Problem 1

(a)

3.1-2:

Let $f(n) = (n+a)^b$, and we define $f(n) = \Theta(g(n))$. From the definition of Θ . We know that there exists c_1 , c_2 and n_0 , which let any $n \ge n_0$, $0 \le c_1 g(n) \le f(n) \le c_2 g(n)$.

By Binomial theorem, we can know that,

$$(n+a)^b = C_b^b n^b + C_b^{b-1} n^{b-1} a + \dots + C_b^0 a^b$$

So we can easily find that for $c_1 = 1$, $c_2 = b$, $g(n) = n^b$, the definition is true. So $(n + a)^b = \Theta(n^b)$.

3-2:

| A | В | О | О | Ω | ω | Θ |
|------------|----------------|-----|-----|-----|----------|-----|
| $lg^k n$ | n^{ϵ} | yes | yes | no | no | no |
| n^k | c^n | yes | yes | no | no | no |
| \sqrt{n} | $n^{sin(n)}$ | no | no | no | no | no |
| 2^n | $2^{n/2}$ | no | no | yes | yes | no |
| n^{lgc} | c^{lgn} | yes | no | yes | no | yes |
| lg(n!) | $lg(n^n)$ | yes | no | yes | no | yes |

3.2-7:

While i = 0, we can easily know that,

$$F_0 = 0$$

While i = 1, we can find that,

$$F_1 = \frac{\frac{1+\sqrt{5}}{2} - \frac{1-\sqrt{5}}{2}}{\sqrt{5}}$$
$$= \frac{\sqrt{5}}{\sqrt{5}} = 1$$

While i = k, we can find that,

$$F_k = \frac{\phi^k - (\overline{\phi})^k}{\sqrt{5}}$$

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While i = k + 1, we can find that,

$$\begin{split} F_{k+1} &= F_k + F_{k-1} \\ &= \frac{\phi^k - (\overline{\phi})^k}{\sqrt{5}} + \frac{\phi^{k-1} - (\overline{\phi})^{k-1}}{\sqrt{5}} \\ &= \frac{(1 + \frac{\sqrt{5} - 1}{2})\phi^k - (1 - \frac{\sqrt{5} + 1}{2})(\overline{\phi})^k}{\sqrt{5}} \\ &= \frac{\frac{1 + \sqrt{5}}{2}\phi^k - \frac{1 - \sqrt{5}}{2}(\overline{\phi})^k}{\sqrt{5}} \\ &= \frac{\phi^{k+1} - (\overline{\phi})^{k+1}}{\sqrt{5}} \end{split}$$

So we have prove that

$$F_i = \frac{\phi^i - (\overline{\phi})^i}{\sqrt{5}}$$

(b)

$$O(\prod_{k=1}^{n}(1-\frac{1}{k^2})) = O((1-\frac{1}{n})^n) = O(n^{\frac{1}{n}}) \in O(\ln(\ln n)) \in O(\ln n) = O(\sum_{k=1}^{n}\frac{1}{k}) \in O(n^{1+\cos n}) \in O(\log(n!)) \in O(3^{\ln n}) \in O(n^2+3n\ln n+5) = O(n^2+3n+5) = O(n^2+n^{-2}) = O(\sum_{k=1}^{\log n}\frac{n^2}{2^k}) \in O(2^n) \in O(n!) \in O((1+n)^n) \in O(n^{n^2-1}) \in O(n^{n^2}+n!)$$

(c)

In this question, we can just take $T(n) = c_1 n + c_2 n \log_2(n)$ into $T(n) = 2T(\frac{n}{2}) + n$, then we can get,

$$c_1 n + c_2 n \log_2(n) = 2 * (c_1 \frac{n}{2} + c_2 \frac{n}{2} \log_2(\frac{n}{2})) + n$$

$$c_1 n + c_2 n \log_2(n) = c_1 n + c_2 n \log_2(\frac{n}{2}) + n$$

$$c_1 n + c_2 n \log_2(n) = c_1 n + c_2 n (\log_2 n - \log_2 2) + n$$

$$c_1 n + c_2 n \log_2(n) = (c_1 - c_2 + 1)n + c_2 n \log_2(n)$$

From this, we can know that $c_2 = 1$ and c_1 can be any real numbers.

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(d)

For the new case, we just do it again.

$$c_1 n^{\gamma} + c_2 n^k = b^{\gamma} * (c_1 (\frac{n}{b})^{\gamma} + c_2 (\frac{n}{b})^k) + n^k$$
$$= c_1 n^{\gamma} + c_2 b^{\gamma - k} n^k + n^k$$
$$= c_1 n^{\gamma} + (c_2 b^{\gamma - k} + 1) n^k$$

So, c_1 can be any real numbers and $c_2 = \frac{1}{1-b^{\gamma-k}}$. If we set $\gamma = k$, it will be wrong. That is because the denominator of c_2 will be 0.

Then we start over with $T(n) = c_1 n^{\gamma} + c_2 n^{\gamma} log_2 n$.

$$c_{1}n^{\gamma} + c_{2}n^{\gamma}log_{2}n = b^{\gamma} * (c_{1}(\frac{n}{b})^{\gamma} + c_{2}(\frac{n}{b})^{\gamma}log_{2}(\frac{n}{b})) + n^{\gamma}$$

$$= c_{1}n^{\gamma} + c_{2}n^{\gamma}(log_{2}n - log_{2}b) + n^{\gamma}$$

$$= (c_{1} + 1)n^{\gamma} + c_{2}n^{\gamma}(log_{2}n - log_{2}b)$$

$$= (c_{1} + 1 - c_{2}log_{2}b)n^{\gamma} + c_{2}n^{\gamma}log_{2}n$$

So, c_1 can be any real numbers and $c_2 = \frac{1}{\log_2 b}$.

Problem 2

```
Algorithm 1 Nuts and Bolts
Input: Arraynuts, Arraybolts, begin, end
Output: Arraynuts, Arraybolts
0: function SORTNUTANDBOLT(Arraynuts, Arraybolts, begin, end)
  if begin > end then
    return
  end if
  tempb = Arraybolts[end]
  ini = begin
  for i in begin to end - 1 do
    if Test(Arraynuts[i], tempb) < 0 then
      temp = Arraynuts[i]
      Arraynuts[i] = Arraynuts[ini]
      Arraynuts[ini] = temp
      ini + +
    else if Test(Arraynuts[i], tempb) == 0 then
      temp = Arraynuts[i]
      Arraynuts[i] = Arraynuts[end]
      Arraynuts[end] = temp
    end if
  end for
  temp = Arraynuts[ini]
  Arraynuts[ini] = Arraynuts[end]
  Arraynuts[end] = temp
  tempn = Arraynuts[ini]
  ini = begin
  for j in begin to end - 1 do
    if Test(tempn, Arraybolts[j]) < 0 then
      temp = Arraybolts[j]
      Arraybolts[j] = Arraybolts[ini]
      Arraybolts[ini] = temp
      ini + +
    else if Test(tempn, Arraybolts[j]) == 0 then
      temp = Arraybolts[j]
      Arraybolts[j] = Arraybolts[end]
      Arraybolts[end] = temp
      i - -
    end if
  end for
  temp = Arraybolts[ini]
  Arraybolts[ini] = Arraybolts[end]
  Arraybolts[end] = temp
  SortNutAndBolt(Arraynuts, Arraybolts, begin, ini-1)
  SortNutAndBolt(Arraynuts, Arraybolts, ini + 1, end)
```