# DV ROUTING ALGORITHM WITH POISONED REVERSE

# University of Dhaka Department of Computer Science & Engineering

Group No: 18

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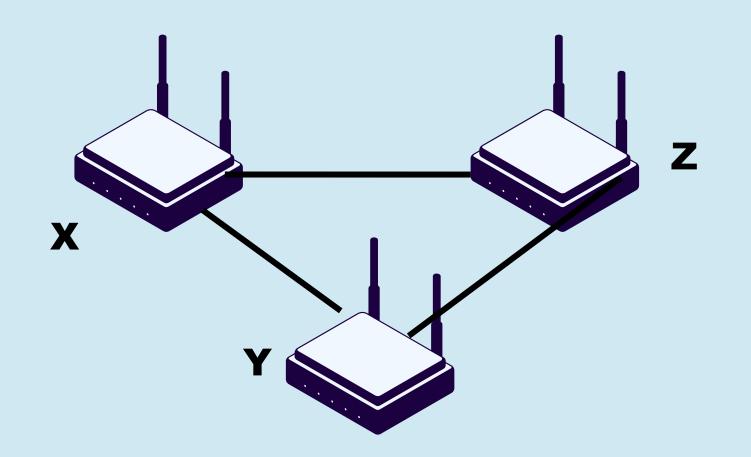
**Engineering** 

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## COUNT TO INFINITY PROBLEM



### **Shortest Path For**

X:

X->Y is 4

X->Z is 5 (X->Y-Z)

### **Shortest Path For Y:**

Y->X is 4

Y->Z is 1

#### case:

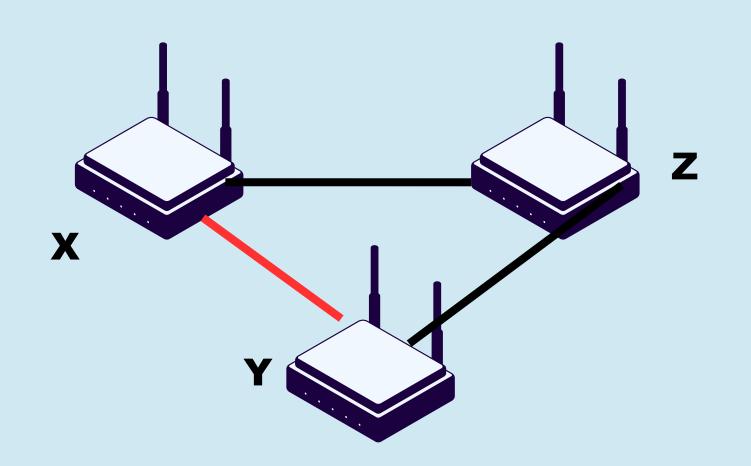
What will happen if the link cost c(X,Y) is increased to 60?

## **Shortest Path For Z:**

Z->X is 5 (Z->Y->X)

Z->Y is 1

# COUNT TO INFINITY PROBLEM



#### **Shortest Path For**

X:

X->Y is 4

 $X \rightarrow Z$  is 5  $(X \rightarrow Y - Z)$ 

#### **Shortest Path For Y:**

Y->X is 4 Y->Z is 1

#### **Shortest Path For Z:**

Z->X is 5 (Z->Y->X) Z->Y is 1

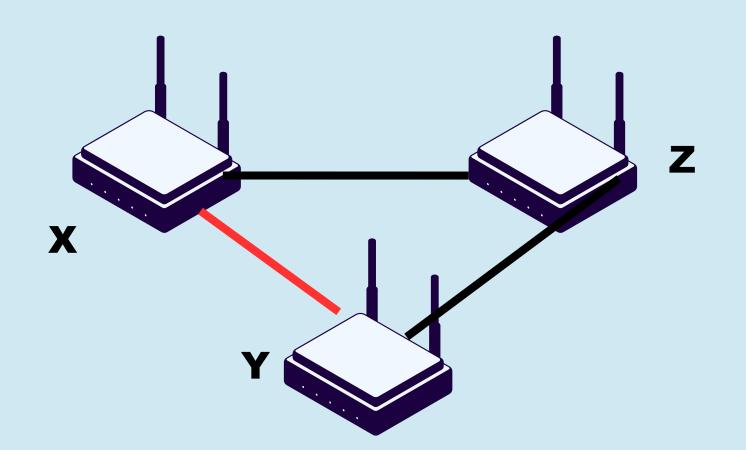
### At t<sub>0</sub>: Y detects cost change

- Y's calculation:  $D_Y(X) = min\{60+0, 1+5\} = 6$
- Y decides: Route to X via Z (cost 6 vs direct cost 60)
- Problem: Creates routing loop! Y→Z→Y→Z...

### At t<sub>1</sub>: Y informs Z

- Y tells Z: "My distance to X is now 6"
- Z calculates: D\_Z(X) = min{50+0, 1+6} = 7
- Z updates: "My distance to X is now 7"

## COUNT TO INFINITY PROBLEM



## **The Counting Loop:**

- $t_2$ : Y gets Z's update  $\rightarrow$  D\_Y(X) = 1+7 = 8
- $t_3$ : Z gets Y's update  $\rightarrow$  D\_Z(X) = 1+8 = 9
- $t_4$ : Y gets Z's update  $\rightarrow$  D\_Y(X) = 1+9 = 10
- ...continues for 44 iterations...

### When Does It Stop?

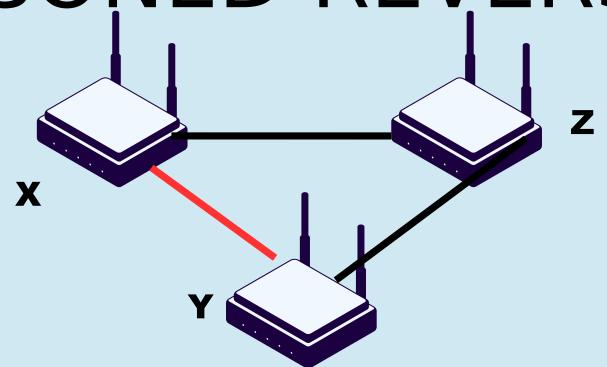
- Stops when: Z's cost via Y > 50 (direct cost)
- After 44 rounds: Z finally chooses direct path to X
- Total time: Until costs count up from 6 to 50+

But we can improve this by using

**Poisoned Reverse** 

**Technique** 

# POISONED REVERSE



## DELIBERA TELIE

If router Z routes to X via Y, Z tells Y that its distance to X is  $\infty$ . (Z lies even though it knows D\_Z(X) = 6)

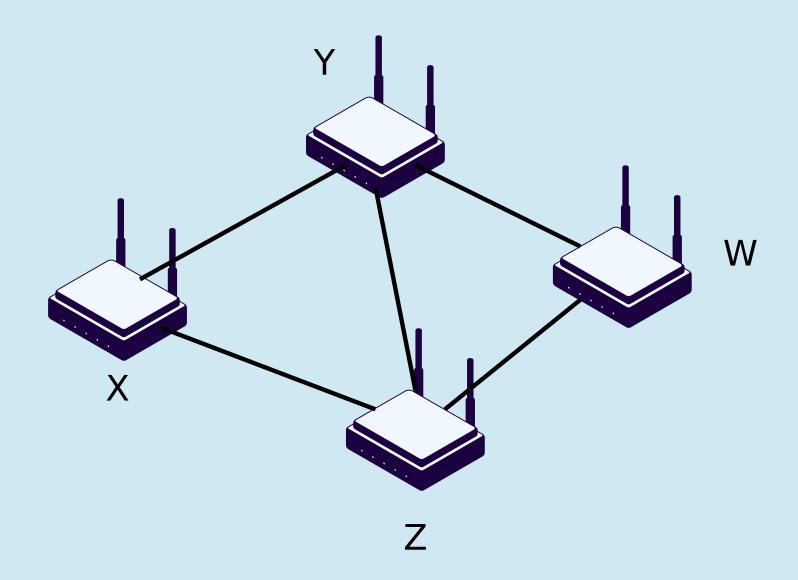
## WHY LIE?

Prevents Y from thinking Z is a valid route \*\*Movement of Youting To X via Z, preventing routing loops

## HOW LONG LIE?

**Z** continues to advertise  $D_Z(X) = \infty$  to Y as long as Z routes to X **Diack** Z finds a better direct path, it stops lying and tells the  $^6$  truth.

# POISONED REVERSE

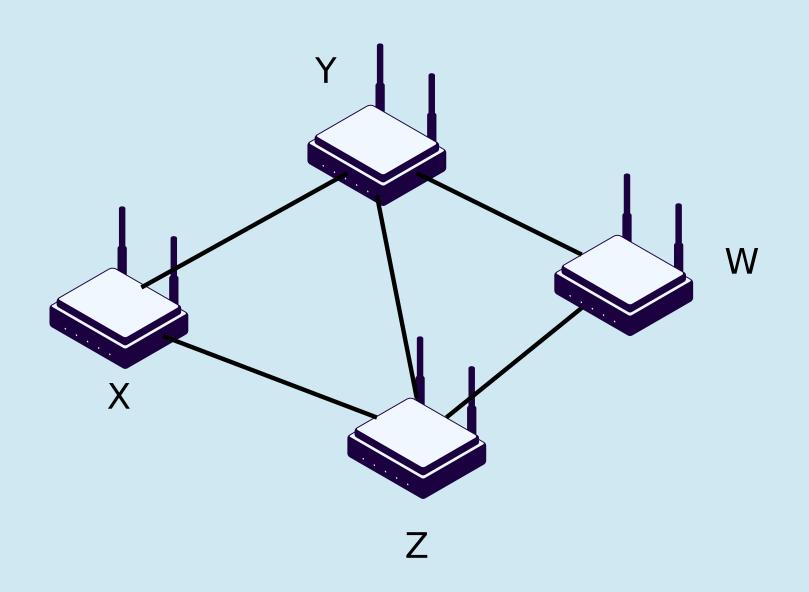


## LIMITATION

It can only stop loops between two directly connected routers.

If the loop involves more than two routers, this method does not work.

## PROBLEM 11.A

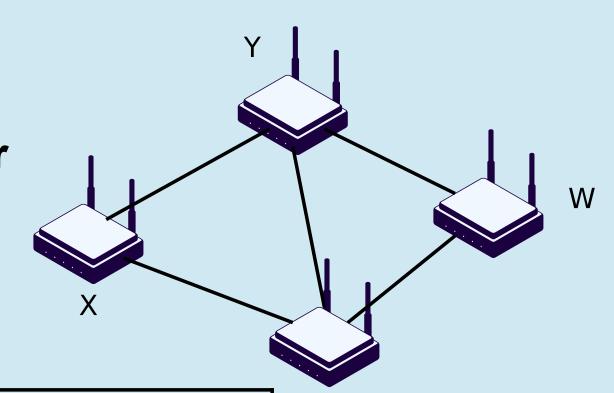


Poison reversed is used in Distance vector algorithm

 When the distance vector routing is stabilized, router w, y, and z inform their distances to x to each other. What distance values do they tell each other?

# SOLUTION 11.A

 When the distance vector routing is stabilized, router w, y, and z inform their distances to x to each other.
 What distance values do they tell each other?

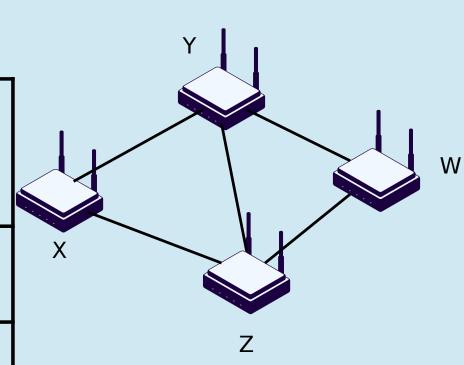


Source Router	Destinatio n	Cost	Next Hop	Shortest Path	
W	X	5	Y	W->Y->X	
Y	X	4	X	Y->X	
Z	X	6	Y	Z->Y->X	

Z

# PROBLEM 11.A

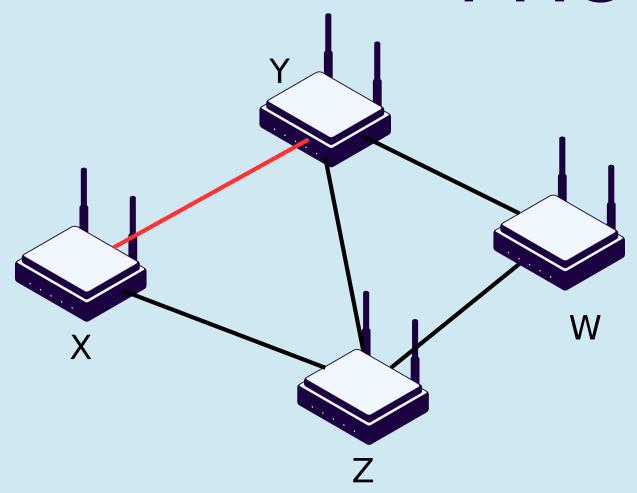
Source Router	Destinatio n	Cost	Next Hop	Shortest Path
W	X	5	Y	W->Y->X
Y	X	4	X	Y->X
Z	X	6	Y	Z->Y->X



## Router w, y, and z inform their distances to x to

Router Z	each other: Informs W: Dz(X)=6	Informs Y: Dz(X)=∞	
Router W	Informs Y, Dw(X)=∞	Informs z, Dw(X)=5	
Router Y	Informs W, Dy(X)=4	Informs z, Dy(X)=4	

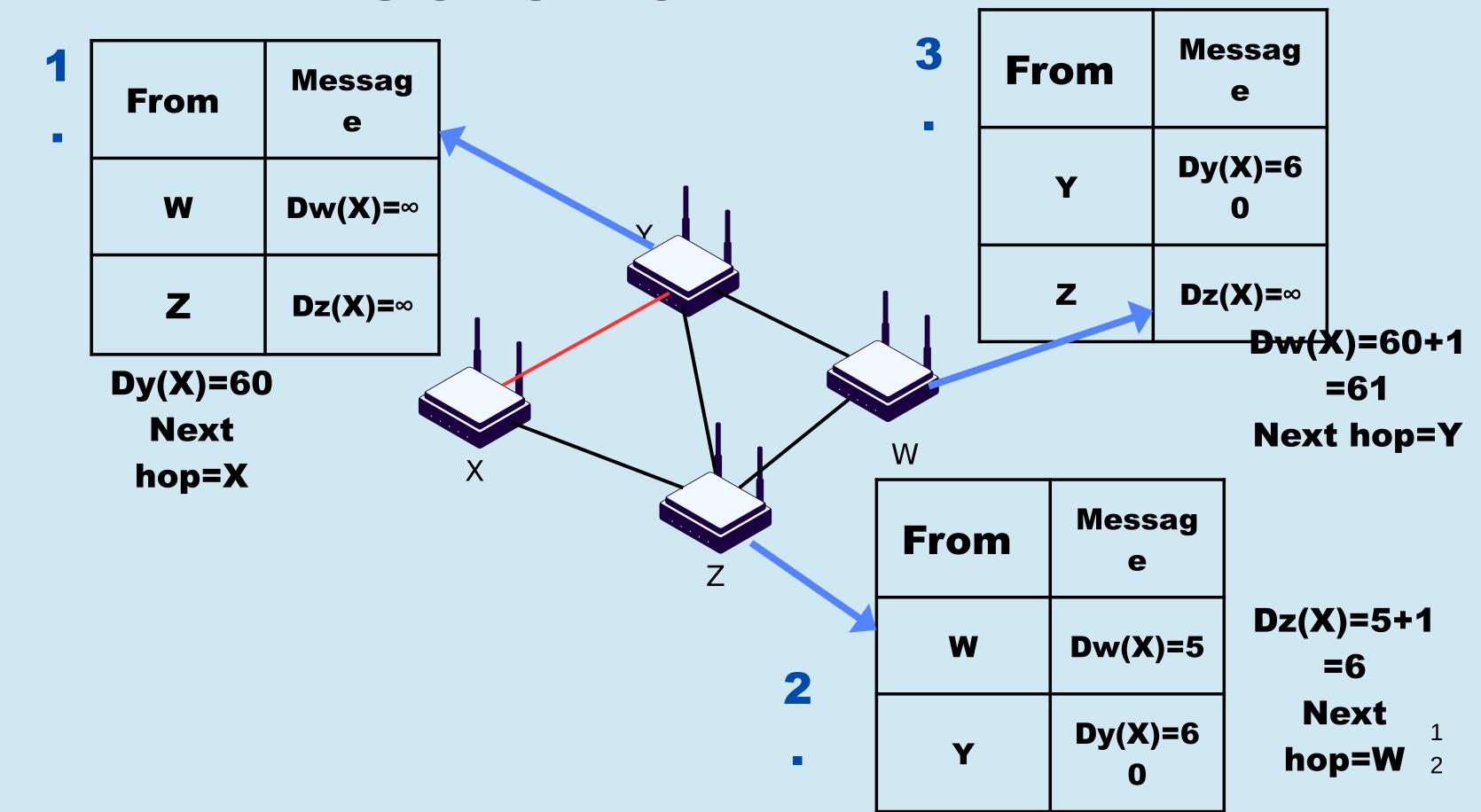
## PROBLEM 11.B



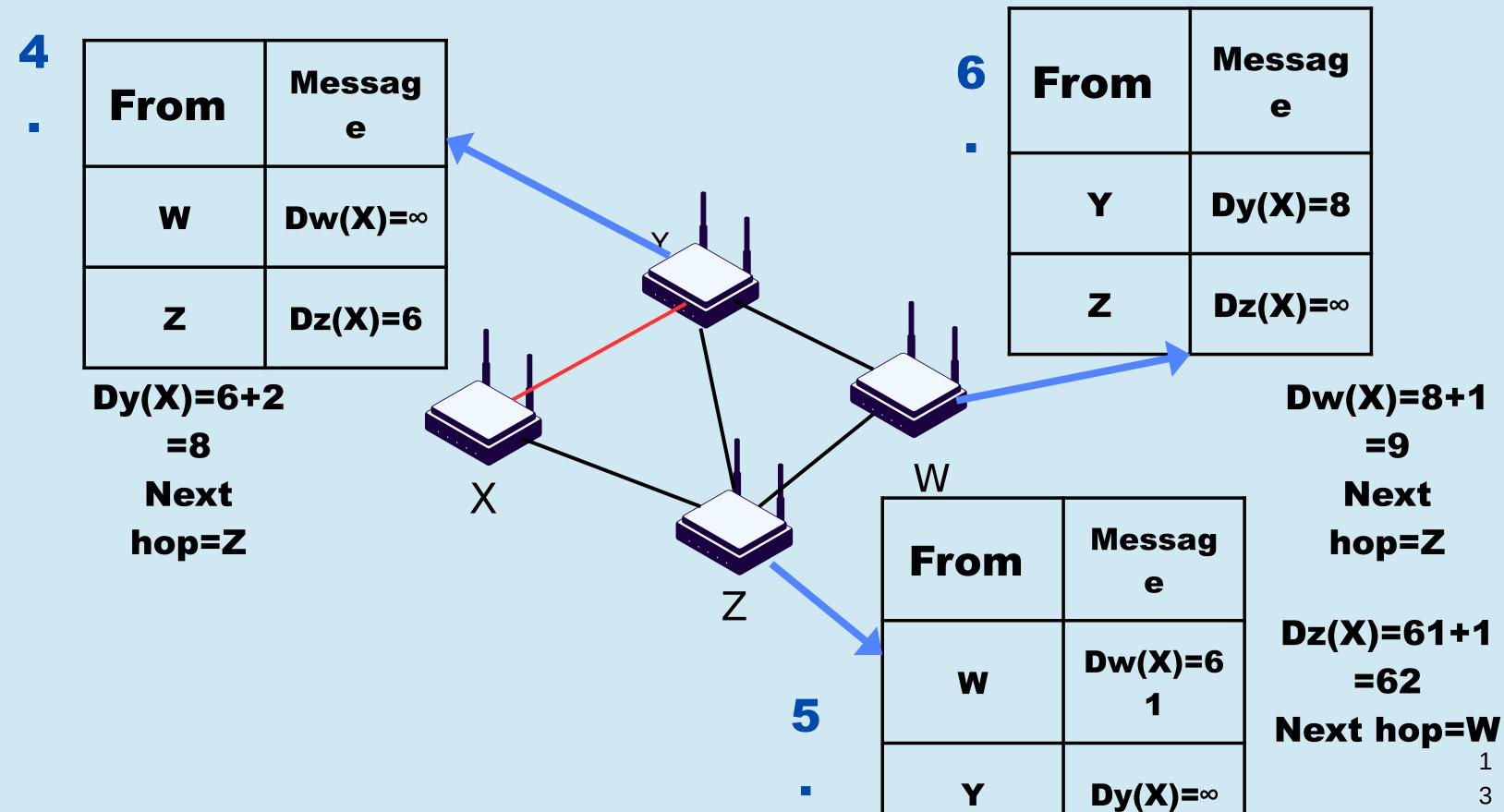
Poison reversed is used in Distance vector algorithm

 Now suppose that the link cost between x and y increases to 60. Will there be a count-to-infinity problem even if poisoned reverse is used?
 Why or why not? If there is a count-to-infinity problem, then how many iterations are needed for the distance-vector routing to reach a stable state again?

# SOLUTION 11.B



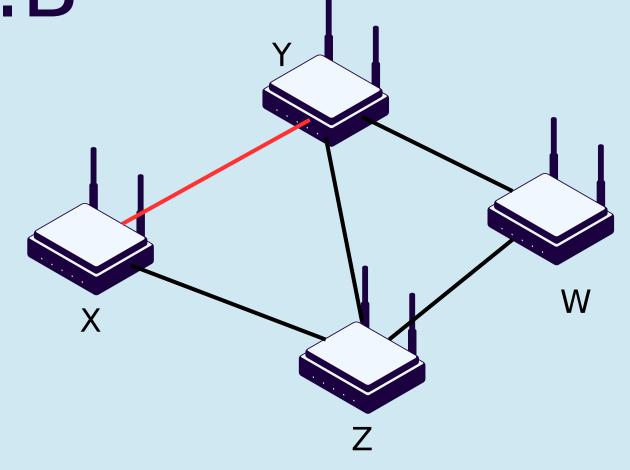
# SOLUTION 11.B



# SOLUTION 11.B

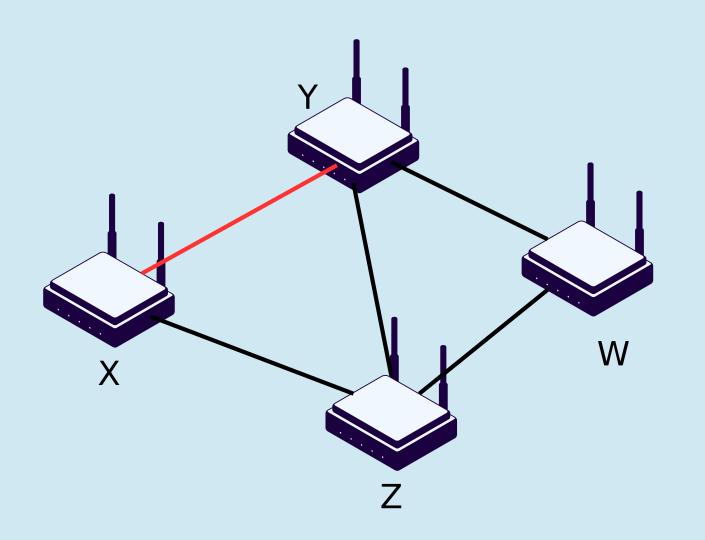
After 44 iterations, table converged:

From/To	X	Y	Z	W
X	0	52	50	51
Y	52	0	2	1
Z	50	2	0	1
W	51	1	1	0



Count to infinity problem occured although poisoned reverse is used!!

# PROBLEM 11.C



Poison reversed is used in Distance vector algorithm

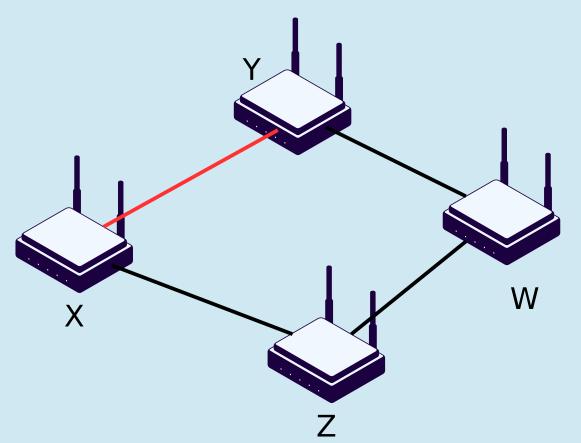
 How do you modify c(y,z) such that there is no count-to-infinity problem at all if c(y,x) changes from 4 to 60?

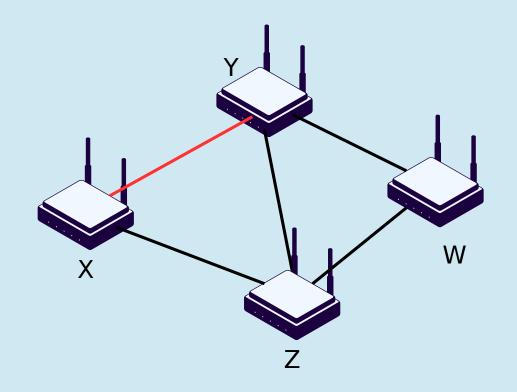
# SOLUTION 11.C

 How do you modify c(y,z) such that there is no count-to-infinity problem at all if c(y,x) changes from 4 to 60?

#### Answer:

We will cut the cut the link between y and z.





# ANY QUESTION?

# THANKYOU