



CS-218

DATA STRUCTURE

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BINARY TREE

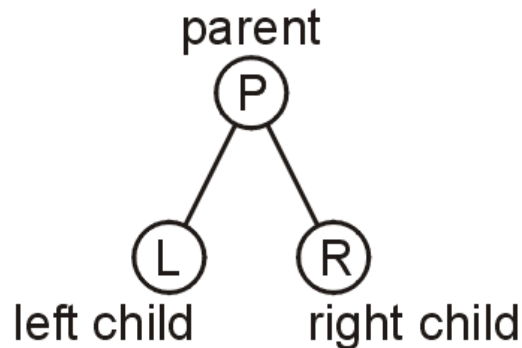


Binary Trees

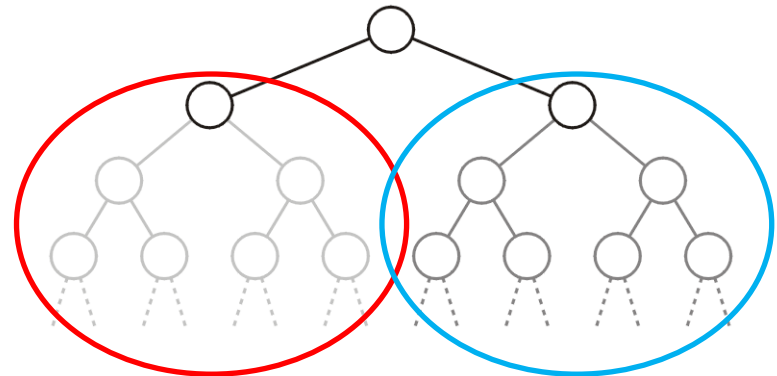
3

Binary Tree

- In a binary tree, **each node has at most two children**
 - ▣ Allows to label the children as left and right



- Likewise, the two sub-trees are referred as
 - ▣ **Left sub-tree**
 - ▣ **Right sub-tree**



Binary Tree ... Implementation

4

- Binary Tree ... Implementation
 - ❖ Array based implementation
 - ❖ Linked List based implementation

Binary Tree ... Implementation

5

- Objects: any type of objects can be stored in a tree data structure
- Accessor methods
 - ▣ `root()` – return the root of the tree
 - ▣ `parent(p)` – return the parent of a node
 - ▣ `children(p)` – returns the children of a node

Binary Tree ... Implementation

6

□ query methods

- `size()` – returns the number of nodes in the tree
- `isEmpty()` - returns true if the tree is empty
- `elements()` – returns all elements
- `isRoot(p)`, `isInternal(p)`, `isExternal(p)`

□ other methods

- Tree traversal, Node addition/deletion, create/destroy

Binary Tree ... Implementation

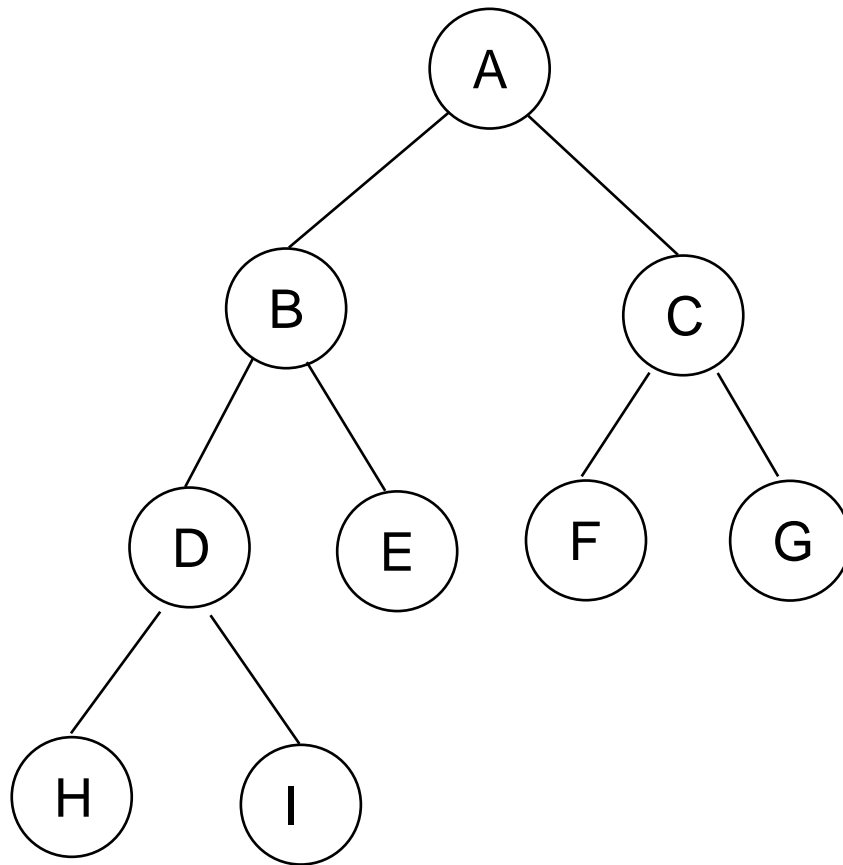
7

Array-based Implementation:

- Value in root node stored first, followed by left child, then right child
- Each successive level in the tree stored left to right; unused nodes in tree represented by a bit pattern to indicate nothing stored there
- Children of any given node n is stored in cells $2n$ and $2n + 1$ (If array index starts at 1)
- *Storage allocated as for full tree, even if many nodes empty*
- For a tree of depth h we need array of $2^{h+1}-1$ cells

Binary Tree ... Implementation

Array-based Implementation:

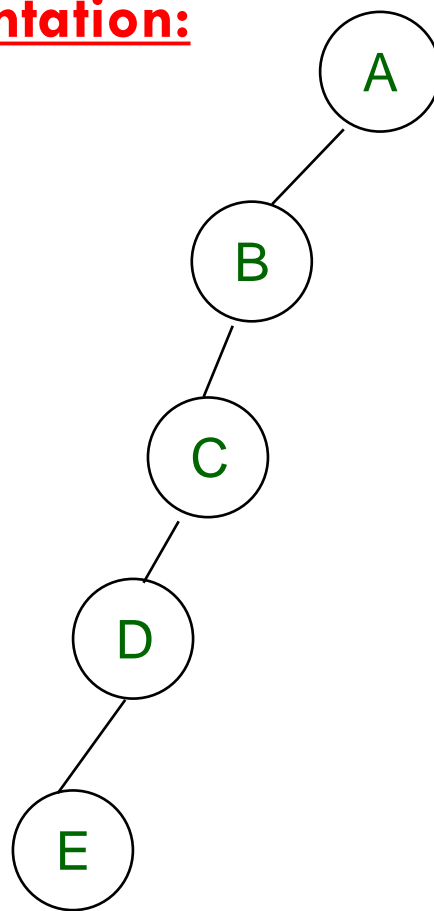
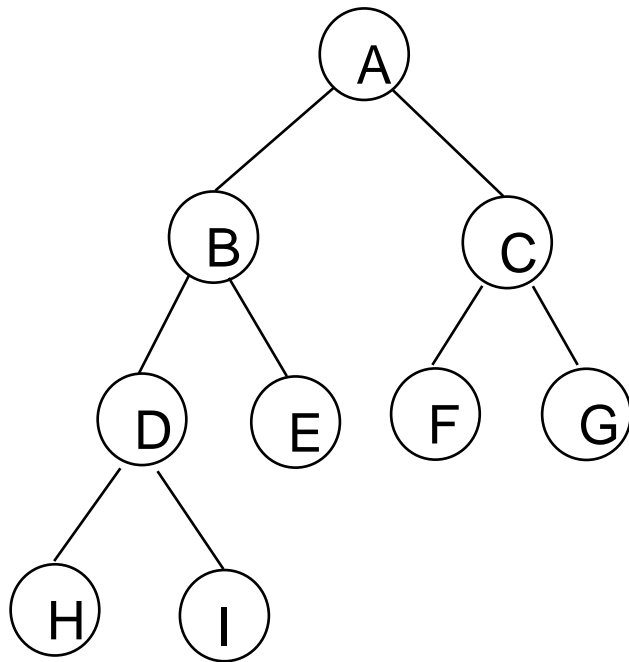


[1]	A
[2]	B
[3]	C
[4]	D
[5]	E
[6]	F
[7]	G
[8]	H
[9]	I

Binary Tree ... Implementation

9

Array-based Implementation:



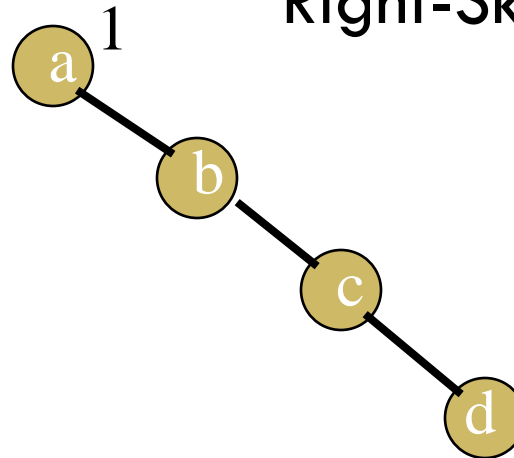
[1]	A
[2]	B
[3]	--
[4]	C
[5]	--
[6]	--
[7]	--
[8]	D
[9]	--
.	.
[16]	E

Binary Tree ... Implementation

10

Array-based Implementation:

Right-Skewed Binary Tree



tree[]

a	-	b	-	-	-	c	-	-	-	-	-	-	-	d
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

Binary Tree ... Implementation

11

Linked List-based Implementation:

- Implementation of Binary tree can be done using dynamic creation of node having data and two pointers.

```
Struct Node {  
    int data;  
    Node* left;  
    Node * right;  
}
```

Binary Tree ... Implementation

12

- A binary tree is composed of **zero or more nodes**
- Each node contains:
 - ▣ A **value** (some sort of data item)
 - ▣ A **reference or pointer** to a left child (may be **null**),
and
 - ▣ A **reference or pointer** to a right child (may be **null**)
- A binary tree may be *empty* (contain no nodes)

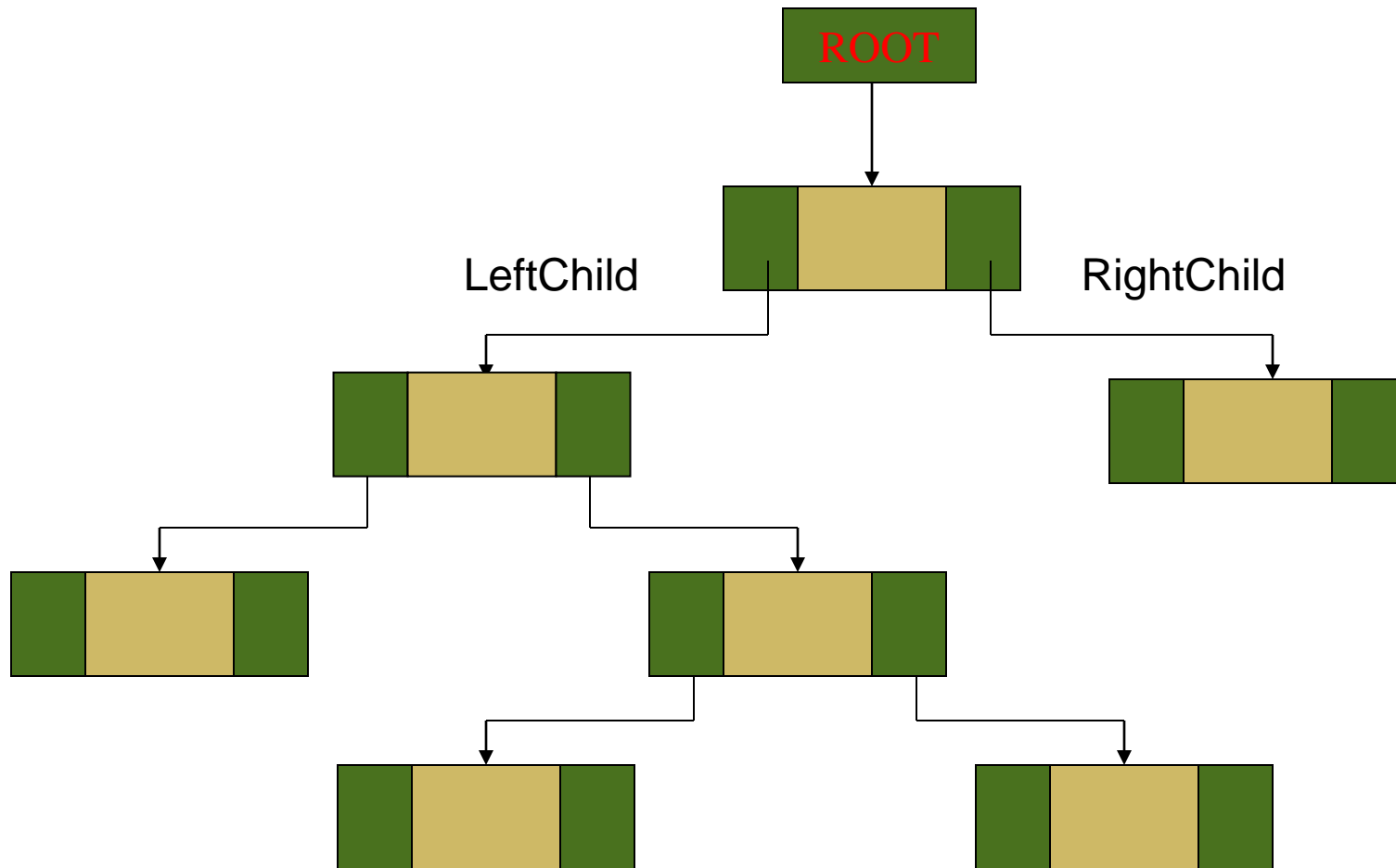
Binary Tree ... Implementation

13

- If not empty, a binary tree has a **root node**
 - ▣ Every node in the binary tree is reachable from the root node by a **unique** path
- A node with neither a left child nor a right child is called a **leaf**
 - ▣ In some binary trees, only the leaves contain a value

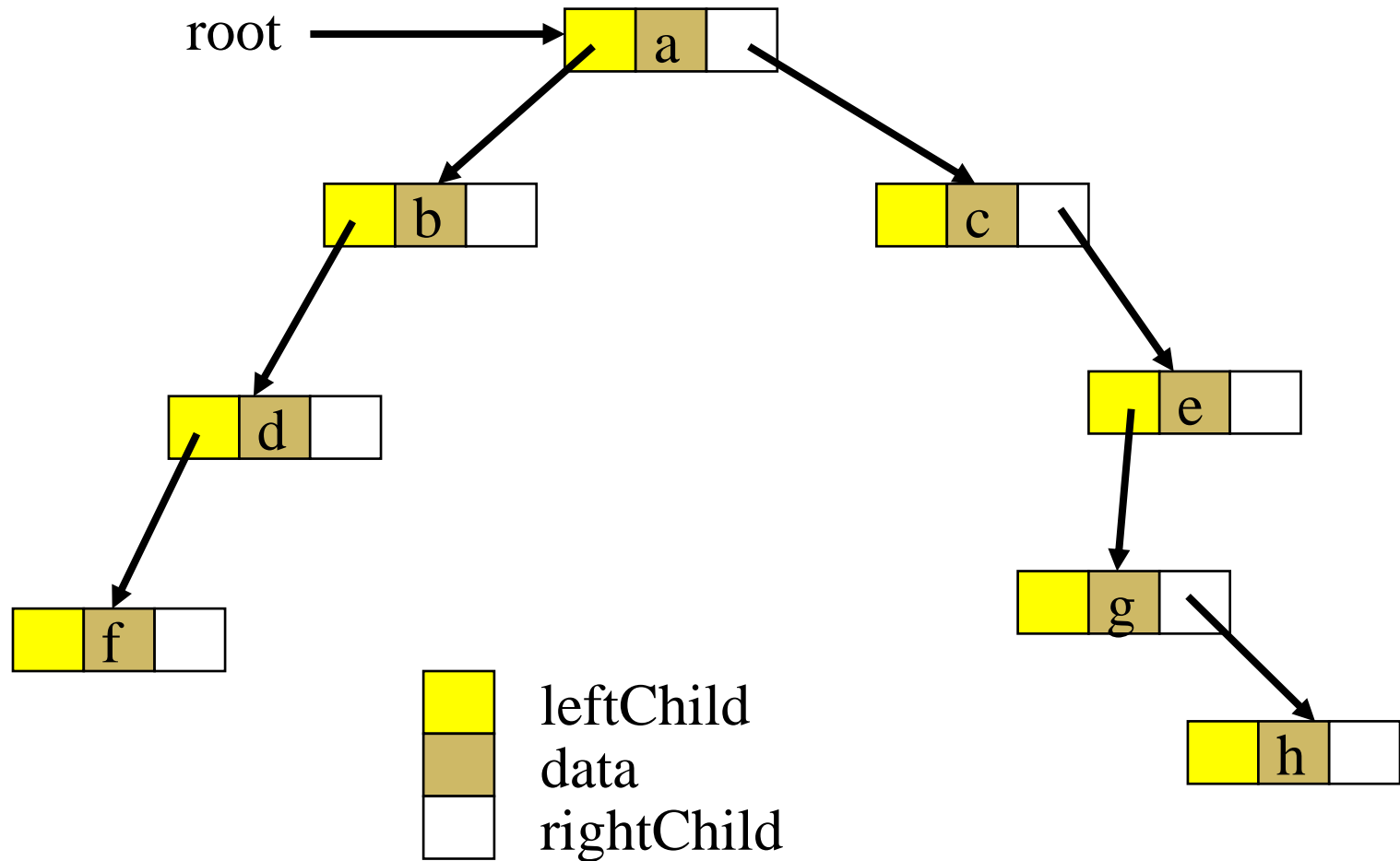
Binary Tree ... Implementation

14



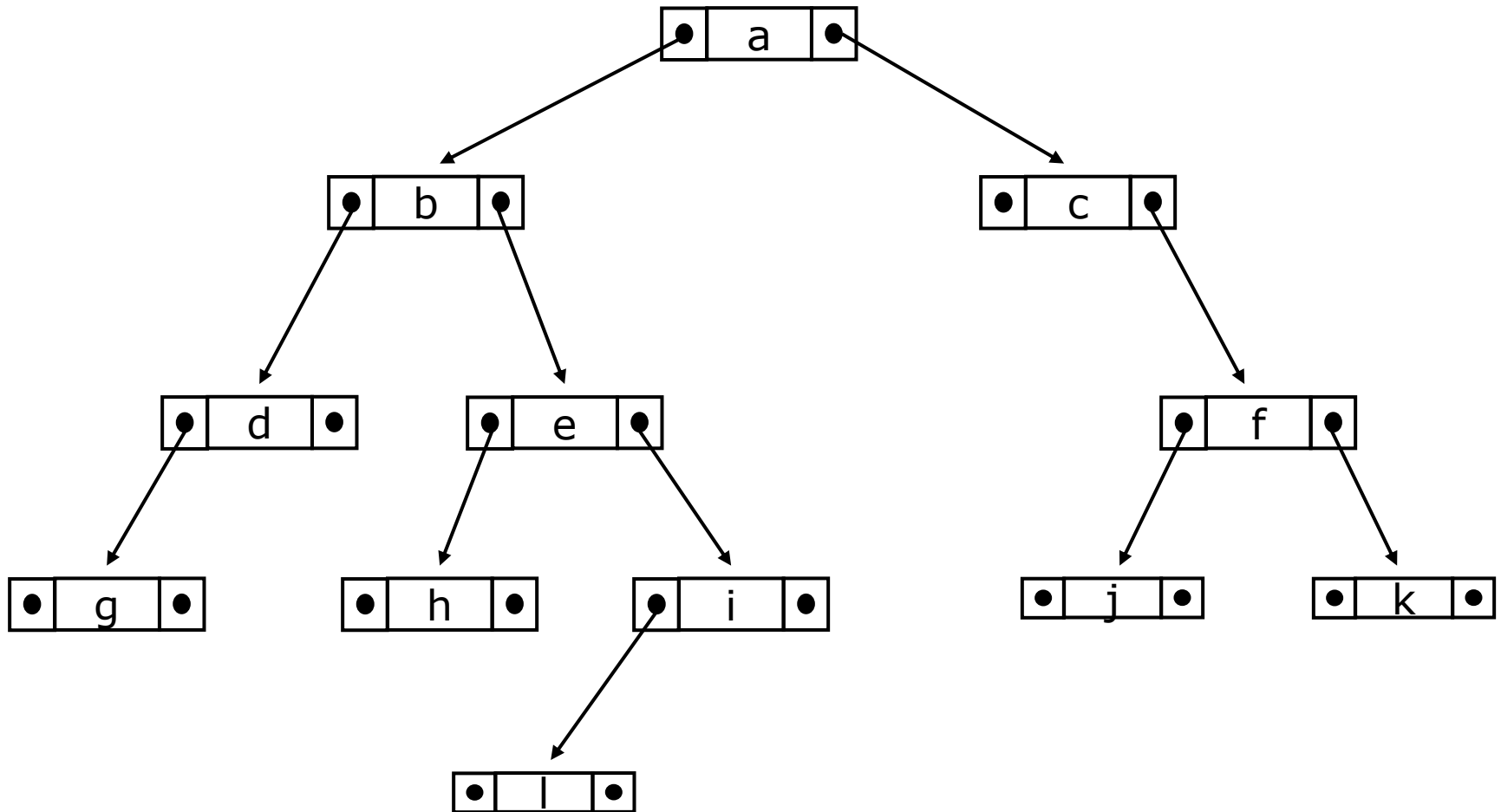
Binary Tree ... Implementation

15



Binary Tree ... Implementation

16



Binary Tree ... Implementation

17

Implementation: with Linked List

- **Each node** in the tree consists of:
 - ▣ The **data**, or value contained in the element
 - ▣ A **left child pointer** (pointer to first child)
 - ▣ A **right child pointer** (pointer to second child)
- A **root pointer** points to the root node
 - ▣ Follow pointers to find every other element in the tree
- *Add and remove nodes by manipulating pointers*
- **Leaf nodes** have child pointers set to null

Binary Tree ... Implementation

18

- **Left(node):** Gives index/pointer of left child
- **Right(node):** Gives index/pointer of right child
- **Parent(node):** Returns index/pointer of parent
- **Brother(node):** Returns index/pointer of brother
- **Root:** Gives index/pointer of root node
- **Info(Node):** Data/Info stored at node
- **IsLeft(node):** Is node left child? Yes/No
- **IsRight(node):** Is node right child? Yes/No

BINARY SEARCH TREE

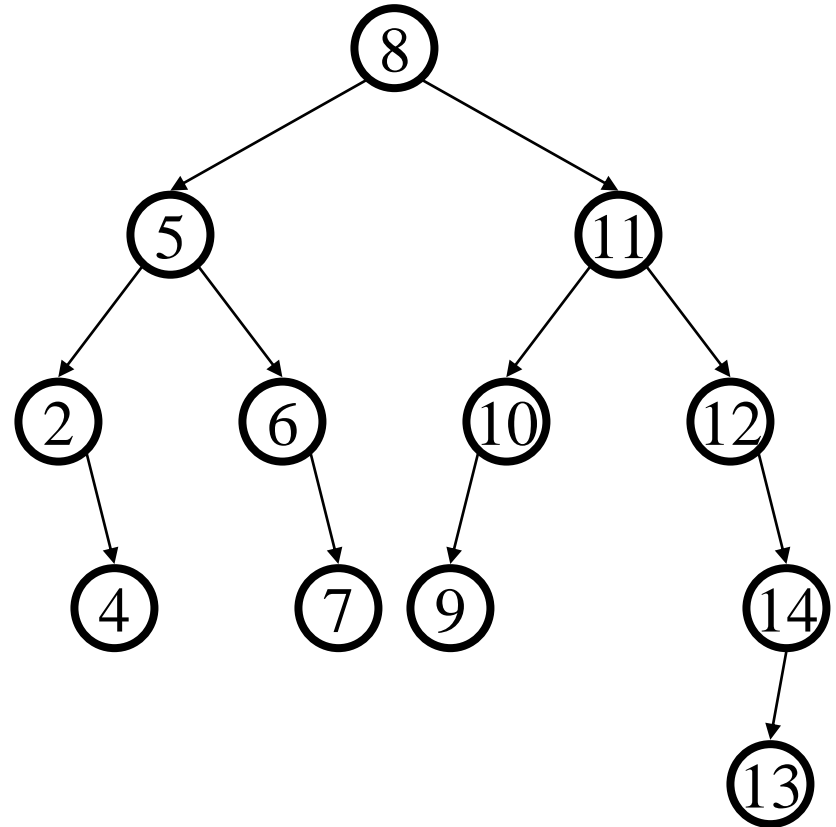
Binary Search Trees (BSTs)

20

□ Binary Search Tree

Property:

The value stored at a node is **greater than the value stored at its left child** and **less than the value stored at its right child**



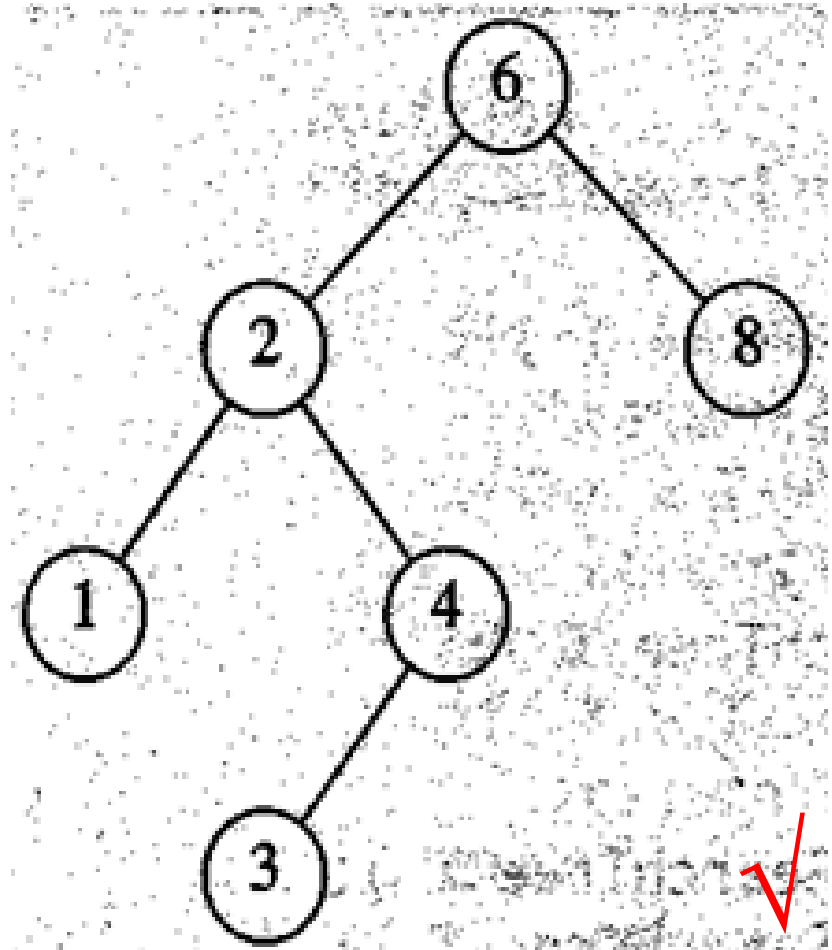
Binary Search Trees (BSTs)

21

□ Binary Search Tree

Property:

The value stored at a node is *greater than* the value stored at its *left child* and *less than* the value stored at its *right child*



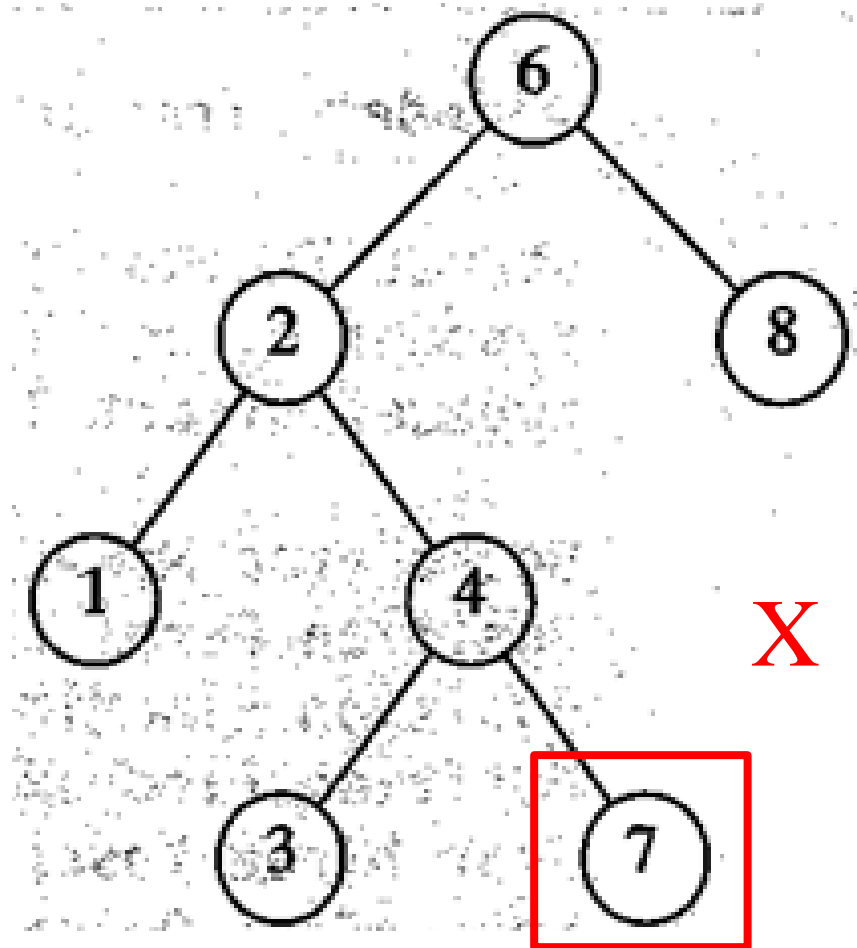
Binary Search Trees (BSTs)

22

□ Binary Search Tree

Property:

The value stored at a node is *greater than* the value stored at its *left child* and *less than* the value stored at its *right child*



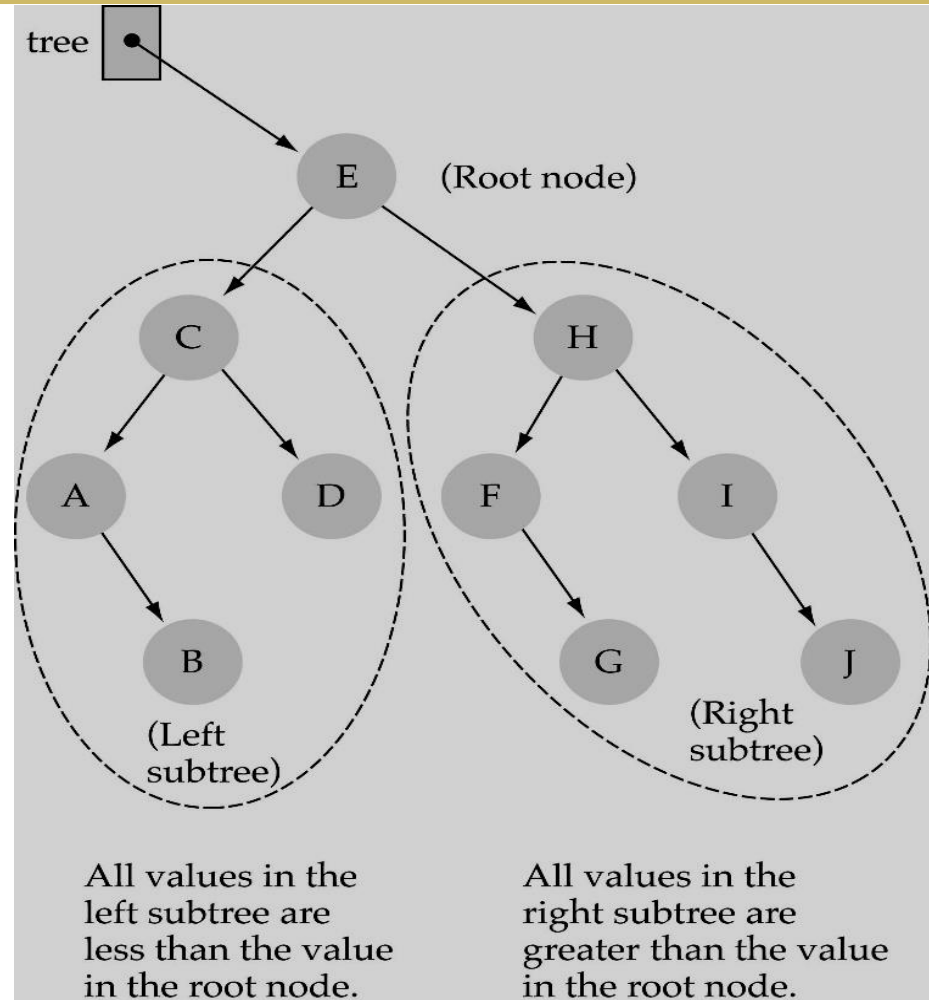
Binary Search Trees (BSTs)

23

□ Binary Search Tree

Property:

The value stored at a node is *greater than* the value stored at its *left child* and *less than* the value stored at its *right child*



Binary Search Trees (BSTs)

24

□ Binary Search Tree

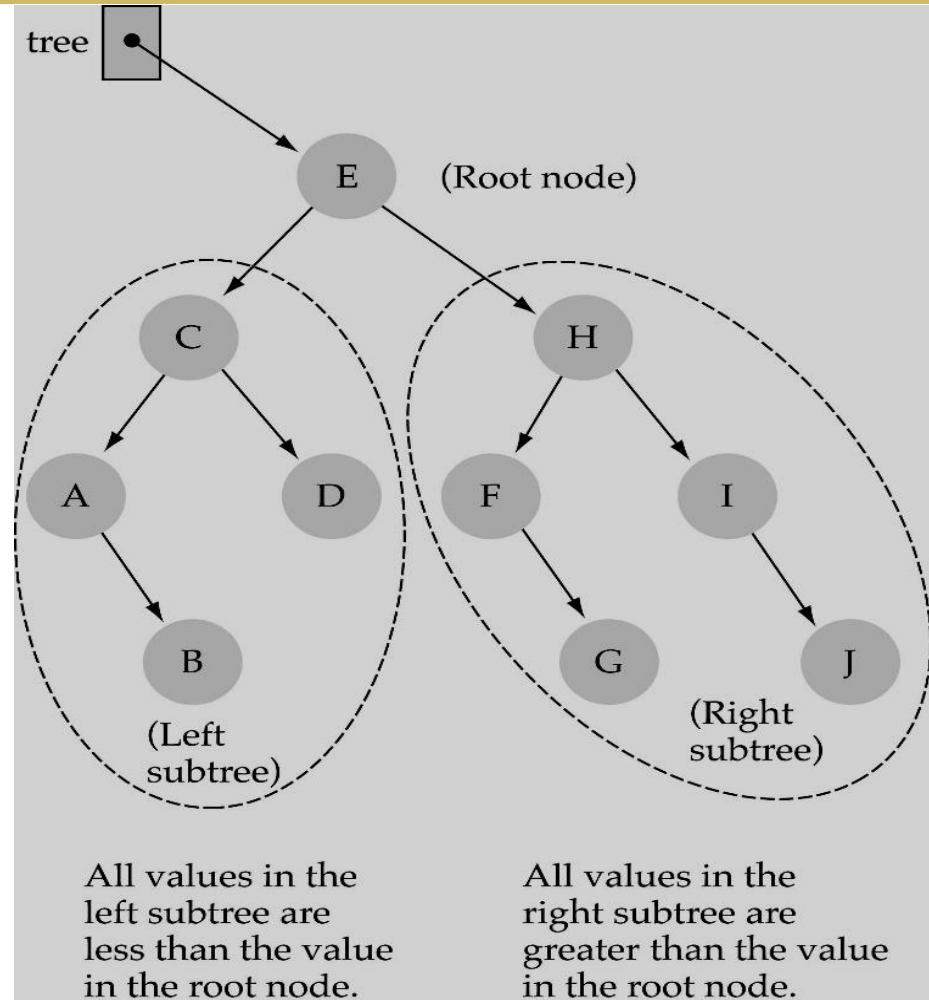
Property:

Where is the smallest element?

Ans: leftmost element

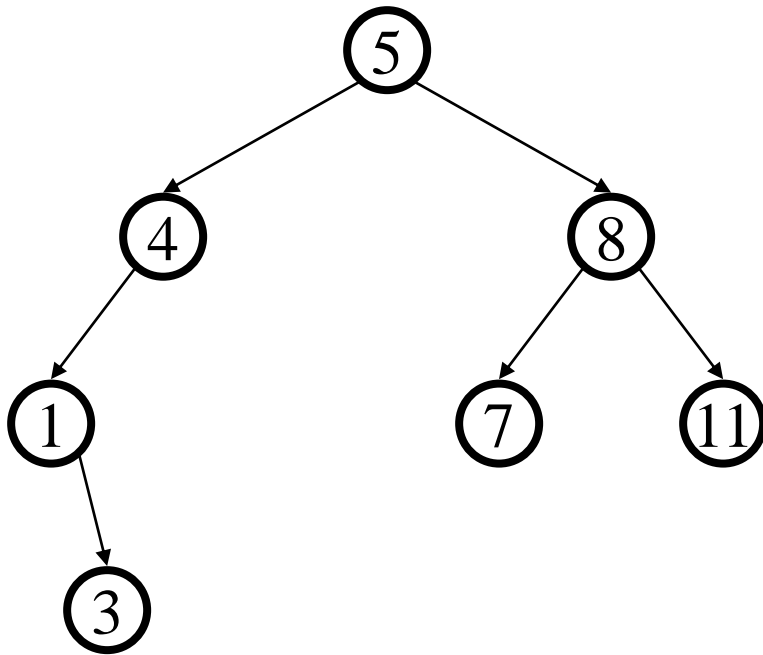
Where is the largest element?

Ans: rightmost element

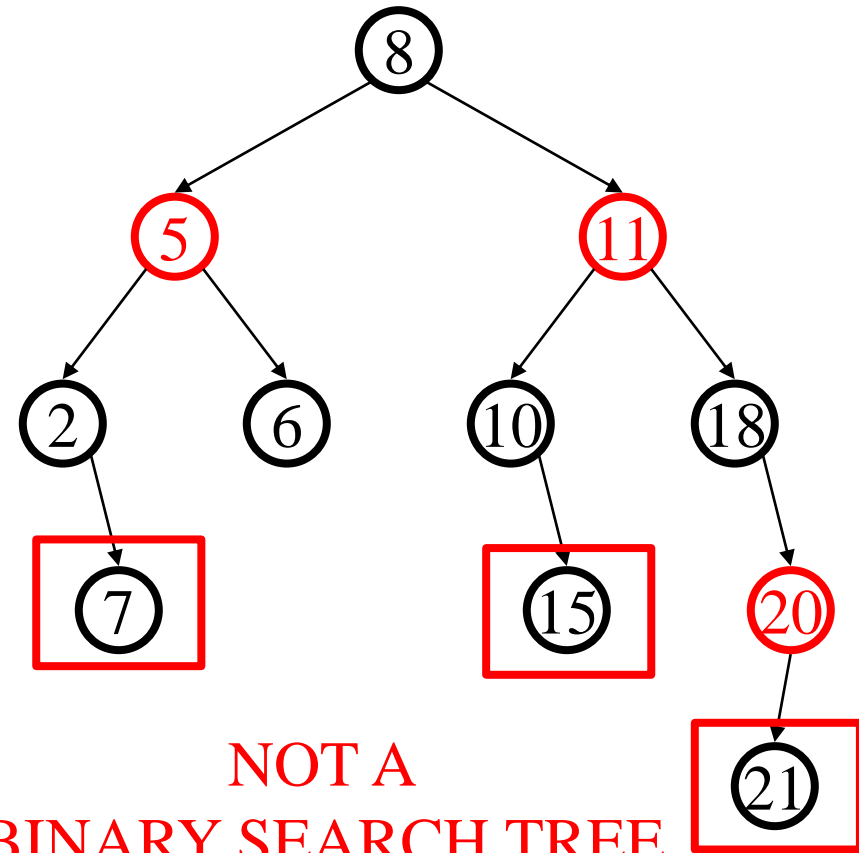


Binary Search Trees (BSTs)

25



BINARY SEARCH TREE



Binary Search Trees (BSTs)

26

Binary tree property

- ▣ each node has ≤ 2 children
- ▣ result:
 - storage is small
 - operations are simple
 - average depth is small

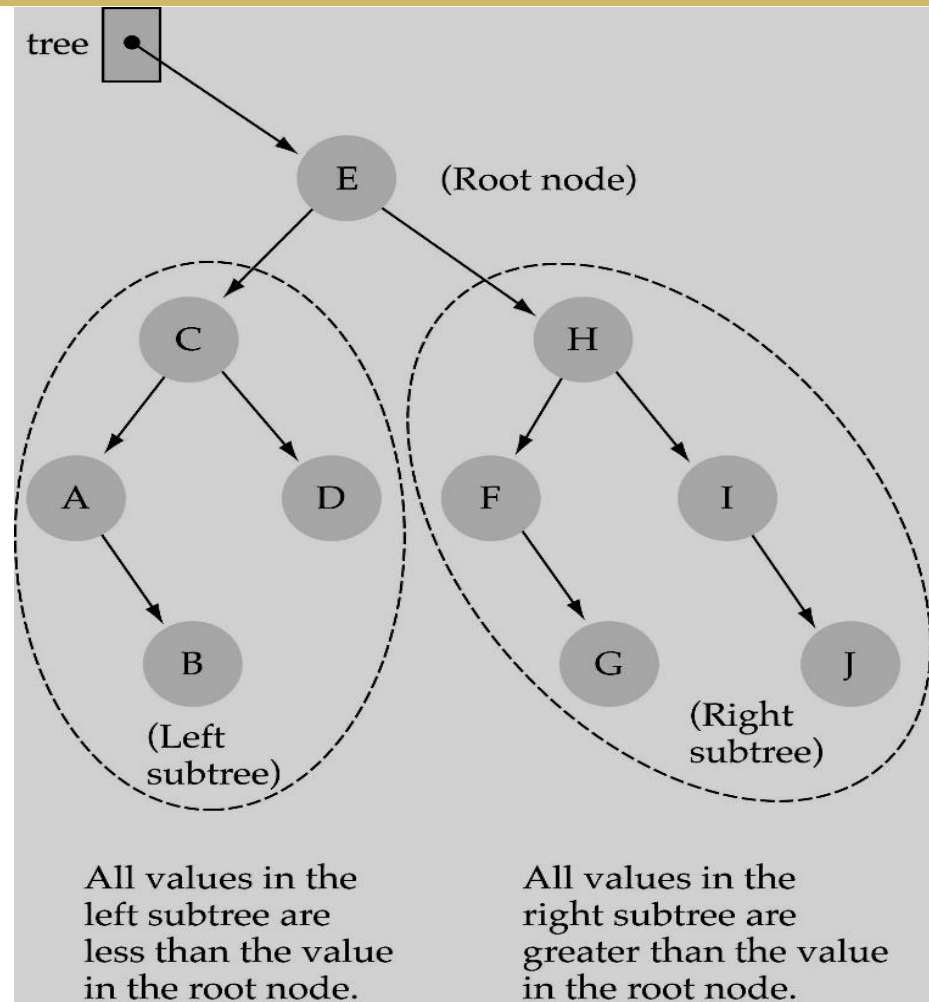
Search tree property

- ▣ all keys in left subtree smaller than root's key
- ▣ all keys in right subtree larger than root's key
- ▣ result:
 - easy to find any given key
 - Insert/delete by changing links

Searching in BST

27

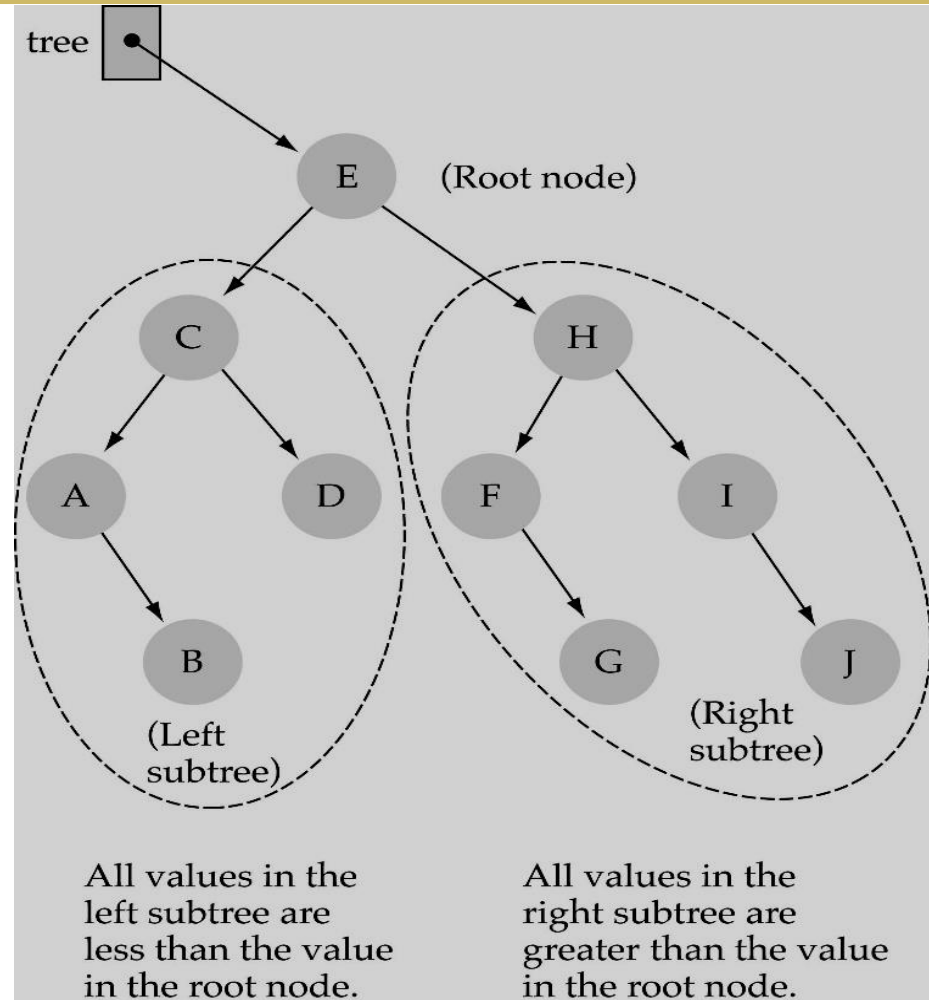
1. Start at the root
2. Compare the value of the item you are searching for with the value stored at the root
3. If the values are equal, then *item found*; otherwise, if it is a leaf node, then *not found*



Searching in BST

28

4. If it is **less than the value** stored at the root, then search the left subtree
5. If it is **greater than the value** stored at the root, then search the right subtree
6. Repeat steps 2-6 for the root of the subtree chosen in the previous step 4 or 5

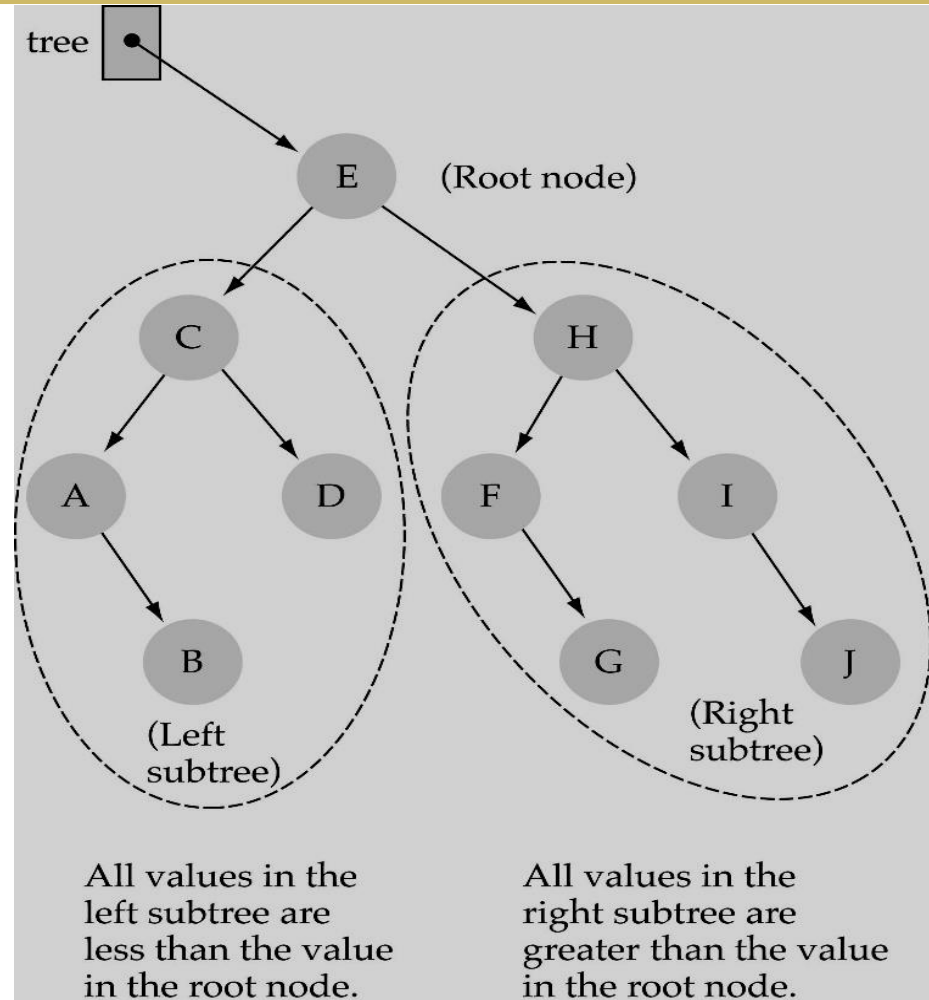


Searching in BST

29

Is this better than
searching a linked list?

Yes !! ---> $O(\log N)$



Why BST

30

Array

- Searching in the Array $O(n)$
- Insertion $O(1)$
- Remove $O(n)$

Linked List

- Searching in the Linked List $O(n)$
- Insertion $O(1)$
- Remove $O(n)$

BST

- Searching in the BST $O(\log n)$
- Insertion $O(\log n)$
- Remove $O(\log n)$

BST Operations

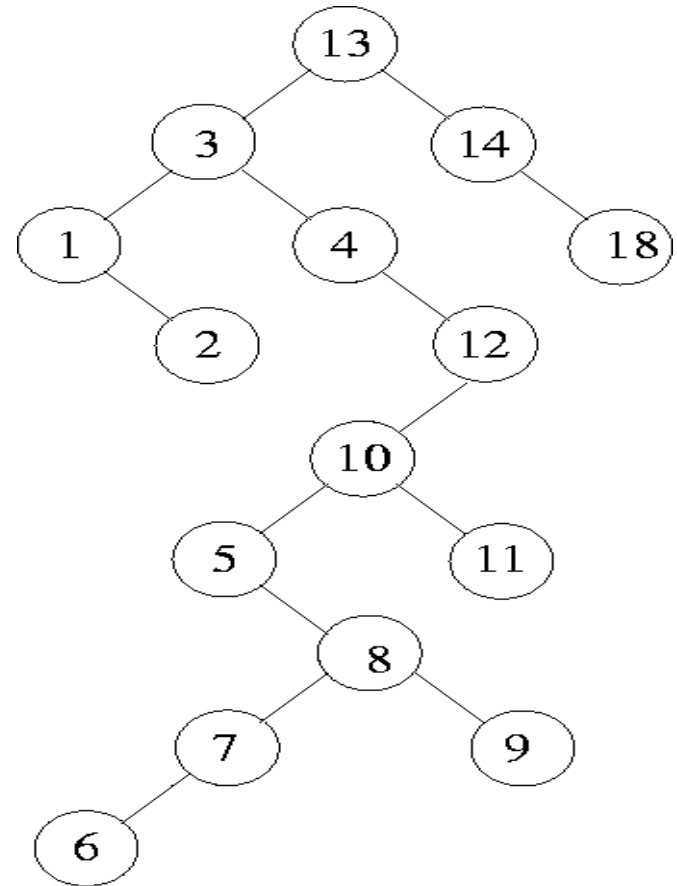
31

- There are many operations, one can perform on a *binary search tree*.
- **Creating** a binary search tree
- **Inserting** a node into a binary search tree
- **Finding** a node in a binary search tree
- **Deleting** a node in a binary search tree.

BST

32

13, 3, 4, 12, 14, 10, 5, 1, 8, 2, 7, 9, 11, 6, 18



Implementation

33

- The basis of our binary tree node is the following struct declaration:

```
struct TreeNode
{
    int value;
    TreeNode *left;
    TreeNode *right;
};
```

```

class IntBinaryTree{
private:
    struct TreeNode{
        int value;
        TreeNode *left;
        TreeNode *right;
    };

    TreeNode *root;
    void destroySubTree(TreeNode *);
    void deleteNode(int, TreeNode *&);
    void makeDeletion(TreeNode *&);
    void displayInOrder(TreeNode *);
    void displayPreOrder(TreeNode *);
    void displayPostOrder(TreeNode *);

public:
    IntBinaryTree() { root = NULL; } // Constructor
    ~IntBinaryTree() { destroySubTree(root); } // Destructor
    void insertNode(int);
    bool searchNode(int);
    void remove(int);
    void showNodesInOrder() { displayInOrder(root); }
    void showNodesPreOrder() { displayPreOrder(root); }
    void showNodesPostOrder() { displayPostOrder(root); }
};

```

Implementation

35

- The **root pointer** is the pointer to the binary tree. This is similar to the *head* pointer in a linked list.
- The **root pointer** will point to the first node in the tree, or to NULL (if the tree is empty).
- It is initialized in the constructor.
- The destructor calls **destroySubTree**, a private member function, that **recursively deletes all** the nodes in the tree.

Insertion

36

- The code to insert a new value in the tree is fairly straightforward.
- First, a new node is allocated, and its *value* member is initialized with the new value.

Insertion

37

Insert a Node

- First, a **new node is allocated** and its value member is initialized with the new value.
- The **left and right child pointers** are set to NULL, because all nodes must be inserted as leaf nodes.
- Next, we determine **if the tree is empty**. If so, we simply make root point to it, and there is nothing else to be done.
- But, if there are nodes in the tree, we must find the new node's proper insertion point.

Insertion

38

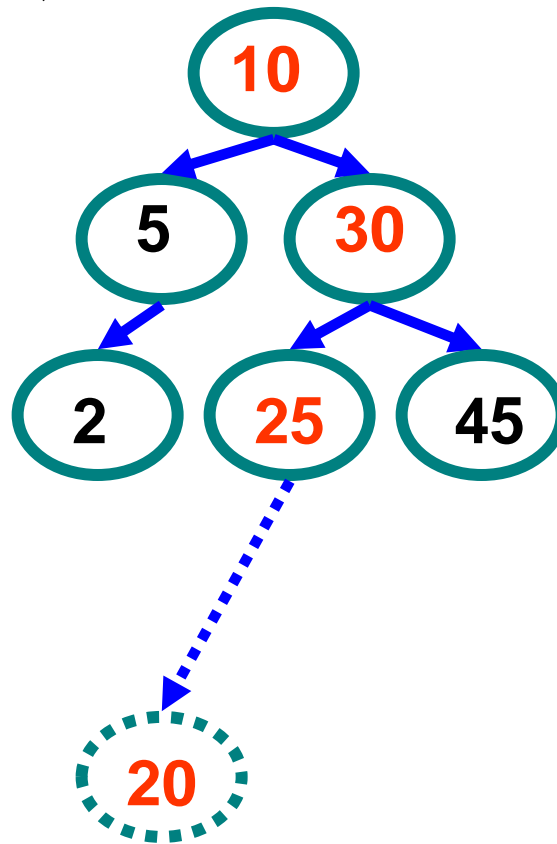
Insert a Node

- If the new value is less than the root node's value, we know it will be inserted somewhere in the left subtree. *Otherwise*, the value will be inserted into the right subtree.
- We simply *traverse the subtree*, comparing each node along the way with the new node's value, and deciding if we should continue to the left or the right.
- *When we reach a child pointer that is set to NULL, we have found out insertion point.*

Insertion

39

Insert (20)



$10 < 20$, right

$30 > 20$, left

$25 > 20$, left

Insert 20 on left

```

void IntBinaryTree::insertNode(int num){
    TreeNode *newNode,          // Pointer to a new node
              *nodePtr;          // Pointer to traverse the tree
    newNode = new TreeNode; // Create a new node
    newNode->value = num;
    newNode->left = newNode->right = NULL;
    if (!root){root = newNode;}    // Is the tree empty?
    else{
        nodePtr = root;
        while (nodePtr != NULL){
            if (num < nodePtr->value){
                if (nodePtr->left)
                    nodePtr = nodePtr->left;
                else {
                    nodePtr->left = newNode;    break;
                }
            }
            else if (num > nodePtr->value) {
                if (nodePtr->right)
                    nodePtr = nodePtr->right;
                else {
                    nodePtr->right = newNode; break;
                }
            }
            else {
                cout << "Duplicate value found.\n";    break;
            }
        }
    }
}
}

```

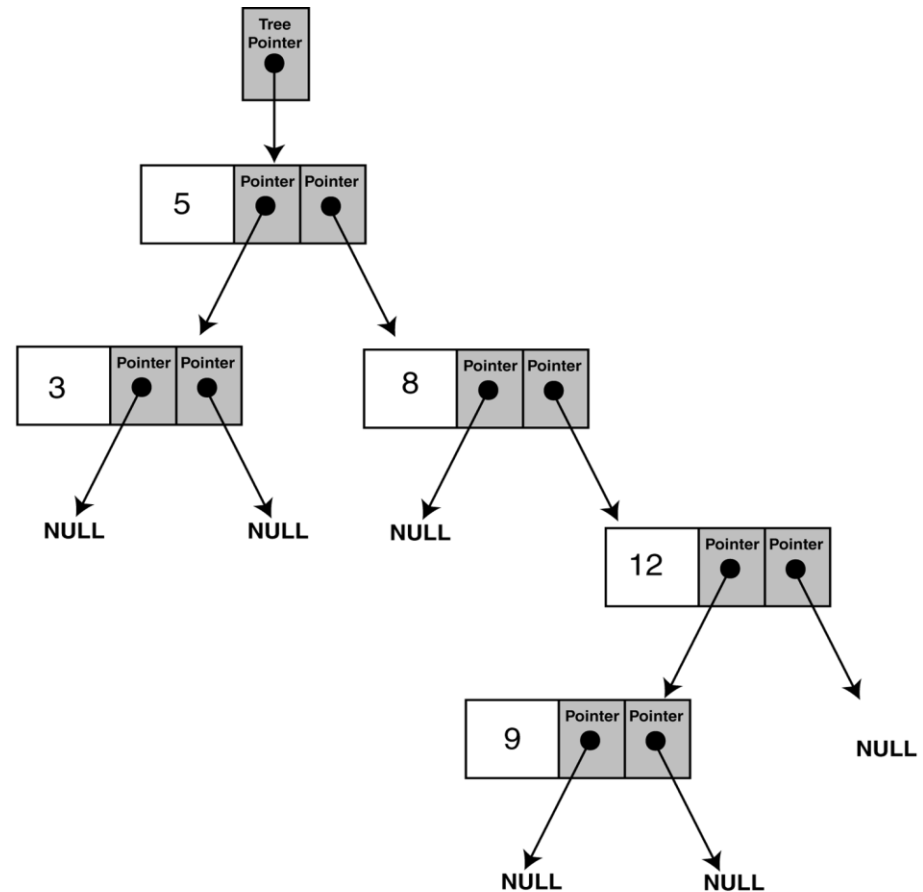


```
// This program builds a binary tree with 5 nodes.
```

```
#include <iostream.h>
#include "IntBinaryTree.h"
```

```
int main(void)
{
    IntBinaryTree tree;

    cout << "Inserting nodes. ";
    tree.insertNode(5);
    tree.insertNode(8);
    tree.insertNode(3);
    tree.insertNode(12);
    tree.insertNode(9);
    cout << "Done.\n";
    return 0;
}
```



Note: The shape of the tree is determined by the **order** in which the values are inserted. The root node in the diagram above holds the value 5 because that was the first value inserted.

Traversal

42

- The IntBinaryTree class can *display all* the values in the tree using all 3 of these algorithms.
- The algorithms are initiated by the following inline public member functions -

```
void showNodesInOrder(void)
    {          displayInOrder(root); }
void showNodesPreOrder()
    {          displayPreOrder(root); }
void showNodesPostOrder()
    {          displayPostOrder(root); }
```

Traversal

43

- Each of these public member functions **calls a recursive private member function and** passes the root pointer as argument.

```
void showNodesInOrder(void)
    {      displayInOrder(root) ; }
void showNodesPreOrder()
    {      displayPreOrder(root) ; }
void showNodesPostOrder()
    {      displayPostOrder(root) ; }
```

Traversal

44

```
void IntBinaryTree::displayInOrder(TreeNode *nodePtr)
{
    if (nodePtr)
    {
        displayInOrder(nodePtr->left);
        cout << nodePtr->value << endl;
        displayInOrder(nodePtr->right);
    }
}
```

Traversal

45

```
void IntBinaryTree::displayPreOrder(TreeNode *nodePtr)
{
    if (nodePtr)
    {
        cout << nodePtr->value << endl;
        displayPreOrder(nodePtr->left);
        displayPreOrder(nodePtr->right);
    }
}
```

Traversal

46

```
void IntBinaryTree::displayPostOrder(TreeNode *nodePtr)
{
    if (nodePtr)
    {
        displayPostOrder(nodePtr->left);
        displayPostOrder(nodePtr->right);
        cout << nodePtr->value << endl;
    }
}
```

```
// This program builds a binary tree with 5 nodes.
```

```
#include <iostream.h>
#include "IntBinaryTree.h"
```

```
int main(void)
{
```

```
    IntBinaryTree tree;
```

```
    cout << "Inserting nodes. ";
```

```
    tree.insertNode(5);
```

```
    tree.insertNode(8);
```

```
    tree.insertNode(3);
```

```
    tree.insertNode(12);
```

```
    tree.insertNode(9);
```

```
    cout << "Done.\n";
```

```
    cout << "Inorder traversal:\n";
```

```
    tree.showNodesInOrder();
```

```
    cout << "\nPreorder traversal:\n";
```

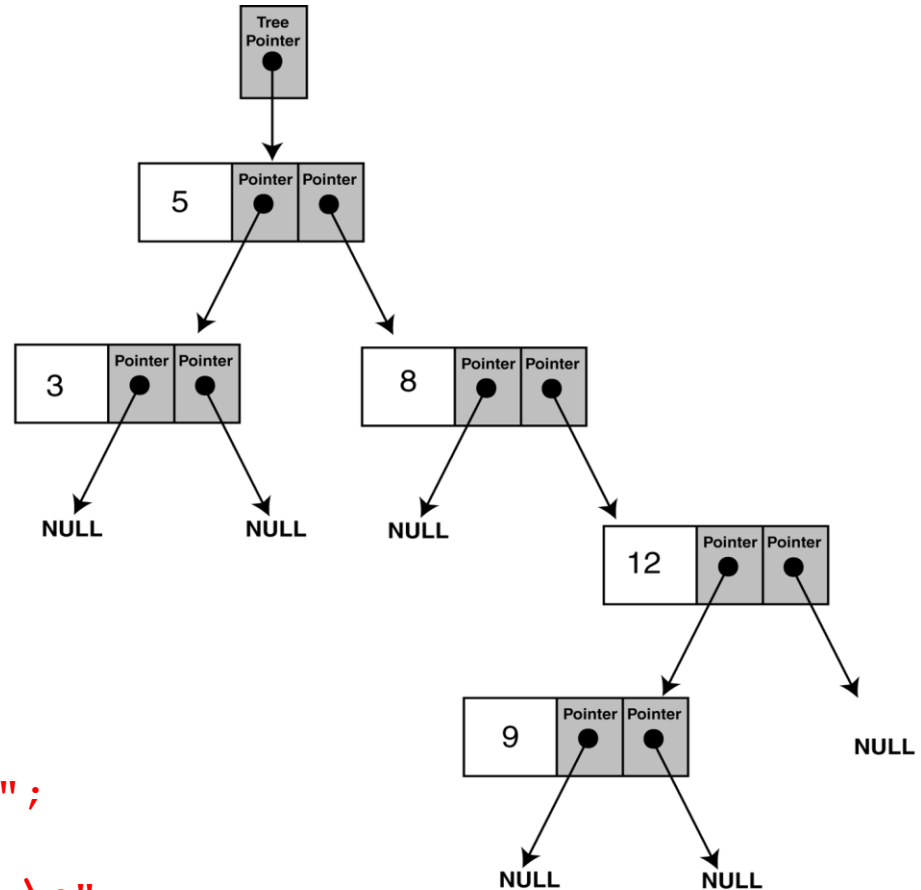
```
    tree.showNodesPreOrder();
```

```
    cout << "\nPostorder traversal:\n";
```

```
    tree.showNodesPostOrder();
```

```
    return 0;
```

```
}
```



```
// This program builds a binary tree with 5 nodes.
```

```
#include <iostream.h>
#include "IntBinaryTree.h"
```

```
int main(void)
{
```

```
    IntBinaryTree tree;
```

```
    cout << "Inserting nodes. ";
```

```
    tree.insertNode(5);
```

```
    tree.insertNode(8);
```

```
    tree.insertNode(3);
```

```
    tree.insertNode(12);
```

```
    tree.insertNode(9);
```

```
    cout << "Done.\n";
```

```
    cout << "Inorder traversal:\n";
```

```
    tree.showNodesInOrder();
```

```
    cout << "\nPreorder traversal:\n";
```

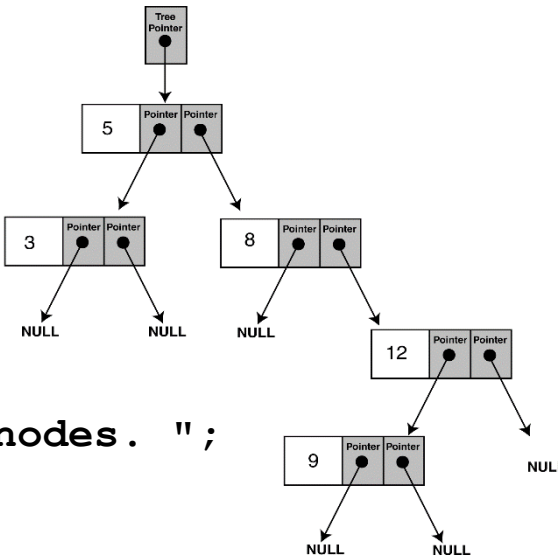
```
    tree.showNodesPreOrder();
```

```
    cout << "\nPostorder traversal:\n";
```

```
    tree.showNodesPostOrder();
```

```
    return 0;
```

```
}
```



Program Output

Inserting nodes.

Inorder traversal:

3

5

8

9

12

Preorder traversal:

5

3

8

12

9

Postorder traversal:

3

9

12

8

5

Recursive Search of Binary Tree

49

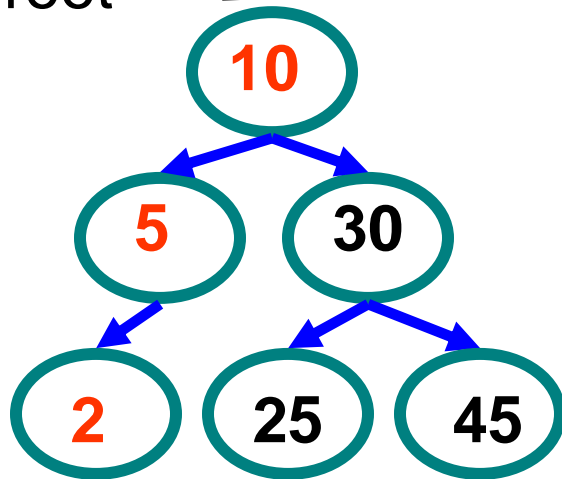
```
Node *Find( Node *n, int key) {  
    if (n == NULL)                // Not found  
        return( n );  
    else if (n->data == key)       // Found it  
        return( n );  
    else if (n->data > key)        // In left subtree  
        return Find( n->left, key );  
    else                          // In right subtree  
        return Find( n->right, key );  
}  
Node * n = Find( root, 5);
```

Examples

50

Find (root, 2)

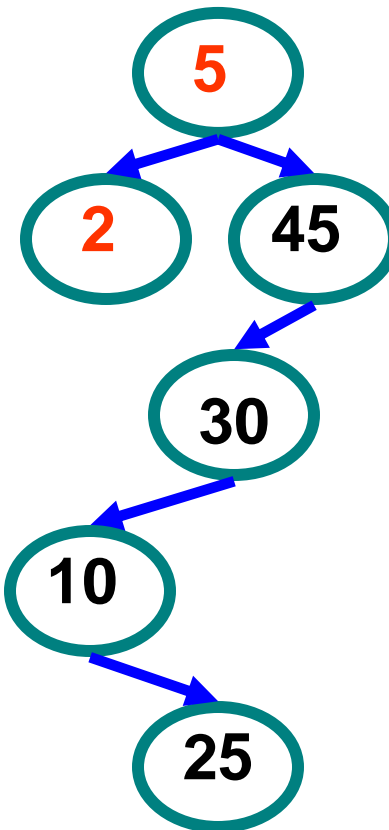
root
→



$10 > 2$, left

$5 > 2$, left

$2 = 2$, found



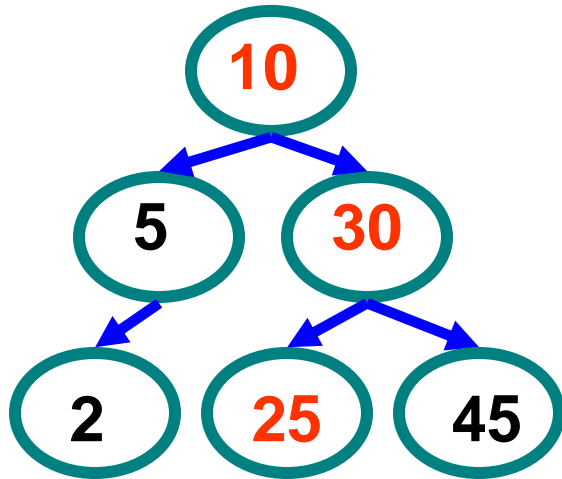
$5 > 2$, left

$2 = 2$, found

Examples

51

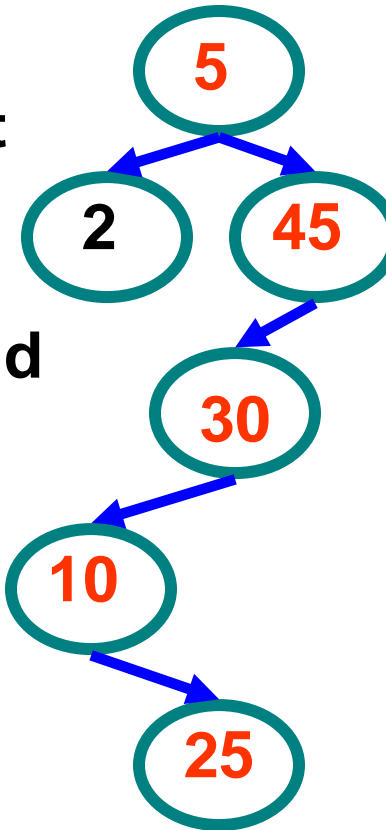
Find (root, 25)



$10 < 25$, right

$30 > 25$, left

$25 = 25$, found



$5 < 25$, right

$45 > 25$, left

$30 > 25$, left

$10 < 25$, right

$25 = 25$, found

Reading Materials

52

- Schaum's Outlines: Chapter # 7
- D. S. Malik: Chapter # 11
- Nell Dale: Chapter # 8
- Allen Weiss: Chapter # 4
- Tenebaum: Chapter # 5

