

# THE JONES POLYNOMIAL AND AFT

## - PART I -

### FEYNMAN DIAGRAMS, AN ORIENTED STATE MODEL FOR VR(t)

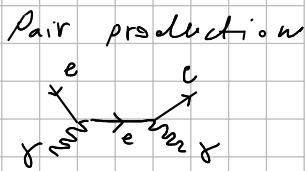
#### § 1 Motivation from physics - Feynman diagrams (1948)

Quantum Electrodynamics (QED) simplest of the dynamical theories



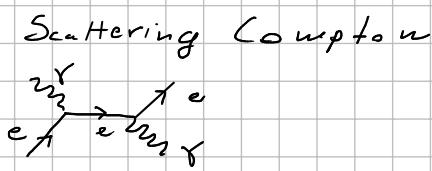
$$\text{Scattering Møller} \quad e^- + e^+ = e^- + e^+$$

Any lepton does the same as they don't carry colours!



$$e^- + e^+ = \gamma + \gamma$$

$$\gamma + \gamma = e^+ + e^-$$



$$\gamma + \gamma = \gamma + e^+$$

Physicists are interested in these three values:  
(processes in nature)

This answers the question: why are these diagrams useful?

- cross-section (1)
- decay rates times (2)
- bound state (3)

(1) scattering - when particles collide

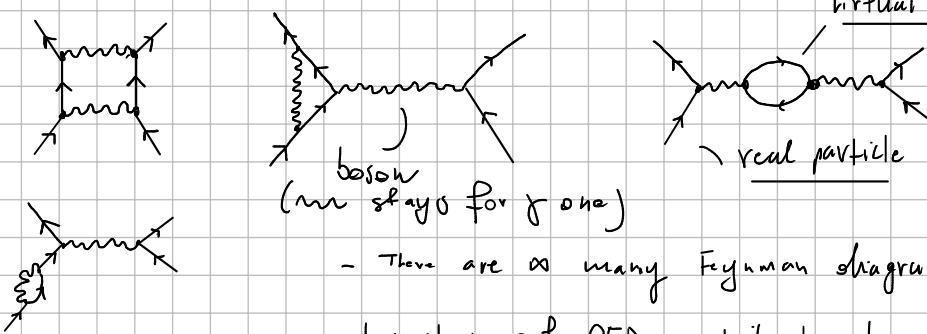
(2) when a particle naturally loses energy and changes during time (\*) see cut off this talk notes

(3) Quantum states of a particle subject to a potential such that the particle has the tendency to remain localised in one or more regions of space.

One can increase vertices --- for example 4

they do not lie on their mass shell

virtual particles (not observed particles lie inside loops)



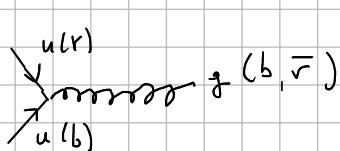
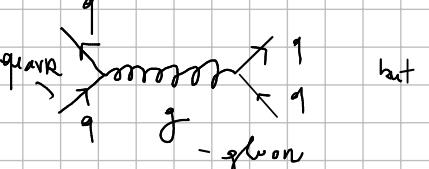
- There are as many Feynman diagrams for a given process, BUT  
as each vertex of QED contributes by a factor of  $\alpha_e = \frac{e^2}{\hbar c}$  (fine structure constant)

$\Rightarrow$  diagrams with more & more vertices contribute less & less to the final result and can be therefore safely ignored (usually  $\propto V^4$  in QED).

- There are conservation laws at vertex (mass, energy, spin) to be imposed -

\* What if we have colours? QCD (Quantum Chromodynamics)

colour (like charge)  
is conserved



2 gluons (bicoloured)

r = red b = blue g = green

as quarks carry colour ↴

Here is possible to have glueballs too: bound states of interacting gluons



However, observe that now  $\alpha_g > 1$  (not  $\alpha_c < 1$ )  
as above

&  $\alpha_g$  is not constant, but depends on  
distance between particles! ASYMPTOTIC  
FREEDOM

→ Diagrams with more vertices will contribute more & more -

In QED vacuum is a dielectric, the vacuum polarisation partially screens the charge and reduces its field -

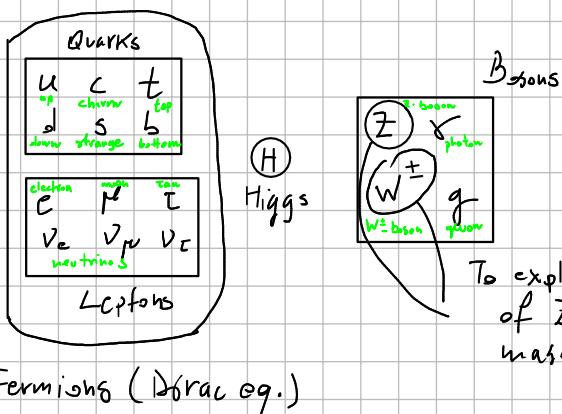
In QCD we have both quark polarisation & gluon polarisation diagrams -  
drives  $\alpha_s \nearrow$  at short distances      ↴ drives it ↓

Let  $f := \#$  flavours &  $n := \#$  colours  $\Rightarrow \alpha = 2f - 4pn -$

For the standard model,  $f=6$ ,  $n=3 \Rightarrow \alpha = -21 -$

### The Standard Model

Quarks and leptons acquire masses (amount depends on the coupling strength) interacting with the H field -



To explain the non-zero masses of Z and W, bosons & fermions masses are generated by the Higgs mechanism

Fermions (Dirac eq.)

→ QED vs QCD: electric charge is a true feature of particles, no colour is carried by elementary ones -

( $\infty$  mesons)

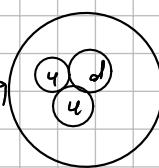
( $\infty$  baryons)

$W^\pm, Z$

weak interactions  
(bosons)

Bound states of quarks

hydrogen atom  
(nucleus)



proton p

CONSERVATION LAWS: charge, colours, baryon & lepton spin numbers (mass & momenta) -

### § 2. Back to knots: what we already know -

Reidemeister moves

~ ambient isotopy

✓ ~ ~

0.

Y ~ ~ ~ Y

1.

Y ~ Y

2.

$$\text{Diagram} \sim -\text{Diagram} \quad \& \quad \text{Diagram} \sim \text{Diagram}$$

Just sliding of the top string ✓

$$\rightarrow \text{Remember also the parity } \begin{array}{c} \nearrow \searrow \\ +1 \end{array} \quad \begin{array}{c} \nearrow \searrow \\ -1 \end{array}$$

→ Invariants converge by summing / integrating over the set of states for a given diagram (evolution of states runs fund. group)

$$\rightarrow \text{Splicing} \quad \begin{array}{c} \nearrow \searrow \\ A \\ \downarrow \\ B \end{array} \quad \begin{array}{c} \overbrace{\quad}^B \\ A \end{array} \quad \text{or} \quad \begin{array}{c} \nearrow \searrow \\ B \\ \downarrow \end{array} \quad \Rightarrow \text{for } n \text{ splicings, we have } 2^n \text{ configurations!}$$

→ If  $\bar{G}$  is a state of  $K$ ,  $\langle K | \bar{G} \rangle$  is the commutative product of labels attached to  $\bar{G}$  &  $\|\bar{G}\| = \# \text{ of loops in } \bar{G} - 1$

$$= d^{\|\bar{G}\|}$$

$$\rightarrow \text{BRACKER POLYNOMIAL} \quad \langle K \rangle = \langle K \rangle(A, b, d) = \sum_{\bar{G}} \langle K | \bar{G} \rangle \langle \bar{G} \rangle$$

This is not a topological invariant: The bracket with  $B = A^{-1}$  &  $d = -\left(A^2 + \frac{1}{A^2}\right)$  is invariant under type 2 & 3 Reidemeister moves, but the invariance under 1. is gained via normalisation!

Let  $K$  be an oriented link

$$\omega(K) = \sum_{p \in CK(K)} \varepsilon(p)$$

$$(1) \quad \omega(\text{Diagram}) = 1 + \underbrace{\omega(\text{Diagram})}_{=0} \quad \& \quad \omega(\text{Diagram}) = -1 + \underbrace{\omega(\text{Diagram})}_{=0} \quad (2)$$

NORMALISED BRACKET

$$\mathcal{L}_K := (-A^3)^{-\omega(K)} \langle K \rangle$$

$$(1) \quad \langle \text{Diagram} \rangle = (-A^3) \langle \text{Diagram} \rangle \quad \& \quad \langle \text{Diagram} \rangle = (-A)^{-3} \langle \text{Diagram} \rangle \quad (2)$$

$$\text{Property: } \mathcal{L}_{K^*}(A) = \mathcal{L}_K(A^{-1})$$

⇒ If  $\mathcal{L}_K(A) \neq \mathcal{L}_K(A^{-1})$ ,  $K \not\sim K^*$  ⇒  $\mathcal{L}_K$  DETECTS CHIRALITY ✓

Property: If  $T$  is a knot,  $\omega(T)$  does not depend on the orientation

$$\Rightarrow \mathcal{L}_T = (-A^3)^{-\omega(T)} \langle T \rangle \rightarrow \mathcal{L}_L = \langle L \rangle \text{ for a linked diagram} \quad \checkmark$$

$$\text{Theorem: } \mathcal{L}_K\left(\frac{1}{\sqrt{E}}\right) = V_K(t) \quad \checkmark \text{ Jones polynomial}$$

normalised bracket

Def.  $V_K(t)$  is a Laurent polynomial in  $\sqrt{t}$  assigned to an oriented  $K$ :

- $K \sim K' \Rightarrow V_K(t) = V_{K'}(t);$

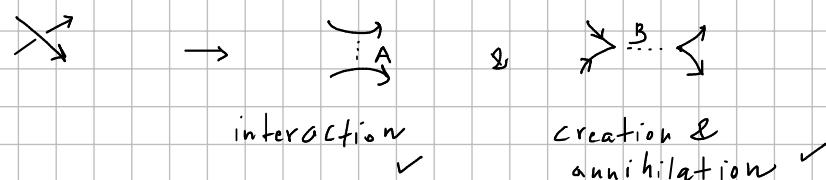
- $V_\emptyset = 1;$

- $\frac{1}{t} V_{\nearrow\searrow} - t V_{\nearrow\searrow} = (\sqrt{t} - 1/\sqrt{t}) V_{\nearrow\searrow}.$

### § 3. A new perspective [why section 1. makes sense]

Jones polynomial  $V_K(t) = (-t^{3/4})^{\omega(K)} \langle K \rangle \left(\frac{1}{\sqrt{t}}\right)$

- What about splicing a Feynman diagram?



! Time here is local  
✓ coincides with  
the arrows on the  
diagram!

Observe that:  $\text{---} \sim \text{---}$ ,  $\nearrow \nwarrow \sim \nearrow \nwarrow$ ,  $\nearrow \nearrow \sim \nearrow \nearrow$ ,  $\circlearrowleft \sim \circlearrowright \sim \circlearrowright$

- topological diagrams do not know about field effect  $\nearrow \nearrow \nearrow$  - is ignored ( $\nearrow$ )



• We impose (cross) channel unitarity, i.e.  $\nearrow \nearrow \sim \nearrow \nearrow$  &  $\nearrow \nearrow \sim \nearrow \nearrow$   
(interaction of particles) & triangle invariance

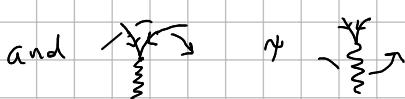


$$\Rightarrow \sqrt{\nearrow \nearrow}^+ = A V_{\nearrow\searrow} + B V_{\nearrow\searrow}^- \quad \& \quad \sqrt{\nearrow\searrow}^- = A' V_{\nearrow\searrow} + B' V_{\nearrow\searrow}^+$$

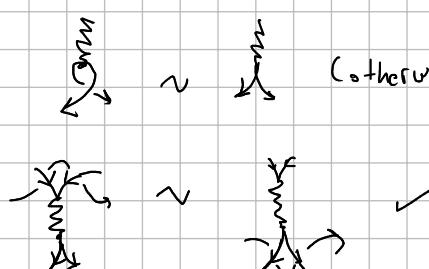
$$V_{\nearrow\searrow} = \delta V_K \quad (\text{loop value})$$

We want to determine  $A, B, \delta$

Remark: It is not true that



However



(otherwise is trivial)

### 3.1 Calculations

## 1. Channel unitarity ( & cross-channel unitarity )

$$V = AV + BV \xrightarrow{\text{unitarity}} = A\{A'V\} + B\{B'V\} + B\{A'V + B'V\} = AA'V + (AB' + BA')V$$

From cross-dimensionality unitarity (left as exercise), we obtain an analogous set of eq.s. Combining both of them, we get:

$$AA' = I = BB'$$

$$\delta = -\left(\frac{A}{B} + \frac{A'}{B'}\right) = -\left(\frac{B'}{A'} + \frac{B}{A}\right) \quad \left\{ \begin{array}{l} A' = A^{-1}, \quad B' = B^{-1}, \quad \delta = -\left(\frac{A}{B} + \frac{B}{A}\right) \\ (1) \end{array} \right.$$

2. Annihilation & crossing (this implies triangle invariance) -

$$\begin{aligned} \sqrt{\frac{K}{K'}} &= A' V_{\downarrow} + B' V_{\rightarrow} = A' \left\{ A' V_{\downarrow} + B' V_{\rightarrow} \right\} + B' \left\{ A' V_{\downarrow} + B' V_{\rightarrow} \right\} = \\ &= \left( (A')^2 + (B')^2 + \delta_{A'B'} \right) V_{\downarrow} + A' B' V_{\rightarrow} \stackrel{(1)}{=} \frac{1}{AB} V_{\rightarrow} \quad (\text{the particle just passed by}) \end{aligned}$$

$$\text{Similarly, } \sqrt{V_{\downarrow k}} = AD V_{\downarrow k} \quad \text{and also} \quad \sqrt{\frac{1}{\pi k}} = AB V_{\downarrow k} \quad (\text{a})$$

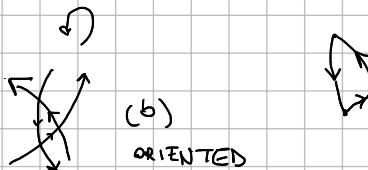
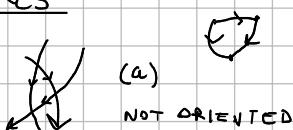
$$\sqrt{\frac{1}{AB}} = \frac{1}{\sqrt{AB}} \quad (\text{b}) \quad \text{and} \quad \sqrt{\frac{1}{k}} = \frac{1}{\sqrt{k}} \quad (\text{c})$$

It implies triangle invariance - Indeed:

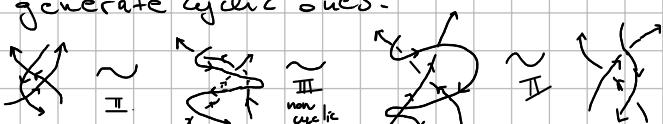
$$V_{\downarrow \downarrow} = A V_{\downarrow \downarrow} + B V_{\downarrow \downarrow} = A V_{\downarrow \downarrow} + B V_{\downarrow \downarrow} = V_{\downarrow \downarrow}$$

### 3. Oriented triangles

Def.



However, we can transform (a) into (b) via R-moves, as non-cyclic triangle moves can generate cyclic ones.



$\Rightarrow$  we found:

$$V_{\nearrow \nwarrow} = AV_{\nearrow} + BV_{\nwarrow}$$

$$V_{\nearrow \nwarrow} = \frac{1}{A} V_{\nearrow} + \frac{1}{B} V_{\nwarrow}$$

#### 4. Type I. invariance (invariance under RI)

$$V_{\nearrow \nwarrow} = AV_{\nearrow} + BV_{\nwarrow} = (A\delta + B) V_{\nearrow}$$

$$A\delta + B \stackrel{(1)}{=} A\left(-\frac{A}{B} - \frac{B}{A}\right) + B = -A^2/B$$

$$\Rightarrow V_{\nearrow \nwarrow} = V_{\nearrow} \Leftrightarrow A\delta + B = 1 = -A^2/B \Leftrightarrow B = -A^2 \checkmark$$

#### § 3.2 The Jones polynomial

We have just concluded  $B = -A^2$ . Let  $A = -\sqrt{t} \Rightarrow B = -t$ , this implies

$$\begin{cases} V_{\nearrow \nwarrow} = -\sqrt{t} V_{\nearrow} - t V_{\nwarrow} \\ V_{\nearrow \nwarrow} = -\frac{1}{\sqrt{t}} V_{\nearrow} - \frac{1}{t} V_{\nwarrow} \end{cases}$$

$$t^{-1} V_{\nearrow \nwarrow} + (-t) V_{\nearrow \nwarrow} = \left(-\frac{\sqrt{t}}{t} + \frac{t}{\sqrt{t}}\right) V_{\nearrow} = \left(\sqrt{t} - \frac{1}{\sqrt{t}}\right) V_{\nearrow}$$

$\Rightarrow V_R(t)$  is the 1-variable Jones polynomial.

$$\text{Lemma: } V_{\nearrow \nwarrow} = t\sqrt{t} V_{\nearrow} \quad V_{\nearrow \nwarrow} = \frac{t}{t\sqrt{t}} V_{\nwarrow}$$

$$V_{\nearrow \nwarrow} = t\sqrt{t} V_{\nearrow} \quad V_{\nearrow \nwarrow} = \frac{t}{t\sqrt{t}} V_{\nwarrow}$$

$$V_{\nearrow \nwarrow} = -\sqrt{t} V_{\nearrow} - t V_{\nwarrow} \Leftrightarrow \left(-\sqrt{t} - t \left(-\sqrt{t} - \frac{1}{\sqrt{t}}\right)\right) V_{\nearrow} = t\sqrt{t} V_{\nwarrow} \checkmark$$

Analogously for the other cases.

• Let's call  $\hat{V}$  the renormalised Jones poly., then:

$$t\hat{V}_{\nearrow \nwarrow} - \frac{t}{t} \hat{V}_{\nearrow \nwarrow} = \left(\sqrt{t} - \frac{1}{\sqrt{t}}\right) \hat{V}_{\nearrow}$$

$$\hat{\delta} = \left(\sqrt{t} + \frac{1}{\sqrt{t}}\right) \text{ the loop variable}$$

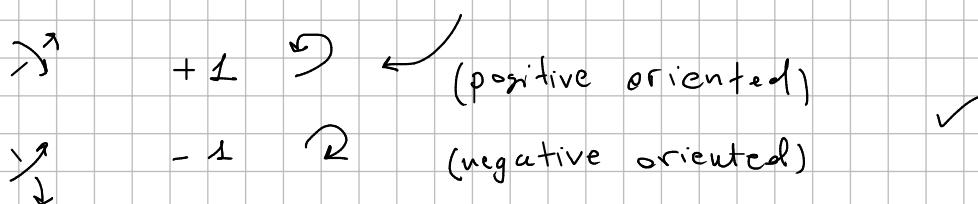
physical interpretation: For corresponding state expansion, we have:

$$V_K = \sum_{\vec{G}} \langle K | \vec{G} \rangle \left( \sqrt{t} + \frac{1}{\sqrt{t}} \right)^{\| \vec{G} \|}$$

where  $\langle K | \vec{G} \rangle$  is the product of the vertex weights

$\| \vec{G} \| :=$  # of loops in the oriented state

For the sign is  $(-1)^\pi$  where  $\pi$  is the parity of the number of creation - annihilation splines in the state  $\vec{G}$

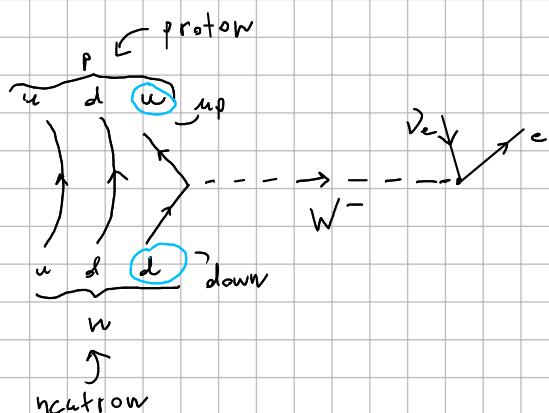


\* On the decays (some examples):

-  $\alpha$ -decay: does not change the structure or composition of an atom.

Some unstable nuclei dissipate excess of energy via electromagnetic radiation. One

-  $\beta$ -decay:  $n = p + e^- + \bar{\nu}_e$  antineutrino



In this process the quark changes its flavour - A neutron transforms into a proton by the emission of an electron and an antineutrino.

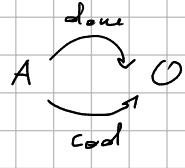
# THE JONES POLYNOMIAL AND QFT

- PART II -

## STRINGS & BRAIDED CATEGORIES: THE JONES POLYNOMIAL REVEALED

### § 1. A primer in category theory

Def. A directed graph is a set of objects  $\mathcal{O}$ , a set of morphisms (arrows)  $A$  and two functions



$$\text{ex. } f: a \rightarrow b, f \in A$$

$$\begin{aligned} \text{dom}(f) &= a \{ \in \mathcal{O} \\ \text{cod}(f) &= b \{ \in \mathcal{O} \end{aligned}$$

In this graph, the set of composable arrows is the set:

$$A \times_{\mathcal{O}} A = \{ \langle g, f \rangle : g, f \in A \text{ & } \text{dom } g = \text{cod } f \}$$

called the product over  $\mathcal{O}$ .

Def. A category is a graph with two additional functions:

$$\begin{array}{ccc} \mathcal{O} & \xrightarrow{\text{id}} & A \\ (\text{identity}) & c \rightarrow \text{id}_c & A \times_{\mathcal{O}} A \xrightarrow{\circ} A \\ & & \langle g, f \rangle \rightarrow g \circ f \end{array} \quad (\text{composition})$$

such that  $\text{dom}(\text{id}_a) = a = \text{cod}(\text{id}_a)$

$$\text{dom}(g \circ f) = \text{dom } f$$

$$\text{cod}(g \circ f) = \text{cod } g$$

$$\forall a \in \mathcal{O}$$

$$\forall \langle g, f \rangle \in A \times_{\mathcal{O}} A$$

and - associativity       $\overline{\text{For } a \xrightarrow{f_1} b \xrightarrow{f_2} c \xrightarrow{f_3} d, (f_2 \circ f_1) \circ f_3 = f_2 \circ (f_1 \circ f_3)}$

- unit law       $\forall f: a \rightarrow b, g: b \rightarrow c \quad f \circ \text{id}_b = f \text{ & } g \circ \text{id}_c = g$

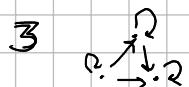
Remark (notation): If  $\mathcal{C}$  is a category, the set of objects is denoted also by  $\mathcal{C}^0$ , the set of morphisms by  $\mathcal{C}^1$  or  $\mathcal{C}^1(a, b)$  when one wants to stress  $\text{dom}$  &  $\text{cod}$ . We also write this last one has

$$\text{hom}(a, b) = \{f \mid f \in \mathcal{C}^1, \text{dom } f = a, \text{cod } f = b\}$$

Examples:  $\emptyset$   $\neq$  category, i.e.  $\emptyset^0 = \emptyset^1 = \emptyset$

1)  $\bullet$  one obj. and one (identity) arrow.

2)  $\bullet \xrightarrow{\quad} \bullet$



and so on...

↓ two obj & two non-trivial arrows ( $e.g. g, f: a \rightarrow b$ ) called parallel arrows

Monoids: A monoid is a category with one object. It can be described as a set  $M$  with a binary operation  $M \times M \rightarrow M$  associative and unital. It is a semigroup with identity.

For any category  $\mathcal{C}$  and  $a \in \mathcal{C}^\circ$ ,  $\text{hom}(a, a)$  is a monoid.

$\text{Set}, \text{Set}_*, \text{Top}, \text{Top}_*, \text{Top}_\text{h}$  (morphisms are homotopy classes of maps)

$\text{Ab}, \text{Grp}, R\text{-Mod}, \text{Rng}, \dots$

Def. A morphism of categories is called Functor, e.g.  $\Phi: \mathcal{C}_1 \rightarrow \mathcal{C}_2$

$$\forall c \in \mathcal{C}_1 \quad \Phi(c) \in \mathcal{C}_2^\circ$$

$$\forall f \in \mathcal{C}_1(a, b) \quad \Phi(f) \in \mathcal{C}_2(\Phi(a), \Phi(b))$$

in such a way that  $\Phi(1_c) = 1_{\Phi(c)}$   $\forall c \in \mathcal{C}_1^\circ$

$$\text{if covariant} \quad \Phi(g \circ f) = \Phi(g) \circ \Phi(f) \quad \forall (g, f) \in \mathcal{C}_1^\circ$$

(if contra-variant) (or  $\Phi(g \circ f) = \Phi(f) \circ \Phi(g)$ )

Examples:  $\text{Top}_* \xrightarrow{\pi_n} \text{Grp} \quad (X, x_*) \rightarrow \pi_n(X, x_*)$

$$\text{Grp} \xrightarrow{\phi} \text{Ab} \quad G \rightarrow G/[G, G]$$

forgetful  $\rightarrow \text{Grp} \xrightarrow{\cup_1} \text{Set} \quad G \rightarrow U(G)$  forgetting group structure

$$\text{Set} \xrightarrow{\mathcal{P}} \text{Set} \quad X \rightarrow \mathcal{P}(X)$$

what if we erase the inverse image?  $f^{-1}$

$$\{x \rightarrow y\} \rightarrow \{\mathcal{P}(x) \rightarrow \mathcal{P}(y)\}$$

$$\begin{aligned} \Phi(1_X) &= 1_{\mathcal{P}(X)} \\ \Phi(fg) &= \mathcal{P}f \circ \Phi(g) \end{aligned}$$

covariance

Def. An arrow  $m: a \rightarrow b$  is monic in  $\mathcal{C}$  iff  $\forall f, f_a: d \rightrightarrows a$ ,   
  $m \circ f_1 = m \circ f_2 \Rightarrow f_1 = f_2$  (left cancellable). monic  
in Set?

Def. An arrow  $h: a \rightarrow b$  is epic (epi) in  $\mathcal{C}$  iff  $\forall g, g_b: b \rightrightarrows c$    
  $g_1 \circ h = g_2 \circ h \Rightarrow g_1 = g_2$  (right cancellable). epi  
in Set?  
surj?

Def. Two objects  $a, b \in \mathcal{C}^\circ$  are isomorphic ( $a \cong b$ ) iff  $\exists$  an isomorphism  $e: a \rightarrow b$ , i.e. iff  $\exists$  an arrow  $e: a \rightarrow b$  invertible (i.e.  $\exists e^{-1}: b \rightarrow a$  |  $e^{-1}e = 1_a$  &  $ee^{-1} = 1_b$ ). iso in Set? bijections ✓

## fd - Monoidal categories\* (a functor of functors is a natural transfo-)

Roughly, a monoidal category, is a category equipped with a product.

Def. A strictly monoidal category  $\langle \mathcal{B}, \square, e \rangle$  is a category  $\mathcal{B}$  with a

bifunctor (functor in both variables)  $\square: \mathcal{B} \times \mathcal{B} \rightarrow \mathcal{B}$  which is

- associative  $\square(\square \times 1) \stackrel{(")}{=} \square(1 \times \square): \mathcal{B} \times \mathcal{B} \times \mathcal{B} \rightarrow \mathcal{B}$

- has a unit, i.e. an object  $e: \square(e \times 1) \stackrel{\alpha}{=} \text{id}_{\mathcal{B}} = \square(1 \times e)$

$$\text{where } (\text{ex } 1)(c) \rightarrow (e, c) \quad e \square c = c = c \square e \quad \forall c \in \mathcal{B}^0$$

$$\mathcal{B} \rightarrow \mathcal{B} \times \mathcal{B}$$

$$1_e \square f = f = f \square 1_e \quad \forall f \in \mathcal{B}^1$$

How it works?

On objects:  $(a, b) \in (\mathcal{B}^0)^2 \rightarrow a \square b \in \mathcal{B}^0$

On morphisms:  $(f, g) \in (\mathcal{B}^1)^2 \rightarrow f \square g: a \square b \rightarrow a' \square b'$  for  $f: a \rightarrow a'$   
 $g: b \rightarrow b'$

such that

$$\ell_a \square \ell_b = \ell_{a \square b}$$

$$(f' \square g')(f \square g) = f'f \square g'g$$

Def. A monoidal category  $\mathcal{B} = \langle \mathcal{B}, \square, e, \alpha, \lambda, \rho \rangle$  is a category  $\mathcal{B}$  with a bifunctor  $\square: \mathcal{B} \times \mathcal{B} \rightarrow \mathcal{B}$ ,  $a \in \mathcal{B}$  as above and three natural isomorphisms  $\alpha, \lambda, \rho$  as follows:

(1) associativity

$$\alpha = \alpha_{a,b,c}: a \square (b \square c) \xrightarrow{\cong} (a \square b) \square c \quad \forall a,b,c \in \mathcal{B}^0$$

$$\begin{array}{ccc} a \square (b \square (c \square d)) & \xrightarrow{\alpha} & (a \square b) \square (c \square d) \xrightarrow{\alpha} ((a \square b) \square c) \square d \\ \downarrow 1 \square \alpha & & \downarrow \alpha \square 1 \\ a \square ((b \square c) \square d) & \xrightarrow{\alpha} & (a \square (b \square c)) \square d \end{array}$$

(2) unit law (left & right)

$$\lambda_a: e \square a \xrightarrow{\cong} a \quad \rho_a: a \square e \xrightarrow{\cong} a$$

$$: a \square (e \square c) \xrightarrow{\alpha} (a \square e) \square c$$

$$\forall a \in \mathcal{B}^0$$

$$\begin{array}{ccccc} 1 \square \lambda & \searrow & \swarrow & \rho \square 1 & \quad \& \quad \lambda_e = \rho_e: e \square e \rightarrow e \\ & & & & \checkmark \end{array}$$

Examples:  $(\text{Set}, \times)$   $(\text{Ring}, \otimes)$

They are usually called tensor categories rather than monoidal in some textbooks -  $(\text{Vect}_{\mathbb{F}}, \otimes, 1\text{-dim sub.})$

Def. A monoid  $m \in \mathcal{B}$  (monoidal) is an object together with two arrows  $\mu : m \square m \rightarrow m$   $\eta : e \rightarrow m$  such that

$$\begin{array}{ccccc} m \square (m \square m) & \xrightarrow{\alpha} & (m \square m) \square m & \xrightarrow{\mu \square 1} & m \square m \\ 1 \square \mu \downarrow & & \Downarrow & & \downarrow \mu & (2) \\ m \square m & \xrightarrow{\mu} & m & & \end{array}$$

$$\begin{array}{ccccc} e \square m & \xrightarrow{\eta \square 1} & m \square m & \xleftarrow{1 \square \eta} & m \square e \\ & \searrow & \downarrow \mu & \swarrow & \\ & & m & & \end{array} \quad (2)$$

• An arrow between monoids  $f : \langle m, \mu, \eta \rangle \rightarrow \langle m', \mu', \eta' \rangle$  is a morphism  $f : m \rightarrow m'$  such that

$$f\mu : \mu' (f \square f) : m \square m \rightarrow m'$$

$$f\eta = \eta' : e \rightarrow m'$$

→ The monoids in  $\mathcal{B}$  form therefore a category denoted by  $\text{Mon}_{\mathcal{B}}$  and  $\exists$  a forgetful functor  $U : \text{Mon}_{\mathcal{B}} \rightarrow \mathcal{B}$

$$\langle m, \mu, \eta \rangle \mapsto m$$

that just forgets the multiplication  $\mu$  -

$$\begin{array}{c|ccc|c} \text{Examples:} & \langle R\text{-mod}, \otimes, R \rangle & \longleftrightarrow & \mathbb{Z}\text{-algebras} & \hookrightarrow \text{Mon}_{\mathcal{B}} \\ \mathcal{B} \nearrow & \langle \text{Ab}, \otimes, \mathbb{Z} \rangle & \longleftrightarrow & \mathbb{Z}\text{-modules} & \end{array}$$

To see: The simplicial category as UNIVERSAL MONOID -



$$\Delta \rightarrow \text{Top}$$

\* The coherence theorem part is kept for time reasons (please see Mac Lane) -

## 3. J. Braiding

(Ch. XI Mac Lane)

Def. A monoidal category  $\mathcal{B}$  is called symmetric if the commutativity property holds up to a natural isomorphism  
 (3.1)  $\gamma : a \square b \cong b \square a$ . (see def 3.5)

- Is it always possible to strictify, i.e.  $\gamma^\vee = \text{id}$ ? No

3 examples of categories with a twist, i.e.  $\gamma^2 \neq 1$  ↴ Braided cat. II

Def. A braiding for a monoidal category  $\mathcal{M}$  consists of a family  
 (3.2) of iso's  $\gamma_{a,b} : a \square b \cong b \square a$  natural in  $a, b \in \mathcal{M}^\circ$ :

$$a \square e \xrightarrow{\gamma} e \square a$$

$$\begin{matrix} g \downarrow & \curvearrowright \\ a & \curvearrowleft \end{matrix}$$

and  $\gamma^*$

$$(ab)c \xrightarrow{\gamma} c(ab)$$

$$\begin{matrix} \alpha^{-1} \downarrow & \curvearrowright \\ a(bc) & \curvearrowleft \end{matrix}$$

$$(ca)b \downarrow \gamma^{\square 1}$$

$$1 \square \gamma \downarrow$$

$$a(cb) \xrightarrow{\gamma} (ac)b$$

$$\begin{matrix} & \curvearrowright \\ \alpha & \curvearrowleft \end{matrix}$$

$$a(bc) \xrightarrow{\gamma} (bc)a$$

$$\begin{matrix} & \curvearrowright \\ \alpha & \curvearrowleft \end{matrix}$$

$$b(ca) \downarrow \alpha^{-1}$$

$$(ab)_c \curvearrowright$$

$$b(ac) \downarrow 1 \square \gamma$$

$$(ba)_c \xrightarrow{\alpha} b(ac)$$

\* dropping  $\square$

Remark: If  $\gamma$  is a braiding on  $\mathcal{M}$ ,  $\gamma^{-1}$  also is -

Def. (3.3) A symmetric monoidal category, is a category  $\mathcal{M}$  with a braiding  $\gamma$ :

$$ab \xrightarrow{\gamma_{a,b}} ba$$

$$\begin{matrix} id_{ab} \searrow & \curvearrowright \\ ab & \downarrow \gamma_{b,a} \end{matrix}$$

Remark: For symmetric monoidal cat. all formal diagrams involving  $\alpha$  &  $\gamma$  commute (coherence) -

For braided cat. - 3  $\infty$  number of canonical automorphisms

$$1, \gamma^2, \gamma^4, \dots, \gamma^{-2}, \gamma^{-4}, \dots : a \square b \rightarrow a \square b.$$

↙ ↘

### § 4. The Artin braid groups $B_n$ and the braid category

A braid is formed by strings fixed at the extremes of the diagram as follows:



The one in the example is an element of  $B_3$ , the Artin braid group on 3 strings.

Def. The  $n$ -th Artin braid group is defined as

$$B_n = \{ \text{braids on } n \text{ strings} \} / \sim$$

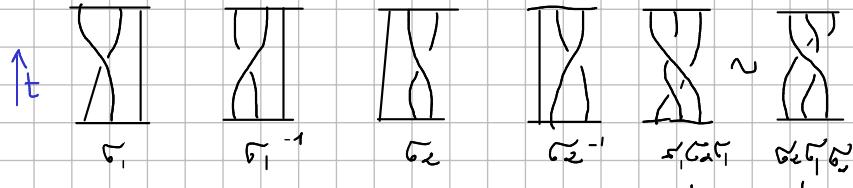
where  $b_1, b_2 \in B_n$  iff  $b_1$  can be continuously deformed into  $b_2$  without crossing or cutting the component strings.

- On  $B_n$  we define the composition operation simply composing the single paths (strings) in the braids (fuxtaposition).

Example : 

#### 4.1 Generators of $B_n$ & $S_n$

On  $B_3$ , consider the following



$t_i := i\text{-th string under the } (i+1)\text{-th}$  - (\*)

$\xrightarrow{\text{Composition reads way}}$

$$\Rightarrow B_3 = \langle t_1, t_2 \mid t_1 t_2 t_1 = t_2 t_1 t_2 \rangle \text{ as presentation.}$$

It is quite natural to think about the existence of relation between the  $n$ -th Artin group  $B_n$  and the group  $S_n$  of permutations over  $n$  elements  $B_n \rightarrow S_n$ .

or can also be omitted, as it is a group.

$$B_n = \langle t_1, \dots, t_{n-1} : t_i t_i^{-1} = 1, t_i t_j = t_j t_i \text{ for } |i-j| > 1, t_i t_{i+1} t_i = t_{i+1} t_i t_{i+1} \rangle$$

$t_i$  as in (\*) above.

- Little observation:  $B_n$  can be also seen as fundamental group of  $T_n$ , where we fix  $P = \mathbb{E}^n$  and define  $T_n = \{(x_1, \dots, x_n) \mid x_i \in P \forall i\}$

$$\Rightarrow B_n = \text{fundamental group of } T_n. \quad (\text{to become clear later})$$

Remember: every permutation  $\sigma \in S_n$  can be written as a finite product of transpositions.

$$\text{Ex. } (123) \rightarrow (321) = (13) \quad \checkmark$$

$$(12345) \rightarrow (54321) = (12)(24)(15) \quad \checkmark$$

The permutation group  $S_n$  is presented as follows:

$$S_n = \langle \tau_i := (i, i+1) \mid \tau_i^2 = 1, \tau_i \tau_j = \tau_j \tau_i \text{ for } i \neq j, \tau_i \tau_{i+1} \tau_i = \tau_{i+2} \tau_i \tau_{i+1} \rangle$$

$$\Rightarrow \text{It is apparent that } \exists \text{ an homomorphism } \begin{matrix} B_n & \xrightarrow{\delta_i} & S_n \\ \tau_i & \xrightarrow{\delta_i} & \tau_i \end{matrix} \quad i=1, \dots, n-1$$

### § 4.2 The braid category $B$

All the braid groups  $B_n$  can be assembled into a category  $B$ , called the braid category, whose:

- objects are the natural numbers  $\mathbb{N}^{\geq 0}$
- arrows are the braids  $n \rightarrow n$  ( $\nexists$  arrow  $n \rightarrow m$  for  $n \neq m$   
and only the identity arrow  $\emptyset \rightarrow \emptyset$ )

$\Rightarrow$  This defines a monoidal category  $B$  with product  $+ : B \times B \rightarrow B$ .

The addition operation  $+$ :

- on objects, it is just addition of natural numbers  $\checkmark$
- On morphisms,

$$\text{ex: } \begin{array}{c} \text{Diagram for } B_3 \\ \text{with 3 strands} \end{array} + \begin{array}{c} \text{Diagram for } B_2 \\ \text{with 2 strands} \end{array} = \begin{array}{c} \text{Diagram for } B_5 \\ \text{with 5 strands} \end{array}$$

$$(3 \rightarrow 3) + (2 \rightarrow 2) = (5 \rightarrow 3+2=5)$$

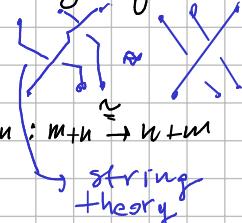
as defined for standard  $\square$  operation  $\checkmark$

$$\text{In general, } (n \rightarrow n) + (m \rightarrow m) = (n+m) \rightarrow (n+m)$$

This operation can be shown to be associative and with a unit given by the empty braid  $\emptyset \rightarrow \emptyset \Rightarrow \langle B, +, \emptyset \rightarrow \emptyset \rangle$  is a strict mon. category.

$\rightarrow$  What about commutativity?  $\gamma^2 \neq 1 \Rightarrow$  not symmetric.

The addition is commutative and one can define a natural trafo  $\delta_{m,n} : m+n \rightarrow n+m$  just crossing  $m$  strings over  $n$  strings.  $\rightarrow$  we have a braiding  $\gamma$



## § 5. Braided groups & The Jones Polynomial

### § 5.1 Von Neumann algebras

Def. Given a commutative ring  $R$ , an  $R$ -algebra is an  $R$ -module with an associative  $R$ -bilinear operation which also contains a multiplicative identity. It can be viewed as a monoid in  $\langle R\text{-mod}, \otimes, R \rangle$ .

Def. A free  $R$ -algebra on a set  $E$  is an algebra of polynomials with coefficients in  $R$  and noncommuting indeterminates taken from the set  $E$ .

Consider a sequence  $\{A_n\}_{n=2}^{\infty}$  of free additive algebras on multiplicative generators  $\{e_i\}_{i=1}^{n-1}$  (for each  $A_n$ ), i.e. modules over the ring  $\mathbb{C}[\tau, \tau^{-1}]$ , where  $\tau$  as in the relations :

1.  $e_i^2 = e_i$
2.  $e_i e_{i+1} e_i = \tau e_i$
3.  $e_i e_j = e_j e_i \quad |i-j| > 2$

von Neumann  
algebras

and the following tower generated by inclusions of algebras as follows:

$M_0 := N$ ,  $M_1 := \mathbb{C}$ ,  $M_2 := \langle M_1, e_1 \rangle$  where  $e_1: M_1 \rightarrow M_0$  is the projection &  $\langle M_1, e_1 \rangle$  denotes the algebra generated by  $M_1$  plus  $e_1$  (allowing 1 gen.) -  
 $\Rightarrow M_0 \subset M_1 \subset M_2 \subset M_3 \subset \dots \subset M_n \subset M_{n+1} \subset \dots$

$$e_i: M_i \rightarrow M_{i-1} \quad \forall i \quad \& \quad M_{i+1} := \langle M_i, e_i \rangle.$$

### § 5.2 Jones algebras

Jones constructed such a tower of algebras with  $\{e_i\}$ :

$$e_i e_{i+1} e_i = \tau e_i \quad \forall i \quad \& \quad e_i e_j = e_j e_i \quad \forall i, j : |i-j| > 1$$

where  $\tau := [M_1 : M_0]$  (complex) -

Moreover, he defined a trace  $\text{tr}: M_n \rightarrow \mathbb{C}$  satisfying the Hurkacz Property:  $\text{tr}(w e_i) = \tau \cdot \text{tr}(w) \quad \forall w \in \langle M_0, e_1, \dots, e_{i-1} \rangle$   $\subseteq$  span symbol

Def. A function  $f: \mathcal{A} \rightarrow \mathbb{C}$ , where  $\mathcal{A}$  is an algebra, is called trace iff  $\text{tr}(ab) = \text{tr}(ba) \quad \forall a, b \in \mathcal{A}$

Example: The ordinary trace of matrices is indeed a trace -

So far we have then the following:

$$\begin{aligned} e_i^2 &= e_i \\ e_i e_{i+1} e_i &= e_i \\ e_i e_j &= e_j e_i \quad |i-j| > 1 \\ \text{gen.: } e_1, \dots, e_{n-1} \end{aligned}$$

Jones Algebra  $\mathcal{A}_n$

$$\begin{aligned} \overline{e_i} \overline{e_i}^{-1} &= 1 \quad (\text{it's a group}) \\ \overline{e_i} \overline{e_{i+1}} \overline{e_i} &= \overline{e_{i+1}} \overline{e_i} \overline{e_{i+1}} \\ \overline{e_i} \overline{e_j} &= \overline{e_j} \overline{e_i} \\ \text{gen.: } \overline{e_1}^{\pm 1}, \dots, \overline{e_{n-1}}^{\pm 1} \end{aligned}$$

$n$ -th braid group  $B_n$

### § 5.3 Jones construction of a representation of $B_n$

Def. Given a group  $G$ , a representation of  $G$  is a finite dimensional complex vector space  $E^G$  along with a morphism  $\rho: G \rightarrow \text{GL}_n(E)$    
 $\text{GL}_n(E) = \{ \text{C-linear maps } \Phi: E^n \rightarrow E^n \}$ . It is denoted by  $(E, \rho)$ ,  $E$  is called support of the representation.

→ In particular, if  $G \subseteq \text{GL}_n(E) \Rightarrow$  the representation induced by  $i: G \hookrightarrow \text{GL}_n(E)$  is called fundamental representation.

Consider the following representation  $g_a(\overline{e_i}) \mapsto a e_{i+1}$ ,  $g_b: B_n \rightarrow \mathcal{A}_n$  for some  $a, b$  appropriately chosen. We already defined  $\text{tr}: \mathcal{A}_n \rightarrow \mathbb{C}[t, t^{-1}]$  a trace on  $\mathcal{A}_n \Rightarrow$  it is natural to compose them, i.e.:

$$\text{tr} \circ g_b: B_n \rightarrow \mathbb{C}[t, t^{-1}]$$

We want to prove this gives indeed the Jones polynomial  $V_k(t)$  upon normalization.

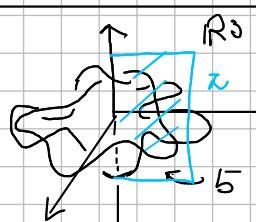
Remark: We know it is a braid isotopy invariant for oriented links.

Thanks to theorems by Harer & Alexander, we will show  $V_k(t)$  is not only well-defined on braids, but for arbitrary knots and links.

## § 5.4 Alexander's & Markov's Theorems

Def. Let  $b \in B_n$  for some  $n \in \mathbb{N}^{\geq 0}$  a braid, then its closure  $\bar{b}$  is obtaining by gluing together the initial and terminal points of  $b$  in order.

ALEXANDER'S THEOREM: Each link in 3D-space is ambient isotopic to a link in the form of a closed braid.



It is always possible to arrange  $\bar{b}$  in such a way that  $\exists$  an axis through the center as in the figure on the left -

If we take a plane  $\pi$ , it is rotated  $\odot$  or  $\odot$  wrt the vertical axis, as we follow the orientation on  $\bar{b}$ .

It is also apparent that  $\# \pi \cap \bar{b} = \# \text{ of strands} = n$ .

Now follow along the knot or link and when the strand starts circulating incorrectly, just throw it over the axis. The knot/link will be a braid  $b$  after this procedure.  $\square$

How many different ways are there to represent a link as closed braid?

i.e. to modify  $b, \tilde{b}$  in such a way that their closures are ambient isotopic links?

(1) Using the Markov move: roll a generator  $B_n \rightarrow B_{n+1}$  and close  $b^+$

$$b \rightarrow b^+$$

More specifically, if  $\beta \in B_n \rightarrow \beta, B_{n+1}, B_{n+1}^{-1}$  are such that  $\bar{\beta} \sim \bar{\beta}_{n+1} \sim \bar{\beta}_{n+1}^{-1}$  where  $\sim$  means ambient isotopic.  $\checkmark$

(2) Compose the conjugate braid  $g\beta g^{-1}$  for some  $g \in B_n$  and take the closure

$$g\beta g^{-1} \Rightarrow \bar{g}\beta g^{-1} \sim \bar{\beta}. \quad \checkmark$$

MARKOV'S THEOREM: Let  $\beta_n \in B_n$  and  $\beta'_m \in B_m \Rightarrow L := \bar{\beta}_n, L' := \bar{\beta}'_m$  are ambient isotopic  $\Leftrightarrow \beta'_m$  is obtained from  $\beta_n$ , via a series of:

- a) equivalences
- b) conjugations
- c) Markov moves

Using this theorem, we can use the presentations of the braid group to extract topological information about knots and links.

Let  $R$  be a com. ring and  $\{J_n : B_n \rightarrow R\}_{n \geq 2}$  family of functions  $\Rightarrow$   
by Markov's theorem, they can be used to construct link invariants if the  
following holds:

$$(i) b \sim b' \Rightarrow J_n(b) = J_n(b') \quad \forall b, b' \in B_n$$

$$(ii) gbg \in B_n \Rightarrow J_n(b) = J_n(gbg^{-1})$$

$$(iii) b \in B_n \Rightarrow \exists \alpha \in R (\text{if } n) : J_{n+1}(b\bar{\tau}_n) = \alpha^{+1} J_n(b)$$

$$J_{n+1}(b\bar{\tau}_n^{-1}) = \alpha^{-1} J_n(b)$$

Remark: for  $\overline{b'} = \overline{bb_n}$ , the Markov move  $b \rightarrow b' \leftrightarrow \text{RI on } \overline{b}$ .  
eq.

( $\overline{b\tau_n}$  RI. positive,  $\overline{b\tau_n^{-1}}$  RI. negative)

Def. let  $b = \overline{a_1} \dots \overline{a_k}$  be a braid, then we define its write as

$$\sum_{k=1}^K a_k =: w(b)$$

It is apparent that  $w(b) = w(\overline{b})$ .

Def. Let  $\{J_n\}_{n \geq 2}$  be the family of functions above, we call it  
a Markov trace on  $\{B_n\}$ .

Def. Let  $L \sim \overline{b}, b \in B_n$  for any link  $L$  (Alexander's theorem)  
existence of  $\overline{b}$  follows by

$\Rightarrow$  we define the link invariant  $J(L)$  for the Mar. trace  $\{J_n\}$  the  
following:

$$R \ni J(L) = \alpha^{-w(b)} J_n(b)$$

Notice that the function  $J_n$  is fixed by the choice of the representative  
 $b \in B_n$  for the link.

Proposition: If  $L \sim L' \Rightarrow J(L) = J(L')$

Let  $L \sim \overline{b}, L' \sim \overline{b'}$   $b, b' \in B_n$  (via Alexander's theorem as usual)  
as  $L \cup L'$  by hypothesis,  $\overline{b} \sim \overline{b'}$  via a sequence of moves of type a), b), c)  
in Markov's theorem  $\Rightarrow w(\overline{b}) = w(\overline{b'}) \Rightarrow J(L) = J(L')$   $\square$

## S 6. Bracket For Drawers

Def. Let  $\langle \rangle : B_n \rightarrow \mathbb{Z}[A, A^{-1}]$  be defined as  $\langle b \rangle = \overline{\langle \bar{b} \rangle}$  the bracket for braids or equivalently, if  $J_n : B_n \rightarrow \mathbb{Z}[A, A^{-1}]$ , then  $J_n(b) = \overline{\langle b \rangle}$ .  
 $\rightarrow J_n$  is a Markov trace with  $\alpha = -A^3$ .

## S. 6.1 states of a braid

Remember  $\langle \cdot \rangle = A \langle \cdot \rangle + A^{-1} \langle \cdot \rangle$  & generalise

$$\begin{aligned} \downarrow \langle \| \dots | y_1 | \dots \| \rangle &= A \langle \| \dots | \rangle + A^{-1} \langle \| \dots | \underbrace{y_i}_{\text{ith position}} | \dots | \rangle \\ \Rightarrow \langle \tilde{v}_i \rangle &= A \langle 1_n \rangle + A^{-1} \underbrace{\langle U_i \rangle}_{(i+1)th} \end{aligned}$$

$v_i$  is a new element in the braid, a composition of cup  $\cup$  and cap  $\cap$  at the  $i$ -th &  $(i+1)$ -th strand

$\Rightarrow$  Each state of  $S$  can be written as the closure of products of  $U_i$

Define  $\mathbf{G}_i := \mathbf{A} + \mathbf{A}^{-1} \mathbf{U}_i$ ,  $\mathbf{F}_i^{-1} := \mathbf{A}^{-1} + \mathbf{A} \mathbf{U}_i \Rightarrow \mathbf{b} = \mathbf{U}(\mathbf{b})$

just inserting for each generator  $b_i$  of the braid word, the new definitions in terms of the  $U_i$  ✓

$$\Rightarrow \langle b \rangle = \langle u(b) \rangle = \langle u(\bar{b}) \rangle = \langle \bar{b} \rangle = \delta^{\parallel u \parallel} \quad \parallel u \parallel = \# \text{ loops in } \bar{u} - 1$$

$$\delta = -\left(\frac{1}{A^2} + A^2\right)$$

$$\Rightarrow u(b) = \sum_s \langle b | s \rangle u_s$$

where  $\sigma$  indexes all terms in the products and  $\langle b/s \rangle$  is as usual the product of the labels  $A, A^{-1}$ .

$$\Rightarrow \langle u(b) \rangle = \langle \sum_s \langle b|s \rangle u_s \rangle = \sum_s \langle b|s \rangle \langle u_s \rangle = \sum_s \langle b|s \rangle \quad \mathcal{S}^{\parallel s \parallel}$$

§ 6.2 Temperley-Lieb algebra  $\mathcal{A}_n = TL_n(\delta) \leftarrow$  statistical mechanics

In terms of  $U_i$ , we have a free multiplicative algebra on  $U_1, \dots, U_{n-1}$  multiplicative generators and relations, called Temperley Lieb algebra  $\mathcal{D}_n$ .  $\mathcal{D}_n$  is a  $\mathbb{Z}[A, A^{-1}]$ -module with  $\delta = -\left(\frac{1}{A^2} + A^2\right) \in \mathbb{Z}[A, A^{-1}]$  loop value.

The relation for an arc

$$[\alpha] \quad \left\{ \begin{array}{l} U_i U_{i+1} U_i = U_i \\ U_i^2 = \delta U_i \\ U_i U_j = U_j U_i \quad |i-j| > 1 \end{array} \right.$$

Visualization:

$$\underbrace{\text{U}}_{\cap} \sim \frac{\text{U}}{\cap} \quad U_1 = U_1 U_2 U_1$$

$$\underbrace{\text{U}}_{\cap} \sim \circ \frac{\text{U}}{\cap} \quad U_1^2 = \delta U_1$$

$$\underbrace{\text{U}}_{\cap} \underbrace{\text{U}}_{\cap} \sim \frac{\text{U}}{\cap} \underbrace{\text{U}}_{\cap} \quad U_2 U_3 = U_3 U_2$$

→ Remember Jones algebra? Now define

$$e_i := \frac{U_i}{\delta} + i$$

$$\Rightarrow \begin{cases} e_i^2 = e_i \\ e_i e_{i+1} e_i = \bar{c} e_i \quad \text{with } \bar{c} = \frac{1}{\delta^2} \end{cases}$$

from the bracket for braids we recovered the Jones algebra!

!!

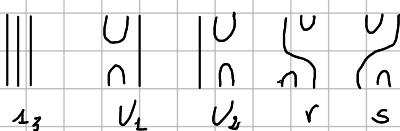
### § 6.3 $\mathcal{D}_n(\delta)$ as diagram monoid on $n$ points $\mathbb{D}_n$

Def. let  $\mathbb{D}_n$  denote the collection of all topological equivalence classes of diagrams obtained by connecting pairs of points in two parallel rows of  $n$  points.

The arc connecting these points must be such that:

- (a) All arcs are drawn between the two rows;
- (b) No two arcs cross one another;
- (c)  $a, b \in \mathbb{D}_n$  are equivalent if  $\exists$  a planar isotopy between them.

$\mathbb{D}_n$  is called diagram monoid on  $n$  points.

Example:  elements of  $\mathbb{D}_3$  (6 points)

$(\mathbb{D}_n, \cdot)$  comes equipped with a multiplication. If  $a, b \in \mathbb{D}_n$   $a \cdot b = \delta_c^b$  for some  $c \in \mathbb{D}_n$ ,

↪ # of loops in the  $a \cdot b$ .

$$\text{Example: } \xrightarrow{rs} \underbrace{\text{U}}_{\cap} = \delta \frac{\text{U}}{\cap} = \delta U_2 \quad \checkmark$$

Proposition: The elements  $\{1, U_1, \dots, U_{n-1}\}$  generate  $\mathbb{D}_n$ . If  $\mathbb{D}_n \ni x \sim p := \prod U_i$   $\sim q := \prod U_j$   $\Rightarrow p$  is obtained by applying the relations in [6].

## §7. The Jones Polynomial revealed

Finally we have all the elements in order to investigate the true nature of  $V_k(t)$  -

- Define a map  $\rho_n : B_n \rightarrow TL_n(\delta)$  via
 
$$\begin{cases} \rho(\sigma_i) = A + A^{-1}U; \\ \rho(\sigma_i^{-1}) = A^{-1} + AU; \end{cases}$$
 (in fixed by choice of  $B_n$ )

We know already that for  $b \in B_n$ ,  $\langle \bar{b} \rangle = \sum_s \langle b|_s \rangle \langle \bar{u}_s \rangle = \langle b \rangle$  and

$$\rho(b) = \sum_s \langle b|_s \rangle \langle u_s \rangle \text{ as } \rho(xy) = \rho(x) \cdot \rho(y) \quad \forall x, y \in B_n.$$

$$\langle \bar{u}_s \rangle := \# \text{ loops in } \bar{u}_s - 1$$

- Define  $\text{tr} : TL_n(\delta) \rightarrow \mathbb{Z}[A, A^{-1}]$  realisation of Jones trace on the von Neumann algebra  $A_n$  (by loop counts)

$$U \rightarrow \text{tr}(U) = \langle \bar{U} \rangle$$

for  $U \in \mathbb{D}_n$  and extend it by linearity to  $TL_n(\delta)$ .

$$\Rightarrow \boxed{\langle b \rangle = \text{tr}(\rho(b))}$$

brace

Jones Polynomial

Conclusion:  $V_k(t)$ , the Jones polynomial in one variable is a trace on a representation of the braid group  $B_n$  to the Temperley-Lieb algebra  $TL_n(\delta)$ .

To prove (exercise): The map  $\rho : B_n \rightarrow TL_n(\delta)$  defined above is indeed a representation, i.e. is a morphism:

$$(i) \quad \rho(\sigma_i \sigma_i^{-1}) = \rho(\sigma_i) \rho(\sigma_i^{-1}) = \rho(\sigma_i) \rho(\sigma_i)^{-1} = \text{id} \quad \forall i$$

$$(ii) \quad \rho(\sigma_i \sigma_j) = \rho(\sigma_j \sigma_i) \quad \forall i, j : |i-j| > 1$$

$$(iii) \quad \rho(\sigma_i \sigma_{i+1} \sigma_i) = \rho(\sigma_{i+1} \sigma_i \sigma_{i+1}) \quad \forall i$$