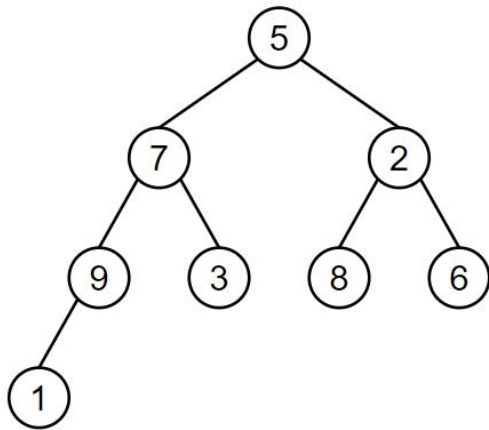


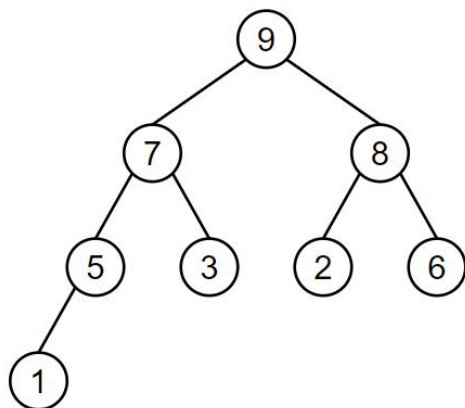
作业 8 参考答案 (by 况鸿翔)

P320-26

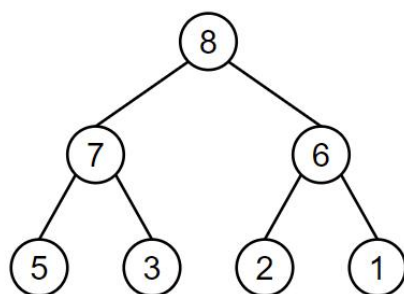
1° 原序列对应的完全二叉树

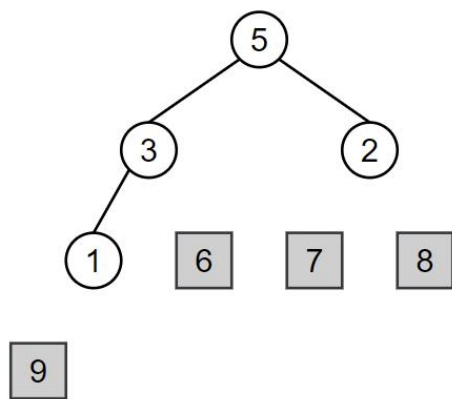
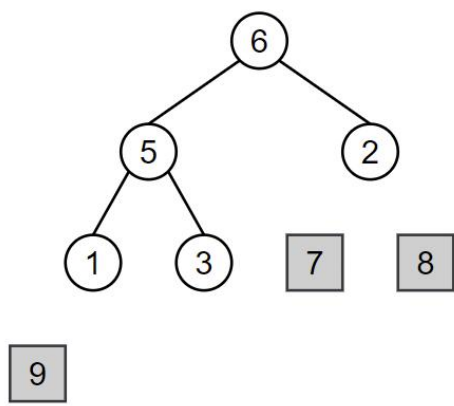
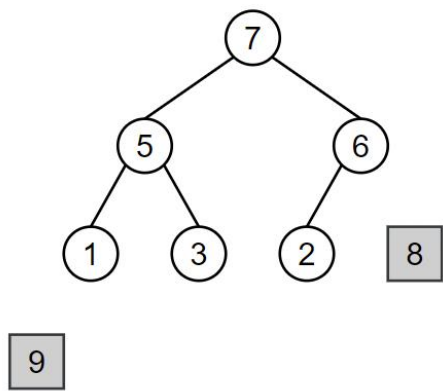


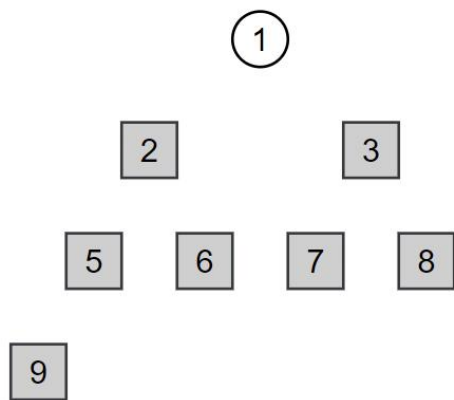
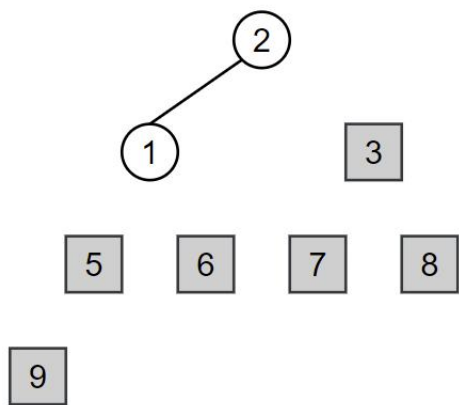
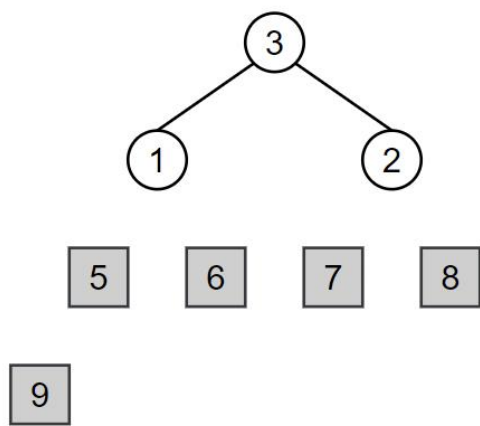
2° 堆化后的树

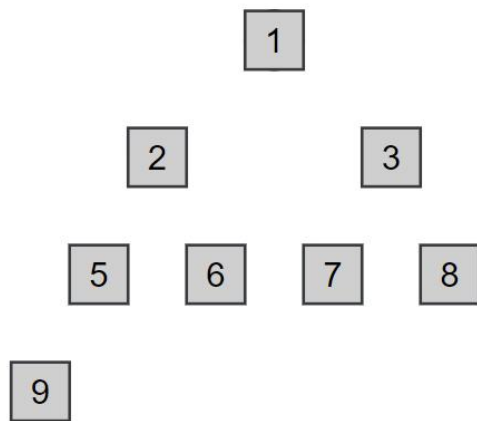


3° 大根堆的删除





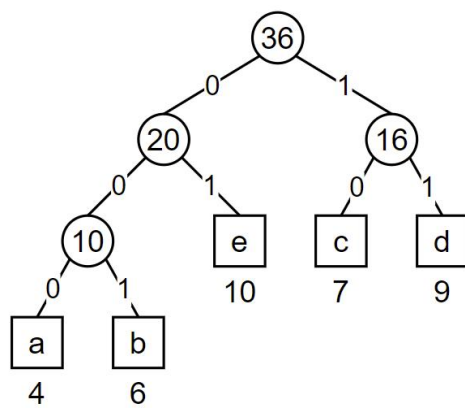




P321-40

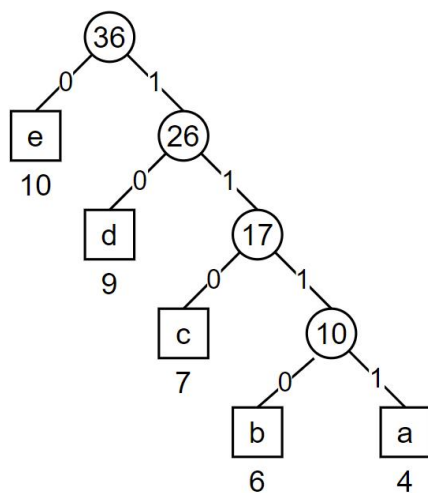
(1)

1° 霍夫曼树



代码: $a=000$, $b=001$, $c=10$, $d=11$, $e=01$
 长度: $WEP=3 \times 4 + 3 \times 6 + 2 \times 10 + 2 \times 7 + 2 \times 9 = 82$

2° 右偏树



代码: a=1111, b=1110, c=110, d=10, e=0

长度: WEP=4×4+4×6+3×7+2×9+1×10=89

(2)

设每个符号的频率为 p .

一方面, n 为 2 的幂时, 霍夫曼树是一颗高度为 $\log_2 n$ 的满二叉树, 故每个符号的霍夫曼编码长度均为 $\log_2 n$, 于是有

$$\text{WEP}_{\text{Huffman}} = p \cdot n \cdot \log_2 n$$

另一方面, 右偏树的高度为 $n-1$, 最低层有 2 个外部节点, 其余各层中除最高层外每层有 1 个外部节点, 于是有

$$\text{WEP}_{\text{RBT}} = p \cdot [\sum_{i=1}^{n-2} i + 2(n-1)] = p \cdot \frac{n^2+n-2}{2}$$

所以 Huffman 编码和 RBT 编码的 WEP 比率为

$$r = \frac{2n \log_2 n}{n^2+n-2}$$