

## 4- Tower Of Hanoi

### Problem Description

There are eight disks of different sizes and four pegs. Initially, all the disks are on the first peg in order of size, the largest on the bottom and the smallest on the top. It's required to use divide and conquer algorithm to transfer all the disks to another peg by a sequence of moves. The constraints are that only one disk can be moved at a time, and it is forbidden to place a larger disk on top of a smaller one.

Also it's required to use dynamic programming algorithm to transfer all the disks to another peg in 33 moves

### Detailed Assumptions

- The disks are labeled with consecutive integers starting from 1. Each disk is represented by an object of Disk class, which has a disk\_num attribute.
- The moves variable is used to keep track of the total number of moves performed during solving the Tower of Hanoi Problem.
- The Tower of Hanoi problem is being solved optimally, it aims to minimize the number of moves required to transfer the disks from the starting peg to the end peg.
- The line that assigns k to the optimal value of k that is resulted in the function of getOptimalDiskNum is commented in the code if divide and conquer method is used and the value of k will be assigned to  $n/2$  where n is the number of disks
- The line that assigns k to  $n/2$  is commented in case of using dynamic programming and the value of k will be assigned to the value resulted in OptimalK array in the getOptimalDiskNum function

### Detailed Solution (pseudocode and steps description)

1. Initialize an array dp of size  $n+1$  to store the number of moves required to transfer the disks from the starting peg to the end peg
2. Initialize an array optimalK of size  $n+1$  to store the optimal value of k that will be used to divide the disks in dynamics programming technique while in divide and conquer the disks will be divided into halves
3. The code keep divide the disks into sub-problems until reaches the base case to transfer the disks to another peg

Both divide and conquer algorithm and dynamic programming algorithm are implemented in the following pseudocode

```

1  /*
2  * Description : Function to print the number of moves required to solve the puzzle
3  * Input : number of disks, initial peg, last peg
4  * Output : none
5  */
6  function printMove(disk_num, from, to):
7      print "Move Disk " + disk_num + " from peg " + from + " to peg " + to
8
9  /*
10 * Description : Function to save the number of moves to solve the puzzle in an array
11 * Input : number of disks
12 * Output : none
13 */
14 function getOptimalDiskNum(n):
15     // Arrays to store the minimum moves and corresponding k values
16     create an array dp of size n + 1
17     create an array optimalK of size n + 1
18
19     // Base cases for dp array
20     dp[0] = 0 //number of moves in case of zero disks
21     dp[1] = 1 //number of moves in case of one disks
22     dp[2] = 3 //number of moves in case of two disks
23
24     //the case of more than two disks
25     // Filling the dp and optimalK arrays
26     for i = 3 to n:
27         minMoves = infinity
28
29         for k = 1 to i - 1:
30             moves = 2 * dp[k] + (2^(i-k)) - 1
31
32             if moves < minMoves:
33                 minMoves = moves
34                 optimalK[i] = i - k // Update the optimal k for this i using
35                                     // the result of k stored previously in the
36                                     // optimalK array */
37
38             dp[i] = minMoves //storing the min number of moves in the corresponding index that
39                             // indicates the number of disks */
40
41 /*
42 * Description : Function to solve the tower of hanoi using only three pegs
43 * Input : disks - disks array
44 *         start - starting rod
45 *         aux_peg - the auxiliary rod
46 *         end - the target rod
47 *
48 * Output : none
49 */
50 function hanoi3(disks, start, aux_peg, end):
51     //base case of number of disks equals 0
52     if length of disks == 0:
53         return
54     //base case of number of disks equals 1
55     if length of disks == 1:
56         printMove(disks[0].disk_num, start, end)
57         return
58     //the case of number of disks more than one disk
59     n = length of disks
60     rem_disks = copy of the first n-1 disks
61
62     hanoi3(rem_disks, start, end, aux_peg)
63     printMove(disks[n-1].disk_num, start, end)
64     hanoi3(rem_disks, aux_peg, start, end)

```

```

65
66
67  /*
68  * Description : Function to solve the tower of hanoi using four pegs
69  * Input : n - number of disks
70  *         start - starting rod
71  *         aux_peg_1 - the first auxiliary rod
72  *         aux_peg_2 - the second auxiliary rod
73  *         end - the target rod
74  * Output : none
75  */
76
77 function hanoi4(disks, start, aux_peg_1, aux_peg_2, end):
78     n = length of disks
79     //base case of number of disks equals 0
80     if n == 0:
81         return
82     //base case of number of disks equals 1
83     if n == 1:
84         printMove(disks[0].disk_num, start, end)
85         return
86     //base case of number of disks equals 2
87     if n == 2:
88         printMove(disks[0].disk_num, start, aux_peg_2)
89         printMove(disks[1].disk_num, start, end)
90         printMove(disks[0].disk_num, aux_peg_2, end)
91         return
92
93     // one of the two following values of k is used according to the chosen running algorithm that the two cases are separate
94     // the value of k in case of divide and conquer algorithm is that the number of disks is divided into two sub-
95     // problems
96     k=n/2
97
98     // the value of k in case of dynamic programming
99     if optimalK is null:
100         getOptimalDiskNum(n)
101
102     k = optimalK[n]
103     rem_disks = copy of the first n-k disks
104     fixed_disks = copy of the last k disks
105
106     hanoi4(rem_disks, start, aux_peg_2, end, aux_peg_1)
107     hanoi3(fixed_disks, start, aux_peg_2, end)
108     hanoi4(rem_disks, aux_peg_1, start, aux_peg_2, end)
109
110
111 function main():
112     n = 8
113     create an array disks of size n
114
115     for i = 0 to n-1:
116         disks[i] = new Disk(i + 1)
117
118     hanoi4(disks, 'A', 'B', 'C', 'D')
119     print "no of moves = " + moves
120
121 class Disk:
122     //attribute of object from class Disk to determine the number of the disk
123     int disk_num
124     Disk(disk_num):
125         this.disk_num = disk_num

```

## Complexity Analysis

1. getOptimalDiskNum function:
  - Time Complexity:  $O(n^2)$
  - Space Complexity:  $O(n)$
2. hanoi3 function:
  - Time Complexity:  $O(2^n)$
  - Space Complexity:  $O(n)$
3. hanoi4 function:
  - Time Complexity:  $O(2^n)$
  - Space Complexity:  $O(n)$
4. main function:
  - Time Complexity:  $O(2^n)$
  - Space Complexity:  $O(n)$

The resulted Time Complexity:  $O(2^n)$

The resulted Space Complexity:  $O(n)$

## Comparison

The solution using decrease and conquer:

```
class TowerOfHanoi {
    static int moves = 0;

    static void towerOfHanoi_FourRods(int n, char from_rod, char to_rod, char aux_rod1, char aux_rod2)
    {
        if (n == 0)
            return;
        if (n == 1)
        {
            System.out.println("\n Move disk" + n + " from rod " + from_rod + " to rod " + to_rod);
            moves++;
            return;
        }

        towerOfHanoi_FourRods(n - 2, from_rod, to_rod, aux_rod1, aux_rod1, aux_rod2, to_rod);
        int m = n-1;

        System.out.println("\n Move disk" + m + " from rod " + from_rod + " to rod " + aux_rod2);
        moves++;

        System.out.println("\n Move disk" + n + " from rod " + from_rod + " to rod " + to_rod);
        moves++;

        System.out.println("\n Move disk" + m + " from rod " + aux_rod2 + " to rod " + to_rod);
        moves++;
        towerOfHanoi_FourRods(n - 2, from_rod, aux_rod1, to_rod, aux_rod1, from_rod, aux_rod2);
    }
}
```

	Divide and Conquer	Decrease and Conquer
Advantages	Efficient for large problem size in case the number of disks increase	Simplicity as it reduces the problem by one until reaches the base case Lower Memory requirements as it often operate on the problem in place
Disadvantages	Overhead of combining sub-problems Increase memory usage	May not be efficient for large problems
Number of moves	Transfers 8 disks in 33 moves	Transfers 8 disks in 41 moves

## Sample of Output

### - Case of five disks

The five disks are divided into two sub-problems of size three and two disks based on the value of optimal k

The upper three disks are ordered using the four rods in the first call of function hanoi4

The lower two disks are ordered using the remaining three rods in the call of function hanoi3

The upper three disks move back again to be ordered above the bottom two disks using the four rods in the second call of function hanoi4

The number of moves is calculated based on dynamic programming that stores the result of the base case which are 0, 1, and 2 disks in

```
int n = 5; // Number of disks

Output - TowerOfHanoiTest (run) X
run:
Move Disk 1 from peg A to peg B
Move Disk 2 from peg A to peg C
Move Disk 1 from peg B to peg C
Move Disk 3 from peg A to peg B
Move Disk 1 from peg C to peg D
Move Disk 2 from peg C to peg B
Move Disk 1 from peg D to peg B
Move Disk 4 from peg A to peg C
Move Disk 5 from peg A to peg D
Move Disk 4 from peg C to peg D
Move Disk 1 from peg B to peg D
Move Disk 2 from peg B to peg A
Move Disk 1 from peg D to peg A
Move Disk 3 from peg B to peg D
Move Disk 1 from peg A to peg C
Move Disk 2 from peg A to peg D
Move Disk 1 from peg C to peg D
no of moves = 17
```

Figure(): Output of 5 disks

0, 1, and 3 moves then storing the number of moves in the dp array to be used in the next number of disks

### -Case of six disks

The six disks are divided into two sub-problems of equal size of three disks each based on the value of optimal k

The upper three disks are ordered using the four rods in the first call of function hanoi4

The lower three disks are ordered using the remaining three rods in the call of function hanoi3

The upper three disks move back again to be ordered above the bottom three disks using the four rods in the second call of function hanoi4

The number of moves is calculated based on

dynamic programming that stores the result of the base cases which are 0,1, and 2 disks in

0, 1, and 3 moves then storing the number of

moves in the dp array to be used in the next number of disks

```
int n = 6; // Number of disks

Output - TowerOfHanoiTest (run) X

run:
Move Disk 1 from peg A to peg B
Move Disk 2 from peg A to peg C
Move Disk 1 from peg B to peg C
Move Disk 3 from peg A to peg B
Move Disk 1 from peg C to peg D
Move Disk 2 from peg C to peg B
Move Disk 1 from peg D to peg B
Move Disk 4 from peg A to peg D
Move Disk 5 from peg A to peg C
Move Disk 4 from peg D to peg C
Move Disk 6 from peg A to peg D
Move Disk 4 from peg C to peg A
Move Disk 5 from peg C to peg D
Move Disk 4 from peg A to peg D
Move Disk 1 from peg B to peg D
Move Disk 2 from peg B to peg A
Move Disk 1 from peg D to peg A
Move Disk 3 from peg B to peg D
Move Disk 1 from peg A to peg C
Move Disk 2 from peg A to peg D
Move Disk 1 from peg C to peg D
no of moves = 21
```

Figure(): Output of six disks

### -Case of eight disks

The six disks are divided into two sub-problems of equal size of four disks each based on the value of optimal k

In the first call of hanoi4

The upper four disks are ordered by recursively calling hanoi4 and divide the four disks into two sub-problems of the same size of two disks each based on the value of optimal k

The upper two disks (disk 1 and 2) ordered on peg C using the four rods in the call of function hanoi4 while the lower two disks (disk 3 and 4) are ordered on peg B using the three rods in the call of function hanoi3 then the upper two disks are ordered above the lower two disks on peg B using the four rods in the call of function hanoi4

The lower four disks are ordered by calling hanoi3 and recursively reduce the problem by one into sub-problem until it reaches the base case which is two disks (disk 5 and disk 6) then order disk 7 then order disk 8 on peg D

In the second call of hanoi4

```
Move Disk 1 from peg A to peg B
Move Disk 2 from peg A to peg C
Move Disk 1 from peg B to peg C
Move Disk 3 from peg A to peg D
Move Disk 4 from peg A to peg B
Move Disk 3 from peg D to peg B
Move Disk 1 from peg C to peg D
Move Disk 2 from peg C to peg B
Move Disk 1 from peg D to peg B
Move Disk 5 from peg A to peg C
Move Disk 6 from peg A to peg D
Move Disk 5 from peg C to peg D
Move Disk 7 from peg A to peg C
Move Disk 5 from peg D to peg A
Move Disk 6 from peg D to peg C
Move Disk 5 from peg A to peg C
Move Disk 8 from peg A to peg D
Move Disk 5 from peg C to peg D
Move Disk 6 from peg C to peg A
Move Disk 5 from peg D to peg A
Move Disk 7 from peg C to peg D
Move Disk 5 from peg A to peg C
Move Disk 6 from peg A to peg D
Move Disk 5 from peg C to peg D
Move Disk 1 from peg B to peg D
Move Disk 2 from peg B to peg A
Move Disk 1 from peg D to peg A
Move Disk 3 from peg B to peg C
Move Disk 4 from peg B to peg D
Move Disk 3 from peg C to peg D
Move Disk 1 from peg A to peg C
Move Disk 2 from peg A to peg D
Move Disk 1 from peg C to peg D
no of moves = 33
```

Figure(): Output of eight disks

The upper four disks are ordered by recursively calling hanoi4 and again divide the four disks into two sub-problems of the same size of two disks each based on the value of optimal k

The upper two disks (disk 1 and 2) ordered on peg A using the four rods in the call of function hanoi4 while the lower two disks (disk 3 and 4) are ordered on peg D using the three rods in the call of function hanoi3 then the upper two disks are ordered above the lower two disks on peg D using the four rods in the call of function hanoi4.

### Conclusion

In Conclusion, using Divide and Conquer with the Dynamic Programming approach to divide the problem to two sub-problems and calculate the minimum number of moves required for different number of disks and storing them in the memory leads to optimizing the solution that offers a modified approach to solving the Tower of Hanoi Problem using four pegs, allowing more efficient solution compared to the traditional approach.

### References

The Four- Peg Tower of Hanoi Puzzle by Richard Johnsonbaugh.