Reference	Motivation	Contribution(s)	Limitation(s)	Methodology
				Summary
[3]	DAG tasks scheduling is getting more popular and fluid scheduling performs great theoretically	Provide a DAG- fluid scheduling algorithm that per- forms way better in terms of acceptance ratio then previous algorithms	Fluid scheduling is unpractical and introduces a lot of overhead and task migrations, also only for implicit deadlines tasks	fluid-based algorithm where it decomposes a DAG's subtasks into multiple sequential segments
[5]	DAGs are popular but no one looked at the intra-task execution order to leverage the graph structure	proposes a priority list scheduling al- gorithm for a single DAG task which performs better than SOTA in terms of makespan	no comparison with opti- mal priority assignment algorithms / optimal sched- ules.	uses the length (in terms of weets) of each paths passing through the current vertex to assign the priority to the current vertex, the higher the length, the higher priority
[12]	Federated scheduling for DAG tasks is has proved efficient but for tasks where the difference between the critical path and the deadline is small, it can lead to overallocating cores.	proposed a fedrated and bundled-based scheduling algorithm to avoid this problem and enhanced the schedulability of DAG tasks using their algorithms	They only compare their method with an example of a dag task set comprised of 3 dag tasks.	Uses federated scheduling for tasks with high critical path to deadline ratio and bundled scheduling for tasks with low critical path to deadline ratio.

so one can only approximate the optimal algorithm (when considering polynomial timed algorithms) and not a lot has been done on scheduling parrallel reccuring tasks The use of multicore systems can induce contentions because of shared memory / cache. This can lead to non-determinim and unpredictable diversed approximate the that shortens the makespan of recurring DAG tasks compared to EDF amakespans amakespans duplication makespans and also only compares with it on an act to the makespan, it is less scalable than comparable algorithms. The use of multicore contentions ing algorithm based that avoids and unpredictable due to LET imple- The use of multicore contentions while approximate the than sample of a DAG task for makespans ample of a DAG task for makespans and duplication rithm and extended to predict the makespan, it is less scalable to the measurement and unpredictable due to LET imple- The use of multicore contentions while ture to evaluate their scheduling algorithm but assigning to cores. Cores.	ombines from chedul- task algo- raluates hal sim- ct (the 78) e mea-	
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real-time systems, the makesp		
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ecution time being	11.	
added because of	11.	
the implementation	11.	
overhead.		

	The capacity-bound is a bound used as a performance measurement but also as a schedulability test for DAG scheduling. However, it uses the same bound to bound both the normalized utilization and the tensity (critical path over the deadline) of a DAG task which can exclude DAG tasks that actually are schedulable but not according to this capacity-bound.	Introduces a new bound called the util-tensity bound which proves to be a better schedulability test for GEDF with federated scheduling.	only looks at GEDF with federated scheduling and not other scheduling algorithms.	The scheduling algorithm applied uses GEDF for very low tensity tasks, tasks with high utilization and relatively high tensities are scheduled using federated scheduling, and low utilization tasks with relatively high tensities are scheduled by partitioned EDF
[8]	Decomposition-based scheduling can improve schedulability for DAG task scheduling but can also worsen it. It is, along with global scheduling, one of the main method to schedule DAG tasks.	Develop a decomposition strategy as well as a metric / schedulability test and the decomposition strategy proves to be the most efficient one according to the defined metric. The scheduling algorithm derived from this decomposition strategy (using GEDF variants) shows promising results in terms of acceptance ratios	Only looks at GEDF variants which is based on the EDF heuristics for priority assignments.	The decomposition works by first defining execution segments and then assigning subtasks to those segments based on their laxity so that there are no segments overloaded with workload.

[6]	The notion of degree of parallelism has been used for DAG task scheduling but lacks a clear definition in the research community.	Proposes a new response-time bound for DAG tasks as well as a new scheduling algorithm based on federated scheduling that outperforms the SOTA by more than 18% on average	They don't say which intratask scheduling algorithm is used (just that it's workconserving) and they don't consider intra-task scheduling.	Using the defined notion of degree of parallelism, they modify the federated scheduling approach by having a better way of choosing the number of cores for heavy tasks, which is based on the degree of parallelism of the heavy DAG task.
[15]	Although several DAG intra-task scheduling algorithm have been proposed in the literature, most of them ignore the communication delays, specifically the inter-core communications between the subtasks of a DAG tasks. In the Robotic / automotive industry, most communications can be done using the L1 cache of a processor, hence having a way to group subtasks into execution groups to execute on a single processor to remove inter-core communication delays.	Extend the DAG task model to EG-DAG (execution group dag) which binds groups of subtasks to a single physical core, thus reducing inter-core communication delays. They also introduce a scheduling algorithm and a wert for the algorithm while comparing the makespan to existing methods such as federated scheduling and critical-path based sub-tasks scheduling. Their method shows comparable performance while minimizing the communication overheads.	The evaluation has been done using only 100 DAG tasks which is quite low to cover all different types of DAG tasks. They propose a way to schedule multiple DAGs but do not offer evaluation results for that.	They use list scheduling with one list per execution group and use worst-fit heuristic to map the execution groups to the processors.

	Federated scheduling has shown great potential for constrained deadline tasks but for arbitrary deadline DAG tasks, especially those where the WCET is longer than the period, processor assignment is tricky and existing work have shown limitations by letting jobs migrate between the assigned processors, which produces a more pessimistic schedulability analysis.	Propose a new federated scheduling algorithm for arbitrary deadline DAG tasks with WCETs longer than their periods. Evaluate the proposed algorithm and comparing it to other scheduling algorithms, effectively outperforming most of them in terms of acceptance ratio	Doesn't tackle the problem of resource wasting when using feder- ated scheduling or their new version of it.	The new federated algorithm is used when the deadlines are bigger than the periods, and the tasks have high densities (according to the classic federated scheduling algorithm). For the high densities tasks that have a deadline lower than their period, they use standard federated scheduling and for the low density tasks, they use EDF-FF.
[18]	The NP-hard aspect of DAG multi-core scheduling makes the optimal solutions (using ILP) time-consuming. Hence the researchers have looked at heuristic to have scalable solutions.	Uses Deep reinforcement learning to construct a model that attempts to learn the optimal (in terms of makespan) scheduling policy for DAG tasks and compares it to the mathematically optimal ILP method.	Doesn't compare the machine learning method with SOTA heuristics, but only compares with ILP.	They use a combination of Graph neural network with attention layers to better capture the structure and dependency information. They use the negative value of the makespan as the reward function to be maximized.

[17]	Current methods to allocate shared resources on multi-core real-time systems use either static analysis of tasks or heuristics which cannot represent all possible system usage scenarios, thus potentially producing higher WCETs and worse system schedulability.	They use Deep Reinforcement learning to propose a holistic scheduling and allocation framework and their model shows better schedulability than existing methods.	Only considers independent periodic tasks and also only considers even-EDF and even-RM when comparing schedulability performance.	The plaform model is a LLC architecture with a shared memory bus and the DRL model uses a dense network with the proximal policy optimization algorithm for training. The DRL model produces a timetriggered schedule table for each tasks' execution and each tasks' memory allocation.
[19]	A previous paper [20] introduced a fixed-priority scheduling algorithm for DAG intra-task scheduling which performed better than SOTA but didn't extend the method to multi-DAG scheduling.	Extends the Concurrent Producer and Consumer (CPC) model from [20] to multi-DAG task scheduling by proposing a new multi-DAG scheduling algorithm based on a Parallelism-aware workload distribution model which outperforms existing methods in terms of system schedulability.	Considers constrained deadlines DAG tasks for the analysis but only consider implicit dead- lines DAG tasks for the experiment evaluation.	Uses a critical path first execution model by assigning the highest priorities to the node in the critical path and treating them as providers of parallel execution time for the 'consumers' which basically are the parrallelizable subtasks for each section of the critical path. For multi-DAG scheduling, they use a method similar to federated scheduling but instead of heavy and light tasks, they look at what they call the degree of parallelism of the DAG tasks.

[19]	Carranal harmi-ti	Duanaga a D	No samanani	The network :-
[13]	Several heuristics for DAG intratask scheduling have been used but no scalable optimal scheduling algorithm exists.	Propose a Deep Reinforcement Learning based machine learning model that computes an intra-task priority list for single DAG task scheduling which outperforms SOTA by up to 3% in terms of makespan	No comparison is made with ILP methods that lead to the mathematically minimum makespan. Also, it doesn't show the scalability of the GoSu DRL model when increasing the amount of cores/subtasks in the system/DAG task.	The network is comprised of a Graph Convolutional network to encode the graph information as well as a sequential decoder based on the attention-mechanism to produce a priority list. The model is trained using REINFORCE with the negative makespan as a reward function.
[9]	Virtual scheduling, using threads as virtual processors, has been considered in the past but never for DAG task scheduling. The similar federaded scheduling method suffers from resource wasting.	Use the concept of virtual processors to provide a virtually-federated scheduling algorithm which significantly outperforms other federated scheduling methods in terms of acceptance ratio.	Only considers implicit and constrained deadline DAG tasks and doesn't consider the running time overhead that the proposed method induces.	Introduce the concept of an active Virtual processor and a passive virtual processor (VP) and assign one of each on each physical core. The active VP execute the high priority (equivalent of high density in federated scheduling) tasks and the unused processing time of the active VP is treated as a passive VP on which low priority (equivalent of low density in federated scheduling) or high priority task can execute.

[1]	Most studies on	Propose a fluid	As for every	They first decom-
	DAG use the im-	scheduling based	fluid scheduling	pose each DAG
			O O	1
	plicity deadline but	algorithm for	based algo-	task into segments
	few are looking	constrained and	rithm the issue	of sequential tasks
	at DAGs with	arbitrary deadlines	of runtime	and then assign
	arbitrary deadline,	DAG tasks. In-	overhead is not	execution rates
	especially in the	troduce the first	entirely con-	to each tasks or
	case where the	capacity bound for	sidered as they	threads. Those
	deadline is greater	DAG with dead-	don't evaluate	two steps aim at
	than the task's	lines greater than	this metric.	producing a fluid
	period. Fluid	their periods. Their		schedule so that it
	scheduling has	algorithm performs		appears as though
	showed promising	better in terms of		each DAG task
	results but has only	acceptance ratio		is continuously
	been applied to	than SOTA (at the		running on the
	DAG tasks with	time of 2022).		cores.
	implicit deadlines.	,		
[4]	Most scheduling al-	Propose a DRL-	The reinforce-	It uses I/O usage
	gorithm that con-	based algorithm	ment learning	and ram allocation
	sider resource use	for task schedul-	algorithm uses	to construct a cost
	consider it as con-	ing on real-time	the previous	function which is
	straints rather than	simulation system	tasks' execution	used as a reward
	considering them as	called FRTDS. The	as experience	for the RL process.
	part of the schedul-	proposed algorithm	which implies	The current cost is
	ing decision pro-	performs better	a lot of mem-	combined with a fu-
	cess.	than existing al-	ory usage,	ture cost, predicted
		gorithm on the	thus affecting	using the successors
		FRTDS platform in	the speed of	of the current sub-
		terms of makespan	execution.	tasks.
		for single DAG task		-
		scheduling.		
		,		

[11] Federated sche	dul Dropogo a vintualla	Only considers	Introduce two vir-
' '	1 1	•	
	own federated schedul-	heavy tasks	tual processor per
great potential		and constrained	physical core, the
	AG efficiently tack-	deadline DAG	Active VP and
	fers les the resource	tasks.	Passive VP that
from a reso			are complementary.
wasting prol			The active VP
	peen ing while having		has the priority in
addressed but			terms of processing
limited extent.	tages as federated		capacity and for the
	scheduling. The		processing capacity
	algorithm performs		not used by an ac-
	better than existing		tive VP is given to
	algorithms in terms		the corresponding
	of acceptance ratio.		passive VP. Next,
			the active VP are
			allocated to tasks
			using the difference
			between their dead-
			line and the critical
			path's length as
			well as the minimal
			number of proces-
			sor to schedule a
			task in federated
			scheduling. Then
			the passive VPs
			are allocated to
			tasks according to
			how useful they are
			for scheduling the
			specific task.

[14]	The LET paradigm	Propose a DMA-	Doesn't con-	For the protocol, for
	is great at deal-	based protocol	sider the	each core, an LET
	ing with memory	to handle LET	scalability of	task is responsible
	contentions and	communications on	their method	for programming
	I/O determinism.	multicore systems	in terms of	the DMA engine
	But typical im-	that minimizes	the number	so that LET com-
	plementations of	the read/write	of tasks/data	munications can
	LET require local	latency of each	transfers to	happen. For the
	buffers for each core	tasks. Also, they	be scheduled.	scheduling al-
	which write/copy	provide a schedul-	Also, they	gorithm and data
	data to/from global	ing algorithm for	doesn't pro-	allocation, a Mixed-
	memory. This can	scheduling the com-	vide evaluation	Ingeteger Linear
	be costly when	munications using	of known or	Programming
	dealing with huge	MILP to provide	existing DAG	(MILP) problem is
	amount of data,	an optimal schedule	scheduling al-	solved to minimize
	like sensor data	which improved	gorithm using	either the number
	in autonomous	by up to 98% the	their LET	of DMA data com-
	driving system.	communication	communication	munications, or the
	Direct Memory	delays compared to	protocol com-	maximum commu-
	Access (DMA) is	the classic Giotto	pared to the	nication delay to
	a possible solution	approach of LET.	same algorithm	period ratio of each
	to such a problem		not using it.	tasks.
	but LET hasn't			
	yet been considered			
	with the DMA			
	protocol.			

Table 1: SLR summary table

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