

# CSCI361

*Fair Exchange and Zero Knowledge Proofs*

# Fair Exchange

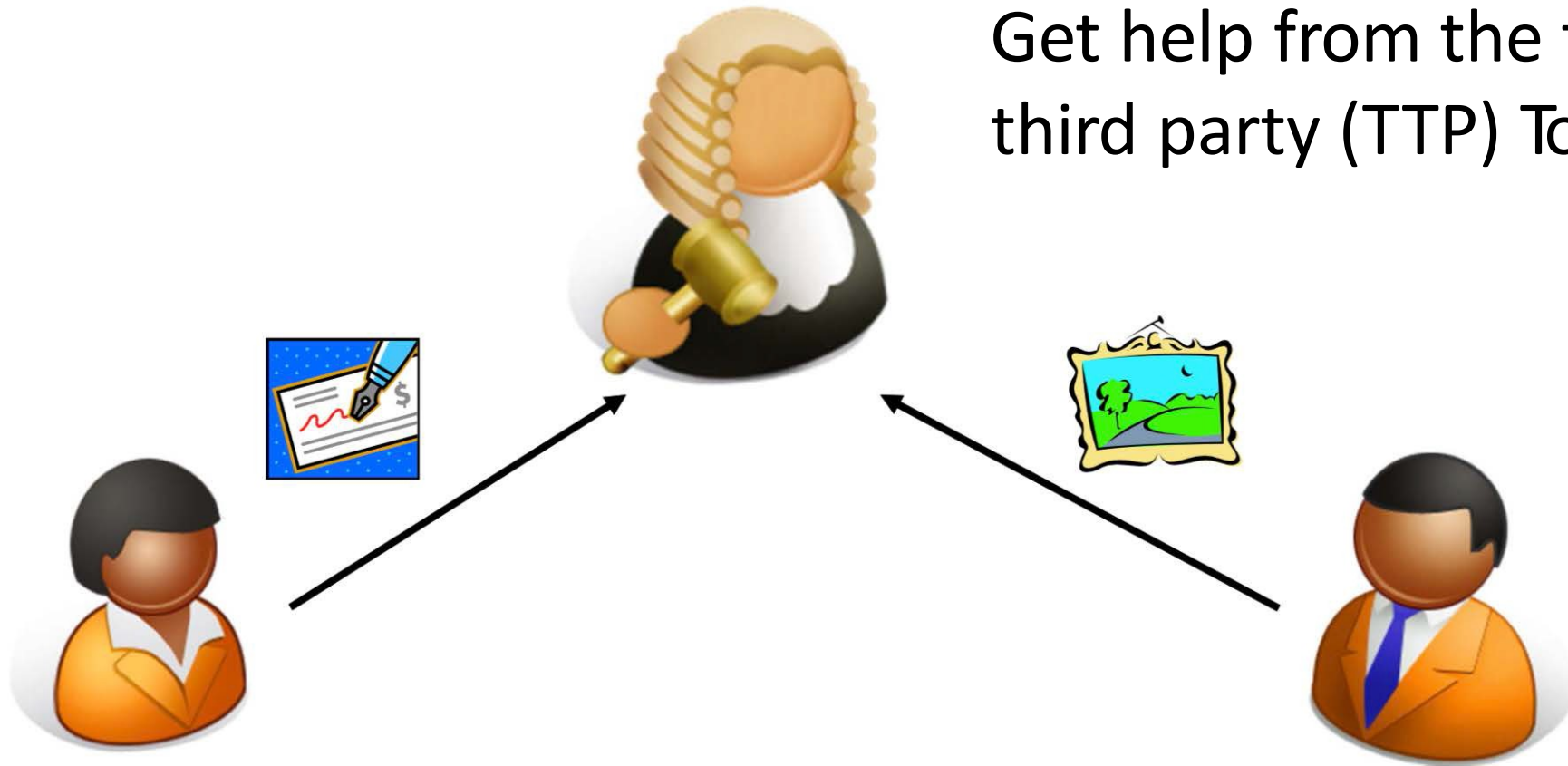
- Two parties, Alice and Bob, would like to exchange something

Let's say Alice wants to get a digital photo from Bob and Bob wants to get Alice's eCheque (which Alice signed) in exchange. We want to make sure that this transaction is **fair**.

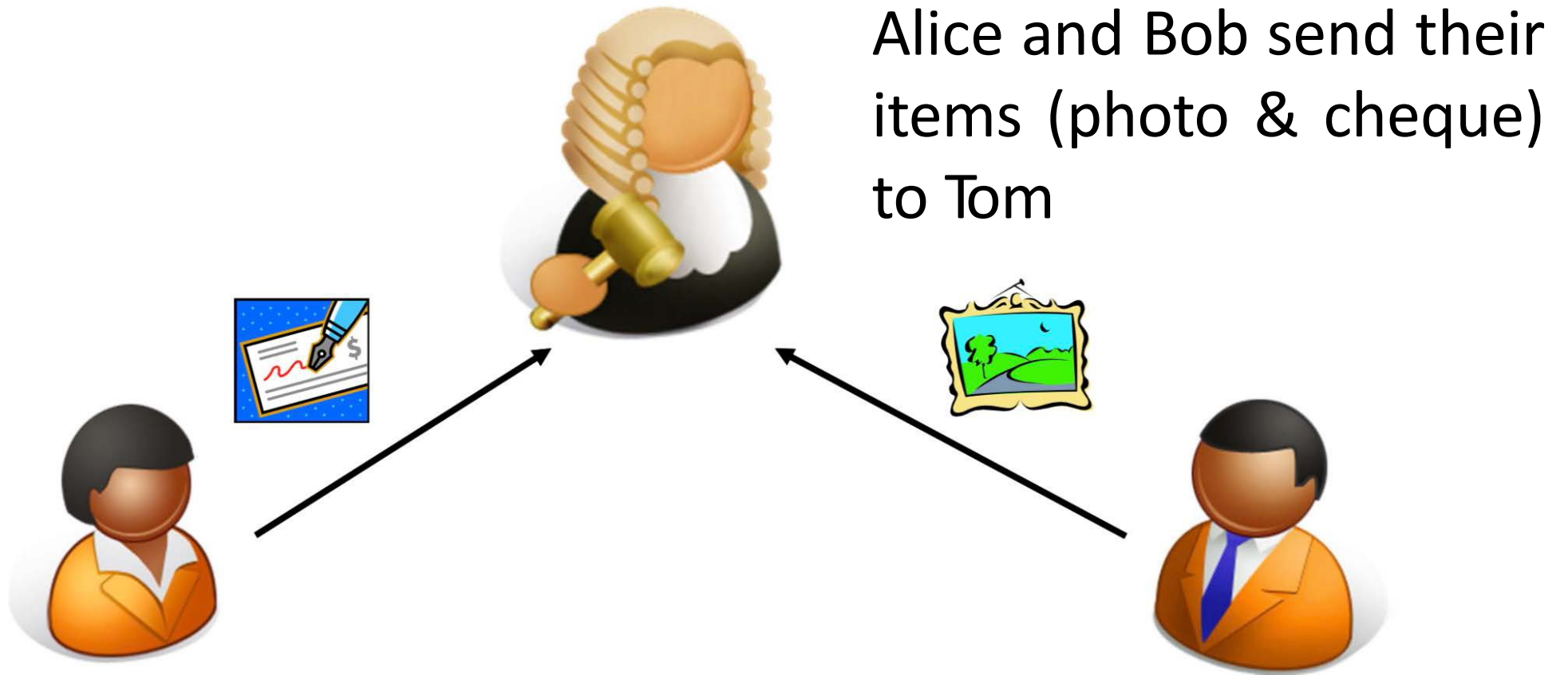
- What is **fairness** then?
  - (In the context of electronic commerce), Participants shouldn't have advantages over each other.
  - For example: It wouldn't be fair if one party can avoid their obligations in a contract if the other party has completed their obligations in a contract.

# Possible Solution

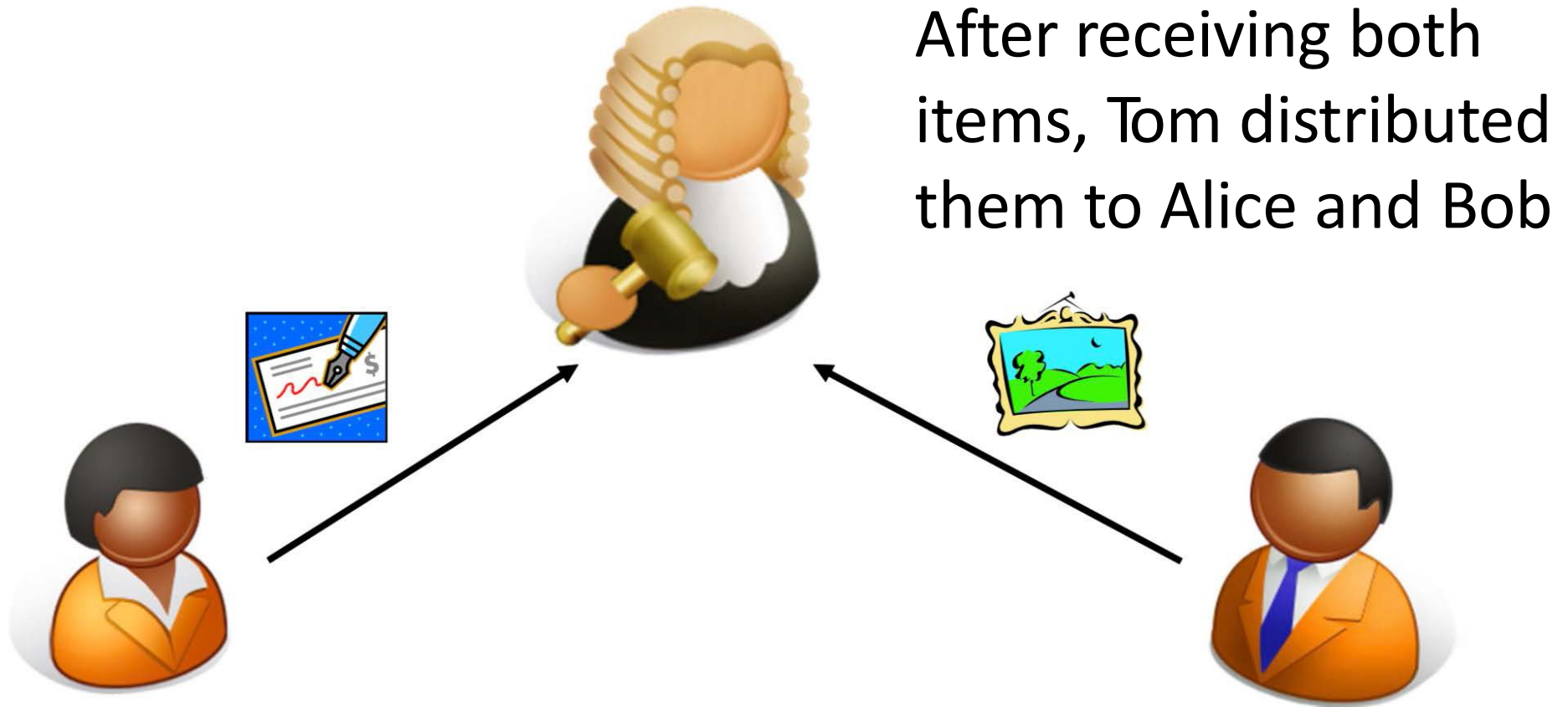
Get help from the trusted third party (TTP) Tom!



# Possible Solution



# Possible Solution



# Good...

- But Tom must be accessible all the time
- Can Alice and Bob conduct the exchange without the *active* participation from Tom? (Tom has to receive and send something from Alice and Bob.)

# Optimistic Fair Exchange

Scenario: Alice wants to get a digital item from Bob and Bob wants to get Alice's signature in exchange.

- (Conceptual) Solution

- Alice creates something called a **partial signature** (P)
- Alice sends P to Bob.
- Bob sends Alice his item.
- Alice sends her full signature (S) to Bob
  - ❖ If Alice runs away after getting Bob's item, Bob can go to Tom and asks Tom to convert P into S.

# Optimistic Fair Exchange - Realisation

1. Alice generates her signature,  $S_A$ .
2. Instead of sending  $S_A$  directly to Bob, Alice encrypts  $S_A$  under Tom's public key. The resulting ciphertext  $P = E_{pk}(S_A)$  is the partial signature.
3. Alice sends  $P$  to Bob.
4. Bob sends his item to Alice.
5. Alice sends  $S_A$  to Bob.
  - If Alice does not send  $S_A$  to Bob, Bob can contact Tom with  $P$  and ask him to decrypt  $P$  to  $S_A$ . (Bob can finally obtain  $S_A$ .)



# Optimistic Fair Exchange - Realisation

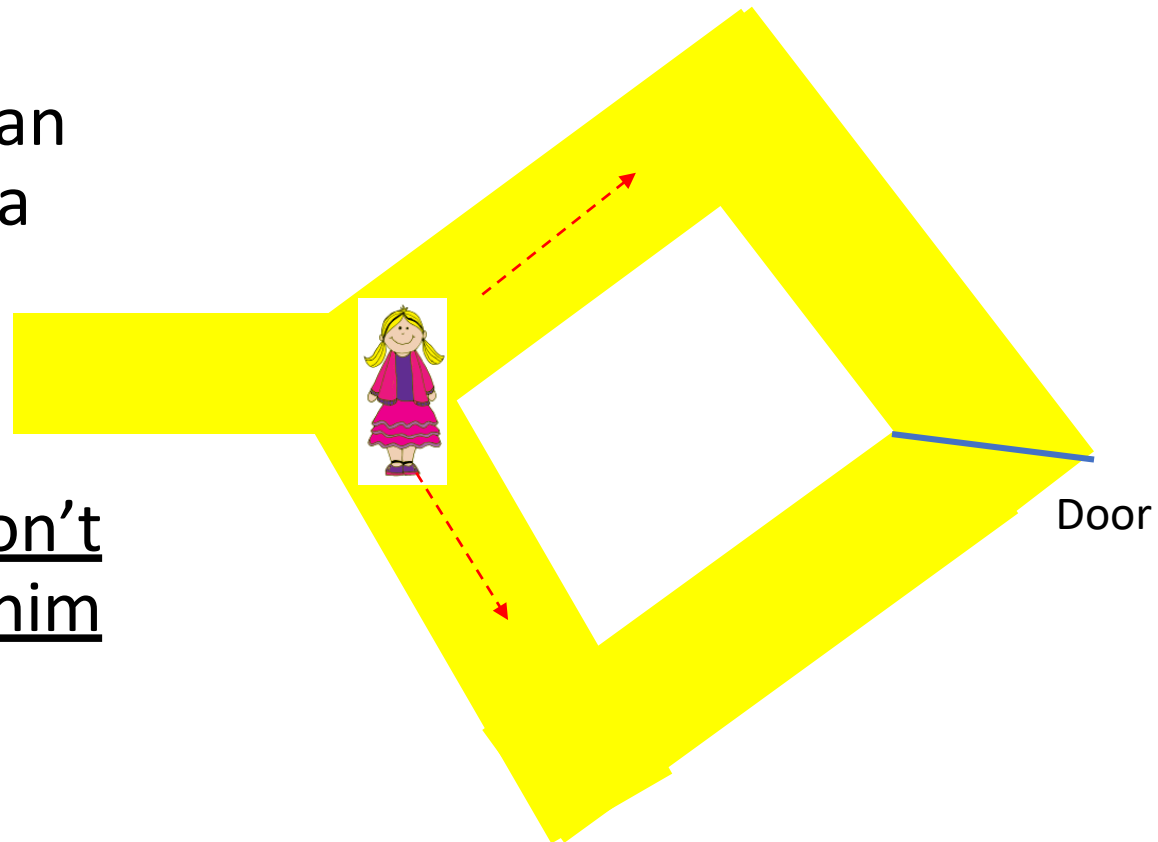
- Additional step
  - ✓ How can Bob make sure C contains a valid signature of Alice?
  - ✓ Using zero-knowledge proof → Prove that P is the encryption of a valid  $S_A$  **without revealing  $S_A$**

# Beyond Optimistic Fair Exchange

- Optimistic Fair Exchange can solve the problem of having to require active TTP
- There is a Fair Exchange scheme that does not require TTP at all!

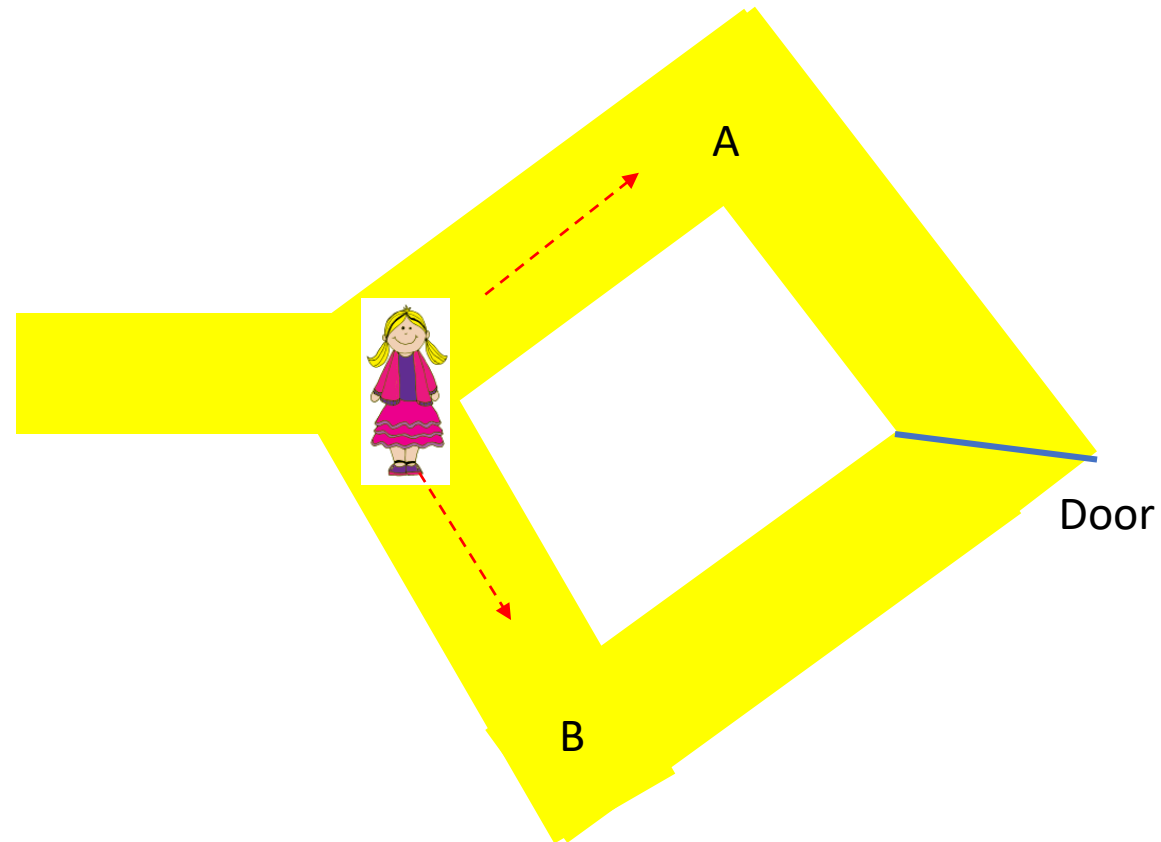
# Zero Knowledge Proof: Ali Baba Cave Problem

- Setting: There is a ring-shape cave where the path is blocked by a door. Peggy can open the door if she uses a right magic word as a key. Victor wants to know whether Peggy knows the magic word. But Peggy won't reveal the magic word to him (or anyone else).



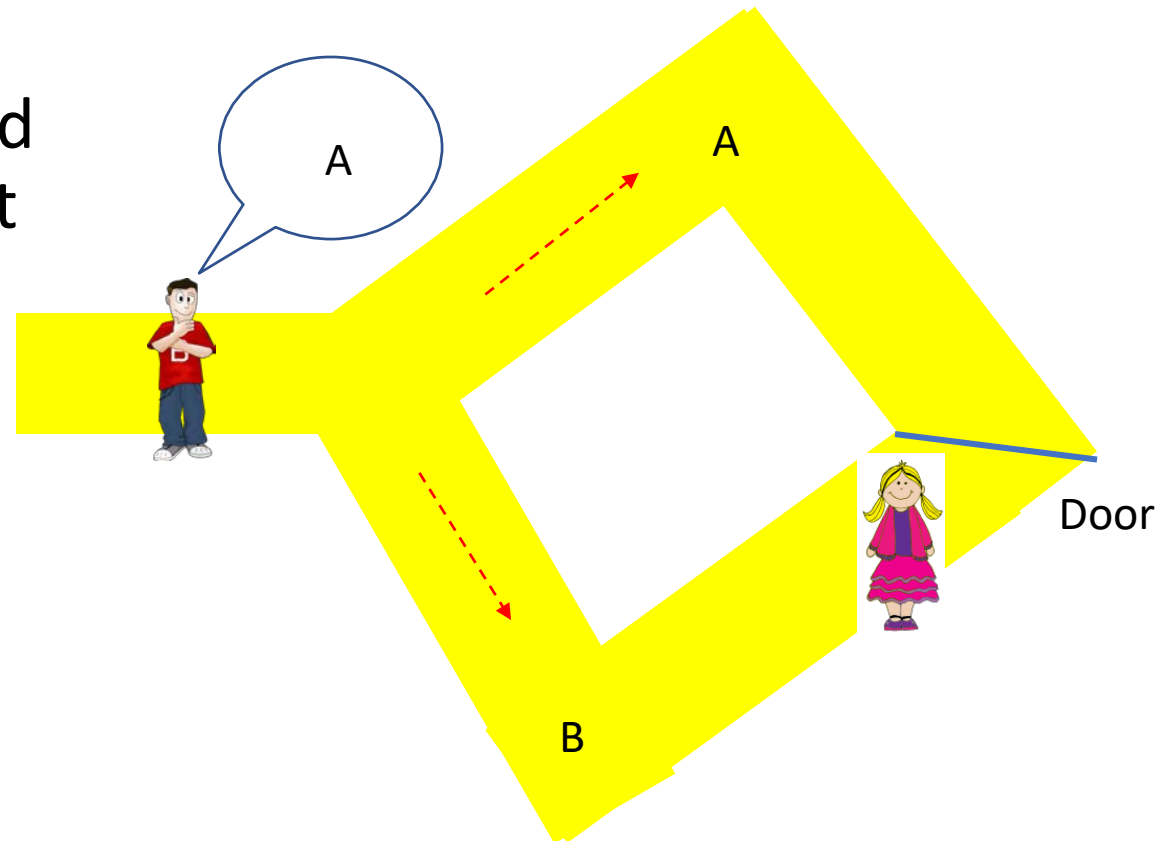
# Zero Knowledge Proof: Ali Baba Cave Problem

- Step 1: Peggy chooses either path A or B and moves towards the door while Victor waits outside



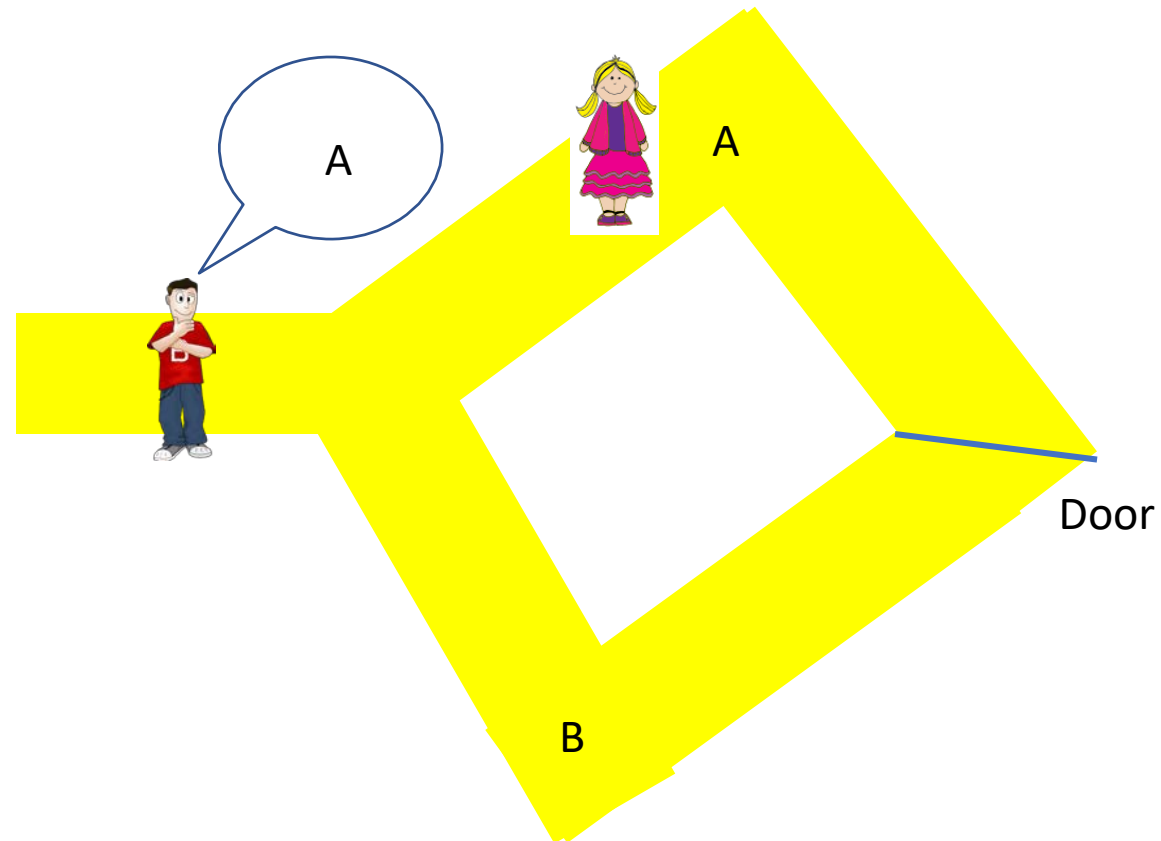
# Zero Knowledge Proof: Ali Baba Cave Problem

- Step 2: Now Victor enters the cave, picks at random the name of path Peggy should use to return and shouts it out!



# Zero Knowledge Proof: Ali Baba Cave Problem

- Step 3: Peggy returns to the entrance using the path Victor named



# Zero Knowledge Proof: Ali Baba Cave Problem

- If Peggy knows the magic word, she can reliably return to the entrance using the path Victor named
- If Peggy does not know the magic word, she can return to the entrance only if Victor named the path she chose in Step 2 with probability  $\frac{1}{2}$
- Victor repeat Step 1 to 3 many times to see whether Peggy really knows the magic word
  - In  $n$  trials, the probability that Peggy returns through the path without knowing the magic word is  $(\frac{1}{2})^n$