# **Cloud Computing** (cloud)

**Module 2: Virtualization** 

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### Input of this Module

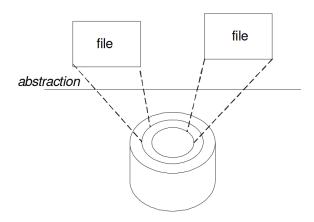
- <a href="https://learning.oreilly.com/playlists/43827098-7a87-435e-823e-736586b5694c">https://learning.oreilly.com/playlists/43827098-7a87-435e-823e-736586b5694c</a> (02-virtualization):
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  - Chisnall, The Definitive Guide to the Xen Hypervisor (2007)
  - Portnoy, Virtualization Essentials, 3rd Edition (2023)
- Research papers
  - Popek and Goldberg, Formal requirements for virtualizable third generation architectures (1974), Communications of the ACM, (Class Materials)
  - Randal, The Ideal Versus the Real: Revisiting the History of Virtual Machines and Containers (2020), ACM Computing Surveys, (Class Materials)
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  - VMware, Understanding Full Virtualization, Paravirtualization, and Hardware Assist (2019), (Class Materials)
  - De Gelas, Hardware Virtualization: the Nuts and Bolts (2008) https://www.anandtech.com/show/2480/9
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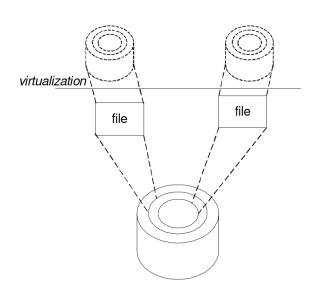
### Agenda

- 1. Introduction to Virtualization
- 2. Virtualization Approaches
  - Type I Hypervisors (native)
  - Type II Hypervisors (hosted)
- 3. Virtualization Technology
  - Intel VT-x
  - Examples (AWS EC2, XEN, KVM)
- 4. Virtual Resources
  - CPU / Memory
  - Storage
  - Networking

### Abstraction vs Virtualization

- Abstraction: systems are built on multiple abstraction levels
  - Higher levels hide details of lower levels
  - Example: files are an abstraction of the disk
- Virtualization: creating a virtual representation of a resource (on a physical machine)
  - Does not necessarily hide any low-level details
  - Example: virtual memory
  - >> A cloud (service) is typically created by adding management and automation layers (abstraction) to multiple layers of virtualization.

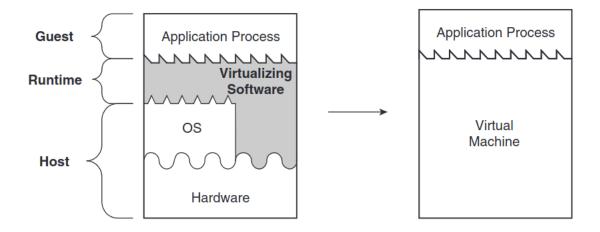




### **Virtual Machines (VM)**

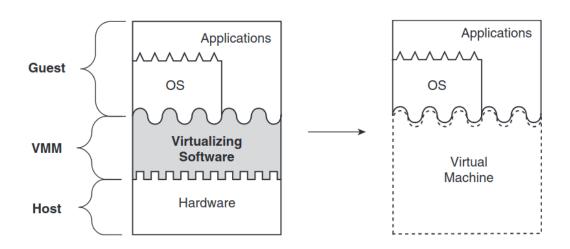
#### **Process** Virtual Machines

- Typically using intermediate language (Java bytecode) and make their execution platform agnostic
- Example: Java Virtual Machine (JVM)



#### **System** Virtual Machines

- Try to mimic the complete hardware system (CPU, memory, ...)
- Example: Xen, KVM, VMware

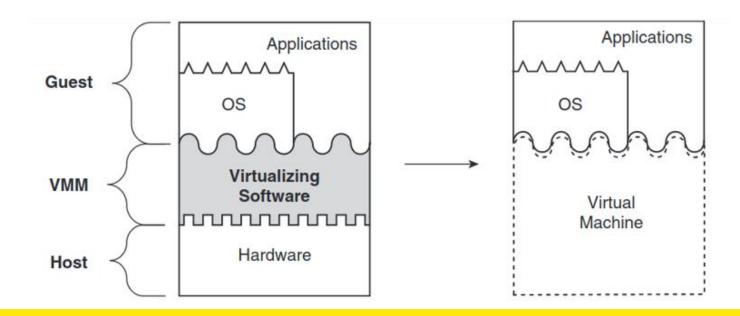


**VMM** = Virtual Machine Monitor or Hypervisor (thin layer of software that supports virtualization)

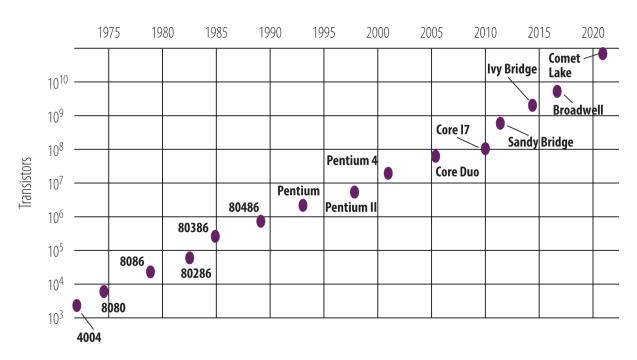
### Requirements of Virtualization

According to Popek and Goldberg (1974) there are three properties of interest when any arbitrary program is run on a virtual machine:

- *Efficiency* minimal overhead, good performance
- Resource control (safety) isolation between guest and VMM / host, isolation between guests
- Equivalence same results with and without VMM

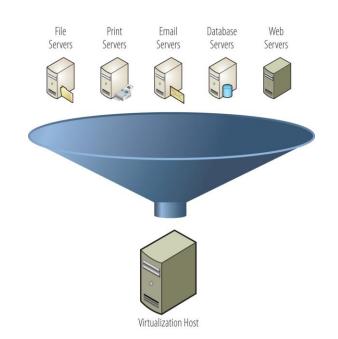


### Understanding the Importance of Virtualization



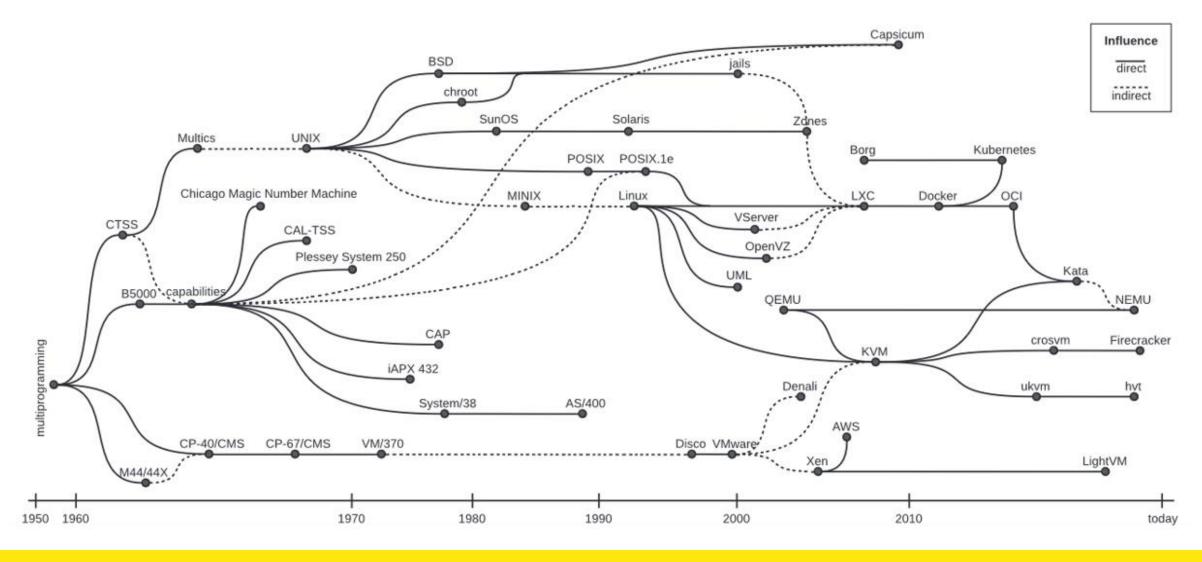
**Moore's law** says that processing power (the speed and capacity of computer chips) roughly doubles every 18 months.

Combination of the effect of Moore's law and the "one server, one application" model → servers do less and less work

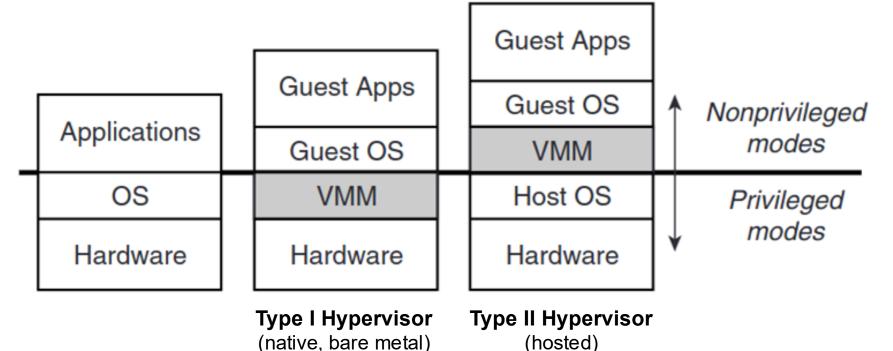


>> Virtualization helped to overcome this by reducing footprint and total cost of ownership per server: TCO is 3-10 times the costs of the server itself in 3 year

### **Evolution of Compute Virtualization**



### Different Types of Virtualization











#### Examples:

- VMWare ESX
- Citrix Xen Server
- Microsoft Hyper-V

#### Examples:

- Oracle VirtualBox
- VMWare Workstation



### Type II Hypervisor (Hosted Virtualization)

- Hypervisor runs a regular application on top of the host operating system.
- Software-level representation (emulation) of physical resources.

#### **Advantages**

- Guest applications run within process boundaries (instructions aren't executed directly on host's hardware) → complete isolation
- Easy to handle privileged instructions.

#### **Disadvantages**

- Emulating hardware resources, e.g., modern processors, is difficult.
- Execution overhead (~ 100x slower than direct execution)



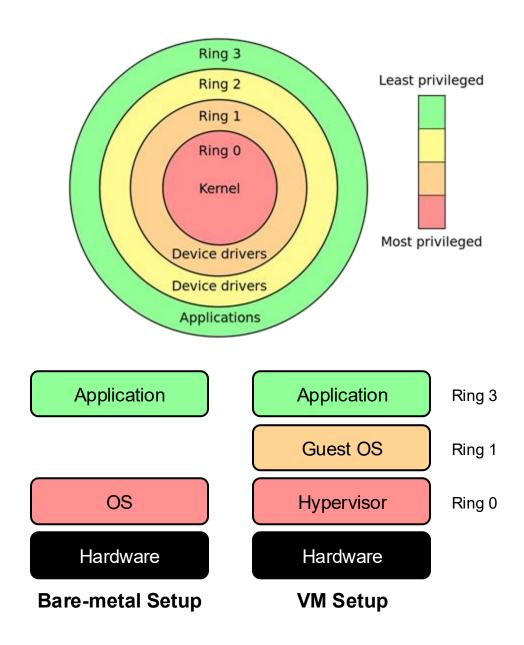
This virtualization approach is not suitable for offering compute (VM) at scale in the cloud because it results in less efficient resource utilization and higher latency.

### Type I Hypervisor

- Full virtualization using binary translation
- Paravirtualization (OS assisted virtualization)
- Hardware assisted virtualization

### x86 Privilege rings

- x86 CPU hardware provides four protection rings.
- Only code in more privileged rings can read and write memory of lower rings.
- Ring 0 is the level with the most privileges.
  - Interaction with physical hardware (e.g., reset processor, perform I/O operations)



### Full Virtualization using Binary Translation (Type I)

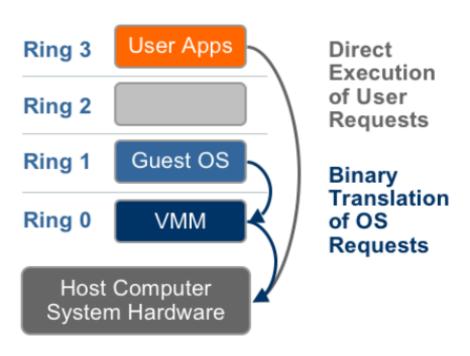
Translates kernel code to replace non-virtualizable instructions with new sequences of instructions that have the intended effect on the virtual hardware. User level code is directly executed on the processor.

#### **Advantages**

- Allowing the guest OS and apps to run without modification or awareness of the virtualization.
- Direct execution for user-level code
  - → high performance

#### **Disadvantages**

- Difficult to implement VMM
- Translation impacts performance
  (e.g. virtual memory → physical memory → machine memory)



### Paravirtualization (Type I)

Modifying the guest OS kernel to replace non-virtualizable instructions with hypercalls that communicate directly with the virtualization layer hypervisor. It is therefore also referred to as OS Assisted Virtualization.

#### **Advantages**

 Reduces the virtualization overhead due to fewer context switches and less bookkeeping

#### **Disadvantages**

- Poor compatibility and portability
  - → cannot support unmodified OS
- Requires substantial modification of the OS

User Apps Ring 3 Direct Ring 2 Execution of User Requests Ring 1 'Hypercalls' to the Paravirtualized Ring 0 Virtualization **Guest OS** Layer replace Non-virtualizable Virtualization Layer OS Instructions Host Computer System Hardware

Typically used in virtual device drivers (e.g., Vmxnet) because CPU and memory paravirtualization is much more difficult.

### Hardware-assisted Virtualization (Type I)

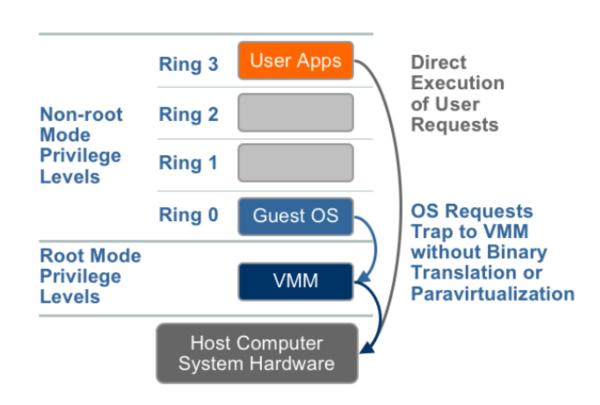
- Hardware extension for virtualization added to the CPU (e.g., AMD-V, Intel VT-x)
- CPU execution mode feature that allows the VMM to run in a new root mode (e.g., VT root mode) below ring 0.
- Allows direct execution of VM on processor until a privileged or sensitive instruction is executed, which are set to automatically trap to the hypervisor.

#### **Advantages**

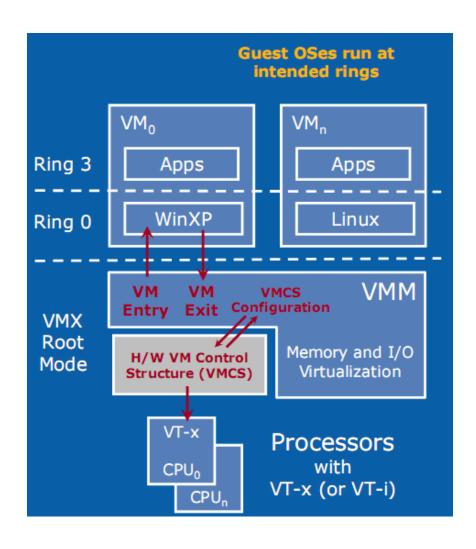
Fast and supported on most CPUs today.

#### **Disadvantages**

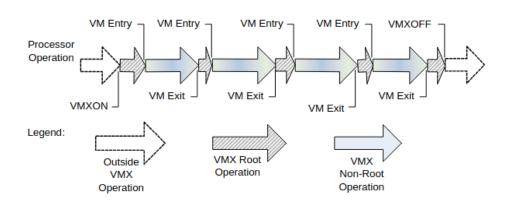
Focus is mainly on the CPU and memory, but IO virtualization still requires translation.



### Intel Virtualization Technology (VT-x)



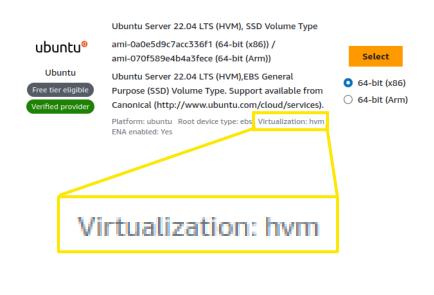
- Virtual Machine Monitor (VMM) configures on the Virtual Machine
  Control Structure (VMCS) which privileged instructions (e.g., exceptions, interrupt, IO) trigger a vmexit.
- VMCS stores context info which contains e.g., CPU register values to save and restore during transitions between root and non-root.
- Using vmentry the VMM can transfer control to VM (one at a time) until vmexit.



### Example: Amazon EC2

Amazon Machine Images (AMI) use one of two types of virtualization: **paravirtual** (PV) or **hardware virtual machine** (HVM).

Characteristic	HVM	PV
Description	HVM AMIs are presented with a fully virtualized set of hardware and boot by executing the master boot record of the root block device of your image. This virtualization type provides the ability to run an operating system directly on top of a virtual machine without any modification, as if it were run on the bare-metal hardware. The Amazon EC2 host system emulates some or all of the underlying hardware that is presented to the guest.	PV AMIs boot with a special boot loader called PV-GRUB, which starts the boot cycle and ther chain loads the kernel specified in the menu.lst file on your image. Paravirtual guests can run on host hardware that does not have explicit support for virtualization. For more information about PV-GRUB and its use in Amazon EC2, see User provided kernels.
Supported instance types	All current generation instance types support HVM AMIs.	The following previous generation instance types support PV AMIs: C1, C3, M1, M3, M2, and T1. Current generation instance types do not support PV AMIs.

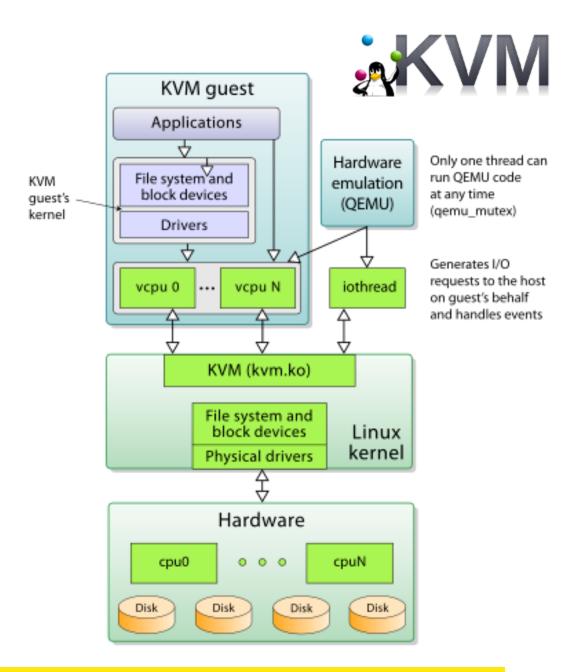


#### PV on HVM

PV drivers are available for HVM guests (performance advantages in storage and network I/O) >> HVM guests can get the same, or better, performance than PV guests.

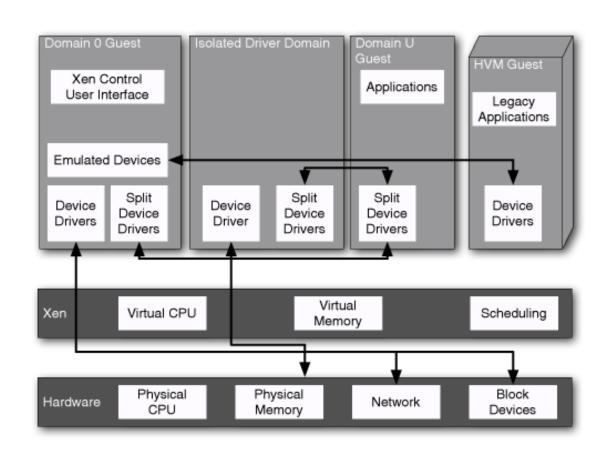
## Example: KVM (Kernel Virtual Machine)

- KVM is a virtualization technology which uses Linux as a hypervisor (kernel module).
- Part of the kernel since version 2.6.20 (2007) → Open Source
- Supports hardware-assisted (Intel VT and AMD-V) and in addition paravirtualization for IO (VirtIO API).
- QEMU (Quick EMUlator) is a generic machine emulator and virtualizer.
- Emulates machine's processor including the devices using dynamic binary translation.
- Often used in combination with KVM.
  - Which hypervisor type is KVM / QEMU?



### Example: Xen

- DomU (unprivileged domain): Guest systems running on Xen hypervisor
- Dom0 (privileged guest): Management System governing access to hardware for DomU
  - User Interface of the XEN system
  - Handles all I/O, driver support, etc. for individual guests
  - If Dom0 is unavailable, all DomU are unavailable
- Different virtualization techniques:
  - PV: Paravirtualization, specialized Kernel necessary, very lightweight
  - HVM: Hardware Virtual Machine, full virtualization



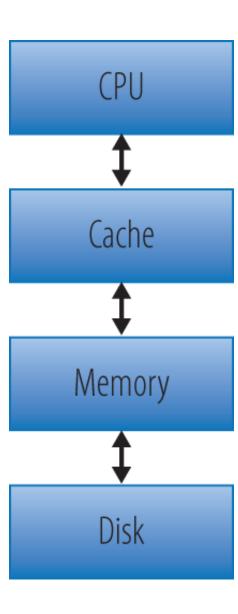
### Virtual Resources: CPU

- Host CPUs are virtualized by granted slices of CPU-time to guests.
- There is no fixed mapping between virtual CPU of the guest (vCPU) and CPU on the host
  - Work might be schedule on any core of the host
  - Gang Scheduling (all-or-nothing) might apply
- Hyperthreading, multiple Cores per CPU must be taken into account
  - Not doubling resources but might add 30% more performance (remember Conway's Law)
  - Makes work not faster but easier to schedule
- How many vCPUs can be provision?
- (Threads x Cores) x Physical CPU = Number vCPU
- Ratios vCPU:CPU between 1:1-3:1 (computing intensive) to 6:1 (non computing)

### Virtual Resources: Memory

### Recap sysad, bsys:

- Primary Storage / Secondary Storage
- If too much RAM is consumed, memory is paged to disk
  - Paging can harm the entire host since paging is expensive (swap-partition anyone?)
- If too less RAM is consumed, memory is "blocked" normally by the process
- Things get more complicated as host sees the guest as blackbox.
  - → **sophisticated methods** (ballooning, overcommitment, memory sharing) need to be applied to get most out of the physical memory.

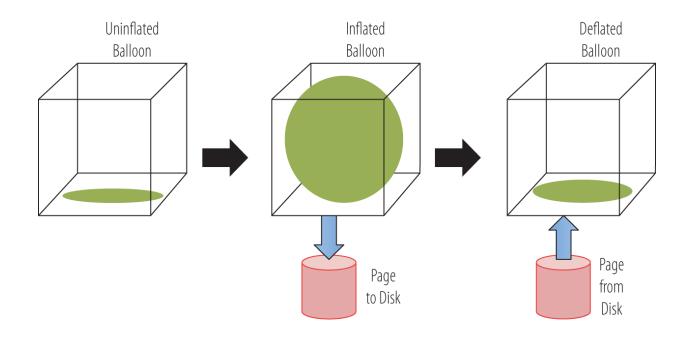


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### Virtual Memory: Ballooning

Pattern to be applied when there is discussion about memory.

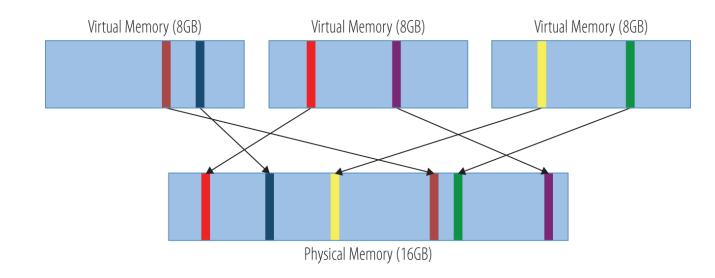
- 1. Memory is needed → inflating the balloon
- OS of Host analyses slices of memory to be paged and pages them to disk
- Memory gets free and can be provided to the demanding guest
- If memory need decreases → deflate the balloon.
  The stale slices are loaded from disk again.



### Virtual Memory: Overcommitment

Pattern to map only used / active memory to host memory

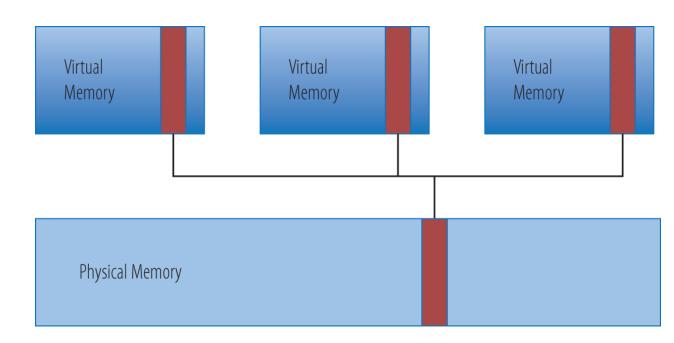
- Analyzing the memory consumption of the guests, the "inactive" parts are not reserved in the host
- Only actively used blocks are mapped to the host
- Memory patterns of guest must be known
- Ratios up to 1.5:1 or 2:1 are possible



### Virtual Memory: Sharing

Often, machines there are many similar guest on the same host

- If the system and the application is the same, the memory consumption might be (partially) the same
- In such a scenario, the memory of multiple guests is mapped once on the slice on the host
- 10%-40% less consumption on practice
- Works best with horizontal scaled infrastructure
- If memory is altered, copy-on-write is applied



### Virtual Memory: Features used by virtualization technologies

Techniques	VMware vSphere (vSphere 8.0)	Microsoft Hyper-V (Server 2022 R2)	KVM
Overcommit	Yes	Yes	Yes
Ballooning	Yes	Yes	Yes
Page sharing	Yes	No	Yes
Compression	Yes	No	No

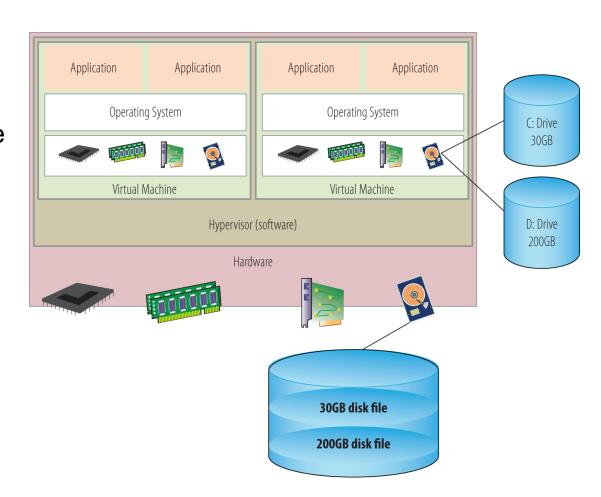
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### Virtual Resources: Storage

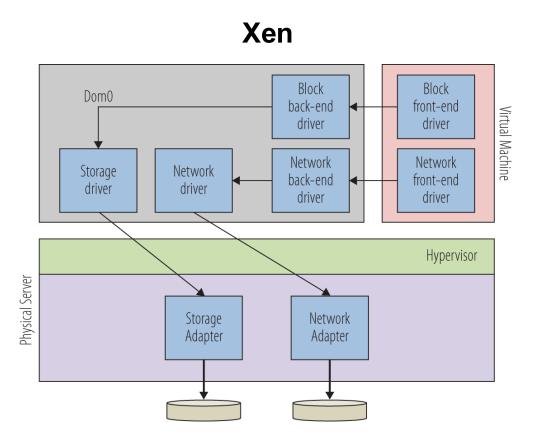
Disks are mapped as **Blockdevices** into the machine:

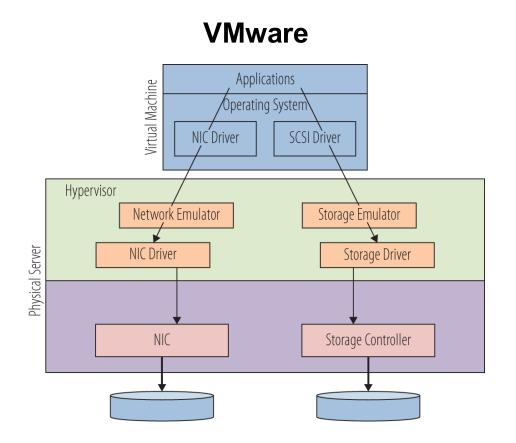
- Large Disk of any kind is bound to Hypervisor
- In the Machine, one or more disks are associated to a single guest
- Since devices are bound as blockdevices, filesystem is generated by the guest
- the Hypervisor / Host is unware about the filesystem



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### Virtual Resources: Storage



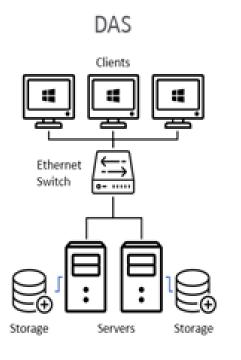


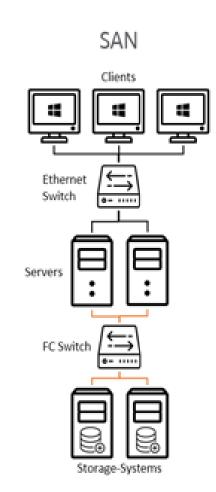
**Storage Controller** → Direct attached Storage (DAS)

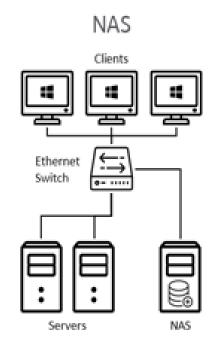
NIC (Network Interface Controller)→ Storage Area Network (SAN) or Networked Attached Storage (NAS)

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### Storage Types: DAS, SAN, NAS







#### **DAS (Direct attached Storage)**

- Disk directly attached to server via SCSI / IDE/ SATA
- · Cheap, easy, fast setup
- If server goes down, disk go down

#### **SAN (Storage Area Network)**

- Disk attached via FibreChannel to one/multiple servers
- Complex, second stack, more expensive
- Failover / Scaling possible

#### **NAS (Networked Attached Storage)**

- Disk attached directly as device to network
- Easier technical setup, cheaper, redudant
- Blockaccess needs workaround e.g. iSCSI

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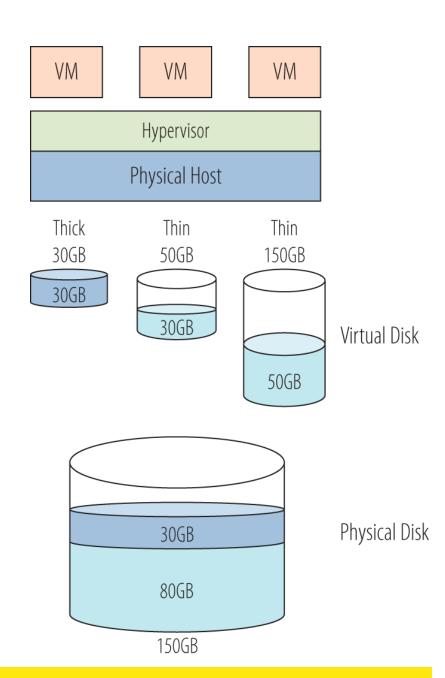
### Thin vs Thick Provisioning

### **Thick Provisioning**

- Reserve the space for the guest, the guest allocate with its disk
- If the guest uses less, space is wasted

#### **Thin Provisioning**

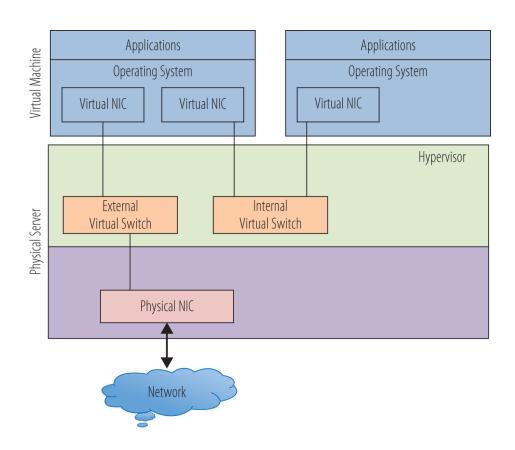
- Reserve only the space, the guest really uses and allocate the size the guest think it has dynamically
- If all guests consume all disks, problems occur.



### Virtual Resources: Network

Network connection takes place via virtual network interface card (NIC):

- No direct connect to physical NIC of host (per default)
- Virtual Switches are applied to route traffic from VMs to NIC of host
- VM-to-VM traffic stays on same host





WHY AREN'T MY

ARMS GROWING

CHY DO GOOD PEOPLE DIE

WHY AREN'T THERE ANY FOREIGN MILITARY BASES IN AMERICA

WHYARENT

THERE GUNS IN

HARRY POTTER

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### **Summarizing Questions**

- Explain why Moore's Law is important regarding virtualization
- Describe the difference between
  - Type1 and Type2 virtualization
  - Hardware-assisted, Full, and Paravirtualization
- Why are the properties and requirments for virtualization also important for cloud computing?
- Which virtualization approach provides the best performance?
- Compute the number of vCPUs for a machine with 4 Quadcores and hyperthreading enabled
- What can be a problem with Balooning?
- What is thick provision? When should you use it?
- Why is generating backups with Snapshots only not a good idea?
- What is the primary advantage of using Open vSwitch (OVS) with VxLAN in a cloud environment?