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General Guidelines

Pumping Lemma is more of a 'proof', therefore, proper steps and assumptions ARE necessary. If you end up missing proper statements, or have a confusing proof, it will obviously lead to deduction of marks (a zero).

You can't write a half-hearted incomplete inconclusive proof and expect the reader to just agree to your claims. You have to show them! Your proof shall be CONCRETE IRREFUTABLE piece of evidence that the reader can't find holes in. (I will be the reader)

I am again letting you know, marking will indeed be tight.

Make sure you study it carefully and properly. There are some clues and hints in the solutions I solved that will give you an idea of how close cut checking I will be doing. So keep a sharp eye!

This preparation will help you in your Mid-2 as well as understanding Pumping Lemma 2 which will come in Finals.

You have been notified of the binary checking well before time. So prepare well!

Best of Luck!

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$L = \{x \mid x \in \Sigma^* \text{ and } x=a^ib^j \text{ where } i\neq j \text{ and } 2i\neq j\}$

Let there be a Language $\alpha = a^pb^{3p}$ which has a pumping length "p" & has a DFA with p states & it fulfills the criteria of Language, L, described above ∵ $p\neq 3p$ & $2p\neq 3p$

Assume α is regular and that $\alpha \in L$

Step I) Prove $|\alpha| \geq p \rightarrow p + 3p \geq p \rightarrow 4p \geq p$ ∴ proved

Step II) Now divide α into 3 substrings u v w so let's take it as follows:

$uv= a^p \quad w= b^{3p}$

- Now further divide u, v and w to get:

$u = \wedge \quad v = a^p \quad w = b^{3p}$

  - Fulfills case $|v| > 0 \rightarrow p > 0$ ∴ proved
  - Fulfills case $|uv| \leq p \rightarrow p \leq p$ ∴ proved

Now pump it as uv^xw // NOTE: we didn't use uv^iw... why? ;) [1]

Step III) For all $x \geq 0$, $\alpha \in L$

$uv^xw = (\wedge)(a^p)^x(b^{3p})$

Let x = 3

$uv^3w = (\wedge)(a^p)^3(b^{3p})$

For x=3; $\{a^{3p}b^{3p}\} \notin L$ as it violates the condition $i\neq j$

→ Thus, this contradicts our assumption

→We assumed α was regular but it is proved to be non-regular

→ Hence this L is non-regular

Note: ∵ symbol means "because", ∴ symbol means "therefore"

Note: we can use either of $|v| > 0$, $|v| \geq 1$ or $|v| \neq 0$. Keep an eye out in future examples!

[1] We can't use variable 'i' as it's already being used in the description of L i.e. a^ib^j etc.

$$\text{Funny\_Maths} = \left\{1^{(n+m)!}0^{n!} \mid n,\ m \in N\right\}$$

Let there be a Language β = $1^{2e!}0^{e!}$ which has a pumping length "e" & has a DFA with e states & it fulfills the criteria of Language, Funny_Maths, described above ∵ letting n=m // NOTE: the name of language used

Assume β is regular and that β ∈ Funny_Maths

Step I) Prove |β| ≥ e → 2e! + e! ≥ e ∴ proved

Step II) Now divide β into 3 substrings x y z so let's take it as follows:

xy= 1 z= $1^{2e!-1}0^{e!}$

- Now further divide x, y and z to get:

  x = ^ y = 1 z = $1^{2e!-1}0^{e!}$

  - Fulfills case |y| ≠ 0 → 1 ≠ 0 ∴ proved
  - Fulfills case |xy| ≤ e → 1 ≤ e ∴ proved

  Now pump it as xy^fz

Step III) For all f ≥ 0, β ∈ Funny_Maths

$xy^fz = (\text{^})(1)^f(1^{2e!-1}0^{e!})$

Let f = 33

$xy^{33}z = (\text{^})(1)^{33}(1^{2e!-1}0^{e!})$

For f=33; $\{1^{32+2e!}0^{e!}\}$ ∉ Funny_Maths as it violates the condition of power of 0 being e! whilst the power of 1 is twice factorial of power of 0// NOTE: Power of 1 [2]

→ Thus, this contradicts our assumption

→We assumed β was regular but it is proved to be non-regular

→ Hence Funny_Maths is non-regular

[2] Power of 1 is (32+2e!) because it was simplified from (33+2e!-1)

QWERTY = {www | w ∈ (01)*}

Let there be a Language x = $0^n10^n10^n1$ which has a pumping length "n" & has a DFA with n states & it fulfills the criteria of Language, QWERTY, described above ∵ w=0^n1

Assume x is regular and that x ∈ QWERTY

Step I) Prove $|x| \geq n \rightarrow (n+1)3 \geq n \rightarrow 3n+3 \geq n$ ∴ proved

Step II) Now divide x into 3 substrings a b c so let's take it as follows:

ab= 0^n c= 10^n10^n1

- Now further divide a, b and c to get:

a = ^ b = 0^n c = 10^n10^n1

  - Fulfills case $|b| \geq 1 \rightarrow n \geq 1$ ∴ proved
  - Fulfills case $|ab| \leq n \rightarrow n \leq n$ ∴ proved

Now pump it as ab^tc

Step III) For all $t \geq 0$, x ∈ QWERTY

ab^tc = (^)$(0^n)^t(10^n10^n1)$

Let t = 50

$ab^{50}c$ = (^)$(0^n)^{50}(10^n10^n1)$

For t=50; {$0^{50n}10^n10^n1$} ∉ QWERTY as it violates the condition of 'n' 0s separated by one 1

⟶ Thus, this contradicts our assumption

⟶We assumed x was regular but it is proved to be non-regular

⟶ Hence QWERTY is non-regular // NOTE: the name of language used is QWERTY not x

My_Automata_Lang = $\{a^i b^j a^k \mid k > i+j\}$

Let there be a Language $\Omega = a^c b^c a^{3c}$ which has a pumping length "c" & has a DFA with c states & it fulfills the criteria of Language, My_Automata_Lang, described above $\because 3c > c + c \rightarrow 3c > 2c$

Assume Ω is regular and that $\Omega \in$ My_Automata_Lang

Step I) Prove $|\Omega| \geq c \rightarrow c+c+3c \geq c \rightarrow 5c \geq c$ $\therefore$ proved

Step II) Now divide Ω into 3 substrings μ δ λ so let's take it as follows: // NOTE: variable names[3]

$$\mu\delta = a^c \qquad \lambda = b^c a^{3c}$$

- Now further divide μ, δ and λ to get:

$$\mu = \wedge \quad \delta = a^c \qquad \lambda = b^c a^{3c}$$

  - Fulfills case $|\delta| \geq 1 \rightarrow c \geq 1$ $\therefore$ proved
  - Fulfills case $|\mu\delta| \leq c \rightarrow c \leq c$ $\therefore$ proved

Now pump it as $\mu\delta^q\lambda$

Step III) For all $q \geq 0$, $\Omega \in$ My_Automata_Lang

$$\mu\delta^q\lambda = (\wedge)(a^c)^q(b^c a^{3c})$$

Let q = 100

$$\mu\delta^{100}\lambda = (\wedge)(a^c)^{100}(b^c a^{3c})$$

For q=100; $\{a^{100c} b^c a^{3c}\} \notin$ My_Automata_Lang as it violates the condition $3c > 100c + c$

$\rightarrow$ Thus, this contradicts our assumption

$\rightarrow$ We assumed Ω was regular but it is proved to be non-regular

$\rightarrow$ Hence My_Automata_Lang is non-regular

[3] Notice the use of μ, δ and λ variables instead of our usual alphabetic variables. This 'twist' is done intentionally just so that you can solidify your conceptual learning and understanding. This is not meant to intimidate you, it is just to make you open your vision to perceive things from different angles. Take some time to understand these examples and wrap your head around this. Be fearless when using variable names! I would love to check your quizzes or assignments with such interesting 'twist(s)'. **BUT** make sure you don't make any foolish mistakes!

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DISCLAIMER

The Pumping Lemma 1 questions below are solved ones that I found over the internet. **For your quizzes, assignments, Mids and Finals; You will follow the proof steps as outlined in the above examples NOT the ones below.** I do not own any of the questions or their solutions below. They are simply here as reference and as a guidance to you as to how you can use Pumping Lemma 1 by thinking up of faulty cases. You can also use the question statements for your practice and check the solution here to see if your logic was correct or not.

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1. L= $\{a^k b^k \mid k \geq 0\}$

   See Notes or solve on your own… its simple really…

2. L = $\{a^k \mid k \text{ is a prime}\}$

   Proof by contradiction:

   Let us assume L is regular. Clearly L is infinite (there are infinitely many prime numbers). From the pumping lemma, there exists a number n such that any string w of length greater than n has a "repeatable" substring generating more strings in the language L. Let us consider the first prime number $p \geq n$. For example, if n was 50 we could use p = 53. From the pumping lemma the string of length p has a "repeatable" substring. We will assume that this substring is of length $k \geq 1$. Hence:

   $a^p \in L$ and

   $a^{p+pk} \in L$ as well as

   $a^{p+2k} \in L$, etc.

   It should be relatively clear that p + k, p + 2k, etc., cannot all be prime but let us add k p times, then we must have:

   $$a^{p+pk} \in L\text{, of course } a^{p+pk} = a^{p(k+1)}$$

so this would imply that $(k + 1)p$ is prime, which it is not since it is divisible by both p and $k + 1$.

Hence L is not regular.

3. $L = \{a^n b^{n+1}\}$

Assume L is regular. From the pumping lemma there exists a p such that every w 0 L such that $|w| \geq p$ can be represented as x y z with $|y| \neq 0$ and $|xy| \leq p$. Let us choose a b . Its length is $2p + 1 \geq p$. Since the length of xy cannot exceed p, y p p+1 must be of the form a for some $k > 0$. From the pumping lemma a b must also k p-k p+1 be in L but it is not of the right form. Hence the language is not regular.

Note that the repeatable string needs to appear in the first n symbols to avoid the following situation:

assume, for the sake of argument that $n = 20$ and you choose the string $a^{10} b^{11}$ which is of length larger than 20, but $|xy| \leq 20$ allows xy to extend past b, which means that y could contain some b's. In such a case, removing y (or adding more y's) could lead to strings which still belong to L.

4. $L = \{a^n b^{2n}\}$

Assume L is regular. From the pumping lemma there exists a p such that every $w \in L$ such that $|w| \geq p$ can be represented as x y z with $|y| \neq 0$ and $|xy| \leq p$. Let us choose $a^p b^{2p}$. Its length is $3p \geq p$. Since the length of xy cannot exceed p, y must be of the form a^k for some $k > 0$. From the pumping lemma $a^{p-k} b^{2p}$ must also be in L

but it is not of the right form. Hence the language is not regular.

5. TRAILING-COUNT as any string s followed by a number of a's equal to the length of s.

Assume L is regular. From the pumping lemma there exists a p such that every $w \in L$ such that $|w| \geq p$ can be represented as x y z with $|y| \neq 0$ and $|xy| \leq p$. Let us choose $b^p a^p$. Its length is $2p \geq p$. Since the length of xy cannot exceed p, y must be of the form b^k for some $k > 0$. From the pumping lemma $b^{p-k} a^p$ must also be in L but it is not of the right form. Hence the language is not regular.

6. EVENPALINDROME = { all words in PALINDROME that have even length}

Same as #2 above, choose $a^n bba^n$.

7. ODDPALINDROME = { all words in PALINDROME that have odd length}

Same as #2 above, choose $a^n ba^n$.

8. DOUBLESQUARE = { $a^n b^n$ where n is a square }

Assume DOUBLESQUARE is regular. From the pumping lemma there exists a p such that every w ∊ L such that |w| ≥ p can be represented as x y z with |y| ≠ 0 and |xy| ≤ p. Let us choose $a^{p*p}b^{p*p}$. Its length is $2p^2 \geq$ p. Since the length of xy cannot exceed p, y must be of the form a^k for some k > 0. Let us add y p times. From the pumping lemma $a^{p*p+pk}b^{p*p} = a^{p(p+k)}b^{p*p}$ must also be in L but it is not of the right form. Hence the language is not regular.

9. L = { w | w ∊ {a, b}*, w = w^R}

Proof by contradiction:
Assume L is regular. Then the pumping lemma applies.
From the pumping lemma there exists an n such that every w ∊ L longer than n can be represented as x y z with |y| ≠ 0 and |x y| ≤ n.

Let us choose the palindrome $a^n b a^n$.

Again notice that we were clever enough to choose a string which:

a. *has a center mark which is not a (otherwise when we remove or add y we would be left with an acceptable string)*
b. *has a first portion on length n which is all a's (so that when we remove or add y it will create an imbalance).*

Its length is 2n + 1 ≥ n. Since the length of xy cannot exceed n, y must be of the form a^k for some k > 0. From the pumping lemma a^{n-k} b a^n must also be in L but it is not a palindrome.

Hence L is not regular.

10. L = { w ∊ {a, b}* | w has an equal number of a's and b's}

Let us show this by contradiction: assume L is regular. We know that the language generated by a*b* is regular. We also know that the intersection of two regular languages is regular. Let M = ($a^n b^n$ | n ≥ 0} = L(a*b*) ∩ L. Therefore if L is regular M would also be regular. but we know tha M is not regular. Hence, L is not regular.

11.L={ w w^R | w ∊ {a, b}* }
see # 7

12. L = { 0^n | n is a power of 2 }

Assume L is regular. From the pumping lemma there exists a p such that every w ∊ L such that |w| ≥ p can be represented as x y z with |y| ≠ 0 and |xy| ≤ p. Let us choose

$n = 2^p$. Since the length of xy cannot exceed p, y must be of the form 0^k for some $0 < k \leq p$. From the pumping lemma 0^m where $m = 2^p + k$ must also be in L. We have

$$2^p < 2^p + k \leq 2^p + p < 2^{p+1}$$

Hence this string is not of the right form. Hence the language is not regular.

13. $L = \{a^{2k}w \mid w \in \{a, b\}^*, |w| = k\}$

Assume L is regular. From the pumping lemma there exists a p such that every $w \in L$ such that $|w| \geq p$ can be represented as x y z with $|y| \neq 0$ and $|xy| \leq p$. Let us choose $a^{2p}b^p$. Its length is $3p \geq p$. Since the length of xy cannot exceed p, y must be of the form a^k for some $k > 0$. From the pumping lemma $a^{2p-k}b^p$ must also be in L but it is not of the right form since the number of a's cannot be twice the number of b's (Note that you must subtract not add , otherwise some a's could be shifted into w). Hence the language is not regular.

14. $L = \{a^k w \mid w \in \{a, b\}^*, |w| = k\}$

Assume L is regular. From the pumping lemma there exists a p such that every $w \in L$ such that $|w| \geq p$ can be represented as x y z with $|y| \neq 0$ and $|xy| \leq p$. Let us choose $a^p b^p$. Its length is $2p \geq p$. Since the length of xy cannot exceed p, y must be of the form a^k for some $k > 0$. From the pumping lemma $a^{p-k}b^p$ must also be in L but it is not of the right form since the number of a's cannot be equal to the number of b's (Note that you must subtract not add , otherwise some a's could be shifted into w). Hence the language is not regular.

15. $L = \{a^n b^l \mid n \leq l\}$

Assume L is regular. From the pumping lemma there exists a p such that every $w \in L$ such that $|w| \geq p$ can be represented as x y z with $|y| \neq 0$ and $|xy| \leq p$. Let us choose $a^p b^p$. Its length is $2p \geq p$. Since the length of xy cannot exceed p, y must be of the form a^k for some $k > 0$. From the pumping lemma $a^{p+k} b^p$ must also be in L but it is not of the right form since the number of a's exceeds the number of b's (Note that you must add not subtract, otherwise the string would be OK). Hence the language is not regular.

16. $L = \{a^n b^l a^k \mid k = n + l\}$

Assume L is regular. From the pumping lemma there exists a p such that every $w \in L$ such that $|w| \geq p$ can be represented as x y z with $|y| \neq 0$ and $|xy| \leq p$. Let us choose $a^p b a^{p+1}$. Its length is $2p+2 \geq p$. Since the length of xy cannot exceed p, y must be of the form a^m for some $m > 0$. From the pumping lemma $a^{p-m}ba^{p+1}$ must also be in L but it is not of the right form. Hence the language is not regular.

17. $L = \{va^{k+1} \mid v \in \{a, b\}^*, |v| = k\}$

Assume L is regular. From the pumping lemma there exists an n such that every $w \in L$ such that $|w| \geq n$ can be represented as x y z with $|y| \neq 0$ and $|xy| \leq p$. Let us choose $b^n a^{n+1}$. Its length is $2n+1 \geq n$. Since the length of xy cannot exceed n, y must be of the form b^k for some $k > 0$. From the pumping lemma if we add two y to the original string $b^{n+2k}a^{n+1}$ must also be in L but that string is of length $2n+2k+1$ and v would have to be b^{n+k} to fit the pattern the rest of the string would then be $b^k a^{k+1}$ which is not of the right form. Hence the language is not regular.

18. $L = \{va^{2k} \mid v \in \{a, b\}^*, |v| = k\}$

Assume L is regular. From the pumping lemma there exists a n such that every $w \in L$ such that $|w| \geq n$ can be represented as x y z with $|y| \neq 0$ and $|xy| \leq n$. Let us choose $b^n a^{2n}$. Its length is $3n \geq n$. Since the length of xy cannot exceed n, y must be of the form b^k for some $k > 0$. From the pumping lemma $b^{n+k} a^n$ must also be in L but it is not of the right form since the number of a's exceeds the number of b's and we cannot move any b's on the a side (Note that you must add not subtract, otherwise the string would be OK by shifting a's to the b side). Hence the language is not regular.

19. $L = \{ww \mid w \in \{a, b\}^*\}$

Assume L is regular. From the pumping lemma there exists an n such that every $w \in L$ such that $|w| \geq n$ can be represented as x y z with $|y| \neq 0$ and $|xy| \leq n$. Let us choose $a^n b^n a^n b^n$. Its length is $4n \geq n$. Since the length of xy cannot exceed n, y must be of the form a^k for some $k > 0$. From the pumping lemma $a^{n+k}b^n a^n b^n$ must also be in L but it is not of the right form since the middle of the string would be in the middle of the b which prevents a match with the beginning of the string. Hence the language is not regular.

20. $L = \{ a^{n!} \mid n \geq 0\}$

Proof by contradiction:

Let us assume L is regular. From the pumping lemma, there exists a number p such that any string w of length greater than p has a "repeatable" substring generating more strings in the language L. Let us consider $a^{p!}$ (unless $p < 3$ in which case we chose $a^{3!}$). From the pumping lemma the string w has a "repeatable" substring. We will assume that this substring is of length $k \geq 1$.

From the pumping lemma $a^{p!-k}$ must also be in L. For this to be true there must be j such that $j! = m! - k$ But this is not possible since when $p > 2$ and $k \leq m$ we have $m! - k > (m - 1)!$

Hence L is not regular.

21. $L = \{ a^n b^l \mid n \neq l\}$

Proof by contradiction:

Let us assume L is regular. From the pumping lemma, there exists a number p such that any string w of length greater than p has a “repeatable” substring generating more strings in the language L. Let us consider n = p! and l = (p+1)! From the pumping lemma the resulting string is of length larger than p and has a “repeatable” substring. We will assume that this substring is of length $k \geq 1$.

From the pumping lemma we can add y i-1 times for a total of i ys. If we can find an i such that the resulting number of a’s is the same as the number of b’s we have won. This means we must find i such that:

$m! + (i - 1)*k = (m + 1)!$ or

$(i - 1) k = (m + 1) m! - m! = m * m!$ or

$i = (m * m!) / k + 1$

but since k < m we know that k must divide m! and that (m * m!) / k must be an integer. This proves that we can choose i to obtain the above equality.

Hence L is not regular.

22. $L = \{a^n b^l a^k \mid k > n + l\}$

Assume L is regular. From the pumping lemma there exists a p such that every $w \in L$ such that $|w| \geq p$ can be represented as x y z with $|y| \neq 0$ and $|xy| \leq p$. Let us choose $a^p b a^{p+2}$. Its length is $2p+3 \geq p$. Since the length of xy cannot exceed p, y must be of the form a^m for some m > 0. From the pumping lemma $a^{p+2m} b a^{p+2}$ must also be in L but it is not of the right form since p+2m+1 > p+2. Hence the language is not regular.

23. $L = \{a^n b^l c^k \mid k \neq n + l\}$

Assume L is regular. From the pumping lemma there exists a p such that every $w \in L$ such that $|w| \geq p$ can be represented as x y z with $|y| \neq 0$ and $|xy| \leq p$. Let us choose $a^{p!} b^{p!} a^{(p+1)!}$. Its length is $2p!+(p+1)! \geq p$. Since the length of xy cannot exceed p, y must be of the form a^m for some m > 0. From the pumping lemma any string of the form $xy^i z$ must always be in L. If we can show that it is always possible to choose i in such a way that we will have k = n + l for one such string we will have shown a contradiction. Indeed we can have

$$p!+(i-1)m + p! = (p+1)!$$

if we have $i = 1 + ((p+1)! - 2 p!)/ m$ Is that possible? only if m divides

((p+1)! -2 p!

$((p + 1)! - 2 * (p)! = (p + 1 - 2)\ p!$ and since $m \leq p$ m is guaranteed to divide $p!$. Hence

i exists and the language is not regular.

24. $L = \{a^n b^l a^k \mid n = l \text{ or } l \neq k\}$

Proof by
contradiction:
Let us assume L is regular. From the pumping lemma, there exists a number p such that any string w of length greater than p has a "repeatable" substring generating more strings in the language L. Let us consider $w = a^p b^p a^p$.
From the pumping lemma the string w, of length larger than p has a "repeatable" substring. We will assume that this substring is of length $m \geq 1$. From the pumping lemma we can remove y and the resulting string should be in L. However, if we remove y we get $a^{p - m} b^p a^p$. But this string is not in L since $p-m \neq p$ and $p = p$.

Hence L is not regular.

25. $L = \{a^n b a^{3n} \mid n \geq 0\}$

Assume L is regular. From the pumping lemma there exists a p such that every $w \in L$ such that $|w| \geq p$ can be represented as x y z with $|y| \neq 0$ and $|xy| \leq p$. Let us choose $a^p b a^{3p}$. Its length is $4p+1 \geq p$. Since the length of xy cannot exceed p, y must be of the form a^k for some $k > 0$. From the pumping lemma $a^{p-k} b a^{3p}$ must also be in L but it is not of the right form. Hence the language is not regular.

26. $L = \{a^n b^n c^n \mid n \geq 0\}$

Assume L is regular. From the pumping lemma there exists a p such that every $w \in L$ such that $|w| \geq p$ can be represented as x y z with $|y| \neq 0$ and $|xy| \leq p$. Let us choose $a^p b^p c^p$. Its length is $3p \geq p$. Since the length of xy cannot exceed p, y must be of the form a^k for some $k > 0$. From the pumping lemma $a^{p-k} b^p a^p$ must also be in L but it is not of the right form. Hence the language is not regular.

27. $L = \{a^i b^n \mid i, n \geq 0, i = n \text{ or } i = 2n\}$

Assume L is regular. From the pumping lemma there exists a p such that every $w \in L$ such that $|w| \geq p$ can be represented as x y z with $|y| \neq 0$ and $|xy| \leq p$. Let us choose $a^p b^p$. Its length is $2p \geq p$. Since the length of xy cannot exceed p, y must be of the form a^k for some $k > 0$. From the pumping lemma $a^{p-k} a^p$ must also be in L but it is not of the right form. Hence the language is not regular.

28. $L = \{0^k 1 0^k \mid k \geq 0\}$

Assume L is regular. From the pumping lemma there exists an n such that every $w \in L$ such that $|w| \geq n$ can be represented as x y z with $|y| \neq 0$ and $|xy| \leq n$. Let us choose $0^n 1 0^n$. Its length is $2n+1 \geq n$. Since the length of xy cannot exceed n, y must be of the form 0^p for some $p > 0$. From the pumping lemma $0^{n-p} 1 0^n$ must also be in L but it is not of the right form. Hence the language is not regular.

29. $L = \{0^n 1^m 2^n \mid n, m \geq 0\}$

Assume L is regular. From the pumping lemma there exists a p such that every $w \in L$ such that $|w| \geq p$ can be represented as x y z with $|y| \neq 0$ and $|xy| \leq p$. Let us choose $0^p 1 2^p$. Its length is $2p+1 \geq p$. Since the length of xy cannot exceed p, y must be of the form 0^p for some $p > 0$. From the pumping lemma $0^{n-p} 1 2^n$ must also be in L but it is not of the right form. Hence the language is not regular.