

## Assignment 5 – CS 301 - Theory of Automata – Fall 2020

Due: Thursday, December 3, 2020 (online – 8 AM)

If any assignment is deemed to be copied from any other student or internet, you may be awarded an F in this course.

1. [50 = 10 \* 5]

For the following grammars, determine whether the language is a CFL or non-CFL. If the language is CFL, you are required to give a CFG. Otherwise, use pumping lemma to show that the language is a non-CFL language.

- a.  $A = \{a^n ww^R a^n \mid n > 0 \text{ and } w \in \{a, b\}^*\}$  Note that  $w^R$  stands for reverse of word  $w$
- b.  $B = \{w \mid w \in \{a, b\}^* \text{ and } w \text{ is not a palindrome}\}$
- c.  $C = \{w \mid w \in \{a, b, c\}^* \text{ where } na(w) = 2nb(w) \text{ and } n(w) = na(w) + nb(w)\}$  Note:  $na(w)$  stands for number of a's in  $w$ .
- d.  $D = \{a^n b^m c^k \mid n \geq m \text{ and } k = n - m\}$
- e.  $E = \{a^{2n} b^{2m} c^n d^m \mid n, m \geq 0\}$
- f.  $F = \{a^n b^n c^k \mid n \leq k \leq 2n\}$

Part a and b from <http://www.cs.williams.edu/~tom/courses/361/notes/Lect17.pdf>

Part c from final fall 2018

Part d from ...



HW5Sol.pdf

Part e from lums

Part f from lums HW8 SOL.pdf

2. [15] [https://courses.engr.illinois.edu/cs373/fa2009/hws/hw\\_04.pdf](https://courses.engr.illinois.edu/cs373/fa2009/hws/hw_04.pdf)

In class we have seen that  $a^n b^n c^n$  is not a CFL as we have just one stack and 2 comparisons cannot be done with one stack only. Assume that we have a special kind of PDA which has two stacks. Prove that a PDA with two stacks will be able to accept some non-CFL languages as well using  $a^n b^n c^n$  as an example. Write down in steps how would the new PDA accept this language.

List another example that can be accepted by 2-stack PDA but not by one-stack PDA

How about PDAs with queue? Can it also accept our example language? Show complete working

3. [10]

In class, we have seen CNF form of CFG, according to which all rules must be of the following form

$S \rightarrow BC$

$S \rightarrow a$

Assume new-CNF of the following format

$S \rightarrow ABC$

$S \rightarrow a$

All other CNF rules are still valid. For example, start variable symbol cannot appear on RHS of the rule.

Will it be possible to find grammar for every CFL?

**Counter example  $A \rightarrow aa$  is a CFG but cannot be converted into the desired form. Infact any string with length two cannot be converted into the form given in question.**

4.

Given two NPDA's, is it possible to verify whether both NPDA's accept the same language. If yes then how? If no then why?

Hint: consider how would someone check if two sets are equal or not.

Also, consider the properties of CFL to do this question.