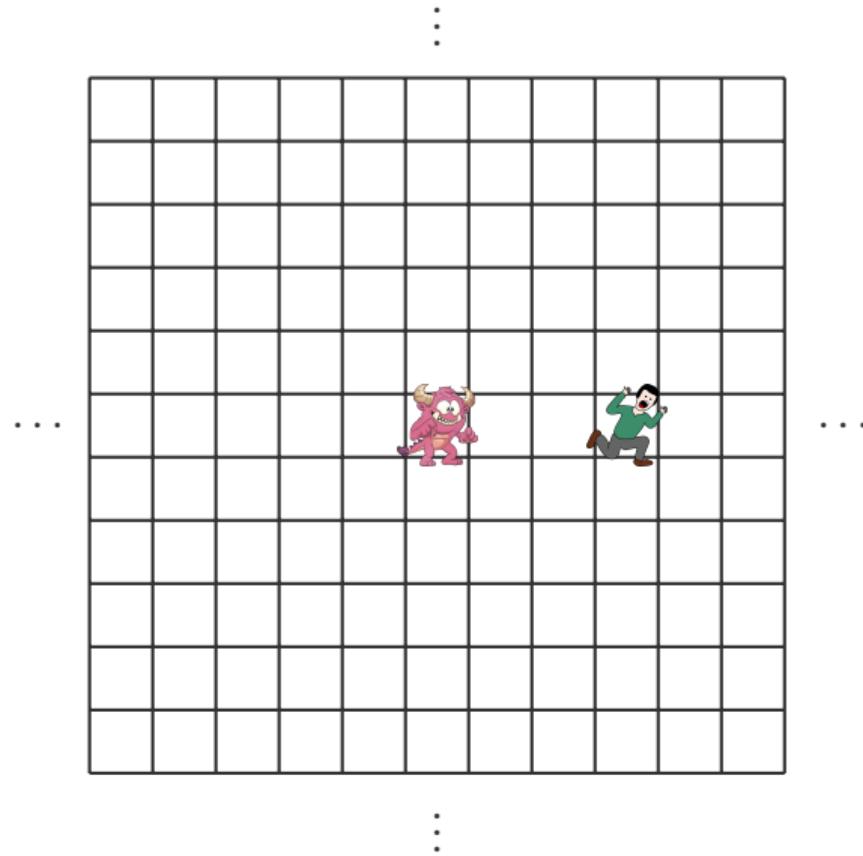
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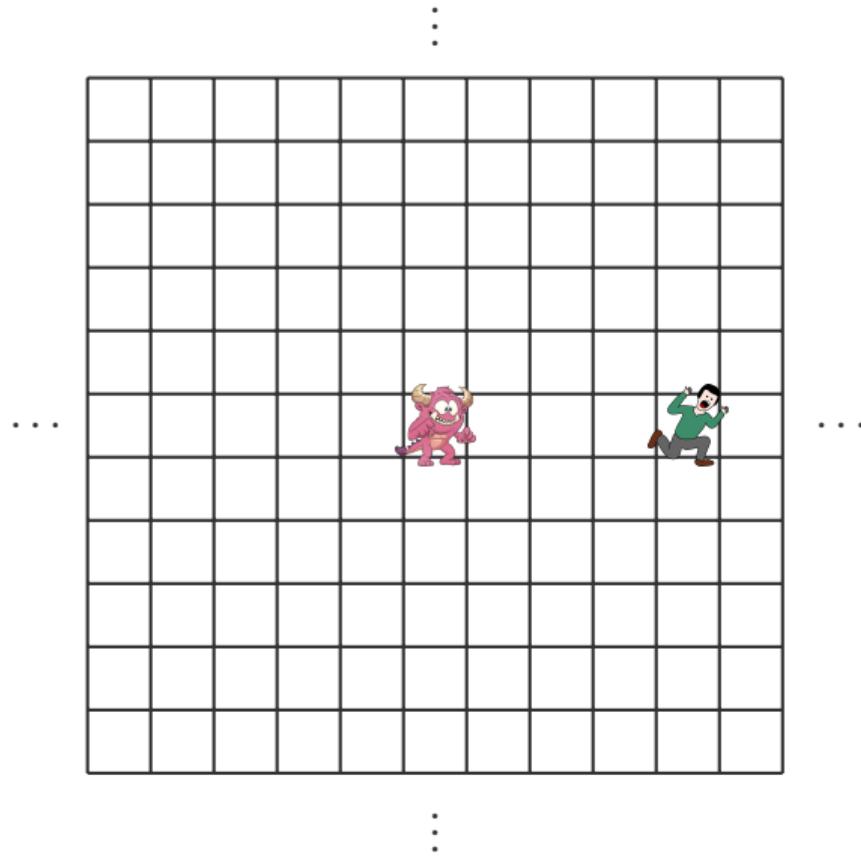
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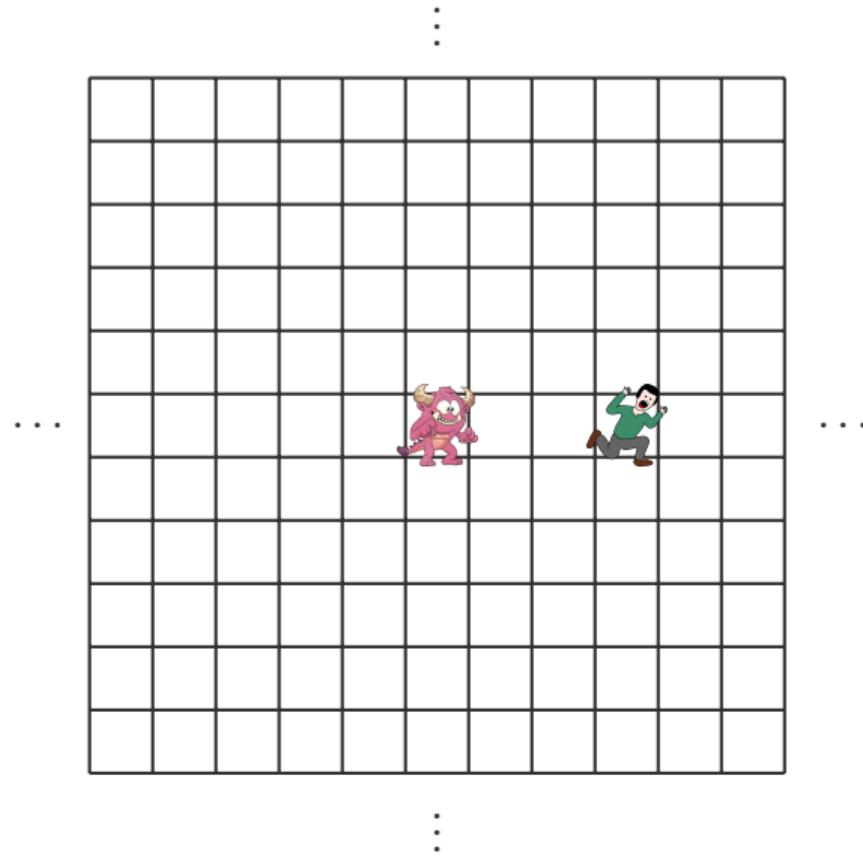
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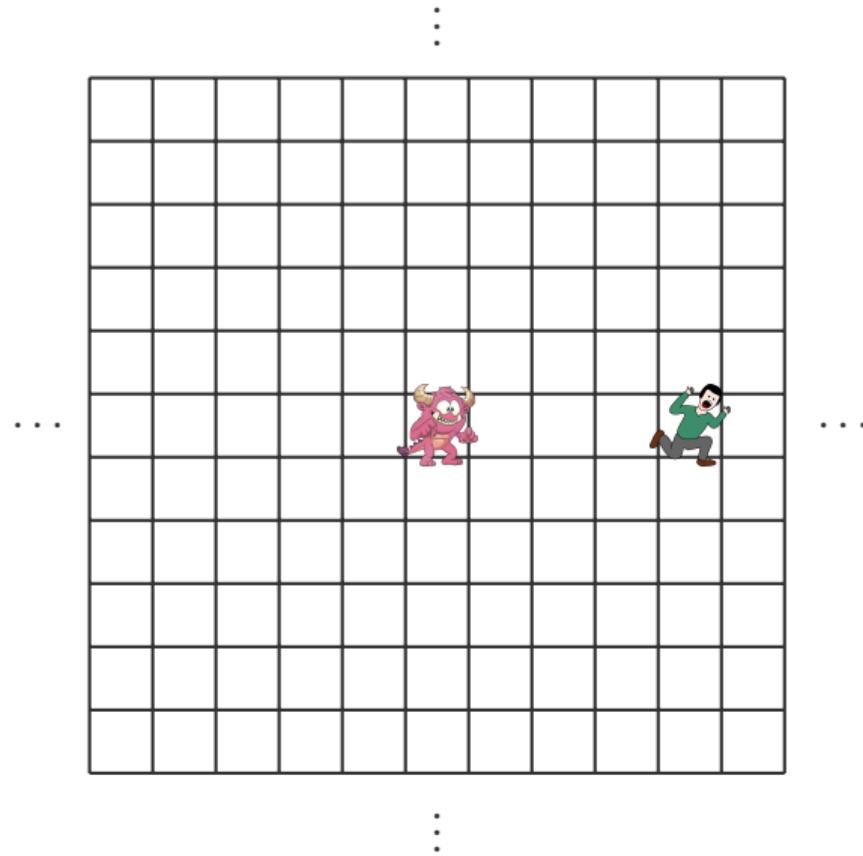
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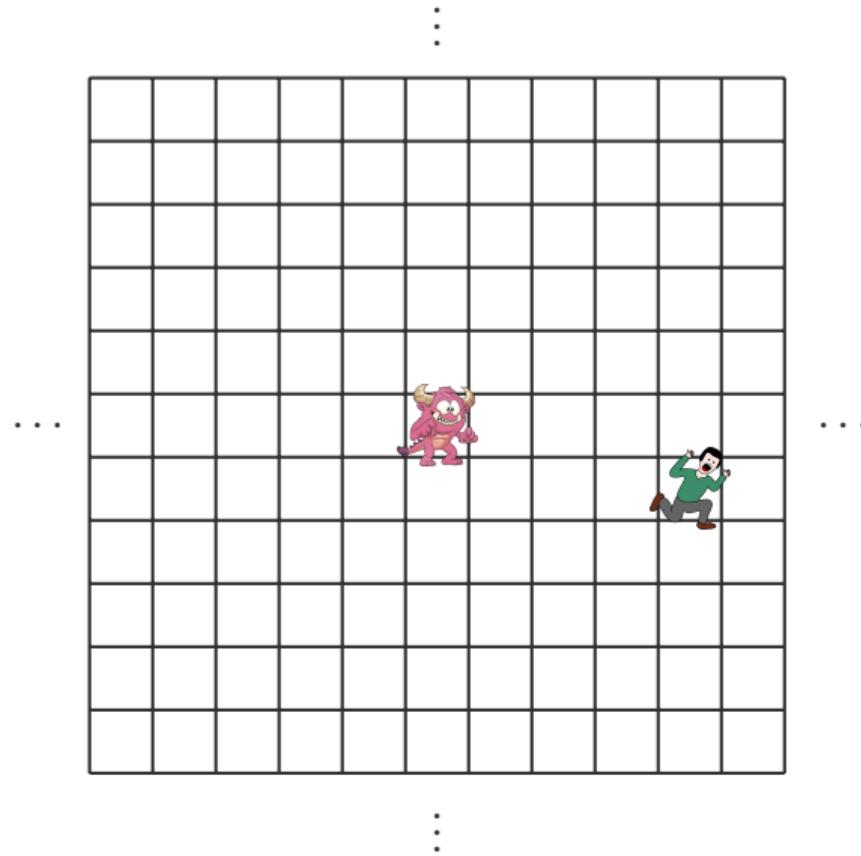
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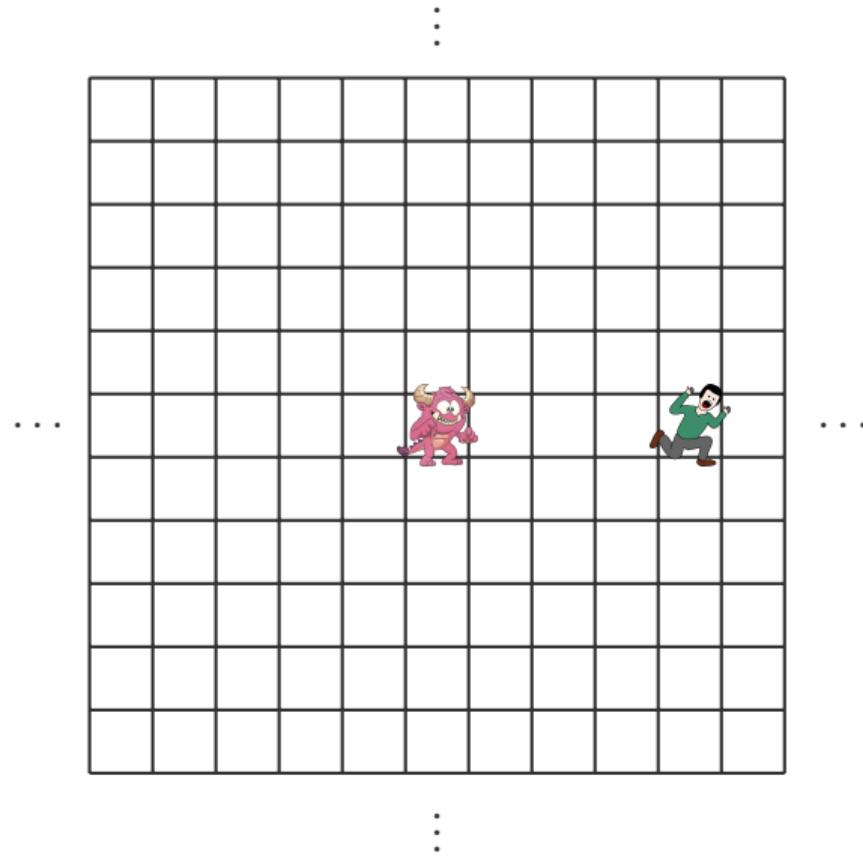
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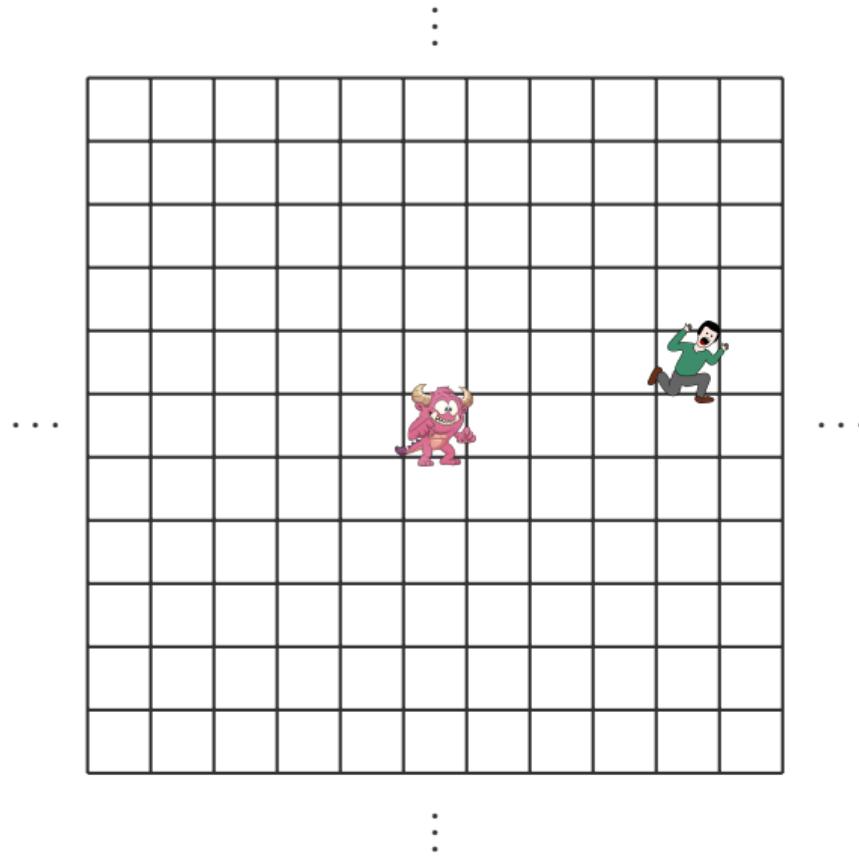
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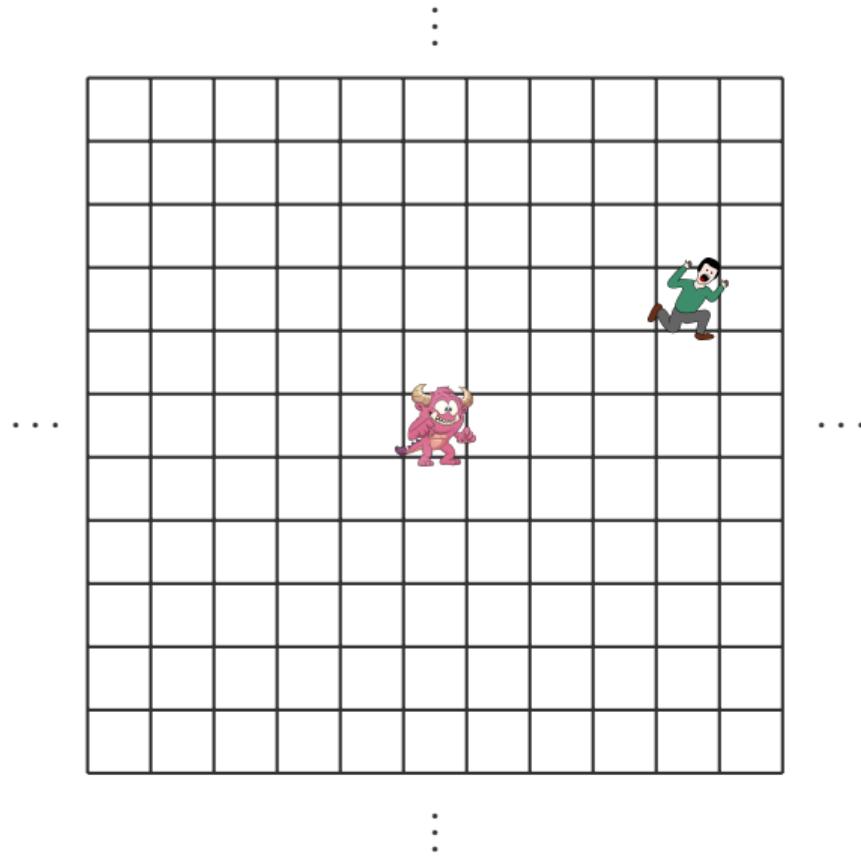
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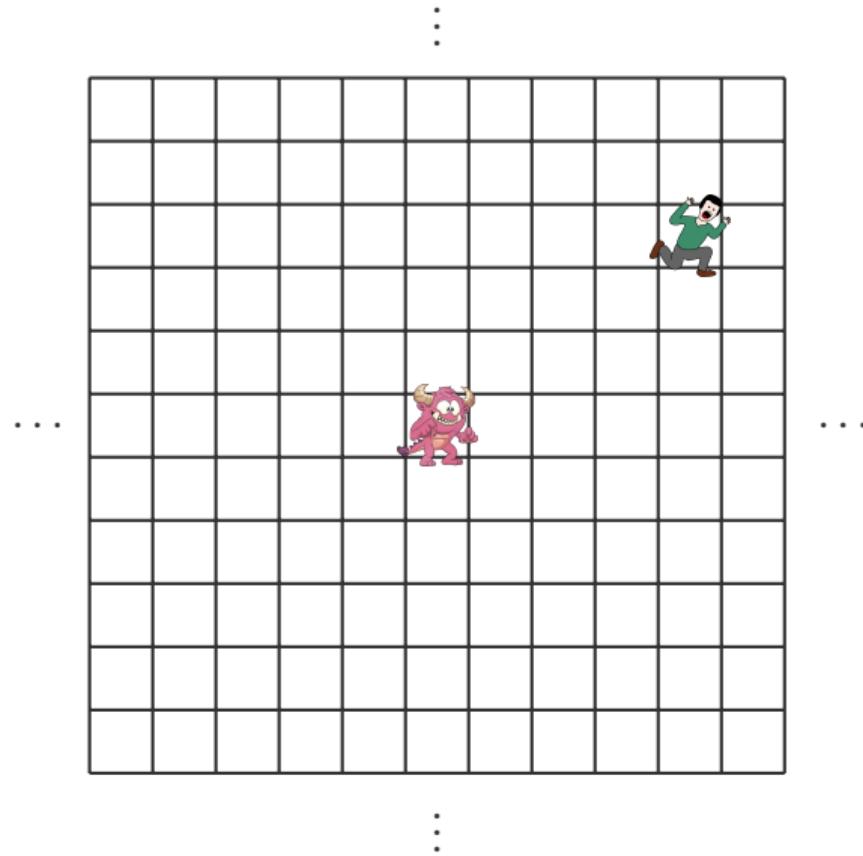
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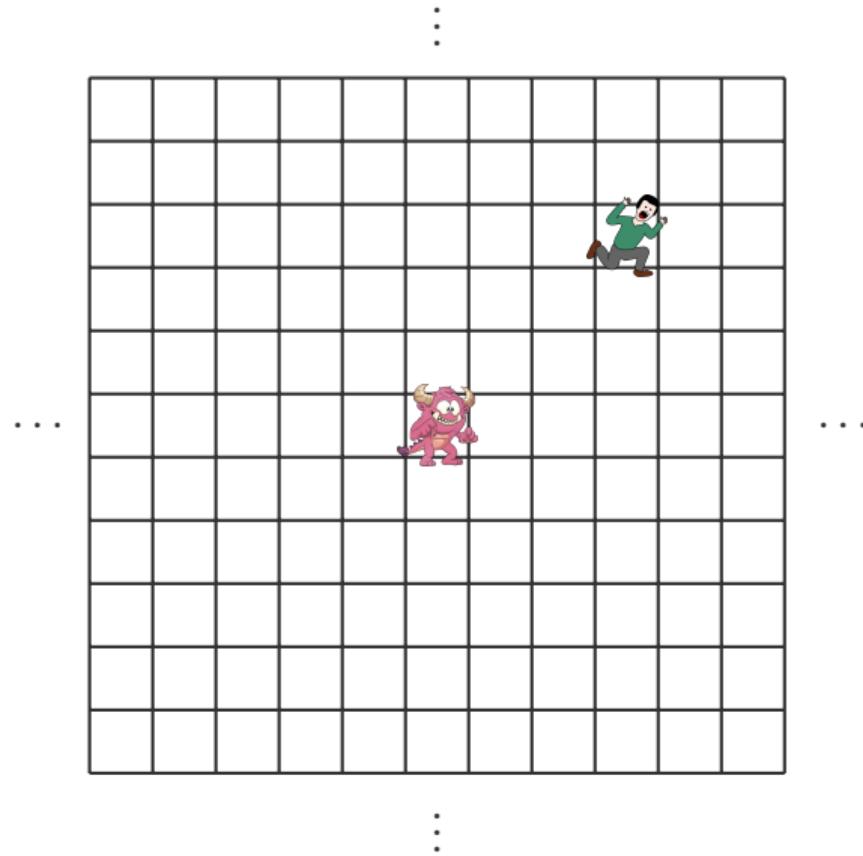
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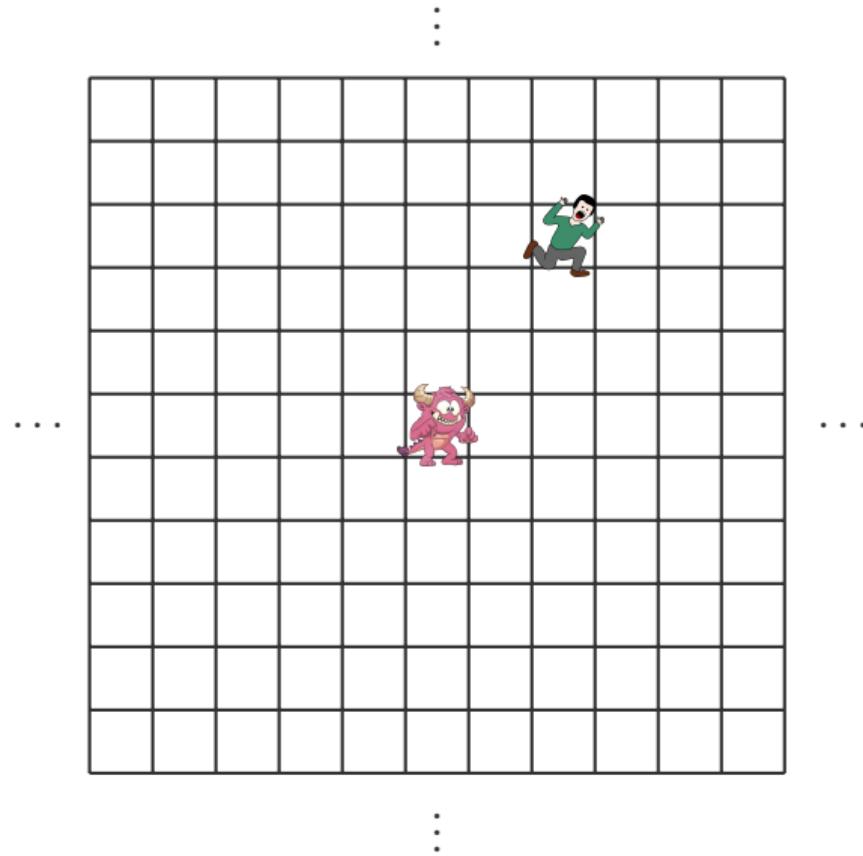
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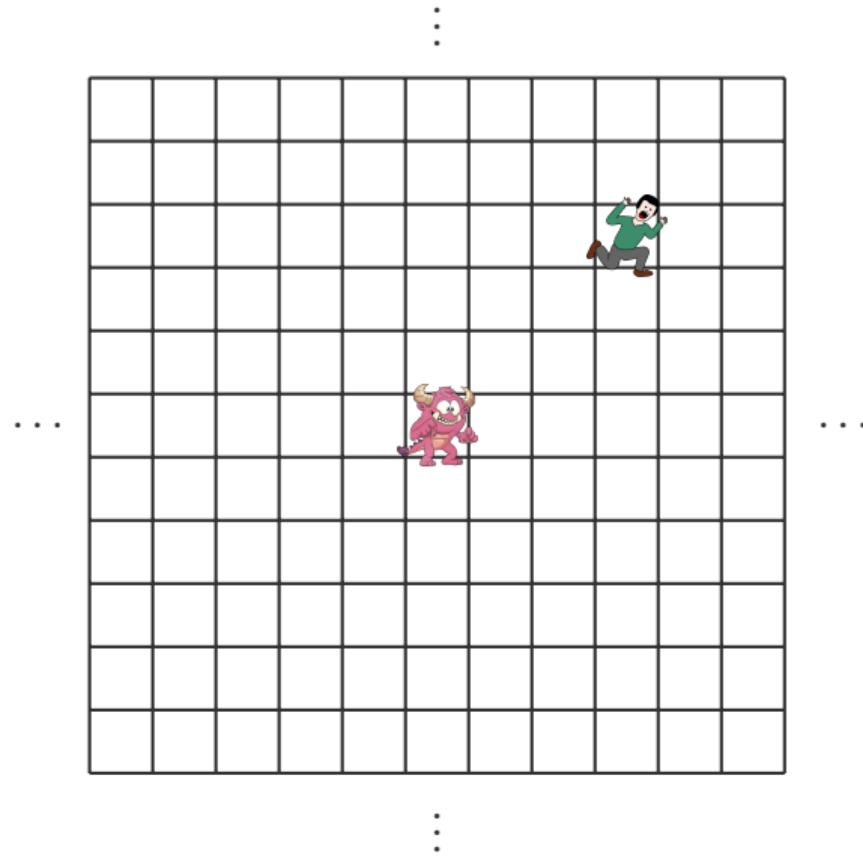
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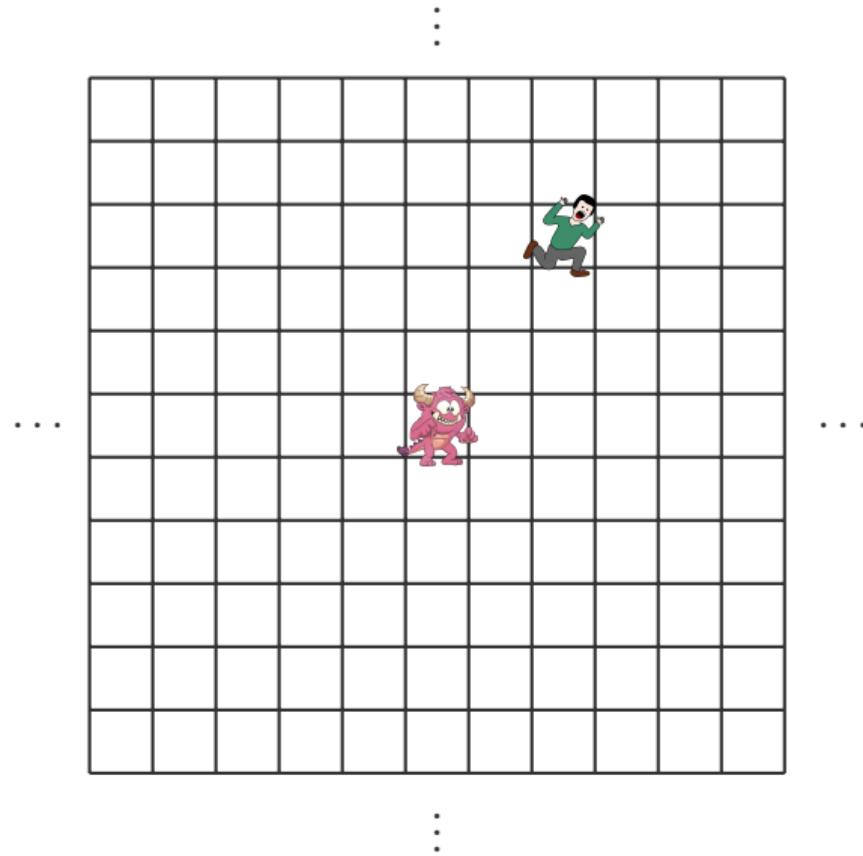
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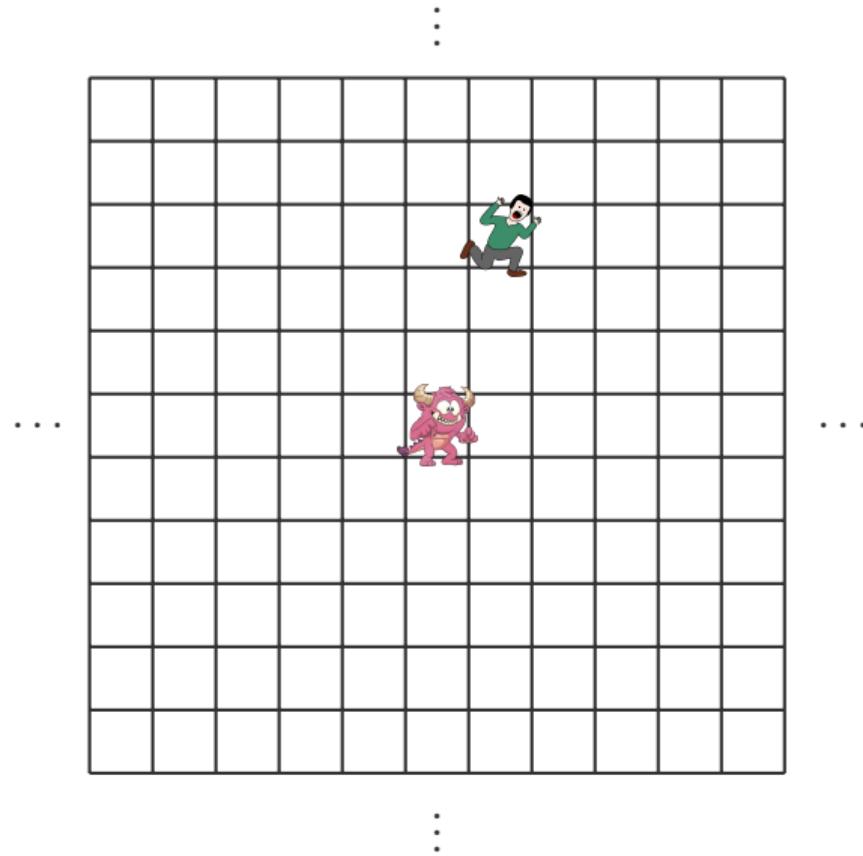
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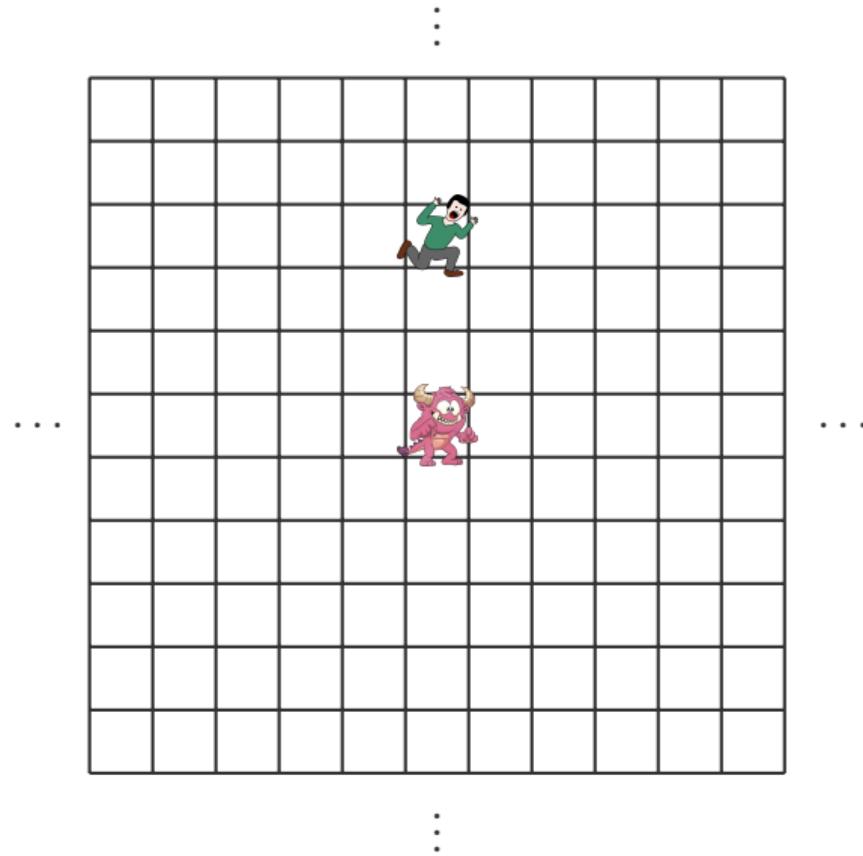
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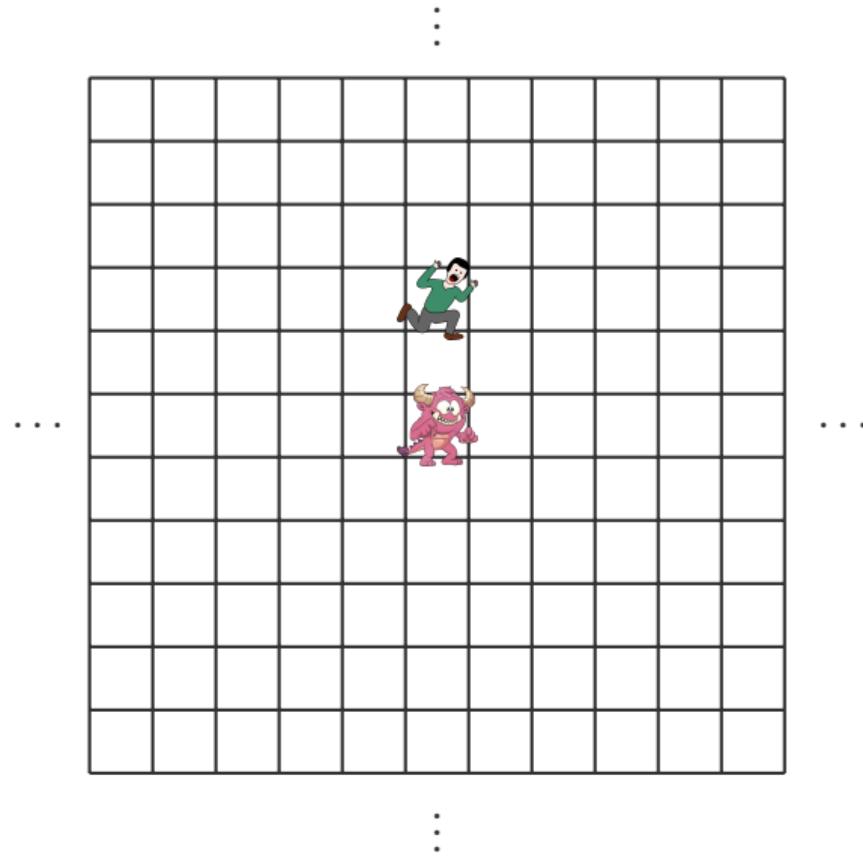
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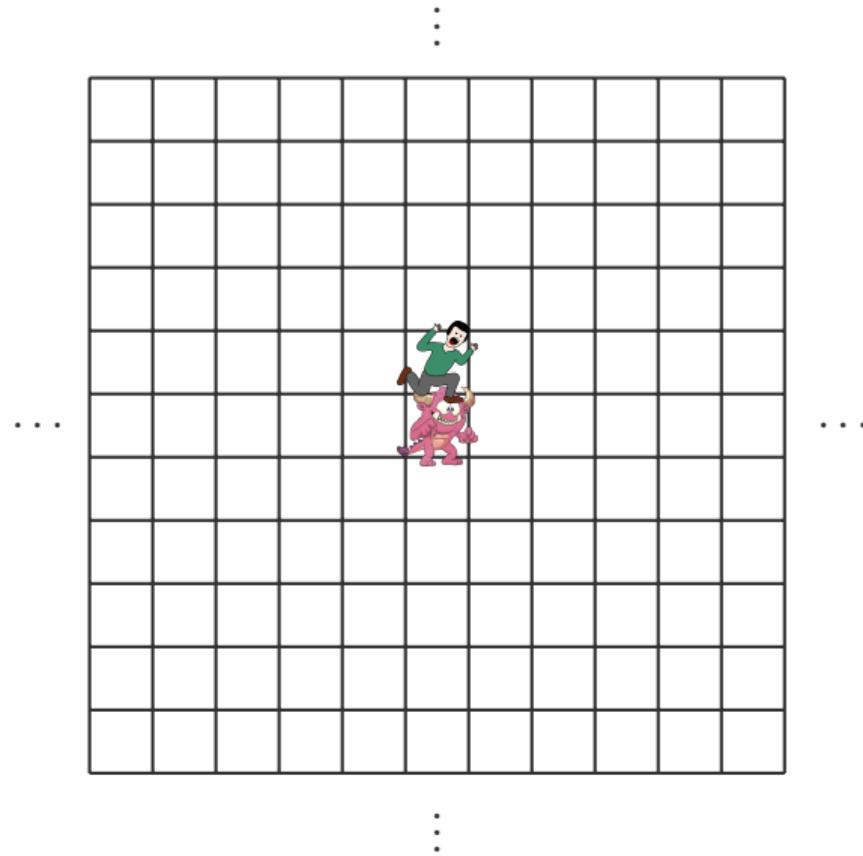
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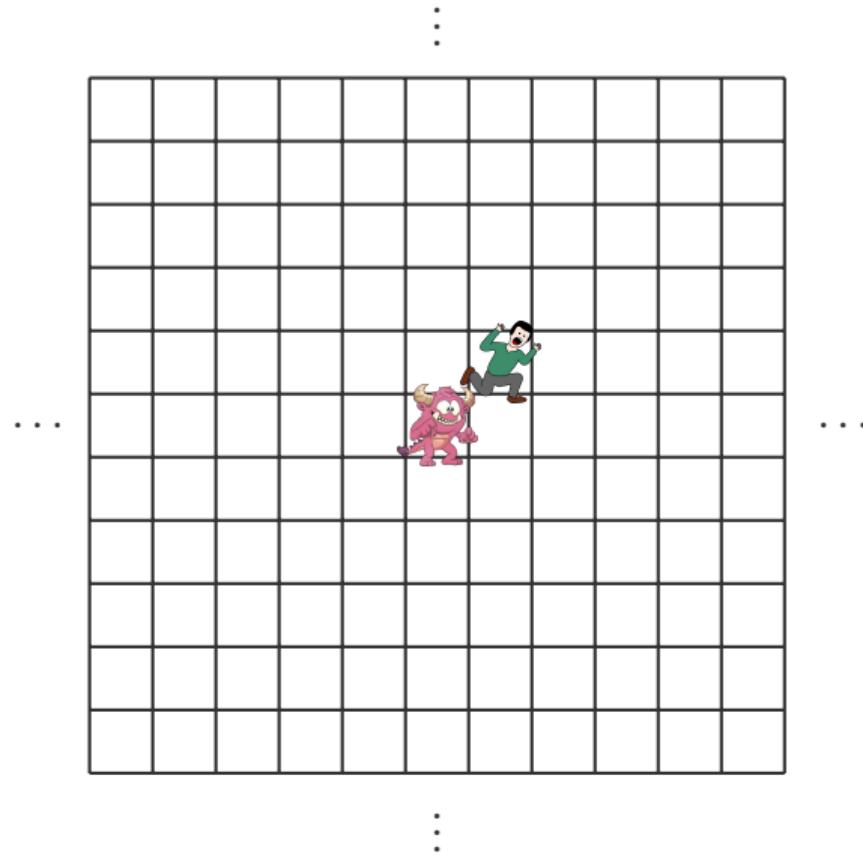
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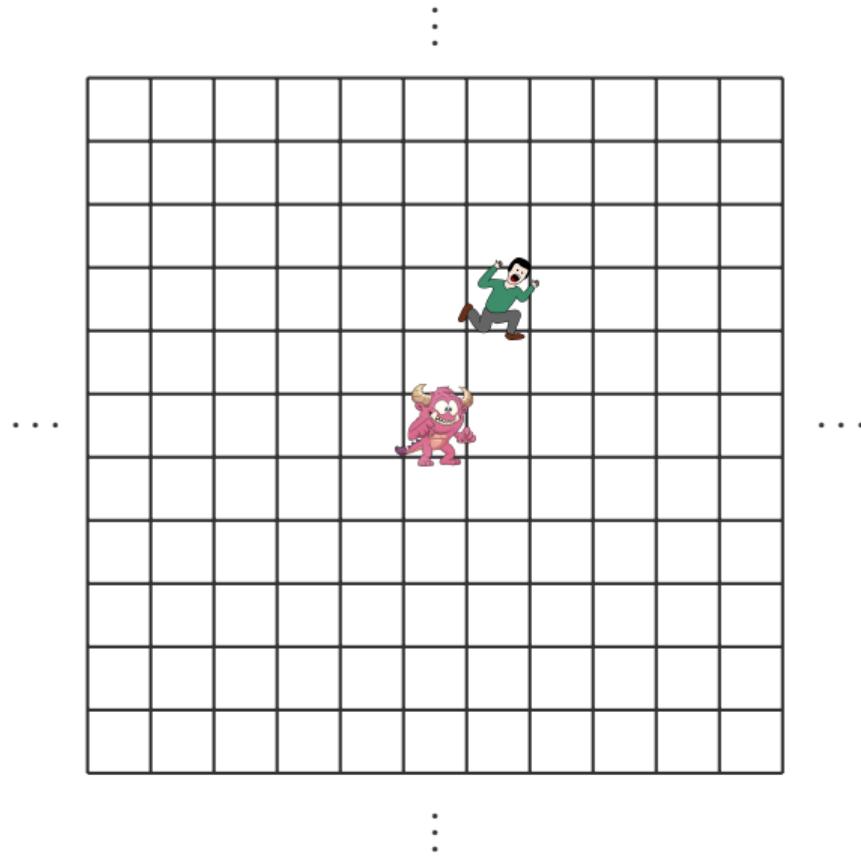
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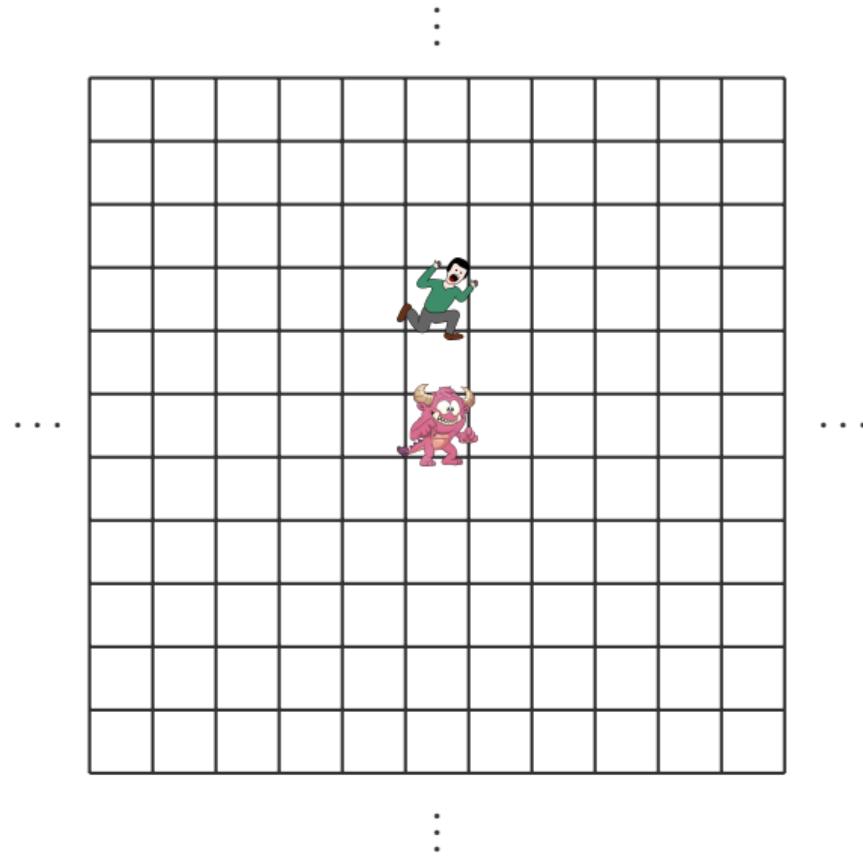
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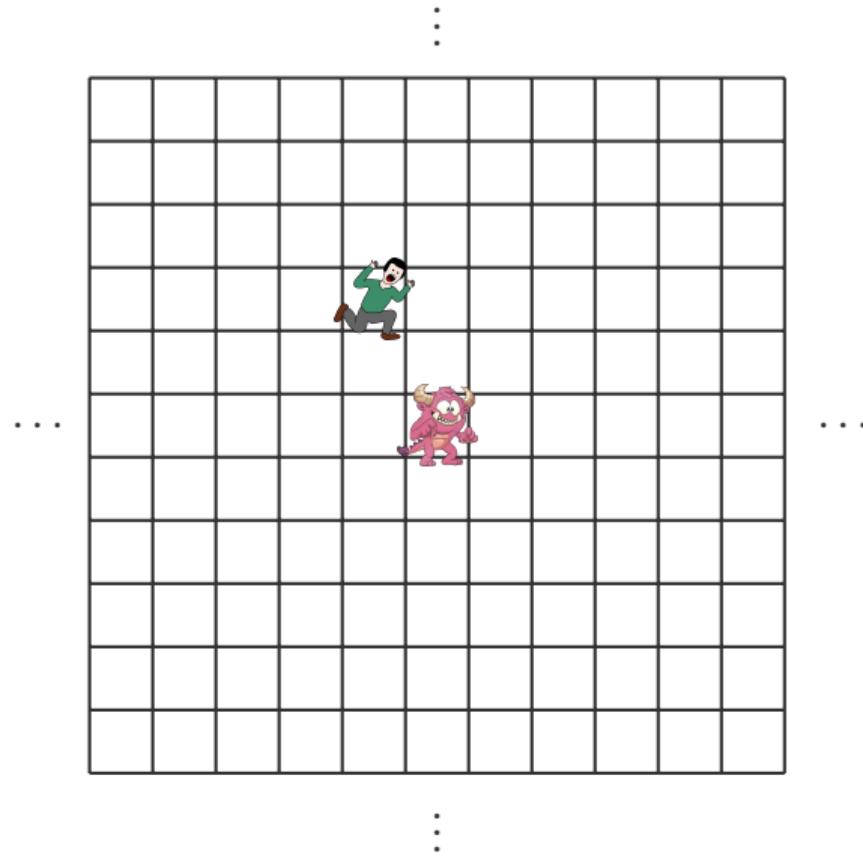
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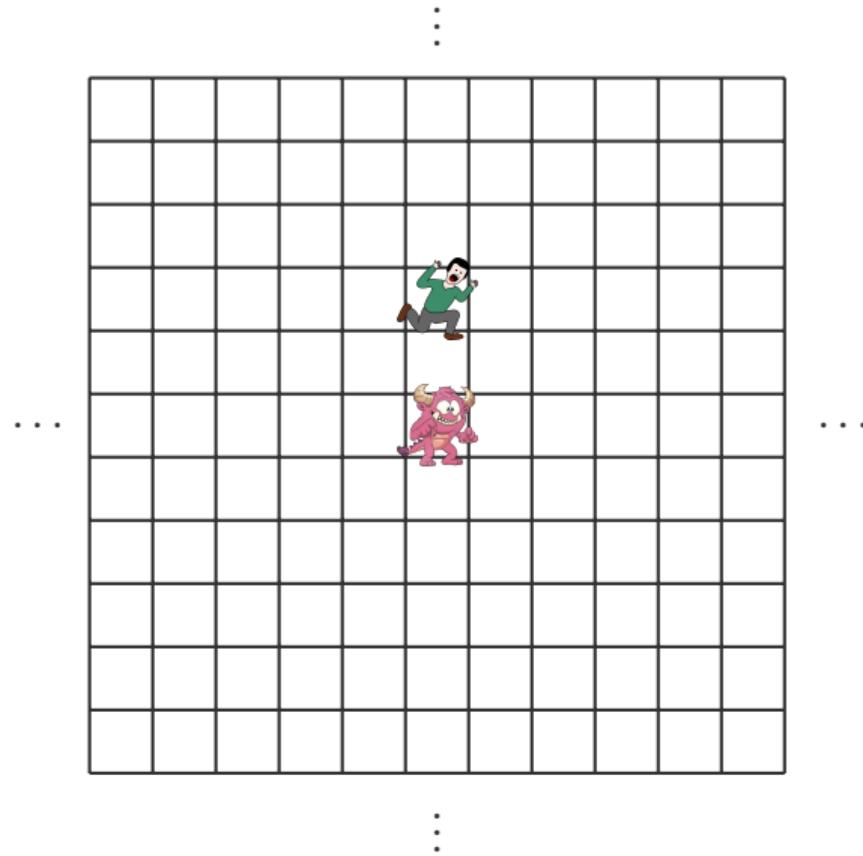
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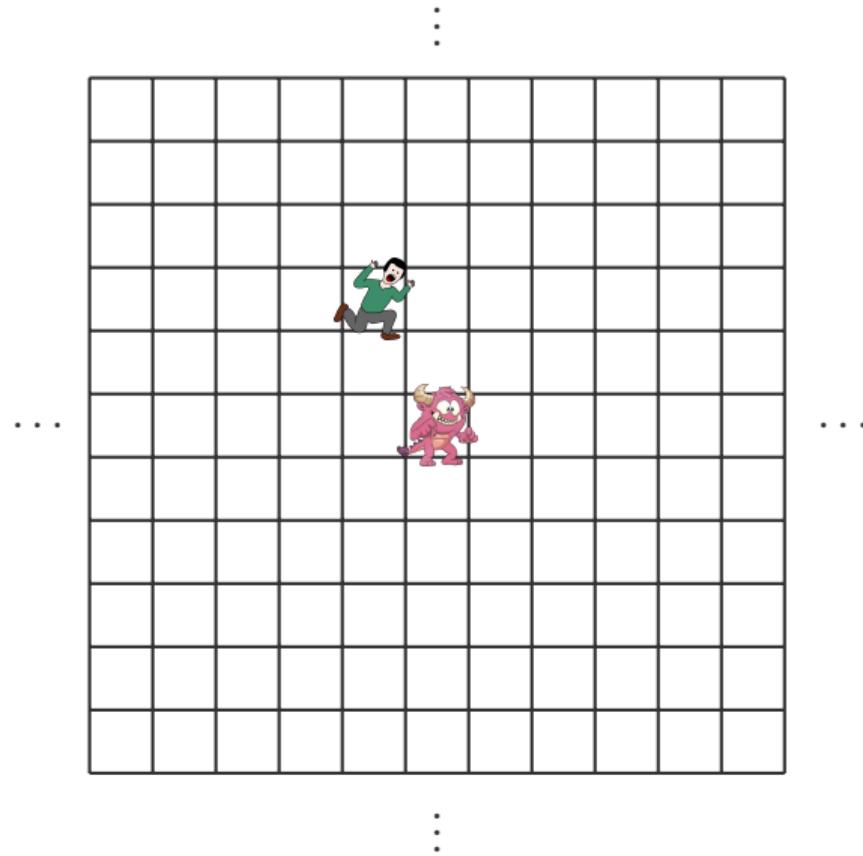
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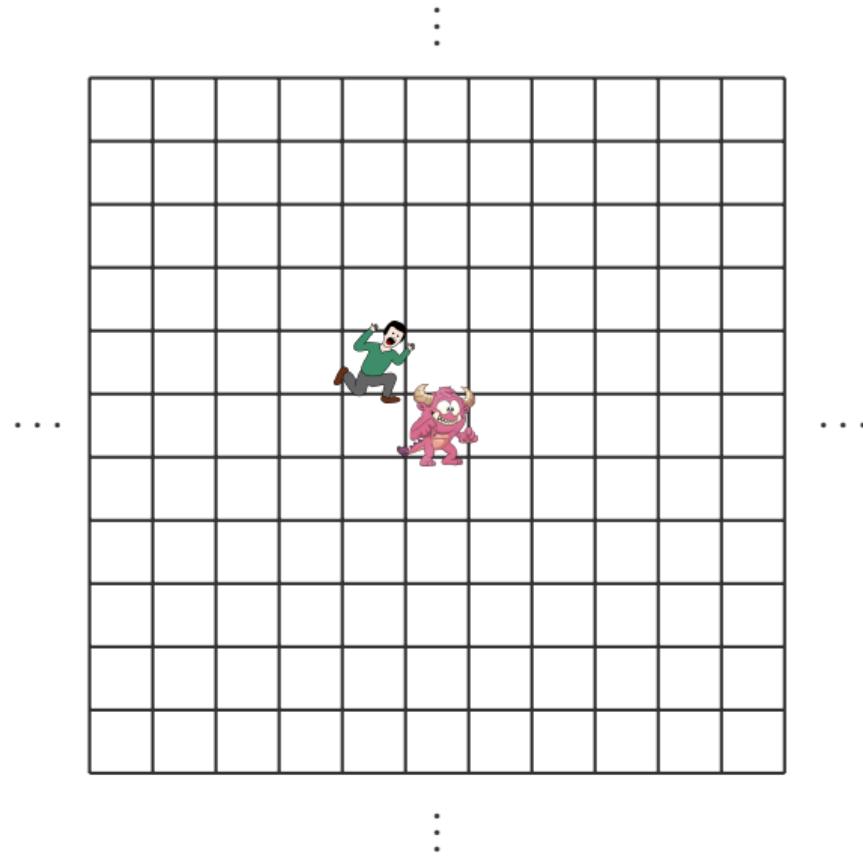
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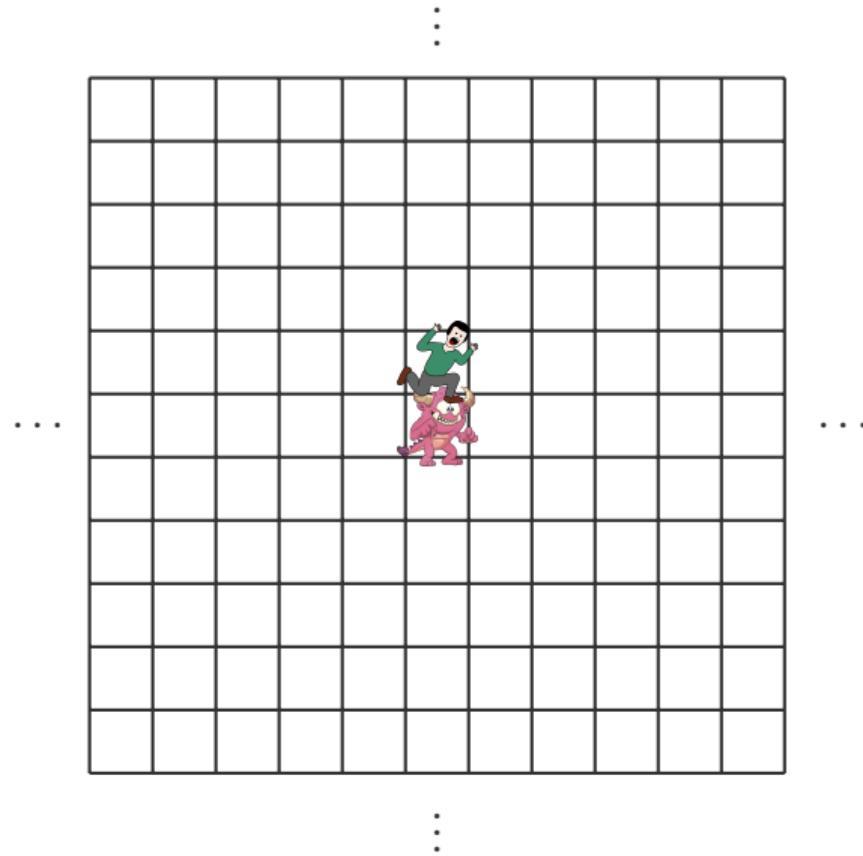
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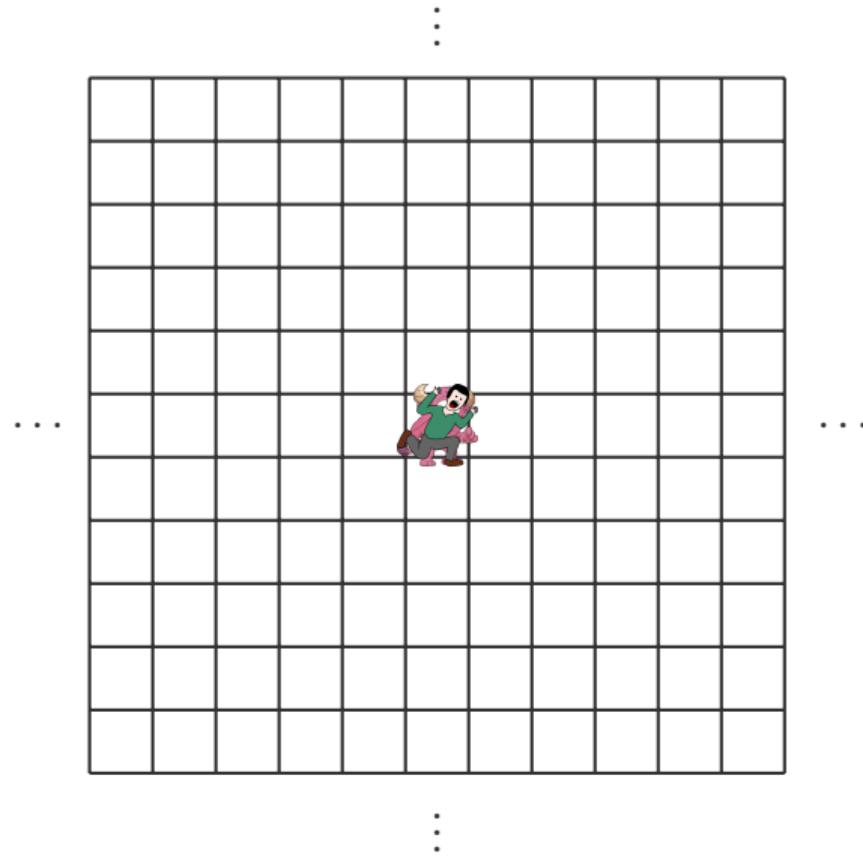
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Recurrence vs. Transience of Random Walks

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- ◎ All previously-known constructions of planar graphs with uniform polynomial growth are *recurrent*.

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Thm. [E - Lee '20]

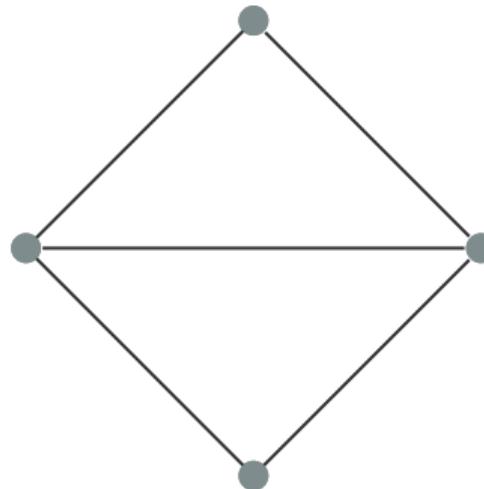
For any $\alpha > 2$ we can construct a *transient* planar graph with growth R^α .

Effective Resistance

$\text{Reff}(u, v)$ is the energy consumed by the 1-unit electrical flow from u to v , where every edge is a 1Ω resistor.

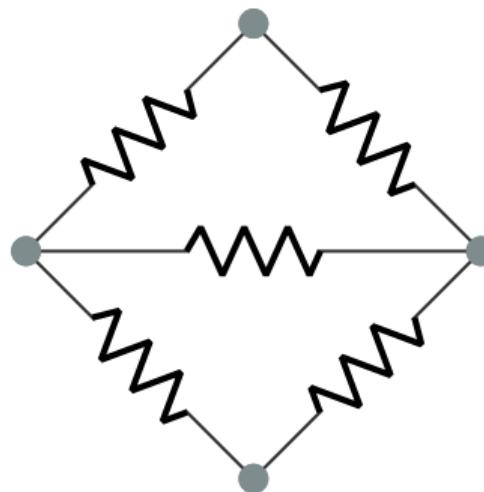
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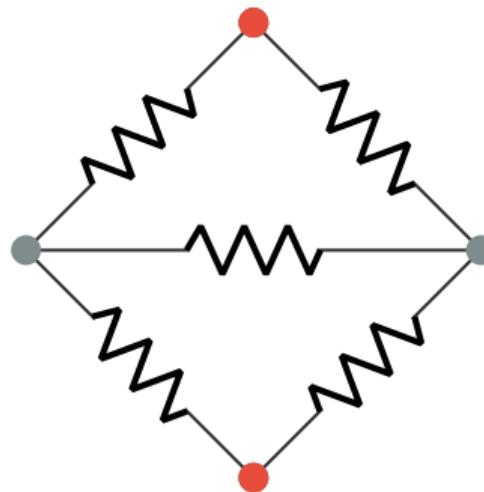
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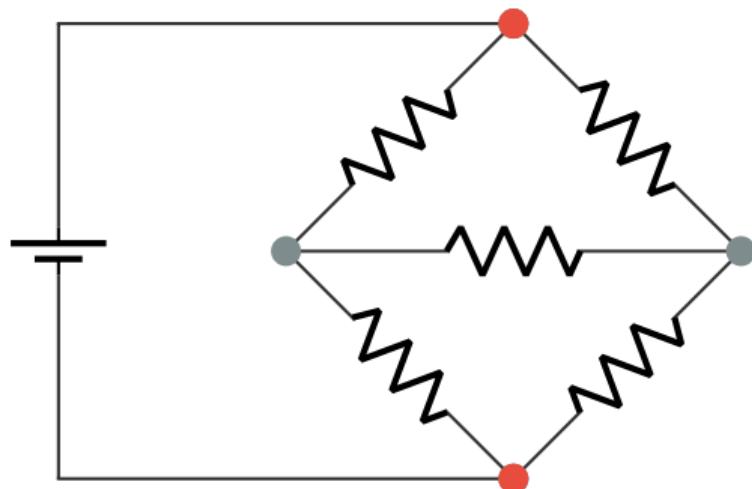
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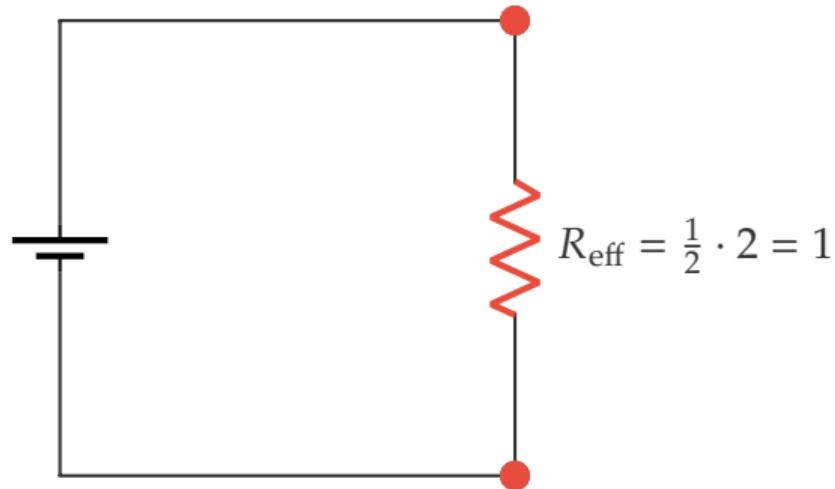
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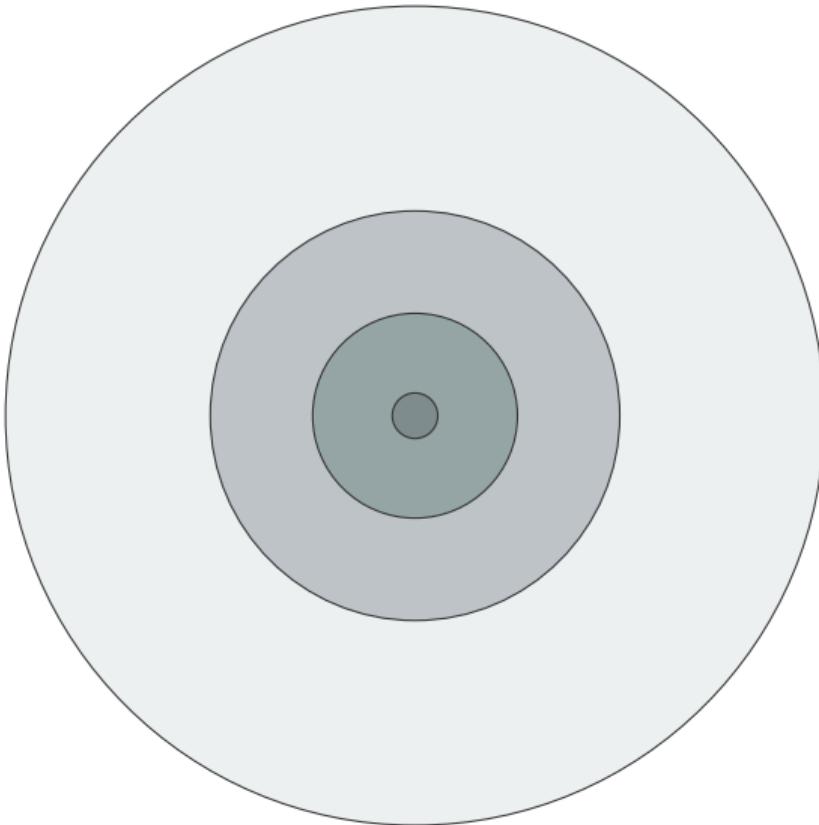
Effective Resistance and Recurrence

Thm. [Doyle - Snell '84]

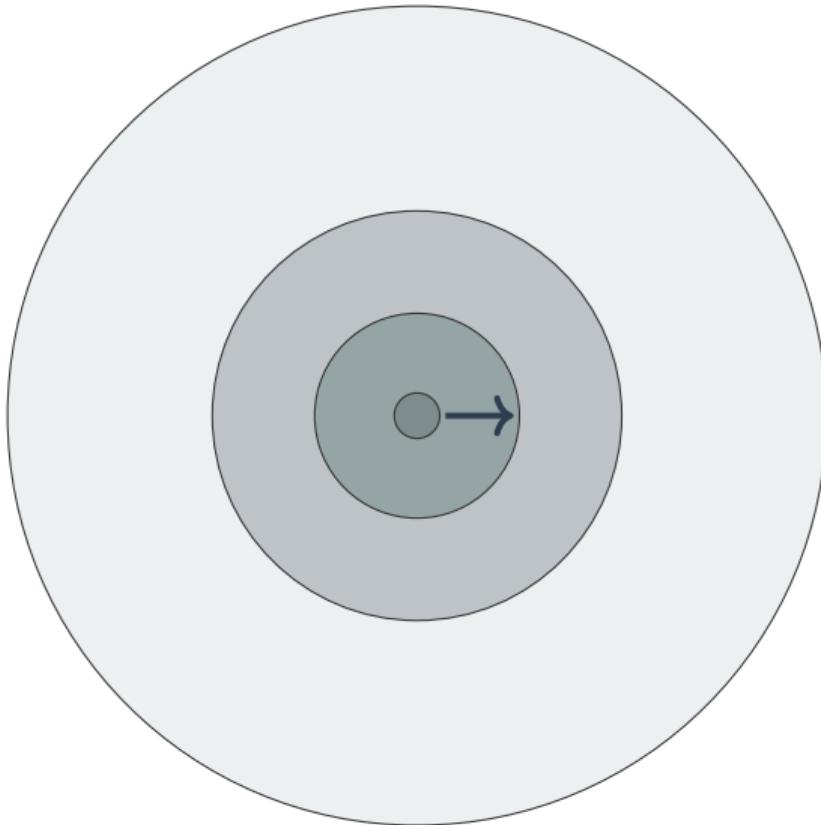
A graph G is recurrent if and only if $\text{Reff}(v, \infty) = \infty$ for $v \in V(G)$. Where

$$\text{Reff}(v, \infty) = \lim_{R \rightarrow \infty} \text{Reff}(v, V(G) \setminus B(v, R)).$$

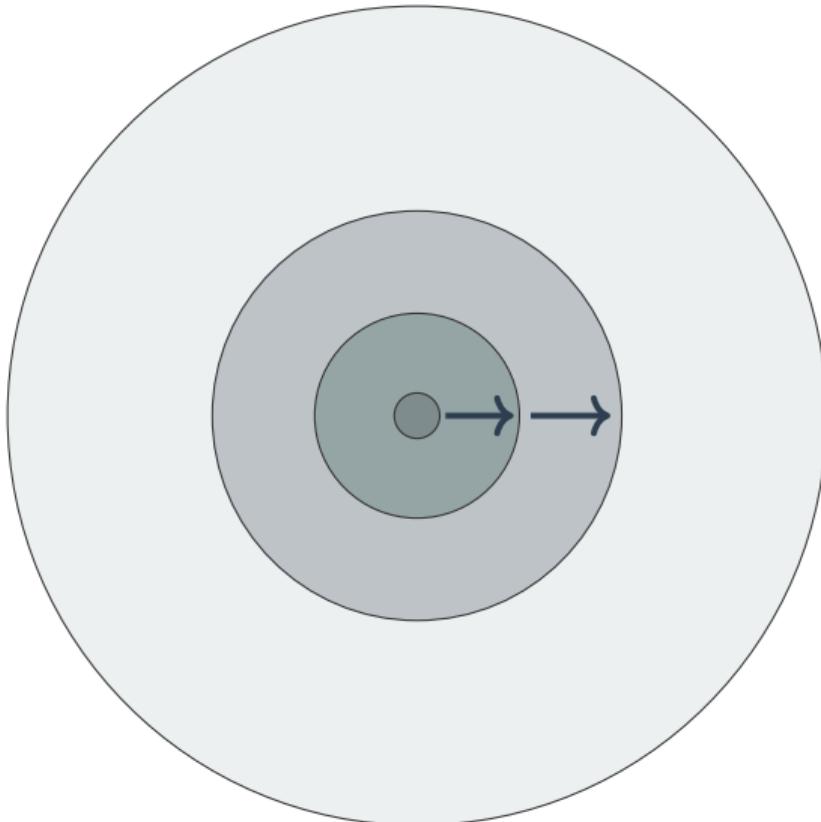
One Approach...



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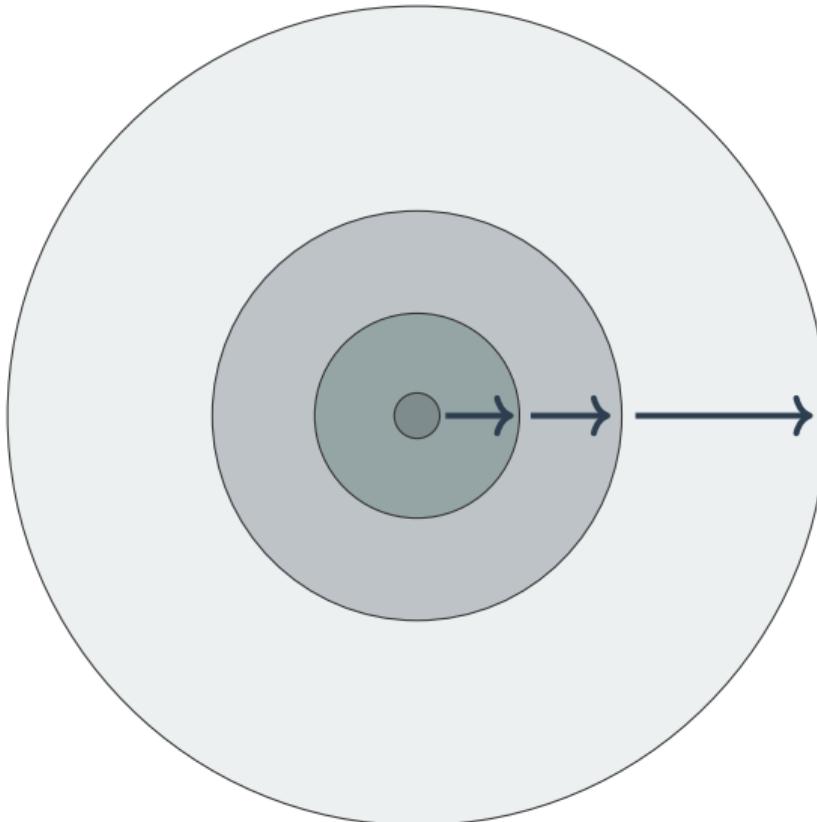


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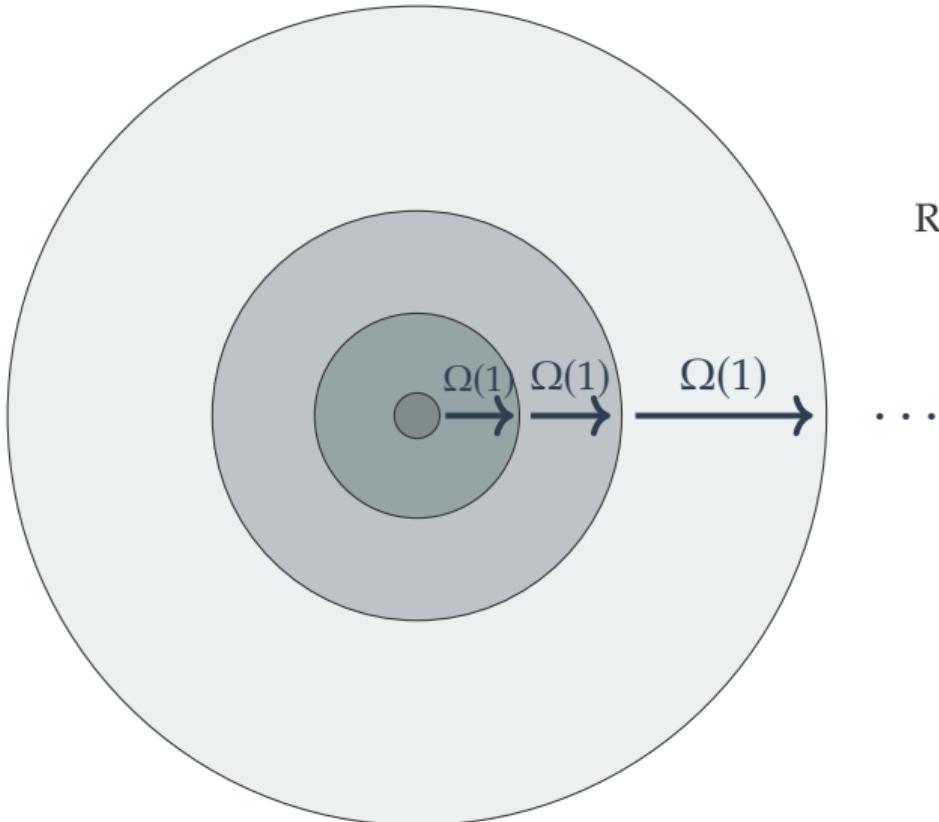
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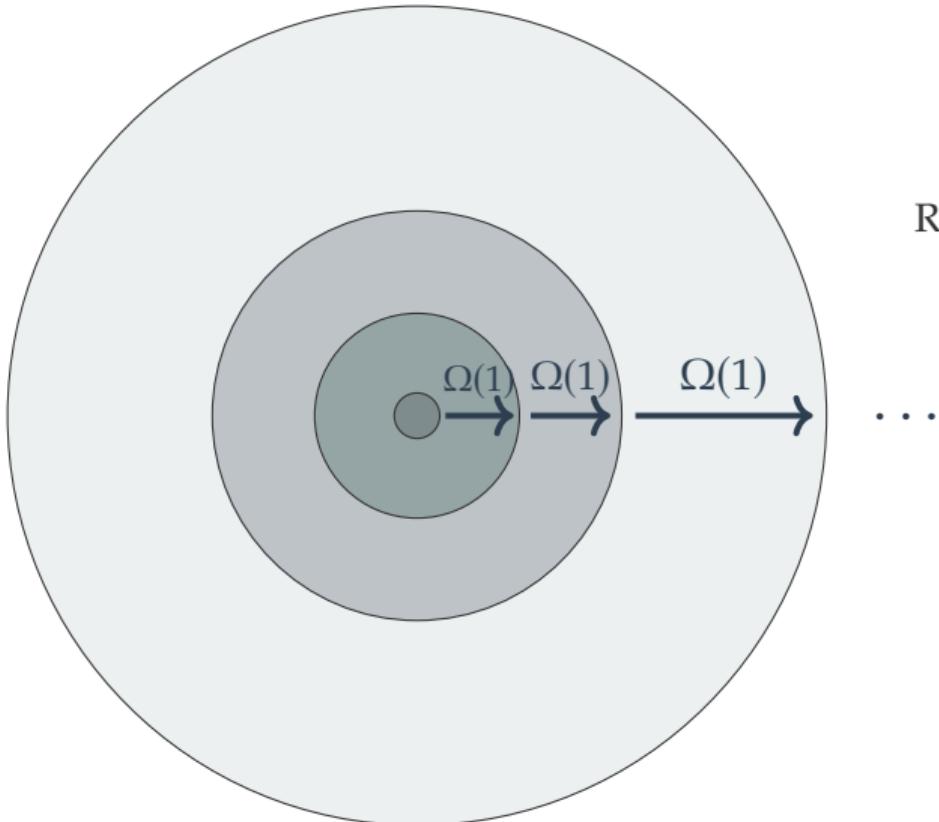
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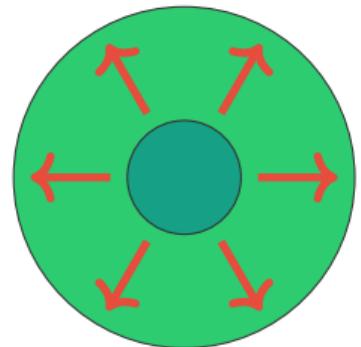
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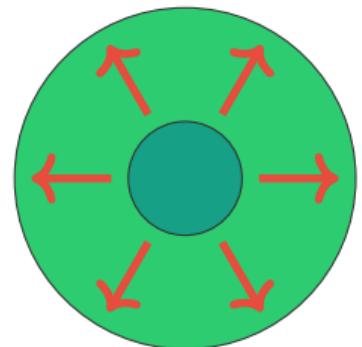


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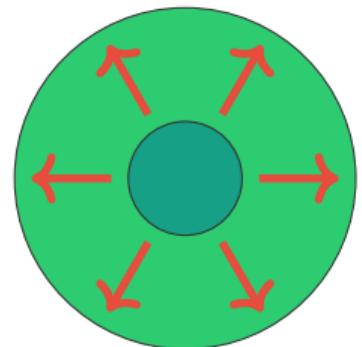


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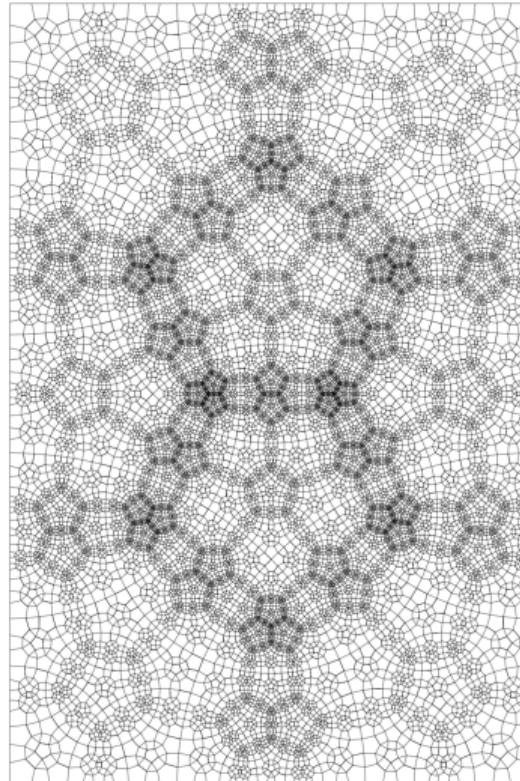
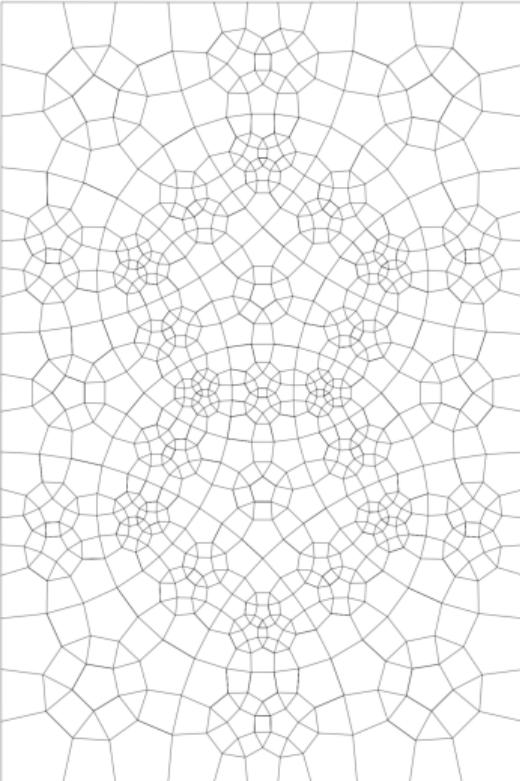
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Cannon's Conjecture



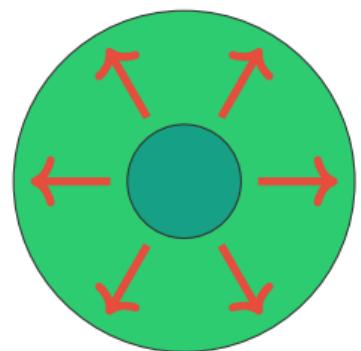
Source: Cannon - Floyd - Parry '01

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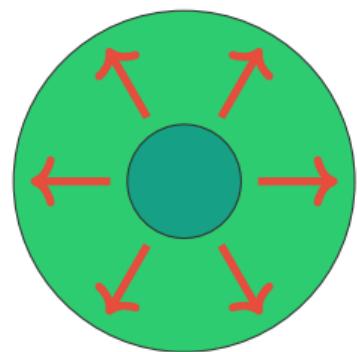


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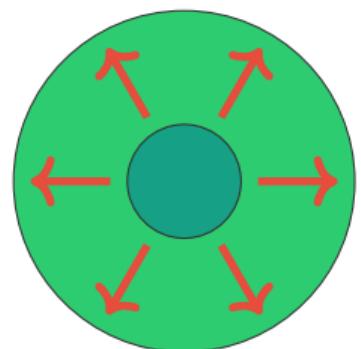
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Thm. [E - Lee '20]

For any $\epsilon > 0$ we can construct a planar graph with uniform polynomial growth having

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So far I have covered ...

- ✓ Planarity.
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- ✓ Planar graphs with uniform polynomial growth > 2 .
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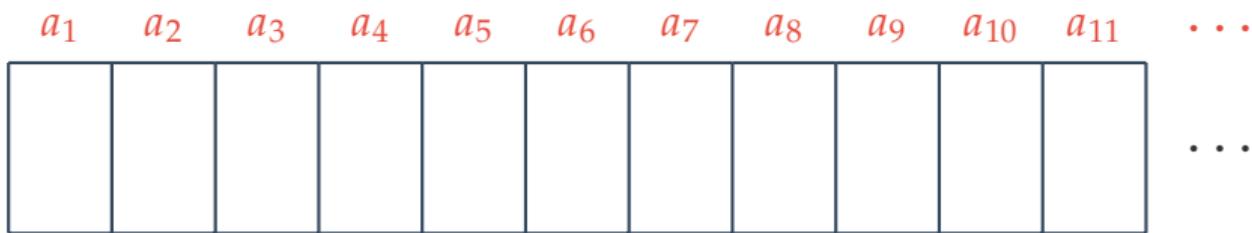
Up next ...

- Our construction.

The Construction

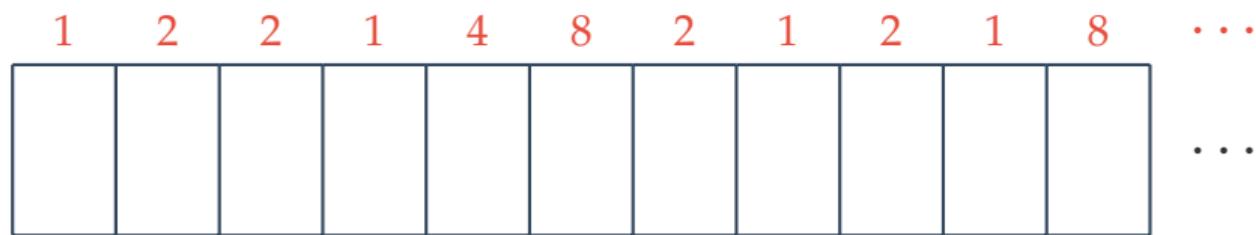
A Generalization of the Grid

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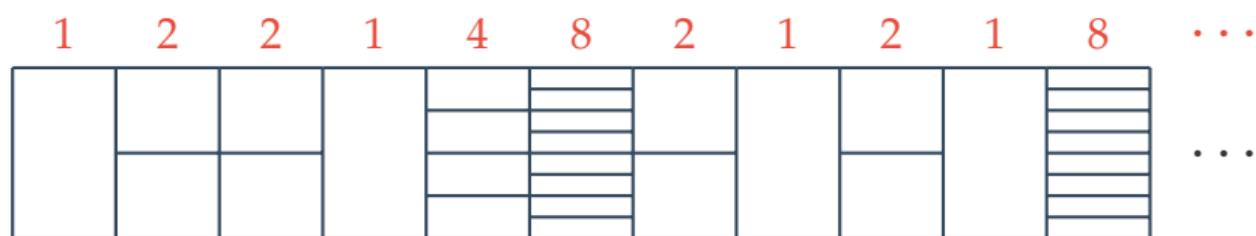
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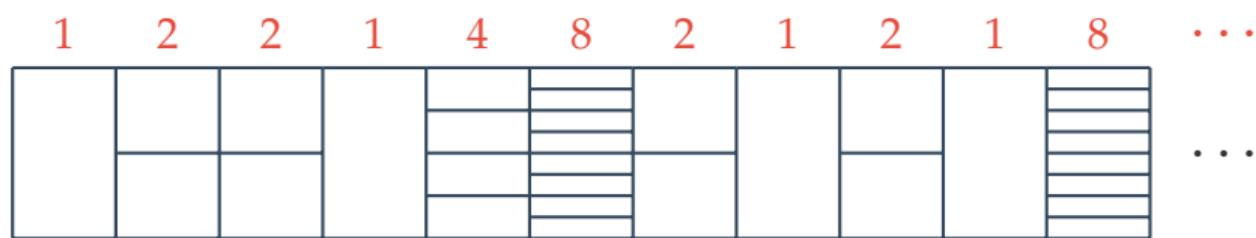
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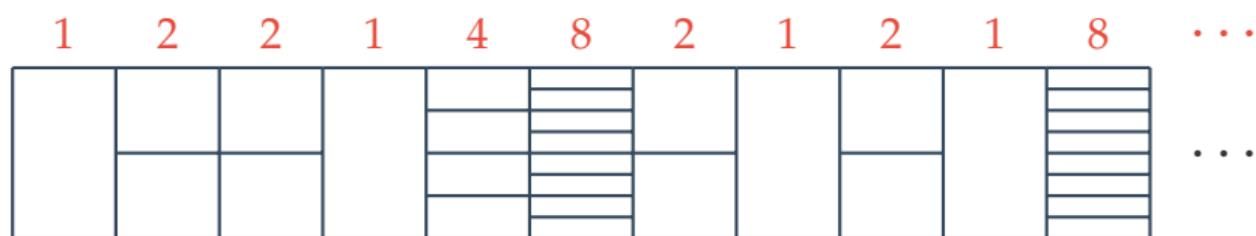
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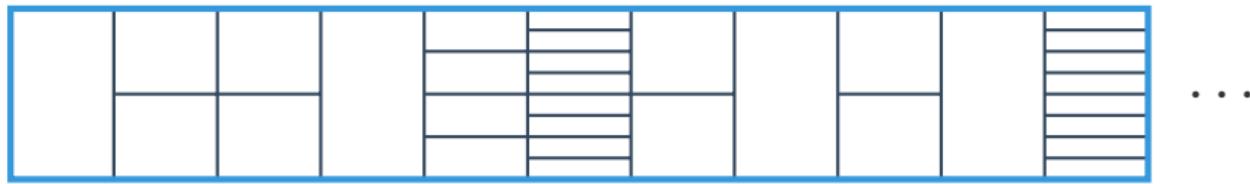
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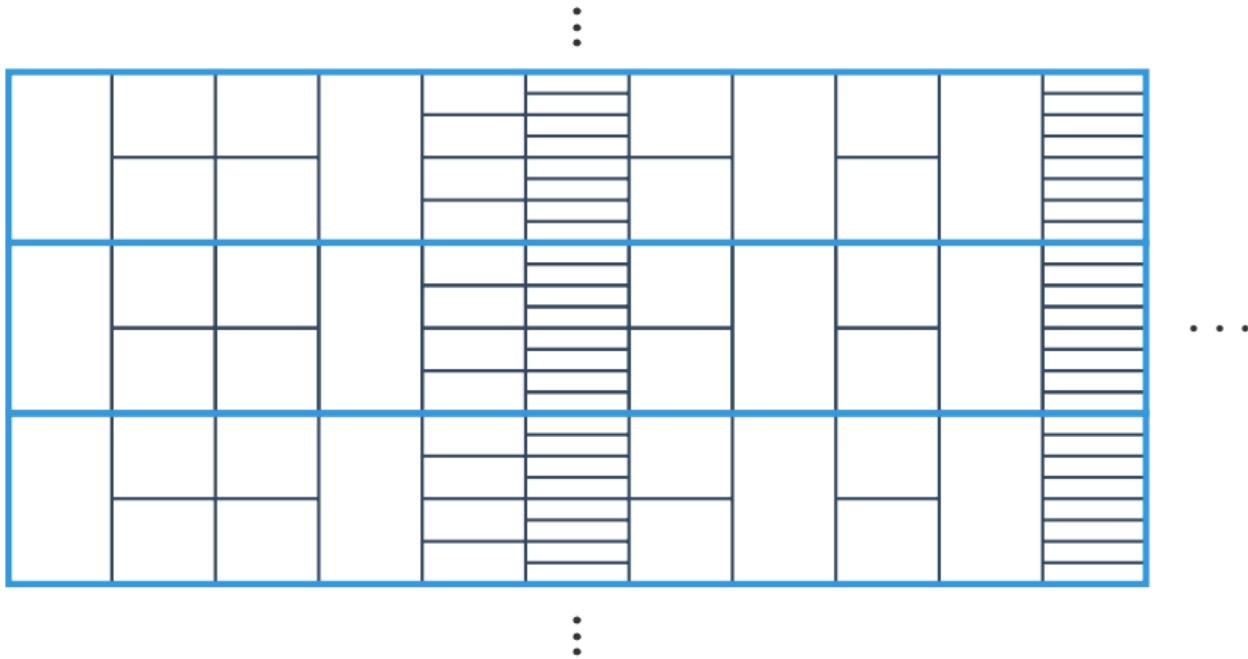
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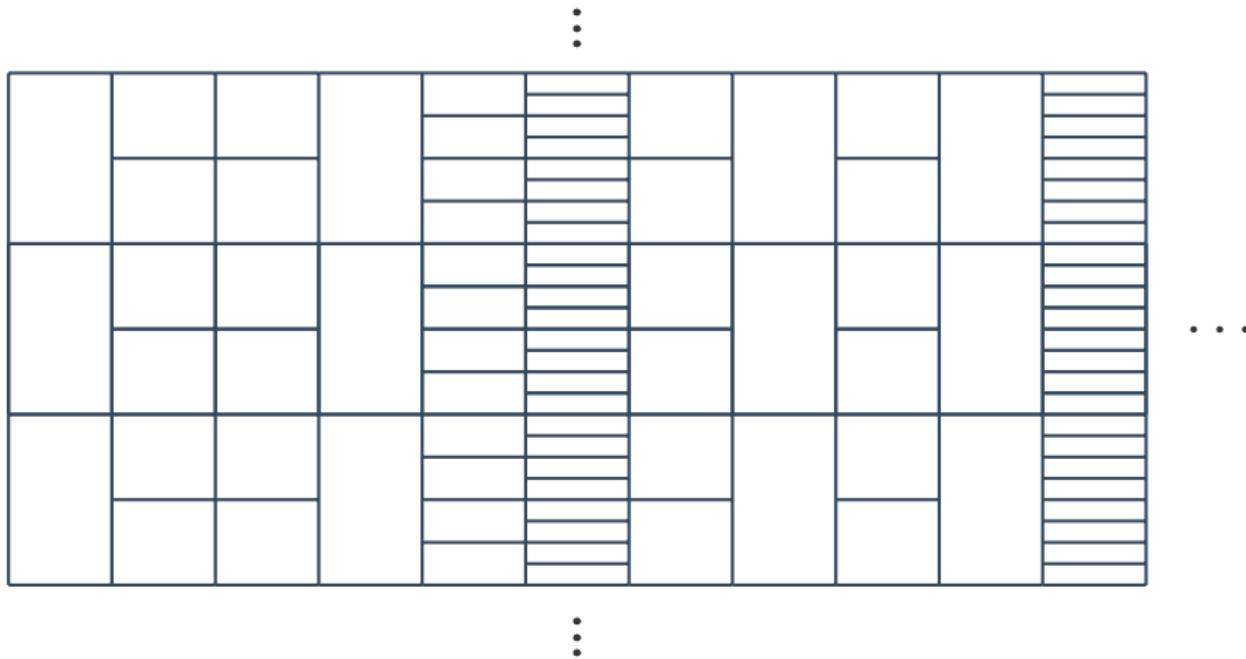
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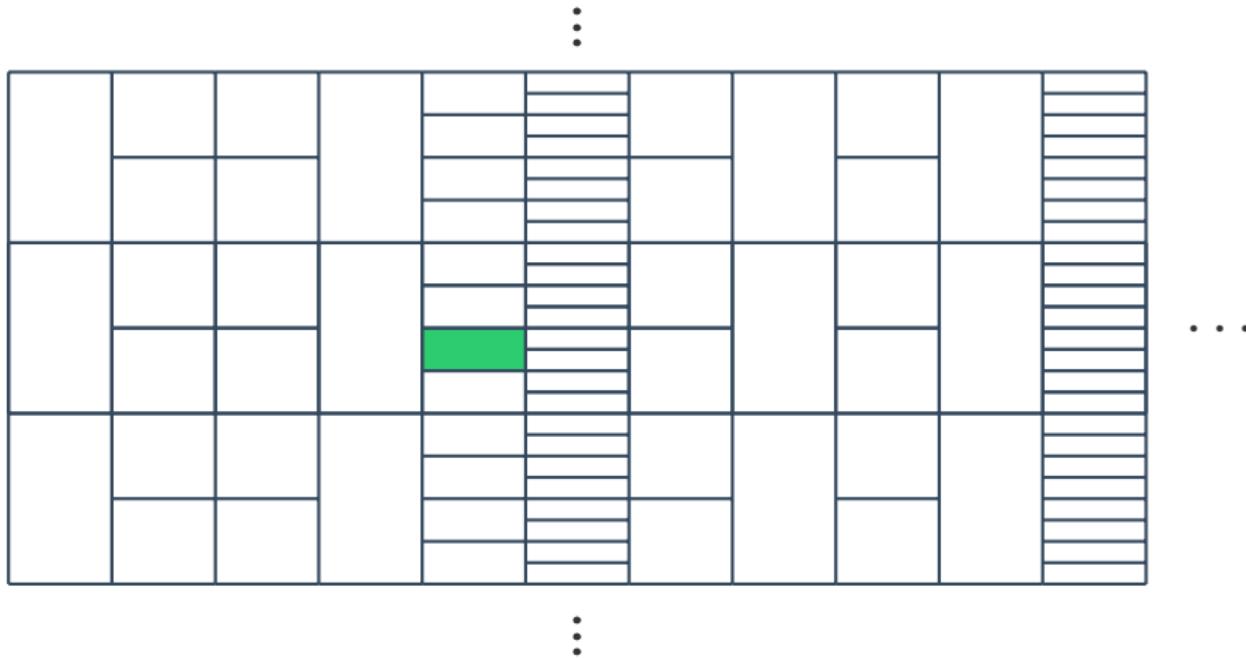
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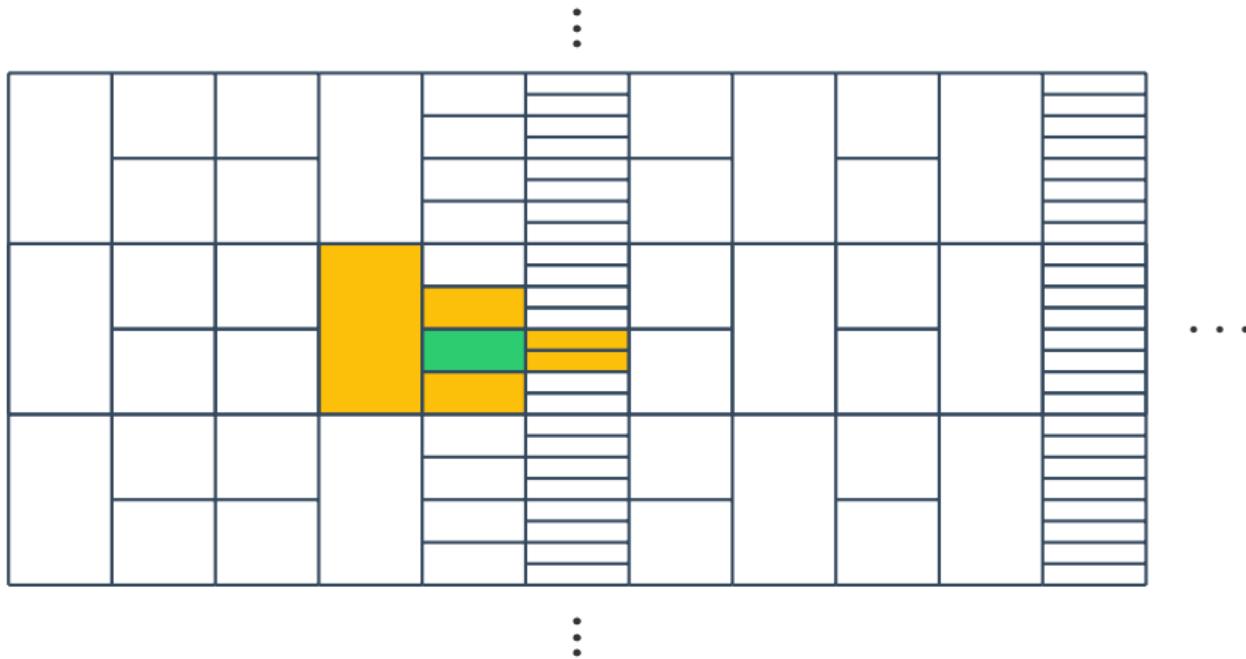
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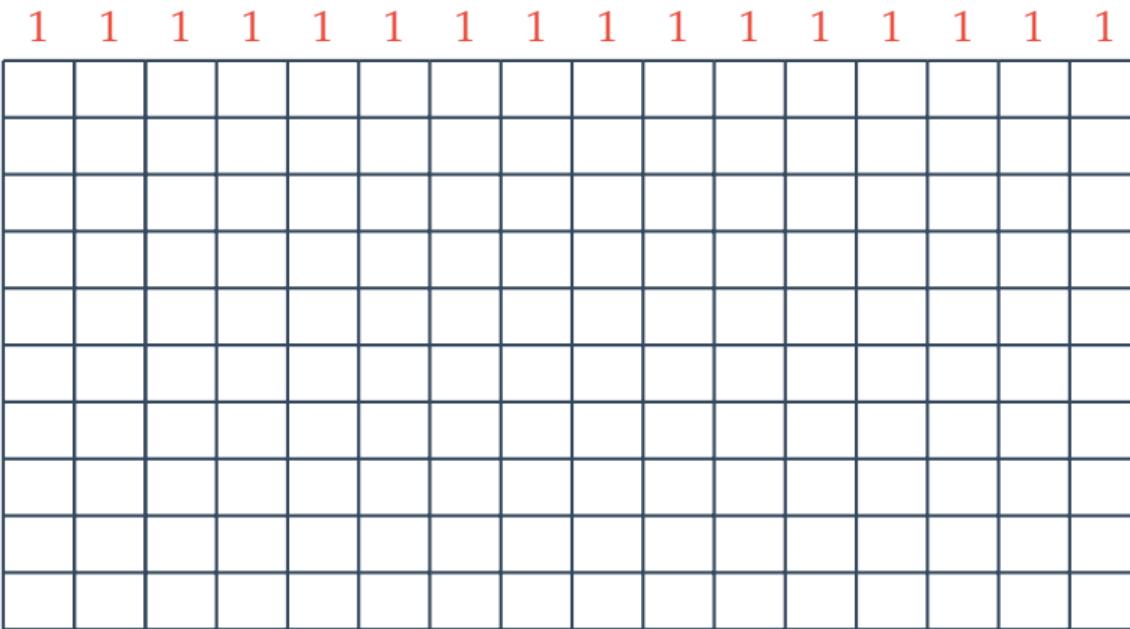


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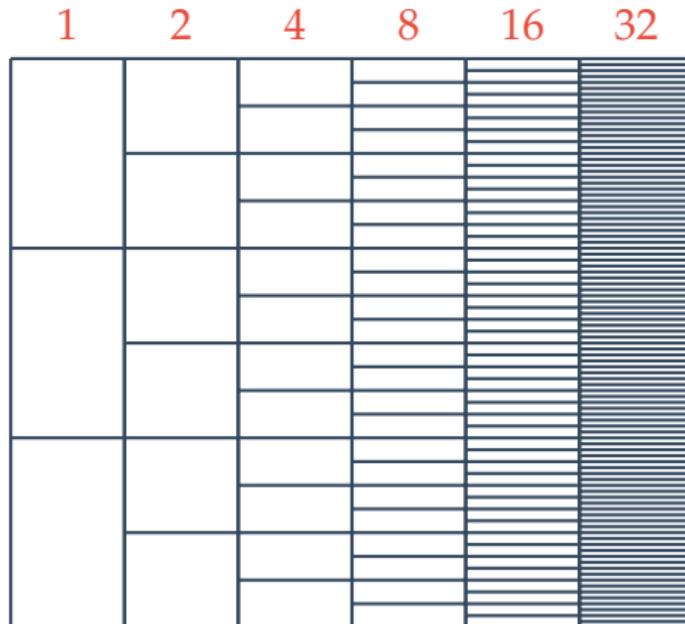
Examples – The Grid



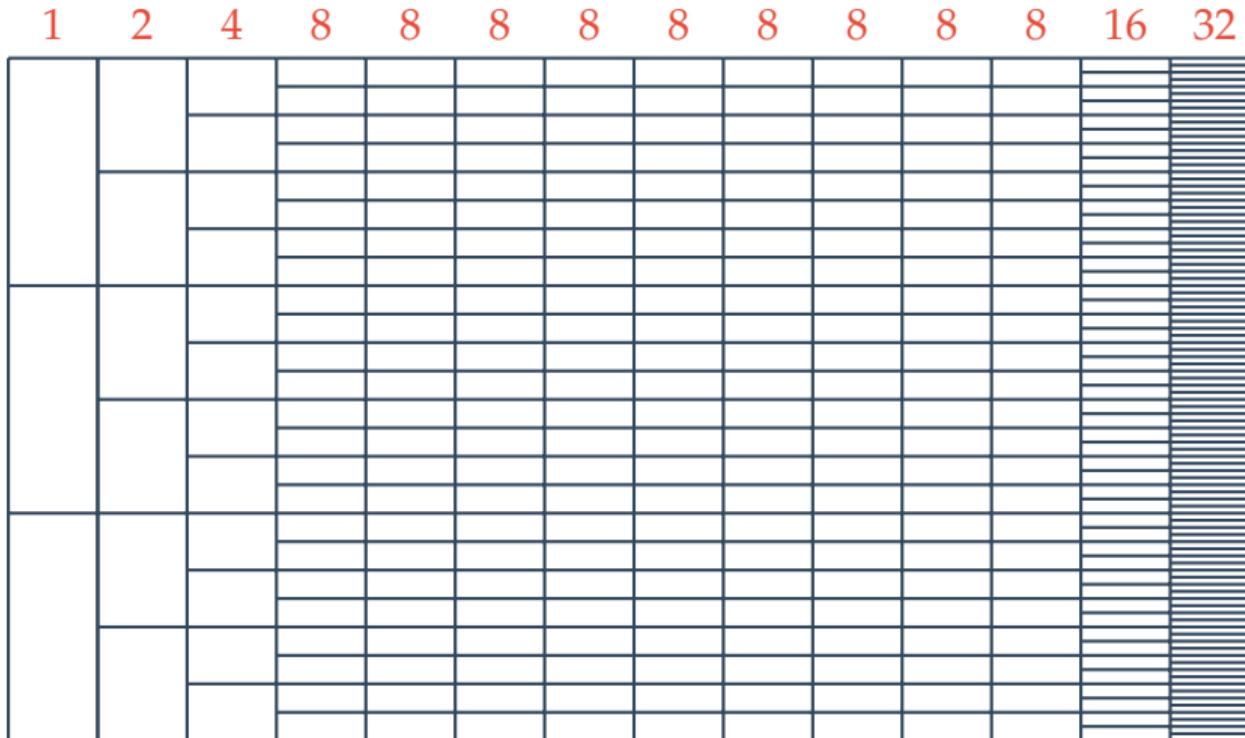
Examples – Scaling Doesn't Matter

10 10 10 10 10 10 10 10 10 10 10 10 10 10 10 10 10 10

Examples – Exponential Growth



Examples – Sub-exponential Growth?



Constructing the Sequence

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Let us set $\tau_0 = (1)$. We recursively define

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Constructing the Sequence

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Finally we set $\vec{\tau} = \lim_{n \rightarrow \infty} \tau_n$.

$$\tau_0 = (1)$$

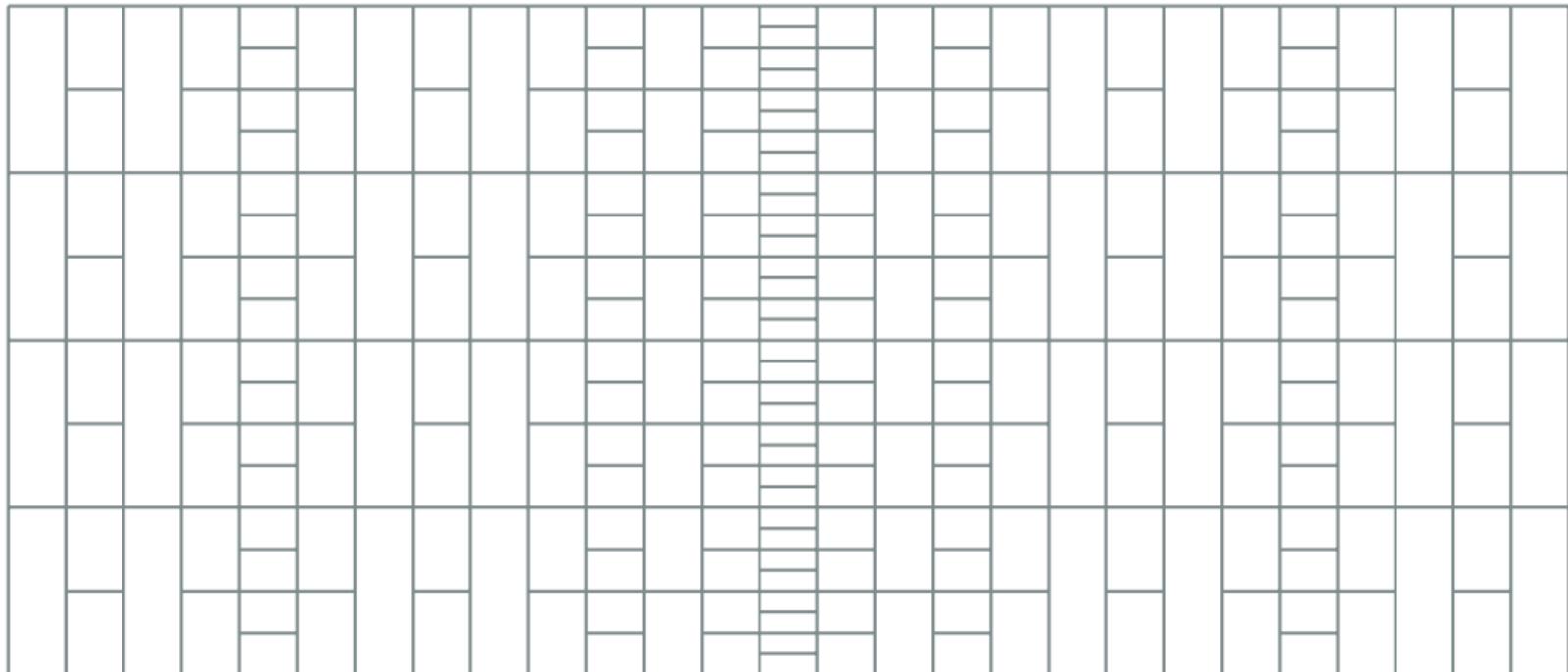
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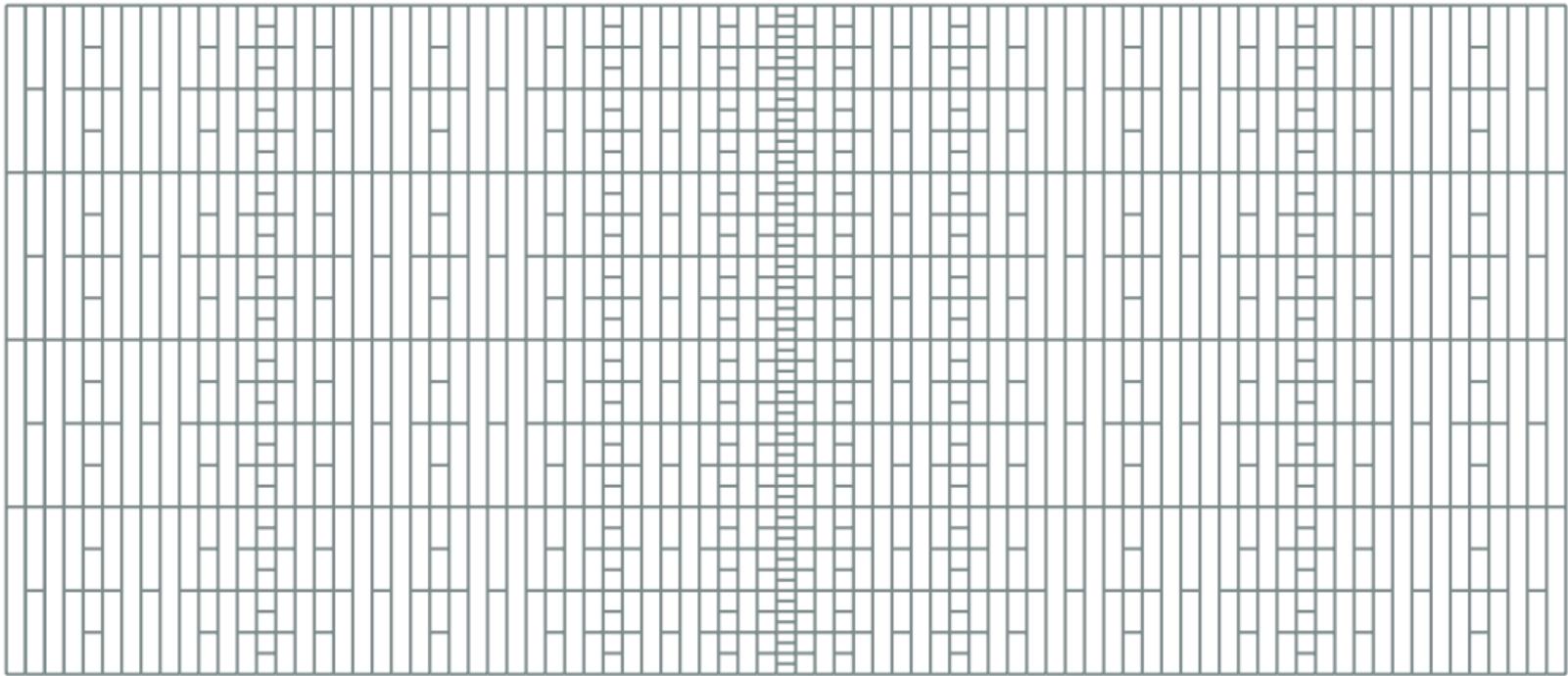
$$\tau_3 = (\underbrace{1, 2, 1, 2, 4, 2, 1, 2, 1}_{\tau_2}, \underbrace{2, 4, 2, 4, 8, 4, 2, 4, 2}_{2\tau_2}, \underbrace{1, 2, 1, 2, 4, 2, 1, 2, 1}_{\tau_2})$$

How Does It Look Like?

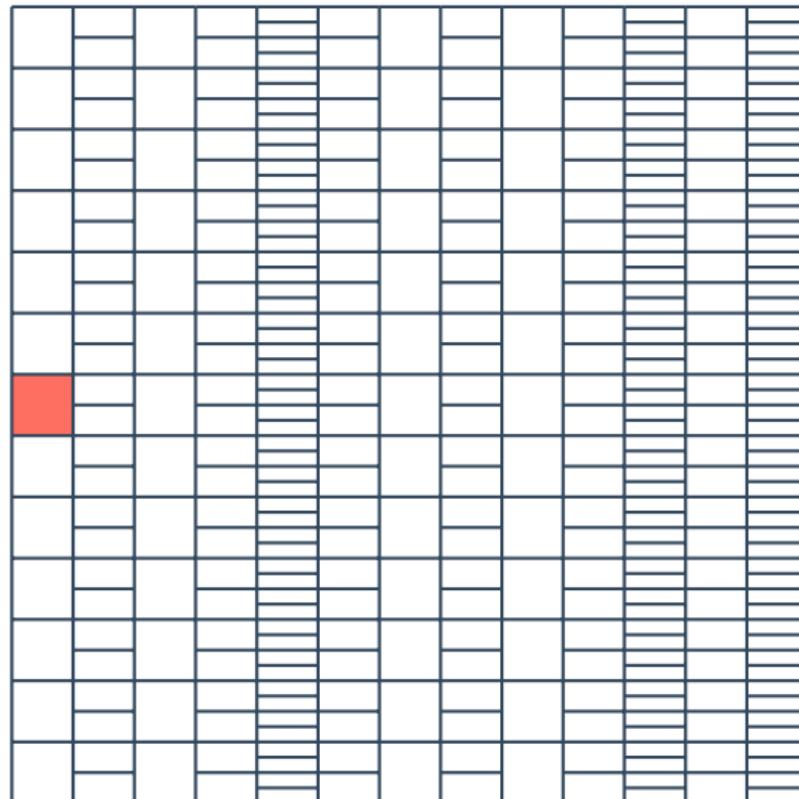
1 2 1 2 4 2 1 2 1 2 4 2 4 2 4 8 4 2 4 2 1 2 1 2 4 2 1 2 1



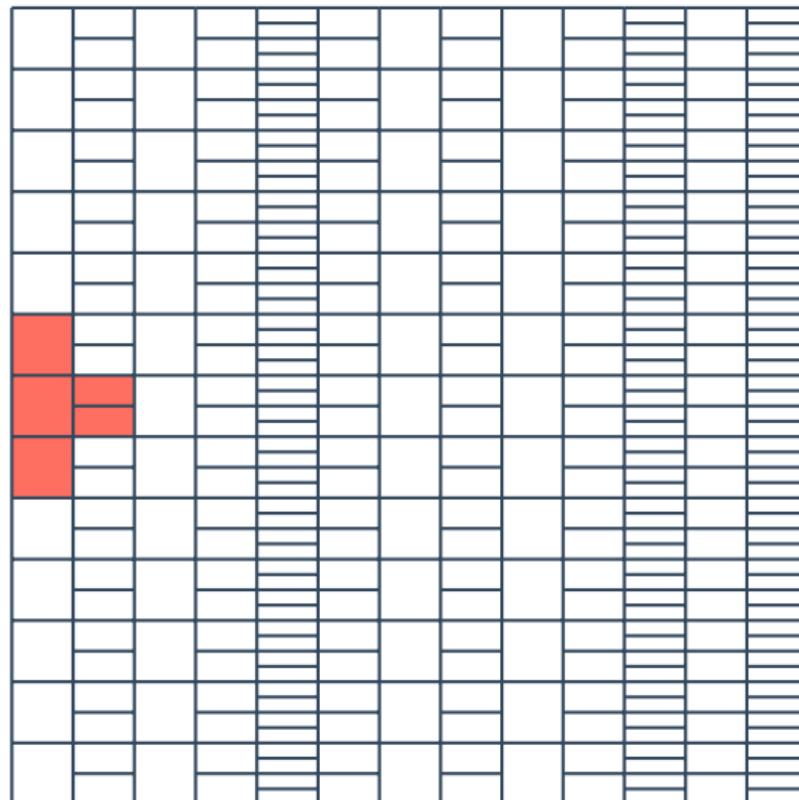
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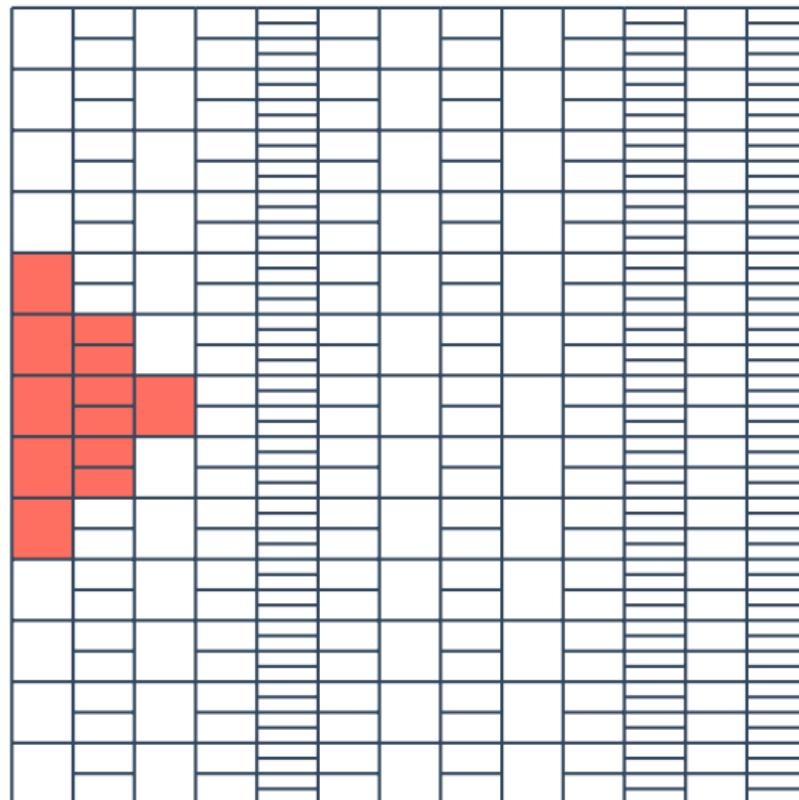
Volume Growth



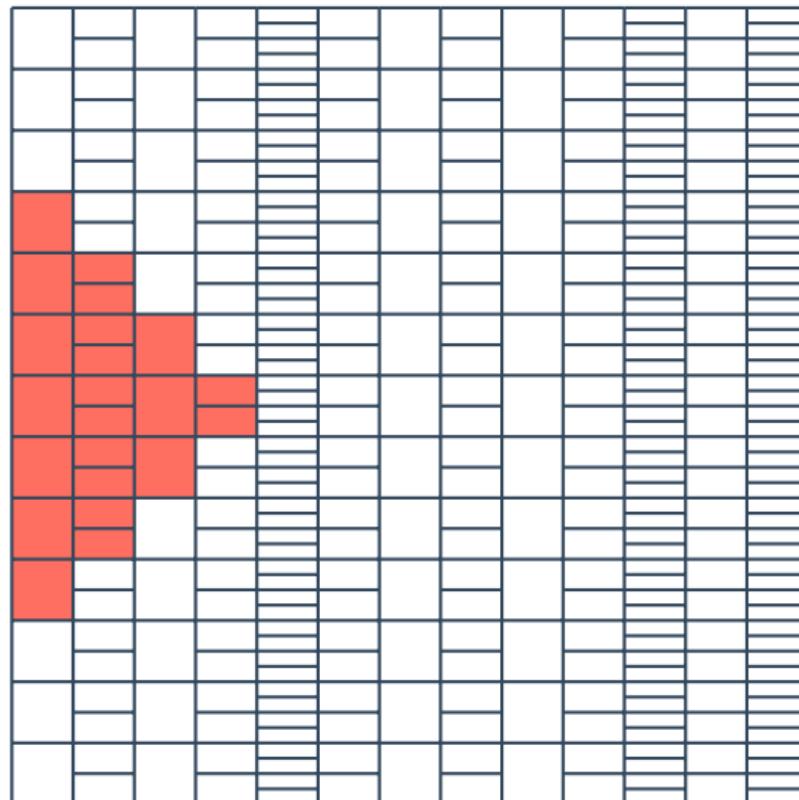
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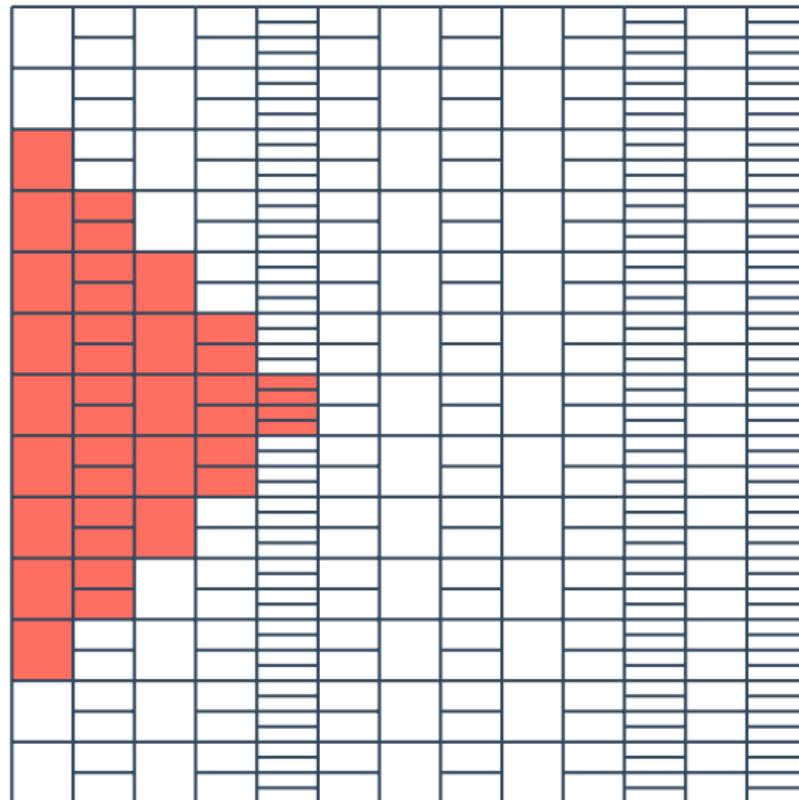
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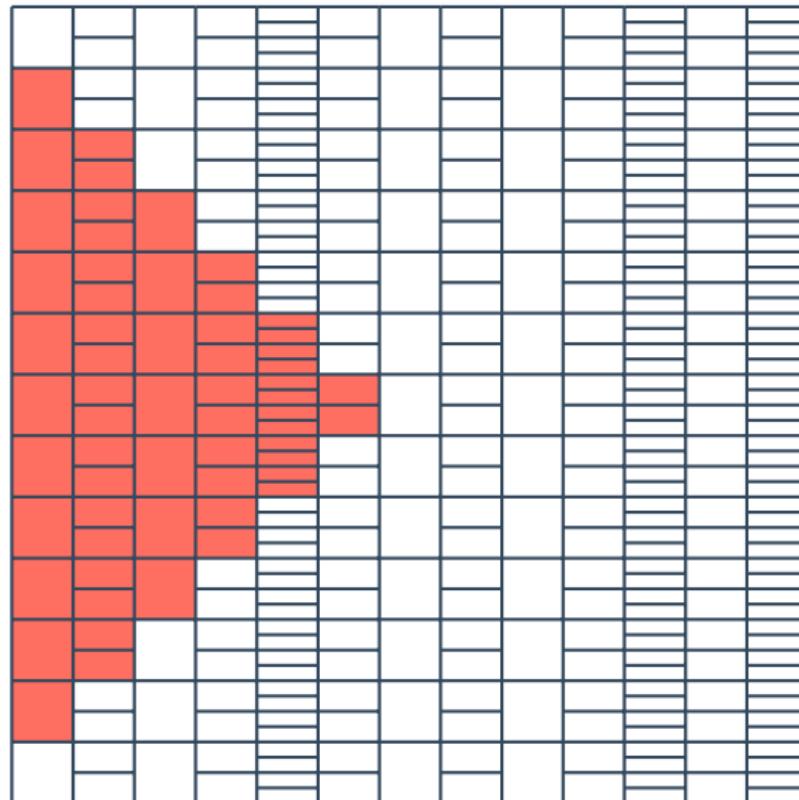
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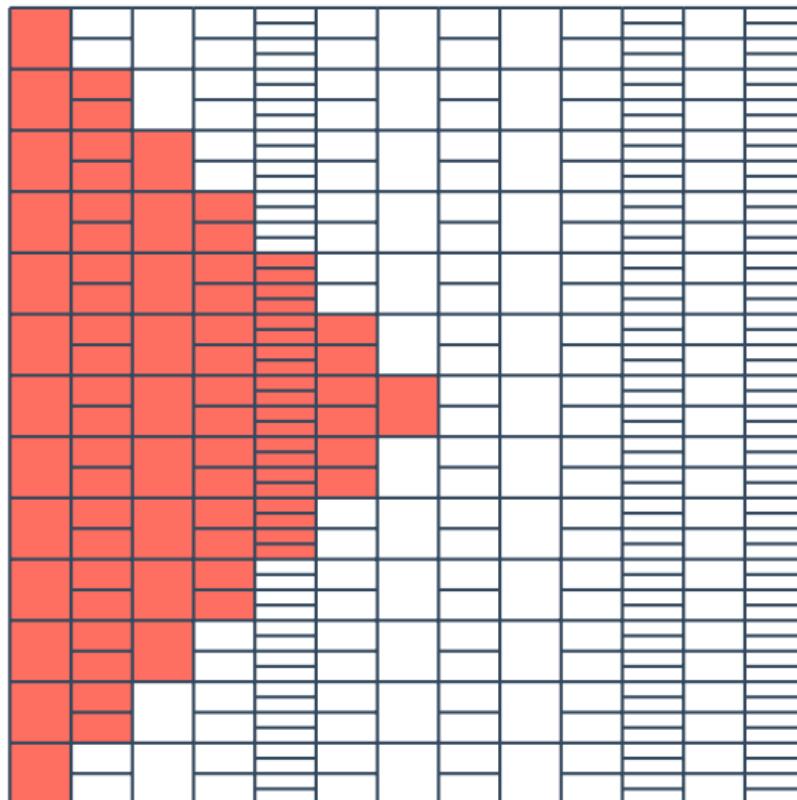
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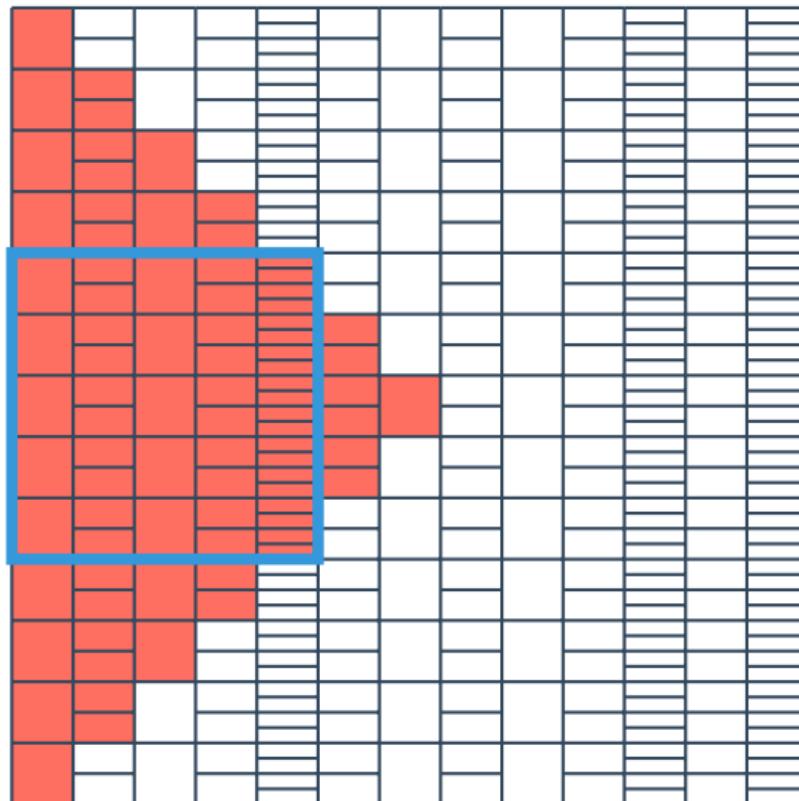
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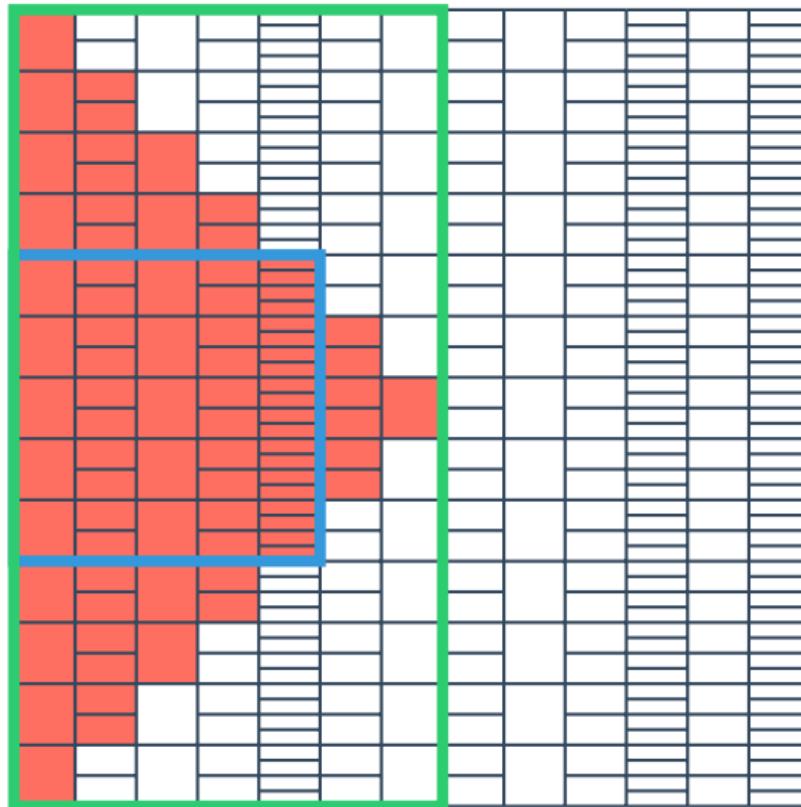
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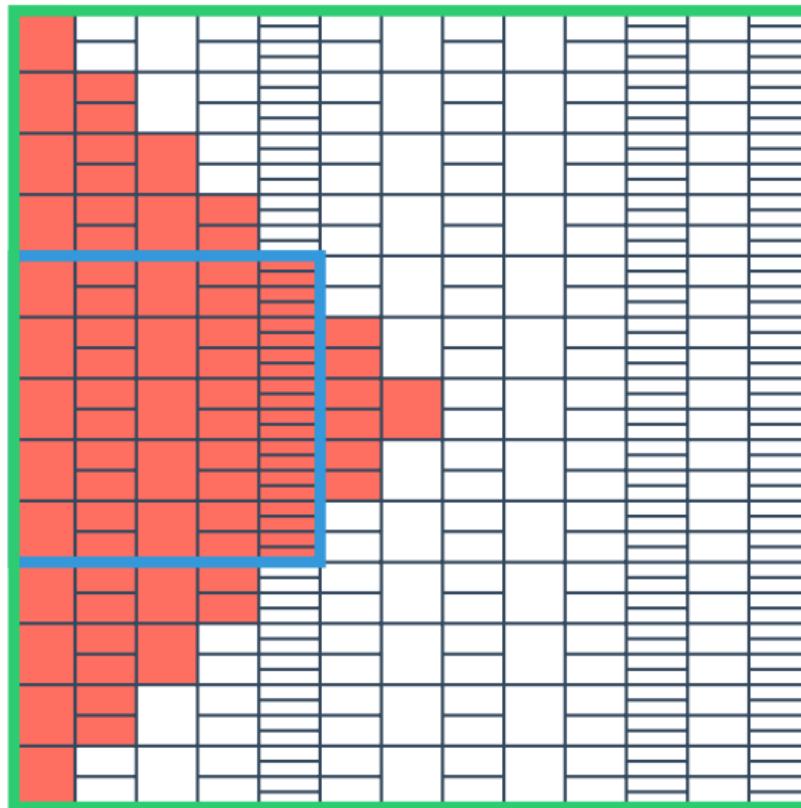
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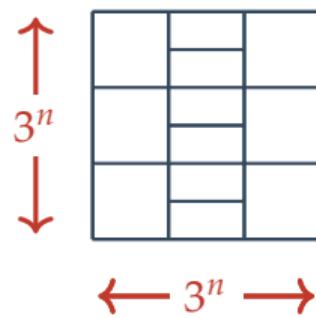
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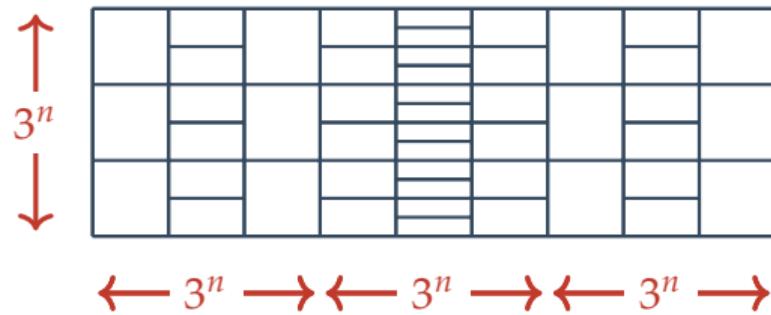
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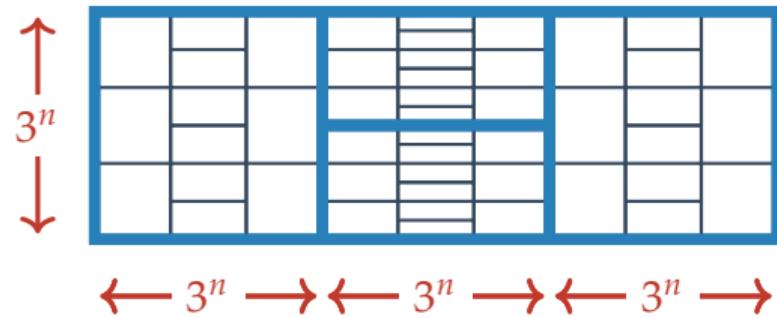
A Subdivision Rule



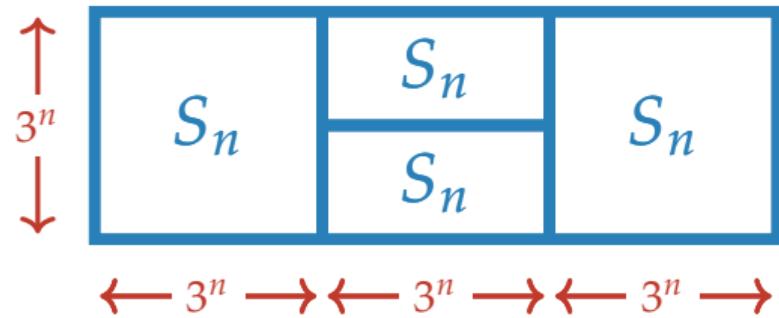
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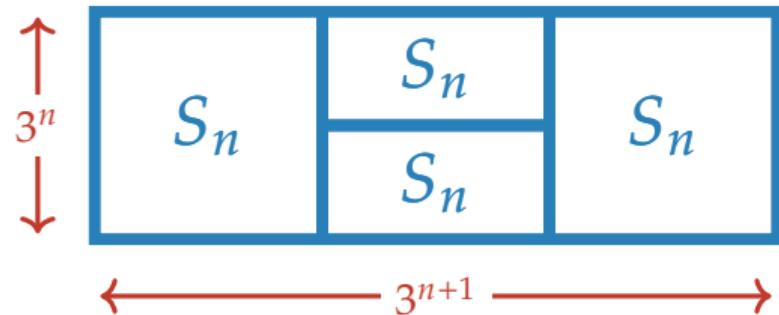
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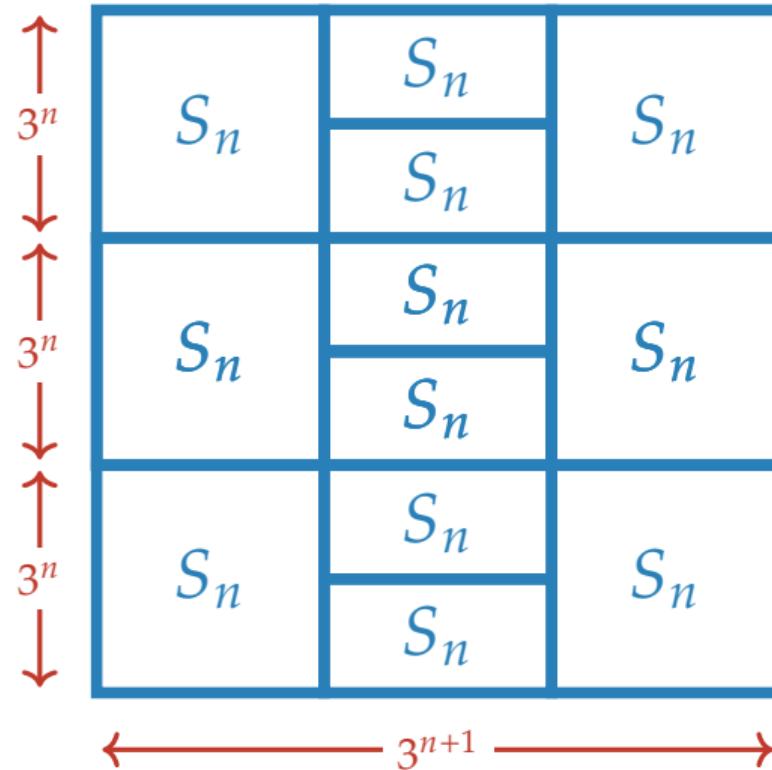
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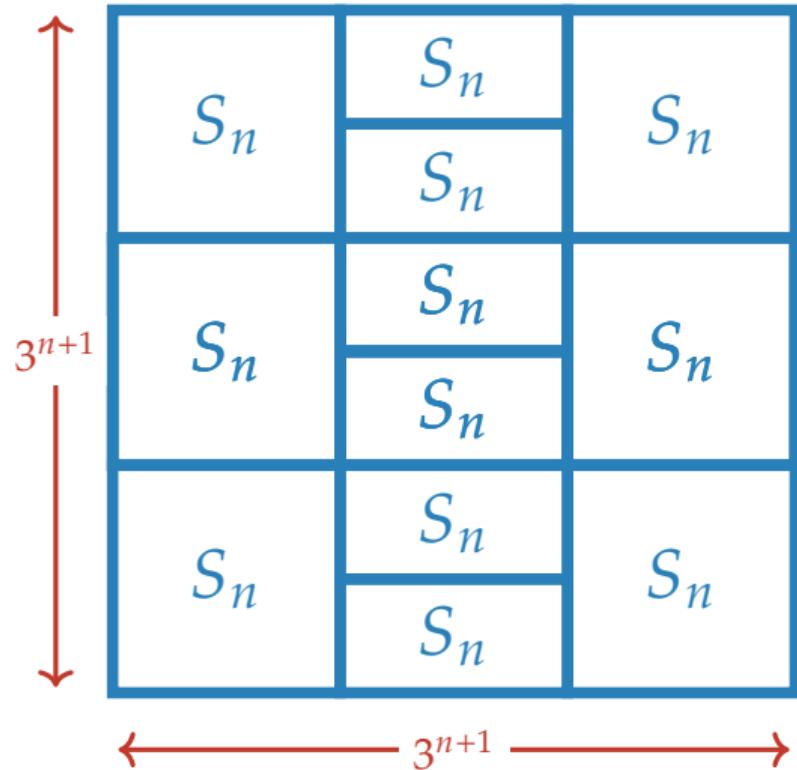
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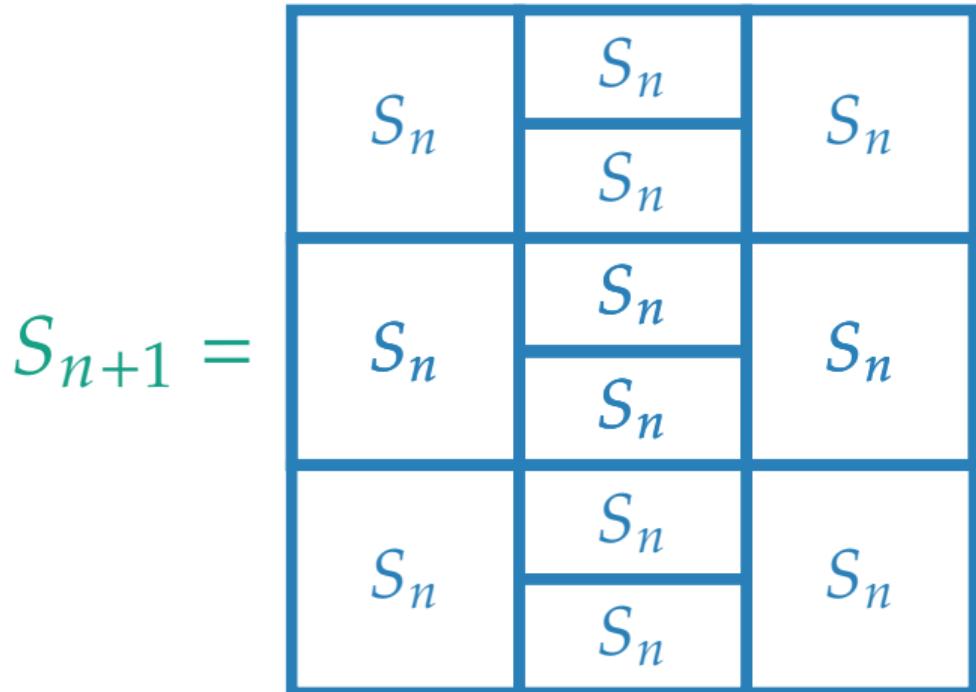
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The Transient Construction

$T =$  \dots

The diagram shows a sequence of five rectangular boxes arranged horizontally. Each box contains a blue mathematical expression: s_0 , s_1 , s_2 , s_3 , and s_4 . To the left of the first box, there is a green mathematical expression $T =$. To the right of the fifth box, there are three black dots \dots , indicating that the sequence continues indefinitely.

So far I have covered ...

- ✓ Planarity.
- ✓ Uniform volume growth.
- ✓ Planar graphs with uniform polynomial growth > 2 .
 - ✓ Their *structure*.
 - ✓ *Random walks* and *effective resistances* in these graphs.
 - ✓ Cannon's conjecture.
 - ✓ Our construction.

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- ◎ Cannon’s conjecture remains wide open.

Questions?

Thank you!