Lecture 6 - Tree (Part 2)

CPE112 - Programming with Data Structures 8 March 2024

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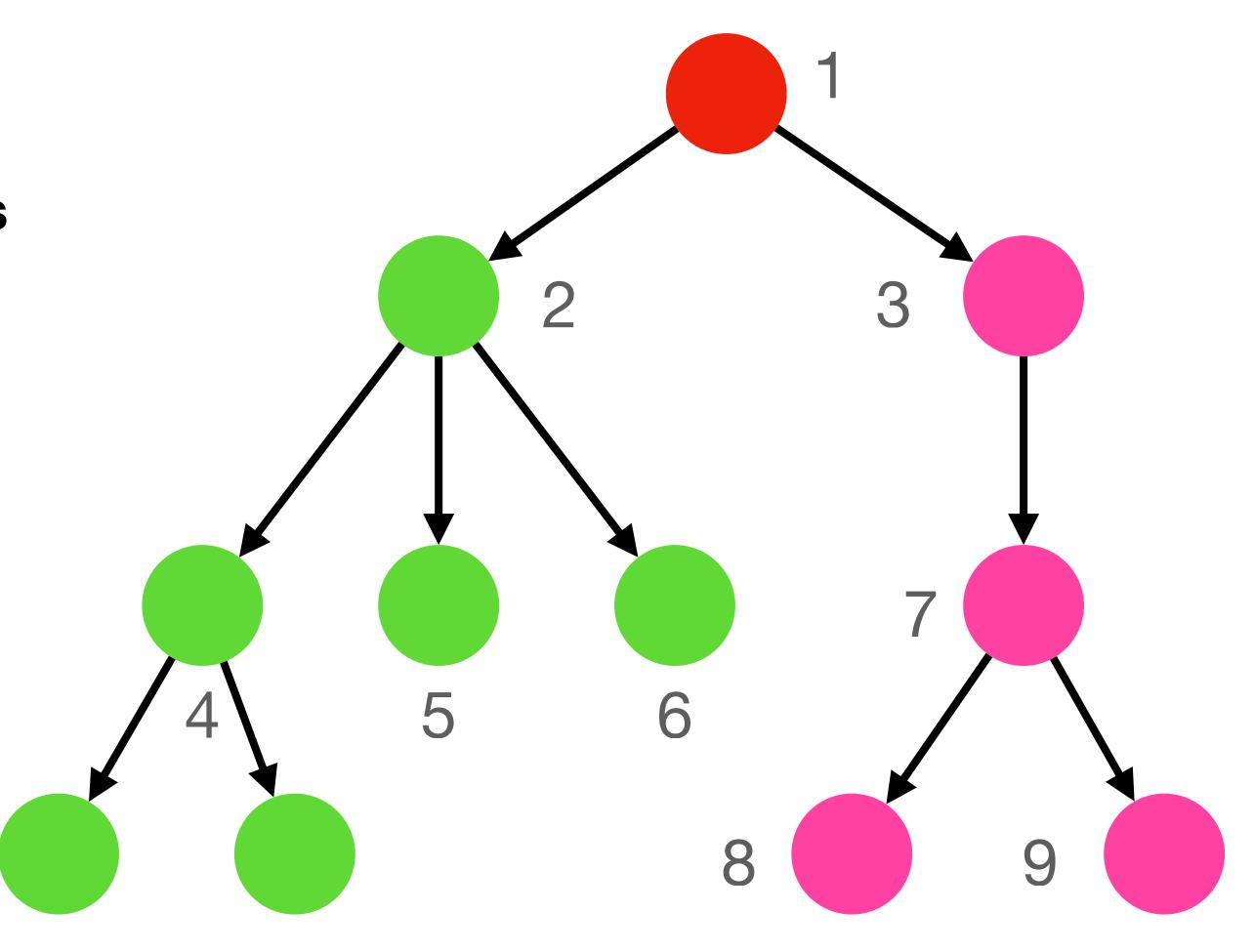
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Review (1)

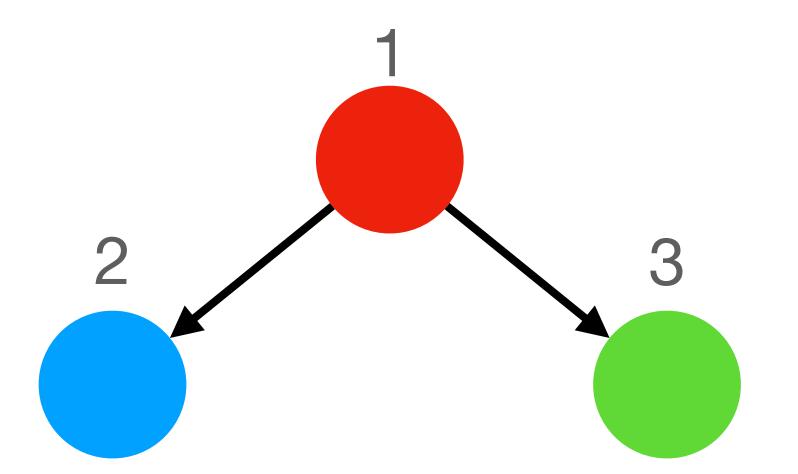
Linear vs Hierarchical Data Structures

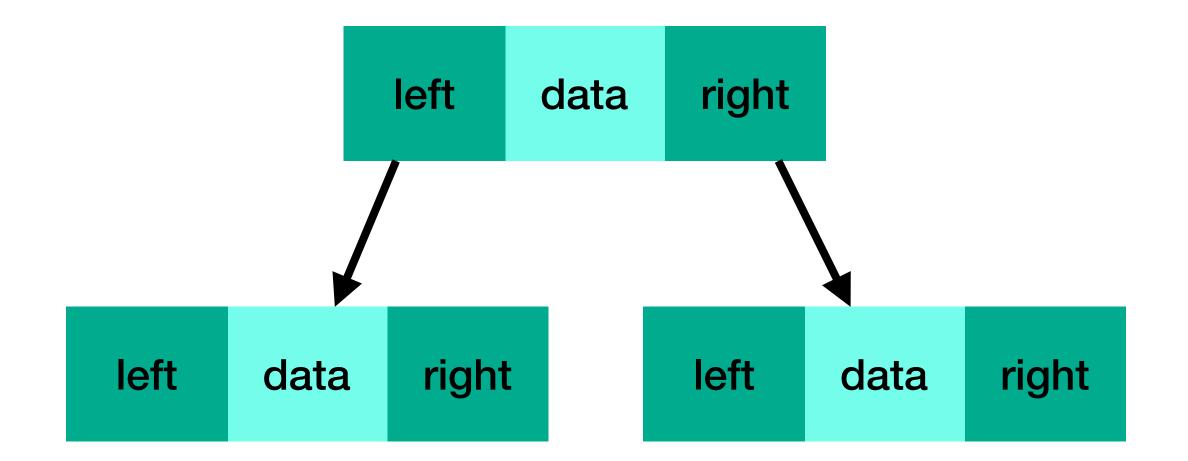
- Tree
 - Node / Edge / Path
 - Parent / Children / Sibling
 - Ancestor / Descendant
 - Root / Leaf
 - Sub-tree
 - Level / Path Length / Height / Depth



Review (2)

- Binary Tree
- Tree traversal
 - Depth-first search: pre-order, in-order, post-order
 - Breadth-first search: Root -> Leaf L -> Leaf R



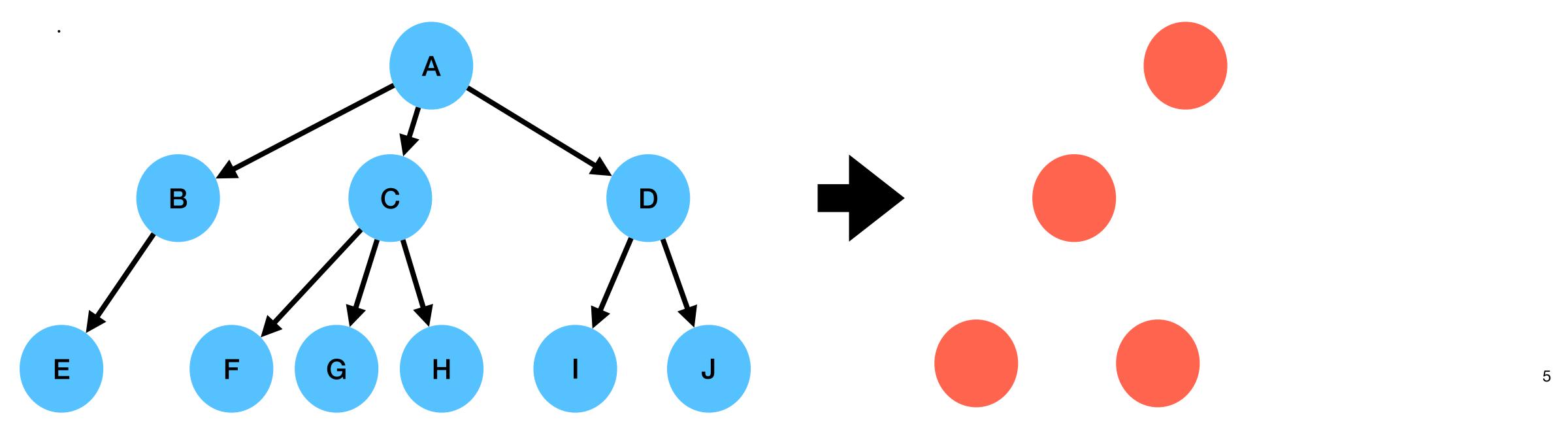


Today's Goal

- General Tree -> Binary Tree
- Binary Search Tree
- AVL Tree

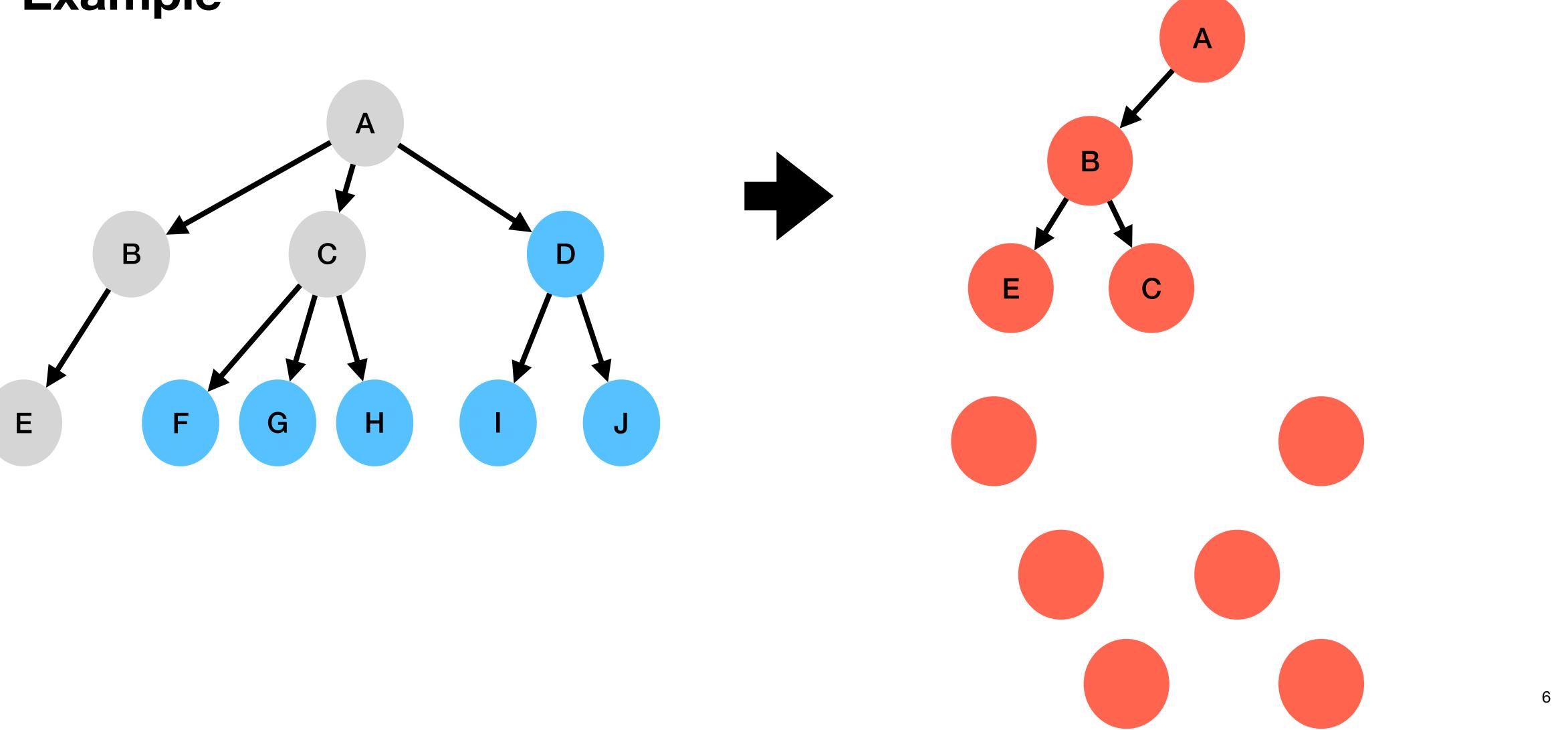
Creating a Binary Tree from a General Tree Rules

- Root of the BT = Root of the GT
- **Left child** of a node in the BT = **Left most child** of the node in the GT
- Right child of a node in the BT = Right sibling of the node in the GT



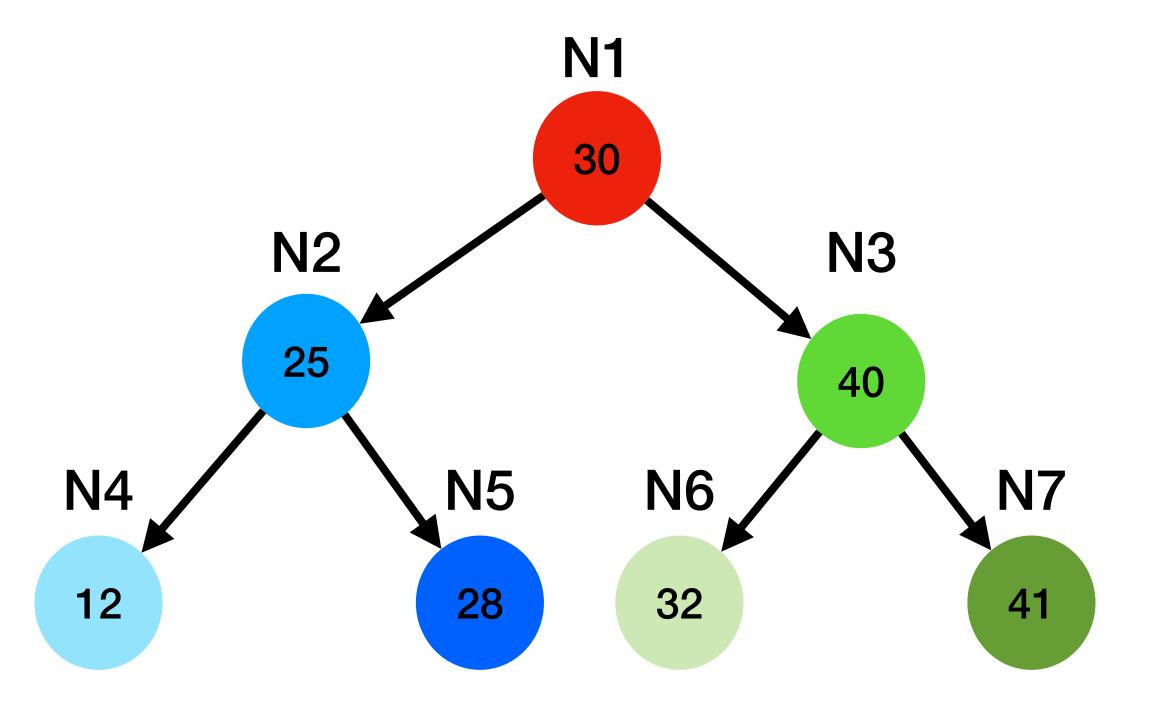
Creating a Binary Tree from a General Tree

Example



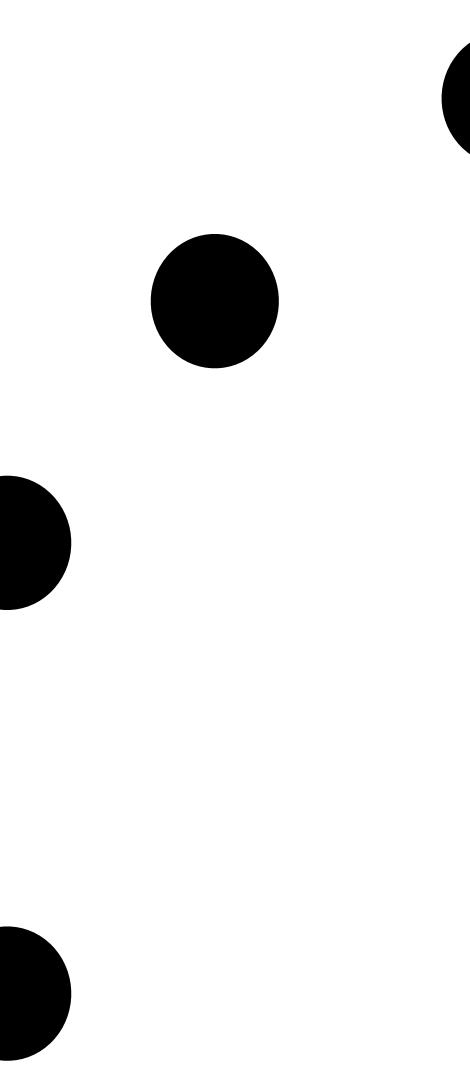
Definition

• If T is a binary search tree and N is a node in T, the value of N must be greater than all nodes in the left sub-tree and less than all nodes in the right sub-tree



Create a binary tree

• 40, 32, 41, 30, 25, 12, 28

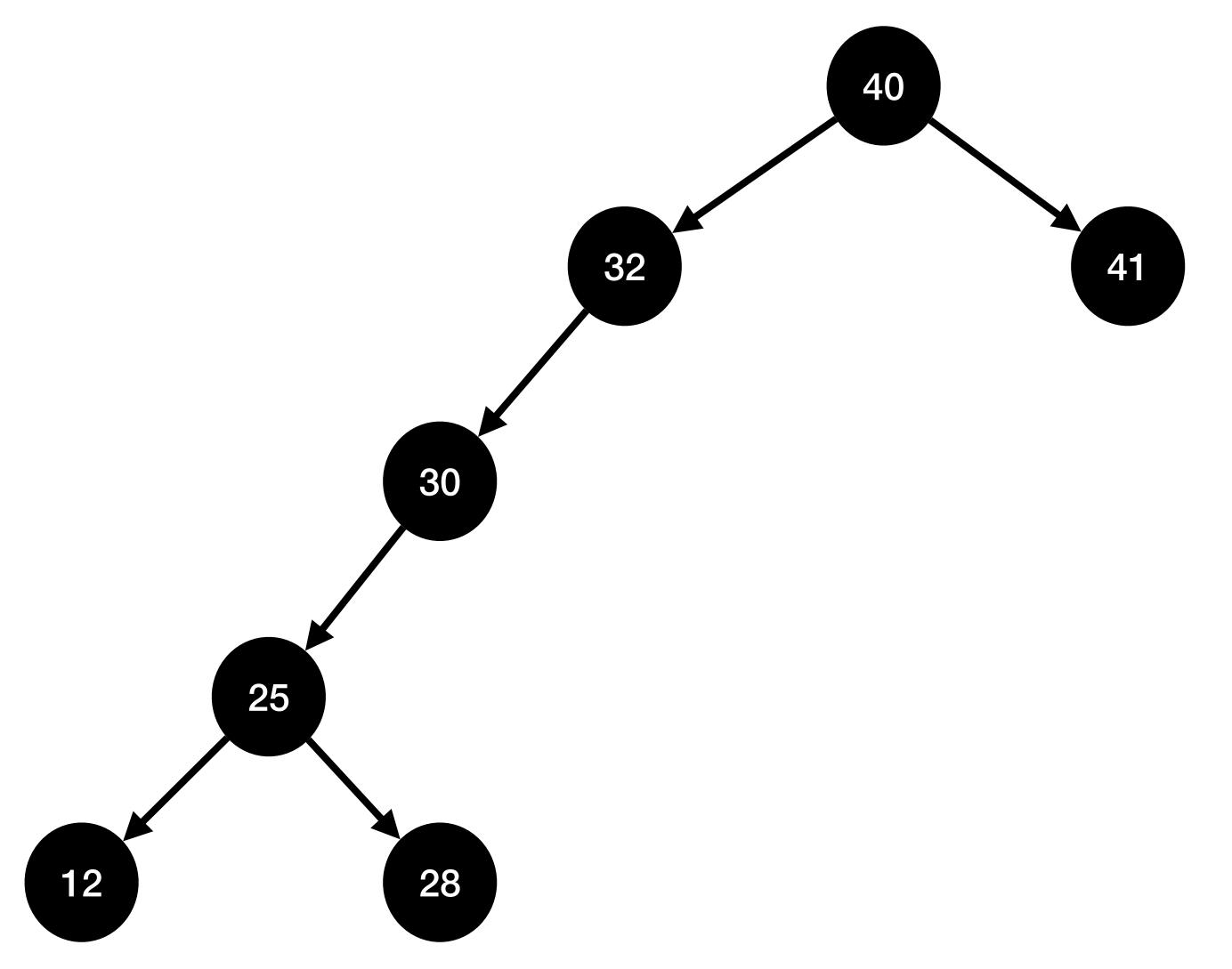


Create a binary tree

```
TREENODE_T* insertNodeToBinaryTree (TREENODE_T* root, TREENODE_T* newNode)
 TREENODE_T *p, *previous;
 //CASE 1: Root is null THEN root = newNode
 //CASE 2: Root is not null THEN
      (a) find a suitable node for insertion (hint: use a while loop then compare values)
      (b) link newNode with the suitable node
 //RETURN Root
```

Search a node

• Find 28



Search a node

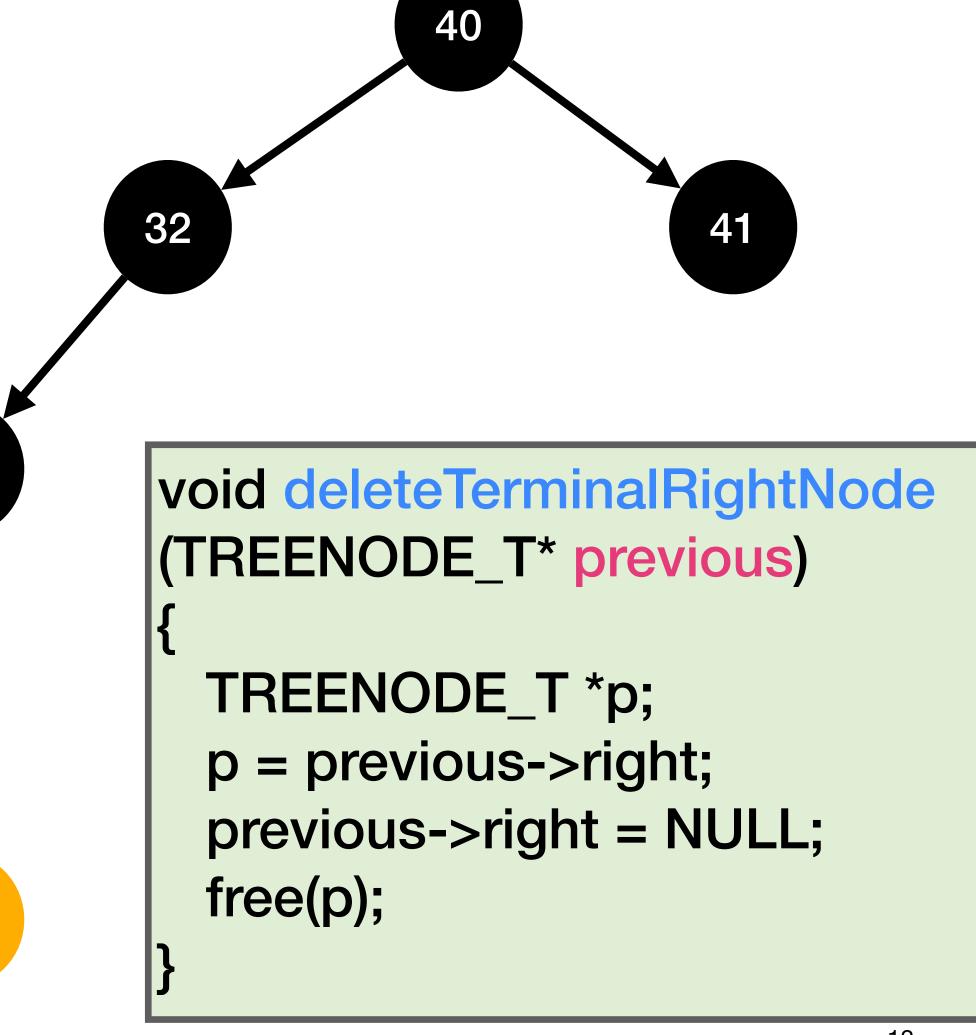
```
TREENODE_T* binarySearch (TREENODE_T* root, int key)
 TREENODE_T *p, *node;
 p = root; node = NULL;
 do {
     if (p->info == key) node = p; /*Search found*/
     else if (p-\sin b) = p-\sec b;
     else p = p->right;
 } while ((p != NULL) && (node == NULL));
 return (node);
```

Delete a node

Delete 28

Need to know the <u>parent node</u>

 Delete a node in any tree is complicated since there are several cases

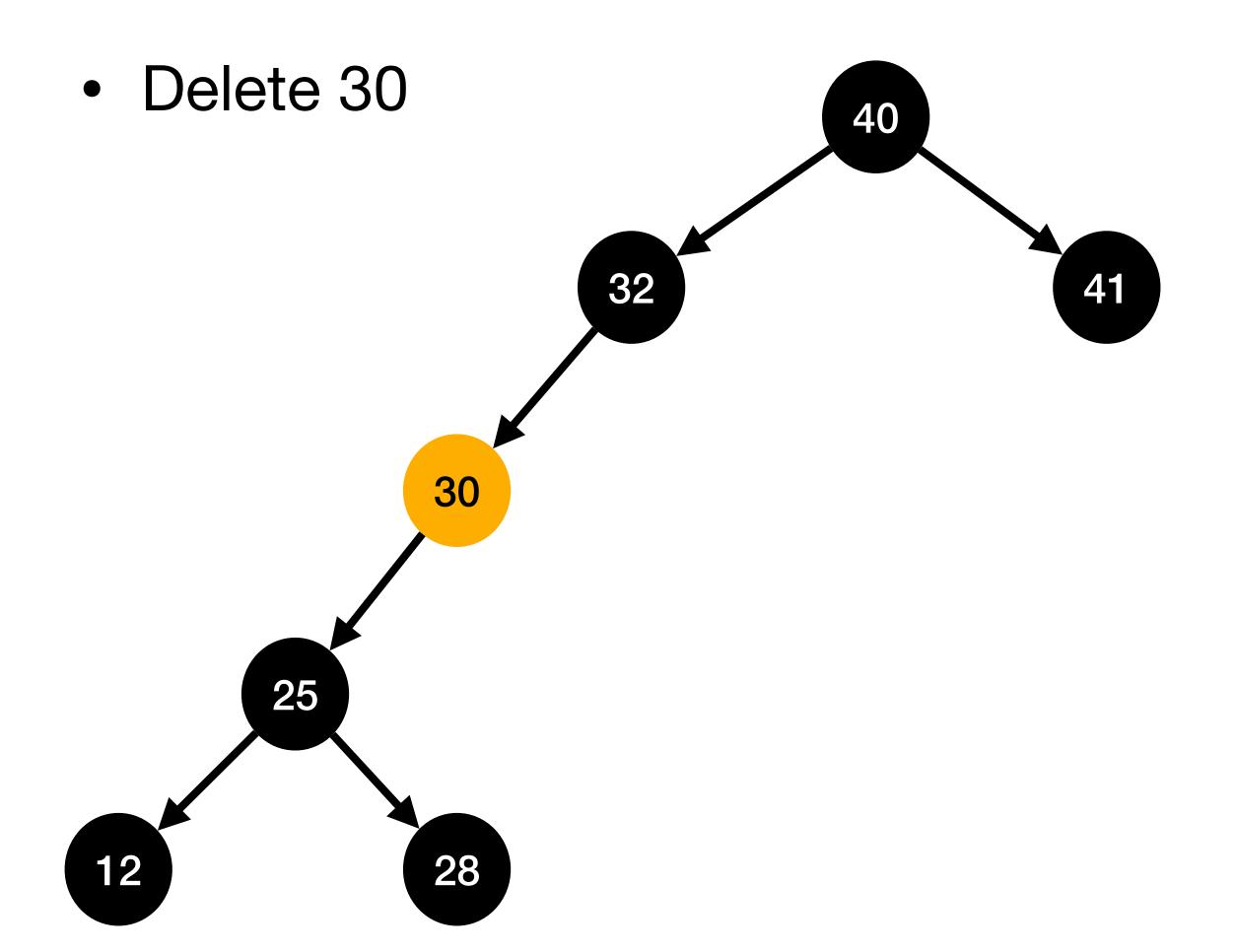


30

28

25

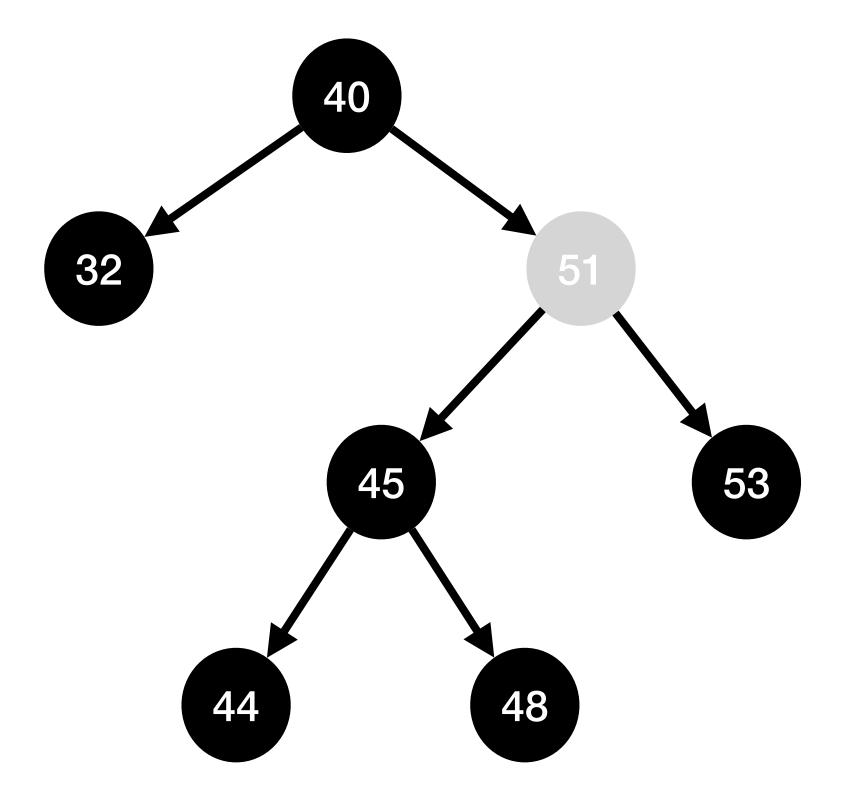
Delete a node



```
void deleteNonTerminalWithOnlyLeftNode
(TREENODE_T* previous)
{
   TREENODE_T *p;
   p = previous->left;
   previous->left = p->left;
   free(p);
}
```

Delete a node

• More case: delete 51



Delete a node

• More case: delete 51

```
40
                 p
                          32
                                 q(1)
q(1)
                    53
    45
```

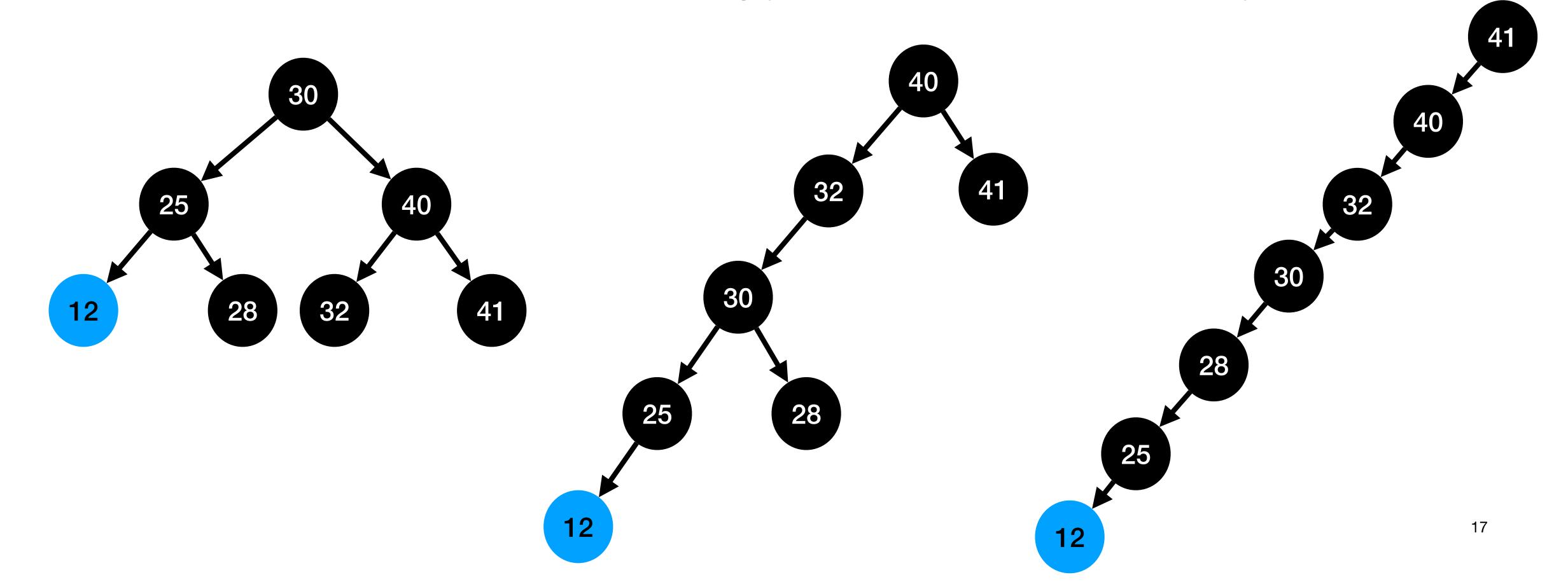
```
void deleteNonTerminalRightNode
(TREENODE_T* previous)
 TREENODE_T *p, *q;
 p = previous->right;
  previous->right = q = p->left;
 while (q->right != NULL)
   q = q->right;
 q->right = p->right;
 free(p);
```

Other operations

- Find the height of a Binary Search Tree
- Find the height of a node
- Find the smallest/largest node
- Delete a Binary Search Tree (delete left sub-tree then right sub-tree)

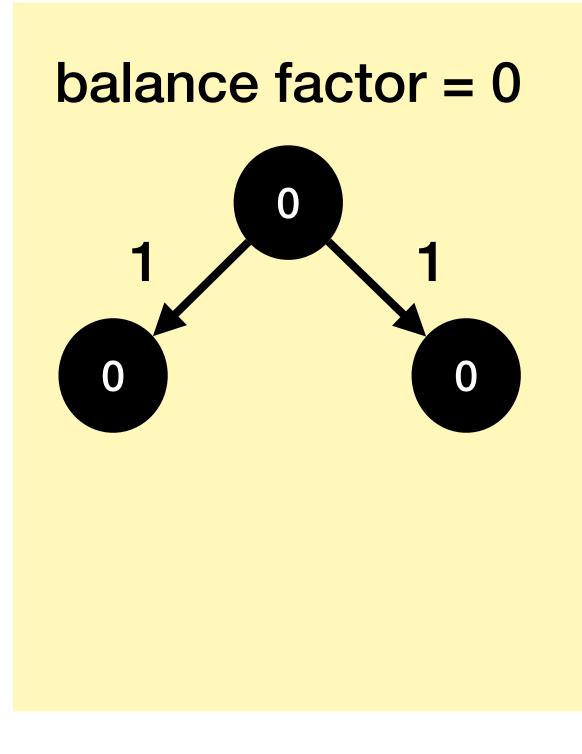
Balance Trees

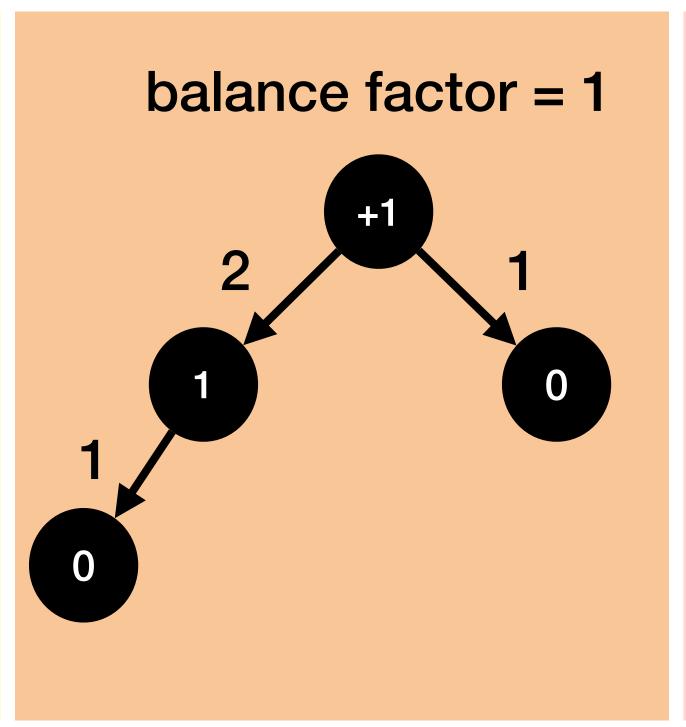
• The sequence of information is strongly affected search efficiency.

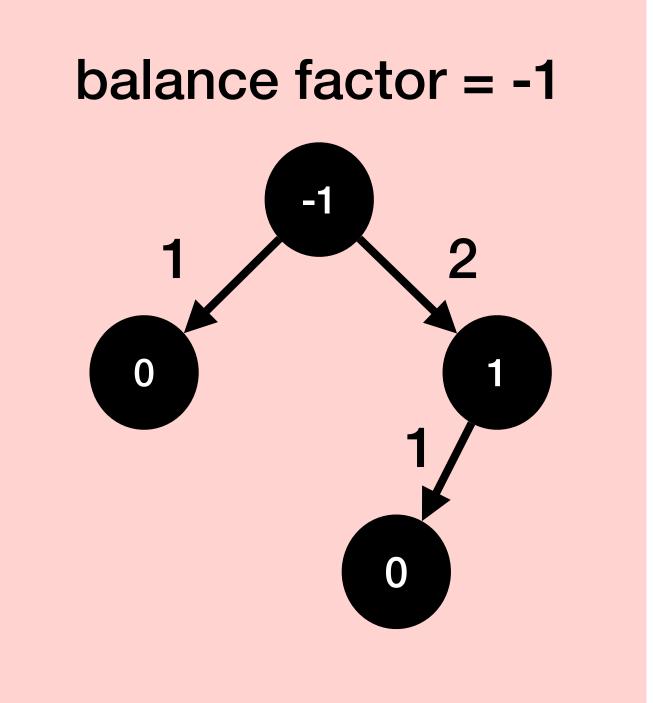


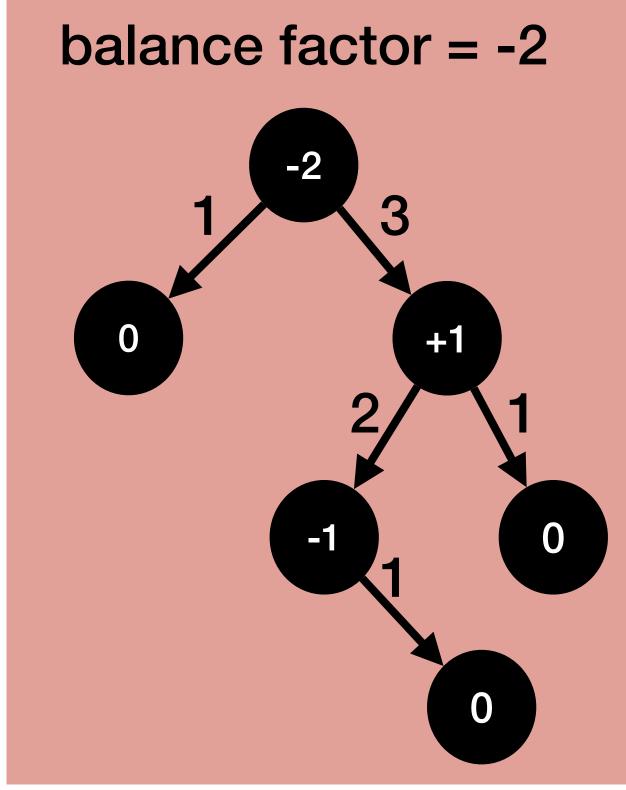
Balance Factor

- Balance factor = Height (left sub-tree) Height (right sub-tree)
- Note: height is the longest path length to the leaf



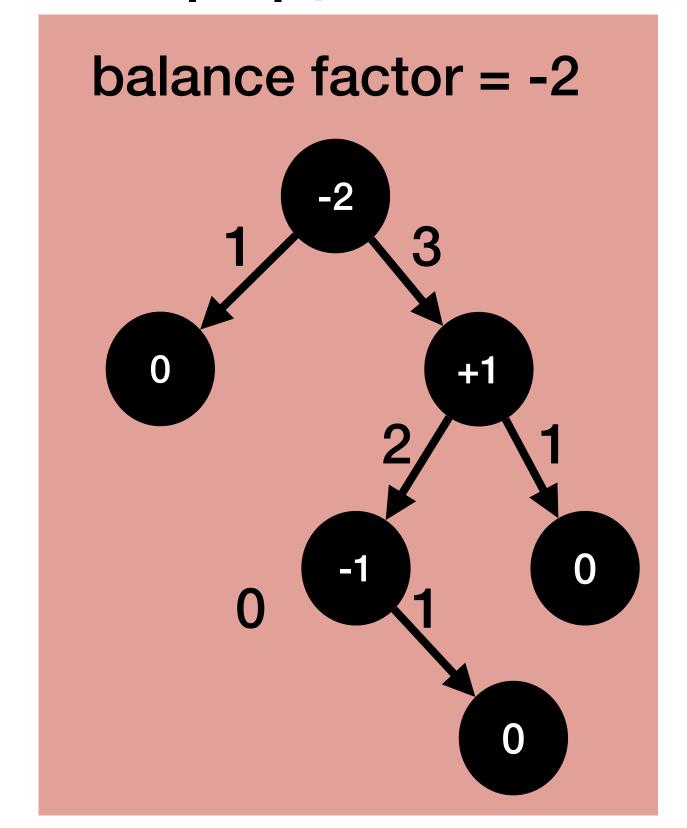






Balanced Binary Search Tree

 A binary search tree is balance if the balance factor of any node is less than | 1 | (balance factor = -1, 0, or 1)



Q: How to find imbalanced points?

A: Need to know the height and balance of each node

```
typedef struct _treenode
{
  int data ;
  struct _treenode* left;
  struct _treenode* right;
} TREENODE_T;
```

Is a tree balanced?

Concept: post-order traversal: LT -> RT -> Root

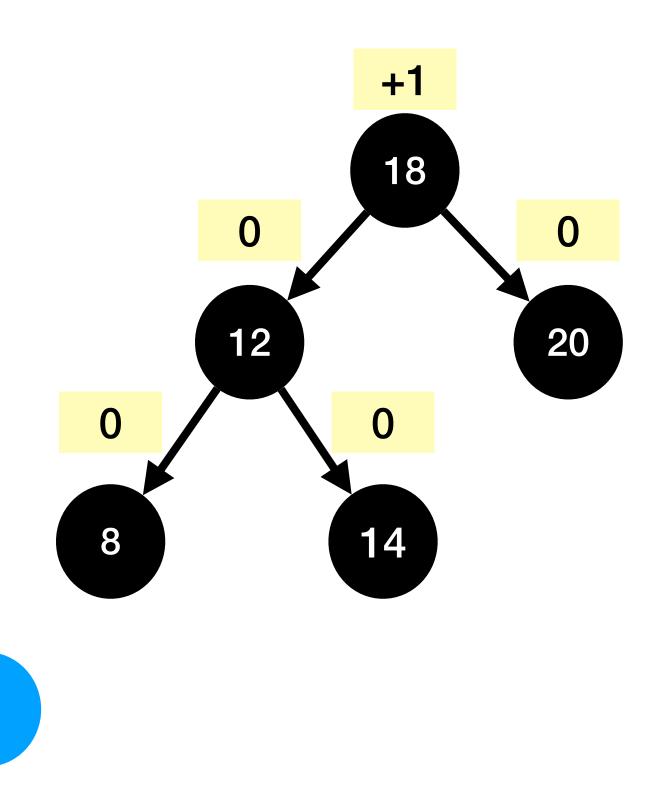
```
void fillHeight (TREENODE_T *node)
 int height = 0;
 if (node->left != NULL) fillHeight(node->left);
 if (node->right != NULL) fillHeight(node->right);
 if (node->left == NULL && node->right == NULL) height = 0; //Case1
 else if (node->left == NULL) height = node->right->height +1; //Case2
 else if (node->right == NULL) height = node->left->height +1; //Case3
 else if (node->right->height > node->left->height) height = node->right->height +1; //Case4
 else height = node->left->height +1; //Case5
 node->height = height;
                                                                                           20
```

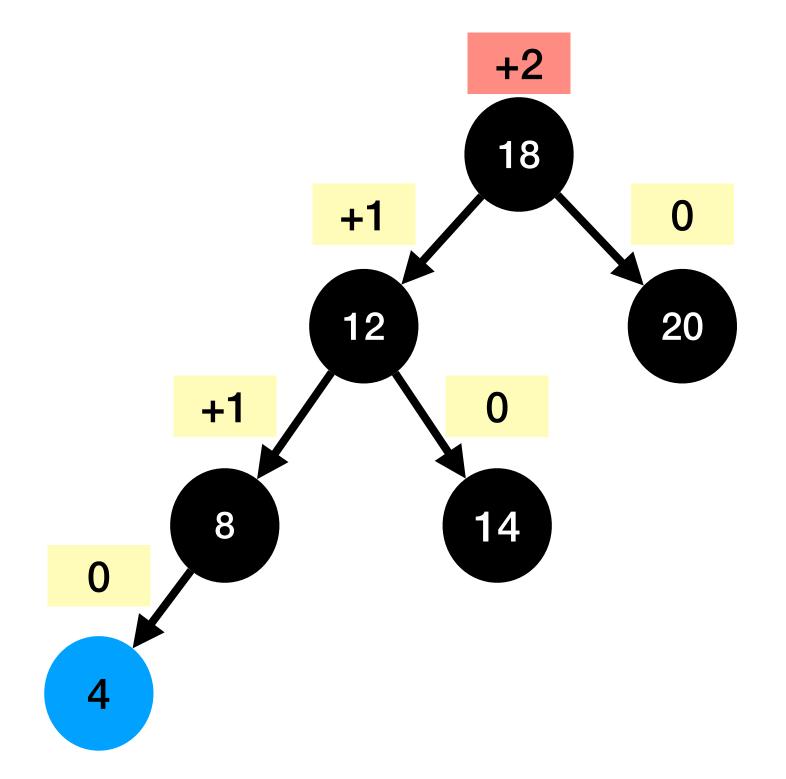
Is a tree balanced?

```
void fillBalanceFactor (TREENODE_T *node)
  int leftHeight = 0, rightHeight = 0;
 if (node->left != NULL) fillBalanceFactor(node->left);
 if (node->right != NULL) fillBalanceFactor(node->right);
 //Get the height of the left sub-tree
 if (node->left == NULL) ...
  else ...
 //Get the height of the right sub-tree
 if (node->right == NULL)
  else ...
 node->balance = ...;
```

Lost Balance

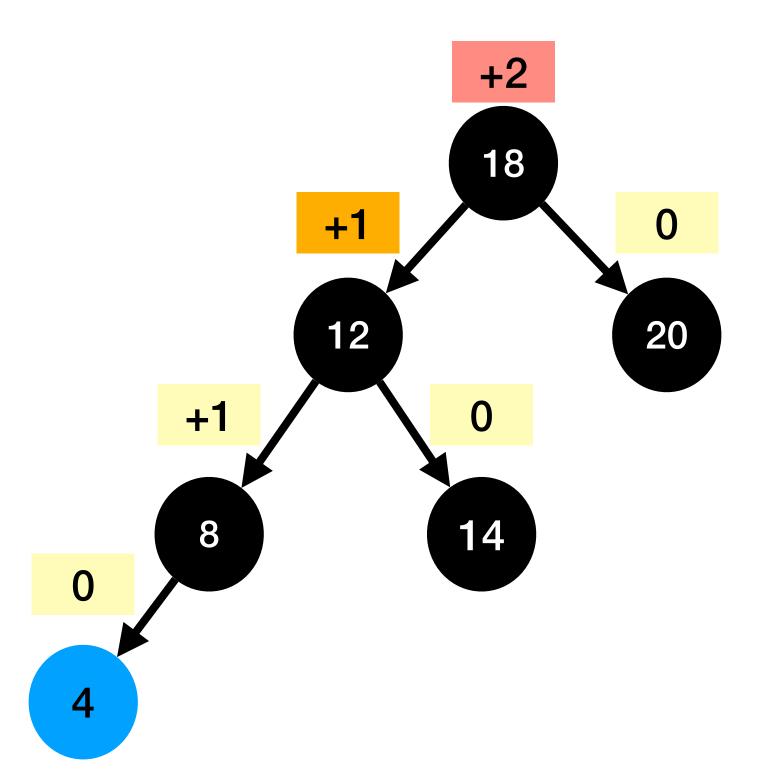
Insert or delete a node





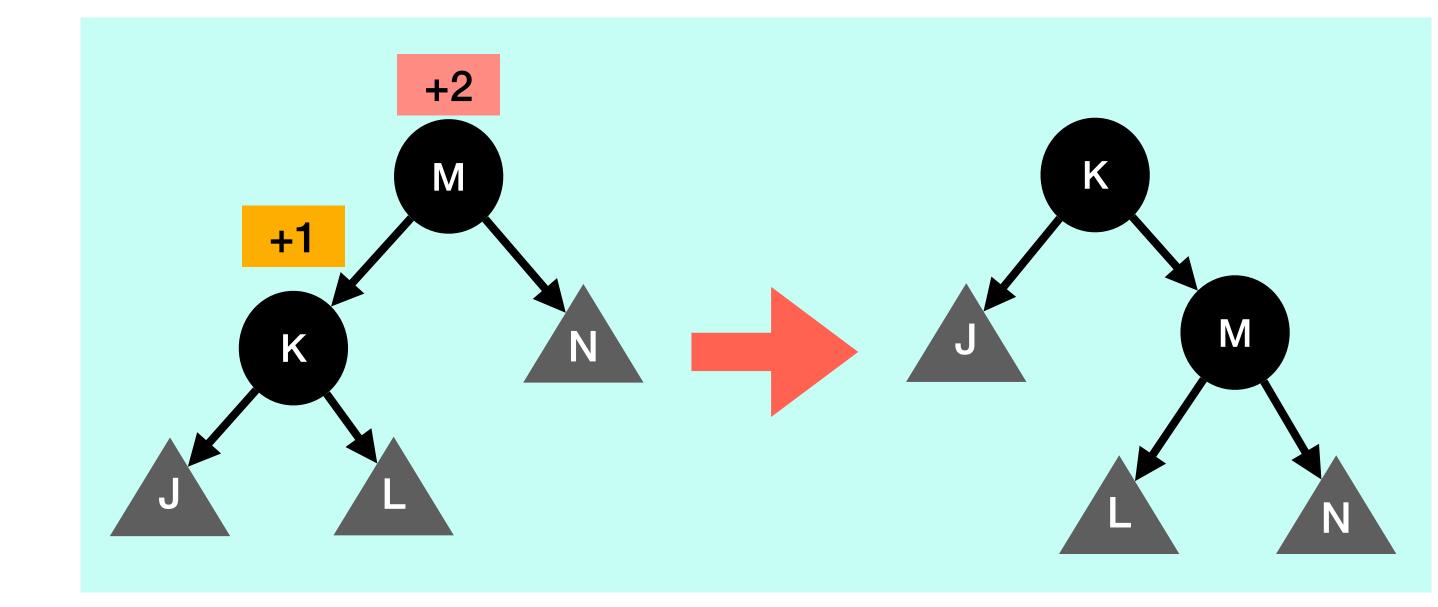
Rebalancing Rotation

1. **Left of Left**: The new node is inserted in the left sub-tree of the left sub-tree of the critical node



Rebalancing Rotation

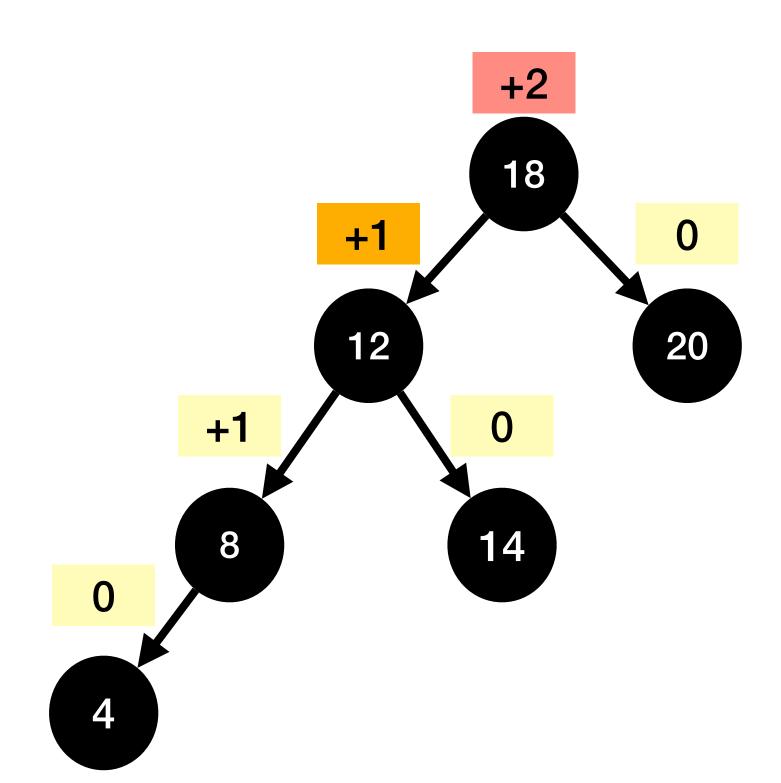
1. LL Rotation (Rotate Right)

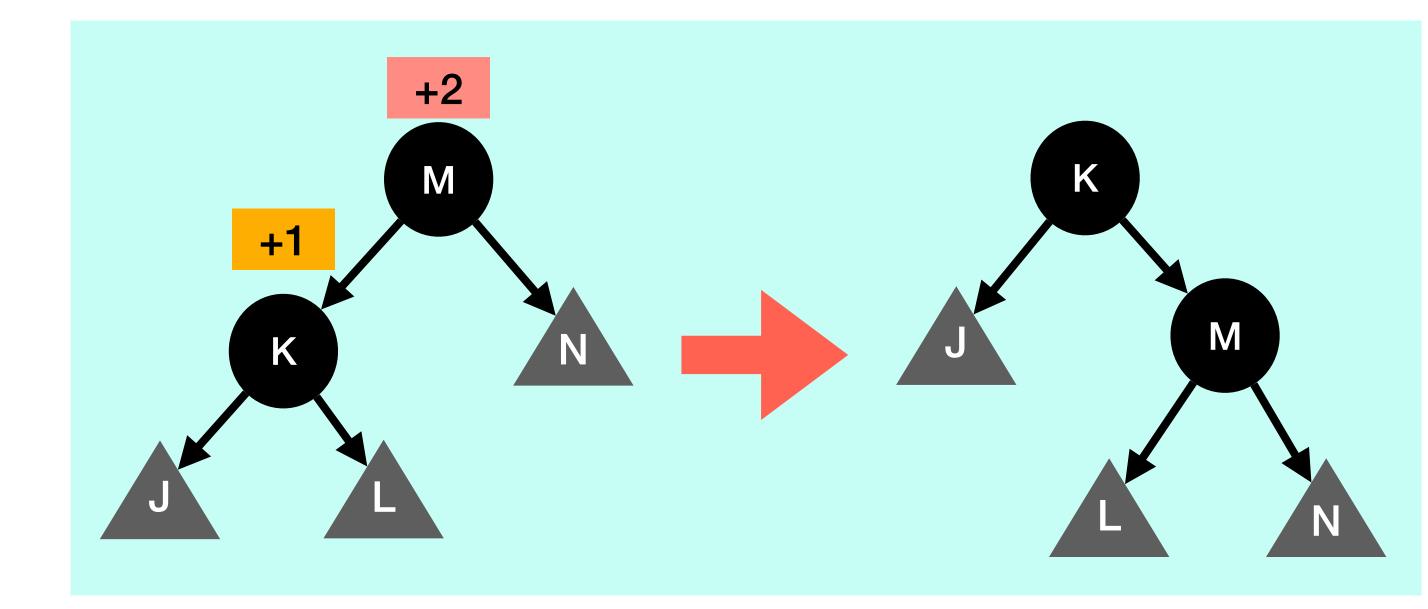


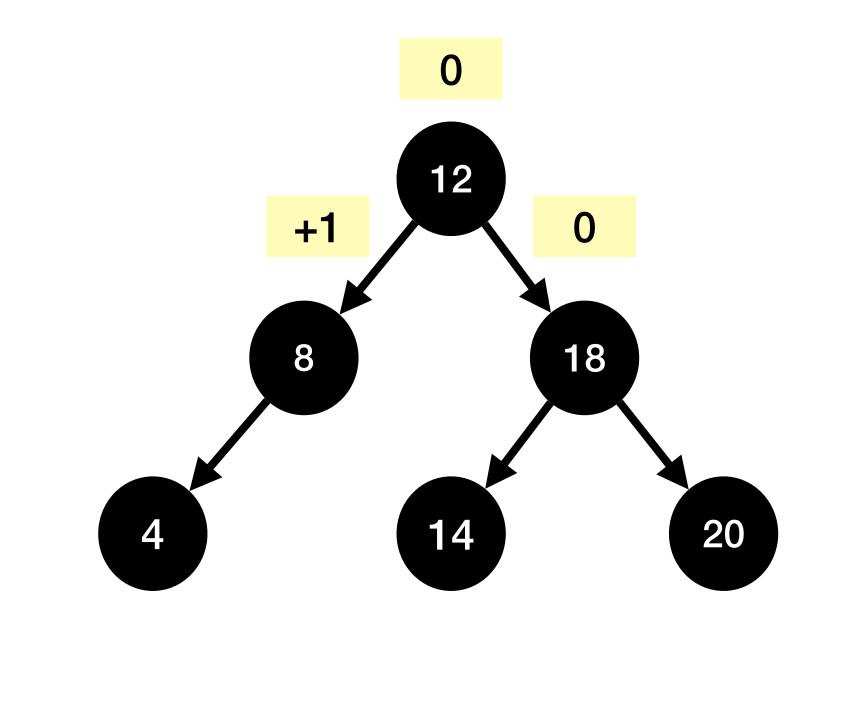
```
TREENODE_T *rotateRight (TREENODE_T *node)
{
   TREENODE_T *temp;
   temp = node->left;
   node->left = temp->right;
   temp->right = node;
   node = temp;
   return(node);
}
```

Rebalancing Rotation

1. LL Rotation (Rotate Right)

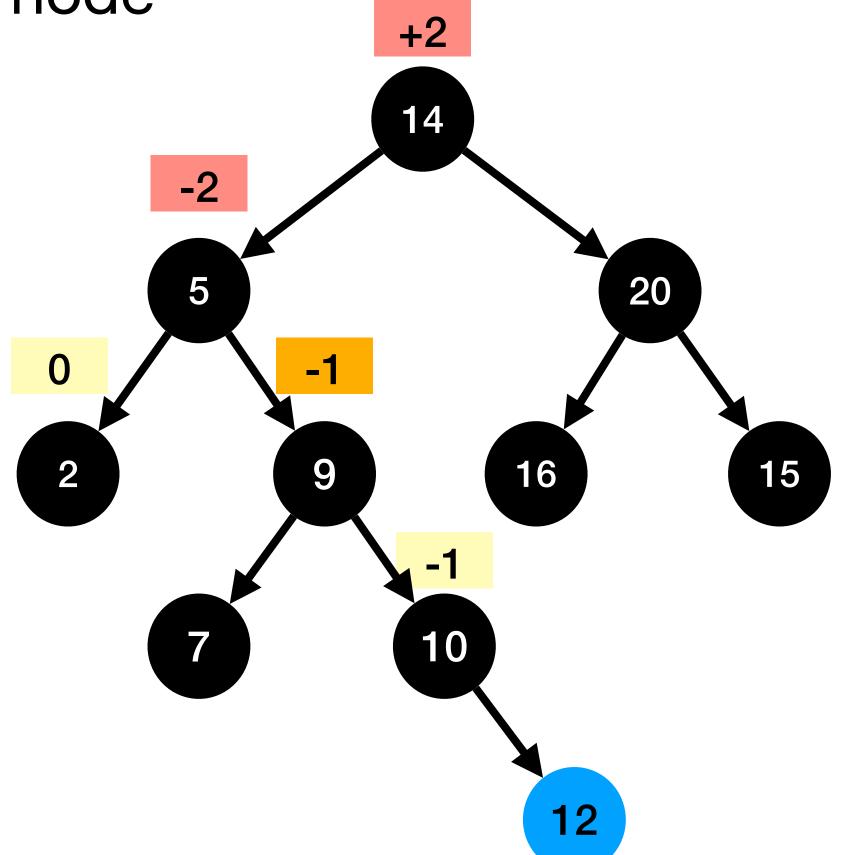






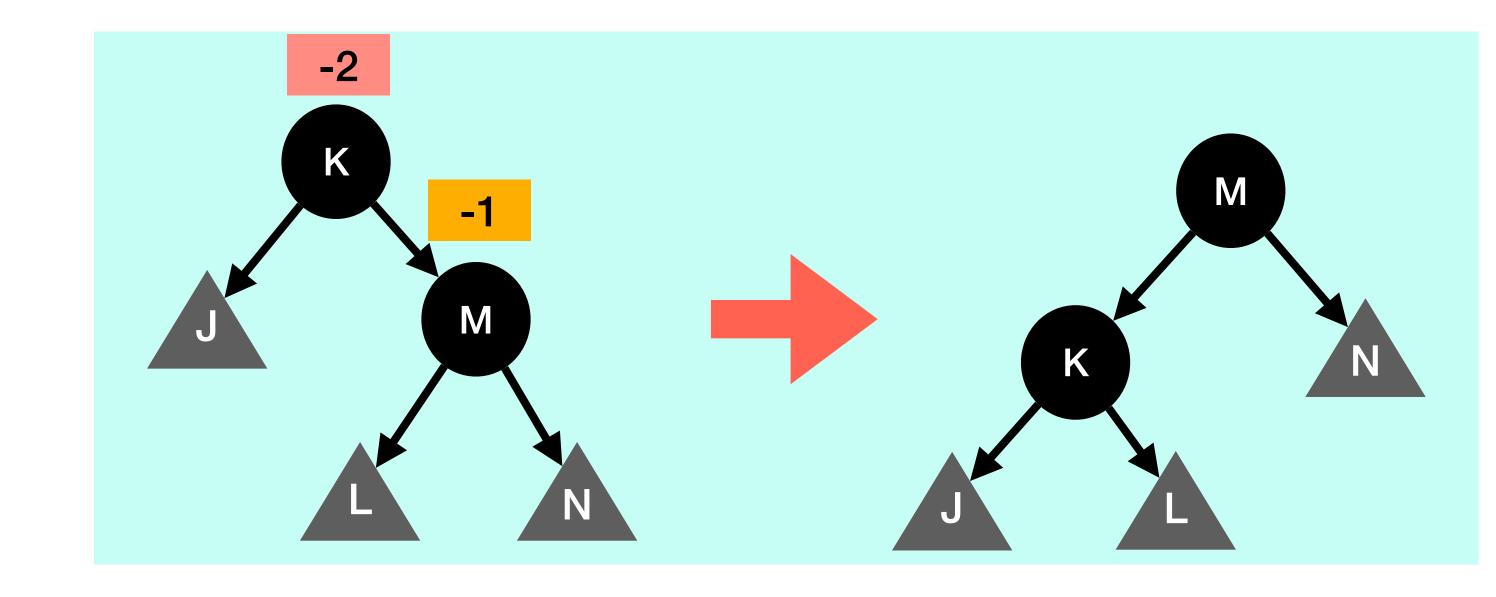
Rebalancing Rotation

2. **Right of Right:** The new node is inserted in the right sub-tree of the right sub-tree of the critical node



Rebalancing Rotation

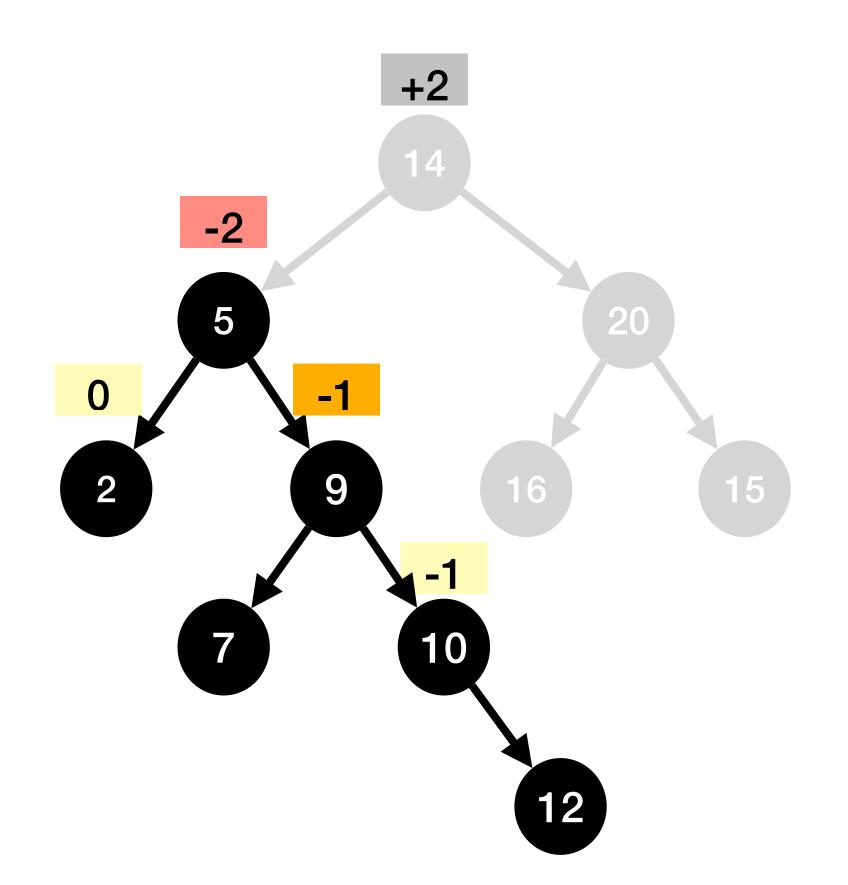
2. RR Rotation (Rotate Left)

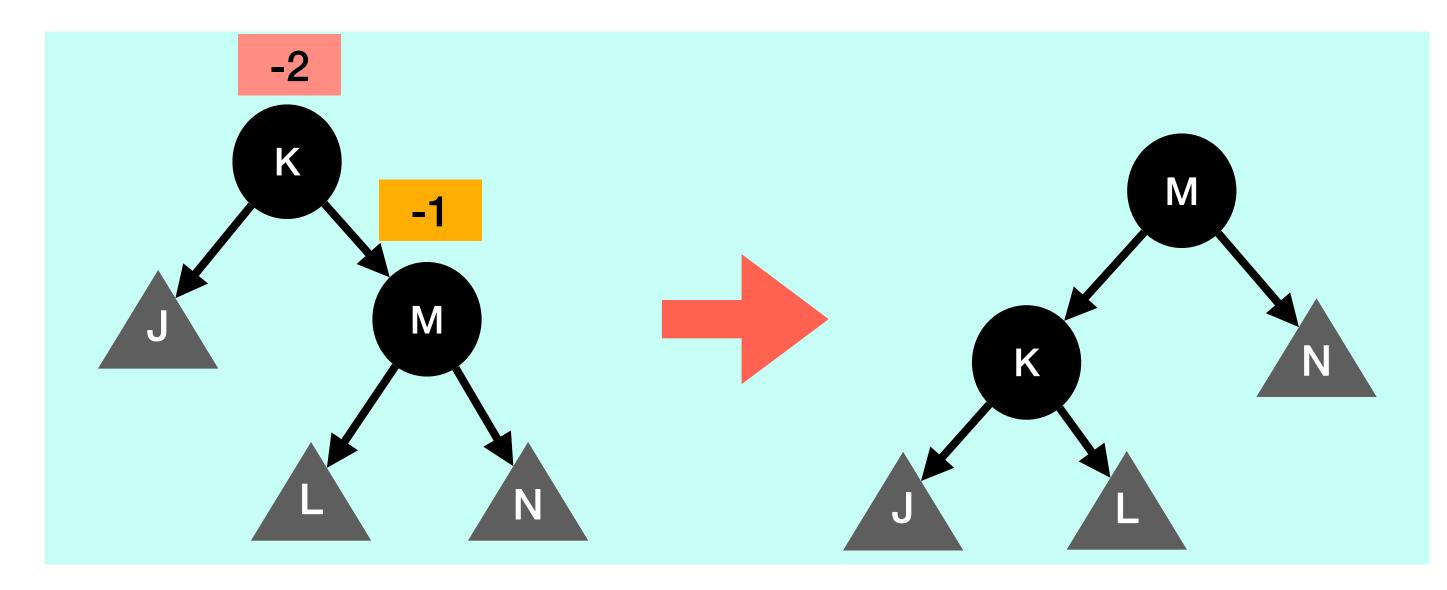


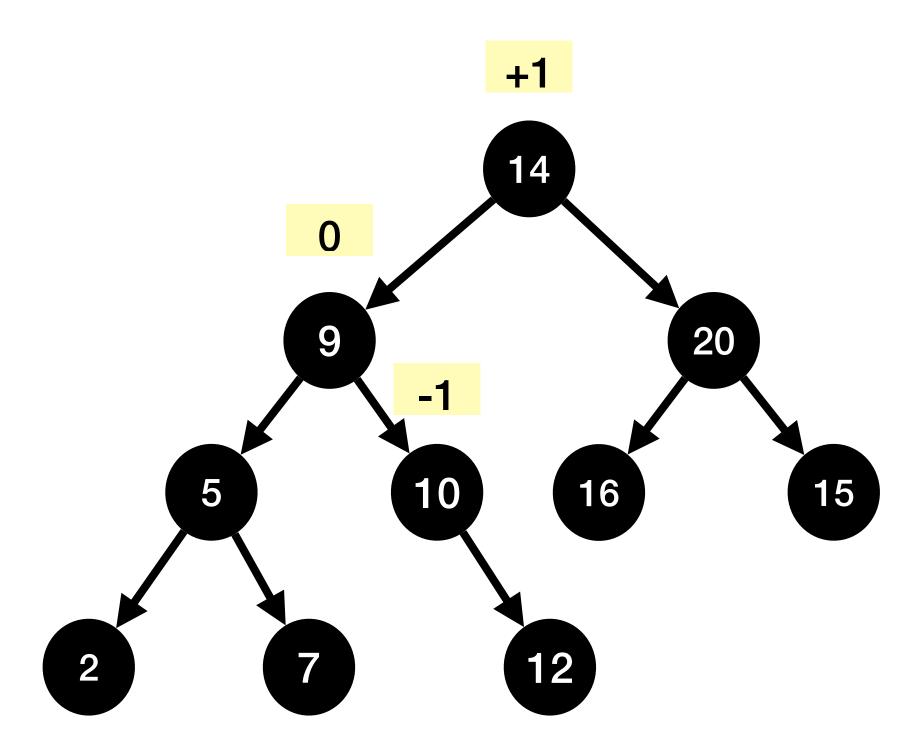
```
TREENODE_T *rotateLeft (TREENODE_T *node)
{
   TREENODE_T *temp;
   temp = node->right;
   node->right = temp->left;
   temp->left = node;
   node = temp;
   return(node);
}
```

AVL TreeRebalancing Rotation

2. RR Rotation (Rotate Left)

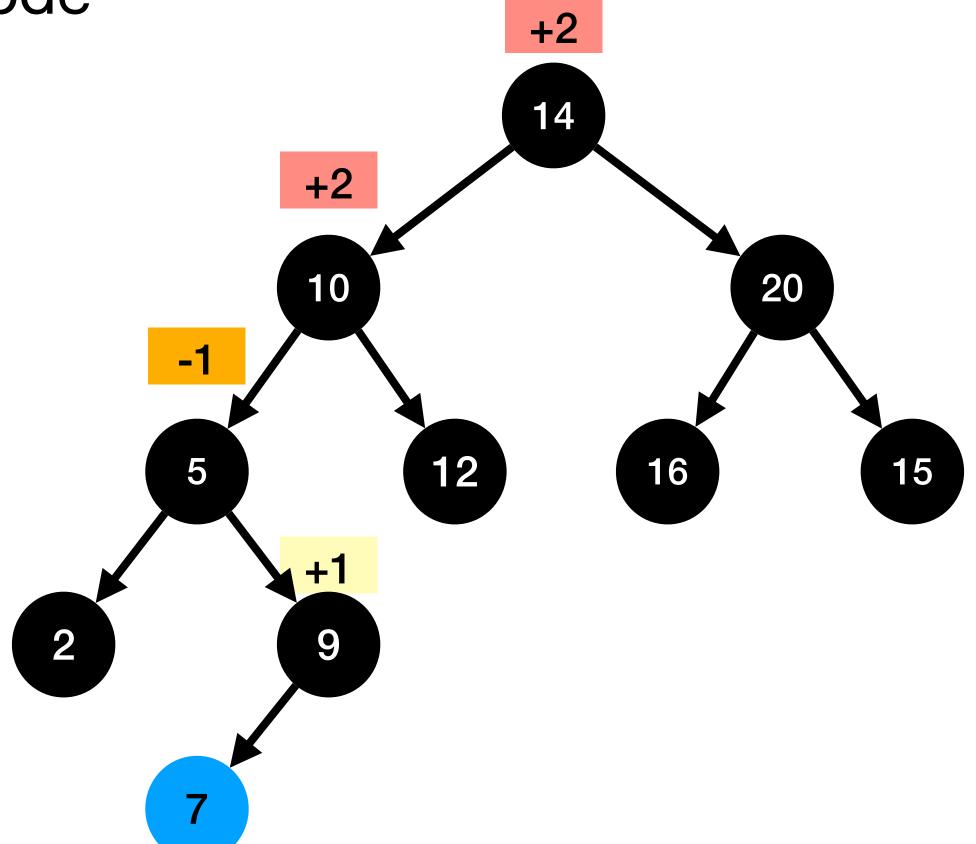






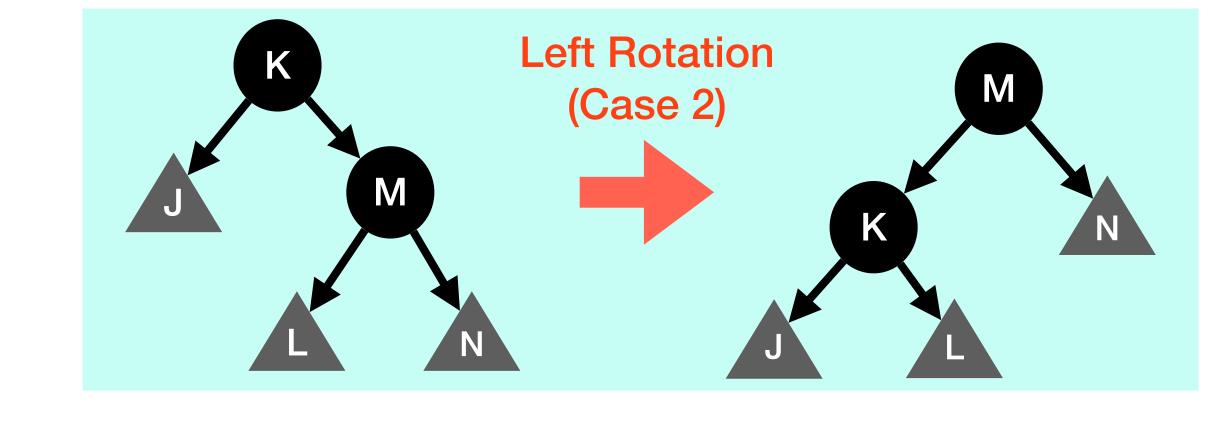
Rebalancing Rotation

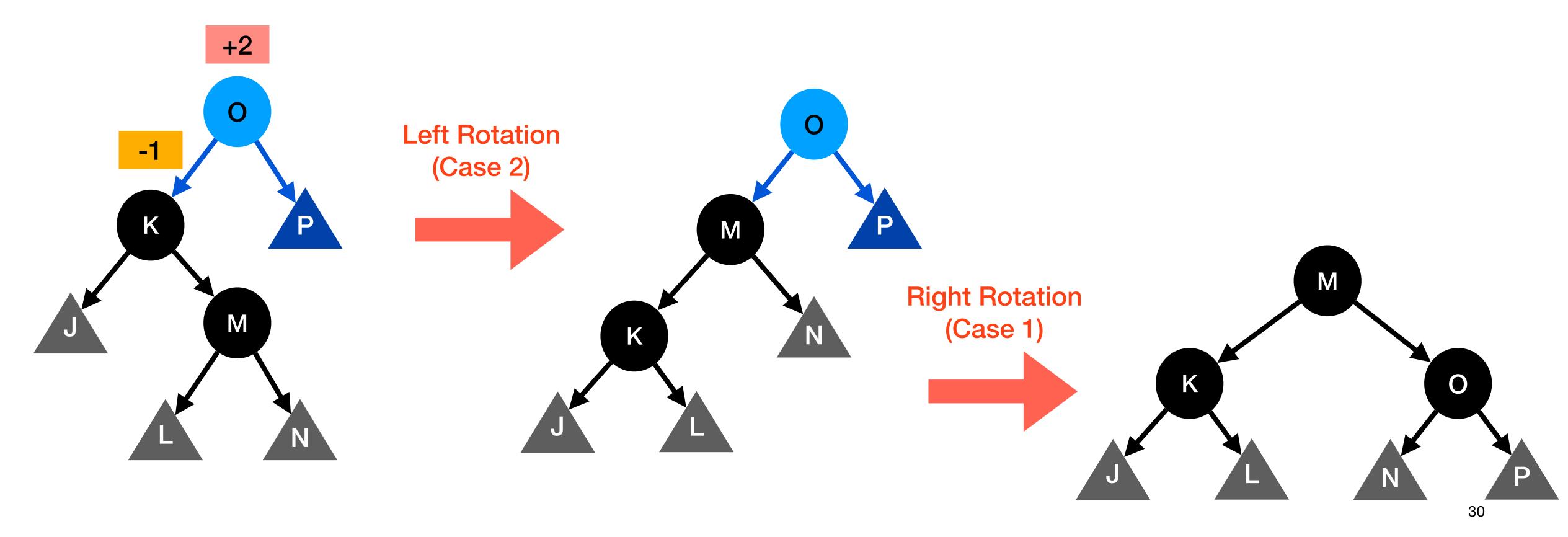
3. **Right of Left**: The new node is inserted in the right sub-tree of the left sub-tree of the critical node



Rebalancing Rotation

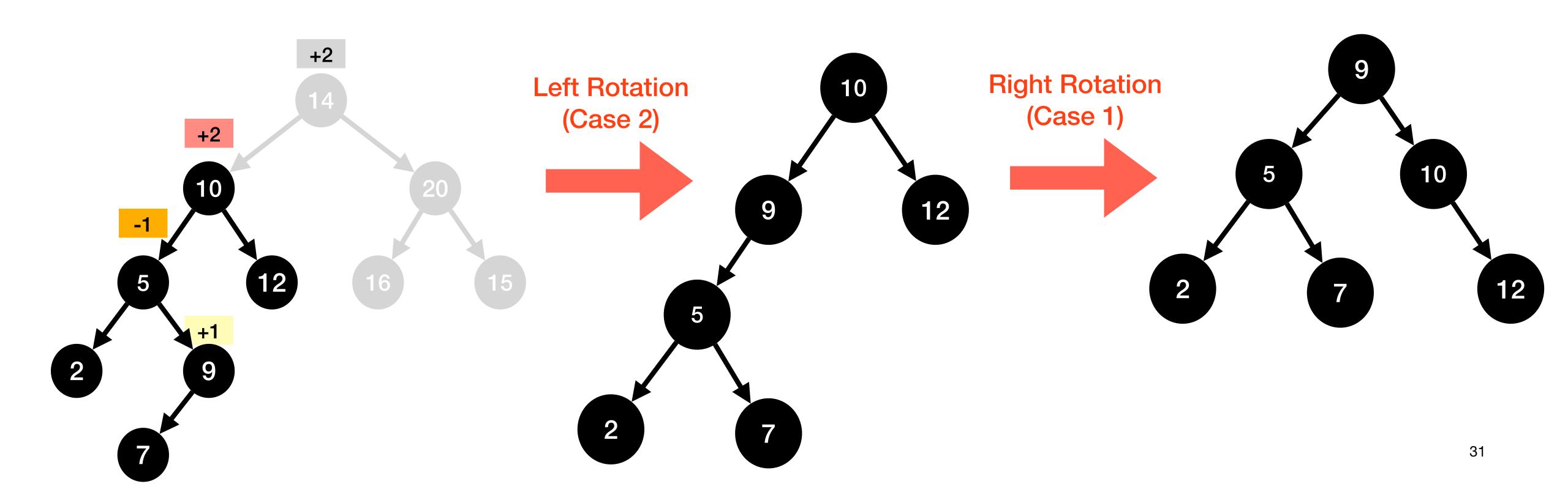
3. LR Rotation (Rotate Left to Right)





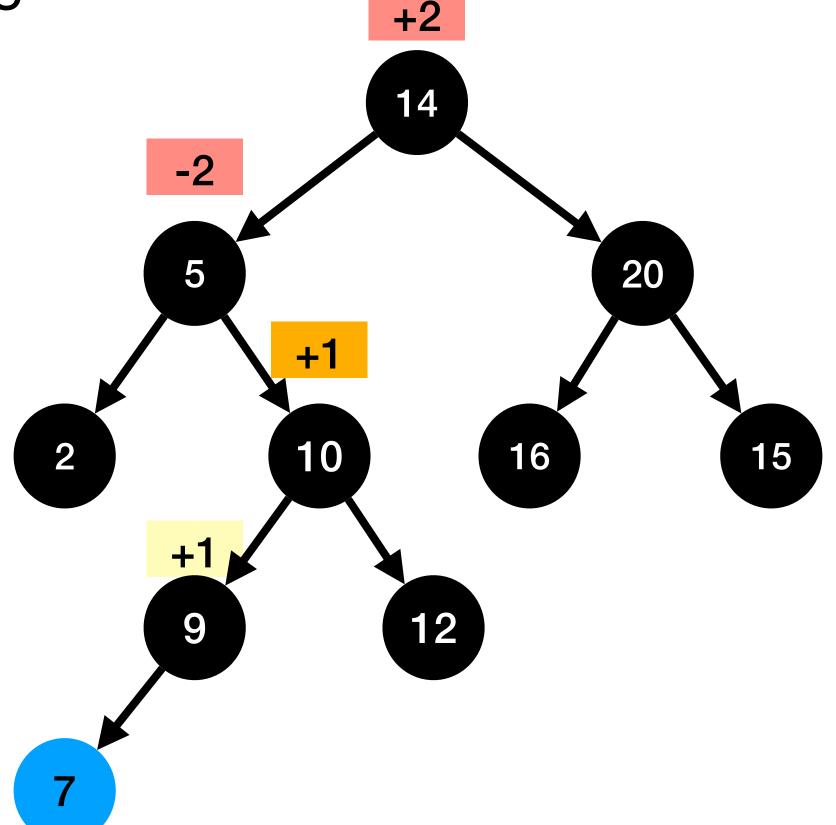
Rebalancing Rotation

3. LR Rotation (Rotate Left to Right)



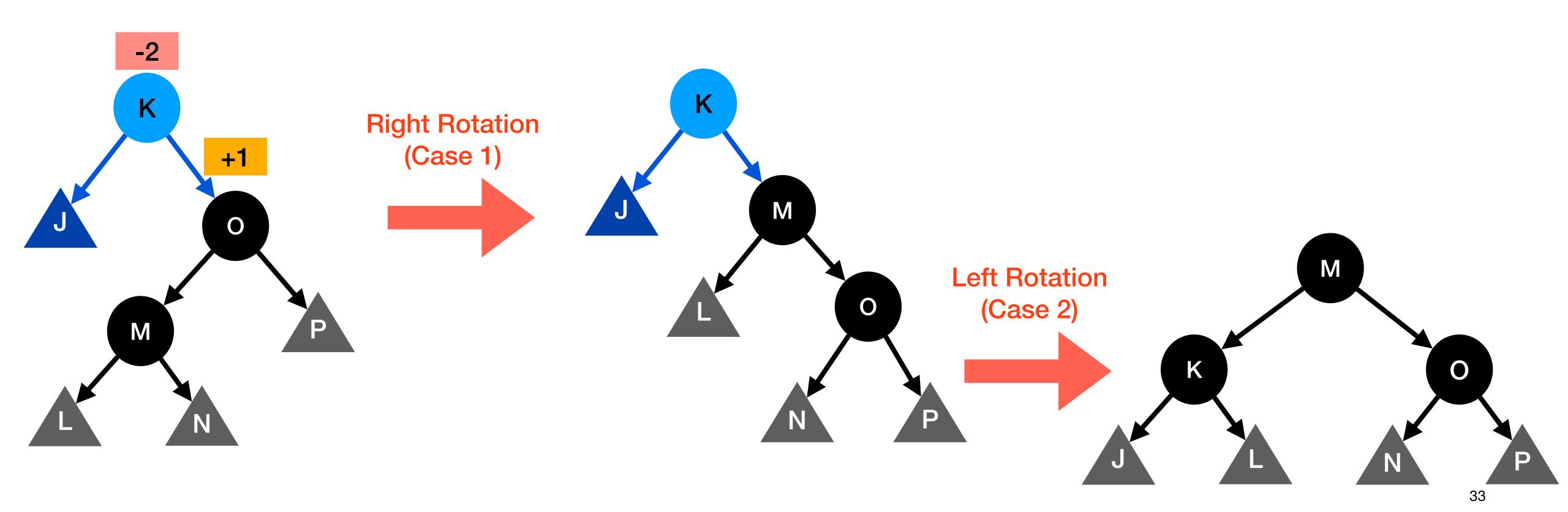
Rebalancing Rotation

4. **Left of Right**: The new node is inserted in the left sub-tree of the right sub-tree of the critical node



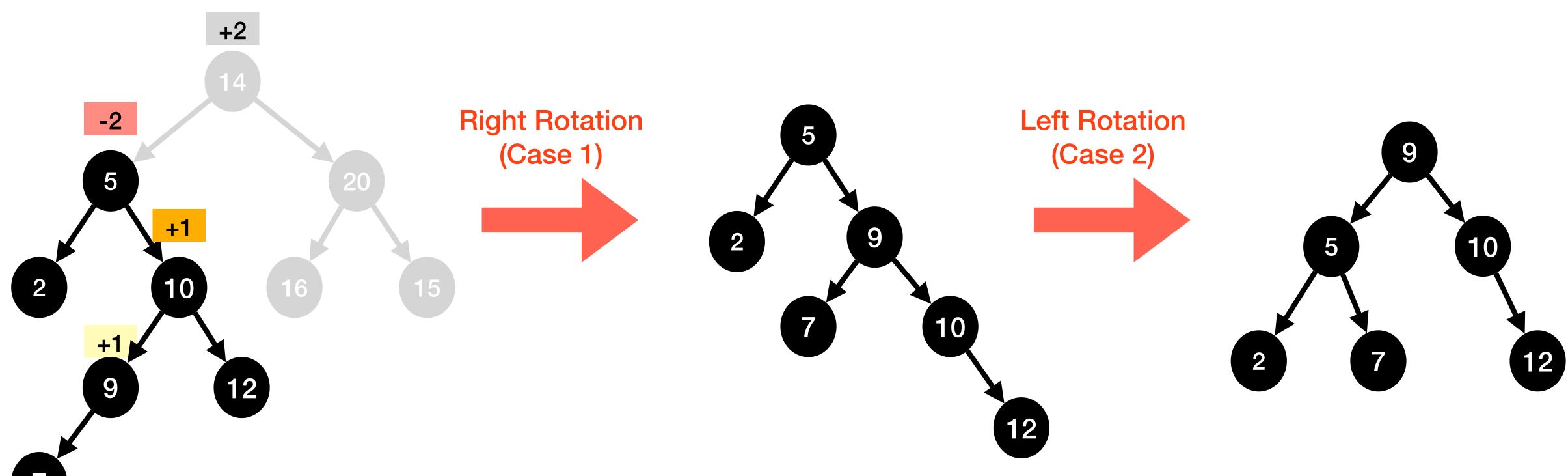
Rebalancing Rotation

4. RL Rotation (Rotate Right to Left)



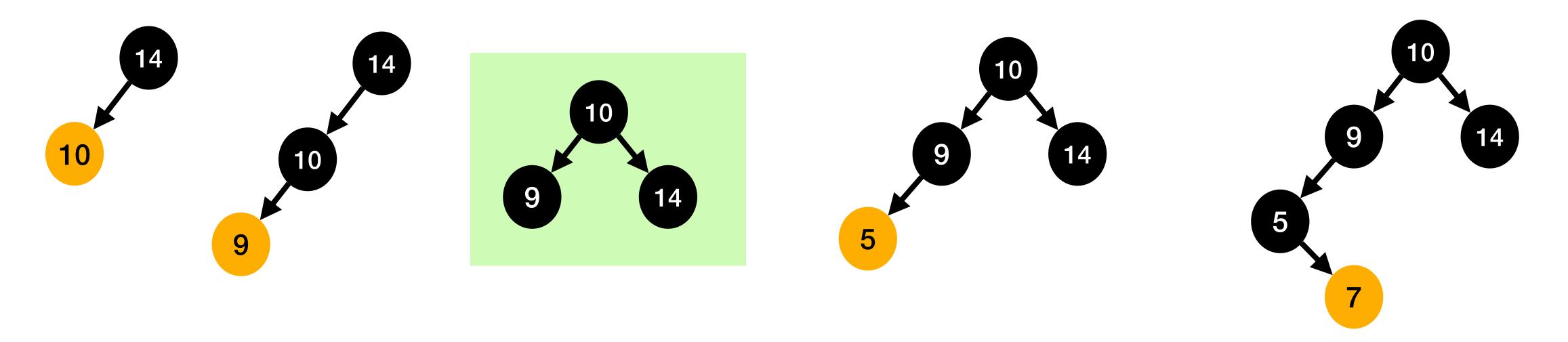
Rebalancing Rotation

4. RL Rotation (Rotate Right to Left)



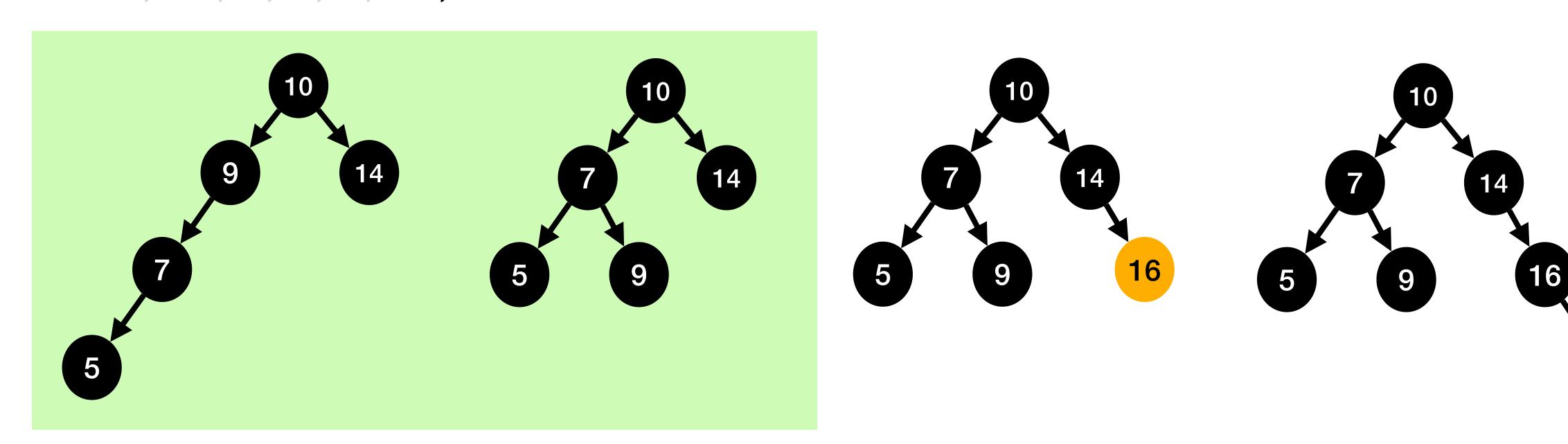
Create an AVL Tree

• 14, 10, 9, 5, 7, 16, 25



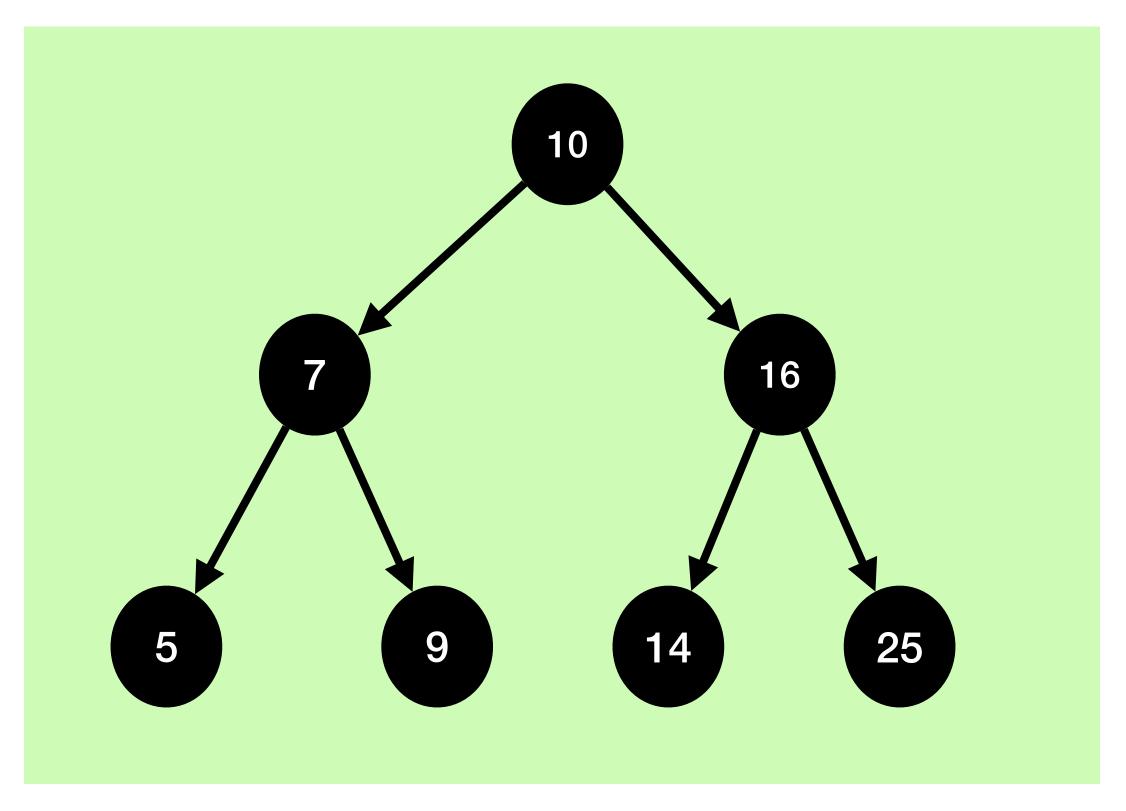
Create an AVL Tree

• 14, 10, 9, 5, 7, 16, 25



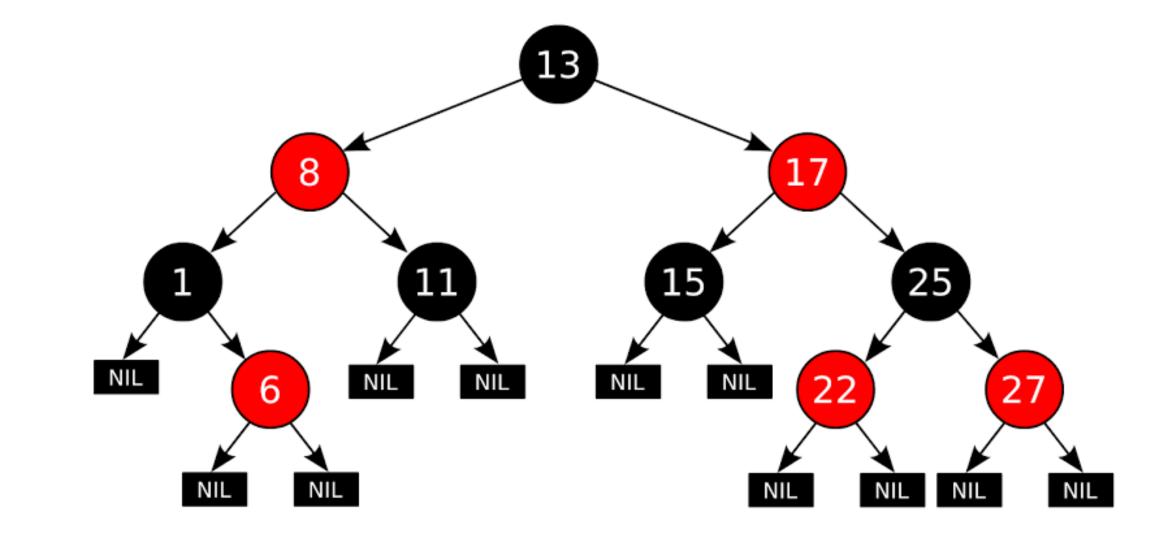
Create an AVL Tree

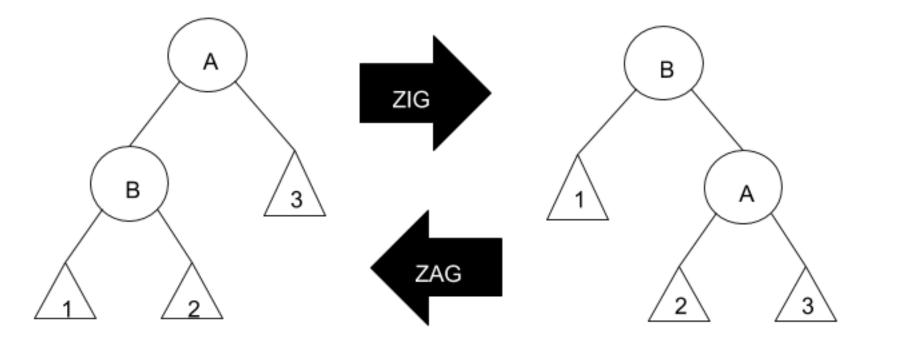
• 14, 10, 9, 5, 7, 16, 25

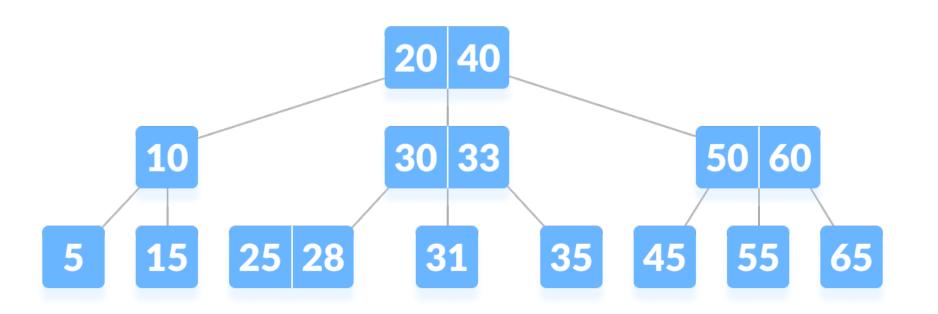


More trees

- Red-black tree
- Splay tree
- B-Tree
- Etc.







Wrap up

- Binary Tree from General Tree
- Binary Search Tree
 - Operations: Create, Search, Delete
- AVL Tree
 - Balance Factor
 - Rotation cases: LL, RR, LR, RL