



# **CS-2001**

## **DATA STRUCTURE**

**Dr. Hashim Yasin**

**National University of Computer  
and Emerging Sciences,  
Faisalabad, Pakistan.**

GRAPH

# Graph

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- A **graph** is a collection of **nodes** (or **vertices**) and **edges** (or **arcs**)
  - ▣ Each **node** contains an element
  - ▣ Each **edge** connects two nodes together (or possibly the same node to itself) and may contain an **edge attribute**

# Graph

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**Graph:** A data structure that consists of a set of nodes and a set of edges that relate the nodes to one another.

**Vertex:** A node in a graph is called vertex

**Edges(arc):** A pair of vertices representing a connection between two nodes in a graph

**Undirected Graph:** A graph in which the edges have no direction

**Directed Graph:** A graph in which each edge is directed from one vertex to another( or the same) vertex

# Graph

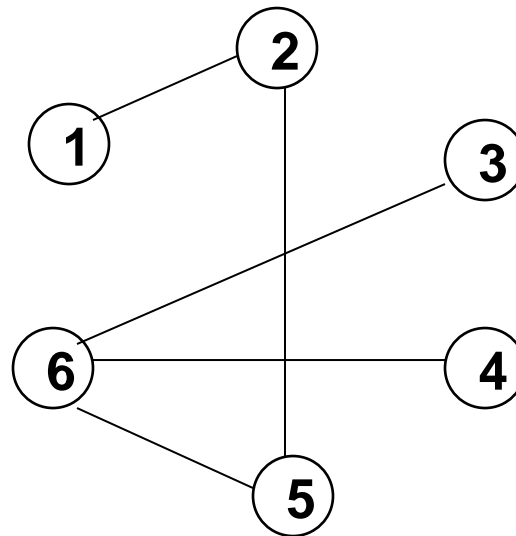
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- There are two kinds of graphs:
  - Undirected graphs
  - Directed graphs
- A **directed graph** is one in which the edges do not have direction
- An **undirected graph** is one in which the edges have a direction.

# Graph ... Examples

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## Undirected Graph



$$V = \{1, 2, 3, 4, 5, 6\}$$

$$E = \{(1, 2) (2, 5) (3, 6) (4, 6) (5, 6)\}$$

# Graph ... Examples

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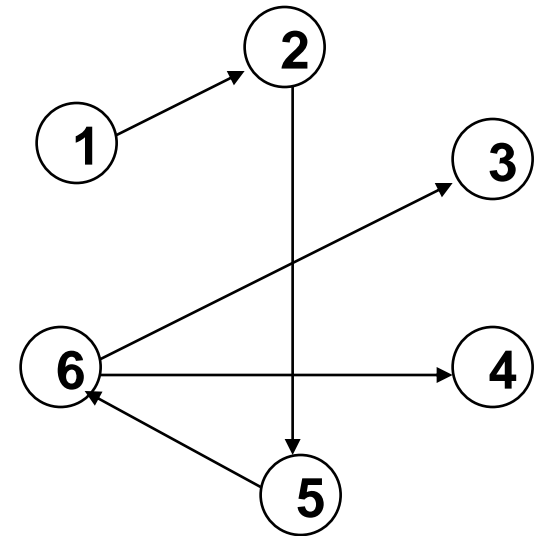
## Directed Graph

If the pair of vertices is ordered then the graph is a **directed graph**,

$G=(V,E)$ :

$V = \{1, 2, 3, 4, 5, 6\}$

$E = \{(1, 2) (2, 5) (5, 6) (6, 3) (6, 4)\}$



# Terminologies

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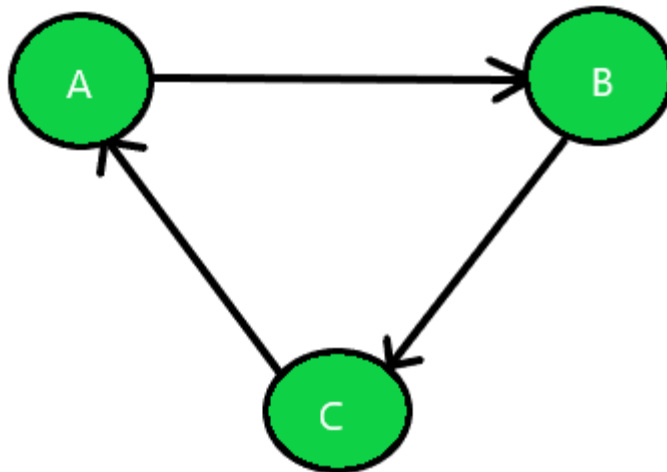
- A graph is said to be **connected** if there is a path from every vertex to every other vertex
- **Strongly Connected:** A graph is said to be **strongly connected** if *every pair of vertices( $u, v$ ) in the graph contains a path between each other.*
- **Unilaterally Connected:** A graph is said to be **unilaterally connected** if *it contains a directed path from  $u$  to  $v$  OR a directed path from  $v$  to  $u$  for every pair of vertices  $u, v$ .*
- **Weakly Connected:** A graph is said to be **weakly connected** if *there doesn't exist any path between any two pairs of vertices.*



# Terminologies

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## *Strongly Connected Graph*

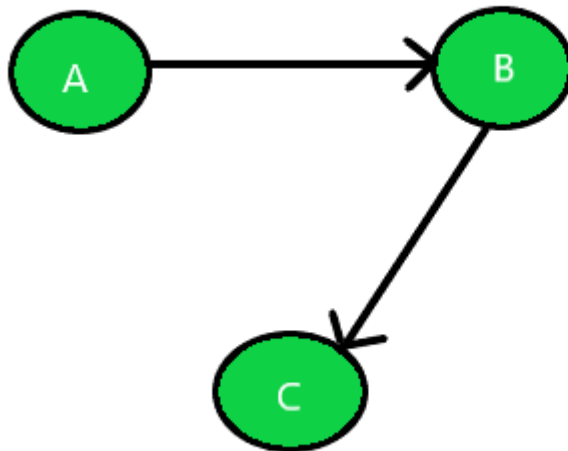


Path Matrix :    A B C  
A 1 1 1  
B 1 1 1  
C 1 1 1

# Terminologies

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## *Unilaterally Connected Graph*

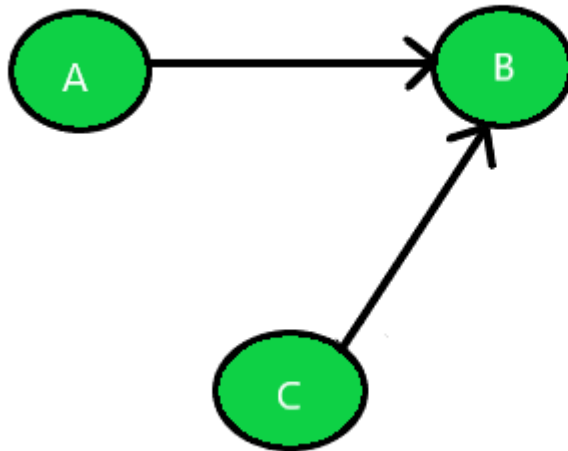


Path Matrix :    A B C  
A 0 1 1  
B 0 0 1  
C 0 0 0

# Terminologies

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## *Weakly Connected Graph*

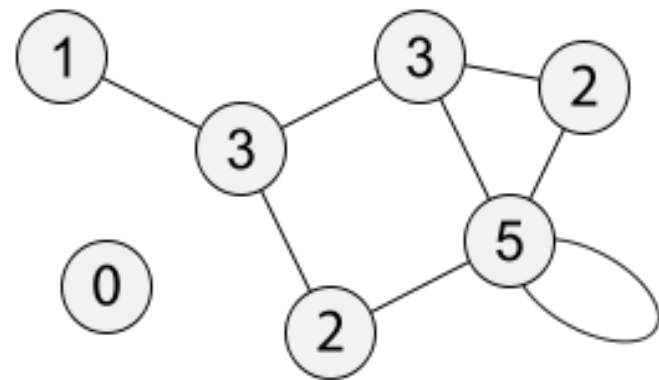


Path Matrix :    A B C  
A   0 1 0  
B   0 0 0  
C   0 1 0

# Adjacency

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- Vertices connected by an edge are **adjacent**
- **Degree** of a vertex  $v$ ,  $\deg(v)$ 
  - ▣ number of edges incident on  $v$
- The **degree (or valency)** of a vertex of a graph is the number of edges incident to the vertex, with loops counted twice



IMPLEMENTATION

# Graph Representations

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- Adjacency-matrix Representation
- Adjacency Lists Representation

# Adjacency Matrix

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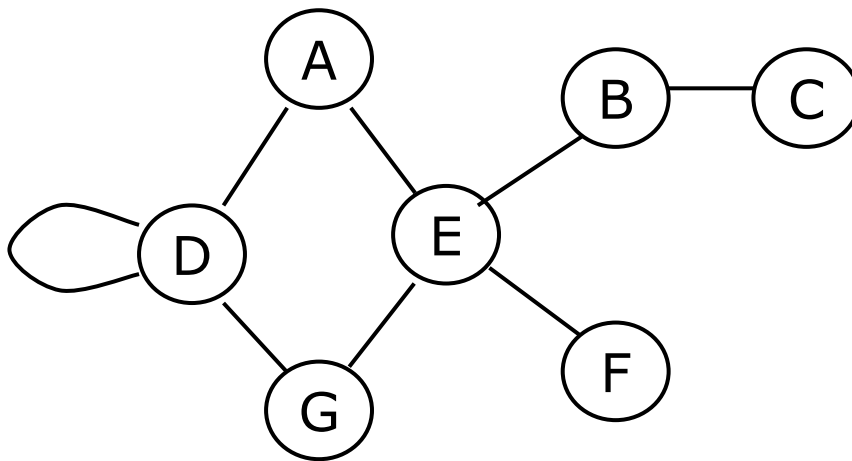
## Undirected Graphs:

- One simple way of representing a graph is the **adjacency matrix**
- A 2-D array has a mark at  $[i][j]$  if there is an edge from node  $i$  to node  $j$
- The adjacency matrix is **symmetric** about the main diagonal.
- This representation is only suitable for small graphs!

# Adjacency Matrix

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## Undirected Graphs:



	A	B	C	D	E	F	G
A				●	●		
B			●		●		
C		●					
D	●			●			●
E	●	●				●	●
F					●		
G				●	●		



# Adjacency Matrix

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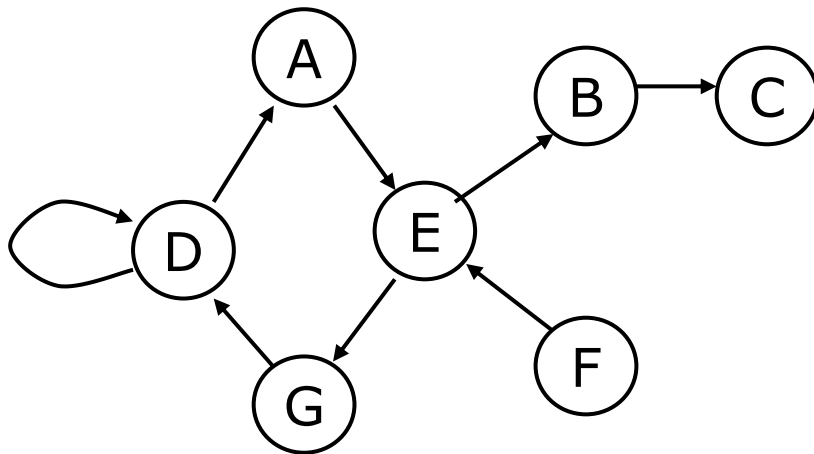
## Directed Graphs:

- An **adjacency matrix** can equally well be used for directed graphs.
- A 2-D array has a mark at  $[i][j]$  if there is an edge from node  $i$  to node  $j$ .

# Adjacency Matrix

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## Directed Graphs:

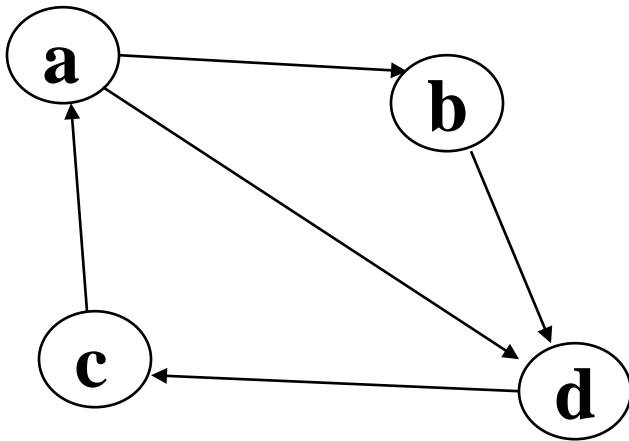


	A	B	C	D	E	F	G
A					●		
B			●				
C							
D	●			●			
E		●					●
F					●		
G				●			

# Adjacency Matrix

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## Directed Graphs:

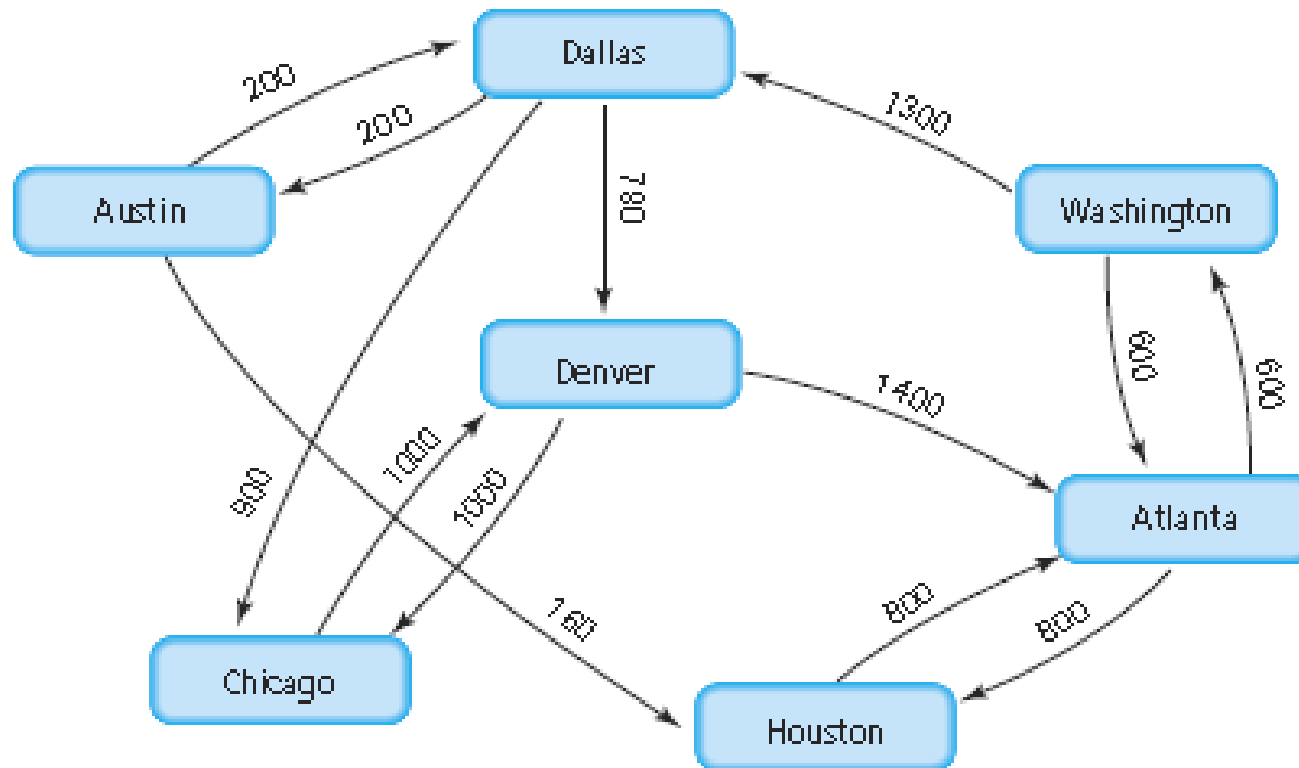


	<b>a</b>	<b>b</b>	<b>c</b>	<b>d</b>
<b>a</b>	0	1	0	1
<b>b</b>	0	0	0	1
<b>c</b>	1	0	0	0
<b>d</b>	0	0	1	0

# Adjacency Matrix

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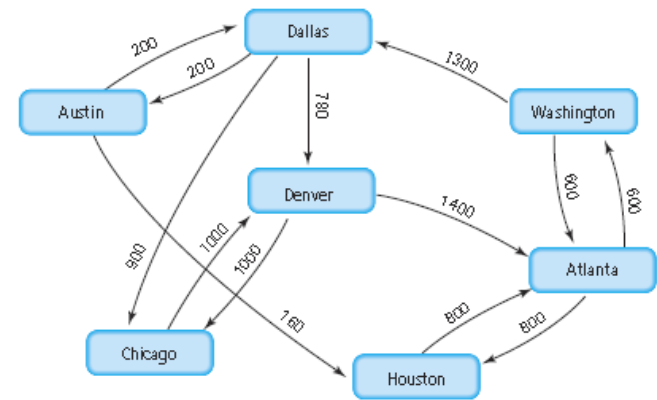
## Weighted Directed Graphs:



# Adjacency Matrix

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## Weighted Directed Graphs:



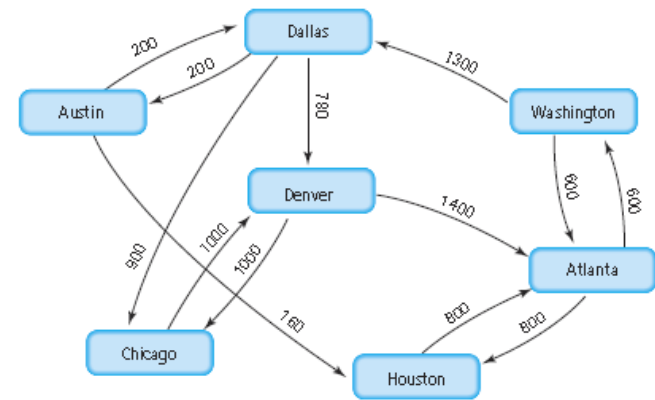
Austin	}	160 miles
Houston		
Atlanta	}	800 miles
Washington		

Austin	}	200 miles
Dallas		
Denver	}	780 miles
Atlanta		
Washington	}	1400 miles
	}	600 miles

# Adjacency Matrix

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## Weighted Directed Graphs:



.vertices

.edges

[0]	"Atlanta"	"
[1]	"Austin"	"
[2]	"Chicago"	"
[3]	"Dallas"	"
[4]	"Denver"	"
[5]	"Houston"	"
[6]	"Washington"	"

[0]	0	0	0	0	0	800	600
[1]	0	0	0	200	0	160	0
[2]	0	0	0	0	1000	0	0
[3]	0	200	900	0	780	0	0
[4]	1400	0	1000	0	0	0	0
[5]	800	0	0	0	0	0	0
[6]	600	0	0	1300	0	0	0
	[0]	[1]	[2]	[3]	[4]	[5]	[6]

# Adjacency Lists

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## Adjacency List

- ▣ Array of lists
- ▣ Each vertex *has an array entry*
- ▣ A vertex **w** is inserted in the list for vertex **v** if there is an outgoing edge from **v** to **w**

# Adjacency Lists

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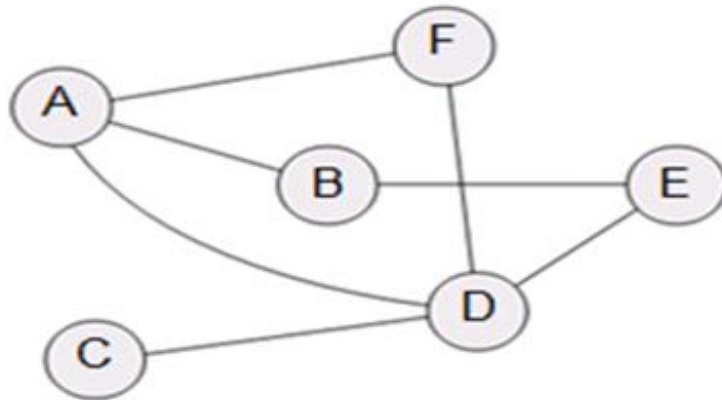
- A graph of  $n$  nodes is represented by a one-dimensional array  $L$  of linked lists, where,
  - ▣  $L[i]$  is the linked list containing all the nodes adjacent to node  $i$
  - ▣ The nodes in the list  $L[i]$  are in **NO particular order**,
  - ▣ An adjacency list for a weighted graph should contain **two elements** in the list nodes – one element for the vertex and the second element for the weight of that edge



# Adjacency Lists

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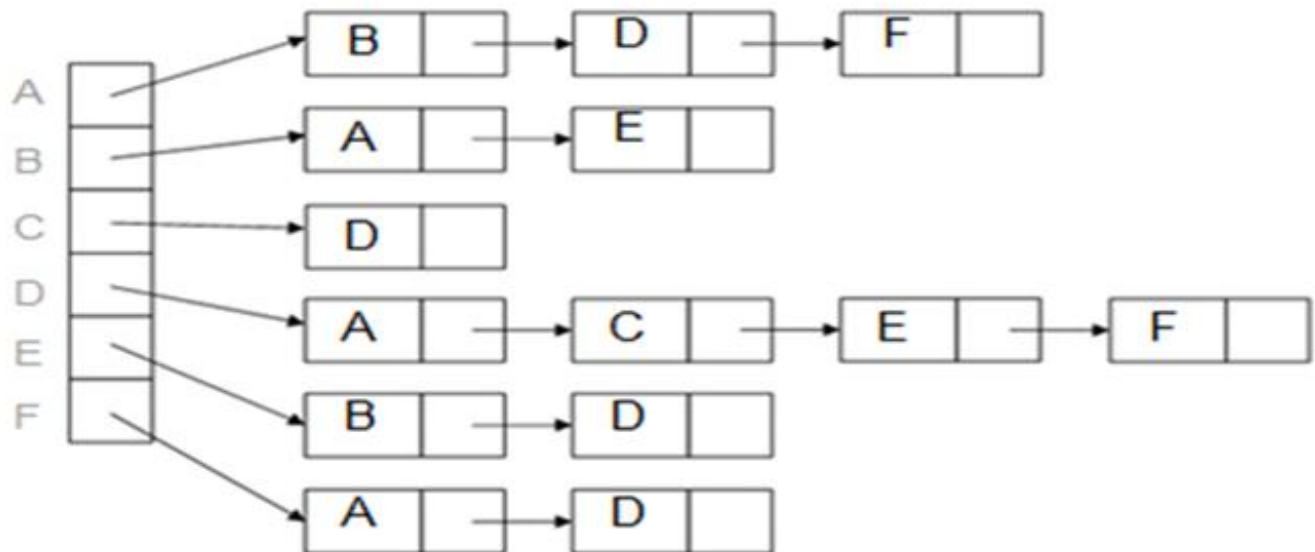
An (undirected) graph  $G = (V, E)$



*adjacency matrix for  $G$*

	A	B	C	D	E	F
A	0	1	0	1	0	1
B	1	0	0	0	1	0
C	0	0	0	1	0	0
D	1	0	1	0	1	1
E	0	1	0	1	0	0
F	1	0	0	1	0	0

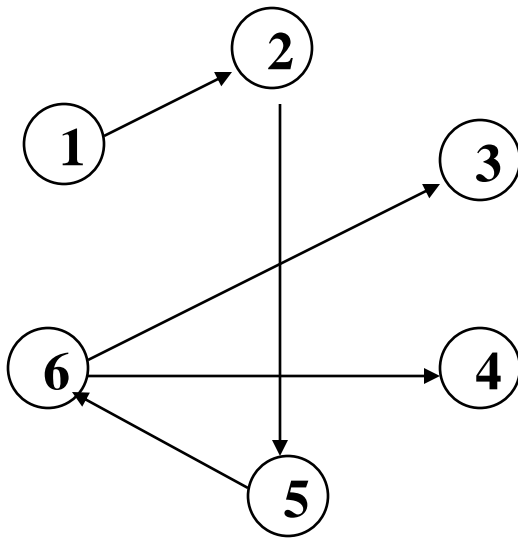
*adjacency list for  $G$*



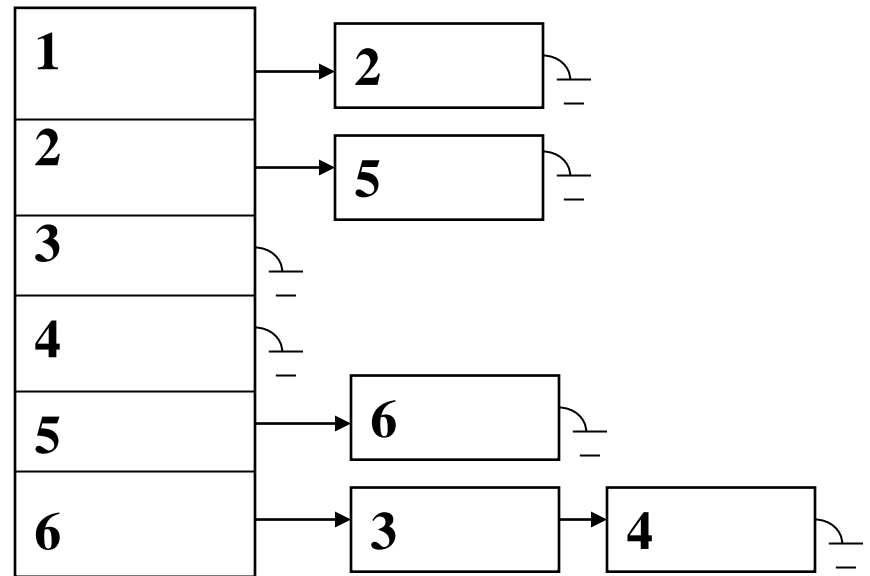
# Adjacency Lists

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## Directed Graphs:



Graph

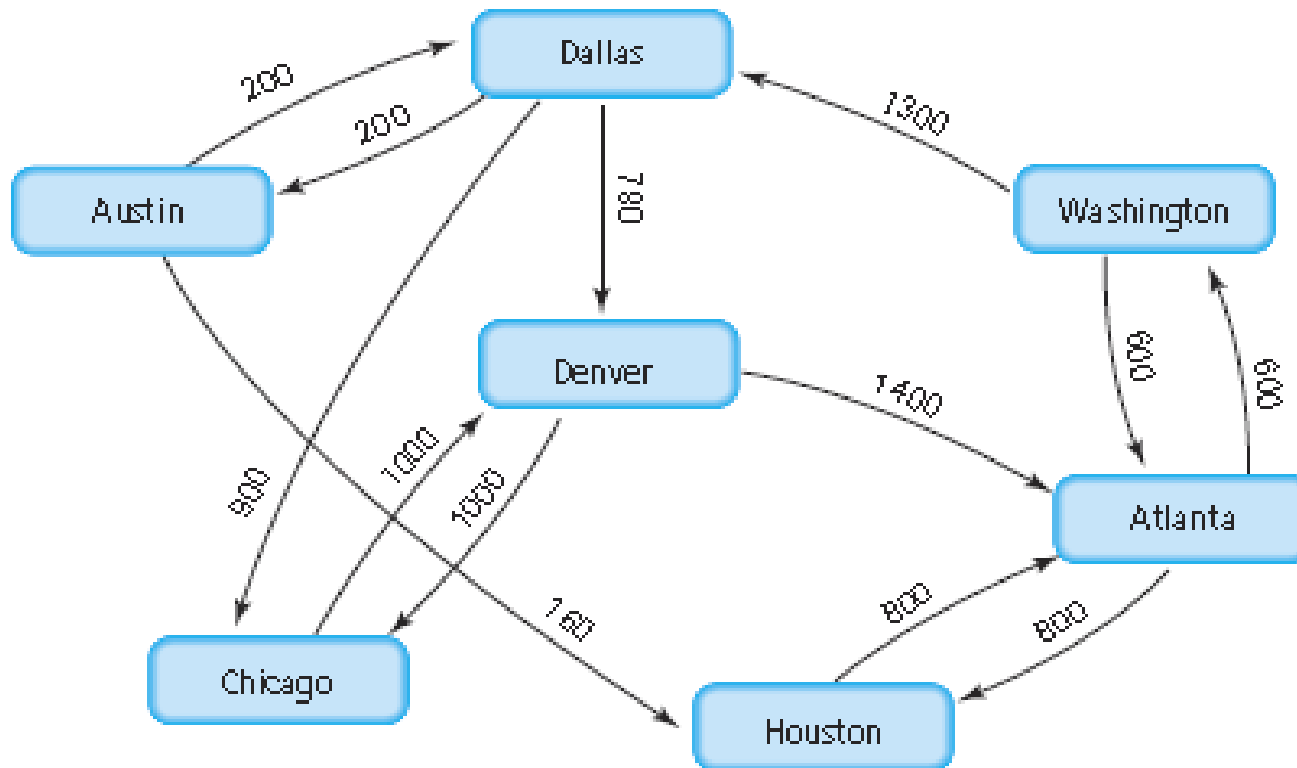


Adjacency List

# Adjacency Lists

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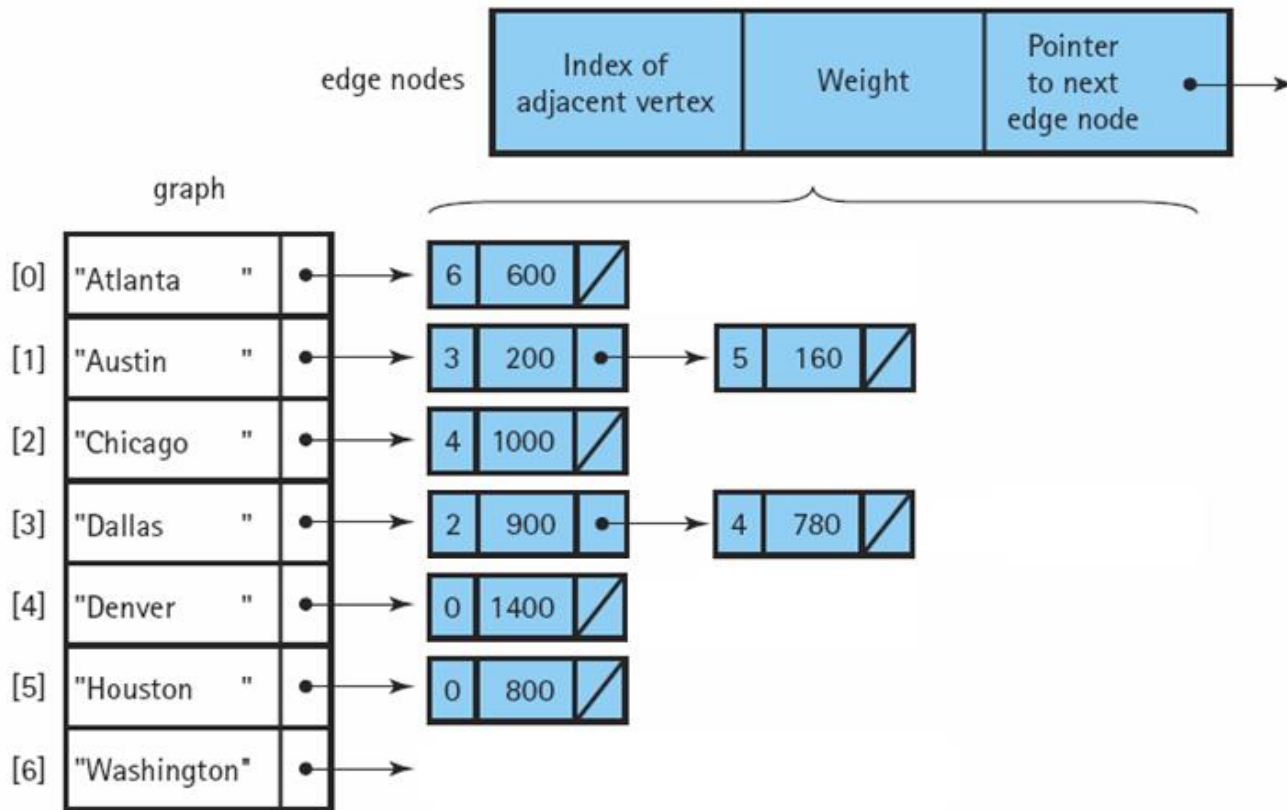
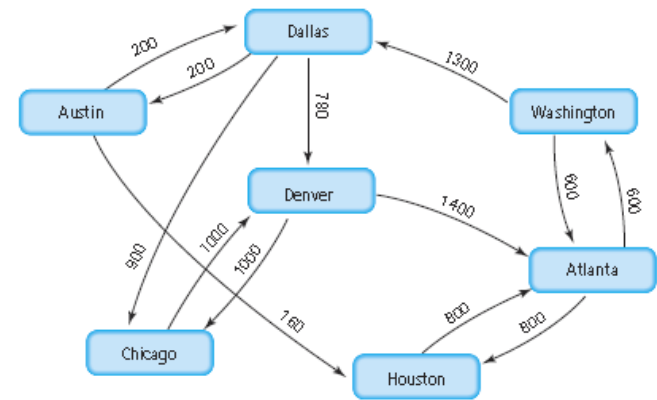
## Weighted Directed Graphs:

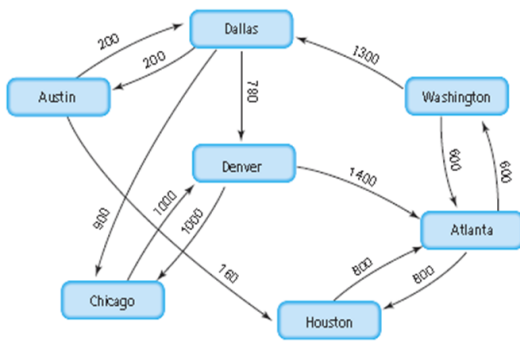


# Adjacency Lists

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## Weighted Directed Graphs:



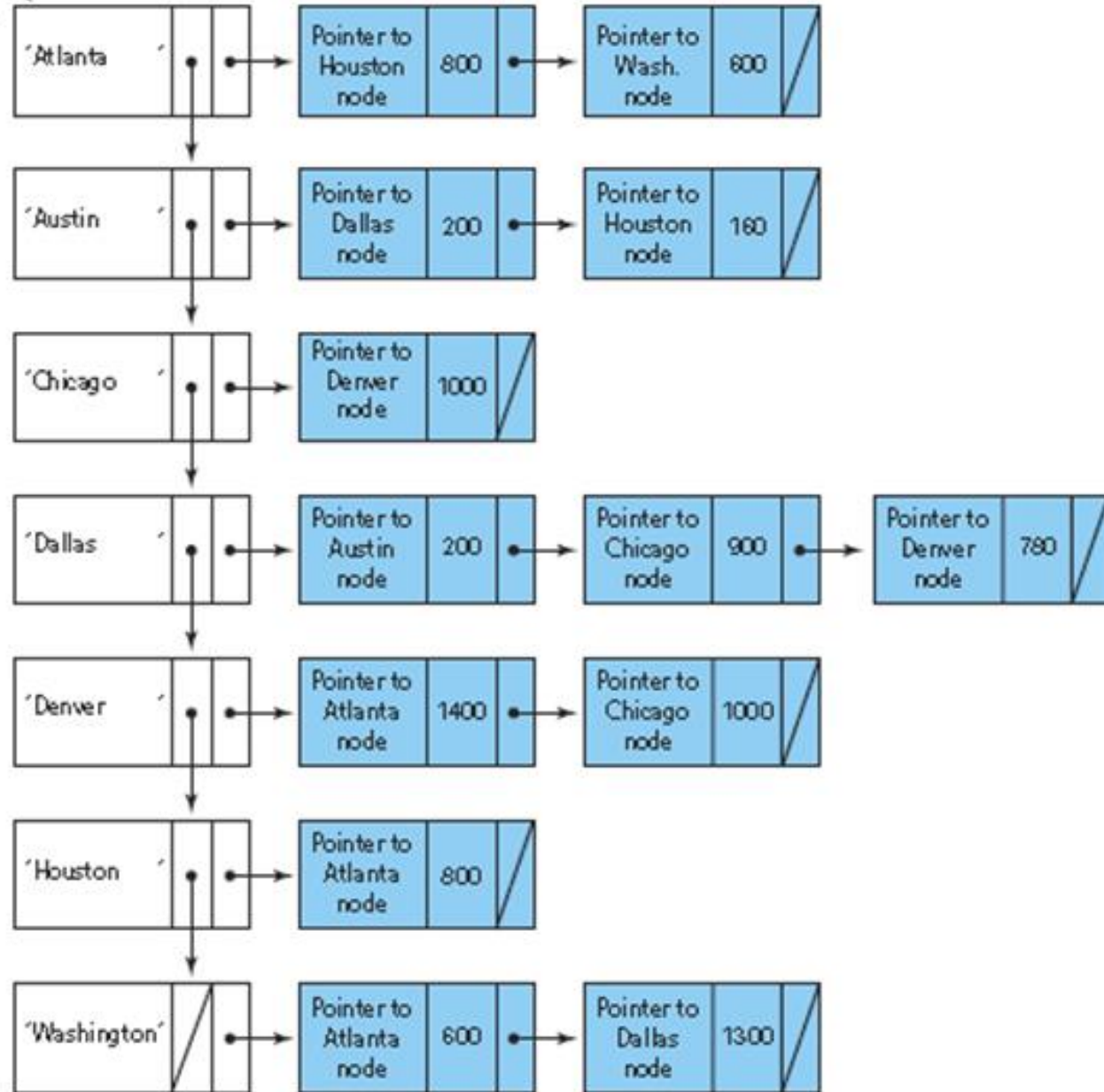
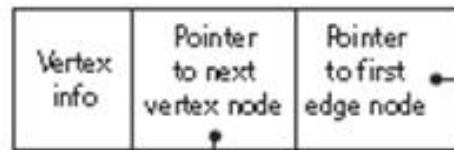


graph

edge nodes



vertices  
(header nodes)



# Adjacency Matrix

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## □ Pros:

- ▣ Simple to implement
- ▣ Easy and fast to tell if a pair  $(i,j)$  is an edge:
  - simply check if  $A[i][j]$  is 1 or 0

## □ Cons:

- ▣ No matter how few edges the graph has, the matrix takes  $O(n^2)$  in memory

# Adjacency Lists

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## Pros:

- ▣ **Save space (memory):** the representation takes as many memory words as there are nodes and edges.
- ▣ For both directed and undirected graphs, the adjacency-list representation has the desirable property that the amount of memory it requires, is  $O(V+E)$ .

## Cons:

- ▣ *It can take up to  $O(n)$  time to determine if a pair of nodes  $(i,j)$  has an edge:*
  - One would have to search the linked list  $L[i]$ , which takes time proportional to the length of  $L[i]$ .

# Graph ... Implementation

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- Operations: The basic functions of graph
  - ▣ **MakeEmpty** →
    - *Initializes the graph to an empty state.*
  - ▣ **Boolean IsEmpty** →
    - *Tests whether the graph is empty*
  - ▣ **Boolean IsFull** →
    - *Tests whether the graph is full.*
  - ▣ **AddVertex(VertexType vertex)** →
    - *Add vertex to graph*



# Graph ... Implementation

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- Operations: The basic functions of graph
  - ▣ **AddEdge** (VertexType fromVertex, VertexType toVertex, EdgeValueType weight) → *Adds an edge with the specified weight from **fromVertex** to **toVertex**.*
  - ▣ **EdgeValueType Weights** (VertexType fromVertex, VertexType toVertex) → *Determines the weight of the edge from **fromVertex** to **toVertex**.*
  - ▣ **GetToVertices** (VertexType vertex, QueType& vertexQ) → *Returns a queue of the vertices that are adjacent from vertex.*

# Reading Materials

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□ Nell Dale: Chapter # 9

