

# Skewing the Sequent Calculus:

## Skew Monoidal Categories $\wedge$ Skew Multicategories

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# Outline

**1** Skew-Monoidal Categories

**2** Skew Multicategories

**3** Sequent Calculus

# Definition

A **skew-monoidal category** is a category  $\mathcal{C}$  equipped with:

- a bifunctor  $\otimes : \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$
- an object  $I \in \mathcal{C}_0$
- three natural transformations, called associator, left and right unitor respectively. Component-wise written as:

$$\alpha_{x,y,z} : (x \otimes y) \otimes z \rightarrow x \otimes (y \otimes z)$$

$$\lambda_x : I \otimes x \rightarrow x$$

$$\rho_x : x \rightarrow x \otimes I$$

natural in  $x, y, z$ . ! They are not required to be natural isos !

such that the following diagrams commute:

# Definition

$$\begin{array}{c}
 \begin{array}{ccccc}
 & & ((x \otimes y) \otimes z) \otimes w & & \\
 & \swarrow \alpha_{x \otimes y, z, w} & & \searrow \alpha_{x, y, z \otimes w} & \\
 (x \otimes y) \otimes (z \otimes w) & & & & (x \otimes (y \otimes z)) \otimes w \\
 & \searrow \alpha_{x, y, z \otimes w} & & \swarrow \alpha_{x, y \otimes z, w} & \\
 & x \otimes (y \otimes (z \otimes w)) & \xleftarrow{x \otimes \alpha_{y, z, w}} & x \otimes ((y \otimes z) \otimes w) & 
 \end{array} \\
 \\
 \begin{array}{ccc}
 x \otimes y & \xlongequal{\quad} & x \otimes y \\
 \rho_x \otimes y \downarrow & & \uparrow x \otimes \lambda_y \\
 (x \otimes I) \otimes y & \xrightarrow{\alpha_{x, I, y}} & x \otimes (I \otimes y)
 \end{array} \\
 \\
 \begin{array}{ccccc}
 (I \otimes x) \otimes y & \xrightarrow{\alpha_{I, x, y}} & I \otimes (x \otimes y) & (x \otimes y) \otimes I & \xrightarrow{\alpha_{x, y, I}} & x \otimes (y \otimes I) & I & \xlongequal{\quad} & I \\
 \swarrow \lambda_x \otimes y & & \swarrow \lambda_{x \otimes y} & \swarrow \rho_{x \otimes y} & \searrow x \otimes \rho_y & & \searrow \rho_I & & \swarrow \lambda_I \\
 & x \otimes y & & x \otimes y & & & I \otimes I & & 
 \end{array}
 \end{array}$$

! In the case of monoidal categories, the final three equations follow from the first two.

# Examples

- Trivially, any monoidal category is a skew-monoidal category.
- Take  $\mathcal{C}$  to be **Ptd**, the category of pointed sets, with  $(X, x_0) \otimes (Y, y_0) := (X + Y, \text{inl}x_0)$  and  $I := (1, *)$ . Then  $\alpha, \lambda, \rho$  can be easily defined (see next slide).
- A skew-monoidal structure can be put on the poset  $\mathbb{N}$ , as we will see shortly.
- Let  $\mathcal{D}$  be cocomplete so that all left Kan extensions exist, and fix a functor  $J : \mathcal{C} \rightarrow \mathcal{D}$ . We define a tensor product on  $[\mathcal{C}, \mathcal{D}]$  by  $F \otimes G := \text{Lan}_J F \circ G$ . The monoidal unit is  $J$ . The universal property of Kan extensions gives maps  $\alpha, \lambda, \rho$  pointing in the right direction, but not necessarily invertible.

# **Ptd** as Skew Monoidal Category 1/3

Consider  $I := (1, *)$ ,  $(X, x_0)$ ,  $(Y, y_0) \in \mathbf{Ptd}$  pointed sets (or pointed topological spaces !):  
with tensor product

$$(X, x_0) \otimes (Y, y_0) := (X + Y, \text{inl}x_0)$$

Define associator, left and right unitors as follows:

$$\begin{aligned} ((X, x_0) \otimes (Y, y_0)) \otimes (Z, z_0) &= ((X + Y) + Z, \text{inl}(\text{inl}x_0)) \\ &\downarrow \alpha_{X,Y,Z} \\ (X, x_0) \otimes ((Y, y_0) \otimes (Z, z_0)) &= (X + (Y + Z), \text{inl}x_0) \end{aligned}$$

# Ptd as Skew Monoidal Category 2/3

$$\lambda_X : (1, *) \otimes (X, x_0) = (1 + X, \text{inl}*) \rightarrow (X, x_0)$$

$$\text{inl}* \mapsto x_0$$

$$\text{inr}x \mapsto x$$

$$\rho_X : (X, x_0) \rightarrow (1 + X, \text{inl}*)$$

$$x \mapsto \text{inl}x$$

! The left unitor is not monic, the right not epic - but the associator is iso !

# Ptd as Skew Monoidal Category 3/3

What we have just defined look likes a co-product in **Ptd**, but it is not ! - it is indeed the tensor product.

Here it is the actual co-product:

$$(X, x_0) + (Y, y_0) = ((X + Y) / \sim, [\text{inl}x_0])$$

with equivalence relation induced by gluing  $\text{inl}x_0$  and  $\text{inr}y_0$ .



# $\mathbb{N}$ as Posetal SMC

Consider the natural numbers with their standard order as posetal category, meaning

$$n \leq m \iff n \rightarrow m$$

for  $n, m \in \mathbb{N}$ .

For a fixed number  $\mathbf{n} \in \mathbb{N}$ , that we set as unit of the monoidal structure, define the monoidal product truncating as follows:

$$h \otimes r := (h - \mathbf{n}) + r$$

The left, right unitors and associator reads:

$$\lambda_h : (\mathbf{n} - \mathbf{n}) + h = 0 + h = h \quad \rho_h : h \leq h \text{ max } \mathbf{n} = (h - \mathbf{n}) + \mathbf{n}$$

$$\alpha_{h,r,t} : (((h - \mathbf{n}) + r) - \mathbf{n}) + t \leq (h - \mathbf{n}) + (r - \mathbf{n}) + t$$

# Coherence

- When  $\alpha$ ,  $\rho$ , and  $\lambda$  are isomorphisms, we have a **monoidal category**, and the morphisms in these are easily characterised by MacLane's *coherence theorem*.
- This can be stated as follows: the free monoidal category on any set of objects is a preorder.
- In the skew case, this is no longer true. For instance, it is not the case that  $\rho_I \circ \lambda_I = \text{id}_{I \otimes I}$ . It would be nice therefore to have some other way to characterise the morphisms in free skew-monoidal categories.
- The paper solves the issue by characterizing morphisms as equivalence classes of deduction trees for a sequent calculus, as we will see at the very end.

# Multicategories

- The sequent calculus presented is best understood as a calculus for *left representable skew multicategories*. We give first some background on multicategories and their relation to monoidal categories.
- A **multicategory**  $\mathcal{C}$  is like a category, except morphisms can have “multiple inputs”: for each list  $\overline{A}$  of objects of  $\mathcal{C}$ , and for each object  $B$ , we have a Hom-set  $\text{Hom}(\overline{A}; B)$ . Elements of these homsets are often drawn as follows:

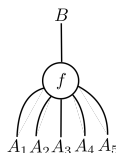
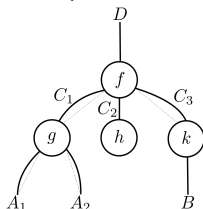


Figure: A multimorphism  $f \in \text{Hom}(A_1, A_2, A_3, A_4, A_5; B)$

# Multicategories

- There is a multi-ary composition operation, and identity morphisms, depicted as:



(a) Graphical representation of  $f(g, h, k)$ .  
 Here  $f \in \text{Hom}(C_1, C_2, C_3; D)$ ,  
 $g \in \text{Hom}(A_1, A_2; C_1)$ ,  $h \in \text{Hom}(\cdot; C_2)$ ,  
 $k \in \text{Hom}(B; C_3)$

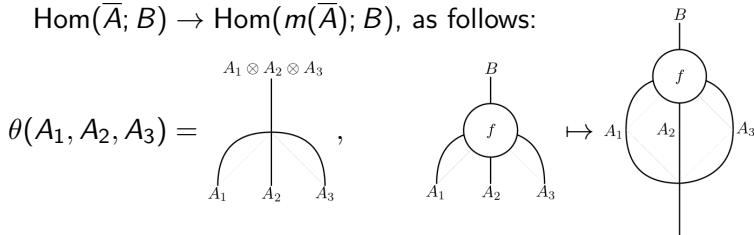


(b) Graphical representation of  $\text{id}_A$

- These are subject to obvious multi-ary associativity and unitality constraints, which follow automatically from the graphical representation.

# Representable Multicategories

- We say a list  $\bar{A}$  of objects is *representable* if there exists an object  $m(\bar{A})$  and a morphism  $\theta(\bar{A}) : \bar{A} \rightarrow m(\bar{A})$  such that precomposing with  $\theta(\bar{A})$  gives a natural isomorphism between  $\text{Hom}(m(\bar{A}), B)$  and  $\text{Hom}(\bar{A}, B)$ , for any  $B$  - that is, the functor  $\text{Hom}(\bar{A}, -)$  is representable.
- We also write  $m(A_1, \dots, A_n)$  as  $(A_1 \otimes \dots \otimes A_n)$ .
- We draw  $\theta$  and the inverse to precomposition, a map  $\text{Hom}(\bar{A}; B) \rightarrow \text{Hom}(m(\bar{A}); B)$ , as follows:



# Representable Multicategories

- A multicategory is **weakly representable** if all lists of objects are, and **representable** if additionally composition with  $(\text{id}_{A_1}, \dots, \text{id}_{A_j}, \theta(\overline{B}), \text{id}_{C_1}, \dots, \text{id}_{C_k})$  induces a bijection between  $\text{Hom}(\overline{A}, m(\overline{B}), \overline{C}; D)$  and  $\text{Hom}(\overline{A}, \overline{B}, \overline{C}; D)$  for all  $\overline{A}, \overline{B}, \overline{C}, D$ .
- Representable multicategories are important because they are in equivalence with monoidal categories.

# Representable Multicategories

- We can reconstruct the coherences of a monoidal category from the representable multicategory structure, e.g.:

$$\alpha_{A_1, A_2, A_3} =$$

- We will now define *skew multicategories*, and a *left representability* condition picking out a subcategory of skew multicategories that are in equivalence with skew-monoidal categories.

# Skew Multicategories

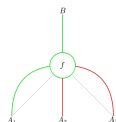
A **skew multicategory** has two types of multimorphism: **tight** and **loose**. Concretely, a skew multicategory is:

- A class of objects.
- For each object  $B$ , a class  $\text{Hom}_l(; B)$  of **loose nullary maps** into  $B$ .
- For each  $n > 0$ , and objects  $A_1, \dots, A_n, B$ , classes  $\text{Hom}_t(A_1, \dots, A_n; B)$  and  $\text{Hom}_l(A_1, \dots, A_n; B)$  of **tight** and **loose** morphisms, and a function  $\lambda_{A_1, \dots, A_n; B} : \text{Hom}_t(A_1, \dots, A_n; B) \rightarrow \text{Hom}_l(A_1, \dots, A_n; B)$  allowing us to view tight maps as loose.
- A chosen element  $\text{id}_A : \text{Hom}_t(A; A)$
- A composition operation on both tight and loose morphisms, such that  $g \circ (f_1, \dots, f_n)$  is tight exactly when  $g$  and  $f_1$  are.

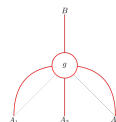


# Skew Multicategories

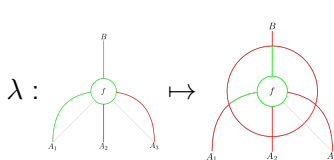
We can graphically represent this structure as follows:



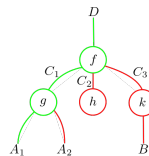
(a) A map in  $\text{Hom}_t(A_1, A_2, A_3; B)$



(b) A map in  $\text{Hom}_l(A_1, A_2, A_3; B)$



(c) Action of  $\lambda$



(d) The composite  $f(g, h, k)$ . Here  $f$  and  $g$  are tight, and thus the composite is.

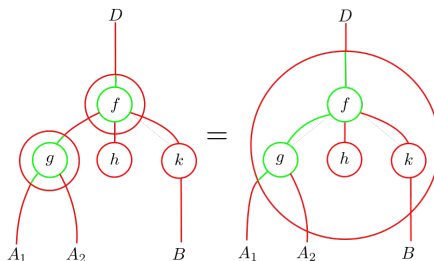
# Skew Multicategories

subject to:

- Analogues of the unitality and associativity laws.
- Composition commutes with  $\lambda$ , i.e.

$$(\lambda g) \circ (\lambda f_1, \dots, f_n) = \lambda(g \circ (f_1, \dots, f_n)).$$

The associativity and unitality are automatic from the graphical representation. The other condition is:



# Representable Skew Multicategories

A skew multicategory is **left representable** if for all lists  $\bar{A}$  of objects, there are objects  $m_t(\bar{A})$  and  $m_l(\bar{A})$ , and multimorphisms  $\theta_t(\bar{A}) \in \text{Hom}_t(\bar{A}; m_t(\bar{A}))$  and  $\theta_l(\bar{A}) \in \text{Hom}_l(\bar{A}; m_l(\bar{A}))$  which induce bijections:

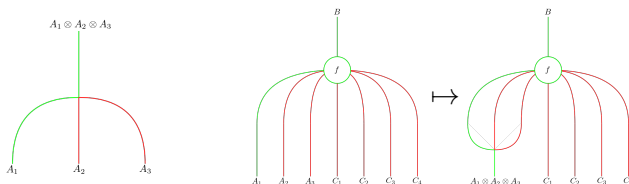
$$\text{Hom}_t(m_t(\bar{A}), \bar{C}; B) \rightarrow \text{Hom}_t(\bar{A}, \bar{C}; B)$$

$$\text{Hom}_t(m_l(\bar{A}), \bar{C}; B) \rightarrow \text{Hom}_l(\bar{A}, \bar{C}; B)$$

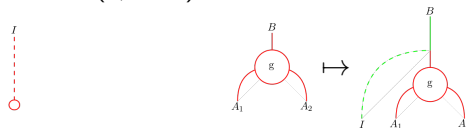
We also write  $(A_1 \otimes \cdots \otimes A_n)$  for  $m_t(A_1, \dots, A_n)$  and  $I$  for  $m_l()$ .

# Representable Skew Multicategories

We graphically represent  $\theta_t$  and the inverse to precomposition with  $\theta_t$ , a map  $\text{Hom}_t(\bar{A}, \bar{C}; B) \rightarrow \text{Hom}_t(m_t(\bar{A}), \bar{C}; B)$ , as:

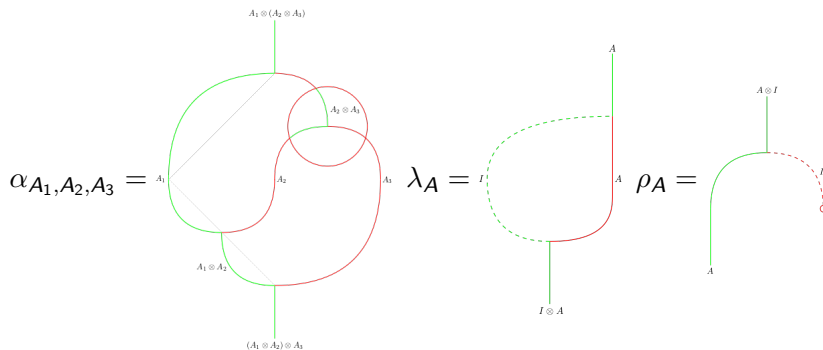


We represent  $\theta_l()$  and the inverse to precomposition with  $\theta_l()$ , a map  $\text{Hom}_l(\bar{A}; B) \rightarrow \text{Hom}_t(I, \bar{A}; B)$ , as:



# Representable Skew Multicategories

Left representable skew multicategories are in equivalence with skew-monoidal categories. For instance, we can build the associator and unitors as tight morphisms using the above structure:

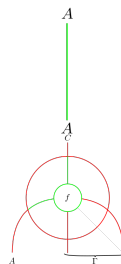


# Sequent Calculus

From the definition of left representable skew multicategories, we can derive a sequent calculus for the coherences in a skew-monoidal category. We present these one by one with justification and the corresponding graphical representation. Sequents are of the form  $S \mid \Gamma \rightarrow A$  where  $S$  is the “stoup”, either empty or a single object,  $\Gamma$  is a list of objects, and  $A$  is an object.

$$\frac{}{A \mid \rightarrow A} \text{ id}$$

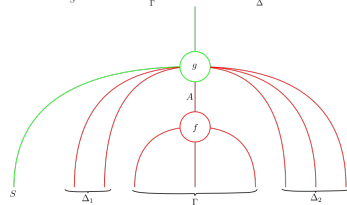
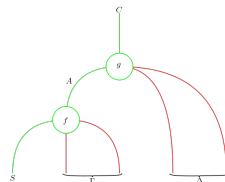
$$\frac{A \mid \Gamma \xrightarrow{f} C}{- \mid A, \Gamma \rightarrow C} \text{ shift}$$



# Sequent Calculus

$$\frac{S \mid \Gamma \xrightarrow{f} A \quad A \mid \Delta \xrightarrow{g} C}{S \mid \Gamma, \Delta \rightarrow C} \text{scut}$$

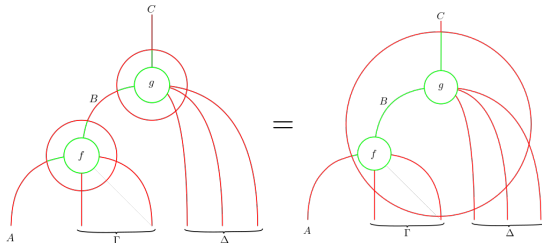
$$\frac{- \mid \Gamma \xrightarrow{f} A \quad S \mid \Delta_1, A, \Delta_2 \xrightarrow{g} C}{S \mid \Delta_1, \Gamma, \Delta_2 \rightarrow C} \text{ccut}$$



# Sequent Calculus

We impose an equational theory on derivations enforcing the axioms of a skew multicategory, for instance:

$$\frac{\frac{A \mid \Gamma \xrightarrow{f} B}{- \mid A, \Gamma \rightarrow B} \text{ shift} \quad B \mid \Delta \xrightarrow{g} C}{- \mid A, \Gamma, \Delta \rightarrow C} \text{ scut} = \frac{A \mid \Gamma \xrightarrow{f} B \quad B \mid \Delta \xrightarrow{g} C}{A \mid \Gamma, \Delta \rightarrow C} \text{ scut} \quad \frac{}{- \mid A, \Gamma, \Delta \rightarrow C} \text{ shift}$$



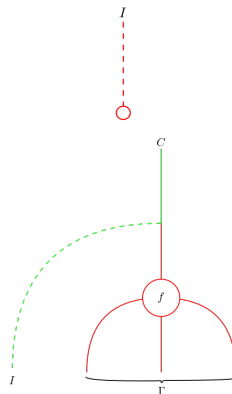


# Sequent Calculus

To capture left representable skew multicategories, we add four additional rules.

$$\frac{}{- \mid \rightarrow I} \text{ IR}$$

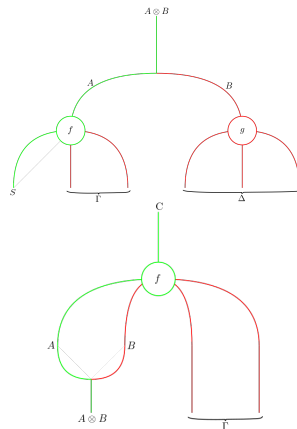
$$\frac{- \mid \Gamma \xrightarrow{f} C}{I \mid \Gamma \rightarrow C} \text{ IL}$$



# Sequent Calculus

$$\frac{S \mid \Gamma \xrightarrow{f} A \quad - \mid \Delta \xrightarrow{g} B}{S \mid \Gamma, \Delta \rightarrow A \otimes B} \otimes R$$

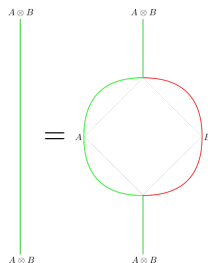
$$\frac{A \mid B, \Gamma \xrightarrow{f} C}{A \otimes B \mid \Gamma \rightarrow C} \otimes L$$



# Sequent Calculus

These are also subject to equations, like:

$$\frac{}{A \otimes B \mid \rightarrow A \otimes B} \text{id} = \frac{\frac{}{A \mid \rightarrow A} \text{id} \quad \frac{\frac{}{B \mid \rightarrow B} \text{id} \quad \frac{}{- \mid B \rightarrow B} \text{shift}}{A \mid B \rightarrow A \otimes B} \otimes R}{A \otimes B \mid \rightarrow A \otimes B} \otimes L$$



# Main Theorem

## Theorem

*Morphisms from  $A$  to  $B$  in a free skew-monoidal category on a set of objects are in bijection with equivalence classes of derivations of  $A \mid \rightarrow B$  in the above sequent calculus.*

# References

- [1] John Bourke and Stephen Lack. “Skew monoidal categories and skew multicategories”. In: *Journal of Algebra* 506 (July 2018), pp. 237–266. ISSN: 0021-8693. DOI: 10.1016/j.jalgebra.2018.02.039. URL: <http://dx.doi.org/10.1016/j.jalgebra.2018.02.039>.
- [2] Tarmo Uustalu, Niccolò Veltri, and Noam Zeilberger. *The Sequent Calculus of Skew Monoidal Categories*. 2020. arXiv: 2003.05213 [cs.LG].
- [3] Paul Wilson, Dan Ghica, and Fabio Zanasi. *String diagrams for non-strict monoidal categories*. 2022. arXiv: 2201.11738 [math.CT].