

\*1. Analyze the logical forms of the following statements.

- Anyone who has forgiven at least one person is a saint.
- Nobody in the calculus class is smarter than everybody in the discrete math class.
- Everyone likes Mary, except Mary herself.
- Jane saw a police officer, and Roger saw one too.
- Jane saw a police officer, and Roger saw him too.

a)  $S(x) = x \text{ é um santo}$       b)  $C = \text{sala de calculo}$   
 $P(x, y) = x \text{ perdoou } y$        $D = \text{sala de matemáticas discritas}$   
 $\forall x (\exists y P(x, y) \rightarrow S(x))$        $I(x, y) = x \text{ é mais inteligente que } y$   
 $\neg \exists x (x \in C \wedge \forall y (y \in D \rightarrow I(x, y)))$

c)  $L(x, y) = x \text{ gosta de } y$

$\forall x (x \neq \text{Mary} \rightarrow L(x, \text{Mary}))$

d)  $P(x, y) = x \text{ viu uma poluição}$

$\exists y P(\text{Jane}, y) \wedge \exists z P(\text{Roger}, z)$

e)  $P(x, y) = x \text{ viu uma poluição}$   
 $\exists y (P(\text{Jane}, y) \wedge P(\text{Roger}, y))$

2. Analyze the logical forms of the following statements.

- Anyone who has bought a Rolls Royce with cash must have a rich uncle.
- If anyone in the dorm has the measles, then everyone who has a friend in the dorm will have to be quarantined.
- If nobody failed the test, then everybody who got an A will tutor someone who got a D.
- If anyone can do it, Jones can.
- If Jones can do it, anyone can.

c)  $F(x) = x \text{ falhou no teste}$

$O(x, y) = x \text{ tem } y$

$T(x, y) = x \text{ instruiu } y$

$$\neg \exists x F(x) \rightarrow \forall y (O(y, A) \rightarrow \exists z (O(x, D(z, y)) \wedge T(z, y)))$$

a)  $C(x, y, z) = x \text{ comprou } y \text{ utilizando } z$   
 $U(x, y) = x \text{ é primo de } y$   
 $R(x) = x \text{ é juiz}$   
 $\forall x (C(x, \text{Rolls Royce}, \text{dinheiro}) \rightarrow \exists y (U(y, x) \wedge R(y)))$

b)  $S(x) = x \text{ é um santo}$   
 $A(x, y) = x \text{ é imagem de } y$   
 $Q(x) = x \text{ sera intrometido}$   
 $D(x) = x \text{ é do dormitório}$   
 $\exists x (D(x) \wedge S(x)) \rightarrow (\forall y (\exists z (D(z) \wedge A(z, y)) \rightarrow Q(y))$

d)  $C(x) = x \text{ pode fumar MA}$   
 $\forall x (C(x) \rightarrow C(\text{fumar}))$

e)  $C(\text{fumar}) \rightarrow \forall x C(x)$

3. Analyze the logical forms of the following statements. The universe of discourse is  $\mathbb{R}$ . What are the free variables in each statement?

- Every number that is larger than  $x$  is larger than  $y$ .
- For every number  $a$ , the equation  $ax^2 + 4x - 2 = 0$  has at least one solution iff  $a \geq 2$ .
- All solutions of the inequality  $x^3 - 3x < 3$  are smaller than 10.
- If there is a number  $w$ , such that  $x^2 + 5x = w$  and there is a number  $y$  such that  $4 - y^2 = w$ , then  $w$  is strictly between -10 and 10.

a)  $\forall n (n > x \rightarrow n > y)$       y e  $x$  não variáveis livres

c)  $\forall x (x^3 - 3x < 3 \rightarrow x < 10)$       não ha variáveis livres

b)  $\forall a \exists x [ax^2 + 4x - 2 = 0 \leftrightarrow a \geq 2]$       não ha variáveis livres

d)  $\exists x (x^2 + 5x = w) \wedge \exists y (4 - y^2 = w) \rightarrow (-10 \leq w \leq 10)$       w é variável livre

4. Translate the following statements into idiomatic English.

- $\forall x [(H(x) \wedge \neg \exists y A(x, y)) \rightarrow O(x)]$ , where  $H(x)$  means "x is a man,"  $M(x, y)$  means "x is married to y," and  $O(x)$  means "x is unhappy."
- $\exists (P(x, x) \wedge S(x, y) \wedge W(y))$ , where  $P(x, x)$  means "x is a parent of x,"  $S(x, y)$  means "x and y are siblings," and  $W(y)$  means "y is a woman."

a) Para todo  $x$ , se  $x$  é homem e não existe  $y$  tal que  $x$  seja casado com  $y$ , então  $x$  é infeliz.

b) Existe um  $y$ , tal que  $y$  é parente de  $x$ ,  $y$  é irmão de  $x$  e  $y$  seja uma mulher.

5. Translate the following statements into idiomatic mathematical English.

- $\forall x [(P(x) \wedge \neg \exists y (x = y)) \rightarrow O(x)]$ , where  $P(x)$  means "x is a prime number" and  $O(x)$  means "x is odd."
- $\exists x [P(x) \wedge \forall y (P(y) \rightarrow y \leq x)]$ , where  $P(x)$  means "x is a perfect number."

a) Para todo  $x$ , se  $x$  é primo e é diferente de dois, então  $x$  é ímpar.

b) Existe um número  $x$  perfeito, tal que para qualquer outro número  $y$  que é menor que  $x$ ,  $y$  não é perfeito. (FIMITEZA DOS NÚMEROS PERFEITOS).

6. Translate the following statements into idiomatic mathematical English. Are they true or false? The universe of discourse is  $\mathbb{R}$ .

- $\neg \exists x (x^2 + 2x + 3 = 0 \wedge Ax^2 + 2x - 2 = 0)$
- $\neg [\exists x (x^2 + 2x + 3 = 0) \wedge \exists x (x^2 + 2x - 3 = 0)]$
- $\neg \exists x (x^2 + 2x + 3 = 0) \wedge \neg \exists x (x^2 + 2x - 3 = 0)$

a) Não existe  $x$  real tal que  $x^2 + 2x + 3 = 0$  e  $x^2 + 2x - 3 = 0$  simultaneamente

$$\frac{-2 \pm \sqrt{4 - 4 \cdot 3}}{2} = -1 \pm \sqrt{-2} = 1 \pm \sqrt{2}i$$

Verdadeiro, já não existe  $x \in \mathbb{R}$  que satisfaça essa exp.

b) Não é verdade que existe  $x$  real tal que  $x^2 + 2x + 3 = 0$  e outro  $x$  tal que  $x^2 + 2x - 3 = 0$

$$\frac{-2 \pm \sqrt{4 - 4 \cdot 1 \cdot (-3)}}{2} = -1 \pm 2 = -3 \text{ ou } 1$$

Verdadeiro, para realmente não existir  $x$  real que satisfaça a primeira exp. mesmo que para a segunda não.

c) Não existe  $x$  real que satisfaz  $x^2 + 2x + 3 = 0$  e não existe  $x$  real que satisfaz  $x^2 + 2x - 3 = 0$

Verdadeiro, para existir  $x$  real que satisfaz a exp. segundo igualdade

7. Are these statements true or false? The universe of discourse is the set of all people, and  $P(x, y)$  means "x is a parent of y."

- $\exists x \forall y P(x, y)$ .
- $\forall x \exists y P(x, y)$ .
- $\neg \exists x \exists y P(x, y)$ .
- $\exists x \neg \exists y P(x, y)$ .
- $\exists x \exists y \neg P(x, y)$ .

a) Existe  $x$  para todo  $y$ , tal que  $x$  é pai de  $y$ . (Falso, não existe um  $x$  que é pai de todo mundo).

b) Para todo  $x$  existe  $y$ , tal que  $x$  é pai de  $y$ . (Falso, nem todo mundo é pai).

c) Não existe  $x$  tal que para algum  $y$ ,  $x$  é pai. (Falso, existem pessoas que não são pais).

d) Existe  $x$  tal que não existe um  $y$  que  $x$  é pai. (Válido, existem pessoas que não são pais).

e) Existe  $x$  tal que existe  $y$ , tal que  $x$  não é pai de  $y$ . (Válido, existem pessoas que não tem relação pai-filho).

8. Are these statements true or false? The universe of discourse is  $\mathbb{N}$ .

- $\forall x \exists y (2x - y = 0)$ .
- $\exists y \forall x (2x - y = 0)$ .
- $\forall x \exists y (x - 2y = 0)$ .
- $\forall x \forall y (x < 10 \rightarrow \forall z (y < x \rightarrow y < 9))$ .
- $\exists y \exists z (y + z = 100)$ .
- $\forall x \exists y (y > x \wedge \exists z (y + z = 100))$ .

a) Válido

b) Falso

c) Se  $x = 3$  o y precisaria ser 1,5; então falso.

d) Válido

e) Válido

f) Falso, se pegarmos  $x = 100$  no mínimo y terá que ser 101 e não existe um  $y$  que satisfaça a condição

9. Same as exercise 8 but with  $\mathbb{Z}$  as the universe of discourse.

- $\forall x \exists y (2x - y = 0)$ .
- $\exists y \forall x (2x - y = 0)$ .
- $\forall x \exists y (x - 2y = 0)$ .
- $\forall x \forall y (x < 10 \rightarrow \forall z (y < x \rightarrow y < 9))$ .
- $\exists y \exists z (y + z = 100)$ .
- $\forall x \exists y (y > x \wedge \exists z (y + z = 100))$ .

a) Válido

b) Falso

c) Válido

d) Falso, no caso de  $x$  pode ser 9,9 e y 9,8.

e) Válido

f) Válido

10. Same as exercise 8 but with  $\mathbb{Z}$  as the universe of discourse.

- $\forall x \exists y (2x - y = 0)$ .
- $\exists y \forall x (2x - y = 0)$ .
- $\forall x \exists y (x - 2y = 0)$ .
- $\forall x \forall y (x < 10 \rightarrow \forall z (y < x \rightarrow y < 9))$ .
- $\exists y \exists z (y + z = 100)$ .
- $\forall x \exists y (y > x \wedge \exists z (y + z = 100))$ .

a) Válido

b) Falso

c) Se  $x = 3$  o y precisaria ser 1,5; então falso.

d) Válido

e) Válido

f) Válido

- \*1. Negate these statements and then reexpress the results as equivalent positive statements. (See Example 2.2.1.)  
 (a) Everyone who is majoring in math has a friend who needs help with his or her homework.  
 (b) Everyone has a roommate who dislikes everyone.  
 (c)  $\forall A \cup B \subseteq C \setminus D$ .  
 (d)  $\exists x \forall y [y > x \rightarrow \exists z (z^2 + 5z = y)]$ .

a)  $\forall x (M(x) \rightarrow \exists y (F(x, y) \wedge H(y)))$   
 $\exists x \neg (\neg M(x) \vee \exists y (F(x, y) \wedge H(y))) \equiv \exists x (M(x) \wedge \forall y \neg (F(x, y) \wedge H(y))) \equiv \exists x (M(x) \wedge \forall y (F(x, y) \rightarrow \neg H(y)))$   
 b)  $\forall x \exists y (R(x, y) \wedge \forall z \neg L(y, z))$   
 $\exists x \forall y \neg (R(x, y) \wedge \forall z \neg L(y, z)) \equiv \exists x \forall y (R(x, y) \rightarrow \exists z L(y, z))$   
 c)  $\forall x (x \in \{x \mid x \in A \wedge x \in B\} \rightarrow x \in \{x \mid x \in C \wedge x \notin D\})$   
 $\exists x \neg (\forall x \in \{x \mid x \in A \wedge x \in B\} \vee x \in \{x \mid x \in C \wedge x \notin D\}) \equiv \exists x (x \in \{x \mid x \in A \wedge x \in B\} \wedge \forall x \in \{x \mid x \in C \wedge x \notin D\})$   
 d)  $\forall x \exists y [\neg (y > x) \vee \exists z (z^2 + 5z = y)] \equiv \forall x \exists y [(y > x) \wedge \forall z (z^2 + 5z \neq y)]$

2. Negate these statements and then reexpress the results as equivalent positive statements. (See Example 2.2.1.)  
 (a) There is someone in the freshman class who doesn't have a room-mate.  
 (b) Everyone likes someone, but no one likes everyone.

- (c)  $\forall a \in A \exists b \in B (a \in C \leftrightarrow b \in C)$ .  
 (d)  $\forall y > 0 \exists x (ax^2 + bx + c = y)$ .

a)  $\exists x (F(x) \wedge \forall y \neg R(x, y))$   
 $\forall x \neg (F(x) \wedge \forall y \neg R(x, y)) \equiv \forall x (\neg F(x) \vee \exists y R(x, y)) \equiv \forall x (F(x) \rightarrow \exists y R(x, y))$   
 b)  $\forall x \exists y L(x, y) \wedge \forall x \forall y L(x, y)$   
 $\neg (\forall x \exists y L(x, y) \wedge \forall x \forall y L(x, y)) \equiv \exists x \forall y \neg L(x, y) \vee \exists x \forall y L(x, y)$   
 c)  $\forall a (\alpha \in A \rightarrow \exists b (\beta \in B \wedge (\alpha \in C \leftrightarrow \beta \in C)))$   
 $\exists a (\alpha \in A \wedge \forall b (\beta \in B \wedge (\alpha \in C \leftrightarrow \beta \in C) \wedge (\alpha \in C \vee \beta \in C))) \equiv \exists a (\alpha \in A \wedge \forall b (\beta \in B \wedge (\alpha \in C \vee \beta \in C) \wedge (\alpha \in C \wedge \beta \in C)))$   
 $\exists a (\alpha \in A \wedge \forall b (\beta \in B \vee \neg (\alpha \in C \vee \beta \in C) \vee \neg (\alpha \in C \wedge \beta \in C))) \equiv \exists a (\alpha \in A \wedge \forall b (\beta \in B \vee (\alpha \in C \wedge \beta \in C) \vee (\alpha \in C \wedge \beta \in C)))$   
 d)  $\forall y \exists x (y > 0 \rightarrow (ax^2 + bx + c = y))$   
 $\exists y \forall x (\neg (y > 0) \vee (ax^2 + bx + c = y)) \equiv \exists y \forall x (y > 0 \wedge (ax^2 + bx + c \neq y))$

3. Are these statements true or false? The universe of discourse is  $\mathbb{N}$ .  
 (a)  $\forall x (x < 7 \rightarrow \exists a \exists b (a^2 + b^2 + c^2 = x))$ .  
 (b)  $\exists a (a^2 + 3 = 4a)$ .  
 (c)  $\exists a (a^2 = 4a + 5)$ .  
 (d)  $\exists x \exists y (x^2 = 4x + 5 \wedge y^2 = 4y + 5)$ .

b)  $\Delta = 16 - 4 \cdot 1 \cdot 3 = 4$        $\frac{-(-4) \pm \sqrt{4}}{2} = 2 \pm 1 = 3 \text{ ou } 1$        $x^2 - 4x + 3 = 0$

a)  $1 = 1^2 + 0^2 + 0^2$        $3 = 1^2 + 2^2 + 2^2$        $5 = 2^2 + 1^2 + 0^2$   
 $2 = 1^2 + 1^2 + 0^2$        $4 = 2^2 + 0^2 + 0^2$        $6 = 2^2 + 1^2 + 1^2$

Verdadeiro, existe dois  $x \in \mathbb{N}$  que satisfazem a expressão.

d)  $x^2 - 4x - 5 = 0$        $y^2 - 4y - 5 = 0$

Verdadeiro, apesar uma das soluções é natural.

- \*4. Show that the second quantifier negation law, which says that  $\neg \forall x P(x)$  is equivalent to  $\exists x \neg P(x)$ , can be derived from the first, which says that  $\neg \exists x P(x)$  is equivalent to  $\forall x \neg P(x)$ . (Hint: Use the double negation law.)

5. Show that  $\neg \exists x \in A P(x)$  is equivalent to  $\forall x \in A \neg P(x)$ .

$\neg \exists x \in A P(x) = \neg \forall x (\neg P(x)) = \neg \forall x \neg (\neg P(x)) = \neg \neg \exists x \neg P(x) = \neg \neg \exists x P(x)$

- \*6. Show that the existential quantifier distributes over disjunction. In other words, show that  $\exists x (P(x) \vee Q(x))$  is equivalent to  $\exists x (P(x) \vee \exists y Q(x))$ . (Hint: Use the fact, discussed in this section, that the universal quantifier distributes over conjunction.)

$\neg \exists x \in A P(x) = \neg \exists x (x \in A \wedge P(x)) = \forall x \neg (x \in A \wedge P(x)) = \forall x (\neg x \in A \vee \neg P(x)) = \forall x (x \in A \rightarrow \neg P(x))$

$\exists x (P(x) \vee Q(x)) = \exists x \neg (\neg P(x) \wedge \neg Q(x)) = \forall x (\neg (\neg P(x)) \wedge (\neg Q(x))) = \neg \forall x (\neg P(x)) \wedge \neg \forall x (\neg Q(x)) =$

$\exists x \neg (\neg P(x)) \vee \exists x \neg (\neg Q(x)) = \exists x P(x) \vee \exists x Q(x) = \exists x P(x) \vee \exists x Q(x)$

7. Show that  $\exists x (P(x) \rightarrow Q(x))$  is equivalent to  $\forall x P(x) \rightarrow \exists x Q(x)$ .

$\exists x (P(x) \rightarrow Q(x)) = \exists x (\neg P(x) \vee Q(x)) = \exists x \neg P(x) \vee \exists x Q(x) = \neg \forall x P(x) \vee \exists x Q(x) = \forall x P(x) \rightarrow \exists x Q(x)$

- \*8. Show that  $\exists x \in A P(x) \wedge \exists x \in B P(x)$  is equivalent to  $\exists x \in (A \cup B) P(x)$ . (Hint: Start by writing out the meanings of the bounded quantifiers in terms of unbounded quantifiers.)

$\exists x (P(x) \wedge Q(x)) = \exists x \neg (\neg P(x) \vee \neg Q(x)) = \forall x (\neg \neg P(x) \wedge \neg \neg Q(x)) = \forall x (P(x) \wedge Q(x))$

$\forall x \in A P(x) \wedge \forall x \in B P(x) = \forall x [(x \in A \rightarrow P(x))] \wedge \forall x [(x \in B \rightarrow P(x))] = \forall x [(x \in A \wedge P(x)) \wedge (x \in B \wedge P(x))] = \forall x [(x \in A \wedge x \in B) \vee P(x)] =$

$\forall x [(\neg (x \in A \vee x \in B)) \vee P(x)] = \forall x [(x \in A \vee x \in B) \rightarrow P(x)] = \forall x (x \in A \vee B \rightarrow P(x)) = \forall x_{(A \cup B)} P(x)$

9. Is  $\forall x (P(x) \vee Q(x))$  equivalent to  $\forall x P(x) \vee \forall x Q(x)$ ? Explain. (Hint: Try assigning meanings to  $P(x)$  and  $Q(x)$ .)

Não é equivalente, por exemplo, se tomarmos:

$P(x) = x \text{ é macho}$

$Q(x) = x \text{ é fêmea}$

E considerando o universo de discussão todos os cães do mundo o significado seria:

$\forall x (P(x) \vee Q(x)) = \text{Sócio não é macho ou fêmea. (Verdadeiro)}$

$\forall x P(x) \vee \forall x Q(x) = \text{Sócio é macho ou todos são fêmeas. (Falso).}$

10. Show that  $\exists x \in A P(x) \vee \exists x \in B P(x)$  is equivalent to  $\exists x \in (A \cup B) P(x)$ .

- (b) Is  $\exists x \in A P(x) \wedge \exists x \in B P(x)$  equivalent to  $\exists x \in (A \cap B) P(x)$ ? Explain.

a)  $\exists x \in A P(x) \vee \exists x \in B P(x) = \exists x (x \in A \rightarrow P(x)) \vee \exists x (x \in B \rightarrow P(x)) = \exists x [(\neg x \in A \wedge P(x)) \vee (\neg x \in B \wedge P(x))] = \exists x (\neg x \in (A \cup B) \rightarrow P(x))$

b) Não é equivalente, pois existem um  $x$  de  $A$  que satisfaça  $P(x)$  e um  $x$  de  $B$  que não satisfaça  $P(x)$  não garante que elas tenham um elemento  $x$  em comum e que satisfaça  $P(x)$  simultaneamente.

- \*11. Show that the statements  $A \subseteq B$  and  $A \setminus B = \emptyset$  are equivalent by writing each in logical symbols and then showing that the resulting formulas are equivalent.

$A \subseteq B = \forall x (x \in A \rightarrow x \in B) = \forall x (\neg x \in A \vee x \in B) = \forall x \neg (x \in A \wedge \neg x \in B) = \neg \exists x (x \in A \wedge x \notin B) = \neg \exists x [x \in (A \setminus B)] = \neg \exists x (A \setminus B) = [A \setminus B = \emptyset]$

12. Show that the statements  $C \subseteq A \cup B$  and  $C \setminus A \subseteq B$  are equivalent by writing each in logical symbols and then showing that the resulting formulas are equivalent.

$C \subseteq A \cup B = \forall x (x \in C \rightarrow (x \in A \vee x \in B)) = \forall x (\neg x \in C \vee (x \in A \vee x \in B)) = \forall x \neg (x \in C \wedge \neg (x \in A \vee x \in B)) = \forall x \neg (x \in C \wedge x \notin (A \cup B)) = \forall x (x \in C \wedge x \in (B \setminus A)) = C \setminus A \subseteq B$

a)  $A \subseteq B = \forall x (x \in A \rightarrow x \in B) = \forall x ((x \in A \vee x \in B) \leftrightarrow x \in B) = (A \cup B = B)$

Relação de igualdade

b)  $A \subseteq B = \forall x (x \in A \rightarrow x \in B) = \forall x ((x \in A \wedge x \in B) \leftrightarrow x \in A) = (A \cap B = A)$

$\forall x (x \in P \leftrightarrow x \in Q) \equiv P = Q$

- \*14. Show that the statements  $A \cap B = \emptyset$  and  $A \cup B = A$  are equivalent.

$\exists x \in A \cap B = \forall x \neg (x \in A \wedge x \in B) = \forall x \neg (x \in A \wedge x \in B) = \forall x \neg (x \in A \wedge x \in B) = \forall x \neg (x \in A \wedge x \in B) = [A \cap B = \emptyset]$

a)  $\exists x \in A \cap B = \forall x \neg (x \in A \wedge x \in B) = \forall x \neg (x \in A \wedge x \in B) = \forall x \neg (x \in A \wedge x \in B) = \forall x \neg (x \in A \wedge x \in B) = [A \cap B = \emptyset]$

b) Existe professor que pertence a um aluno. Verdadeiro, como dito acima.

c) Existe um único professor que é professor de alguém. Falso, existem mais professores do que apenas um.

d) Existem alunos que pertencem a um único professor. Verdadeiro, em vez de pensar que tem apenas um aluno.

e) Existe apenas um professor que pertence a um único aluno. Falso, existe mais de um professor particulares,

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MAS APARENTEMENTE  
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f) Existe  $x \neq y$ , tal que  $x$  é professor de  $y$  e também não existe  $w$  tal que  $w$  seja professor de alguém  $v$  e  $v$  seja diferente de  $x$  e  $v$  diferente de  $y$ .

$$\exists x \exists y \{ T(x,y) \wedge [\forall v \forall w (\neg T(v,w) \vee \neg (v \neq x \vee w \neq y))] \} = \exists x \exists y T(x,y) \wedge \forall v \forall w (T(v,w) \rightarrow (v=x \wedge w=y)) = \exists! x \exists! y T(x,y)$$

$$\bullet \{ P(A) | A \in \mathcal{F} \} = \{ y | \exists x \in P(A) | A \in \mathcal{F} \} (y \in Y)$$

$$\bullet x \in \{ y | \exists x \in P(A) | A \in \mathcal{F} \} (y \in Y) = \exists y \in \{ z | \exists x \in \mathcal{F} (z = P(A)) \} (x \in Y) = \exists A \in \mathcal{F} (x \in P(A)) = \exists A \in \mathcal{F} (x \subseteq A) = \exists A \in \mathcal{F} [\forall w (w \in x \rightarrow w \in A)]$$

Example 2.3.6. Analyze the logical forms of the following statements.

4.  $x \in \bigcup \{\mathcal{P}(A) | A \in \mathcal{F}\}$ .

Example 2.3.8. For this example our universe of discourse will be the set  $S$  of all students. Let  $L(x, y)$  stand for ' $x$  likes  $y$ ' and  $A(x, y)$  for ' $x$  admires  $y$ '. For each student  $x$ , let  $L_x$  be the set of all students that  $x$  likes. In other words  $L_x = \{t \in S | L(t, x)\}$ . Similarly, let  $A_x = \{t \in S | A(t, x)\}$  be the set of all students that  $x$  admires. Describe the following sets.

1.  $\bigcap_{x \in S} L_x$ .
2.  $\bigcup_{x \in S} L_x$ .
3.  $\bigcup_{x \in S} L_x \cup \bigcup_{x \in S} A_x$ .
4.  $\bigcup_{x \in S} (L_x \cup A_x)$ .
5.  $(\bigcap_{x \in S} L_x) \cap (\bigcap_{x \in S} A_x)$ .
6.  $\bigcap_{x \in S} (L_x \cap A_x)$ .
7.  $\bigcup_{B \in \mathcal{B}} L_B$ , where  $B = \bigcap_{x \in S} A_x$ .

- 1) Conjunto dos alunos que não gostam de todos os alunos.
  - 2) Conjunto dos alunos que não gostam de todos os alunos.
  - 3) Conjunto dos alunos que não gostam de todos os alunos, porém não admirados por ninguém.
  - 4) Conjunto dos alunos que alguém gosta, porém não admira.
  - 5) Conjunto dos alunos que é gostado por todo mundo e também admirado por todo mundo.
  - 6) Conjunto dos alunos que não tanto gostam quanto admirados por todo mundo.
  - 7) Conjunto dos alunos que algum aluno admira por todo mundo gosta.
- $B = \text{Conjunto dos alunos admirados por todo mundo}$

\*1. Analyze the logical forms of the following statements. You may use the symbols  $\in, \notin, =, \neq, \wedge, \vee, \rightarrow, \leftrightarrow, \forall, \exists$  in your answers, but not  $\subseteq, \mathcal{P}, \cap, \cup, \setminus, \setminus$ , or  $\neg$ . (Thus, you must write out the definitions of some set theory notation, and you must use equivalences to get rid of any occurrences of  $\neg$ )

- (a)  $\mathcal{F} \subseteq \mathcal{P}(A)$ .
- (b)  $A \subseteq \{2n+1 | n \in \mathbb{N}\}$ .
- (c)  $\{n^2 + n + 1 | n \in \mathbb{N}\} \subseteq \{2n+1 | n \in \mathbb{N}\}$ .
- (d)  $\mathcal{P}(\bigcup_{i \in I} A_i) \not\subseteq \bigcup_{i \in I} \mathcal{P}(A_i)$ .

$$a) \mathcal{F} \subseteq \mathcal{P}(A) \equiv \forall X (X \in \mathcal{F} \rightarrow X \in \mathcal{P}(A)) = \forall X (X \in \mathcal{F} \rightarrow X \subseteq A) \equiv \forall X (X \in \mathcal{F} \rightarrow \forall x (x \in X \rightarrow x \in A))$$

$$b) A \subseteq \{2n+1 | n \in \mathbb{N}\} \equiv A \subseteq \{m | \exists n \in \mathbb{N} (m = 2n+1)\} \equiv \forall x (x \in A \rightarrow x \in \{m | \exists n \in \mathbb{N} (m = 2n+1)\}) \equiv \forall x (x \in A \rightarrow \exists n \in \mathbb{N} (x = 2n+1))$$

$$c) \{n^2 + n + 1 | n \in \mathbb{N}\} \subseteq \{2n+1 | n \in \mathbb{N}\} \equiv \{p | \exists n \in \mathbb{N} (p = n^2 + n + 1)\} \subseteq \{q | \exists n \in \mathbb{N} (q = 2n+1)\} \equiv \forall x [x \in \{p | \exists n \in \mathbb{N} (p = n^2 + n + 1)\} \rightarrow x \in \{q | \exists n \in \mathbb{N} (q = 2n+1)\}] \equiv \forall x [x \in \{n \in \mathbb{N} | n^2 + n + 1 \leq 2n+1\} \rightarrow \exists n \in \mathbb{N} (x = 2n+1)]$$

$$d) \mathcal{P}(\{x | \exists X \in A_i (x \in X)\}) \not\subseteq \{y | \exists Y \in \mathcal{P}(A_i) (y \in Y)\} \equiv \neg [\forall Z (Z \subseteq \{x | \exists X \in A_i (x \in X)\}) \rightarrow Z \subseteq \{y | \exists Y \in \mathcal{P}(A_i) (y \in Y)\}] \equiv$$

$$\neg [\forall Z (Z \subseteq \{x | \exists X \in A_i (x \in X)\} \rightarrow \exists Y \in \mathcal{P}(A_i) (Z \subseteq Y))] \equiv \neg [\forall Z (\forall x \in Z \rightarrow \exists y \in Y (x \in A_i \wedge y \in A_i)) \rightarrow \exists Y \in \mathcal{P}(A_i) (Z \subseteq Y)] \equiv$$

$$\neg \exists Z [\forall z \in Z \rightarrow \exists X \in A_i (z \in X)] \wedge \neg \exists Y \in \mathcal{P}(A_i) (Z \subseteq Y) \equiv \neg \exists Z [\forall z \in Z \rightarrow \exists X \in A_i (z \in X)] \wedge \neg \exists Y (Y \in \mathcal{P}(A_i) \wedge Z \subseteq Y) \equiv$$

$$\neg \exists Z [\forall z \in Z \rightarrow \exists X \in A_i (z \in X)] \wedge \neg \exists Y (Y \in \mathcal{P}(A_i) \wedge Z \subseteq Y) \equiv \neg \exists Z [\forall z \in Z \rightarrow \exists X \in A_i (z \in X)] \wedge \neg \exists Y (Y \in \mathcal{P}(A_i) \wedge \forall z \in Z (z \in Y)) \equiv$$

$$2. \text{ Analyze the logical forms of the following statements. You may use the symbols } \in, \notin, =, \neq, \wedge, \vee, \rightarrow, \leftrightarrow, \forall, \exists \text{ in your answers, but not } \subseteq, \mathcal{P}, \cap, \cup, \setminus, \setminus, \text{ or } \neg. \text{ (Thus, you must write out the definitions of some set theory notation, and you must use equivalences to get rid of any occurrences of } \neg \text{)}$$

- (a)  $x \in \bigcup F \setminus \bigcup G$ .
- (b)  $\{x \in B | x \notin C\} \in \mathcal{P}(A)$ .
- (c)  $x \in \bigcap_{i \in I} (A_i \cup B_i)$ .
- (d)  $x \in (\bigcap_{i \in I} A_i) \cup (\bigcap_{i \in I} B_i)$ .

$$a) x \in \{p | \exists P \in \mathcal{P}(p \in P)\} \setminus \{q | \exists Q \in \mathcal{G} (q \in Q)\} \equiv x \in \{p | \exists P \in \mathcal{P}(p \in P)\} \wedge \neg x \in \{p | \exists Q \in \mathcal{G} (p \in Q)\} \equiv$$

$$x \in \{p | \exists P \in \mathcal{P}(p \in P)\} \wedge \neg x \in \{p | \exists Q \in \mathcal{G} (p \in Q)\} \equiv \exists P \in \mathcal{P} (x \in P) \wedge \neg \exists Q \in \mathcal{G} (x \in Q) \equiv \exists P \in \mathcal{P} (x \in P) \wedge \forall Q \in \mathcal{G} (x \notin Q)$$

$$b) \{x \in B | x \notin C\} \in \mathcal{P}(A) \equiv \{x | x \in B \wedge x \notin C\} \subseteq A \equiv \forall y [y \in \{x | x \in B \wedge x \notin C\} \rightarrow y \in A] \equiv \forall y [y \in B \wedge y \notin C \rightarrow y \in A]$$

$$c) x \in \bigcap_{i \in I} (A_i \cup B_i) \equiv x \in \{y | \forall i \in I (y \in A_i)\} \vee x \in \{y | \forall i \in I (y \in B_i)\} \equiv \forall i \in I (x \in A_i) \vee \forall i \in I (x \in B_i)$$

$$3. \text{ We've seen that } \mathcal{P}(\emptyset) = \{\emptyset\}, \text{ and } \{\emptyset\} \neq \emptyset. \text{ What is } \mathcal{P}(\{\emptyset\})?$$

$$P(\emptyset) = \{\emptyset, \{\emptyset\}\}$$

\*4. Suppose  $\mathcal{F} = \{\text{red, green, blue}, \text{orange, red, blue}, \text{purple, red, green, blue}\}$ . Find  $\bigcap \mathcal{F}$  and  $\bigcup \mathcal{F}$ .

$$\bigcap \mathcal{F} = \{\text{red, blue}\} \quad \bigcup \mathcal{F} = \{\text{red, green, blue, orange, purple}\}$$

5. Suppose  $\mathcal{F} = \{3, 7, 12\}, \{5, 7, 16\}, \{5, 12, 23\}\}$ . Find  $\bigcap \mathcal{F}$  and  $\bigcup \mathcal{F}$ .

$$\bigcap \mathcal{F} = \emptyset \quad \bigcup \mathcal{F} = \{3, 5, 7, 12, 16, 23\}$$

6. Let  $I = \{2, 3, 4, 5\}$ , and for each  $i \in I$  let  $A_i = \{i, i+1, i-1, 2i\}$ .

- (a) List the elements of all the sets  $A_i$ , for  $i \in I$ .
- (b) Find  $\bigcap_{i \in I} A_i$  and  $\bigcup_{i \in I} A_i$ .

7. Let  $P = \{\text{Johann Sebastian Bach, Napoleon Bonaparte, Johann Wolfgang von Goethe, David Hume, Wolfgang Amadeus Mozart, Isaac Newton, George Washington}\}$  and let  $Y = \{1750, 1751, 1752, \dots, 1759\}$ . For each  $y \in Y$ , let  $A_y = \{p \in P | \text{the person } p \text{ was alive at some time during the year } y\}$ . Find  $\bigcup_{y \in Y} A_y$  and  $\bigcap_{y \in Y} A_y$ .

8. Let  $I = \{2, 3\}$ , and for each  $i \in I$  let  $A_i = \{i, 2i\}$  and  $B_i = \{i, i+1\}$ .

- (a) List the elements of the sets  $A_i$  and  $B_i$  for  $i \in I$ .
- (b) Find  $\bigcap_{i \in I} (A_i \cup B_i)$  and  $(\bigcap_{i \in I} A_i) \cup (\bigcap_{i \in I} B_i)$ . Are they the same?
- (c) In parts (c) and (d) of exercise 2 you analyzed the statements  $x \in \bigcap_{i \in I} (A_i \cup B_i)$  and  $x \in (\bigcap_{i \in I} A_i) \cup (\bigcap_{i \in I} B_i)$ . What can you conclude from your answer to part (b) about whether or not these statements are equivalent?

$$a) A_2 = \{1, 2, 3, 4\} \quad A_3 = \{2, 3, 4, 6\} \quad A_4 = \{3, 4, 5, 8\} \quad A_5 = \{4, 5, 6, 10\}$$

$$b) \bigcap_{i \in I} A_i = \{4\} \quad \bigcup_{i \in I} A_i = \{1, 2, 3, 4, 5, 6, 8, 10\}$$

$$\begin{array}{ll} \text{Bach} [1750] & \text{Hume} [1750-1759] \\ \text{Bonaparte} [ ] & \text{Mozart} [1756-1791] \\ \text{Goethe} [1750-1759] & \text{Newton} [ ] \end{array} \quad \begin{array}{ll} \text{Washington} [1750-1759] & \bigcap_{y \in Y} A_y = \{\text{Goethe, Hume, Washington}\} \\ \bigcup_{y \in Y} A_y = \{\text{Goethe, Hume, Washington, Mozart, Bach}\} & \bigcap_{i \in I} (A_i \cup B_i) \cup (\bigcap_{i \in I} A_i \cup \bigcap_{i \in I} B_i) = \emptyset \cup \{3\} = \{3\} \end{array}$$

$$a) A_2 = \{2, 4\} \quad B_2 = \{2, 3\} \quad b) \bigcap_{i \in \{2, 3\}} (A_i \cup B_i) = \{3, 4\}$$

$$A_3 = \{3, 6\} \quad B_3 = \{3, 4\}$$

$$c) \text{ Não são equivalentes como foi demonstrado.}$$

$$\forall i \in I (x \in A_i) \vee \forall i \in I (x \in B_i) \not\equiv \forall i \in I (x \in A_i \vee x \in B_i)$$

9. (a) Analyze the logical forms of the statements  $x \in \bigcup_{i \in I} (A_i \setminus B_i)$ ,  $x \in (\bigcup_{i \in I} A_i) \setminus (\bigcup_{i \in I} B_i)$ , and  $x \in (\bigcup_{i \in I} A_i) \setminus (\bigcup_{i \in I} B_i)$ . Do you think that any of these statements are equivalent to each other?

(b) Let  $I$ ,  $A_i$ , and  $B_i$  be defined as in exercise 8. Find  $\bigcup_{i \in I} (A_i \setminus B_i)$ ,  $(\bigcup_{i \in I} A_i) \setminus (\bigcup_{i \in I} B_i)$ , and  $(\bigcup_{i \in I} A_i) \setminus (\bigcup_{i \in I} B_i)$ . Now do you think any of the statements in part (a) are equivalent?

Não são equivalentes

$$a) \bullet x \in \bigcup_{i \in I} (A_i \setminus B_i) \equiv x \in \{y \mid \exists i \in I (y \in A_i \wedge y \notin B_i)\} \equiv \exists i \in I (y \in A_i \wedge y \notin B_i \wedge x \in y)$$

$$\bullet x \in (\bigcup_{i \in I} A_i) \setminus (\bigcup_{i \in I} B_i) \equiv x \in \{y \mid \exists i \in I (y \in A_i \wedge y \notin B_i)\} \setminus \{z \mid \exists i \in I (z \in B_i \wedge z \in y)\} \equiv \\ x \in \{y \mid \exists i \in I (y \in A_i \wedge y \notin B_i)\} \wedge x \notin \{z \mid \exists i \in I (z \in B_i \wedge z \in y)\} \equiv \exists i \in I (x \in A_i \wedge x \notin B_i \wedge x \in y) \equiv \\ \exists i \in I (y \in A_i \wedge y \notin B_i) \wedge \forall z \in B_i (x \notin z) \equiv \exists i \in I (y \in A_i \wedge y \notin B_i) \wedge \forall z \in B_i (z \in B_i \rightarrow x \notin z)$$

$$\bullet x \in (\bigcup_{i \in I} A_i) \setminus (\bigcup_{i \in I} B_i) \equiv x \in \{y \mid \exists i \in I (y \in A_i \wedge y \notin B_i)\} \setminus \{z \mid \forall Z \in B_i (z \in Z)\} \equiv \\ x \in \{y \mid \exists i \in I (y \in A_i \wedge y \notin B_i)\} \wedge x \notin \{z \mid \forall Z \in B_i (z \in Z)\} \equiv \exists i \in I (x \in A_i \wedge x \notin B_i \wedge x \notin \{z \mid \forall Z \in B_i (z \in Z)\}) \equiv \\ \exists i \in I (y \in A_i \wedge y \notin B_i) \wedge \exists Z \in B_i (x \notin Z)$$

$$b) \quad A_2 = \{2, 4\} \quad B_2 = \{2, 3\} \quad U_{i \in \{2, 3\}} (A_i \setminus B_i) = \{4, 6\} \\ A_3 = \{3, 6\} \quad B_3 = \{3, 4\} \quad U_{i \in \{2, 3\}} A_i \setminus U_{i \in \{2, 3\}} B_i = \{6\} \\ A_2 \setminus B_2 = \{4\} \quad U_{i \in \{2, 3\}} A_i \setminus U_{i \in \{2, 3\}} B_i = \{2, 4, 6\} \\ A_3 \setminus B_3 = \{6\}$$

$$U_{i \in \{2, 3\}} (A_i \setminus B_i) = \{4, 6\} \\ U_{i \in \{2, 3\}} A_i \setminus U_{i \in \{2, 3\}} B_i = \{6\} \\ U_{i \in \{2, 3\}} A_i \setminus U_{i \in \{2, 3\}} B_i = \{2, 4, 6\}$$

Continuum não é equivalente

10. Give an example of an index set  $I$  and indexed families of sets  $\{A_i \mid i \in I\}$  and  $\{B_i \mid i \in I\}$  such that  $\bigcup_{i \in I} (A_i \cap B_i) \neq (\bigcup_{i \in I} A_i) \cap (\bigcup_{i \in I} B_i)$ .

$$A_0 = \{i, 2\} \quad A_0 = \quad B_0 = \\ B_0 = \{i, i+1\}$$

$i$	$A$	$B$	$A \cap B$	$I$	$U(A \cap B)$	$U(A \cap B)$
0	$\{0\}$	$\{0, 1\}$	$\{0\}$	$\{0\}$	$\{1, 2\}$	$\{1, 2\}$
1	$\{1, 2\}$	$\{1, 2\}$	$\{1, 2\}$	$\{0, 1\}$	$\{0, 1, 2\}$	$\{0, 1, 2\}$

\*Nó caso de  $I = \{0, 1\}$  funciona

11. Show that for any sets  $A$  and  $B$ ,  $\mathcal{P}(A \cap B) = \mathcal{P}(A) \cap \mathcal{P}(B)$ , by showing that the statements  $x \in \mathcal{P}(A \cap B)$  and  $x \in \mathcal{P}(A) \cap \mathcal{P}(B)$  are equivalent. (See Example 2.3.3.)

$$X \in P(A \cap B) \equiv X \subseteq A \cap B \equiv \forall x (x \in X \rightarrow (x \in A \wedge x \in B)) \equiv \forall x (x \in X \vee (x \in A \wedge x \in B)) \equiv$$

\*12. Give examples of sets  $A$  and  $B$  for which  $\mathcal{P}(A \cup B) \neq \mathcal{P}(A) \cup \mathcal{P}(B)$ .

$$A = \emptyset \quad P(A) = \{\emptyset\} \quad A = \{1\} \quad P(A) = \{\emptyset, \{1\}\} \\ B = \emptyset \quad P(B) = \{\emptyset\} \quad B = \emptyset \quad P(B) = \{\emptyset\} \\ A \cup B = \emptyset \quad P(A \cup B) = \{\emptyset\} \quad A \cup B = \{1\} \quad P(A \cup B) = \{\emptyset, \{1\}\} \\ P(A) \cup P(B) = \{\emptyset\} \quad P(A) \cup P(B) = \{\emptyset, \{1\}\}$$

$$A = \{1\} \quad P(A) = \{\emptyset, \{1\}\} \\ B = \{2\} \quad P(B) = \{\emptyset, \{2\}\} \\ A \cup B = \{1, 2\} \quad P(A \cup B) = \{\emptyset, \{1\}, \{2\}, \{1, 2\}\} \\ P(A) \cup P(B) = \{\emptyset, \{1\}, \{2\}\}$$

13. Verify the following identities by writing out (using logical symbols) what it means for an object  $x$  to be an element of each set and then using logical equivalence.

$$(a) \bigcup_{i \in I} (A_i \cup B_i) = (\bigcup_{i \in I} A_i) \cup (\bigcup_{i \in I} B_i).$$

$$(b) (\bigcap \mathcal{F}) \cap (\bigcap \mathcal{G}) = \bigcap (\mathcal{F} \cup \mathcal{G}).$$

$$(c) \bigcap_{i \in I} (A_i \setminus B_i) = (\bigcap_{i \in I} A_i) \setminus (\bigcap_{i \in I} B_i).$$

$$a) \bullet x \in \bigcup_{i \in I} (A_i \cup B_i) \equiv x \in \bigcup_{i \in I} (A_i \cup B_i) \equiv x \in \{y \mid \exists i \in I (y \in A_i \cup B_i)\} \equiv \\ \exists i \in I (x \in A_i \cup B_i) \equiv \exists i \in I (x \in A_i \vee x \in B_i) \equiv \exists i \in I (x \in A_i) \vee \exists i \in I (x \in B_i) \equiv \\ x \in \{y \mid \exists i \in I (y \in A_i)\} \vee x \in \{y \mid \exists i \in I (y \in B_i)\} \equiv x \in (\bigcup_{i \in I} A_i) \cup (\bigcup_{i \in I} B_i)$$

$$b) \quad x \in \bigcap (\mathcal{F} \cup \mathcal{G}) \equiv x \in \{w \mid \forall W \in (\mathcal{F} \cup \mathcal{G}) (w \in W)\} \equiv \forall W \in (\mathcal{F} \cup \mathcal{G}) (x \in W) \equiv \forall W [(\mathcal{F} \cup \mathcal{G}) \rightarrow x \in W] \equiv \forall W [\neg (\mathcal{F} \cup \mathcal{G}) \rightarrow \neg x \in W] \equiv \forall W [\neg \mathcal{F} \wedge \neg \mathcal{G}] \equiv$$

$$\forall W [(\neg \mathcal{F} \vee x \in W) \wedge (\neg \mathcal{G} \vee x \in W)] \equiv \forall W (\neg \mathcal{F} \vee x \in W) \wedge \forall W (\neg \mathcal{G} \vee x \in W) \equiv \forall W \neg \mathcal{F} \vee x \in W \equiv \forall W \neg \mathcal{G} \vee x \in W \equiv$$

$$x \in \{y \mid \forall W \neg \mathcal{F} (y \in W) \wedge \forall W \neg \mathcal{G} (y \in W)\} \equiv x \in \{y \mid \neg (\mathcal{F} \wedge \mathcal{G})\}$$

$$c) \quad x \in \bigcap_{i \in I} (A_i \setminus B_i) \equiv x \in \{y \mid \forall i \in I (y \in A_i \wedge y \notin B_i)\} \equiv \forall i \in I (x \in A_i \wedge x \notin B_i) \equiv \forall i \in I (x \in A_i) \wedge \forall i \in I (x \notin B_i) \equiv \\ \forall i \in I (x \in A_i) \wedge \forall i \in I \exists i \in I (x \in B_i) \equiv (x \in \{y \mid \forall i \in I (y \in A_i)\}) \wedge (x \in \{y \mid \exists i \in I (y \in B_i)\}) \equiv x \in \{(x \in \{y \mid \forall i \in I (y \in A_i)\}) \wedge (x \in \{y \mid \exists i \in I (y \in B_i)\})\} \equiv$$

$$x \in \{x \in \{y \mid \forall i \in I (y \in A_i)\} \wedge x \in \{y \mid \exists i \in I (y \in B_i)\}\}$$

$$d) \quad \text{Sometimes each set in an indexed family of sets has two indices. For this problem, use the following definitions: } I = \{1, 2\}, J = \{3, 4\}. \text{ For each } i \in I \text{ and } j \in J, \text{ let } A_{i,j} = \{i, j, i+j\}. \text{ Thus, for example, } A_{2,3} = \{2, 3, 5\}.$$

$$(a) \text{ For each } j \in J \text{ let } B_j = \bigcup_{i \in I} A_{i,j} = A_{1,j} \cup A_{2,j}.$$

$$(b) \text{ Find } \bigcap_{j \in J} B_j. \text{ (Note that, replacing } B_j \text{ with its definition, we could say that } \bigcap_{j \in J} B_j = \bigcap_{j \in J} (\bigcup_{i \in I} A_{i,j}).)$$

$$(c) \text{ Find } \bigcup_{i \in I} (\bigcap_{j \in J} A_{i,j}). \text{ (Hint: You may want to do this in two steps, corresponding to parts (a) and (b). Are } \bigcap_{j \in J} (\bigcup_{i \in I} A_{i,j}) \text{ and } \bigcup_{i \in I} (\bigcap_{j \in J} A_{i,j}) \text{ equal?)}$$

$$(d) \text{ Analyze the logical forms of the statements } x \in \bigcap_{j \in J} (\bigcup_{i \in I} A_{i,j}) \text{ and } x \in \bigcup_{i \in I} (\bigcap_{j \in J} A_{i,j}). \text{ Are they equivalent?}$$

$$a) \quad B_j = A_{1,j} \cup A_{2,j} \quad A_{1,3} = \{1, 3, 4\} \quad A_{2,3} = \{2, 3, 5\} \quad A_{1,4} = \{1, 4, 5\} \quad A_{2,4} = \{2, 4, 6\}$$

$$B_j = A_{1,3} \cup A_{2,3} = \{1, 2, 3, 4, 5\}$$

$$B_4 = A_{1,4} \cup A_{2,4} = \{1, 2, 4, 5, 6\}$$

$$b) \quad \bigcap_{j \in \{3, 4\}} B_j = \{1, 2, 4, 5\}$$

$$c) \quad \bigcap_{j \in \{3, 4\}} A_{i,j} = \{i, 3, i+3\} \cap \{i, 4, i+4\} = C_i = \{i\}$$

$$U_{i \in \{1, 2\}} C_i = \{1, 4\} \cup \{2, 3\} = \{1, 2, 3, 4\}$$

$$U_{i \in \{1, 2\}} A_{i,j} = \{1, j, 1+j\} \cup \{2, j, 2+j\} = D_j = \{1, 2, j, 1+j, 2+j\}$$

$$\bigcap_{j \in \{3, 4\}} D_j = \{1, 2, 3, 4, 5\} \cap \{1, 2, 4, 5, 6\} = \{1, 2, 4, 5\}$$

\* Não são iguais.

$$d) \quad x \in \bigcap_{j \in J} (\bigcup_{i \in I} A_{i,j}) \equiv x \in \bigcap_{j \in J} \{y \mid \exists i \in I (y \in A_{i,j})\} \equiv$$

$$x \in \{y \mid \forall j \in J (\exists i \in I (y \in A_{i,j}))\} \equiv \forall j \in J (\exists i \in I (y \in A_{i,j})) \equiv$$

$$y \in J (\exists i \in I (y \in A_{i,j})) \equiv y \in J \big[ \exists i \in I (y \in A_{i,j}) \big] \equiv$$

$$y \in J \big[ \exists i \in I (y \in A_{i,j}) \big] \rightarrow \exists i \in I (y \in A_{i,j})$$

$$x \in \bigcup_{i \in I} (\bigcap_{j \in J} A_{i,j}) \equiv x \in \bigcup_{i \in I} \{y \mid \forall j \in J (y \in A_{i,j})\} \equiv$$

$$x \in \{y \mid \exists i \in I (\forall j \in J (y \in A_{i,j}))\} \equiv \exists i \in I (\forall j \in J (y \in A_{i,j})) \equiv$$

$$\exists i \in I (\forall j \in J (y \in A_{i,j})) \equiv \exists i \in I \big[ \forall j \in J (y \in A_{i,j}) \big] \equiv$$

\* Não são equivalentes

15. (a) Show that if  $\mathcal{F} = \emptyset$ , then the statement  $x \in \bigcup \mathcal{F}$  will be false no matter what  $x$  is. It follows that  $\bigcap \emptyset = \emptyset$ .
- (b) Show that if  $\mathcal{F} = \emptyset$ , then the statement  $x \in \bigcap \mathcal{F}$  will be true no matter what  $x$  is. In a context in which it is clear what the universe of discourse  $U$  is, we might therefore want to say that  $\bigcap \emptyset = U$ . However, this has the unfortunate consequence that the notation  $\bigcap \emptyset$  will mean different things in different contexts. Furthermore, when working with sets whose elements are sets, mathematicians often do not use a universe of discourse at all. (For more on this, see the next exercise.) For these reasons, some mathematicians consider the notation  $\bigcap \emptyset$  to be meaningless. We will avoid this problem in this book by using the notation  $\bigcap \mathcal{F}$  only in contexts in which we can be sure that  $\mathcal{F} \neq \emptyset$ .

16. In Section 2.3 we saw that a set can have other sets as elements. When discussing sets whose elements are sets, it might seem most natural to consider the universe of discourse to be the set of all sets. However, as we will see in this problem, assuming that there is such a set leads to contradictions.

Suppose  $U$  were the set of all sets. Note that in particular  $U$  is a set, so we would have  $U \in U$ . This is not yet a contradiction; although most sets are not elements of themselves, perhaps some sets are elements of themselves. But it suggests that the sets in the universe  $U$  could be split into two categories: the unusual sets that, like  $U$  itself, are elements of themselves, and the more typical sets that are not. Let  $R$  be the set of sets in the second category. In other words,  $R = \{A \in U \mid A \notin A\}$ . This means that for any set  $A$  in the universe  $U$ ,  $A$  will be an element of  $R$  iff  $A \notin A$ . In other words, we have  $\forall A \in U(A \in R \leftrightarrow A \notin A)$ .

- (a) Show that applying this last fact to the set  $R$  itself (in other words, plugging in  $R$  for  $A$ ) leads to a contradiction. This contradiction was discovered by Bertrand Russell (1872–1970) in 1901, and is known as *Russell's paradox*.
- (b) Think some more about the paradox in part (a). What do you think it tells us about sets?

$$a) x \in \{y \mid \exists Y \in \mathcal{F} (y \in Y)\} \equiv \exists Y \in \mathcal{F} (x \in Y) \equiv \exists Y (Y \in \mathcal{F} \wedge x \in Y), \text{ falso pois não existe } Y \in \mathcal{F} \text{ ne } \mathcal{F} = \emptyset.$$

Logo  $\bigcup \mathcal{F} = \emptyset$ .

$$b) x \in \bigcap \mathcal{F} \equiv \exists z \{y \mid \forall Y \in \mathcal{F} (y \in Y)\} \equiv \forall Y (Y \in \mathcal{F} \rightarrow y \in Y) \equiv \forall Y (Y \in \emptyset \rightarrow y \in Y), \text{ não sempre verdadeiro}$$

pois  $Y \in \emptyset$  é falso e de uma falácia se seguir que tudo é verdade, logo

$\bigcap \mathcal{F}$  contém todos os elementos quando  $\mathcal{F} = \emptyset$ .

$$a) X \in R \equiv X \in \{Y \mid Y \in U \wedge Y \notin Y\} \equiv X \in U \wedge X \notin X$$

$$R \in R \equiv R \in U \wedge R \notin R$$

*Logo  $R \in U$  ja que é o conjunto de todos conjuntos é inconsistente*

*montrar que  $R \in R \equiv R \in R$  e  $R \notin R \equiv R \in R$ , uma contradição*

b) Significa que a teoria ingenua de conjuntos é insuficiente, precisamos seguir para a teoria ZFC de conjuntos para evitar contradições.