



CS-2001 DATA STRUCTURE

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HEAP

Heap

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- A Heap is a special Tree-based data structure in which the **tree is a complete binary tree**.
- Heap is a special case of **balanced binary tree** data structure where *the root-node key is compared with its children* and arranged accordingly.
- Generally, Heaps can be of two types:
 - **Max-Heap:**
 - **Min-Heap:**

Heap

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Max-Heap:

- In a Max-Heap, the key present at the root node must be maximum among the keys present at all of its children.
- The same property must be recursively true for all sub-trees in that Binary Tree.
- If α has child node β then,
$$\text{key}(\alpha) \geq \text{key}(\beta)$$

Heap

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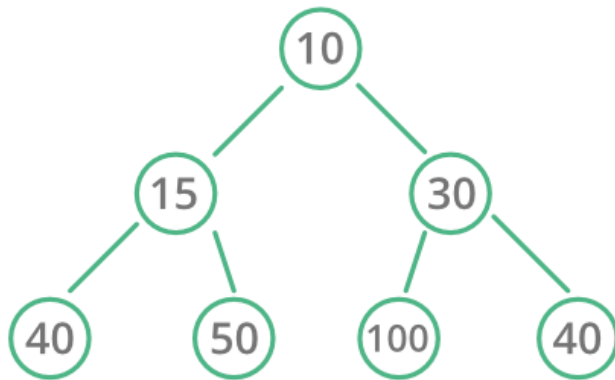
Min-Heap:

- In a Min-Heap, the key present at the root node must be **minimum** among the keys present at all of its children.
- The same property must be **recursively true** for all sub-trees in that Binary Tree.
- If α has child node β then,
$$\text{key}(\alpha) \leq \text{key}(\beta)$$

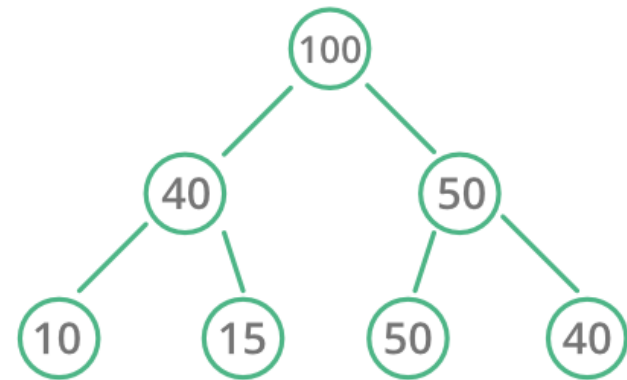
Examples

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Heap Data Structure



Min Heap



Max Heap

GG

MAX HEAP

Max Heap Construction Algorithm

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Step 1 – Create a new node at the end of the heap.

Step 2 – Assign a new value to the node.

Step 3 – Compare the value of this child node with its parent.

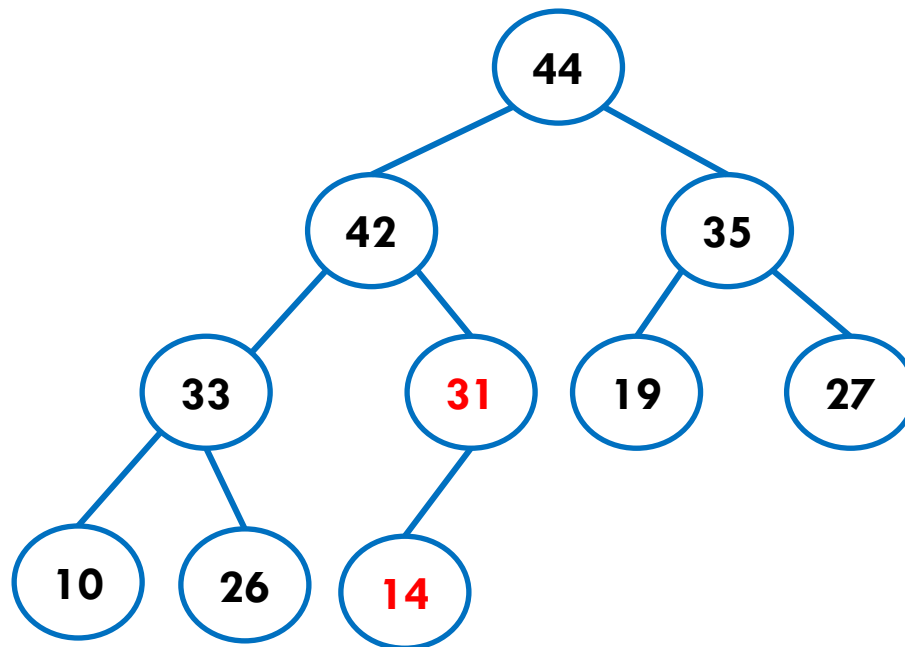
Step 4 – If the value of the parent is less than a child, then swap them.

Step 5 – Repeat steps 3 & 4 until Heap property holds.

Example ... Max Heap

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For Input → 35 33 42 10 14 19 27 44 26 31



Example ... Max Heap

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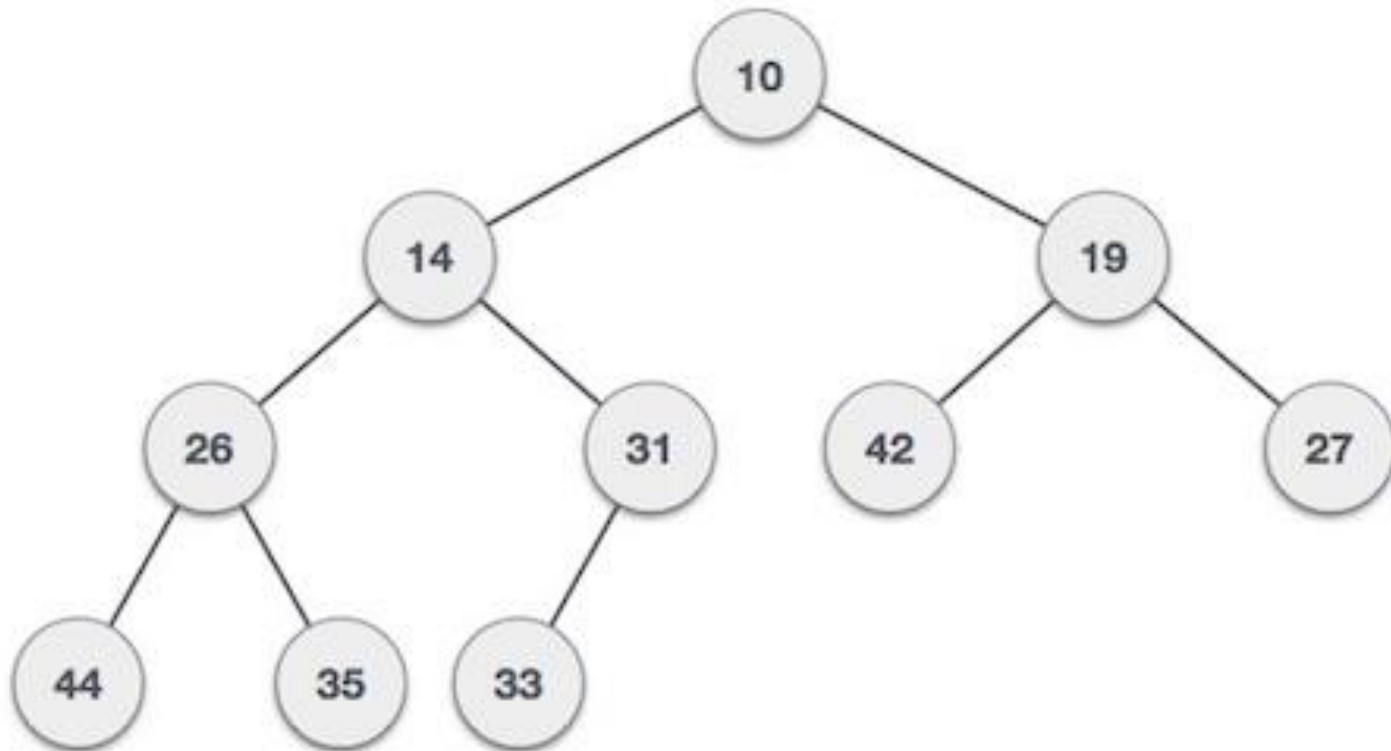
For Input → 35 33 42 10 14 19 27 44 26 31

Input 35 33 42 10 14 19 27 44 26 31

Example ... Min Heap

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For Input → 35 33 42 10 14 19 27 44 26 31



Max Heap ... Deletion

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Step 1 – Remove root node.

Step 2 – Move the last element of the last level to root.

Step 3 – Compare the value of this child node with its parent.

Step 4 – If the value of the parent is less than a child, then swap them.

Step 5 – Repeat steps 3 & 4 until Heap property holds.

HEAP IMPLEMENTATION



Heap ... Array Implementation

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```
#include <iostream>
using namespace std;
// To heapify a subtree rooted with node i which is an index in arr[]. n is size of heap
void heapify(int arr[], int n, int i){
    int largest = i;           // Initialize largest as root
    int l = 2 * i + 1;         // left = 2*i + 1
    int r = 2 * i + 2;         // right = 2*i + 2
    // If left child is larger than root
    if (l < n && arr[l] > arr[largest])
        largest = l;
    // If right child is larger than largest so far
    if (r < n && arr[r] > arr[largest])
        largest = r;
    // If largest is not root
    if (largest != i) {
        swap(arr[i], arr[largest]);
        // Recursively heapify the affected sub-tree
        heapify(arr, n, largest);
    }
}
```

```
// Function to build a Max-Heap from the given array
void buildHeap(int arr[], int n) {
    // Index of last non-leaf node
    int startIdx = (n / 2) - 1;
    // Perform reverse level order traversal from last
    // non-leaf node and heapify each node
    for (int i = startIdx; i >= 0; i--) {
        heapify(arr, n, i);
    }
}
```

```
int main() {
    int arr[] = { 1, 3, 5, 4, 6, 13, 10, 9, 8, 15, 17 };
    int n = sizeof(arr) / sizeof(arr[0]);
    buildHeap(arr, n);
    printHeap(arr, n);
    return 0;
}
```

HEAP SORT



Heap Sort

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- The Heap sort algorithm to arrange a list of elements in ascending order is performed using the following step,
 1. Construct a **Binary Tree**.
 2. Transform the Binary Tree into **Max Heap**.
 3. Delete the root element from Max Heap using the **Heapify** method.
 4. Put the deleted element into the Sorted list.
 5. Repeat the same until Max Heap becomes empty.
 6. Display the sorted list.

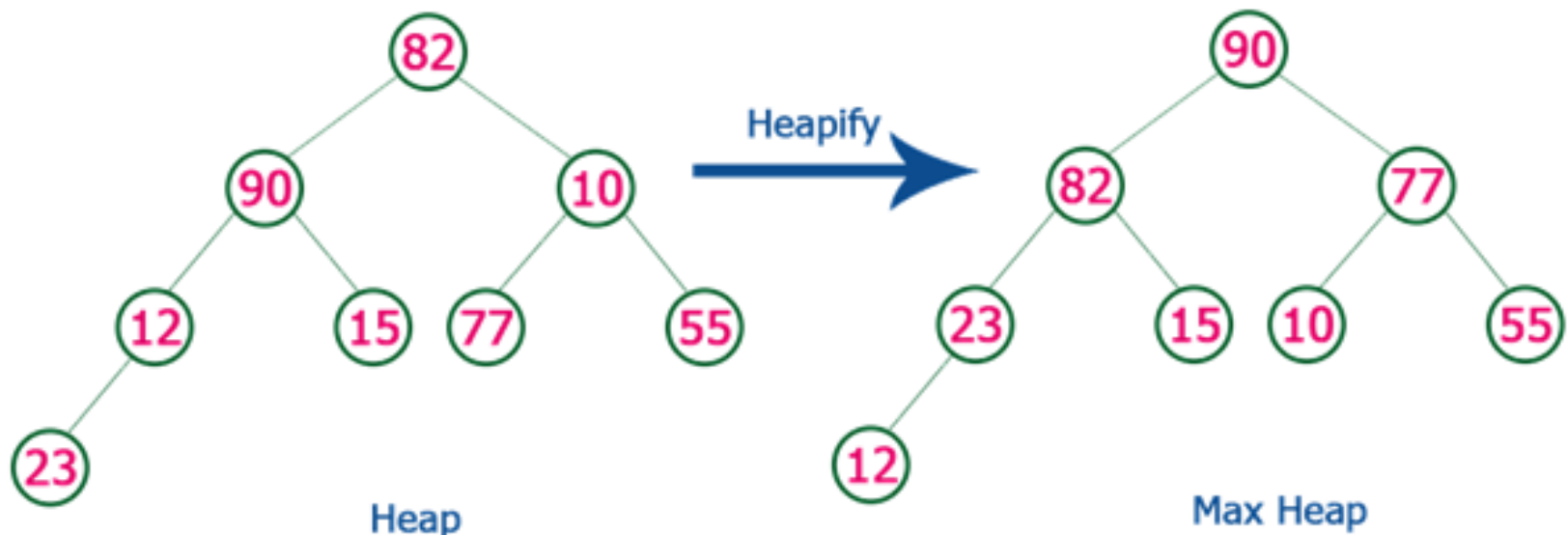
Heap Sort

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Consider the following list of unsorted numbers which are to be sort using Heap Sort

82, 90, 10, 12, 15, 77, 55, 23

Step 1 - Construct a Heap with given list of unsorted numbers and convert to Max Heap



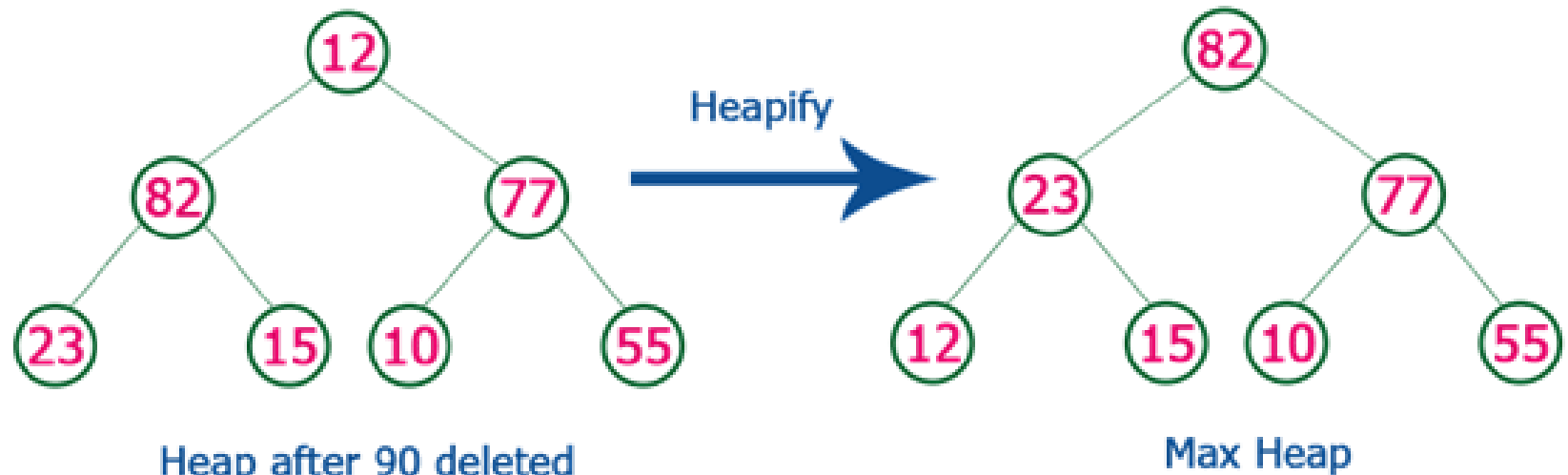
list of numbers after heap converted to Max Heap

90, 82, 77, 23, 15, 10, 55, 12

Heap Sort

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Step 2 - Delete root (**90**) from the Max Heap. To delete root node it needs to be swapped with last node (**12**). After delete tree needs to be heapify to make it Max Heap.



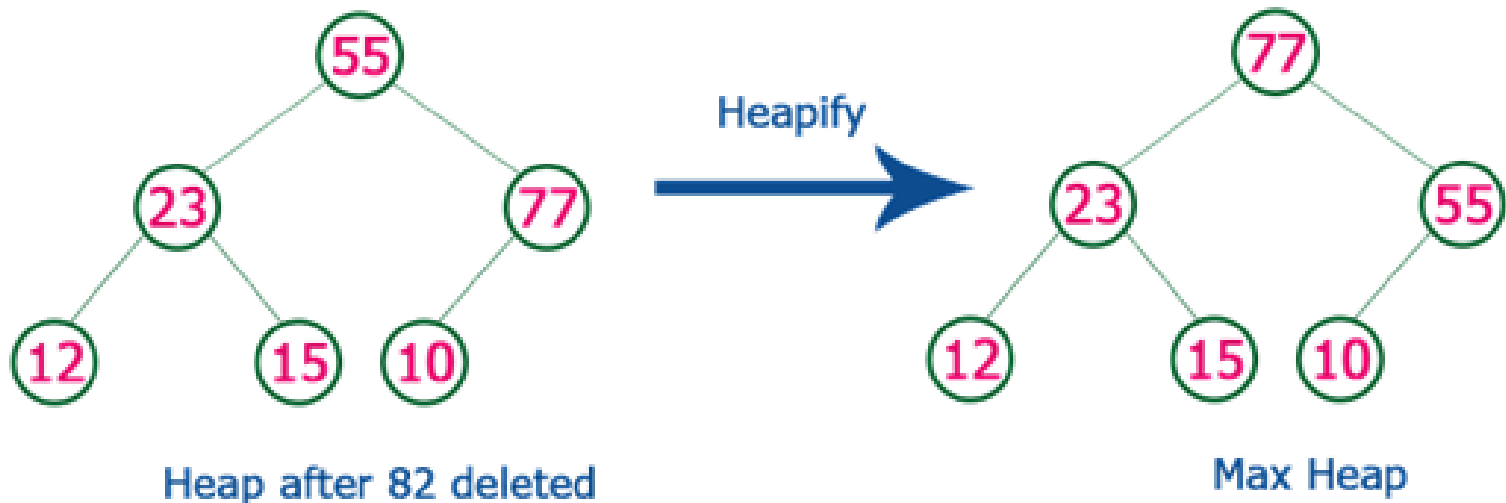
list of numbers after swapping 90 with 12.

12, 82, 77, 23, 15, 10, 55, 90

Heap Sort

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Step 3 - Delete root (**82**) from the Max Heap. To delete root node it needs to be swapped with last node (**55**). After delete tree needs to be heapify to make it Max Heap.



list of numbers after swapping 82 with 55.

12, 55, 77, 23, 15, 10, 82, 90

Heap Sort

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Step 4 - Delete root (**77**) from the Max Heap. To delete root node it needs to be swapped with last node (**10**). After delete tree needs to be heapify to make it Max Heap.



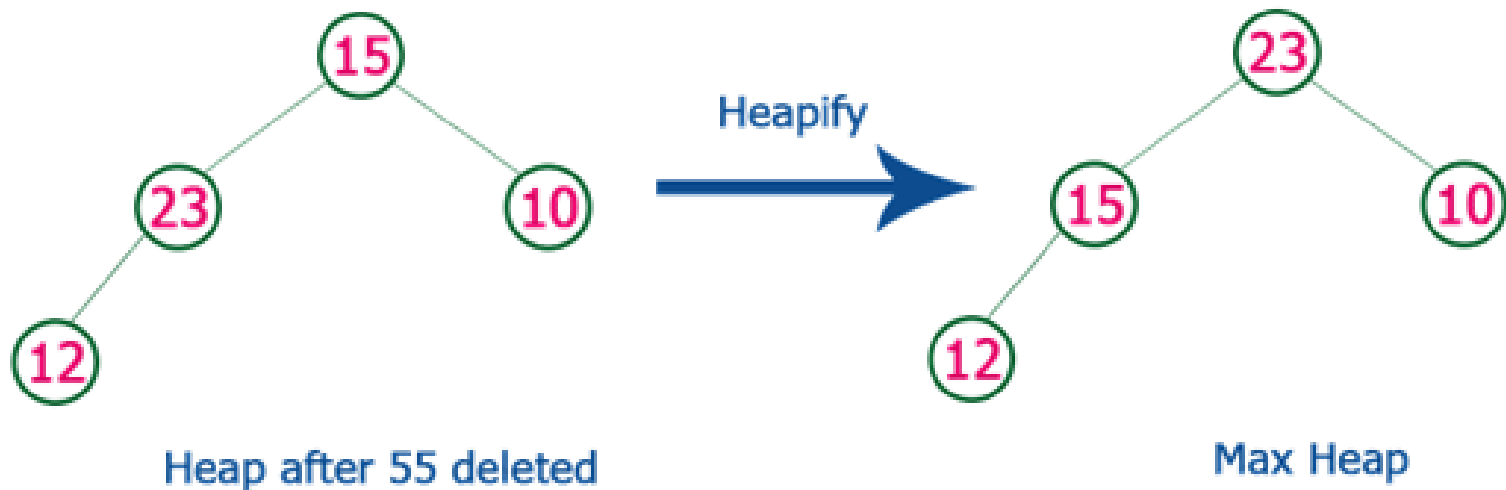
list of numbers after swapping 77 with 10.

12, 55, 10, 23, 15, 77, 82, 90

Heap Sort

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Step 5 - Delete root (**55**) from the Max Heap. To delete root node it needs to be swapped with last node (**15**). After delete tree needs to be heapify to make it Max Heap.



list of numbers after swapping 55 with 15.

12, 15, 10, 23, 55, 77, 82, 90

Heap Sort

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Step 6 - Delete root (**23**) from the Max Heap. To delete root node it needs to be swapped with last node (**12**). After delete tree needs to be heapify to make it Max Heap.



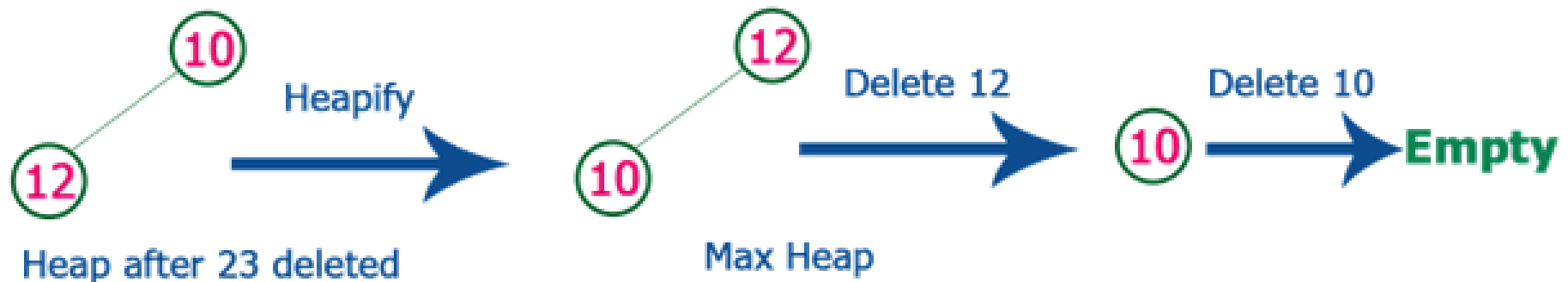
list of numbers after swapping 23 with 12.

12, 15, 10, 23, 55, 77, 82, 90

Heap Sort

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Step 7 - Delete root (**15**) from the Max Heap. To delete root node it needs to be swapped with last node (**10**). After delete tree needs to be heapify to make it Max Heap.



list of numbers after Deleting 15, 12 & 10 from the Max Heap.

10, 12, 15, 23, 55, 77, 82, 90

Whenever Max Heap becomes Empty, the list get sorted in Ascending order

Reading Materials

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- Schaum's Outlines: Chapter # 7
- D. S. Malik: Chapter # 11
- Nell Dale: Chapter # 8

