

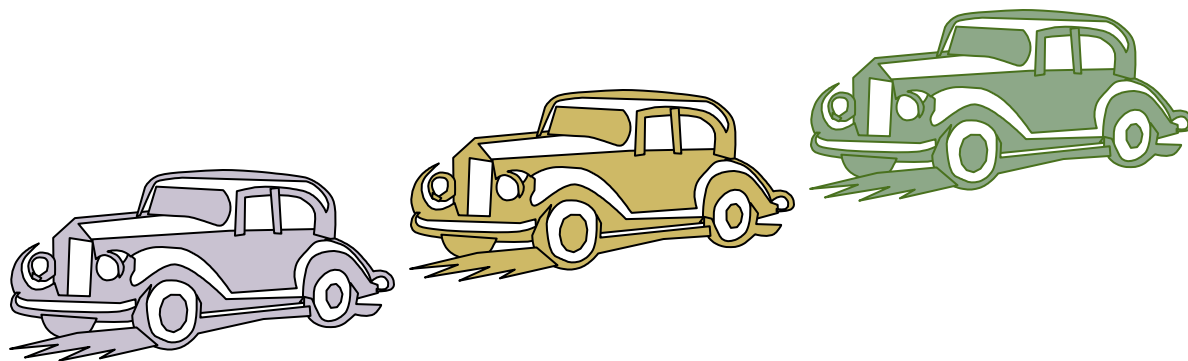


CS-218

DATA STRUCTURE

Dr. Hashim Yasin

**National University of Computer
and Emerging Sciences,
Faisalabad, Pakistan.**



QUEUES

Queues

3

A **Queue** is a special kind of list, where items are,

- inserted at one end (*the rear*) and
- deleted at the other end (*the front*)

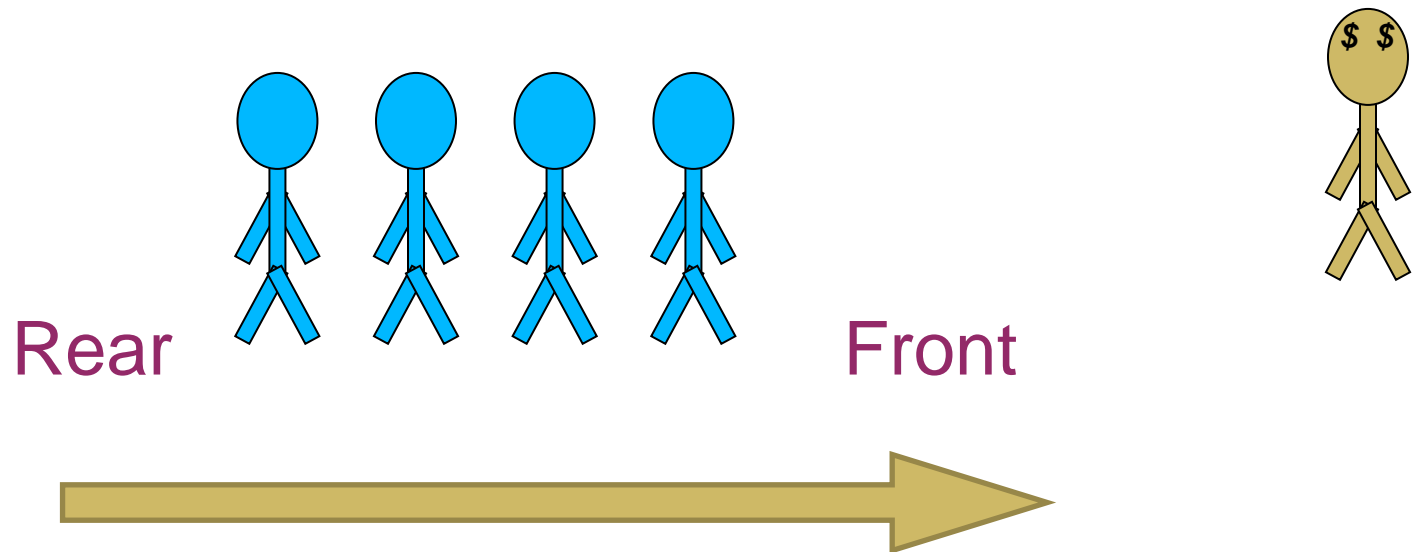
Other Name:

- First In First Out (FIFO)

Queues

4

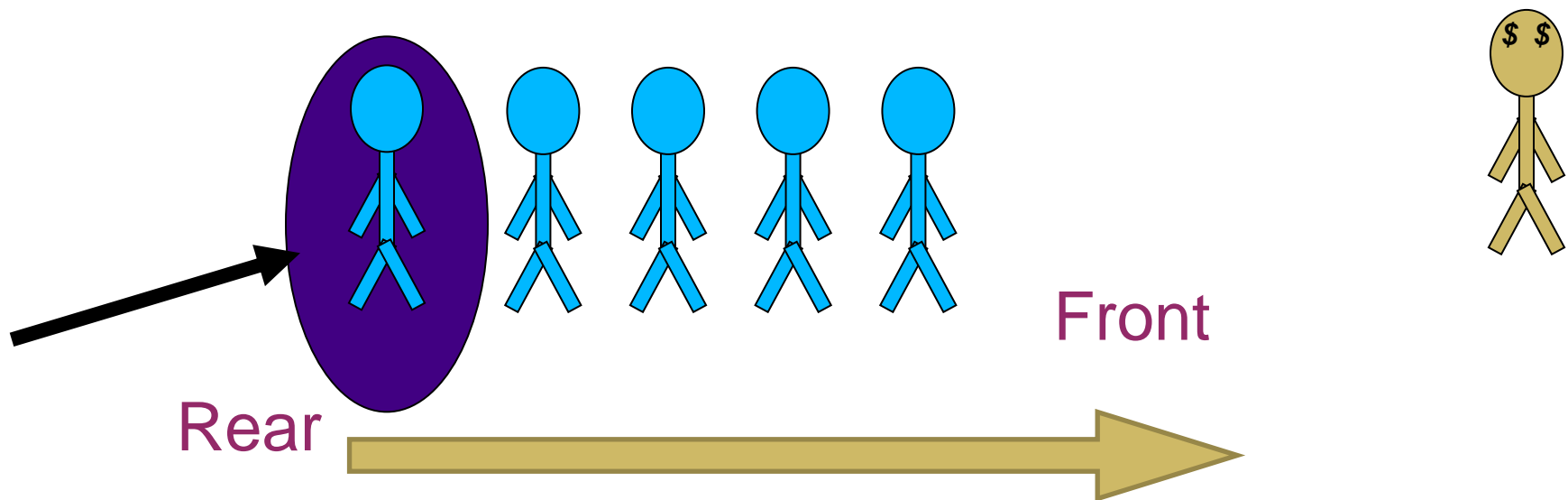
- A queue is like a line of people waiting for a bank teller.
- The queue has a **front** and a **rear**.



Queues

5

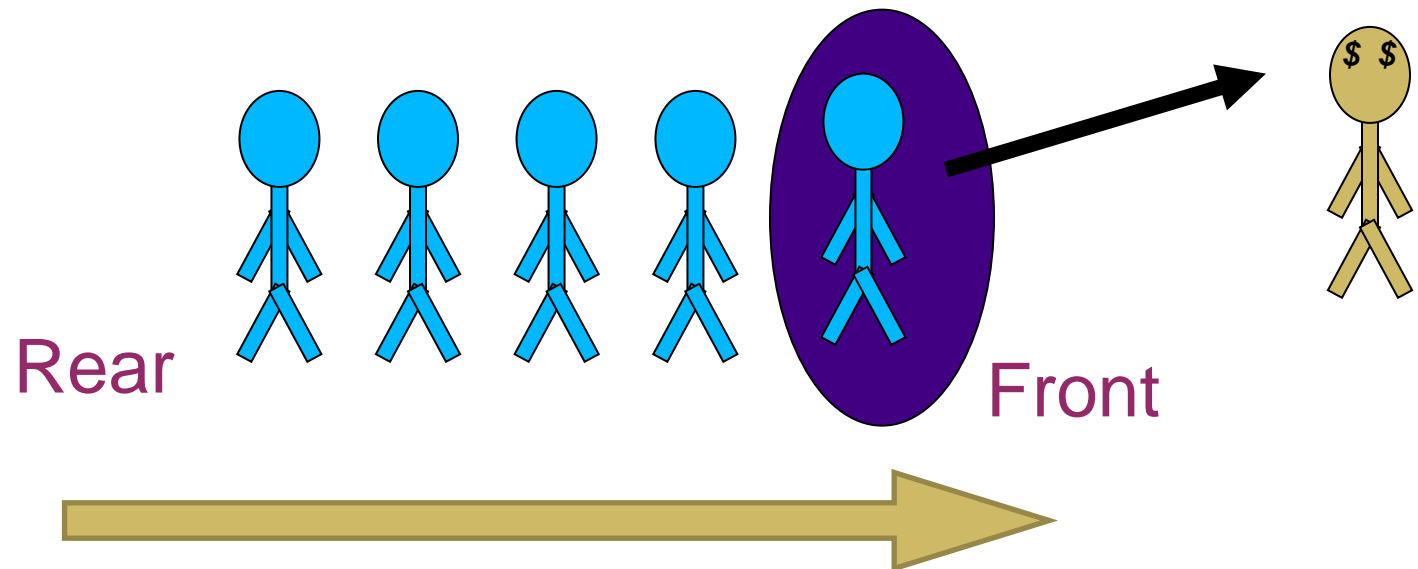
- New people must enter the queue at the **rear**.



Queues

6

- When an item is taken from the queue, it always comes from the **front**.



Examples

7

- Billing counter
 - ▣ Booking movie tickets
 - ▣ Queue for paying bills

- A print queue
- Vehicles on toll-tax bridge
- Luggage checking machine

Applications

8

- **Operating system**

- ▣ **multi-user/multitasking environments**, where several users or tasks may be requesting the same resource simultaneously.

- **Communication Software**

- ▣ queues to hold *information* received over networks and dial up connections.
 - ▣ Information can be transmitted faster than it can be processed, so is placed in a queue waiting to be processed.

Common Operations

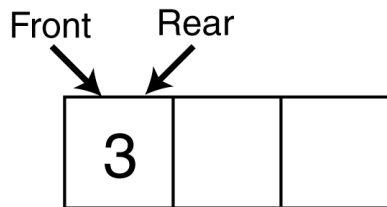
9

1. **MAKENULL(Q)**: Makes Queue Q be an empty list.
2. **FRONT(Q)**: Returns the first element on Queue Q.
3. **ENQUEUE(x, Q)**: Inserts element x at the end of Queue Q.
4. **DEQUEUE(Q)**: Deletes the first element of Q.
5. **EMPTY(Q)**: Returns true if and only if Q is an empty queue.

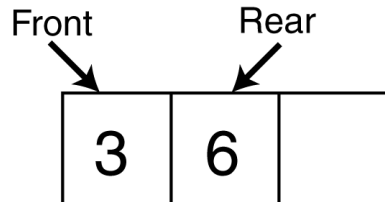
Enqueue & Dequeue

10

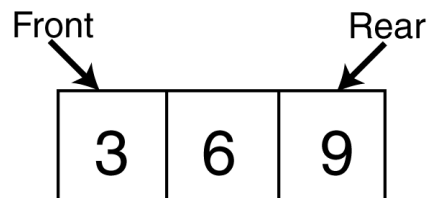
Enqueue(3);



Enqueue(6);



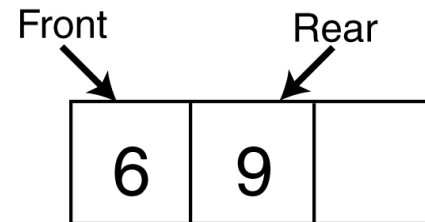
Enqueue(9);



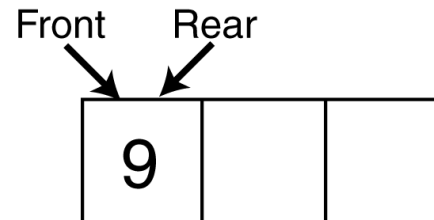
Dr Hashim Yasin

...

Dequeue();

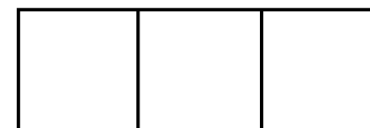


Dequeue();



Dequeue();

Front = -1 Rear = -1



CS-218 Data Structure

Enqueue Operation

11

- **Step 1** – Check if the queue is full.
- **Step 2** – If the queue is full, produce overflow error and exit.
- **Step 3** – If the queue is not full, increment rear pointer to point the next empty space.
- **Step 4** – Add data element to the queue location, where the rear is pointing.
- **Step 5** – Return success.

Dequeue Operation

12

- **Step 1** – Check if the queue is empty.
- **Step 2** – If the queue is empty, produce underflow error and exit.
- **Step 3** – If the queue is not empty, access the data where front is pointing.
- **Step 4** – Increment front pointer to point to the next available data element.
- **Step 5** – Return success.

Implementation

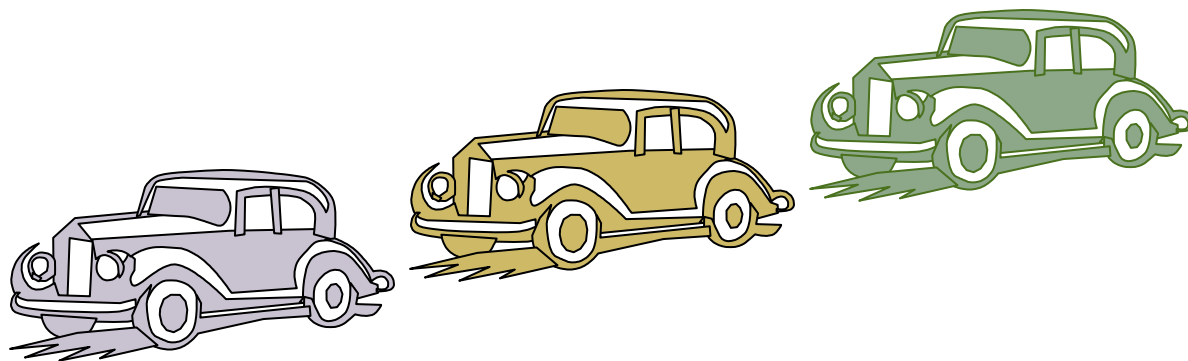
13

□ Static

- Queue is implemented by **an array**, and size of queue remains fix

□ Dynamic

- A queue can be implemented as a **linked list** and *expand* or *shrink* with each *enqueue* or *dequeue* operation.



ARRAY IMPLEMENTATION



Array Implementation

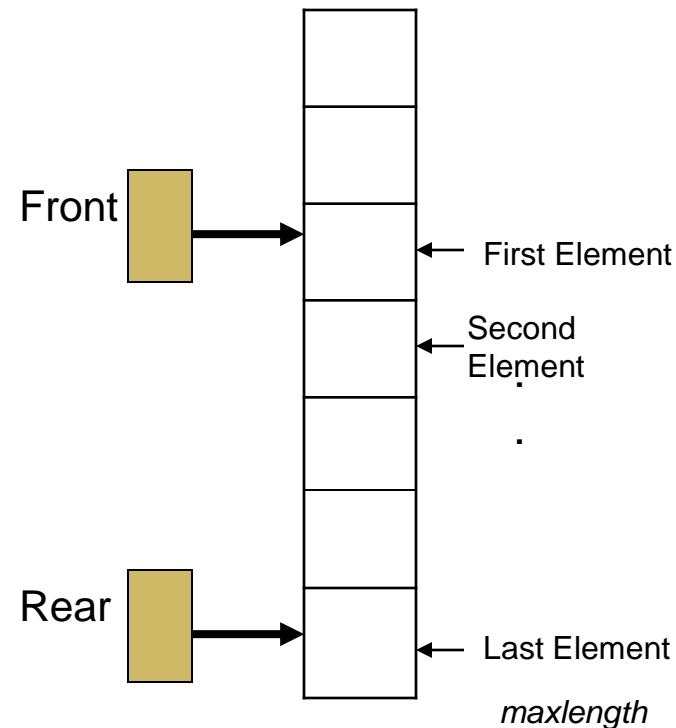
15

- As with Stacks, signify zero index as rear.
- **Enqueue**
 - ▣ Shift elements to the right
 - ▣ As expensive as with stacks
- **Dequeue**
 - ▣ Need to save index of first item inserted
- **On Dequeue, decrement index**
- **On Enqueue, increment index**

Array Implementation

16

- Use two counters that signify **rear** and **front**



- When queue is empty, both front and rear are set to -1
- While enqueueing increment rear by 1, and while dequeuing increment front by 1
- When there is only one value in the Queue, both rear and front have same index

Array Implementation

17

5	4	6	7	8	7	6		
0	1	2	3	4	5	6	7	8

Front=0

Rear=6

				8	7	6		
0	1	2	3	4	5	6	7	8

Front=4

Rear=6

					7	6	12	67
0	1	2	3	4	5	6	7	8

Front=5

Rear=8

Array Implementation

18

5	4	6	7	8	7	6		
0	1	2	3	4	5	6	7	8

Front=0

Rear=6

				8	7	6		
0	1	2	3	4	5	6	7	8

Front=4

Rear=6

					7	6	12	67
0	1	2	3	4	5	6	7	8

Front=5

Rear=8

*How can we insert more elements? Rear index
can not move beyond the last element....*

Array Implementation

19

How can we insert more elements? Rear index can not move beyond the last element....

Solution:

Using Circular Queue

Circular Queue

20

- Allow rear to wrap around the array.

```
if(rear == queueSize-1)
```

```
    rear = 0;
```

```
else
```

```
    rear++;
```

- Or use modular arithmetic

```
rear = (rear + 1) % queueSize;
```

Circular Queue

21

					7	6	12	67
0	1	2	3	4	5	6	7	8

Front=5

Rear=8

Enqueue 39, $\text{Rear} = (\text{Rear} + 1) \bmod \text{Queue Size} = (8 + 1) \bmod 9 = 0$

39					7	6	12	67
0	1	2	3	4	5	6	7	8

Front=5

Rear=0

IMPLEMENTATION

Implementation ... Queue Class

23

```
class IntQueue {
private:
    int *queueArray;
    int queueSize;
    int front;
    int rear;
    int numItems;
public:
    IntQueue(int) ;
    ~IntQueue(void) ;
    void enqueue(int) ;
    int dequeue(void) ;
    bool isEmpty(void) ;
    bool isFull(void) ;
    void clear(void) ;
};
```

Note, the member function **clear, which clears the queue by resetting the front and rear indices and setting the numItems to 0.**

Implementation ... Queue Class

24

```
IntQueue::IntQueue(int s) //constructor
{
    queueArray = new int[s];
    queueSize = s;
    front = -1;
    rear = -1;
    numItems = 0;
}
```

```
IntQueue::~~IntQueue(void) //destructor
{
    delete [] queueArray;
}
```


Implementation ... Enqueue Function

25

```
//*****  
// Function enqueue inserts the value in num *  
// at the rear of the queue. *  
//*****  
void IntQueue::enqueue(int num){  
    if (isFull())  
        cout << "The queue is full.\n";  
    else{  
        // Calculate the new rear position  
        rear = (rear + 1) % queueSize;  
        // Insert new item  
        queueArray[rear] = num;  
        // Update item count  
        numItems++;  
    }  
}
```

Implementation ... Dequeue Function

26

```
//*****  
// Function dequeue removes the value at the      *  
// front of the queue, and copies it into num. *  
//*****  
int IntQueue::dequeue(void) {  
    if (isEmpty())  
        cout << "The queue is empty.\n";  
    else{  
        // Move front  
        front = (front + 1) % queueSize;  
        // Retrieve the front item  
        int num = queueArray[front];  
        // Update item count  
        numItems--;  
    }  
    return num;  
}
```

Implementation ... isEmpty Function

27

```
/** *****  
// Function isEmpty returns true if the queue *  
// is empty, and false otherwise. *  
/** *****  
  
bool IntQueue::isEmpty(void) {  
    if (numItems)  
        return false;  
    else  
        return true;  
}
```

Implementation ... isFull Function

28

```
//*****  
// Function isFull returns true if the queue *  
// is full, and false otherwise. *  
//*****  
  
bool IntQueue::isFull(void){  
    if (numItems < queueSize)  
        return false;  
    else  
        return true;  
}
```

Implementation ... Clear Function

29

```
//*****  
// Function clear resets the front and rear *  
// indices and sets numItems to 0.          *  
//*****  
  
void IntQueue::clear(void) {  
    front = - 1;  
    rear = - 1;  
    numItems = 0;  
}
```

Implementation ... Demonstration

30

```
//Program demonstrating the IntQueue class
void main(void){
    IntQueue iQueue(5);
    cout << "Enqueuing 5 items...\n";
    // Enqueue 5 items.
    for (int x = 0; x < 5; x++)
        iQueue.enqueue(x);
    // Attempt to enqueue a 6th item.
    cout << "Now attempting to enqueue again...\n";
    iQueue.enqueue(5);

    // Dequeue and retrieve all items in the queue
    cout << "The values in the queue were:\n";
    while (!iQueue.isEmpty()){
        int value;
        value = iQueue.dequeue();
        cout << value << endl;
    }
}
```

Implementation ... Output

31

Program Output:

Enqueueing 5 items...

Now attempting to enqueue again...

The queue is full.

The values in the queue were:

0

1

2

3

4

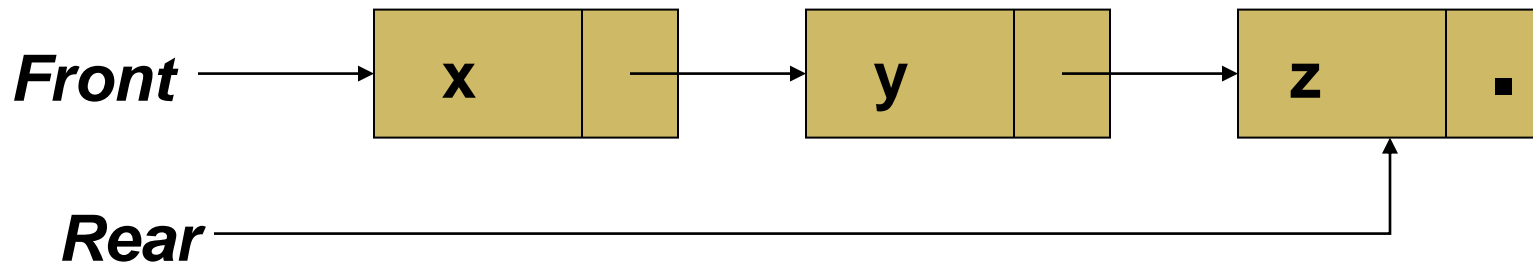
A POINTER-BASED IMPLEMENTATION

A Pointer based Implementation

33

Keep two pointers:

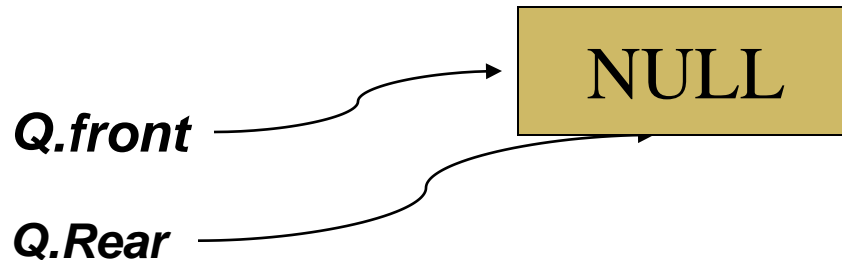
- **FRONT:** A pointer to the first element of the queue.
- **REAR:** A pointer to the last element of the queue.



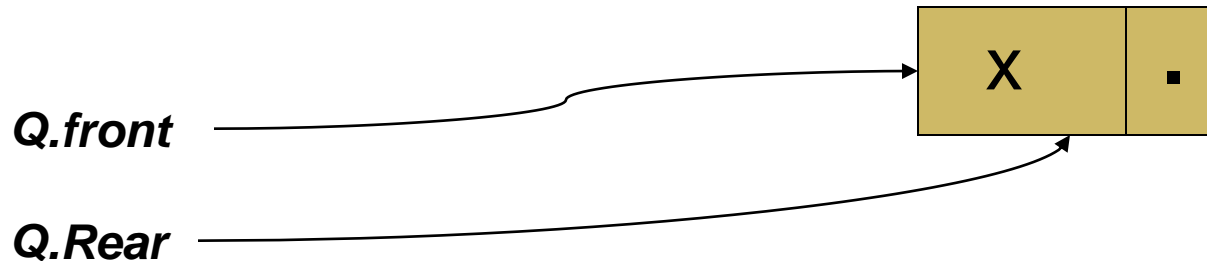
A Pointer based Implementation

34

MAKENULL(Q)



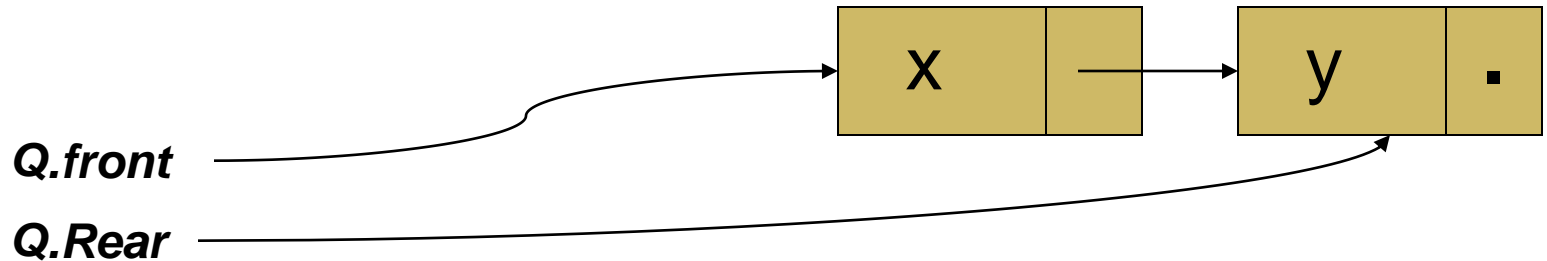
ENQUEUE(x,Q)



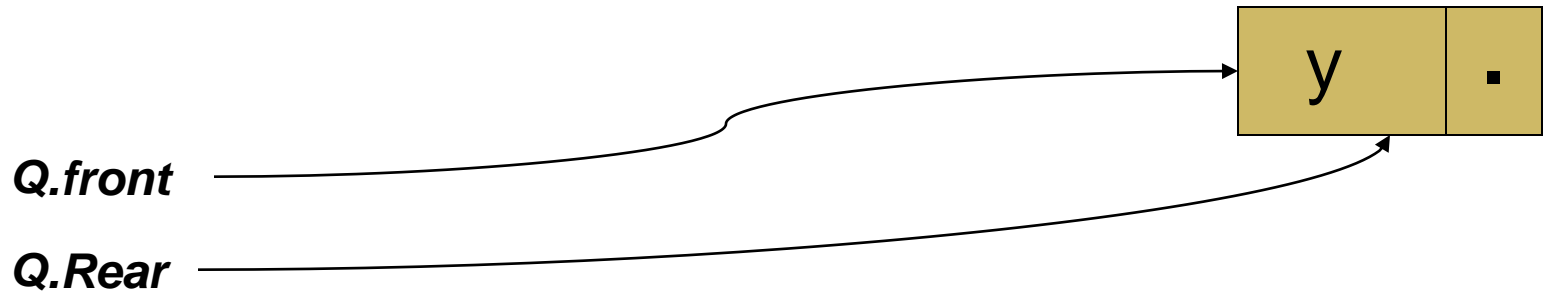
A Pointer based Implementation

35

ENQUEUE(y,Q)



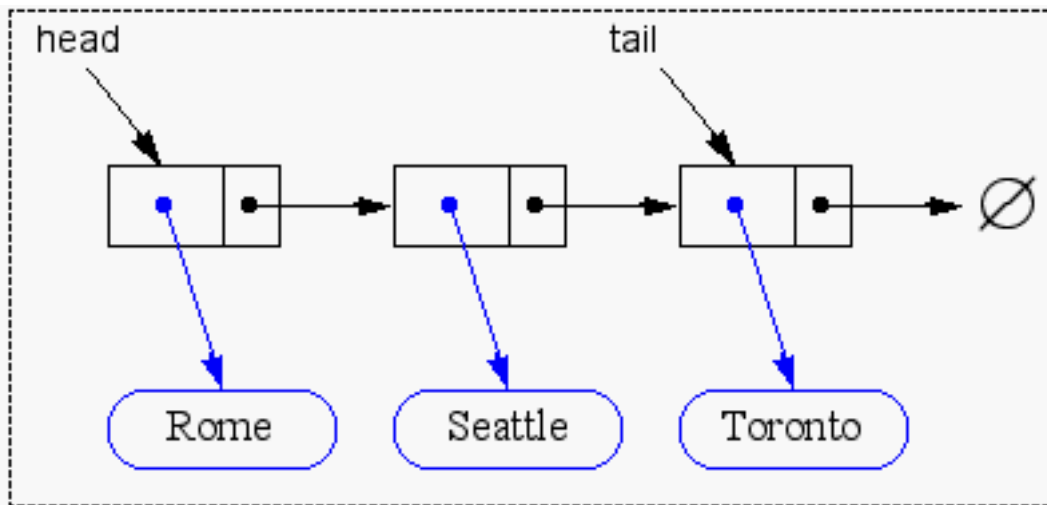
DEQUEUE(Q)



Singly Linked List based Implementation

36

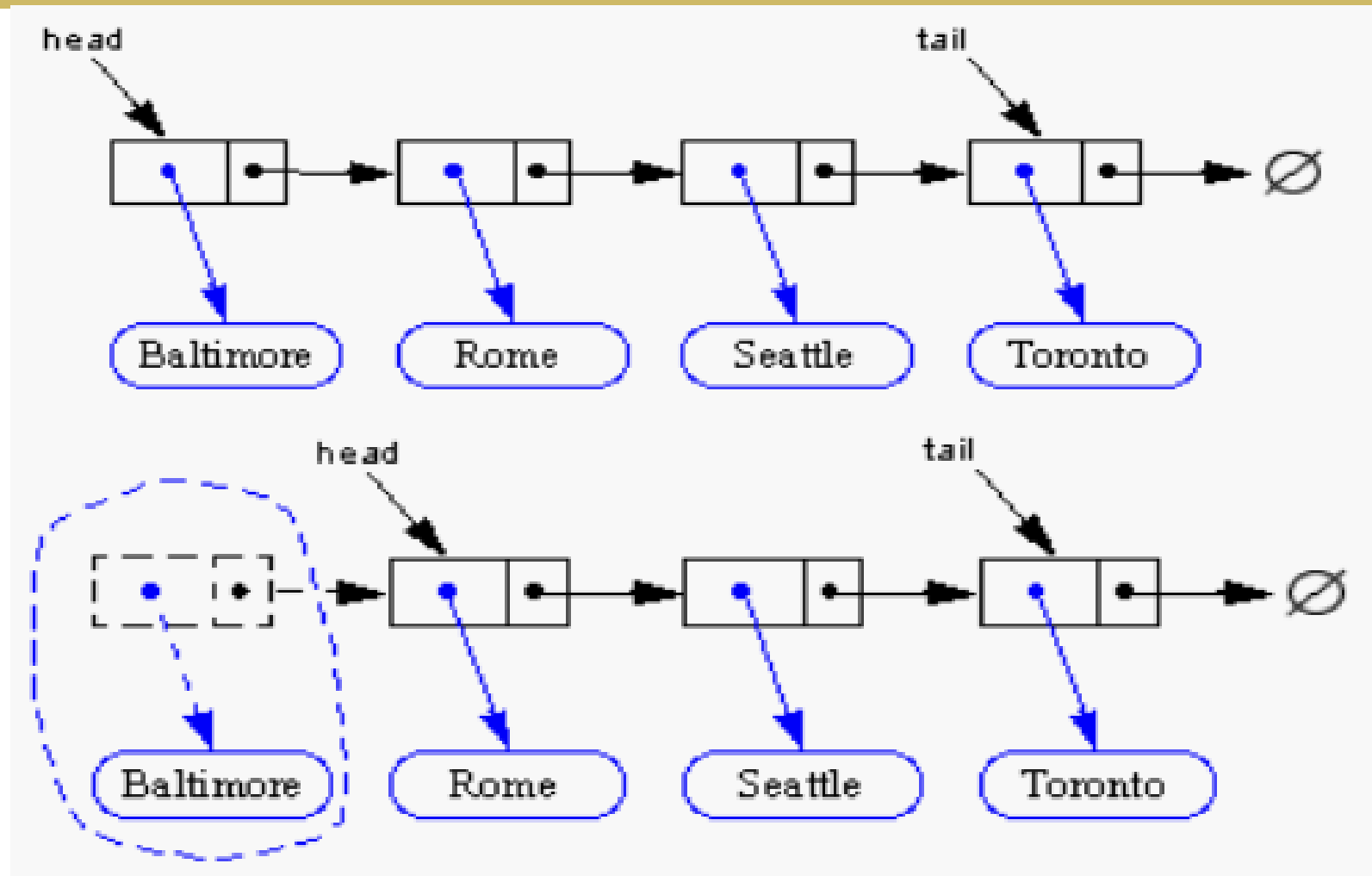
Nodes connected in a chain by links



The *head of the list is the front of the queue*, the *tail of the list is the rear of the queue*.

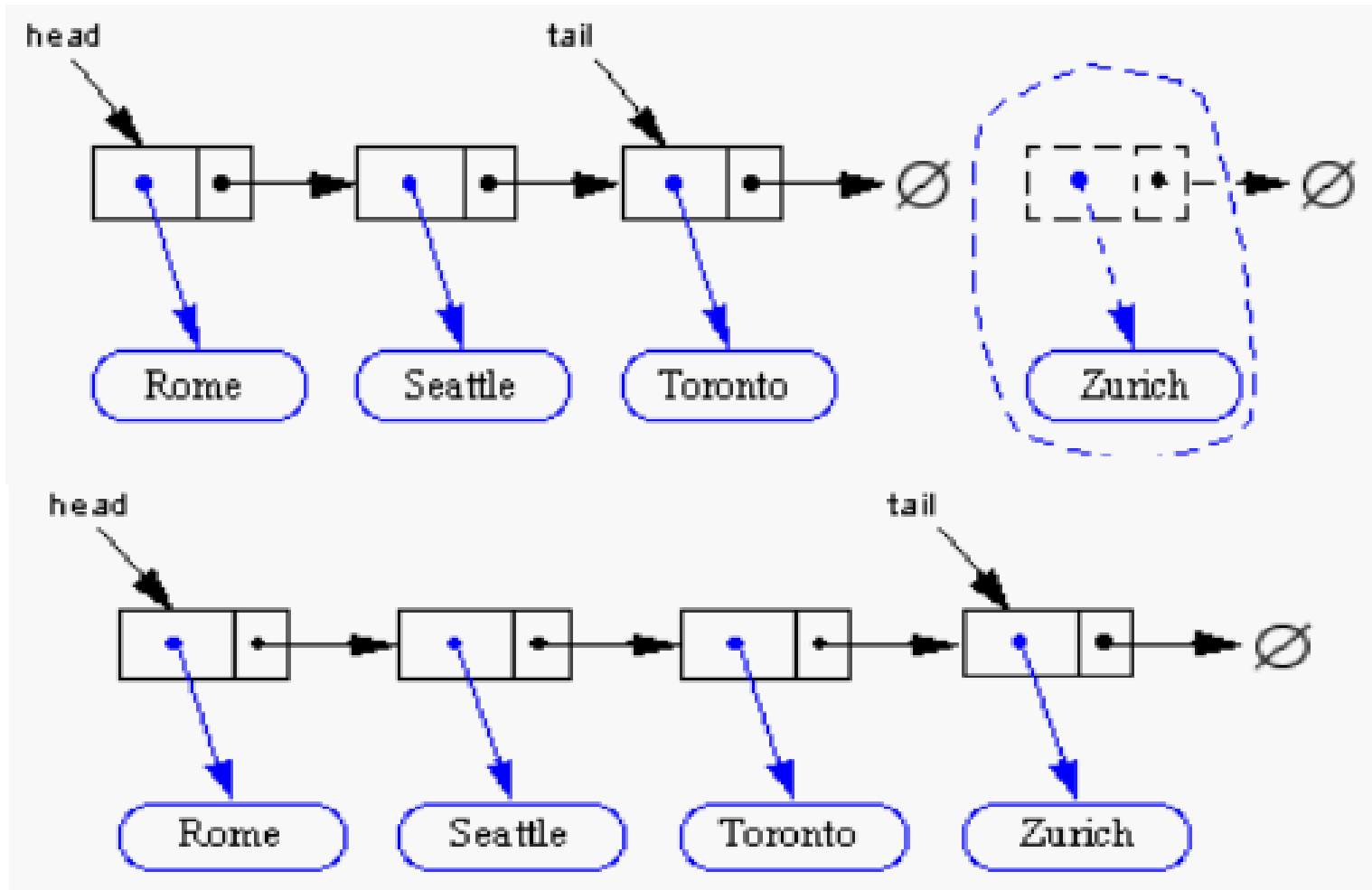
Removing at the Head

37



Inserting at the Tail

38



Homework

39

- Implement the same program given in previous slides with the help of,
 - ▣ Singly Linked List
 - ▣ Doubly Linked List

PRIORITY QUEUE



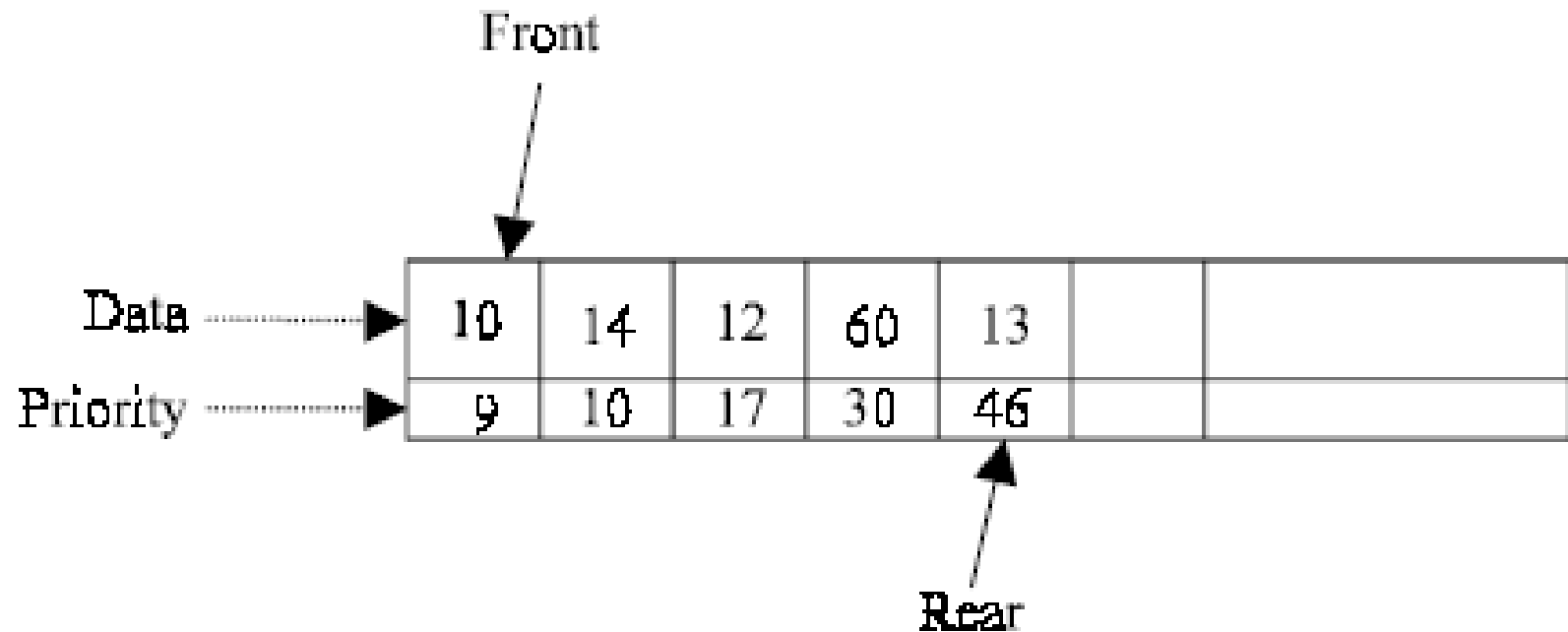
Priority Queue

41

- ✓ Priority Queue is a queue where *each element is assigned a priority.*
- ✓ In priority queue, the elements are deleted and processed by following rules.
 - An element of *higher priority* is processed before any element of lower priority.
 - Two elements with the *same priority* are processed according to the order in which they were inserted to the queue.

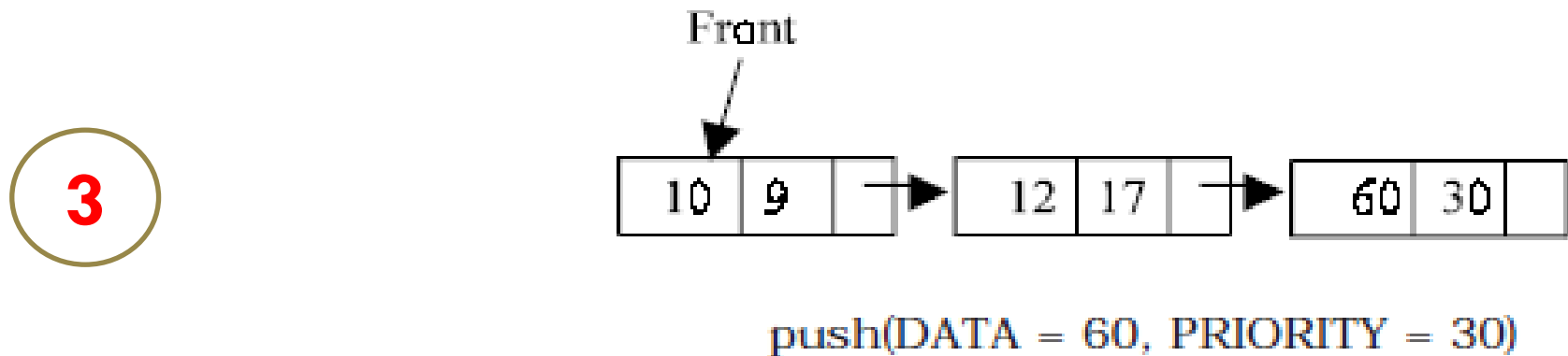
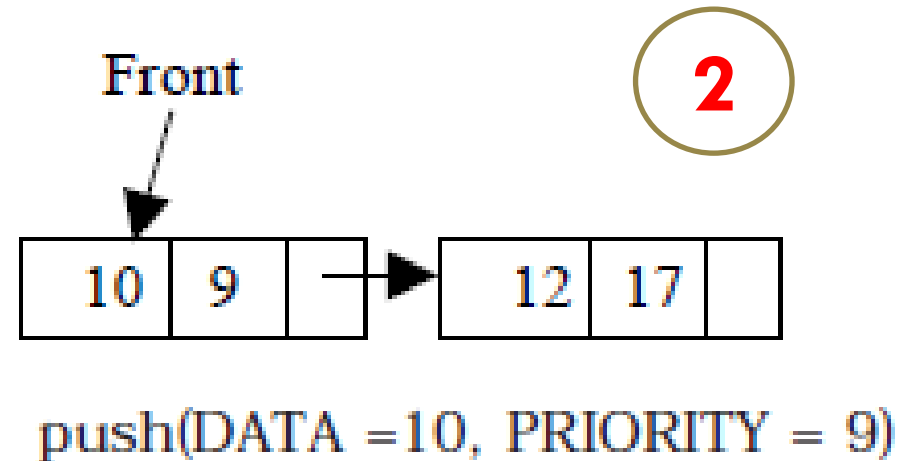
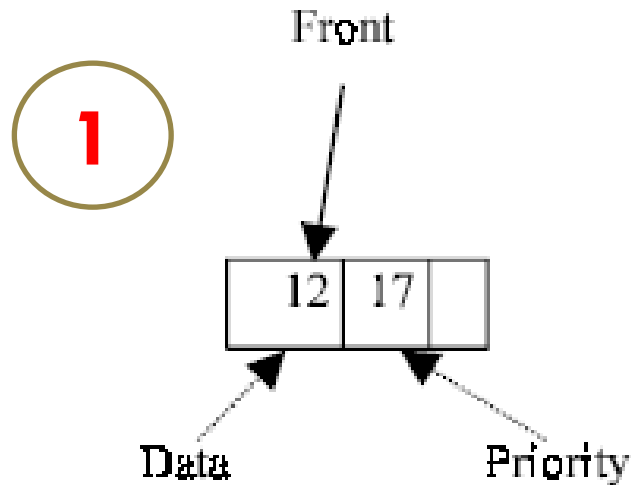
Priority Queue

42



Priority Queue

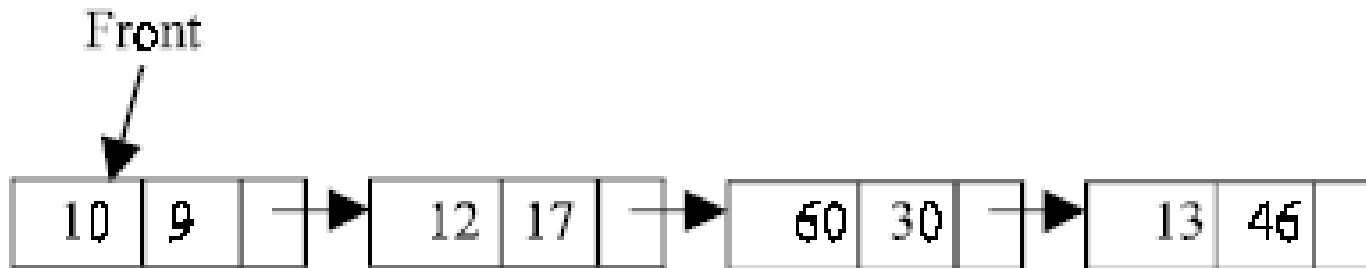
43



Priority Queue

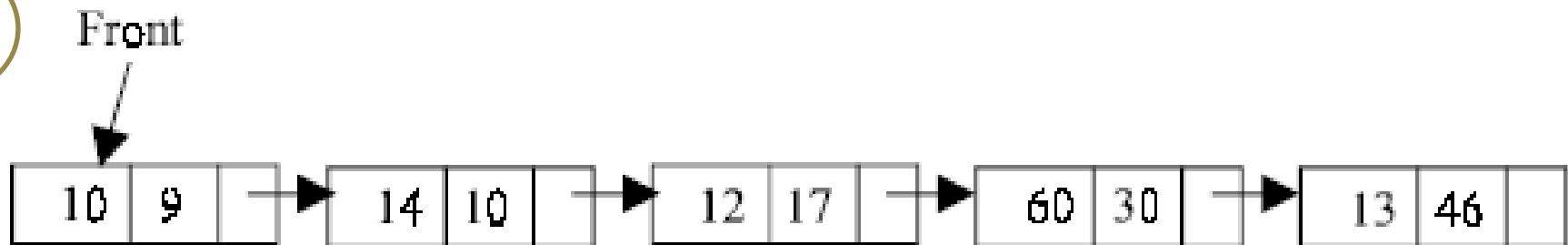
44

4



push(DATA = 13, PRIORITY = 46)

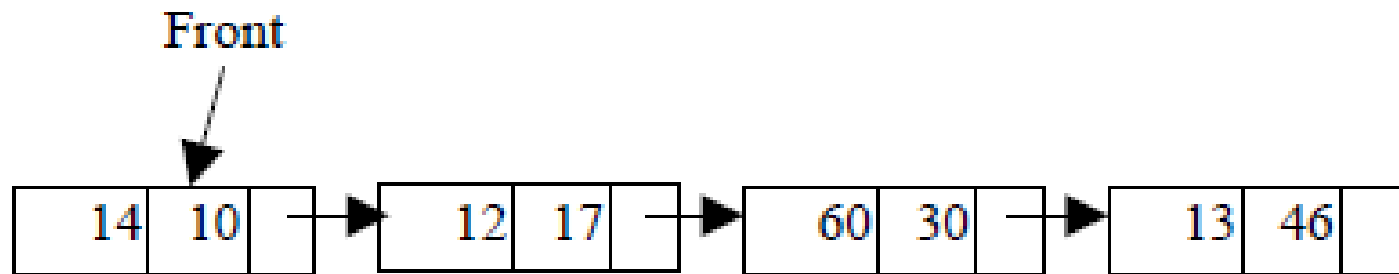
5



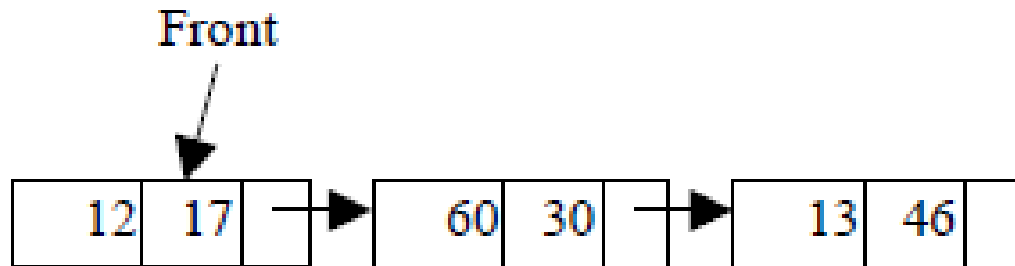
push(DATA = 14, PRIORITY = 10)

Priority Queue

45



$x = \text{pop}()$ (i.e., 10)



$x = \text{pop}()$ (i.e., 14)

Reading Materials

46

- Chapter 8, Data Structures by Larry Nyhoff