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TUTORIAL

CSCI361 – Computer Security

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BLIND SIGNATURE

Blind signature protocol

A blind signature protocol requires the following components:

1. A digital signature mechanism for signer.
2. Functions f and g that are known only to the sender. The function f is known as blinding function, and function g is known as unblinding function.

The blinding function f is defined as: $f(m) = m \cdot k^e \bmod n$

The unblinding function g is defined as: $g(m) = k^{-1}m \bmod n$

Blind signature protocol

Chaum's blind signature protocol:

Sender A receives a signature of B on a blinded message.

From this, A computes B's signature on a message m chosen a priori by A. Note: $0 \leq m \leq n-1$. Note B has no knowledge of m nor the signature associated with m .

Blind signature protocol

Assuming we use RSA signature scheme, then B's RSA public and private keys are (n, e) and d respectively.

A need to choose k , a random secret integer satisfying $0 \leq k \leq n - 1$ and $\gcd(k, n) = 1$.

Blind signature protocol

Protocol:

Blinding: A computes $m' = mk^e \bmod n$ and sends m' to B for signing.

Signing: B compute $s' = (m')^d \bmod n$ which it sends to A.

Unblinding: A computes $s = k^{-1}s' \bmod n$, which is B's signature on m .

Uses s to obtain the original message m .

Blind signature protocol

For example, using the following RSA parameters:

$p = 11$, $q = 13$, $n = p \cdot q = 143$, $e = 37$, $d = 13$, and $m = 74$.

Sender: Bob

Bob chooses $k = 10$. Through extended Euclidean, we obtained $k^{-1} = 43$ (Note, we need k^{-1} to unblind later.)

Bob computes $m' = mk^e \bmod n$ (blinding the message m)

$$= 74 \cdot 10^{37} \bmod 143$$

$$= 74 \cdot 10 \bmod 143 = 25$$

Blind signature protocol

Bob sends $m' = 25$ to Alice, the signatory.

Alice receives the message m' and signs the message:

$$\begin{aligned}s' &= (m')^d \bmod n \\ &= (25)^{13} \bmod 143 \\ &= 38\end{aligned}$$

Alice sends $s'=38$ to Bob.

Blind signature protocol

Bob receives s' and compute:

$$\begin{aligned}s &= k^{-1}s' \bmod n \\ &= (43)(38) \bmod 143 \\ &= 61\end{aligned}$$

Bob can obtain the original message as follow:

$$\begin{aligned}m &= s^e \bmod n \\ &= 61^{37} \bmod 143 \\ &= 74\end{aligned}$$