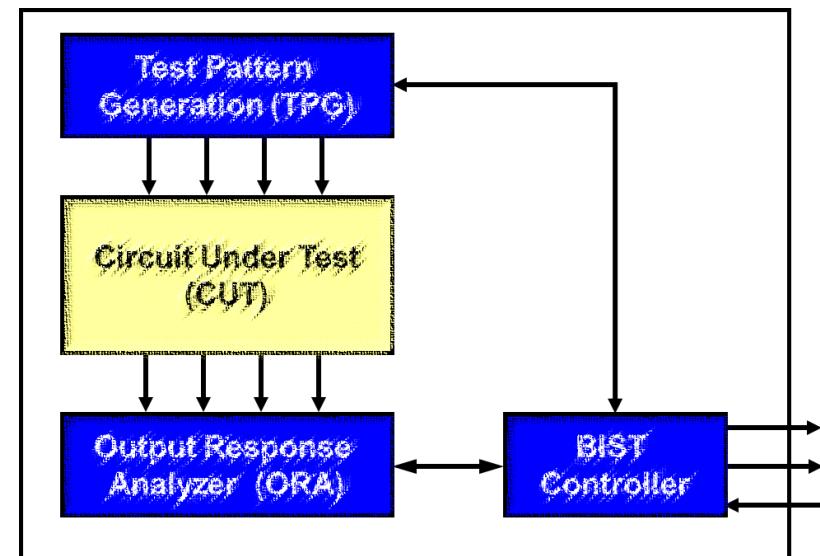


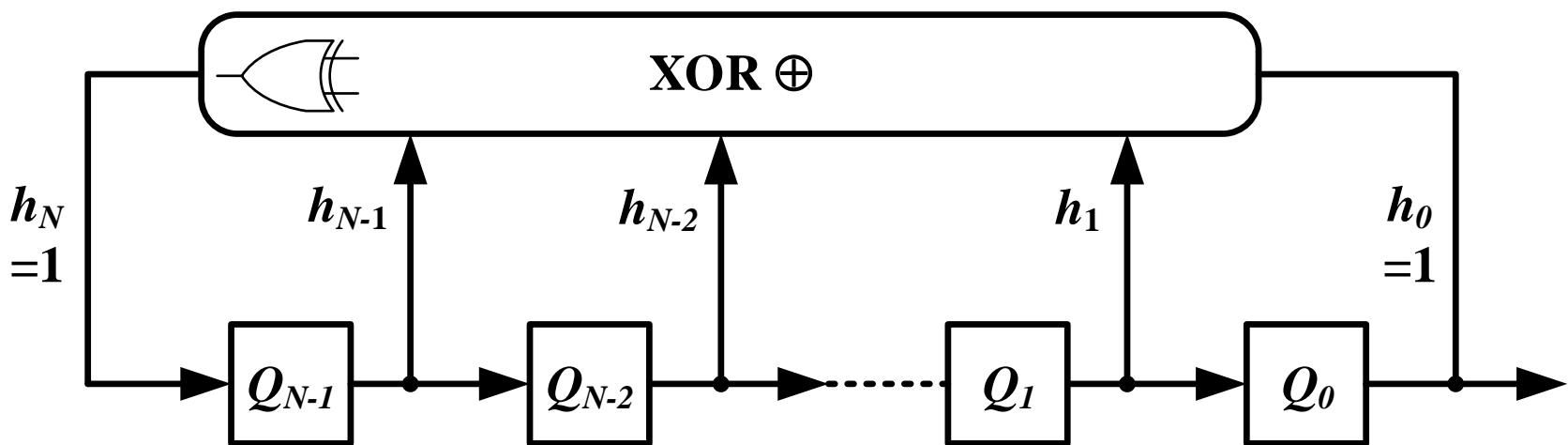
BIST Part1 - TPG

- Introduction
- Test Pattern Generation (TPG)
 - ◆ Deterministic: ROM, Algorithm, Counter
 - ◆ Pseudo Random
 - * Linear Feedback Shift Register, LFSR (1977)
 - Two types of LFSR
 - Design of LFSR
 - * Cellular Automata, CA (1984)



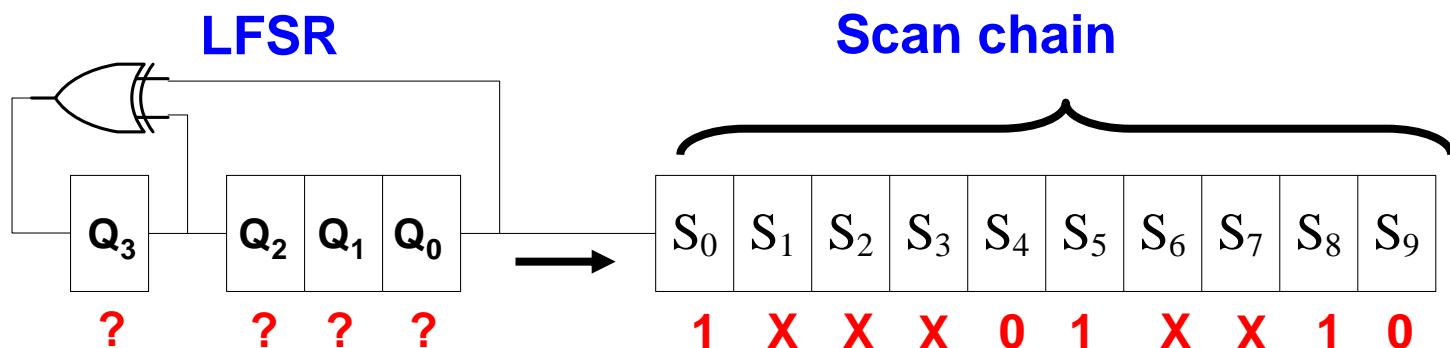
Design of LFSR

- Questions to answer
 - ◆ How to find seed?
 - ◆ What is LFSR degree?
 - ◆ What polynomial?



How to Find a Seed?

- Given LFSR and test pattern, find initial state of LFSR (seed)
- Example:
 - LFSR feeds scan chain inputs
 - $S_0 \sim S_9$ are test pattern (specified by ATPG)
 - $Q_3 \sim Q_0$ are seeds

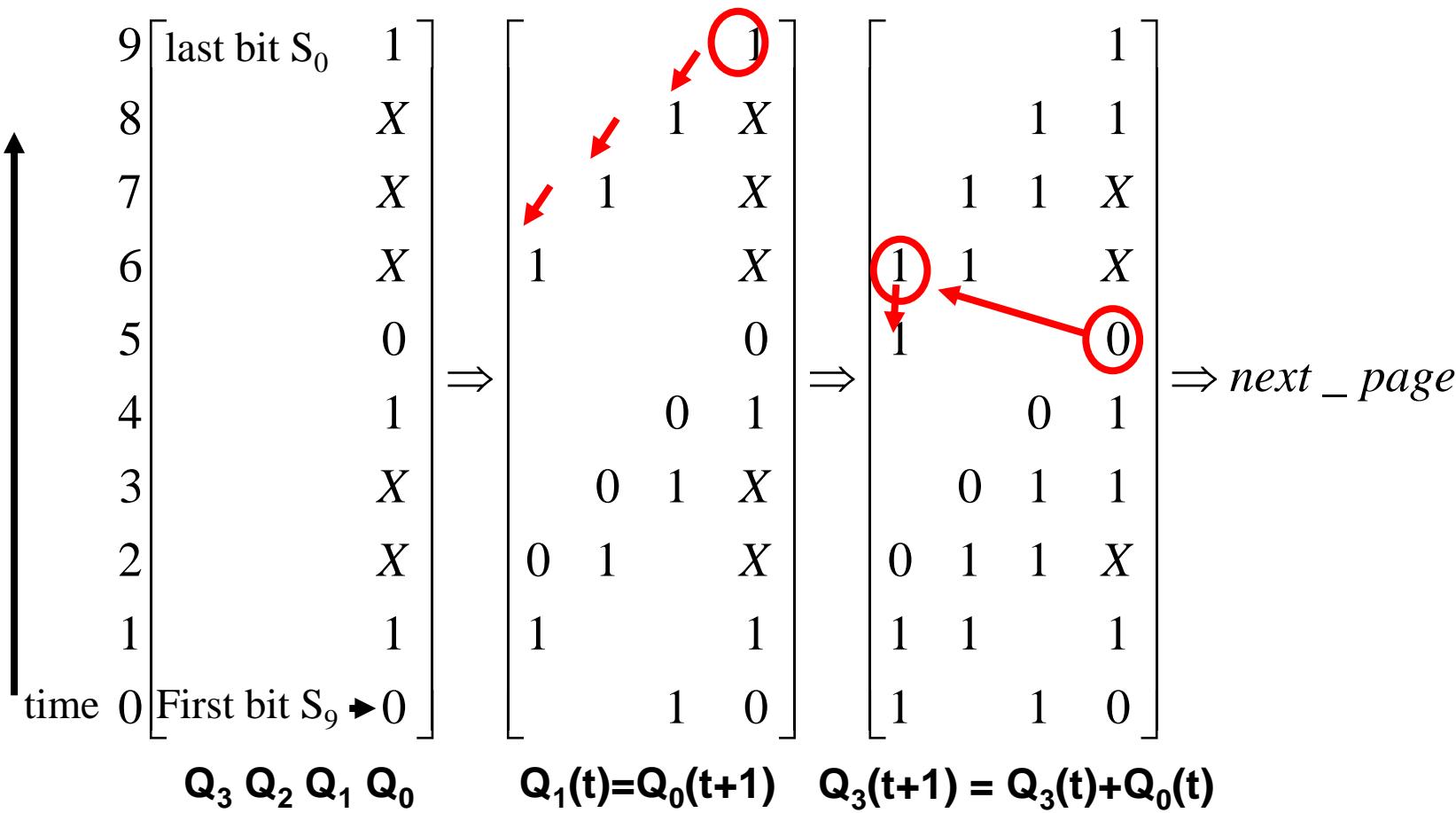
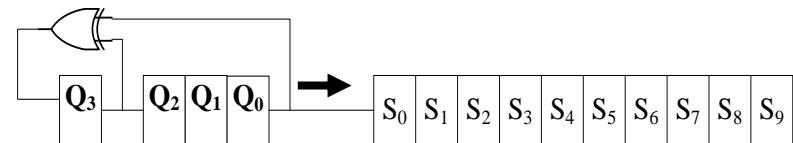


- Two methods:
 - Cycle-by-cycle tracing
 - System of linear equations

Cycle-by-cycle Tracing (1)

- Test pattern desired

◆ $S_0 S_1 \dots S_9 = 1 X X X 0 1 X X 1 0$



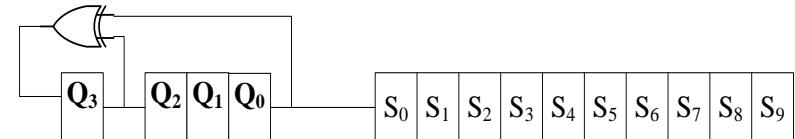
Cycle-by-cycle Tracing (2)

Solve $X_3=0, X_2=1$

$$\Rightarrow \begin{bmatrix} & & & 1 \\ & & 1 & 1 \\ & 1 & 1 & X_3 \\ 1 & 1 & X_3 & X_2 \\ \textcircled{1} & X_3 & X_2 & 0 \\ X_3 & X_2 & 0 & \textcircled{1} \\ X_2 & 0 & 1 & 1 \\ 0 & 1 & 1 & X_1 \\ 1 & 1 & X_1 & 1 \\ 1 & X_1 & 1 & 0 \end{bmatrix} \Rightarrow$$

Solve $X_1=1$

$$\begin{bmatrix} & & & 1 \\ & & 1 & 1 \\ & 1 & 1 & 0 \\ 1 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ \textcircled{1} & 0 & 1 & 1 \\ 0 & 1 & 1 & X_1 \\ 1 & 1 & X_1 & 1 \\ 1 & X_1 & 1 & 0 \end{bmatrix}$$



$$\begin{bmatrix} & & & 1 \\ & & 1 & 1 \\ & 1 & 1 & 0 \\ 1 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ \textcircled{1} & 0 & 1 & 1 \\ 1 & 0 & 1 & X_1 \\ 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \end{bmatrix} \Rightarrow \begin{bmatrix} & & & 1 \\ & & 1 & 1 \\ & 1 & 1 & 0 \\ 1 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ \textcircled{1} & 1 & 1 & 0 \end{bmatrix}$$

Seed = 1110

This is Slow! Can We Do Better?

System of Linear Equations (1)

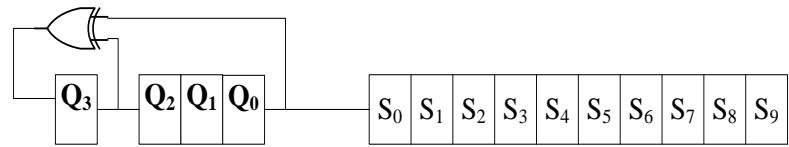
- Initial conditions

- $S_9 = Q_0$
- $S_8 = Q_1$
- $S_7 = Q_2$
- $S_6 = Q_3$

- $S_n = S_{n+4} \oplus S_{n+1}$

- $S_5 = S_9 \oplus S_6 = Q_3 \oplus Q_0$
- $S_4 = S_8 \oplus S_5 = Q_3 \oplus Q_1 \oplus Q_0$
- $S_3 = S_7 \oplus S_4 = Q_3 \oplus Q_2 \oplus Q_1 \oplus Q_0$
- $S_2 = S_6 \oplus S_3 = Q_2 \oplus Q_1 \oplus Q_0$
- $S_1 = S_5 \oplus S_2 = Q_3 \oplus Q_2 \oplus Q_1$
- $S_0 = S_4 \oplus S_1 = Q_2 \oplus Q_0$

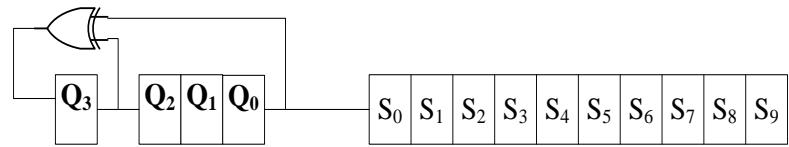
- 4 variables, 10 equations



$$\begin{bmatrix} 0 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \\ 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} Q_3 \\ Q_2 \\ Q_1 \\ Q_0 \end{bmatrix} = \begin{bmatrix} S_0 = 1 \\ S_1 = X \\ S_2 = X \\ S_3 = X \\ S_4 = 0 \\ S_5 = 1 \\ S_6 = X \\ S_7 = X \\ S_8 = 1 \\ S_9 = 0 \end{bmatrix}$$

System of Linear Equations (2)

- $S_9 = 0 = Q_0$
- $S_8 = 1 = Q_1$
- $S_5 = 1 = Q_0 \oplus Q_3 \Rightarrow Q_3 = 1$
- $S_4 = 0 = Q_3 \oplus Q_1 \oplus Q_0$
 - ◆ consistent !
- $S_0 = 1 = Q_2 \oplus Q_0 \Rightarrow Q_2 = 1$
- Solution :
 - ◆ Seed = $[Q_3 \ Q_2 \ Q_1 \ Q_0] = 1110$



$$\left[\begin{array}{cccc} 0 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \\ 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{array} \right] \begin{bmatrix} Q_3 \\ Q_2 \\ Q_1 \\ Q_0 \end{bmatrix} = \begin{bmatrix} S_0 = 1 \\ S_1 = X \\ S_2 = X \\ S_3 = X \\ S_4 = 0 \\ S_5 = 1 \\ S_6 = X \\ S_7 = X \\ S_8 = 1 \\ S_9 = 0 \end{bmatrix}$$

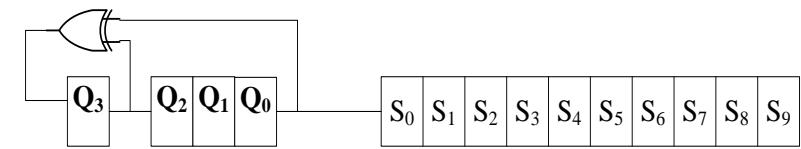
Automate This Process?

Derive Linear Equations (1)

- Start with identity matrix I
 - $t = [h_3 \ h_2 \ h_1 \ h_0]$
- Iteratively bottom up, $k=4 \sim 10$
 - $\text{Row}_k = t \times (\text{Row}_{k-4} \sim \text{Row}_{k-1})$
 - $k++$
- Example
 - $t = [h_3 \ h_2 \ h_1 \ h_0] = [1 \ 0 \ 0 \ 1]$
 - $t \times I_1 = [1 \ 0 \ 0 \ 1]$
 - $t \times I_2 = [1 \ 0 \ 1 \ 1]$
 -
 - Eventually, 10 rows, 4 columns

$$[1 \ 0 \ 0 \ 1] \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix},$$

$$\rightarrow \begin{bmatrix} 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$



$$\dots \rightarrow \begin{bmatrix} 1 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \rightarrow \dots \rightarrow \begin{bmatrix} 0 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \\ 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

Solve Linear Equations (1)

- **Gauss-Jordan elimination**, but **Mod-2 addition**
- Example
 - ◆ Append specified bits to last column (*Augment matrix*)
 - ◆ Remove useless rows with X
 - ◆ Exchange row 1 with row 2

$$\left[\begin{array}{cccc|c} 0 & 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 & X \\ 0 & 1 & 1 & 1 & X \\ 1 & 1 & 1 & 1 & X \\ 1 & 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 & X \\ 0 & 1 & 0 & 0 & X \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \end{array} \right] \xrightarrow{\text{Row Operations}} \left[\begin{array}{ccccc} 0 & 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \end{array} \right] \xrightarrow{\text{Row Operations}} \left[\begin{array}{ccccc} 1 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \end{array} \right]$$

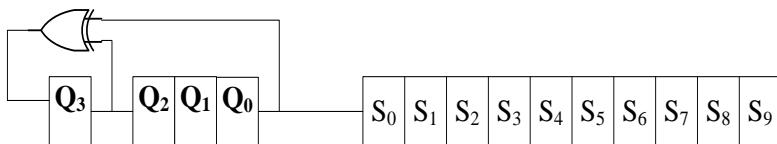
Solve Linear Equations (2)

- Pivot column #1 (add row 1 to row 3)
- Pivot column #2 (do nothing)
- Pivot column #3 (add row3 to row1; add row3 to row4)
- Pivot column #4 (add row5 to row1, add row 5 to row2)
- Solution: seed = $[Q_3 \ Q_2 \ Q_1 \ Q_0] = [1 \ 1 \ 1 \ 0]$

$$\begin{bmatrix} 1 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \end{bmatrix} \xrightarrow{\text{Row 1} \rightarrow R_1 + R_3} \begin{bmatrix} 1 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \end{bmatrix} \xrightarrow{\text{Row 3} \rightarrow R_1 + R_3} \begin{bmatrix} 1 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \end{bmatrix} \xrightarrow{\text{Row 3} \rightarrow R_1 + R_4} \begin{bmatrix} 1 & 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 \end{bmatrix} \xrightarrow{\text{Row 5} \rightarrow R_1 + R_5} \begin{bmatrix} 1 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 \end{bmatrix}$$

Summary

- Number of **equations** depends on **number of care bits**
- Number of **variables** = **LFSR degree**
- If more care bits than degree,
 - ♦ May not be solvable every time
 - ♦ This time we are lucky :)



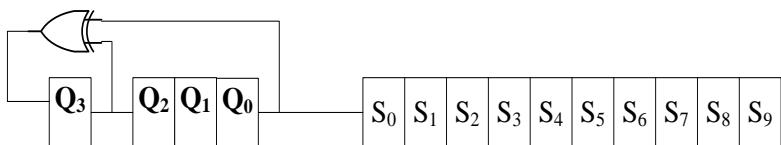
$$\begin{bmatrix} 0 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \\ 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \end{bmatrix} \begin{bmatrix} Q_3 \\ Q_2 \\ Q_1 \\ Q_0 \end{bmatrix} = \begin{bmatrix} S_0 = 1 \\ S_1 = X \\ S_2 = X \\ S_3 = X \\ S_4 = 0 \\ S_5 = 1 \\ S_6 = X \\ S_7 = X \\ S_8 = 1 \\ S_9 = 0 \end{bmatrix}$$

Can We Guarantee Solution Exists?

FFT

- Q: What should we do if we find no solution?

- ◆ 1
- ◆ 2
- ◆ 3



$$\begin{bmatrix} 0 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \\ 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \end{bmatrix} \begin{bmatrix} Q_3 \\ Q_2 \\ Q_1 \\ Q_0 \end{bmatrix} = \begin{bmatrix} S_0 = 1 \\ S_1 = X \\ S_2 = X \\ S_3 = X \\ S_4 = 0 \\ S_5 = 1 \\ S_6 = X \\ S_7 = X \\ S_8 = 1 \\ S_9 = 0 \end{bmatrix}$$