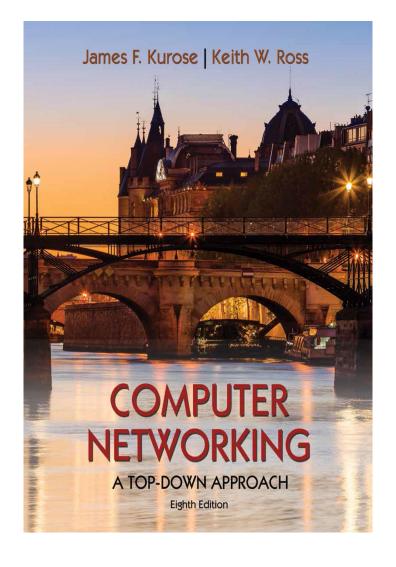
# Network Layer: Control Plane

- introduction
- routing algorithms
  - link state
  - distance vector
- intra-ISP routing: OSPF
- routing among ISPs: BGP
- SDN control plane
- Internet Control Message Protocol
- Network management, configuration

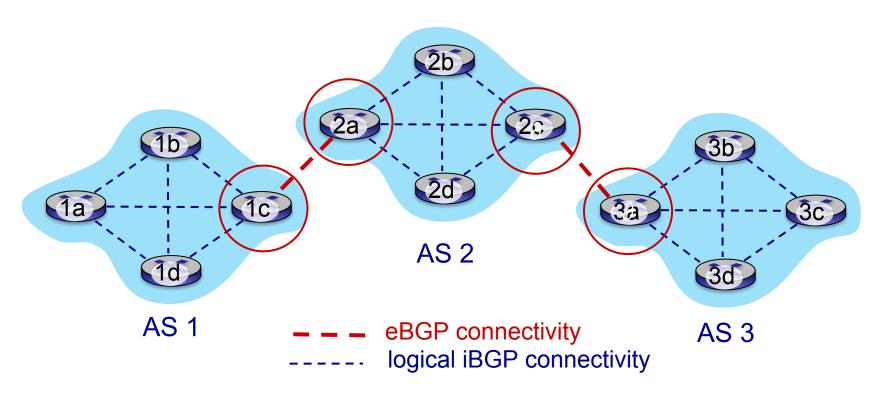
#### Computer Networks



### Internet inter-AS routing: BGP

- BGP (Border Gateway Protocol): the de facto inter-domain routing protocol
  - "glue that holds the Internet together"
- allows subnet to advertise its existence, and the destinations it can reach, to rest of Internet: "I am here, here is who I can reach, and how"
- BGP provides each AS a means to:
  - obtain destination network reachability info from neighboring ASs (eBGP: BGP connection that spans two ASs)
  - determine routes to other networks based on reachability information and policy
  - propagate reachability information to all AS-internal routers (iBGP: BGP connection/session between routers within the same AS)
  - advertise (to neighboring networks) destination reachability info

#### eBGP, iBGP connections

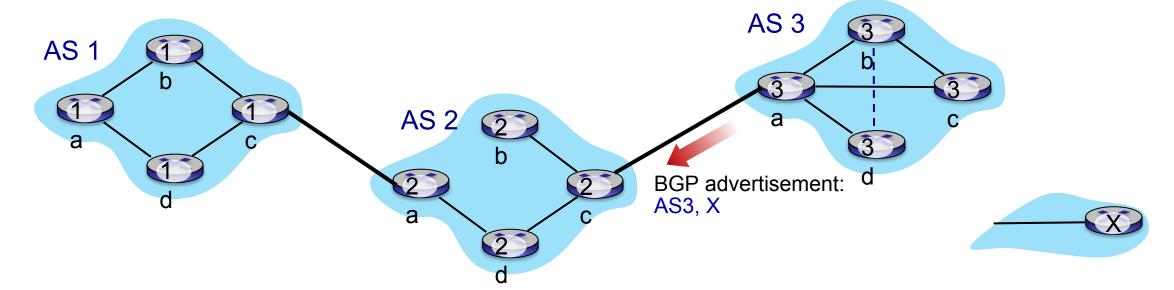




gateway routers run both eBGP and iBGP protocols

#### **BGP** basics

- BGP session: two BGP routers ("peers, speakers") exchange BGP messages over semi-permanent TCP connection:
  - advertising paths to different destination network prefixes (e.g., to a destination /16 network)
  - BGP is a "path vector" protocol
- when AS3 gateway 3a advertises path AS3,X to AS2 gateway 2c:
  - AS3 promises to AS2 it will forward datagrams towards X



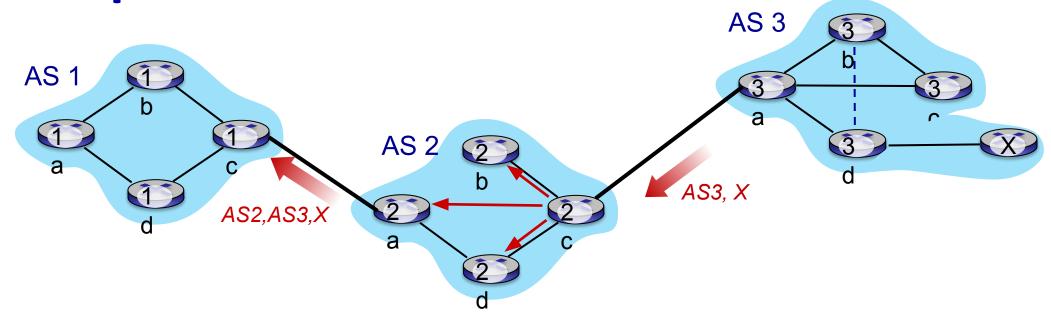
### **BGP** protocol messages

- BGP messages exchanged between peers over TCP connection
- BGP messages [RFC 4371]:
  - OPEN: opens TCP connection to remote BGP peer and authenticates sending BGP peer
  - UPDATE: advertises new path (or withdraws old)
  - KEEPALIVE: keeps connection alive in absence of UPDATES; also ACKs OPEN request
  - NOTIFICATION: reports errors in previous msg; also used to close connection

#### Path attributes and BGP routes

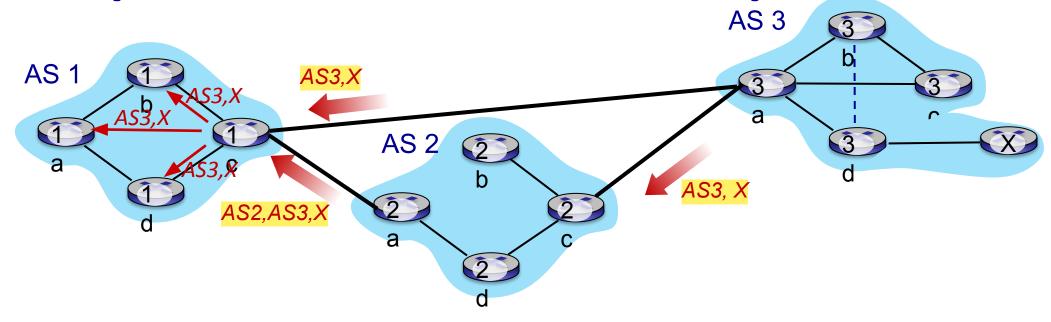
- BGP advertised path: prefix + attributes
  - path prefix: destination being advertised
  - two important attributes:
    - AS-PATH: list of ASes through which prefix advertisement has passed
    - NEXT-HOP: indicates specific internal-AS router to next-hop AS
- policy-based routing:
  - router receiving route advertisement to destination X uses *policy* to accept/reject a path (e.g., never route through AS W, or country Y).
  - router uses policy to decide whether to advertise a path to neighboring AS Z (does router want to route traffic forwarded from Z destined to X?)

#### **BGP** path advertisement



- AS2 router 2c receives path advertisement AS3,X (via eBGP) from AS3 router 3a
- based on AS2 policy, AS2 router 2c accepts path AS3,X, propagates (via iBGP) to all AS2 routers
- based on AS2 policy, AS2 router 2a advertises (via eBGP) path AS2, AS3, X to AS1 router 1c

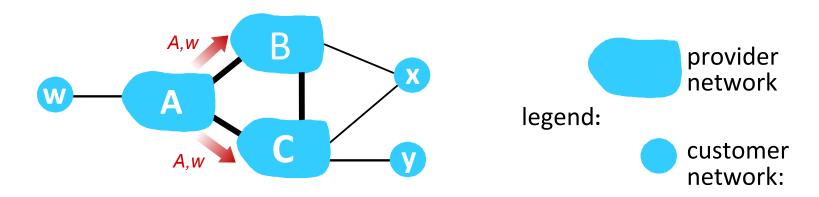
### BGP path advertisement: multiple paths



gateway routers may learn about multiple paths to destination:

- AS1 gateway router 1c learns path AS2,AS3,X from 2a
- AS1 gateway router 1c learns path AS3,X from 3a
- based on policy, AS1 gateway router 1c chooses path AS3,X and advertises path within AS1 via iBGP

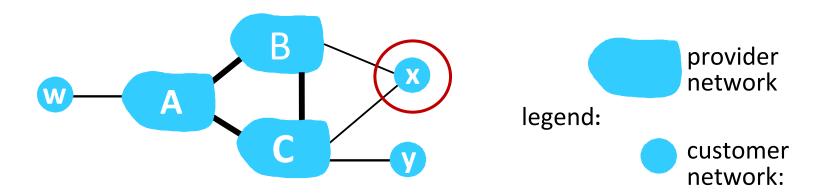
#### **BGP:** achieving policy via advertisements



ISP only wants to route traffic to/from its customer networks (does not want to carry transit traffic between other ISPs – a typical "real world" policy)

- A advertises path Aw to B and to C
- B chooses not to advertise BAw to C!
  - B gets no "revenue" for routing CBAw, since none of C, A, w are B's customers
  - C does not learn about CBAw path
- C will route CAw (not using B) to get to w

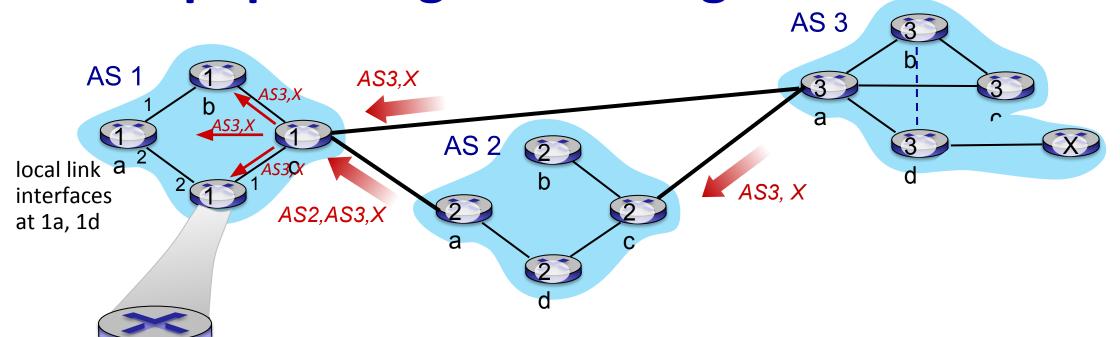
## **BGP:** achieving policy via advertisements (more)



ISP only wants to route traffic to/from its customer networks (does not want to carry transit traffic between other ISPs – a typical "real world" policy)

- A,B,C are provider networks
- x,w,y are customer (of provider networks)
- x is dual-homed: attached to two networks
- policy to enforce: x does not want to route from B to C via x
  - .. so x will not advertise to B a route to C

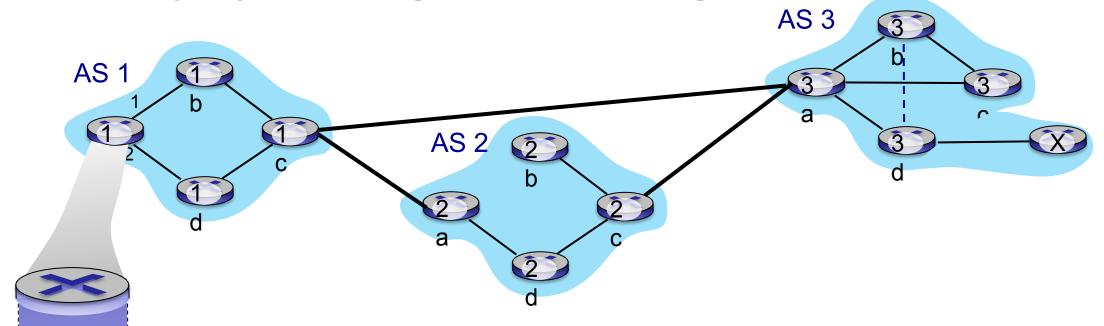
### **BGP:** populating forwarding tables



dest	interface
	•••
1c	1
X	1

- recall: 1a, 1b, 1d learn via iBGP from 1c: "path to X goes through 1c"
- at 1d: OSPF intra-domain routing: to get to 1c, use interface 1
- at 1d: to get to X, use interface 1

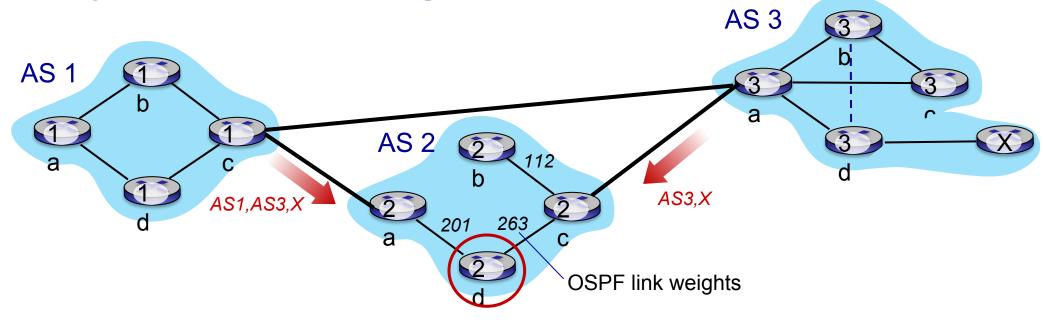
## **BGP:** populating forwarding tables



dest	interface
1c	2
X	2

- recall: 1a, 1b, 1d learn via iBGP from 1c: "path to X goes through 1c"
- at 1d: OSPF intra-domain routing: to get to 1c, use interface 1
- at 1d: to get to X, use interface 1
- at 1a: OSPF intra-domain routing: to get to 1c, use interface 2
- at 1a: to get to X, use interface 2

### Hot potato routing



- 2d learns (via iBGP) it can route to X via 2a or 2c
- hot potato routing: choose local gateway that has least intra-domain cost (e.g., 2d chooses 2a, even though more AS hops to X): don't worry about inter-domain cost!

#### Why different Intra-, Inter-AS routing?

#### policy:

- inter-AS: admin wants control over how its traffic routed, who routes through its network
- intra-AS: single admin, so policy less of an issue

#### scale: reducing forwarding table size, routing update traffic

- hierarchical routing: limiting the scope of full topological information
- BGP routing to CIDRized destination networks (summarized routes)

#### performance:

- intra-AS: can focus on performance
- inter-AS: policy dominates over performance