A farmer with a wolf, a goat, and a cabbage must cross a river by boat. The boat can carry only the farmer and a single item. If left unattended together, the wolf would eat the goat, or the goat would eat the cabbage. How can they cross the river without anything being eaten? https://*en.m.wikipedia.org*/wiki/*Wolf*,_*goat*_*and*_*cabbage*_*problem*

– This spec tought me the difference between CHOICE and with: https://*groups.google.com*/g/tlaplus/c/*QJAulqYgC*0*E*/*m*/*Cm*2*PPJBDCQAJ*
> CHOOSE is definitely a confusing operator, many people have problems with it including myself.
> The idea is that CHOOSE always picks the same value, it does not represent arbitrary values.
> If you're trying to say "sub can be any element of *ClaimsData*", then the "there exists"
> operator (existential quantification) is what you're looking for, *i.e.*
> ∃ sub ∈ *ClaimsData*: *SomePredicate*(*sub*)

And: > What if multiple values satisfy CHOOSE ? In this case the only requirement is that the result is deterministic: the engine must always return the same value, no matter what. In practice this means that *TLC* will always choose the lowest value that matches the set. https://*learntla.com*/core/*operators.html* ≠ choose

CHOOSE has to be deterministic but in my previous implementation (in last commit) I was trying to use CHOOSE to pick any element in subset of side such that removing that item from side would have left it in a valid state. Because I was using CHOOSE , tlc was always picking the same value leading to a very small possible state transitions and a lot of headache to me.

MODULE *wolf*_*goat*_*cabbage*

LOCAL INSTANCE *TLC*
LOCAL INSTANCE *FiniteSets*
LOCAL INSTANCE *Integers*

$WOLF \triangleq$ "Wolf"
$GOAT \triangleq$ "Goat"
$CABBAGE \triangleq$ "Cabbage"
$FinalResult \triangleq \{WOLF, GOAT, CABBAGE\}$
$InvalidStates \triangleq \{\{CABBAGE, GOAT\}, \{WOLF, GOAT\}\}$
$baitinv \triangleq$ TRUE
$baitinv \triangleq TLCGet(\text{“}level\text{”}) < 14$

--algorithm *wolf*_*goat*_*cabbage*{

variables $side\_start = FinalResult$, $side\_end = \{\}$;

define {
$IsValidState(side) \triangleq \neg(side \in InvalidStates)$
$Inv \triangleq side\_end \neq FinalResult$
this operator is used select an item
from side such that removing this item
will leave "side" in a valid state.
$ValidSubsets(side) \triangleq \{s \in \text{SUBSET } (side) : IsValidState(side\_start \setminus s) \wedge Cardinality(s) = 1\}$
}

macro *PickFrom*(*side*, *other*_*side*_*is*_*valid*) {
await $\wedge\ transport = \{\}$
$\wedge\ side \neq \{\} \wedge other\_side\_is\_valid$;

```
      choose an item such that this side is left valid:
    with ( item ∈ ValidSubsets(side) ) {
          transport := item ;
          side := side \ transport ;
     }
 }

macro DropItemTo( side ) {
      leave an item to side_start side. If needed to avoid conflicts, load the other item.
    await transport ≠ {} ;
    side := side ∪ transport ;
    transport := {} ;
 }

process ( Farmer = 1 )
variable transport = {} ; {
W:
    while ( TRUE ) {
        either {
              pick an item from side_start and load it.
            PickFrom(side_start, IsValidState(side_end)) ;
         } or {
              pick an item from side_start and load it.
            PickFrom(side_end, IsValidState(side_start)) ;
         } or {
            DropItemTo(side_start) ;
         } or {
            DropItemTo(side_end) ;
         } ;
     }
 }
 }
  BEGIN TRANSLATION (chksum(pcal) = "7c28162a" ∧ chksum(tla) = "b19f62f8")
VARIABLES side_start, side_end

  define statement
IsValidState(side) ≜ ¬(side ∈ InvalidStates)
Inv ≜ side_end ≠ FinalResult


ValidSubsets(side) ≜ {s ∈ SUBSET (side) : IsValidState(side_start \ s) ∧ Cardinality(s) = 1}

VARIABLE transport

vars ≜ ⟨side_start, side_end, transport⟩

ProcSet ≜ {1}
```

$$
\begin{array}{lll}
\mathit{Init} & \triangleq & \text{Global variables} \\
 & & \land\ \mathit{side\_start} = \mathit{FinalResult} \\
 & & \land\ \mathit{side\_end} = \{\} \\
 & & \text{Process } \mathit{Farmer} \\
 & & \land\ \mathit{transport} = \{\}
\end{array}
$$

$$
\begin{array}{lll}
\mathit{Farmer} & \triangleq & \lor\ \land\ \land\ \mathit{transport} = \{\} \\
 & & \qquad\ \land\ \mathit{side\_start} \neq \{\} \land (\mathit{IsValidState}(\mathit{side\_end})) \\
 & & \quad\ \land \exists\, \mathit{item} \in \mathit{ValidSubsets}(\mathit{side\_start}) : \\
 & & \qquad\ \land\ \mathit{transport}' = \mathit{item} \\
 & & \qquad\ \land\ \mathit{side\_start}' = \mathit{side\_start} \setminus \mathit{transport}' \\
 & & \quad\ \land \text{UNCHANGED } \mathit{side\_end} \\
 & & \lor\ \land\ \land\ \mathit{transport} = \{\} \\
 & & \qquad\ \land\ \mathit{side\_end} \neq \{\} \land (\mathit{IsValidState}(\mathit{side\_start})) \\
 & & \quad\ \land \exists\, \mathit{item} \in \mathit{ValidSubsets}(\mathit{side\_end}) : \\
 & & \qquad\ \land\ \mathit{transport}' = \mathit{item} \\
 & & \qquad\ \land\ \mathit{side\_end}' = \mathit{side\_end} \setminus \mathit{transport}' \\
 & & \quad\ \land \text{UNCHANGED } \mathit{side\_start} \\
 & & \lor\ \land\ \mathit{transport} \neq \{\} \\
 & & \quad\ \land\ \mathit{side\_start}' = (\mathit{side\_start} \cup \mathit{transport}) \\
 & & \quad\ \land\ \mathit{transport}' = \{\} \\
 & & \quad\ \land \text{UNCHANGED } \mathit{side\_end} \\
 & & \lor\ \land\ \mathit{transport} \neq \{\} \\
 & & \quad\ \land\ \mathit{side\_end}' = (\mathit{side\_end} \cup \mathit{transport}) \\
 & & \quad\ \land\ \mathit{transport}' = \{\} \\
 & & \quad\ \land \text{UNCHANGED } \mathit{side\_start}
\end{array}
$$

$$
\mathit{Next} \triangleq \mathit{Farmer}
$$

$$
\begin{array}{lll}
\mathit{Spec} & \triangleq & \land\ \mathit{Init} \land \Box[\mathit{Next}]_{\mathit{vars}} \\
 & & \land\ \mathrm{WF}_{\mathit{vars}}(\mathit{Next}) \\
 & & \land\ \mathrm{WF}_{\mathit{vars}}(\mathit{Farmer})
\end{array}
$$

END TRANSLATION

$$
\begin{array}{lll}
\mathit{TypeOk} & \triangleq & \land \forall\, \mathit{el} \in \mathit{side\_start} : \mathit{el} \in \{\mathit{WOLF},\ \mathit{GOAT},\ \mathit{CABBAGE}\} \\
 & & \land \forall\, \mathit{el} \in \mathit{side\_end} : \mathit{el} \in \{\mathit{WOLF},\ \mathit{GOAT},\ \mathit{CABBAGE}\} \\
 & & \land \forall\, \mathit{el} \in \mathit{transport} : \mathit{el} \in \{\mathit{WOLF},\ \mathit{GOAT},\ \mathit{CABBAGE}\}
\end{array}
$$