### Week 20: Search I

 ${\sf Martha\ Lewis}$  (based on slides from Raul Santos Rodriguez)

Have a look at ...

... Russell and Norvig (Ch. 3)

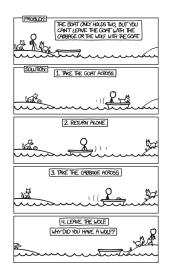
### Question

A farmer wants to get his cabbage, goat, wolf across a river. He has a boat that only holds two. He cannot leave cabbage and goat alone or goat and wolf alone.

#### How many river crossings does he need?

- **4**
- **5**
- **6**
- **7**
- No solution

### One solution...

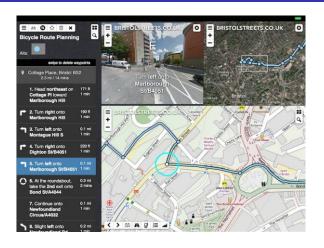


#### Outline

Search or how to find a sequence of actions that achieves a goal when no single action will do. We will cover

- Backtracking search
- Depth-first search
- Breadth-first search
- DFS with iterative deepening

## Application: route finding

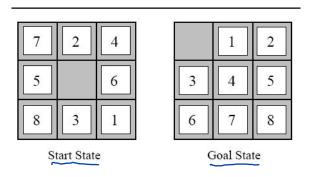


Actions: go straight, turn left, turn right

Objective: shortest? fastest? most scenic?



## Application: solving puzzles

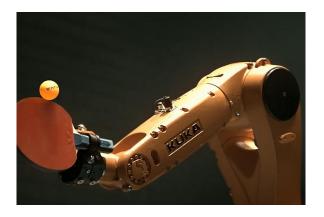


Actions: move pieces

8-puzzle

Objective: reach a certain configuration

## Application: robot planning



Actions: translate and rotate joints

Objective: fastest? most energy efficient? safest?

#### In all of these examples $\rightarrow$ sequence of actions

#### Idea

Decompose our very complex problem into small problems









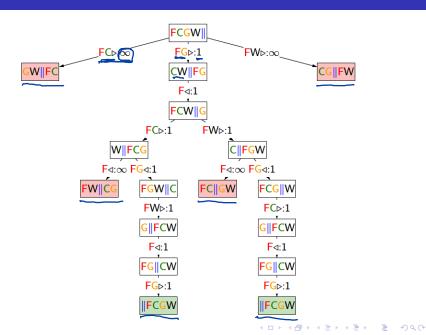
# Tree search: example



Farmer Cabbage Goat Wolf

#### Actions:

# Tree search: example



### Search problem

Search algorithms require a structure to keep track of the search tree that is being constructed:

#### Search

- s<sub>start</sub>: starting state
- Actions(s): possible actions
- Cost(s,a): action cost
- Succ(s,a): successors
- Goal(s): found solution?

### Measuring performance

Completeness: Is the algorithm guaranteed to find a solution when there is one?

Optimality: Does the strategy find the optimal solution?

Time complexity: How long does it take to find a solution?

Space complexity: How much memory is needed to perform the search?

#### Complexity

Complexity is expressed in terms of three quantities:

- b, the branching factor or maximum number of successors of any node;
- d, the depth of the shallowest goal node (i.e., the number of steps along the path from the root);
- D, the maximum length of any path in the state space.

### Backtracking search

#### BacktrackingSearch(s,path)

If Goal(s): update minimum cost path

**Else for each** action  $a \in Actions(s)$ :

Extend path with Succ(s,a) and Cost(s,a)

**Call** BacktrackingSearch(Succ(s,a), path)

Return minimum cost path

If b actions per state, maximum depth is D actions:

- Memory: O(D) (small)
- Time:  $O(b^D)$  (huge)



### Depth-first search

Idea: Backtracking search + stop when find the first goal state

Assumption: Action costs Cost(s,a) = 0

If b actions per state, maximum depth is D actions:

- Memory: O(D) (small)
- Time:  $O(b^D)$  (huge) but that is just for the worst case (could be much better if solutions are easy to find)

DFS is great when there are an abundance of solutions

### Depth-first search

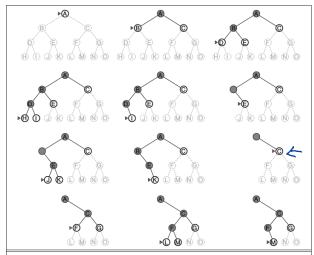


Figure 3.16 Depth-first search on a binary tree. The unexplored region is shown in light gray. Explored nodes with no descendants in the frontier are removed from memory. Nodes at depth 3 have no successors and M is the only goal node.

#### Breadth-first search

Idea: explore all nodes in order of increasing depth

Assumption: Action costs Cost(s,a)=c for some  $c \ge 0$ 

If b actions per state, maximum depth is d actions:

- Memory:  $O(b^d)$  (worse)
- Time:  $O(b^d)$  (depends on d instead of D)

#### Breadth-first search

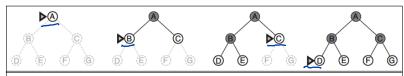


Figure 3.12 Breadth-first search on a simple binary tree. At each stage, the node to be expanded next is indicated by a marker.

# DFS with iterative deepening

Idea: Modify DFS to stop at a maximum depth

Idea: Call DFS for maximum depths 1,2,...

Assumption: Action costs Cost(s, a) = c for some  $c \ge 0$ 

If *b* actions per state, solution size *d*:

- Memory: O(d) (better!)
- Time:  $O(b^d)$  (same as BFS)

# DFS with iterative deepening

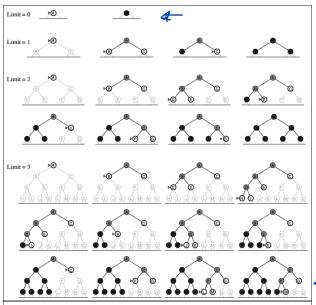


Figure 3.19 Four iterations of iterative deepening search on a binary tree.



# Comparison

Algorithm	Action costs	Space	Time
DFS	zero	O(D)	$O(b^D)$
BFS	constant	$O(b^d)$	$O(b^d)$
DFS-ID	constant	O(d)	$O(b^d)$
Backtracking	any	O(D)	$O(b^D)$

b actions per state, solution depth d, maximum depth D

### Summary

- Backtracking search traverse all nodes of the graph to find the optimal route to the goal
- Depth first search traverse a search tree by going deeper first
- Breadth-first search traverse a tree by going broader first
- Depth first search with iterative deepening increase the depth of the tree step by step
- Think about whether the algorithms are complete, optimal, and their complexity in time and space