Non Custodial Sidechains for Bitcoin utilizing Plasma Cash and Covenants

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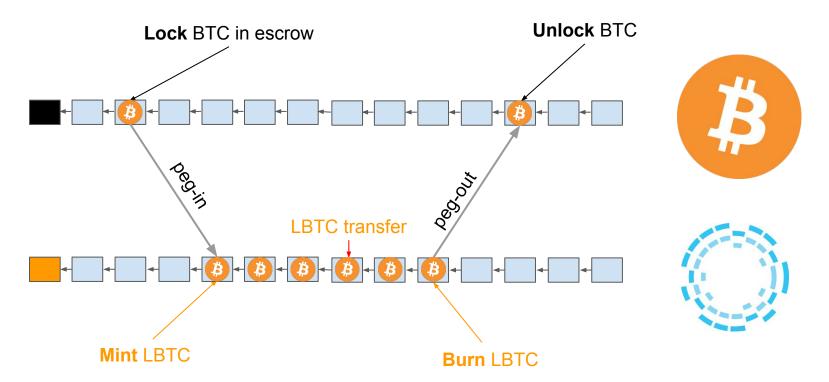
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Slides available: gakonst.com/scalingbitcoin2019.pdf

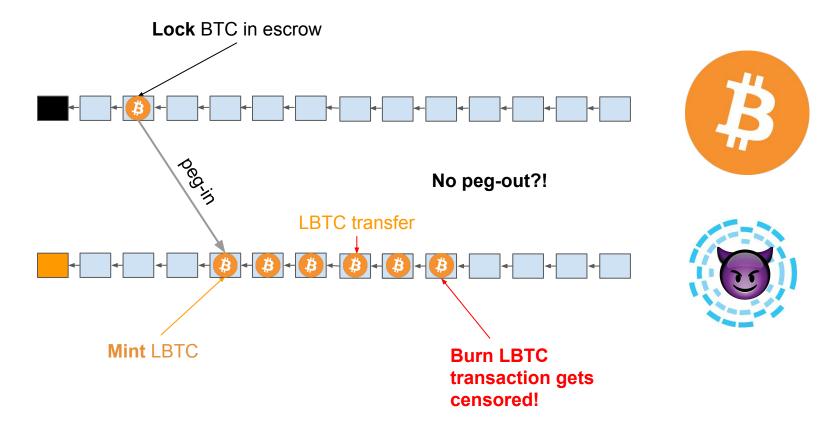
How do we scale?

- 1. Increase semantic density of transactions
 - (Segwit / MAST / Schnorr / Taproot / ... / Layer 2)
- 2. Bigger blocks

Sidechains considered harmful



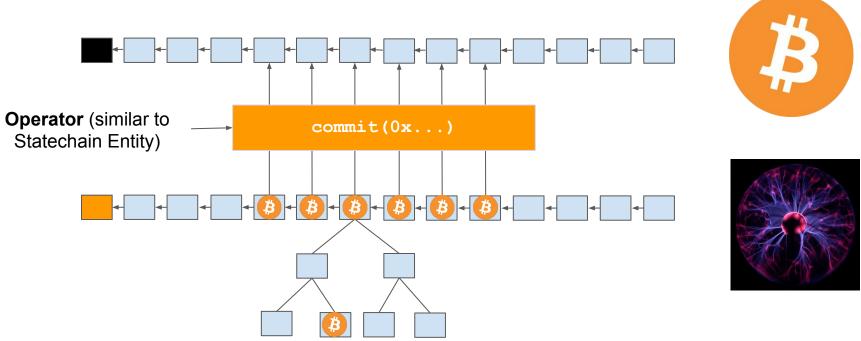
Sidechains considered harmful



Plasma Cash in 1 slide

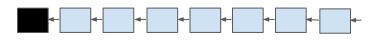
- Untrusted operator notarizes off-chain state
- Operator can censor but cannot steal
- Fixed denomination transfers like cash
- Safe under liveness assumption (~2KB stale state fraud proof)
- Watchtower compatible
- No overcollateralization requirements
- No need to sign to receive a transfer
- Can receive funds without on-chain transaction (no notion of inbound liquidity)

"Operator" commits* each block root to "parent chain"

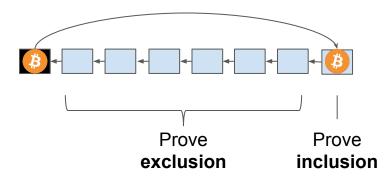


*uses accumulator that supports non-membership proofs e.g. ordered merkle tree

Prove coin history per transfer (off-chain)

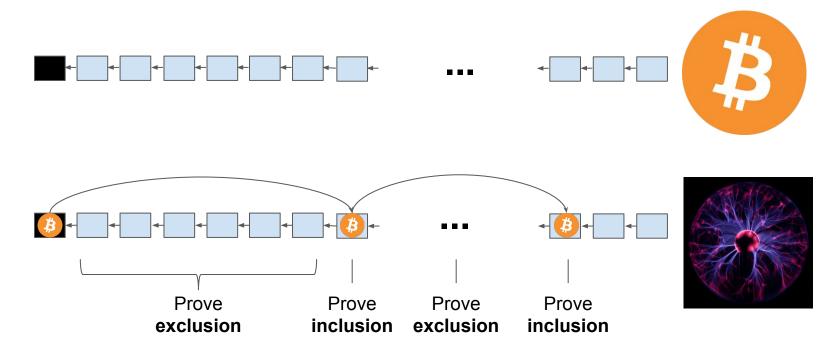






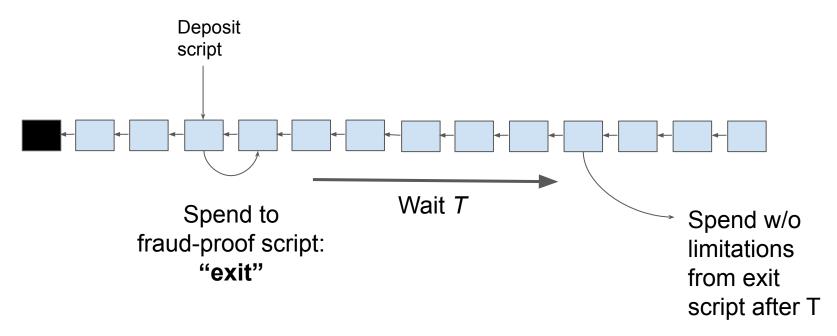


Prove coin history per transfer (off-chain)

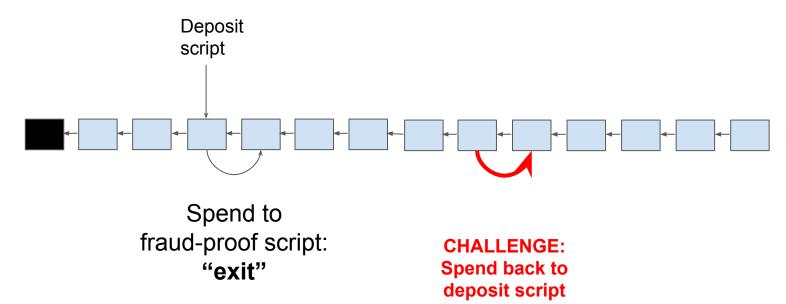


Coin history grows linearly with number of blocks TXO Commitments? RSA Accumulators?

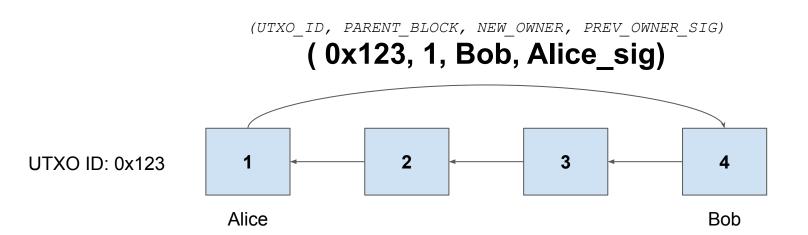
Exit Game: Delayed Withdrawals



Exit Game: Delayed Withdrawals

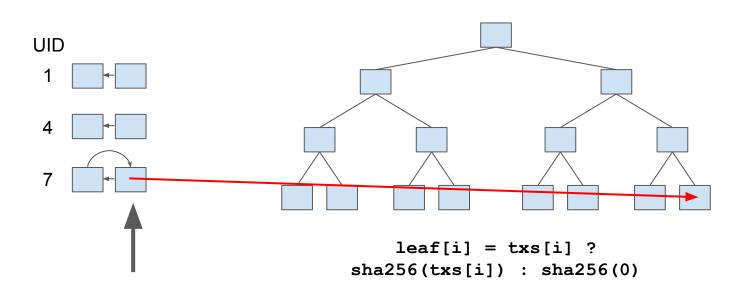


Transaction Format: 1 input 1 output UTXO



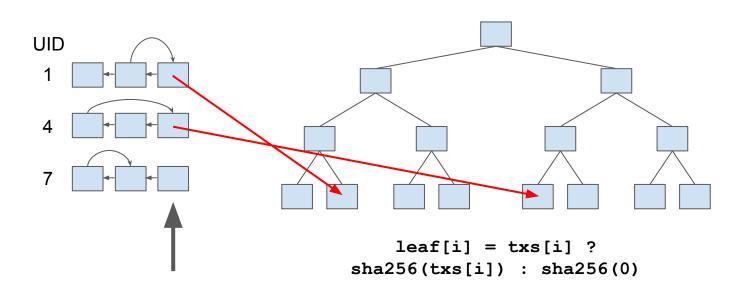
Merkle Tree: TxHash at each UTXO_ID index



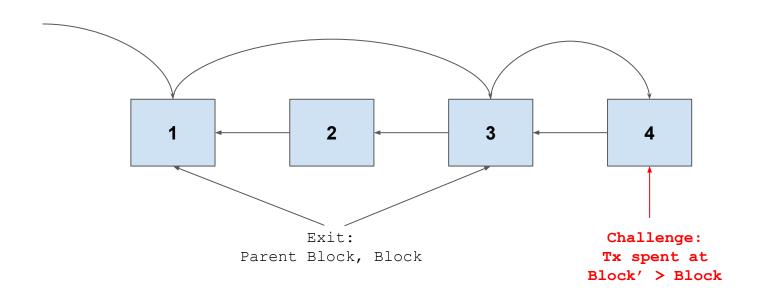


Merkle Tree: TxHash at each UTXO_ID index

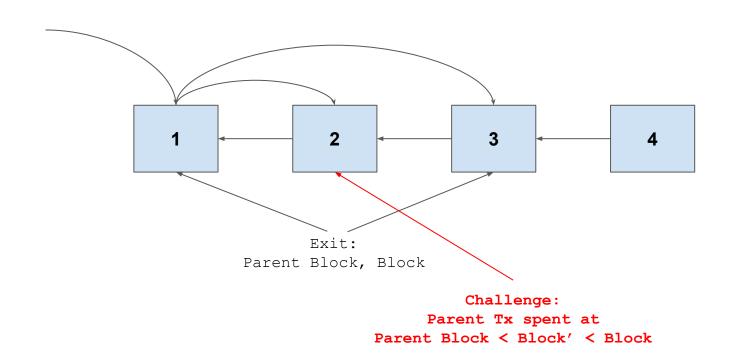
Current Block: 3



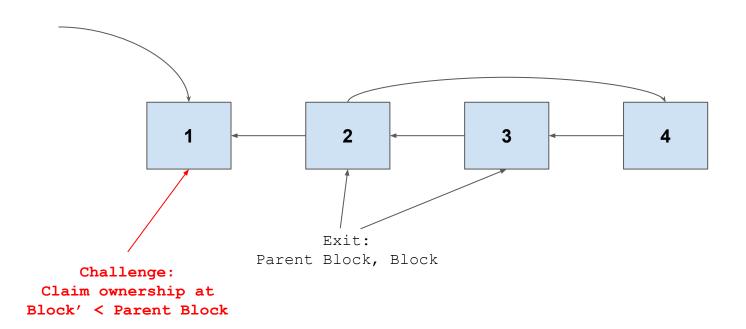
"Exit Spent Coin"



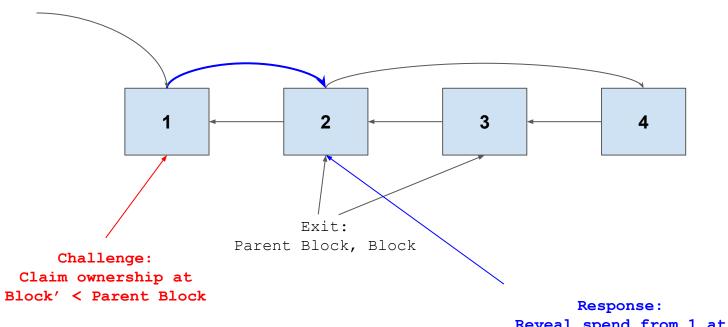
"Exit Double Spend" (malicious operator)



"Invalid History Challenge" (malicious operator)

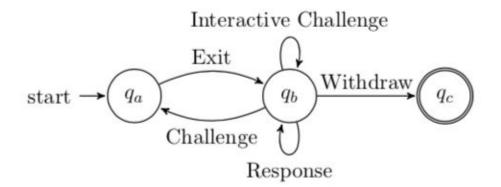


Response to Invalid History Challenge



Reveal spend from 1 at Block' < Block' <= Parent Block

Exit Game state machine (per coin)



Background literature on covenants

What is a covenant?

Restriction on the outputs spending a UTXO.

O'Connor @ Bitcoin Workshop 2017:

- Digital signatures: WHO can spend Bitcoin
- Timelocks: WHEN Bitcoin can be spent



What is a covenant?

Restriction on the outputs spending a UTXO.

O'Connor @ Bitcoin Workshop 2017:

- Digital signatures: WHO can spend Bitcoin
- Timelocks: WHEN Bitcoin can be spent
- Covenants: HOW and WHERE Bitcoin can be spent



Use Cases

- Vaults
- Paralysis Proofs
- Colored Coins (non-fungible tokens)
- Congestion Control
- Fraud proofs → Sidechains with trust-minimized reverse peg
- ...more in the <u>mailing list</u>

Covenant Designs

- OP_CHECKOUTPUT (MES'16)
- OP_CAT + OP_CHECKSIGFROMSTACK (O'Connor, Piekarska '17)
- OP_CHECKOUTPUTSHASH / OP_SECURETHEBAG (Rubin)
- OP_PUSHTXDATA (Lau)
- Presigned Transactions (McElrath / Bishop)

Implementing Plasma Cash on Bitcoin

Merkle Proof Verification

```
VerifyIncluded(UTXO_ID, ROOT, TX_HASH, PROOF):
    ROOT
    TX_HASH
    PROOF
    UTXO_ID
    MERKLEBRANCHVERIFY
```

Verify block root was signed by Operator

```
VerifySignedByOperator(BLOCK_NUM, ROOT, SIG):
    BLOCK_NUM
    ROOT
    CAT
    SIG
    <OPERATOR_ADDRESS>
    CHECKSIGFROMSTACKVERIFY
```

Verify transaction was signed by previous owner

```
VerifyTxSigned(tx)
   UTXO ID
   PARENT BLOCK NUM
   NEW OWNER
   CAT CAT SHA256
   SIG
   <PREV OWNER ADDRESS>
   CHECKSIGFROMSTACKVERIFY
```

Enforce UTXO is spent to next state

```
EnforceSpentTo (ARGS, NEXT_STATE_PATTERN):

ARGS

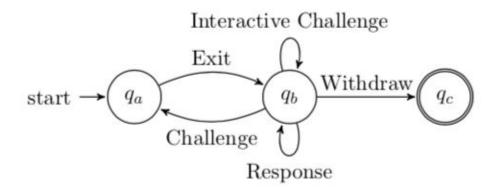
NEXT_STATE_PATTERN

CHECKOUTPUTVERIFY

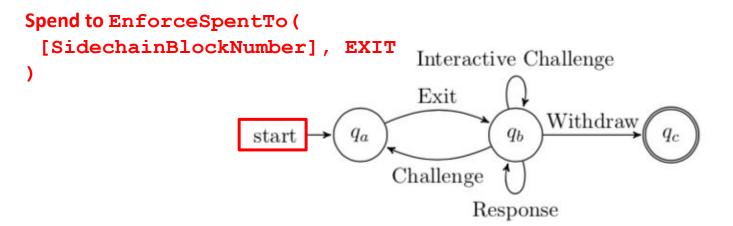
(use PICK etc. to put <ARGS> in OUTPUT_PATTERN / dynamically construct the covenant during redemption)
```

Putting it all together

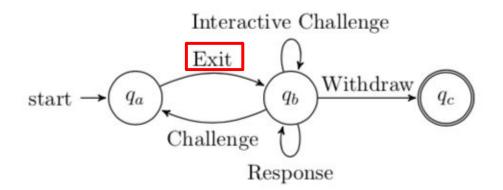
(verify tx merkle proofs, operator & transaction signatures w/ previous scripts)



Deposit = Spend to covenant

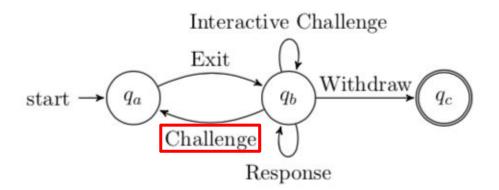


Exit = Spend from Deposit to Exit Script



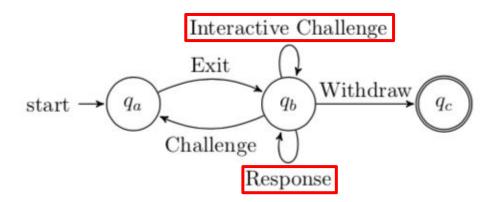
Spend to
EXIT (parentIncludedTx, includedTx)

Challenge Spent Coin / Double Spend = Spend back to Deposit



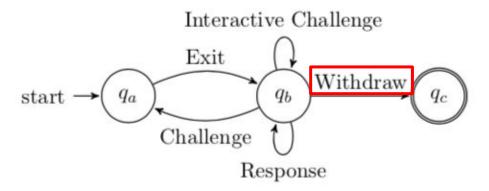
Spend to DEPOSIT, show includedTx according to exit game

Challenge Invalid History = Increment Counter, Response = Decrement Counter



Spend to CHALLENGED', show includedTx according to exit game. New CHALLENGED state = previous state with 1 extra IF condition for the Response.

Witdhraw = Spend wherever after T if counter = 0



CSV 150 BENEFICIARY_ADDRESS CHECKSIG

Summary

- Off-chain fixed-denomination payments
- "Compression" mechanism (more txs settle per block)
- Operator can censor cannot steal (under liveness assumption)
- No on-chain transaction to join
- Can receive payments when offline
- Capital efficient
- Users must audit Bitcoin-chain for fraud (light client side validation)
- Implementation WIP! Complex & Secure scripts are hard. Has been implemented on Ethereum since June '18.

Forked @kalewoof's btcdeb for prototyping (https://github.com/gakonst/btcdeb, rust CLI interpreter for opcode experimentation WIP!)

Thank you for your attention Q & A?

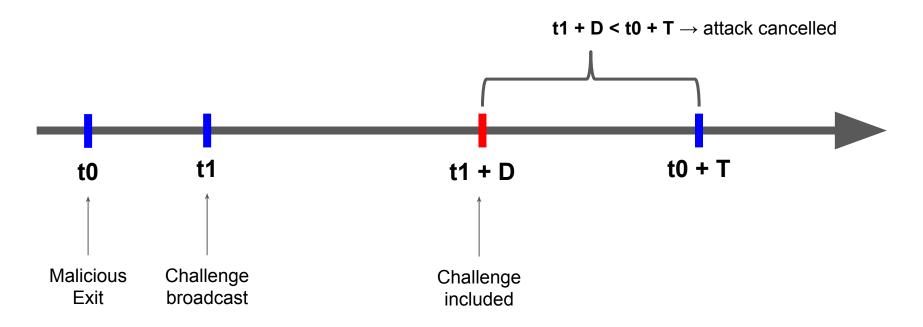
@gakonst / me@gakonst.com
gakonst.com/scalingbitcoin2019.pdf

Appendix

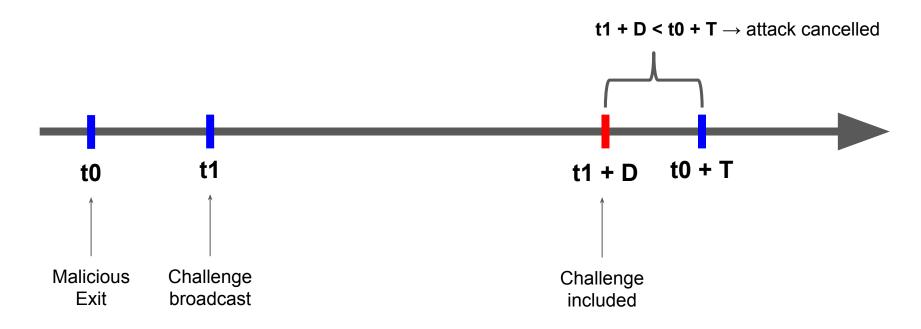
Security & Incentive Compatibility of Layer 2 games requirements*:

- liveness (somebody must challenge)
- expected reward of attacker <=0

Secure iff challenge included before t0 + T



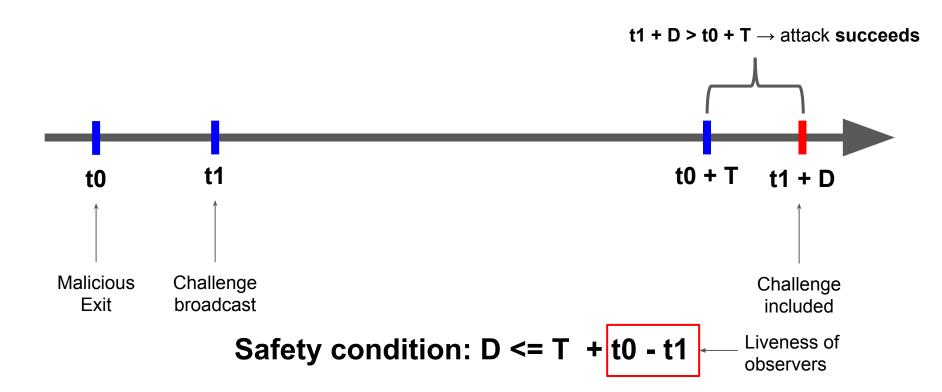
Secure iff challenge included before t0 + T



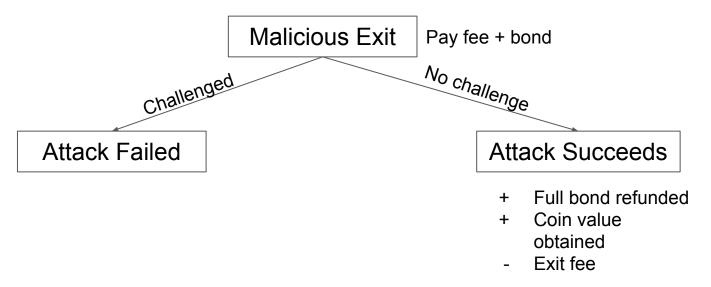
Insecure iff no challenge included before t0 + T



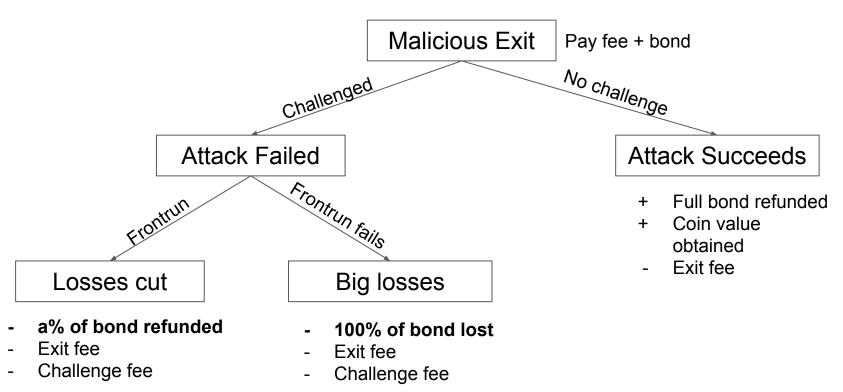
Insecure iff no challenge included before t0 + T



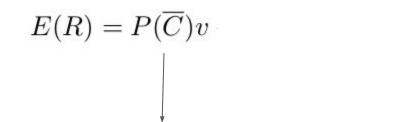
Attacker Decision Flow



Attacker Decision Flow



< 0



No challenges = success:

- ↑ onchain congestion / censorship
- ↑ block withholding
- ↓ liveness of participants
- ↓ challenge period T



$$E(R) = P(\overline{C})v - \underbrace{[gas + P(C) * bond]}_{cost \ to \ attack}$$

No challenges = success:

- † onchain congestion / censorship
- † block withholding
- ↓ liveness of participants
- ↓ challenge period T

Cost to Attack =

- Tx fees (constant)
- Fidelity Bond (goes to challenger)

< 0

Large T = Secure but bad UX!

$$E(R) = P(\overline{C})v - \underbrace{[gas + P(C) * bond]}_{cost \ to \ attack} + \underbrace{P(C)P(F \mid C) * bond}_{reward \ from \ frontrunning} \leq 0$$

$$\begin{array}{c} \text{challenges = success:} \\ \uparrow \text{ onchain congestion / censorship} \\ \uparrow \text{ block withholding} \\ \end{array}$$

$$\begin{array}{c} \text{Cost to Attack =} \\ \bullet \quad \text{Tx fees (constant)} \\ \bullet \quad \text{Fidelity Bond} \\ \end{array}$$
Frontrunning removes be

- No challenges = success:

 - ↓ liveness of participants
 - ↓ challenge period T

Large T = Secure but bad UX!

$$P(F \mid \overline{C}) = 0$$

(goes to challenger)

Attacker won't frontrun if nobody challenged

Frontrunning removes bond from cost if successful

$$E(R) = P(\overline{C})v - \underbrace{[gas + P(C) * bond]}_{cost \ to \ attack} + \underbrace{\alpha P(C)P(F \mid C) * bond}_{reward \ from \ frontrunning} \leq 0$$
No challenges = success:

• ↑ onchain congestion / censorship
• ↑ block withholding
• ↓ liveness of participants
• ↓ challenge period T

Cost to Attack =

• Tx fees (constant)
• Fidelity Bond (goes to challenger)

Frontrunning removes bond from cost if successful

Burn part of the bond.

$$P(F \mid \overline{C}) = 0$$

Attacker won't frontrun

if nobody challenged

Plasma Cash → Fixed-denomination. Arbitrary denomination payments?



https://github.com/plasma-group/plasma-core

Plasma Cash + Fragmentation = Plasma Cashflow

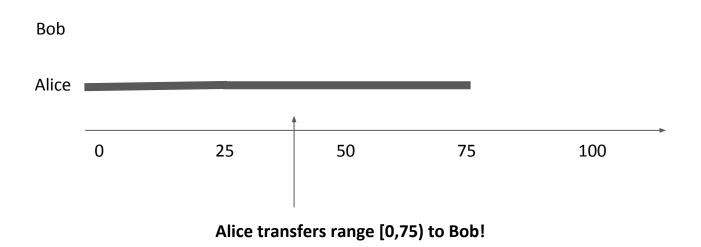


1 Euro

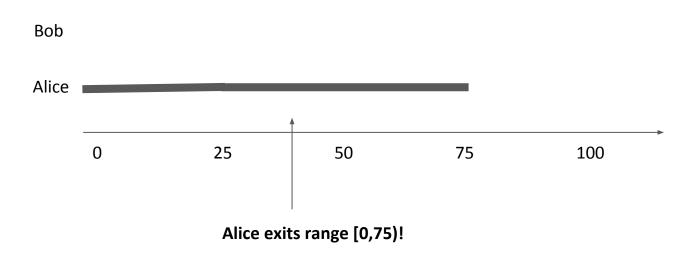
Plasma Cash + Fragmentation = Plasma Cashflow



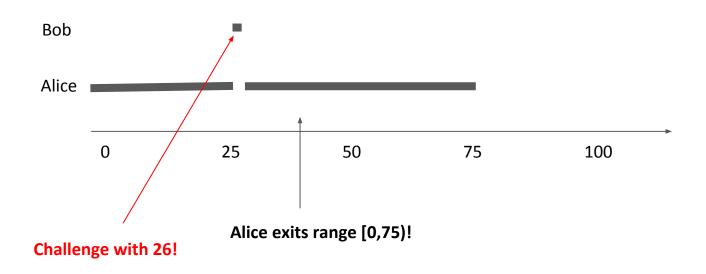
A non-interrupted range can be transferred in 1 tx



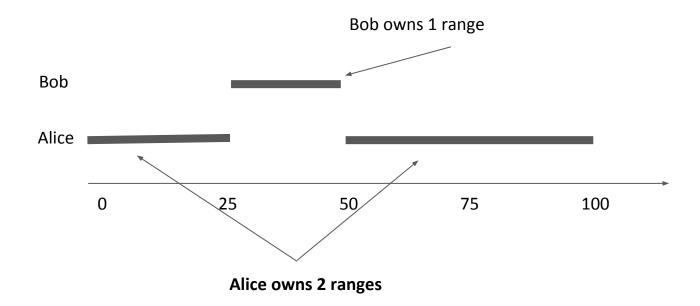
A non-interrupted range can be exited in 1 tx



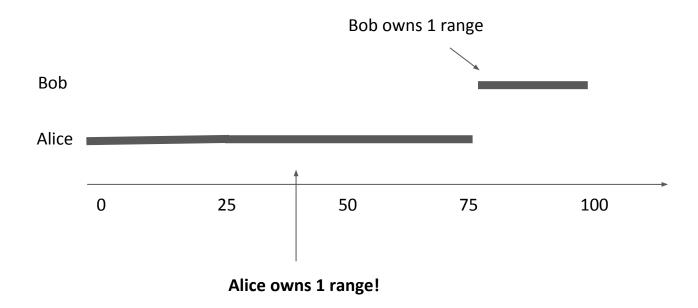
Any 1 coin inside the range is a valid challenge!



Defragmentation of ranges

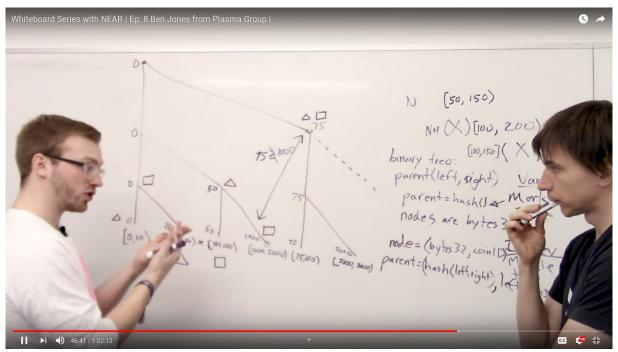


Defragmentation of ranges



Merkle Index Tree

Inclusion / exclusion proofs for ranges w/ light client support!



https://www.youtube.com/watch?v=-8Jp7VjspQE