| Name: Section: _ | |
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National University of Computer and Emerging Sciences, Lahore Campus



| Cours | se: | Computer Organization and Assembly | Course Code: | EE2003 |
|---------|---------|------------------------------------|--------------|-----------|
| | | Language | Semester: | Fall 2021 |
| Progr | am: | BS (CS, DS) | Total Marks: | 100 |
| Durat | tion: | 3 hours | Weightage: | 45 |
| Pape | r Date: | 7-Jan-2022 | Page(s): | 16 |
| Section | on(s): | All | Section: | |
| Exam | : | Final exam | Roll No: | |

Instruction/Notes:

- Exam is Open book, Open notes.
- Properly comment your code.
- You CANNOT use an instruction NOT taught in class.
- If there is any ambiguity, make a reasonable assumption. Questions during the exam are not allowed.
- Write your answer in the space provided. You can take extra sheets BUT they WON'T
 BE ATTACHED WITH THE QUESTION PAPER OR MARKED.
- All other rules pertaining to examinations as per NUCES policy apply.

Question 1 [45 Marks]:

Show your working here:

i. (4 marks): What are the effective and physical addresses generated by the following memory access?

| Memory access | Effective Address | Physical Address |
|---------------|-------------------|------------------|
| [cs: bx + di] | | |

Given: BX=00FFh, CS = 1111h, DS = 3333h, SS = 2526h, IP = 1232h, SP = 1100h, and DI = 0020h

| Method 1: incre | ment bx to advance to | Method 2: use bx with displacements to access each value |
|-----------------|-----------------------|--|
| each value | | |
| | | |
| List db 10h, 20 | 0h, 30h, 40h | |
| sum db 0 | | |
| | | |
| mov bx, List | | |
| mov al, [bx] | ; AL = 10h | |
| inc bx | ; BX points to 20h | |
| add al, [bx] | ; AL = 30h | |
| inc bx | ; BX points to 30h | |
| add al, [bx] | ; AL = 60h | |
| inc bx | ; BX points to 40h | |
| add al, [bx] | ; AL = 0A0h | |
| mov si, sum | ; SI points to sum | |
| mov [si], al | ; SUM = 0A0h | |
| | | |
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ame: _____ Roll Number: ____ Section: _____ ii. (6 marks): The assembly code is provided in method 1 (column 1) to calculate the sum of all the elements of an array?

Optimize (with respect to number of lines) this program by using displacement addressing modes.

iii. (6 marks): Suppose that AX=0x3412, BX=0x7856, CX= 0x1CAB, and SP=0x100. Give the contents of AX, BX, and SP after executing the following instructions:

| | АХ | ВХ | SP |
|-------------|----|----|----|
| push ax | | | |
| push bx | | | |
| xchg ax, cx | | | |
| рор сх | | | |
| push ax | | | |
| рор bx | | | |

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| iv. | (7 marks): AX contains a number between 0-15. Write code to complement the corresponding bit in BX. For example, if AX contains 6; complement the 6th bit of BX. (Note: First bit in BX is at 0th position and last bit is at 15 th position). Hint: Use Bit Manipulation |
| | Time. Ose Die Waliipulation |
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| (4 m | narks): Given bel | ow the listing file of a code. | What is the size of the code i.ecom file? | |
| | 1 | _ | [org 0x100] | |
| | 2 | | [019 0.1100] | |
| | 3 00000000 2 | A1[1700] | mov ax, [num1] | |
| | 4 00000003 | | mov bx, [num2] | |
| | 5 00000007 | 01D8 | add ax, bx | |
| | 6 00000009 | | mov bx, [num3] | |
| | 7 000000D | | add ax, bx | |
| | 8 000000F 2 | A3[1D00] | mov [num4], ax | |
| | 9 | | | |
| | .0 00000012 2 | | mov ax,0x4c00 | |
| | .1 00000015 | CD21 | int 0x21 | |
| | .2 | | | |
| | .3 | | | |
| | .4 | 0.5.0.0 | | |
| | .5 00000017 | | num1: dw 5 | |
| | .6 00000019 (.7 0000001B (| | num2: dw 10 num3: dw 15 | |
| | .8 0000001b (| | num4: dw 0 | |
| 1 | .0 0000001D (| 0000 | Hama. aw o | |
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| (4 m | narks): Find the v | value of the Carry flag after t | the execution of the following code. | |
| • | • | , 3 | G | |
| | | | | |
| 1. | a) mov ax, 85h | | | |
| (0 | • | | | |
| | add ax, 92h | | | |
| | | | | |
| (k | b) mov ax, 15h | | | |
| | add ax, 72h | | | |
| | | | | |

| [org 0x0100] | | | timer: | push ax |
|------------------|---------------|------------------------|-------------|--------------------------------|
| | | | | push bx |
| | jmp s | tart | | push dx |
| | | | | cmp word[cs:timerflag],1 |
| seconds: | dw | 0 | | jne skipall |
| timerflag: | dw | 0 | | |
| oldkb: | dd | 0 | | inc word [cs:seconds] |
| | | | | mov ax, [cs:seconds] |
| printnum: | | | | mov dx, 0 |
| ; copy from Li | sting 9.7 (| lines 9-47) | | mov bx, 2 |
| | | | | div bx |
| kbisr: | push | ax | | cmp dx, 0 |
| | | | | je skipall |
| | in al, | | | push word [cs:seconds] |
| | | al, 0x36 | | call printnum |
| | jne ne | extcmp | | |
| | | | skipall: | mov al, 0x20 |
| | | word [cs:timerflag], 1 | | out 0x20, al |
| | je exit | t | | |
| | | | | pop dx |
| | | word [cs:timerflag], 1 | | pop bx |
| | jmp e | xit | | pop ax |
| | | 1016 | | iret |
| nextcmp: | | al, 0xb6 | | |
| | Jne no | omatch | start: | Listing O. O. //inco O.F. 113) |
| | mou | word [costimorflag] 0 | ; copy from | n Listing 9.8 (lines 95-113) |
| | jmp e | word [cs:timerflag], 0 | | |
| | Jilip e | XIL | | |
| nomatch: | рор а | x | | |
| | | ar [cs:oldkb] | | |
| | Jb | [00:0:0::0::0] | | |
| exit: | mov a | al, 0x20 | | |
| | | x20, al | | |
| | | , | | |
| | рор а | х | | |
| | iret | | | |
| | | | | |
| | | ress code = 0x36 | | |
| ; hint: right sh | nift key's re | elease code = 0xb6 | | |
| | | | | |
| ; (code is cont | tinued in t | he second column) | | |
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| viii. | (7 marks): You are given a piece of code and information pressed during the execution of this code. Considering the write out the sequence in which the instructions are executions are executions are executions. | he code and the occurrences of interrupts/keystroke as | given, |
| | Sample answer: Instructions executed in following order | | |

I6 I10

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| I1 | jmp start | Write your Answer here | |
|-----|---|-------------------------|--|
| | tickcount: dw 0 | that is the sequence in | |
| | | which instructions | |
| | kbisr: | executed | |
| 12 | in al, 0x60 | | |
| 13 | cmp al, 0x2a | | |
| 14 | jne end | | |
| 15 | mov bl, 0x2a | | |
| | | | |
| 16 | end: | | |
| 16 | mov al, 0x20 | | |
| 17 | out 0x20, al | | |
| 18 | iret | | |
| | timerISR: | | |
| | ;assume that keyboard was pressed by user at this point | | |
| 19 | push ax | | |
| 110 | inc word [cs:tickcount]; increment tick count | | |
| 111 | mov al, 0x20 | | |
| 112 | out 0x20, al ; end of interrupt pop ax | | |
| 113 | iret | | |
| | | | |
| | start: | | |
| 114 | xor ax, ax | | |
| 115 | mov es, ax | | |
| 116 | cli | | |
| 117 | mov word [es:9*4], kbisr | | |
| 118 | mov [es:9*4+2], cs | | |
| 119 | mov word [es:8*4], timerISR | | |
| 120 | mov [es:8*4+2], cs | | |
| 121 | ;int 8h occurred here sti | | |
| 121 | mov ax, 20 | | |
| 123 | mov bx,15 | | |
| 124 | add ax, bx | | |
| 125 | mov ax, 0x4C00 | | |
| 126 | int 0x21 | | |
| 120 | III OAZI | <u> </u> | |

| | e: Roll Number: ion 2 [25 Marks] | | Section: | |
|------|--|-------------------------|------------------------------|------------------------|
| i. | (3 marks): Increasing the number of pipeline s | | clock cycle time. However | , give a reason why |
| | processors should not have hundreds or thousan | ids of pipeline stages. | | |
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| ii. | (4 marks): Identify all data dependencies in the find. For example, if I3 depends on register di from | | | |
| | the table entry blank. | · | | |
| | Instruction | Depends on Regis | ter from |] |
| | I1: mov ax, [bx] I2: mov bx, [bp] | | | |
| | I3: add ax, bx | | |] |
| | I4: mov [bx], ax | | |] |
| iii. | (4 marks): Find at least 4 possible data hazards (\ | WAW. RAW. WAR) whi | ich may occur in the instruc | tions given for O2(ii) |
| | Use following method to write hazard between t | | | , , |
| | AW: Instruction X & Instruction Y, Instruction X & I | | | |
| | /AR: Instruction A & Instruction B, Instruction C & I /AW: Instruction X & Instruction L, Instruction M & | | | |
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| and I2) is already filled out. We will assume that there are no data dependencies and no branch hazard detection mechanism is used. Clock Cycle | Name | e: | | | | Roll N | Number: | | | | Se | ction: | | | | |
|---|--|--------------------|---------|----------|-----------|-----------------|------------|-----------|-----------|------------|-------------|----------|---------|-----------|-----------|----------|
| which the third instruction is a conditional branch to instruction 15. The schedule for the first two instructions (i.e., I1 and I2) is already filled out. We will assume that there are no data dependencies and no branch hazard detection mechanism is used. Clock Cycle | | | marks): | Assume | a pipeli | | | | instruc | tion (FI), | _ decode | instruct | ion and | calculate | address | es (DA), |
| which the third instruction is a conditional branch to instruction 15. The schedule for the first two instructions (i.e., 11 and I2) is already filled out. We will assume that there are no data dependencies and no branch hazard detection mechanism is used. Clock Cycle | fetch operand (FO) and execute (EX). Complete the pipeline schedule given below for a sequence of 7 instru | | | | | | | | | | | | | | | |
| and I2) is already filled out. We will assume that there are no data dependencies and no branch hazard detection mechanism is used. Clock Cycle | | | | | | | | | | | | | | | | |
| Took Cycle 1 2 3 4 5 6 7 8 9 10 11 12 13 11 FI DA FO EX 12 FI DA FO EX 13 FI DA FO EX 14 FO EX 15 FI DA FO EX 16 FI DA FO EX 17 FI DA FO EX 18 FI DA FO EX 19 FI DA FO EX 10 FI DA FO EX 10 FI DA FO EX 11 FI DA FO EX 12 FI DA FO EX 13 FI DA FO EX 14 FI DA FO EX 15 FI DA FO EX 16 FI DA FO EX 17 FI DA FO EX 18 FI DA FO EX 19 FI DA FO EX 19 FI DA FO EX 10 FI DA FO EX 10 FI DA FO EX 10 FI DA FO EX 11 FI DA FO EX 12 FI DA FO EX 13 FI DA FO EX 14 FI DA FO EX 15 FI DA FO EX 16 FI DA FO EX 17 FI DA FO EX 18 FI DA FO EX 19 FI DA FO EX 10 FI DA FO EX 11 FI DA FO EX 12 FI DA FO EX 13 FI DA FO EX 14 FI DA FO EX 15 FI DA FO EX 16 FI DA FO EX 17 FI DA FO EX 18 FI DA FO EX 19 FI DA FO EX 10 FI DA | | | | | | | | | | | | | | | | |
| V. (2 marks): There are 128 blocks in a cache memory, which can store one word each. To which block number does main memory word address number 900 would map in the case of a direct mapped cache? | | | | | | t. VVC V | viii assai | iic tiiat | tricic ar | . no dat | a acpen | acricics | ana no | Dianen i | iazara a | ctcction |
| 1 | | mechanism is used. | | | | | | | | | | | | | | |
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| v. (2 marks): There are 128 blocks in a cache memory, which can store one word each. To which block number does main memory word address number 900 would map in the case of a direct mapped cache? | | | | | | | | | | | | | 10 | | 40 | 40 |
| v. (2 marks): There are 128 blocks in a cache memory, which can store one word each. To which block number does main memory word address number 900 would map in the case of a direct mapped cache? | | | | | | | | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 |
| v. (2 marks): There are 128 blocks in a cache memory, which can store one word each. To which block number does main memory word address number 900 would map in the case of a direct mapped cache? | | - | | FI | | | | ΓV | | | | | | | | |
| v. (2 marks): There are 128 blocks in a cache memory, which can store one word each. To which block number does main memory word address number 900 would map in the case of a direct mapped cache? | | 5 | 12 | | FI | DA | FU | EX | | | | | | | | |
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| memory word address number 900 would map in the case of a direct mapped cache? | | اقا | | | | | | | | | | | | | | |
| memory word address number 900 would map in the case of a direct mapped cache? | | } | | | | | | | | | | | | | | |
| memory word address number 900 would map in the case of a direct mapped cache? | | | | | ļ | | | <u> </u> | ļ | | | | ļ | | | |
| memory word address number 900 would map in the case of a direct mapped cache? | ., | (2. | marks). | Thora a | ro 120 hi | ocks in a | s cacha n | aamaru | which co | n store | ano wor | d oach I | o which | block pr | ımbar de | oc main |
| | v. | | | | | | | | | | | | | DIOCK III | iiibei uc | es main |
| vi. (2 marks): Which replacement policy can be used to update an n-way set associative cache? | | me | emory w | ord add | ress nun | nber 900 | would n | nap in th | e case o | r a direct | mappe | d cache: | , | | | |
| vi. (2 marks): Which replacement policy can be used to update an n-way set associative cache? | | | | | | | | | | | | | | | | |
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| VI. (2 marks): which replacement policy can be used to update an n-way set associative cacher | | /2 | | المائمان | | نامم همم | | | | | | !-+! | | | | |
| | VI. | (2) | marks): | wnich r | replacem | ient poli | cy can be | e usea to | update | an n-way | y set ass | ociative | cacne? | | | |
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| vii. | (5 marks): Map following physical addresses of a RAM to 2-way set | associative cache and complete the table given |
| | below. Replacement algorithm is Least Recently Used (LRU) and the | following block access sequence is used: |

0, 24, 0, 3, 24

| Block | Cache index | Hit/miss | Cache content after access | | | |
|---------|----------------|----------|----------------------------|-------|--|--|
| address | | | Set 0 | Set 1 | | |
| 0 | | | | | | |
| 24 | | | | | | |
| 0 | | | | | | |
| 33 | | | | | | |
| 24 | | | | | | |

| Name: | Roll Number: Section: |
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| | n 3 [30 Marks]: You are required to implement a game Bomb_Diffuser with the following requirements. |
| i. | Game total time is 1000 timer ticks. (5 marks) |
| ii. | Random 15 characters from (A-0) will be placed on screen at the start of the game (Assume you have a |
| | characterRandom function that return x and y coordinate in register ah and al. You have to maintain an array of 15 |
| | words. When you run rand function first time it gives coordinates for character A, 2 nd time it gives coordinates for |
| | character B and then so on). |
| iii. | Out of these 15 characters, 3 characters contains bomb (Assume you have an array of 3 random characters between |
| | A-0. You are required name it as bomb_arr . Hardcoded but can be any characters so in code you have to find the bomb |
| | by traversing this 3-character array) |
| i. | |
| iv. | Once game have started when a user presses a key other than A-O nothing will happen on the screen. No need of Scan |
| | key comparison in the code just convert scan key to ASCII and use ASCII for checking valid characters in hooked ISR. |
| | Assume you have a function named ScanToAscii that converts scan key to ASCII and return ASCII in AX. ASCII values A- |
| | 0 in decimal (65-79) and hex (41h-4Fh). |
| ٧. | If user presses key that contains bomb game will end. And an appropriate message along with score will be displayed. |
| vi. | If user presses key that doesn't contain bomb character, that character will be removed by making it black and 10 |
| | points will be incremented in the score. (points ii-vi have 8 marks) |
| vii. | Score needs to be updated live. You have to display timer ticks and score in first row of screen (4 marks) |
| viii. | Game will end in three cases. 1) Timer tick becomes greater than 1000 2) Bomb character pressed 3) Only Bomb |
| | Characters left on screen. (6 marks) |
| ix. | On game end you have to clear screen and display score on the screen. (2 marks) |
| х. | Use proper subroutines and stacks. No marks for code without subroutine and stack implementation. Maintain proper |
| | flow and declare data properly. (5 marks) |
| | |
| | Subroutines Required: |
| | Start, clearScreen, startDisplay(initialize the screen with score, time and characters), timerISR (updated time isr), |
| | scoreUpdate (updates score when correct character pressed), endScreen (call that when game ended clear screen and |
| | display score), kblSR (check valid characters, find bomb, removes valid character. You can also use nested subroutines |
| | like bombCheck, removeCharacter, checkBombLeft) |
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