

Noise Extension: Disco

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Revision 3, 2017-11-14, unofficial/unstable

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1. Introduction

1.1. Motivation

Noise is a framework for crypto protocols based on Diffie-Hellman key agreement. One of its most interesting property is that every new message depends on all

the previous ones. This is done by continuously hashing messages being sent and received, as well as continuously deriving new keys based on the continuous hash and the previous keys. This interesting property stops at the end of the handshake.

Strobe is a protocol framework based on a duplex construction. It naturally benefits from the same property, effectively absorbing every operation to influence the next ones. The Strobe specification is comparable to Noise, but focusing on the symmetric part of a protocol. By merging both protocols into one, Disco achieves the following goals:

- The Noise specification can be greatly simplified by removing all the symmetric cryptographic algorithms and symmetric objects. These can be replaced by a single Strobe object.
- Implementations of Noise with the Disco extension will consequently greatly benefit from this simplification, allowing for a drastic reduction of the codebase, facilitating security audits.
- Messages will continue to rely on every previous messages that were sent or received, even in the symmetric part of the protocol.
- The Strobe functions will allow for more flexible and complex symmetric protocols following the handshake.
- Implementations of Noise with the Disco extension will also benefit from the other Strobe functions, which provide on top of a single primitive the following functions: generation of random numbers, derivation of keys, hashing, encryption and authentication.

1.2. How to Read This Document

This specification is an extension of the Noise protocol framework revision 32. It relies for the most part on Noise's specification, while heavily modifying its foundations. Major changes are listed in the next section.

To implement the Disco extension, a Strobe implementation respecting the functions of the section 3 of this document is required. None of the cipher and hash functions of Noise are required. Furthermore, the CipherState is not necessary while the SymmetricState has been simplified by Strobe calls. When implementing Noise with the Disco extension, simply ignore the CipherState section of Noise and implement the SymmetricState described in section 5 of this document. For advanced features, refer to section 6.

1.3. Change log

this section will be removed in the final document

draft-3:

- The specification now extends Noise draft-33.

- A `isSetKey` boolean value was added to the `SymmetricState`, it is set when `MixKey()` is called.
- `EncryptAndHash()` and `DecryptAndHash()` now look for the value of `isSetKey` and branch to `send_CLR()` and `recv_CLR()` if the boolean value is set to `false`.
- a `StrobeState` object has been introduced to formalize the integration of Strobe in Disco.
- Added the missing `GetHandshakeHash()` function to the `SymmetricState`.
- removed the `TAGLEN` field and set it to 16 everywhere. Following Noise's way of defining the tag length.
- Half-duplex protocols are introduced in Advanced features.
- Out-of-order protocols are introduced in Advanced features, along with `DiscoSecureChannel` objects.

draft-2:

- The `SymmetricState` object has been simplified with Strobe's calls (instead of modifying the `HandshakeState`).

draft-1:

- Protocol names don't have the symmetric algorithms, but instead the version of Strobe.
- The `CipherState` object has been removed.
- The `Handshake` object makes calls to Strobe functions, affecting a unique Strobe state.
- The `Handshake` returns two Strobe states.
- The document extends Noise draft-32.

2. Protocol naming

The name of a Noise protocol extended with Disco follows the same convention, but replaces the symmetric cryptographic algorithms by the version of Strobe used:

`Noise_[PATTERN]_[KEYEXCHANGE]_STROBEvX.Y.Z`

For example, with the current version of Strobe being STROBEv1.0.2:

`Noise_XX_25519_STROBEv1.0.2`

3. The StrobeState object

A `StrobeState` depends on a Strobe object (as defined in section 5 of the Strobe Specification) as well as the following associated constant:

- **StrobeR**: The blocksize of the Strobe state (computed as $N - (2 \cdot \text{sec})/8 - 2$, see section 4 of the Strobe specification).

While a Strobe object responds to many functions (see Strobe’s specification), only the following ones need to be implemented in order for the Disco extension to work properly:

InitializeStrobe(protocol_name): Initialize the Strobe object with a custom protocol name.

KEY(key): Permutes the Strobe’s state and replaces the new state with the key.

PRF(output_len): Permutes the Strobe’s state and removes **output_len** bytes from the new state. Outputs the removed bytes to the caller.

send_ENC(plaintext): Permutes the Strobe’s state and XOR the plaintext with the new state to encrypt it. The new state is replaced by the resulting ciphertext, while the resulting ciphertext is output to the caller.

recv_ENC(ciphertext): Permutes the Strobe’s state and XOR the ciphertext with the new state to decrypt it. The new state is replaced by the ciphertext, while the resulting plaintext is output to the caller.

AD(additionalData): Absorbs the additional data in the Strobe’s state.

send_MAC(output_length): Permutes the Strobe’s state and retrieves the next **output_length** bytes from the new state.

recv_MAC(tag): Permutes the Strobe’s state and compare (in **constant-time**) the received tag with the next 16 bytes from the new state.

RATCHET(length): Permutes the Strobe’s state and set the next **length** bytes from the new state to zero.

The following **meta** functions:

meta_AD(additionalData): XOR the additional data in the Strobe’s state.

The following function which is not specified in Strobe:

Clone(): Returns a copy of the Strobe state.

4. Modifications to the Handshake State

Processing the final handshake message via **WriteMessage()** and **ReadMessage()** now returns two new **StrobeState** objects by calling **Split()**. The first for encrypting transport messages from initiator to responder, and the second for messages in the other direction. At this point the **StrobeState** of the **SymmetricState** should not be deleted as it is the first **StrobeState** object returned by **Split()** (and will be used by the initiator to encrypt messages).

The peers can then encrypt (resp. decrypt) messages by calling `send_ENC` followed by `send_MAC` (resp. `recv_ENC` followed by `recv_MAC`) on the relevant `StrobeState` object.

5. Modifications to the Symmetric State

A `SymmetricState` object contains:

- **StrobeState**: a Strobe state responding to the functions mentioned in the previous section.
- **isKeyed**: a boolean value indicating if a Diffie-Hellman key exchange has already occurred.

A `SymmetricState` responds to the following functions:

InitializeSymmetric(protocol_name): Calls `InitializeStrobe(protocol_name)` on the Strobe state.

MixKey(input_key_material): Calls `AD(input_key_material)` on the Strobe state. It then sets `isKeyed` to `true`.

MixHash(data): Calls `AD(data)` on the Strobe state.

MixKeyAndHash(input_key_material): Calls `AD(input_key_material)` on the Strobe state.

GetHandshakeHash(): Calls `PRF(32)`. This function should only be called at the end of a handshake, i.e. after the `Split()` function has been called. This function is used for channel binding, as described in Section 11.2 of the Noise specification.

EncryptAndHash(plaintext): Returns a ready to be sent payload to the caller by following these steps:

- If `isKeyed` is set to `false`, call `send_CLR(plaintext)` and return the `plaintext`.
- Call `send_ENC(plaintext)` followed by `send_MAC(16)` on the Strobe state. Return the concatenation of both results to the caller.

DecryptAndHash(ciphertext): Returns the received payload by following steps:

- If `isKeyed` is set to `false`, call `recv_CLR(ciphertext)` and return the `ciphertext`.
- Otherwise, check that the length of the received `ciphertext` is at least 16 bytes. If it is not, return an error to the caller and abort the handshake.
- Call `recv_ENC(ciphertext[:-16])` and store the result in a `plaintext` buffer.

- Call `recv_MAC(ciphertext[-16:])` on the Strobe state. If `recv_MAC` returns `false`, the peer must return an error to the caller and abort the connection. Otherwise return the `plaintext` buffer.

Split(): Returns a pair of Strobe states for encrypting transport messages by executing the following steps:

- Let `s1` be the Strobe state and `s2` the result returned by `Clone()`.
- Calls `meta_AD("initiator")` on `s1` and `meta_AD("responder")` on `s2`.
- Calls `RATCHET(StrobeR)` on `s1` and on `s2`.
- Returns the pair `(s1, s2)`.

6. Modifications to Advanced Features

6.1 Channel Binding

Right before calling `Split()`, a binding value could be obtained from the `StrobeState` by calling `PRF()`.

6.2 Rekey

To enable this, Strobe supports a `RATCHET()` function.

6.3 Out-of-order transport messages

In order to build out-of-order protocols out of Disco, the `Split()` function must return nonce-based objects. For this, the `Split()` function is modified in the next section to return a pair of `DiscoSecureChannel` objects which are defined in the section following it.

Transport messages are then encrypted and decrypted by calling `Encrypt()` and `Decrypt()` on the relevant `DiscoSecureChannel`.

6.3.1 Modifications to the `Split()` function

Modify the `Split()` function to add the following steps before returning the pair `(s1, s2)` of `Strobe` objects:

- Create two `DiscoSecureChannel` named `d1` and `d2`.
- Associate `s1` (resp. `s2`) to `d1` (resp. `d2`).
- Return the pair `(d1, d2)`.

6.3.2 The DiscoSecureChannel object

A `DiscoSecureChannel` can encrypt and decrypt data based on its associated `StrobeState` object as well as the following variable:

- `n`: An 8-byte (64-bit) unsigned integer nonce.

A `DiscoSecureChannel` responds to the following functions:

Encrypt(plaintext):

- If `n` is equal to $2^{64}-1$ the function returns an error to the caller and aborts the Disco session.
- Create a new `StrobeState` by calling `Clone()` on the `StrobeState` object.
- Call `AD(n)` on the new `StrobeState`.
- Call `send_ENC(plaintext)` on the new `StrobeState` and add the result to a `ciphertext` buffer.
- Call `send_MAC(16)` on the new `StrobeState` and add the result to the `ciphertext` buffer.
- Increment the nonce `n` and discard the new `StrobeState` object.
- Return the `ciphertext` buffer containing the encrypted data.

Decrypt(ciphertext)::

- Check that the length of the received `ciphertext` is at least 16 bytes. If it is not, return an error to the caller and abort the session.
- If `n` is equal to $2^{64}-1$ the function returns an error to the caller and aborts the Disco session.
- Create a new `StrobeState` by calling `Clone()` on the `StrobeState` object.
- Call `AD(n)` on the new `StrobeState`.
- Call `recv_ENC(ciphertext[:-16])` on the new `StrobeState` and store the result in a `plaintext` buffer.
- Call `recv_MAC(ciphertext[-16:])` on the new `StrobeState`, if the function returns `false`, return an error to the caller and abort the Disco session.
- Increment the nonce `n` and discard the new `StrobeState` object.
- Return the `plaintext` buffer containing the decrypted data.

IntroduceForwardSecrecy(): calls `RATCHET(64)` on the `Strobe` Object.

6.4 Half-duplex protocols

To use Disco in half-duplex mode, modify `Split()` to return the `StrobeState` without modifications.

The same security considerations from the Noise specification applies to this section: if the two peers do not properly take turns to write and read on the channel, the protocol will fail catastrophically.

7. Security Considerations

The same security considerations that apply to both Noise and Strobe are to be considered.

8. Acknowledgements

Thanks to Trevor Perrin and Mike Hamburg for being the foundations and main help in building this specification.