

### Activity #1:

An if statement was used to determine if the rank was odd or even; the even ranks sent first and odds received first to keep from deadlock. This was put in a for loop for 5 iterations. The output is shown below:

```
act1.out
1
2 Rank: 3, my buffer: 2, my counter: 10
3
4
5 Rank: 4, my buffer: 5, my counter: 25
6
7
8 Rank: 0, my buffer: 1, my counter: 5
9
10
11 Rank: 2, my buffer: 3, my counter: 15
12
13
14 Rank: 7, my buffer: 6, my counter: 30
15
16
17 Rank: 6, my buffer: 7, my counter: 35
18
19 Q2
20 Rank: 9, my buffer: 8, my counter: 40
21
22
23 Rank: 5, my buffer: 4, my counter: 20
24
25
26 Rank: 8, my buffer: 9, my counter: 45
27
28
29 Rank: 1, my buffer: 0, my counter: 0
```

Fig. 1: Activity #1 output

Q1: What does process rank 5's counter store at the end of the computation?

- As shown in fig. 1, the counter for rank 5 was 20 at the end.

Q2: How many process ranks are used in the script above?

- The jobscript provided specified 10 ranks with the parameter: `--ntasks = 10`.
- This is also displayed in fig. 1 with 9 ranks displayed.

### Activity #2:

If statements were used to separate rank 0, even ranks, and odd ranks. In the case of rank 0, it would send to rank 1 and receive from the last rank ( $nproc-1$ ). In the case of odd ranks, they would receive first from the rank before them ( $my\_rank-1$ ), and even ranks would send first to the rank after them ( $(my\_rank+1)\%nprocs$ ). The use of  $\%nprocs$  was to have the last rank send to rank 0. The above was put in a loop for 10 iterations.

Originally, I would have used  $(my\_rank-1)\%nprocs$  in the receive to include rank 0 in the even ranks but c does not compute mod of -1 the way I thought it would.

The output is shown below:

```
≡ act2.out
1
2 Rank: 1, my buffer: 0, my counter: 0
3
4
5 Rank: 4, my buffer: 3, my counter: 30
6
7
8 Rank: 5, my buffer: 4, my counter: 40
9
10
11 Rank: 2, my buffer: 1, my counter: 10
12
13
14 Rank: 3, my buffer: 2, my counter: 20
15
16
17 Rank: 0, my buffer: 5, my counter: 50
```

Q3

Fig. 2: Activity #2 output

Q3: What does process rank 5's counter store at the end of the computation?

- As shown in fig. 2, the counter for rank 5 was 40 at the end.

**Activity #3:**

All ranks sent at the same time to  $(\text{my\_rank}+1)\%n\text{procs}$ . An if statement was only needed to separate rank 0 which received from  $n\text{procs}-1$ , and all other ranks which received from  $\text{my\_rank}-1$ . Output in fig. 3.

```
act3.out
1
2 Rank: 1, my buffer: 0, my counter: 0
3
4
5 Rank: 2, my buffer: 1, my counter: 10
6
7
8 Rank: 3, my buffer: 2, my counter: 20
9
10
11 Rank: 4, my buffer: 3, my counter: 30
12
13
14 Rank: 5, my buffer: 4, my counter: 40
15
16
17 Rank: 0, my buffer: 5, my counter: 50
18
```

Fig. 3: Activity #3 output

Q4: Comparing Programming Activities #2 and #3, which was easier to implement? Explain.

- Activity #3 was easier than act Activity #2 because using `MPI_Isend` prevents deadlock when all ranks send at the same time.

#### Activity #4

Rank 0 sends its counter to the first next\_rank determined by generateRandomRank. It broadcasts the current\_rank and the next\_rank to all other ranks.

Each current\_rank after rank 0, receives the counter and then tells rank 0 its current\_rank and next\_rank. Rank 0 broadcasts the new current\_rank and new next\_rank; then the current\_rank sends its new counter to the next\_rank.

Fig. 5 is a visual diagram of the process that may make more sense than the verbal description. Finally, activity #4 was not run on monsoon because a single node could not run 50 tasks and using 2 nodes took too much time.

```

≡ act4.out
1  Master: first rank: 14
2  My rank: 14, old counter: 0
3  My rank: 14, new counter: 14
4  My rank: 14, next rank to recv: 39
5  My rank: 39, old counter: 14
6  My rank: 39, new counter: 53
7  My rank: 39, next rank to recv: 37
8  My rank: 37, old counter: 53
9  My rank: 37, new counter: 90
10 My rank: 37, next rank to recv: 27
11 My rank: 27, old counter: 90
12 My rank: 27, new counter: 117
13 My rank: 27, next rank to recv: 23
14 My rank: 23, old counter: 117
15 My rank: 23, new counter: 140
16 My rank: 23, next rank to recv: 26
17 My rank: 26, old counter: 140
18 My rank: 26, new counter: 166
19 My rank: 26, next rank to recv: 5
20 My rank: 5, old counter: 166
21 My rank: 5, new counter: 171
22 My rank: 5, next rank to recv: 4
23 My rank: 4, old counter: 171
24 My rank: 4, new counter: 175
25 My rank: 4, next rank to recv: 11
26 My rank: 11, old counter: 175
27 My rank: 11, new counter: 186
28 My rank: 11, next rank to recv: 36
29 My rank: 36, old counter: 186
30 My rank: 36, new counter: 222
31 My rank: 36, next rank to recv: 34

```

Fig. 4: Activity #4 output

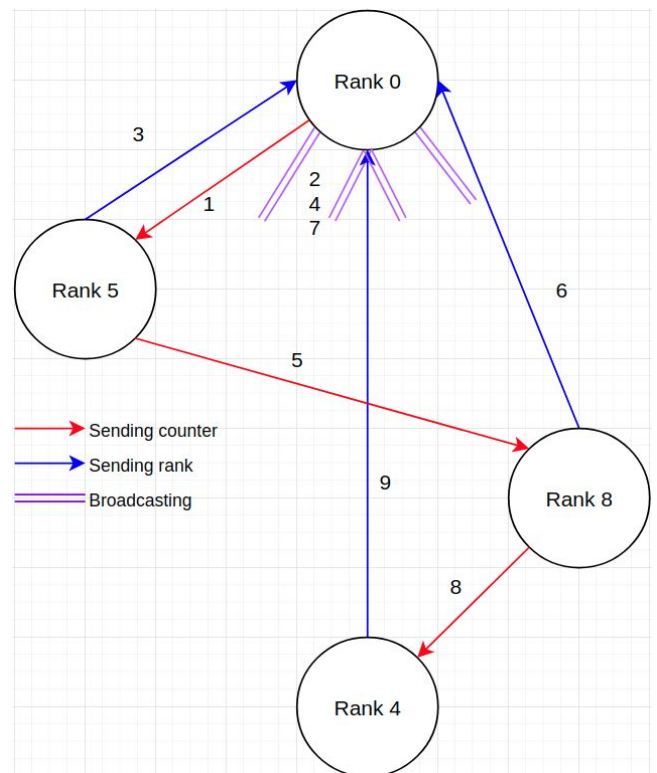


Fig. 5: Visualization of Activity #4

### Activity #5:

The current\_rank broadcasts to all other ranks the next\_rank, and only sends the counter to the next\_rank. The next\_rank uses MPI\_ANY\_SOURCE to receive the counter. Fig. 7 shows a visual representation of this process.

This activity was also not run on monsoon since 1 node could not run 50 tasks and 2 nodes took an unknown amount of time. If run with a smaller number of tasks the end result will not be 222.

```
act5.out
1  My rank: 0, next rank to recv: 14
2  My rank: 14, old counter: 0
3  My rank: 14, new counter: 14
4  My rank: 14, next rank to recv: 39
5  My rank: 39, old counter: 14
6  My rank: 39, new counter: 53
7  My rank: 39, next rank to recv: 37
8  My rank: 37, old counter: 53
9  My rank: 37, new counter: 90
10 My rank: 37, next rank to recv: 27
11 My rank: 27, old counter: 90
12 My rank: 27, new counter: 117
13 My rank: 27, next rank to recv: 23
14 My rank: 23, old counter: 117
15 My rank: 23, new counter: 140
16 My rank: 23, next rank to recv: 26
17 My rank: 26, old counter: 140
18 My rank: 26, new counter: 166
19 My rank: 26, next rank to recv: 5
20 My rank: 5, old counter: 166
21 My rank: 5, new counter: 171
22 My rank: 5, next rank to recv: 4
23 My rank: 4, old counter: 171
24 My rank: 4, new counter: 175
25 My rank: 4, next rank to recv: 11
26 My rank: 11, old counter: 175
27 My rank: 11, new counter: 186
28 My rank: 11, next rank to recv: 36
29 My rank: 36, old counter: 186
30 My rank: 36, new counter: 222
```

Fig. 6: Activity #5 output

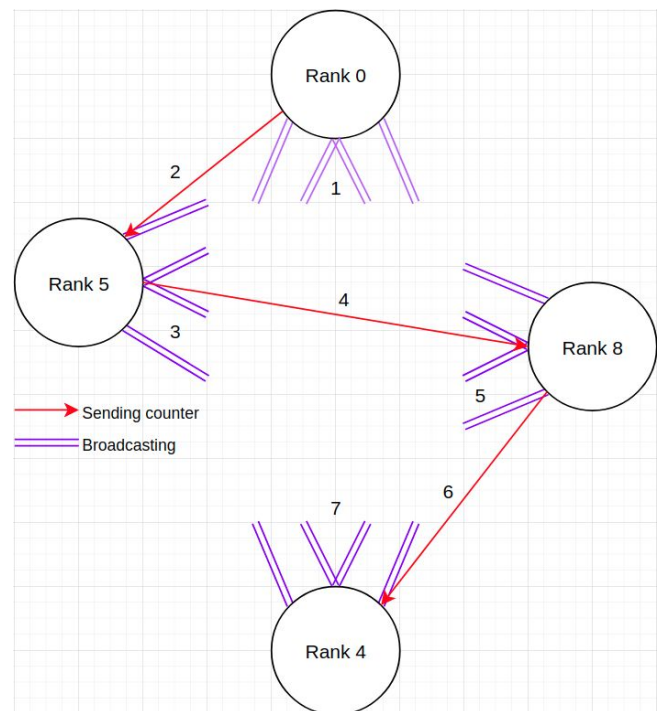


Fig. 7: Visualization of Activity #5

Q5: Comparing Programming Activities #4 and #5, which was easier to implement? Explain.

- Activity 5 was easier to implement because using MPI\_ANY\_SOURCE keeps the current\_rank from deadlocking even if the next\_rank doesn't know the previous rank it's receiving from.