



**National Textile University**  
**Department of Computer Science**

Subject:

Operating System

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Submitted to:

Sir Nasir

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Submitted by:

Ahmad Fawad

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Reg number:

1129

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Assignment No:

01

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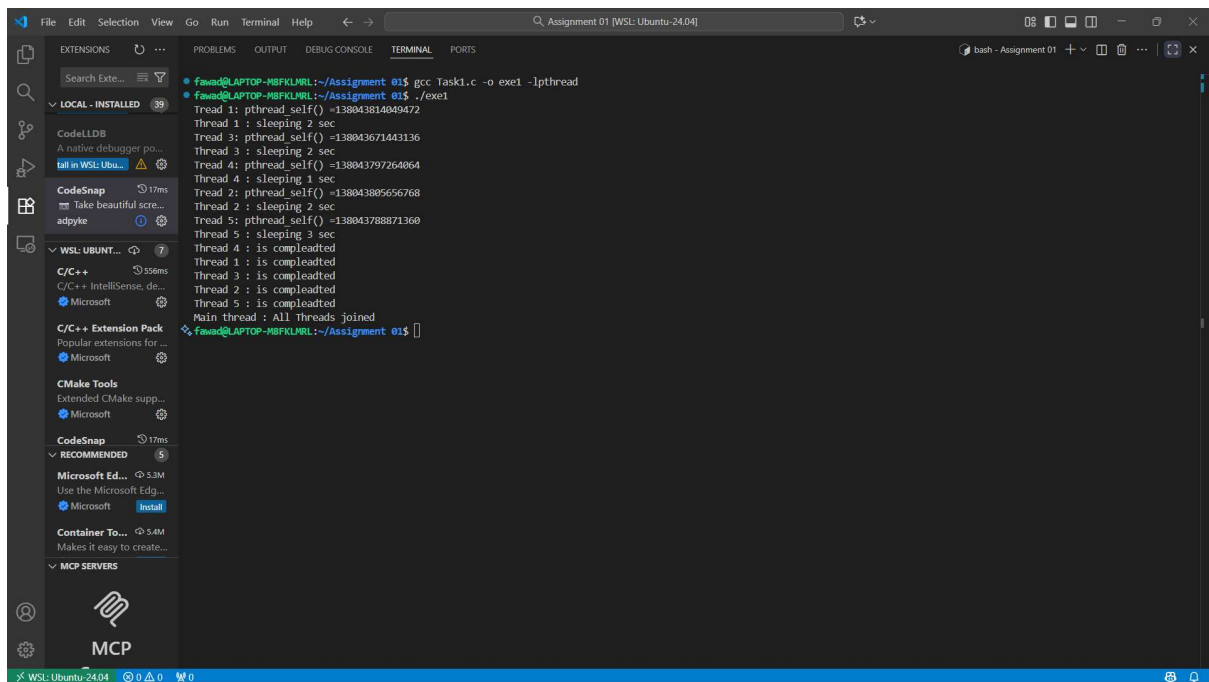
Semester:



## Task 1:

```
1 //Thread Information Display
2 //Name: Ahmad Fawad
3 // RegNo: 1129
4
5 #include <stdio.h>
6 #include <stdlib.h>
7 #include <pthread.h>
8 #include <unistd.h>
9
10 void *thread_function(void *arg)
11 {
12     int *num = (int*)arg;
13     int id = *num;
14     pthread_t tid = pthread_self();
15
16     printf("Tread %d: pthread_self() =%lu\n",id ,(unsigned long)tid);
17
18     int sleep_time = (rand() %3)+1;
19     printf("Thread %d : sleeping %d sec\n",id ,sleep_time);
20     sleep(sleep_time);
21     printf("Thread %d : is compleadted\n",id);
22     free(arg);
23     return NULL;
24 }
25
26
27 int main()
28 {
29     pthread_t threads[5];
30     for (int i =0; i<5;i++)
31     {
32         int *arg = malloc(sizeof(int));
33         *arg = i +1 ;
34         pthread_create(&threads[i] , NULL,thread_function,arg);
35
36     }
37     for (int i = 0; i < 5; i++)
38     {
39         pthread_join(threads[i],NULL);
40     }
41     printf("Main thread : All Threads joined \n");
42     return 0;
43
44 }
```

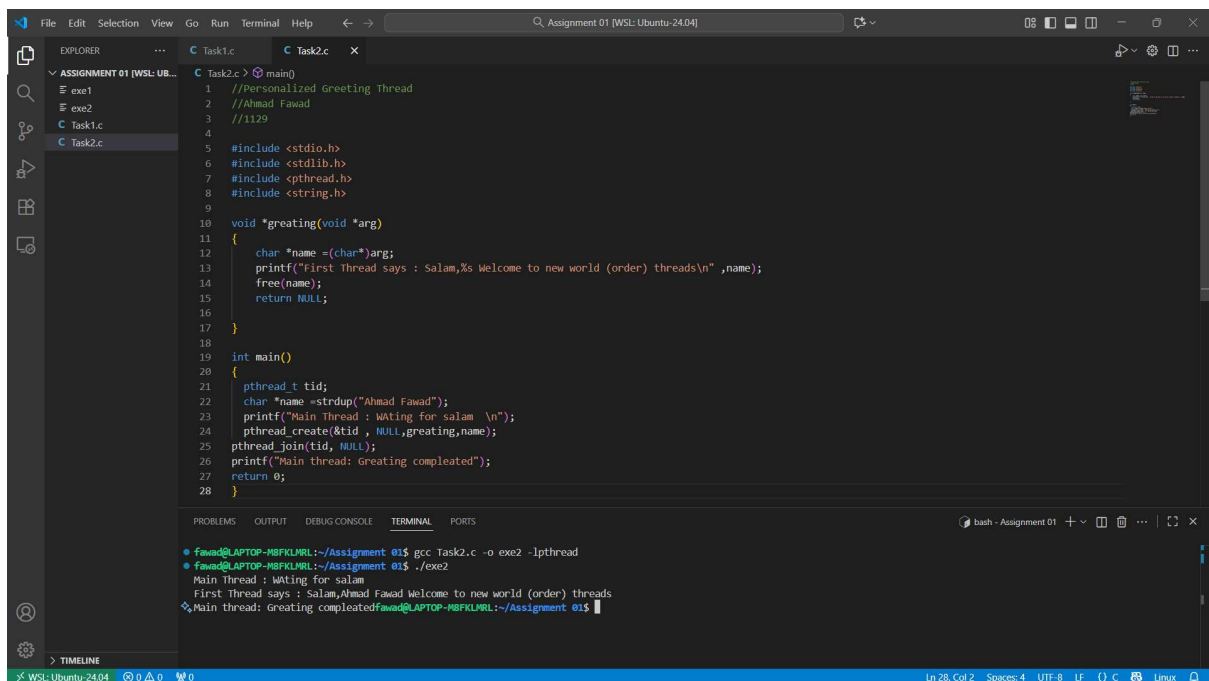
# Output:



The screenshot shows the Visual Studio Code interface with the terminal window open. The terminal displays the output of a C program that creates 5 threads. The output shows the threads sleeping for 2 seconds, then completing. The main thread joins all child threads and prints "Main thread : All Threads joined".

```
fawad@LAPTOP-MBFKLMRL:~/Assignment 01$ gcc task1.c -o exe1 -lpthread
fawad@LAPTOP-MBFKLMRL:~/Assignment 01$ ./exe1
Thread 1 : pthread_self() =138043814049472
Thread 1 : sleeping 2 sec
Thread 3 : pthread_self() =138043671443136
Thread 3 : sleeping 2 sec
Thread 4 : pthread_self() =138043797264064
Thread 4 : sleeping 1 sec
Thread 2 : pthread_self() =138043805656768
Thread 2 : sleeping 2 sec
Thread 5 : pthread_self() =138043788871360
Thread 5 : sleeping 3 sec
Thread 4 : is compleated
Thread 1 : is compleated
Thread 3 : is compleated
Thread 2 : is compleated
Thread 5 : is compleated
Main thread : All Threads joined
fawad@LAPTOP-MBFKLMRL:~/Assignment 01$
```

# Task2:

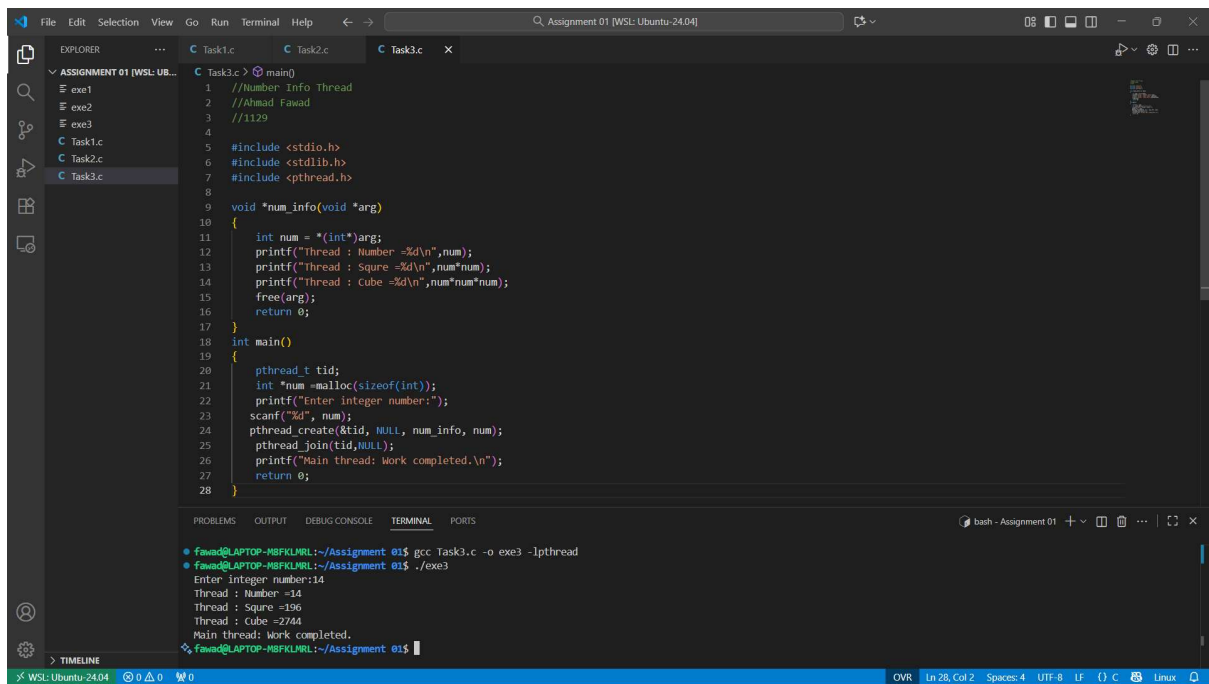


The screenshot shows the Visual Studio Code interface with the Explorer and Terminal windows open. The Explorer window shows the file structure of the project, including Task1.c and Task2.c. The Terminal window shows the output of the program, which creates a main thread and a child thread. The main thread waits for the child thread to complete before printing "Main thread: Greeting completed".

```
1 //Personalized Greeting Thread
2 //Ahmad Fawad
3 //1129
4
5 #include <stdio.h>
6 #include <stdlib.h>
7 #include <pthread.h>
8 #include <string.h>
9
10 void *greeting(void *arg)
11 {
12     char *name =(char*)arg;
13     printf("First Thread says : Salam,%s Welcome to new world (order) threads\n",name);
14     free(name);
15     return NULL;
16 }
17
18
19 int main()
20 {
21     pthread_t tid;
22     char *name =strdup("Ahmad Fawad");
23     printf("Main Thread : WAting for salam \n");
24     pthread_create(&tid , NULL,greeting,name);
25     pthread_join(tid, NULL);
26     printf("Main thread: Greeting compleated");
27     return 0;
28 }
```

```
fawad@LAPTOP-MBFKLMRL:~/Assignment 01$ gcc Task2.c -o exe2 -lpthread
fawad@LAPTOP-MBFKLMRL:~/Assignment 01$ ./exe2
Main Thread : WAting for salam
First Thread says : Salam,Ahmad Fawad Welcome to new world (order) threads
Main thread: Greeting compleatedfawad@LAPTOP-MBFKLMRL:~/Assignment 01$
```

## Task3:

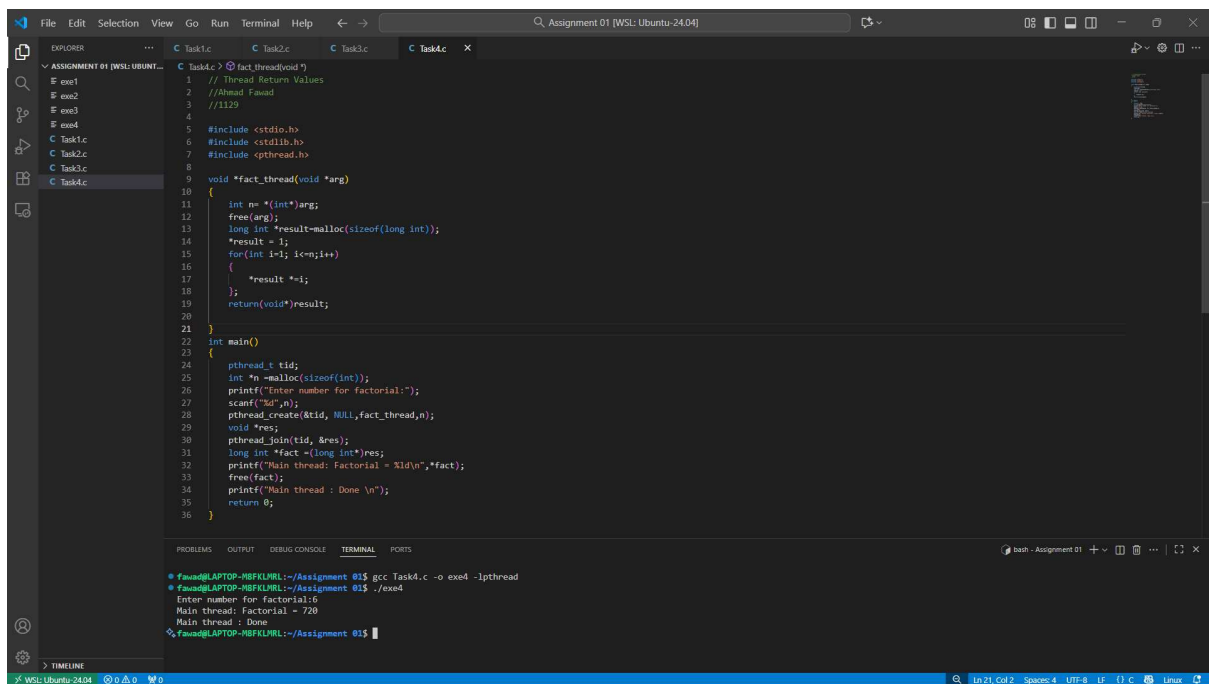


The screenshot shows the Visual Studio Code editor with the file `Task3.c` open. The code defines a `num_info` function that prints the number, its square, and its cube, then returns the number. The `main` function allocates memory for a number, prompts the user to enter an integer, and then creates a thread to call `num_info`. The terminal output shows the program being compiled and executed, with the user entering the number 14. The output displays the thread's calculations and the main thread's completion message.

```
1 //Number Info Thread
2 //Ahmad Fawad
3 //1129
4
5 #include <stdio.h>
6 #include <stdlib.h>
7 #include <pthread.h>
8
9 void *num_info(void *arg)
10 {
11     int num = *(int*)arg;
12     printf("Thread : Number =%d\n",num);
13     printf("Thread : Square =%d\n",num*num);
14     printf("Thread : Cube =%d\n",num*num*num);
15     free(arg);
16     return 0;
17 }
18
19 int main()
20 {
21     pthread_t tid;
22     int *num = malloc(sizeof(int));
23     printf("Enter integer number:");
24     scanf("%d", num);
25     pthread_create(&tid, NULL, num_info, num);
26     pthread_join(tid, NULL);
27     printf("Main thread: Work completed.\n");
28     return 0;
29 }
```

```
bash - Assignment 01
fawad@LAPTOP-MBFKURL:~/Assignment 01$ gcc Task3.c -o exe3 -lpthread
fawad@LAPTOP-MBFKURL:~/Assignment 01$ ./exe3
Enter integer number:14
Thread : Number =14
Thread : Square =196
Thread : Cube =2744
Main thread: Work completed.
fawad@LAPTOP-MBFKURL:~/Assignment 01$
```

## Task4:



The screenshot shows the Visual Studio Code editor with the file `Task4.c` open. The code defines a `fact_thread` function that calculates the factorial of a number and returns the result. The `main` function prompts the user to enter a number for factorial, creates a thread to call `fact_thread`, and then prints the result. The terminal output shows the program being compiled and executed, with the user entering the number 6. The output displays the factorial calculation and the main thread's completion message.

```
1 // Thread Return Values
2 //Ahmad Fawad
3 //1129
4
5 #include <stdio.h>
6 #include <stdlib.h>
7 #include <pthread.h>
8
9 void *fact_thread(void *arg)
10 {
11     int n = *(int*)arg;
12     free(arg);
13     long int *result = malloc(sizeof(long int));
14     *result = 1;
15     for(int i=1; i<=n; i++)
16     {
17         *result *=i;
18     };
19     return(void*)result;
20 }
21
22 int main()
23 {
24     pthread_t tid;
25     int *n = malloc(sizeof(int));
26     printf("Enter number for Factorial:");
27     scanf("%d",n);
28     pthread_create(&tid, NULL, fact_thread, n);
29     void *res;
30     pthread_join(tid, &res);
31     long int *fact = (long int*)res;
32     printf("Main thread: Factorial = %ld\n",*fact);
33     free(fact);
34     printf("Main thread : Done \n");
35     return 0;
36 }
```

```
bash - Assignment 01
fawad@LAPTOP-MBFKURL:~/Assignment 01$ gcc Task4.c -o exe4 -lpthread
fawad@LAPTOP-MBFKURL:~/Assignment 01$ ./exe4
Enter number for factorial:6
Main thread: Factorial = 720
Main thread : Done
fawad@LAPTOP-MBFKURL:~/Assignment 01$
```

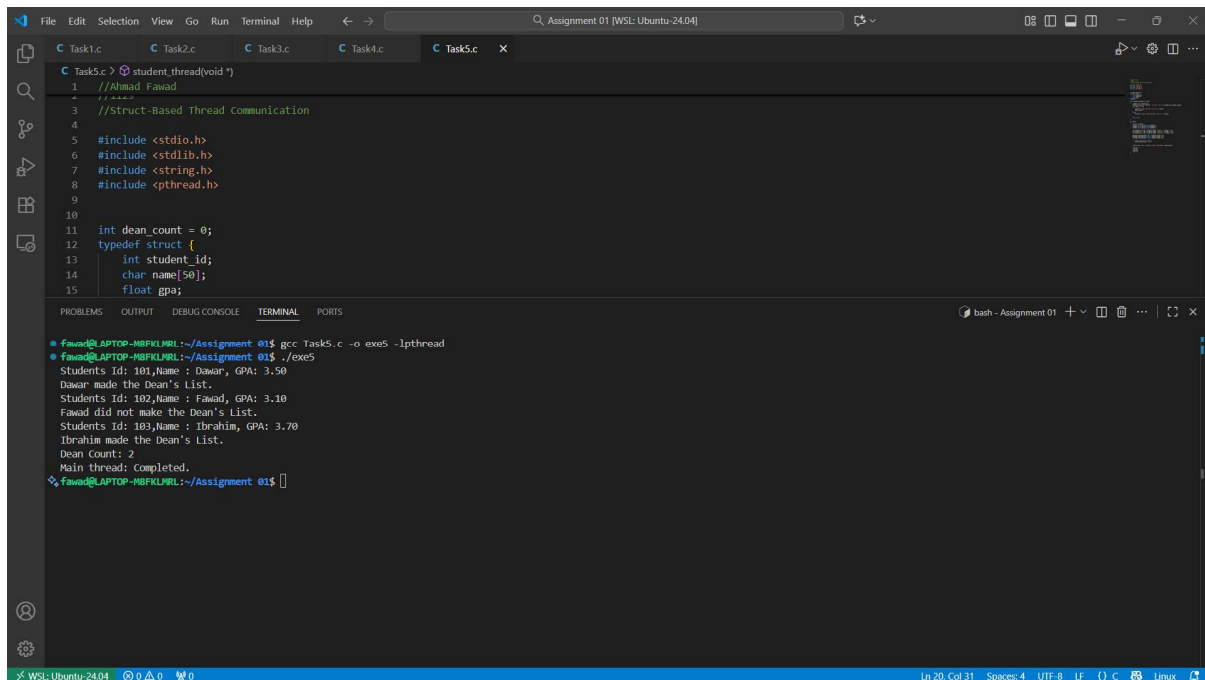
## Task5:

```

1  //Ahmad Fawad
2  //1129
3  //Struct-Based Thread Communication
4
5  #include <stdio.h>
6  #include <stdlib.h>
7  #include <string.h>
8  #include <pthread.h>
9
10
11  int dean_count = 0;
12  typedef struct {
13      int student_id;
14      char name[50];
15      float gpa;
16  }student;
17
18  void *student_thread(void *arg)
19  {
20      student *s =(student*)arg;
21      printf("Students Id: %d,Name : %s, GPA: %.2f\n",s->student_id,s->name,s->gpa);
22      if(s->gpa >= 3.5f)
23      {
24          printf("%s made the Dean's List.\n", s->name);
25          dean_count++;
26      }
27      else{
28          printf("%s did not make the Dean's List.\n", s->name);
29      }
30      return NULL;
31  }
32
33
34  int main()
35  {
36      pthread_t tid[3];
37      student *s1 = malloc(sizeof(student));
38      student *s2 = malloc(sizeof(student));
39      student *s3 = malloc(sizeof(student));
40
41      s1->student_id = 101; strcpy(s1->name, "Dawar"); s1->gpa = 3.5f;
42      s2->student_id = 102; strcpy(s2->name, "Fawad"); s2->gpa = 3.1f;
43      s3->student_id = 103; strcpy(s3->name, "Ibrahim"); s3->gpa = 3.7f;
44
45      pthread_create(&tid[0], NULL, student_thread, s1);
46      pthread_create(&tid[1], NULL, student_thread, s2);
47      pthread_create(&tid[2], NULL, student_thread, s3);
48
49      for (int i = 0; i < 3; i++) {
50          pthread_join(tid[i], NULL);
51      }
52
53      printf("Dean Count: %d \nMain thread: Completed.\n",dean_count);
54
55      free(s1);
56      free(s2);
57      free(s3);
58      return 0;
59  }

```

## Output:



The screenshot shows a Visual Studio Code editor window titled "Assignment 01 [WSL: Ubuntu-24.04]". The editor has several tabs open, with "Task5.c" selected. The code in Task5.c is as follows:

```
1 //Ahmad Fawad
2 //=====
3 //Struct-Based Thread Communication
4
5 #include <stdio.h>
6 #include <stdlib.h>
7 #include <string.h>
8 #include <pthread.h>
9
10
11 int dean_count = 0;
12 typedef struct {
13     int student_id;
14     char name[50];
15     float gpa;
16 } student_t;
```

The terminal output at the bottom shows the execution of the program:

```
Fawad@LAPTOP-MBFKLMRL:~/Assignment 01$ gcc Task5.c -o exe5 -lpthread
Fawad@LAPTOP-MBFKLMRL:~/Assignment 01$ ./exe5
Students Id: 101, Name: Dawar, GPA: 3.50
Dawar made the Dean's List.
Students Id: 102, Name: Fawad, GPA: 3.10
Fawad did not make the Dean's List.
Students Id: 103, Name: Ibrahim, GPA: 3.70
Ibrahim made the Dean's List.
Dean Count: 2
Main thread: Completed.
Fawad@LAPTOP-MBFKLMRL:~/Assignment 01$
```

## Section-B: Short Questions

### Q1. Define an Operating System in a single line

Operating system is a bridge between the user and hardware. It can make sure that each program runs smoothly. It can manage the computer hardware and software resources.

### Q2. What is the primary function of the CPU scheduler?

Primary Function of CPU Scheduler is which process should run next on CPU. It can decide the best process according to the scheduling algorithm.

### Q3. List any three states of a process.

Three main states of processes are:

- 1: ready: its means process is ready to run but wait for CPU
- 2: Running: Process being running in CPU
- 3: waiting: process is blocked due to I/O.

### Q4. What is meant by a Process Control Block (PCB)?

A PCB is a data structure in the OS which can store the information about process like PID, State, Register and memory detail. It helps the OS to switch between processes.



### Q5. Differentiate between a process and a program.

A program is a set of instructions stored on disk while a process is the running state of that program in memory. When a program instruction starts execution it becomes a process.

### Q6. What do you understand by context switching?

When CPU wants to change one process to another, it can save the current process state and load the state of another process. This is known as context switching.

### Q7. Define CPU utilization and throughput.

CPU utilization means how much time the CPU is busy to complete the tasks while throughput is considered the number of tasks CPU completed in given time.

### Q8. What is the turnaround time of a process?

It is a total time of a process that it takes waiting for execution as well as its execution time.

### Q9. How is waiting time calculated in process scheduling?

It can represent the total time spent by the process in ready queue for CPU execution. So we can find it by subtracting the burst time from turnaround time.

### Q10. Define response time in CPU scheduling.

Response time is the first response from the CPU after submission of the task.

### Q11. What is preemptive scheduling?

Preemptive scheduling allows the CPU to be taken from a running process and given to another process with higher priority or shorter job first time.

### Q12. What is non-preemptive scheduling?

In non-preemptive scheduling, once a process starts using the CPU, it runs until it finishes or waits for I/O. No other process can interrupt it.

### Q13. State any two advantages of the Round Robin scheduling algorithm.

Every process gets an equal share of CPU time. It is suitable for time-sharing systems and provides quick responses for small tasks.

#### Q14. Mention one major drawback of the Shortest Job First (SJF) algorithm.

Its big drawback is that it can prefer the shortest jobs which caused starvation for the other processes

#### Q15. Define CPU idle time.

A time in which CPU is free . not a single process is ready to run

#### Q16. State two common goals of CPU scheduling algorithms.

CPU Utilization keep the cpu busy as much as possible  
minimize the waiting time improve system and user effecency

#### Q17. List two possible reasons for process termination.

- 1.Process can complete its task so terminate
2. due to an error, such as memory overflow or killed by OS

#### Q18. Explain the purpose of the wait() and exit() system calls.

The exit() call ends a process and returns its status to the OS. The wait() call allows the parent process to pause until its child process finishes execution.

#### Q19. Differentiate between shared memory and message-passing models of IPC.

Shared Memory	Message Passing
Processes communicate by sharing a common memory space.	Processes communicate by sending and receiving messages.
Faster communication	Slower due to kernel mediation for message transfer.
Requires synchronization .	Easier to synchronize; handled by OS.

#### Q20. Differentiate between a thread and a process.

Process	Thread
Independent program with its own memory space.	A smaller unit of a process sharing the same memory.
Heavyweight (fork)	Lightweight faster to create and switch (pthread).
Communication between processes is slower.	Threads communicate easily via shared memory.

### Q21. Define multithreading.

Running multiple threads in a single process to perform different tasks at the same time to improve the usage of CPU

### Q22. Explain the difference between a CPU-bound process and an I/O bound process.

A CPU-bound process mainly uses the processor for computation like calculations, while an I/O-bound process spends most time waiting for input/output operations to complete.

### Q23. What are the main responsibilities of the dispatcher?

The main task of dispatcher is to give the CPU control to that process which is selected by the scheduler by performing context switching.

### Q24. Define starvation and aging in process scheduling.

Starvation is when the process never gets the CPU for a long time due to low priority. Aging is a technique to increase the priority to avoid starvation.

### Q25. What is a time quantum (or time slice)?

It is a fixed time that is given to each process before switching to another process. It is used in round robin scheduling.

### Q26. What happens when the time quantum is too large or too small?

If too large, it behaves like FCFS and increases waiting time. If too small, too many context switches occur, which can waste CPU time.

### Q27. Define the turnaround ratio (TR/TS).

It is a ratio between total time process served

It shows how much total time a process took compared to its actual CPU service time.

### Q28. What is the purpose of a ready queue?

The ready queue holds all the processes which are ready and waiting for the CPU to be free. The scheduler can select the next process in the ready queue.

### Q29. Differentiate between a CPU burst and an I/O burst.

CPU Burst: the period when a process actively uses the CPU for computation.

I/O Burst: the period when a process waits for or performs input/output operations.

### Q30. Which scheduling algorithm is starvation-free, and why?

Round Robin is starvation free because every process gets an equal time share in a cyclic order so none are ignored.

### Q31. Outline the main steps involved in process creation in UNIX.

Main steps in process creation in UNIX:

1. `fork()` Creates a new child process.
2. `exec()` Replaces the child's memory with a new program.
3. `wait()` Parent waits for the child to finish.
4. `exit()` –Child terminates and releases resources.

### Q32. Define zombie and orphan processes.

A zombie process is which have completed its task and remain in execution list . and the orphan process is whos parents process complete his task before it without wating

### Q33. Differentiate between Priority Scheduling and Shortest Job First (SJF).

Priority scheduling selects the process with the highest priority, while SJF selects the one with the shortest CPU burst time. Both aim to improve efficiency.

### Q34. Define context switch time and explain why it is considered overhead.

Context switch time is the time spent saving and loading process information during switching. It's considered overhead because no actual processing work happens then.

### Q35. List and briefly describe the three levels of schedulers in an OS.

- LongTerm Scheduler Selects which processes are admitted into the ready queue.
- ShortTerm Scheduler Chooses which ready process runs next on the CPU.
- MediumTerm Scheduler Temporarily suspends or resumes processes to control multitasking and memory use

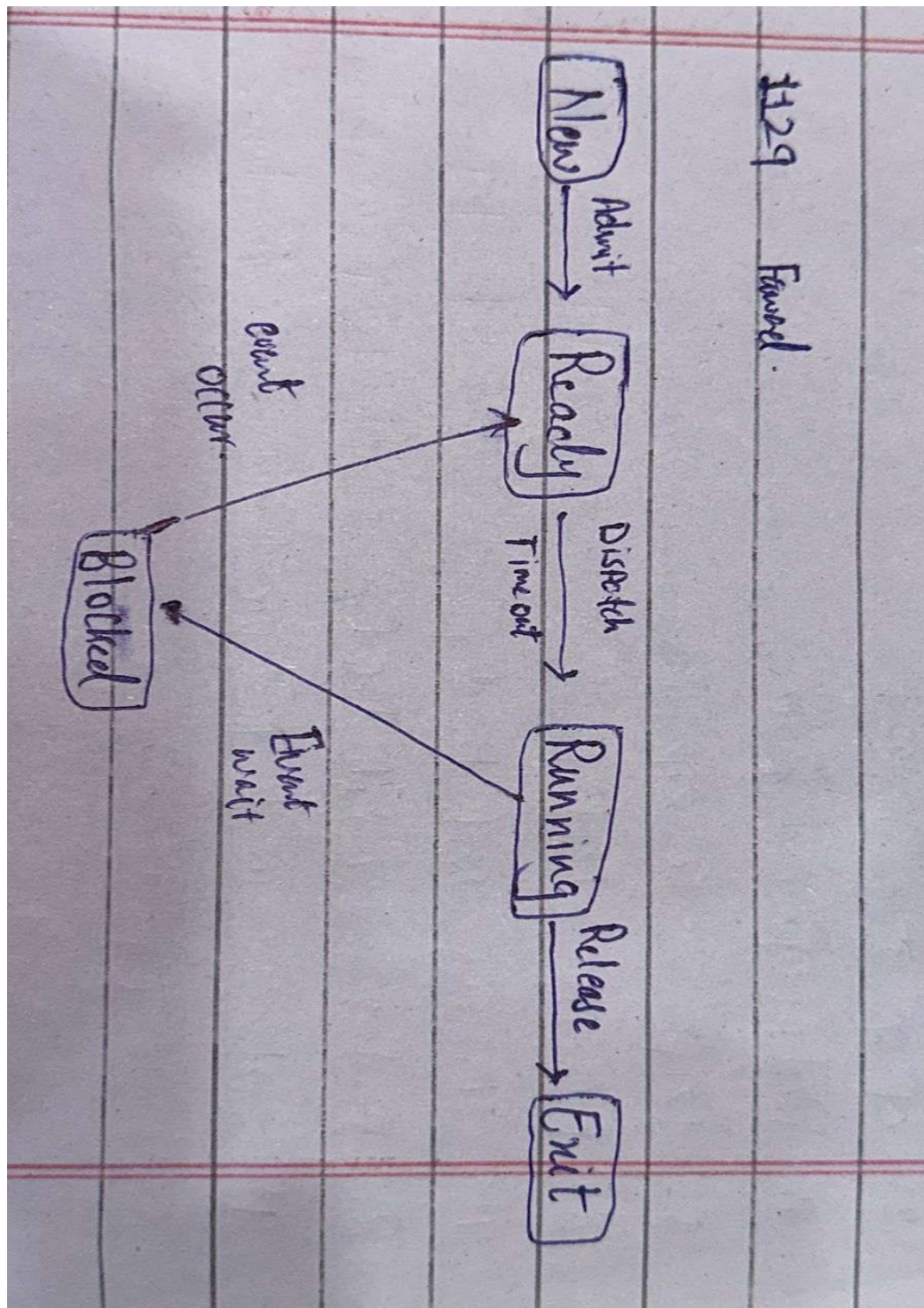
### Q36. Differentiate between User Mode and Kernel Mode in an OS.

User mode is where normal user programs run with limited access. Kernel mode is a privileged mode used by the OS to control hardware and perform critical operations.

## Section-C: Technical / Analytical Questions

1. Describe the complete life cycle of a process with a neat diagram showing transitions between New, Ready, Running, Waiting, and Terminated states.

- New: Process is being created.
- Ready: Process is ready to run and waiting in ready queue.
- Running: CPU is executing the process.
- Waiting/Blocked: Process waits for I/O or event.
- Terminated: Process finished or killed.



## 2. Write a short note on context switch overhead and describe what information must be saved and restored.

Context switch overhead is the extra time the CPU spends saving the state of the current process and loading the state of the next process, during which no productive work of those processes is done. Minimizing unnecessary switches improves performance.

What must be saved / restore

- Context of Processor.

- Process state.
- Memory management.
- Scheduling.

### 3. List and explain the components of a Process Control Block (PCB).

1. Process state: running, waiting
2. Program Id: unique identifier of each process
3. Program counter: Address of next instruction to execute
4. Cpu Register: store process data during execution
5. Memory management: Details like base and limit register
6. Accounting Info: CPU usage, process priority
7. I/O Status: List i/o devices used by process

### 4. Long-Term, Medium-Term, Short-Term Schedulers

Long-Term Scheduler	Short-Term Scheduler	Medium-Term Scheduler
Selects processes from secondary storage (job pool) and loads them into memory.	Selects one ready process for CPU execution.	Temporarily removes and later reintroduces processes to control load.
Slowest of all schedulers.	Fastest; runs most frequently.	Intermediate speed.
Rare or absent in time-sharing systems.	Always present; key in time-sharing systems.	Used mainly in time-sharing systems for swapping.
Admits new jobs into memory.	Dispatches ready processes to CPU.	Suspends and later resumes processes.

### 5. Explain CPU Scheduling Criteria (Utilization, Throughput, Turnaround, Waiting, and Response) and their optimization goals.

1. CPU Utilization: keep CPU busy as much as possible. Goal: high %age
2. Throughput: processes completed per time unit. Goal: maximize
3. Turnaround Time: finish time - arrival time. Goal: minimize average
4. Waiting Time: total time in ready queue. Goal: minimize average
5. Response Time: time from submission to first response. Goal: minimize, important for interactive systems

## Section-D: CPU Scheduling Calculations

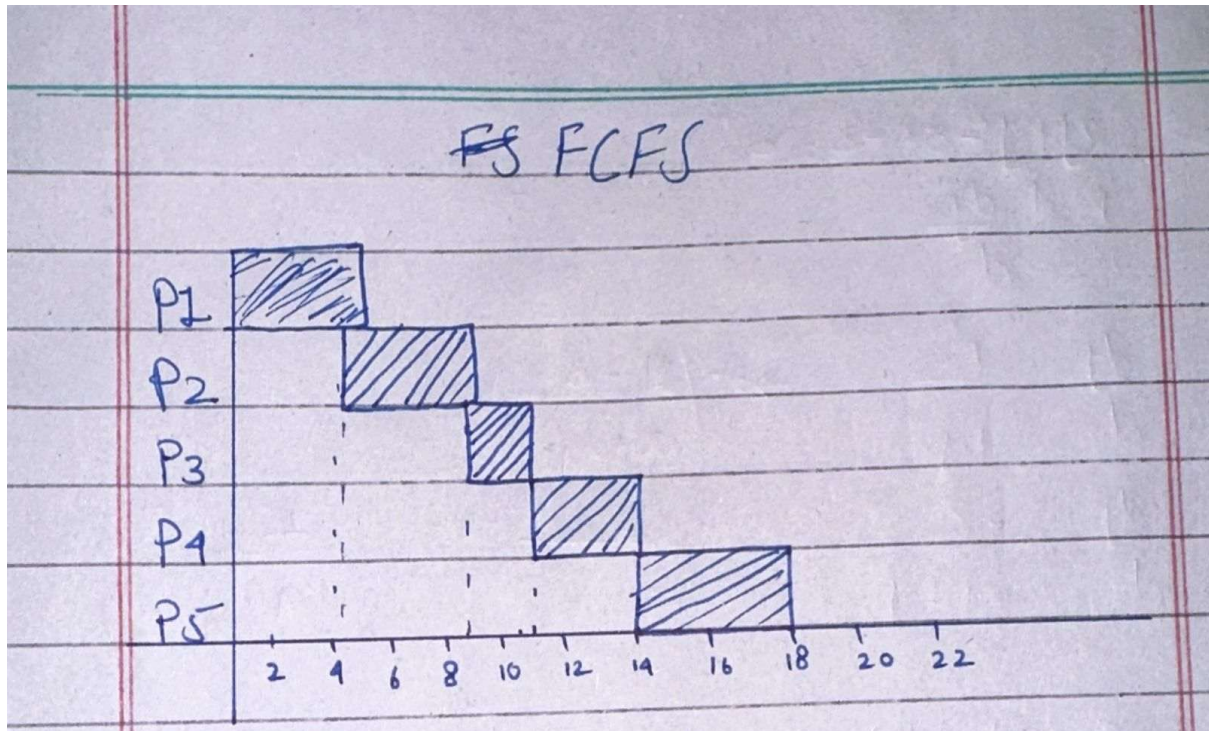
- Perform the following calculations for each part (A–C).
- a) Draw Gantt charts for FCFS, RR (Q=4), SJF, and SRTF.
- b) Compute Waiting Time, Turnaround Time, TR/TS ratio, and CPU Idle Time.
- c) Compare average values and identify which algorithm performs best.



## Part-A:

Process	Arrival Time	Service Time
P1	0	4
P2	2	5
P3	4	2
P4	6	3
P5	9	4

### Gant chart:



### Completion, Turn Around, Waiting:

- P1:** Finish=4, Turn Around=4, Waiting=0
- P2:** Finish=9, Turn Around=7, Waiting=2
- P3:** Finish=11, Turn Around=7, Waiting=5
- P4:** Finish=14, Turn Around=8, Waiting=5
- P5:** Finish=18, Turn Around=9, Waiting=5

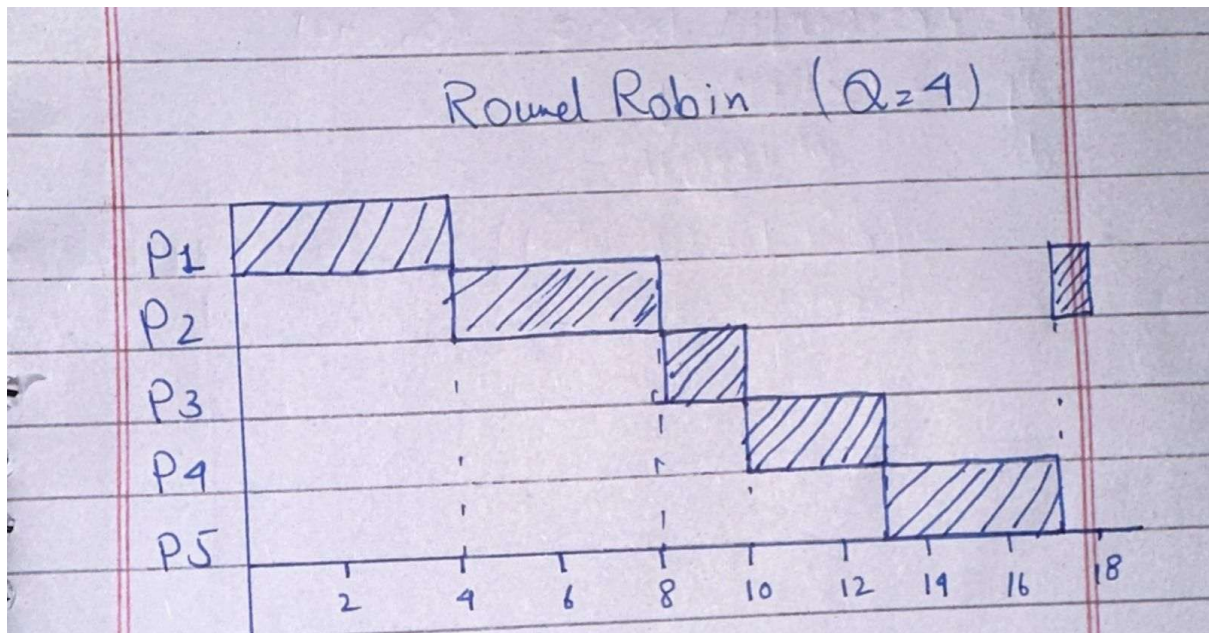
### Averages:

- Avg Wait Time**=  $(0+2+5+5+5)/5 = 3.4$
- Avg Turn Around**=  $(4+7+7+8+9)/5 = 7.0$
- Avg TR/TS** = 2.16
- CPU Idle** = 0

## Round Robin:

Gant Chart:





**Completion, Turn Around, Waiting:**

**P1:** Finish=4, Turn Around=4, Waiting=0

**P2:** Finish=18, Turn Around=16, Waiting=11

**P3:** Finish=10, Turn Around=6, Waiting=4

**P4:** Finish=13, Turn Around=7, Waiting=4

**P5:** Finish=17, Turn Around=8, Waiting=4

**Averages:**

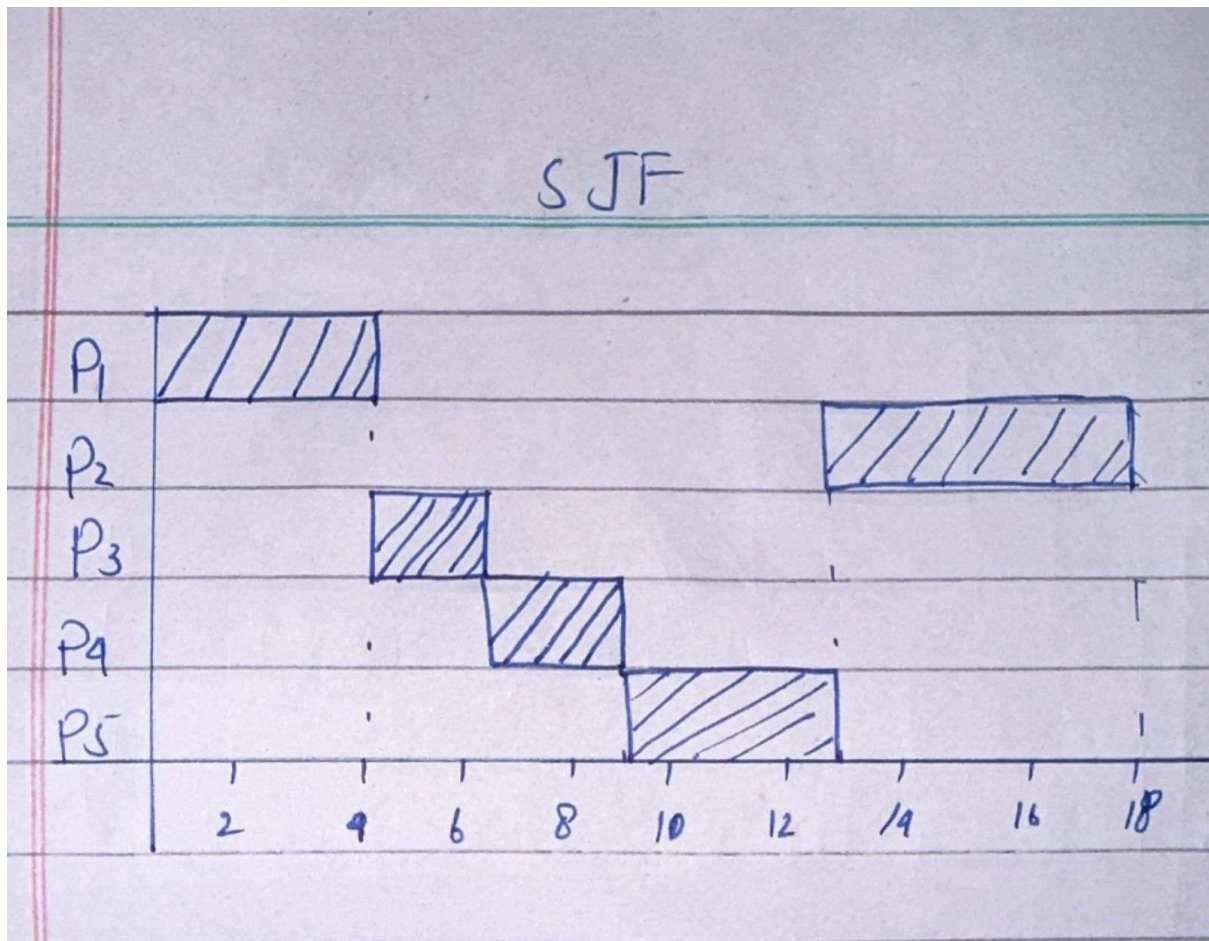
**Avg Wait Time** =  $(0+11+4+4+4)/5 = 4.6$

**Avg Turn Around** =  $(4+16+6+7+8)/5 = 8.2$

**Avg TR/TS** = 2.31

**CPU Idle** = 0

**Shortest Job First:**



**Completion, Turn Around, Waiting:**

**P1:** Finish=4, Turn Around=4, Waiting=0  
**P2:** Finish=18, Turn Around=16, Waiting=11  
**P3:** Finish=6, Turn Around=2, Waiting=0  
**P4:** Finish=9, Turn Around=3, Waiting=0  
**P5:** Finish=13, Turn Around=4, Waiting=0

**Averages:**

**Avg Wait Time** =  $(0+11+0+0+0)/5 = 2.2$   
**Avg Turn Around** =  $(4+16+2+3+4)/5 = 5.8$   
**Avg TR/TS** = 1.44  
**CPU Idle** = 0

**Conclusion (Part-A):**

SJF/SRTF gives best average turnaround and smallest average waiting time for these arrivals. RR improves response but increases average turnaround vs SJF. FCFS is

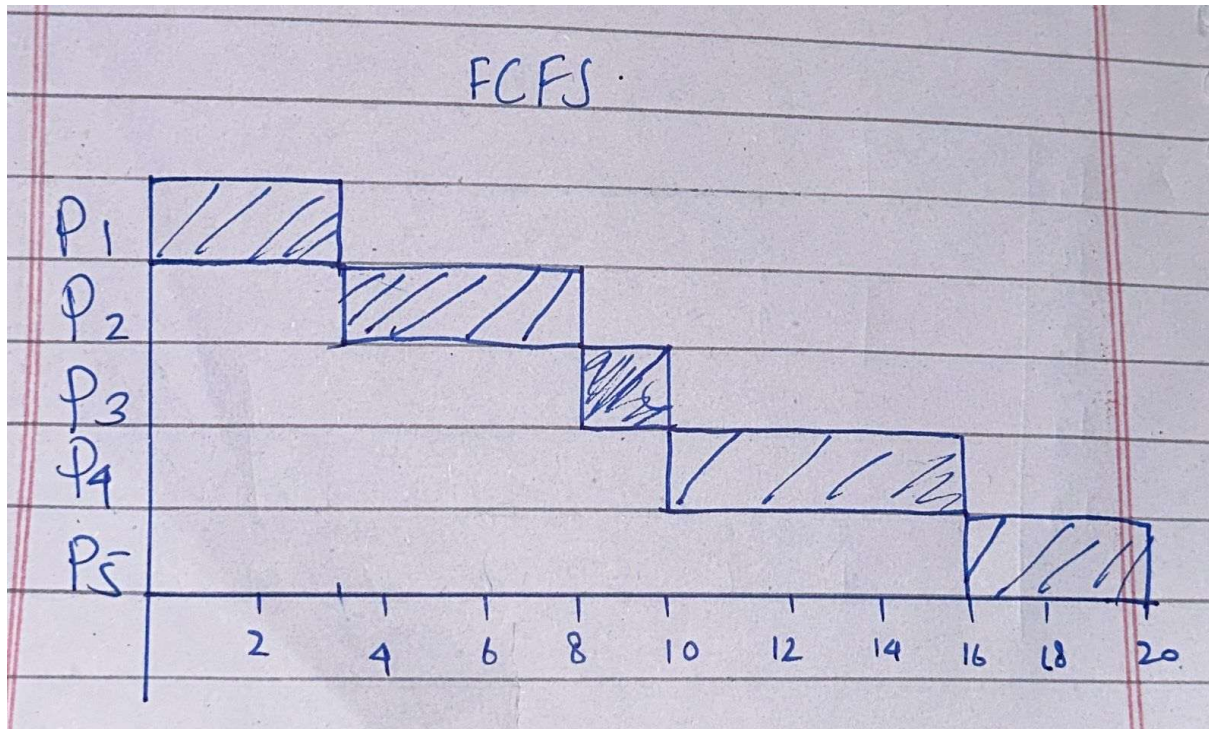
**Part B:**

Process	Arrival Time	Service Time (Burst Time)
P1	0	3
P2	1	5

P3	3	2
P4	9	6
P5	10	4

## FCFS:

Gant chart:



### Completion, Turn Around, Waiting:

**P1:** Finish=3, Turn Around=3, Waiting=0

**P2:** Finish=8, Turn Around=7, Waiting=2

**P3:** Finish=10, Turn Around=7, Waiting=5

**P4:** Finish=16, Turn Around=7, Waiting=11

**P5:** Finish=20, Turn Around=10, Waiting=6

### Averages:

**Avg Wait Time** =  $(0+2+5+11+6)/5 = 4.8$

**Avg Turn Around** =  $(3+7+7+7+10)/5 = 6.8$

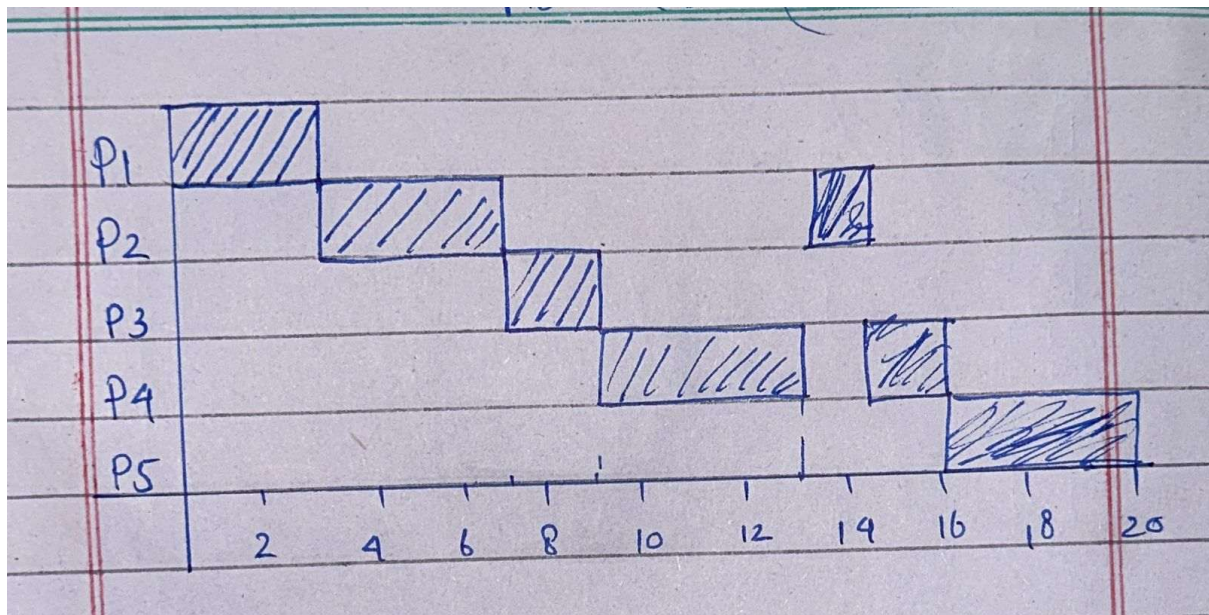
**Avg TR/TS** = 1.91

**CPU Idle** = 0

## Round Robin:

Gant chart:





#### Completion, Turn Around, Waiting:

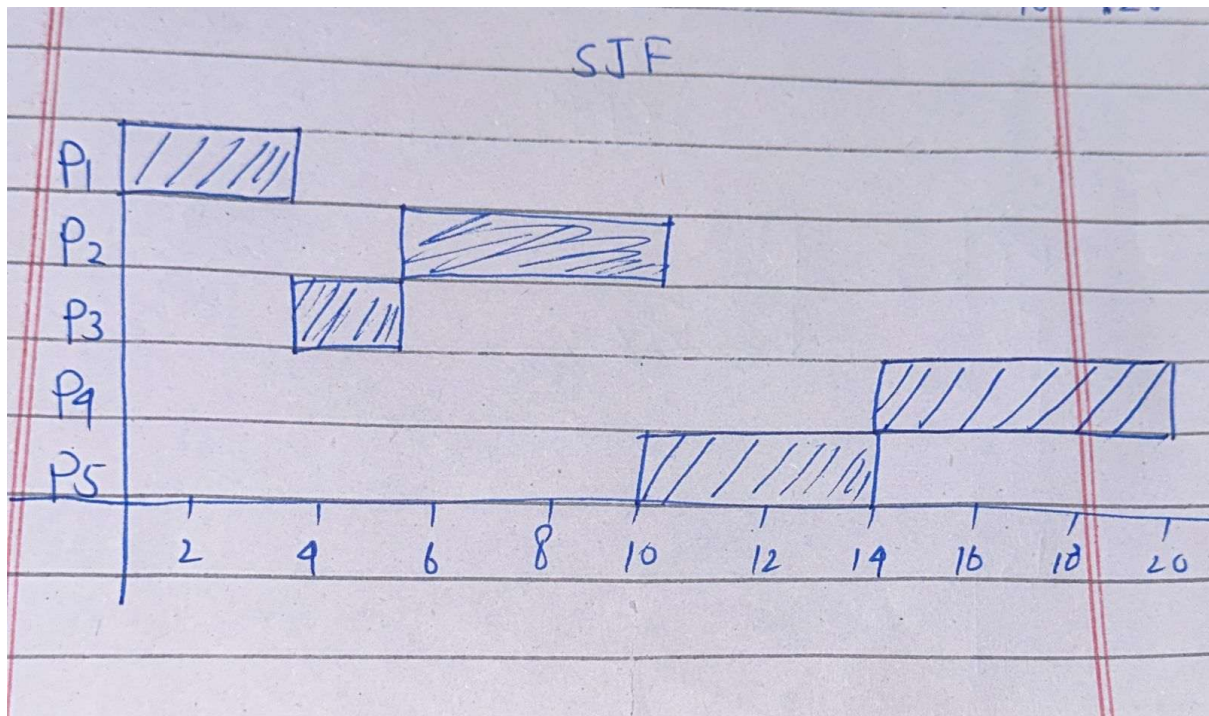
**P1:** Finish=3, Turn Around=3, Waiting=0  
**P2:** Finish=14, Turn Around=13, Waiting=8  
**P3:** Finish=9, Turn Around=6, Waiting=4  
**P4:** Finish=16, Turn Around=7, Waiting=1  
**P5:** Finish=20, Turn Around=10, Waiting=6

#### Averages:

**Avg Wait Time** =  $(0+8+4+1+6)/5 = 3.8$   
**Avg Turn Around** =  $(3+13+6+7+10)/5 = 7.8$   
**Avg TR/TS** = 2.05  
**CPU Idle** = 0

#### Shoertest job First:

Gant chart:



#### Completion, Turn Around, Waiting:

**P1:** Finish=3, Turn Around=3, Waiting=0  
**P2:** Finish=10, Turn Around=9, Waiting=4  
**P3:** Finish=5, Turn Around=2, Waiting=0  
**P4:** Finish=20, Turn Around=11, Waiting=5  
**P5:** Finish=14, Turn Around=4, Waiting=6

#### Averages:

**Avg Wait Time** =  $(0+4+0+5+0)/5 = 1.8$   
**Avg Turn Around** =  $(3+9+2+11+4)/5 = 5.8$   
**Avg TR/TS** = 1.33  
**CPU Idle** = 0

#### Conclusion:

SJF gives best average waiting & turnaround; RR gives fairness but higher average turnaround than SJF.

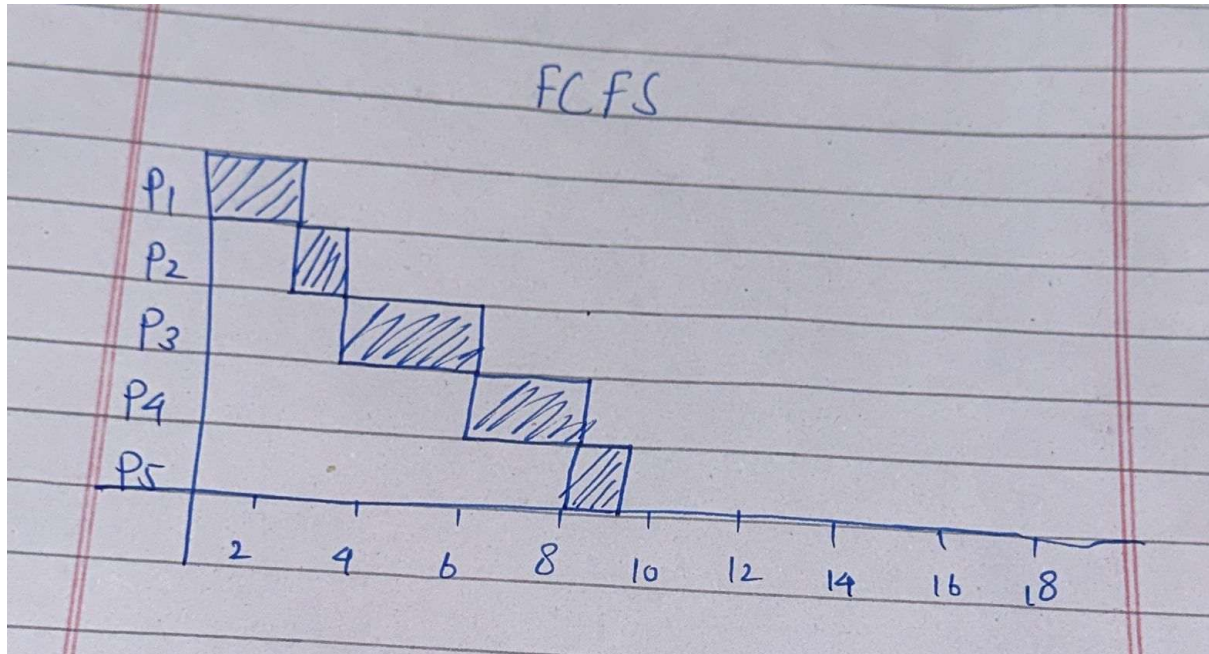
## Part C

#### Process Arrival Time Service Time (Burst)

P1	0	2
P2	1	1
P3	2	3
P4	3	2
P5	4	1

## FCFS:

Gant chart:



### Completion, Turn Around, Waiting:

**P1:** Finish=2, Turn Around=2, Waiting=0

**P2:** Finish=3, Turn Around=3, Waiting=1

**P3:** Finish=6, Turn Around=4, Waiting=1

**P4:** Finish=8, Turn Around=5, Waiting=3

**P5:** Finish=9, Turn Around=5, Waiting=4

### Averages:

**Avg Wait Time**=  $(0+1+1+3+4)/5 = 1.8$

**Avg Turn Around**=  $(2+4+4+5+5)/5 = 3.6$

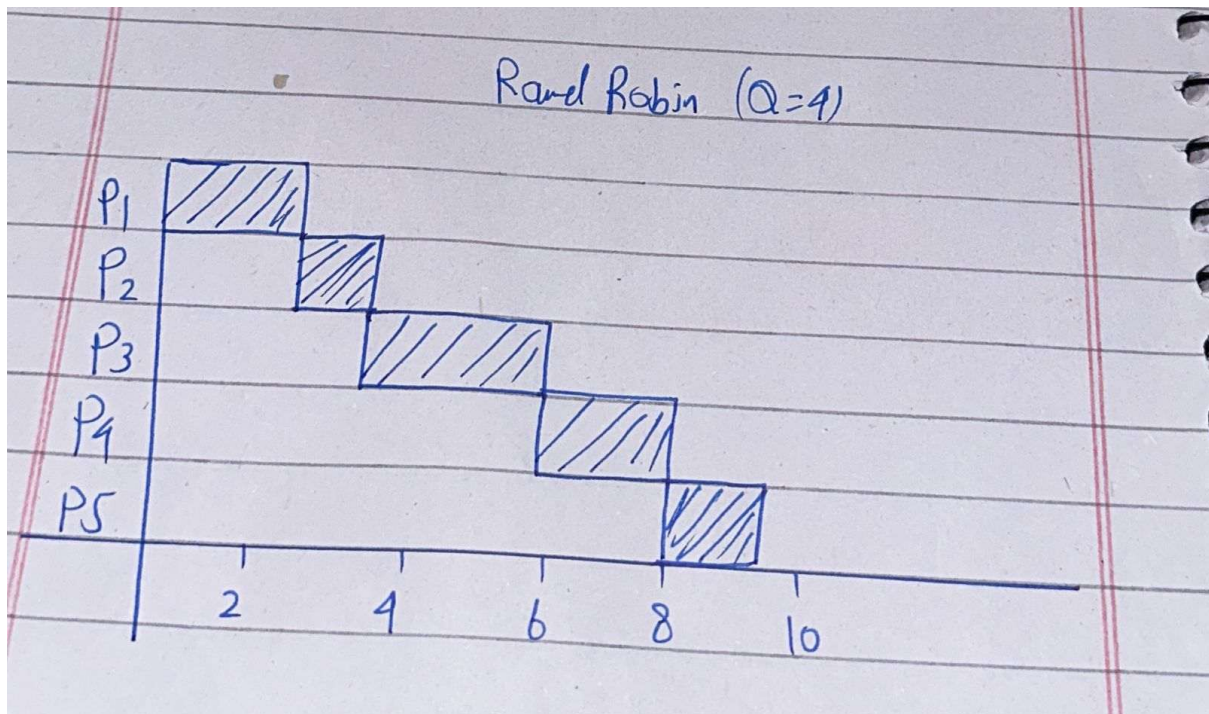
**Avg TR/TS** = 2.37

**CPU Idle** = 0

## Round Robin :

Gant chart:





**Completion, Turn Around, Waiting:**

**P1:** Finish=2, Turn Around=2, Waiting=0

**P2:** Finish=3, Turn Around=3, Waiting=1

**P3:** Finish=6, Turn Around=4, Waiting=1

**P4:** Finish=8, Turn Around=5, Waiting=3

**P5:** Finish=9, Turn Around=5, Waiting=4

**Averages:**

**Avg Wait Time** =  $(0+1+1+3+4)/5 = 1.8$

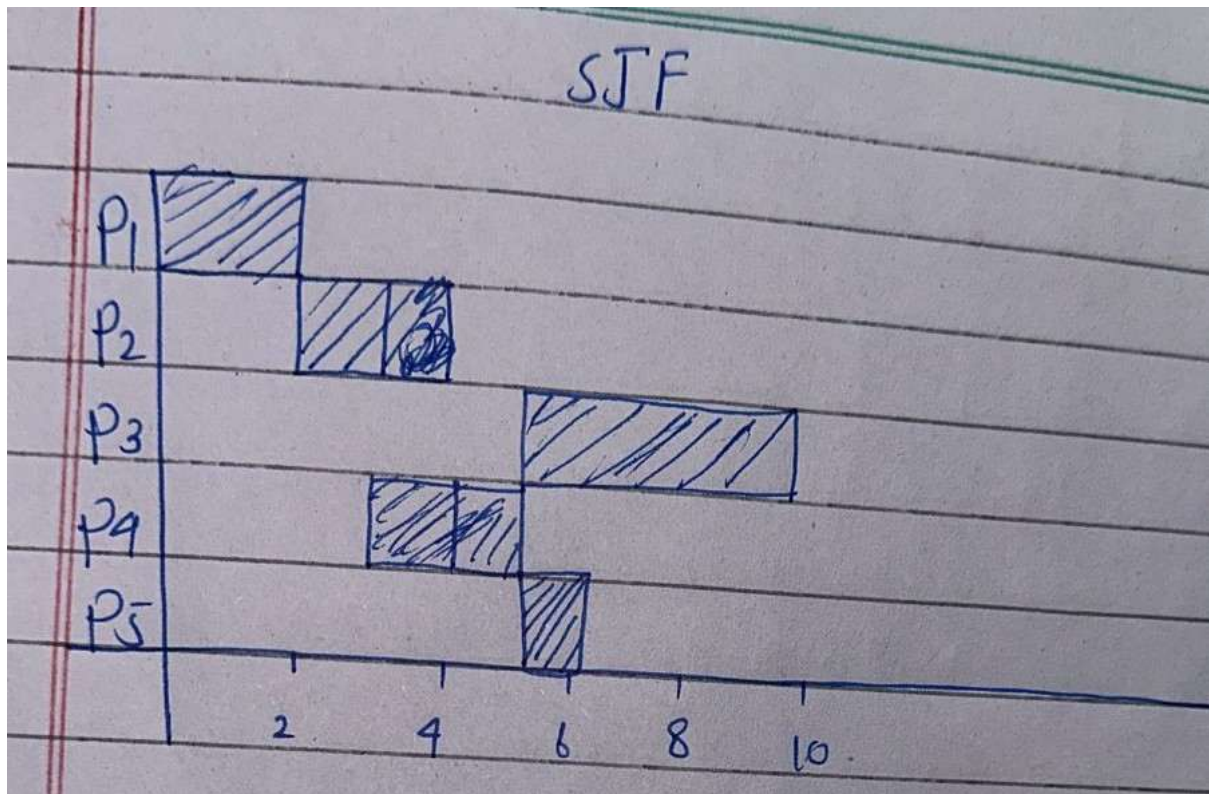
**Avg Turn Around** =  $(2+3+4+5+5)/5 = 3.6$

**Avg TR/TS** = 2.37

**CPU Idle** = 0

**SJF:**

Gant chart:



#### Completion, Turn Around, Waiting:

**P1:** Finish=2, Turn Around=2, Waiting=0

**P2:** Finish=3, Turn Around=2, Waiting=1

**P3:** Finish=9, Turn Around=7, Waiting=4

**P4:** Finish=5, Turn Around=2, Waiting=0

**P5:** Finish=6, Turn Around=2, Waiting=1

#### Averages:

**Avg Wait Time** =  $(0+1+4+0+1)/5 = 1.2$

**Avg Turn Around** =  $(2+2+7+2+2)/5 = 3.0$

**Avg TR/TS** = 2.5

**CPU Idle** = 0

#### Conclusion:

with these sample times SJF gives lower waiting times than FCFS.