

CMSC 141 Automata and Language Theory

Regular Languages

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Regular Expression Examples

RegEx

0^*10^*

$(0+1)^*1(0+1)^*$

$0(0+1)^*0+1(0+1)^*1+1+0$

$((0+1)(0+1))^*$

Regular Language

→ $\{w \mid w \text{ contains a single } 1\}$

→ $\{w \mid w \text{ contains at least one } 1\}$

→ $\{w \mid w \text{ begins and ends with the same symbol}\}$

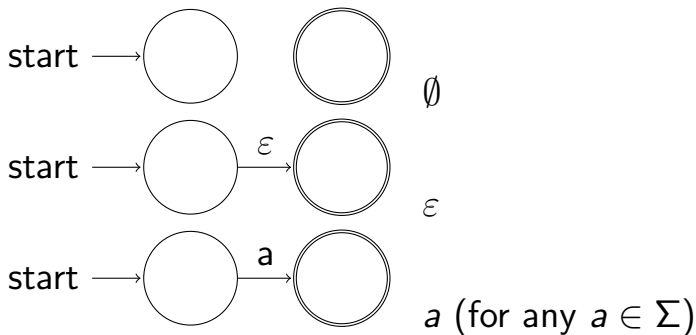
→ $\{w \mid w \text{ has even length}\}$

Regular Expression

- Like DFA and NFA, regular expressions can describe regular languages.
- Therefore, regular expressions can be converted into an equivalent DFA or NFA.

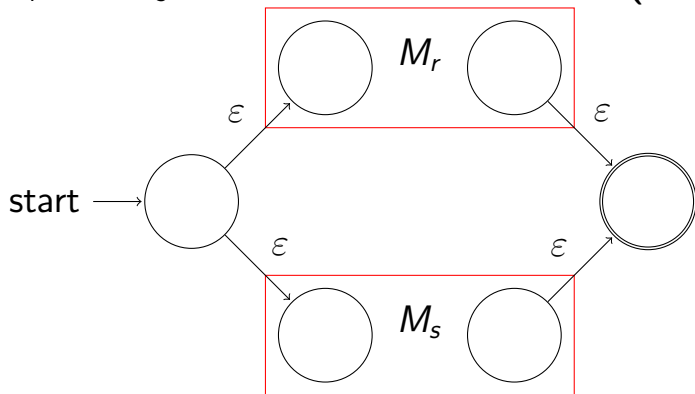
Converting RegEx to ϵ NFA

First, the constants or the basic cases:



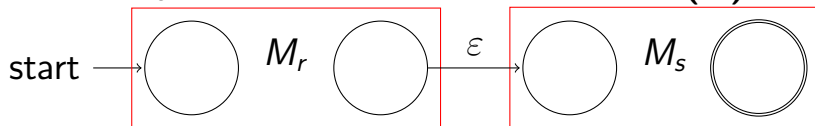
Regular Expression Operations

Let r, s be arbitrary regular expressions with NFAs M_r and M_s . The ϵ NFA for **alternation** ($r+s$):



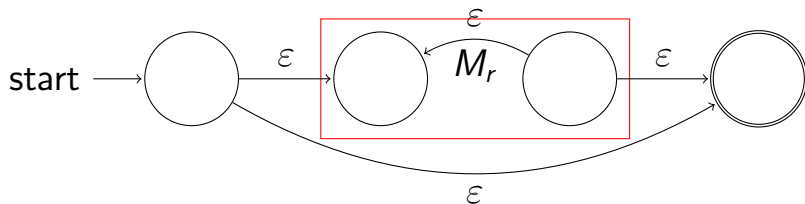
Regular Expression Operations

Let r, s be arbitrary regular expressions with NFAs M_r and M_s . The ε NFA for **concatenation (rs)**:



Regular Expression Operations

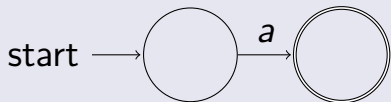
Let r be an arbitrary regular expression with NFA M_r . The ϵ NFA for **Kleene star** (r^*):



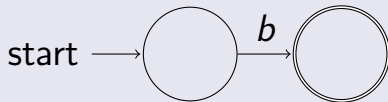
Example conversion of regex to ϵ NFA

$(ab+b)^*$

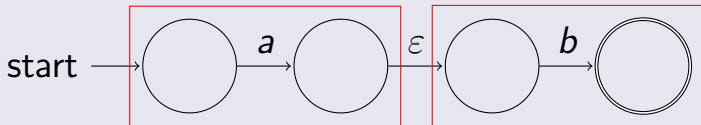
a



b



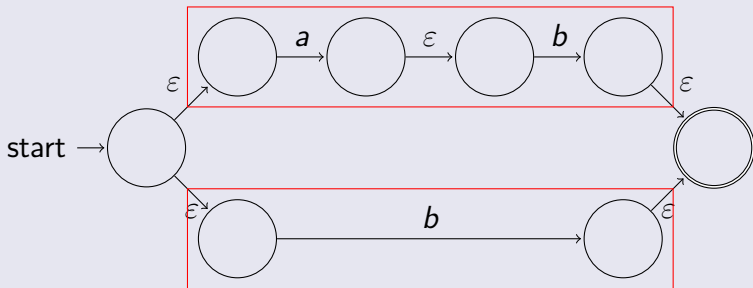
ab



Example conversion of regex to ϵ NFA

$(ab+b)^*$

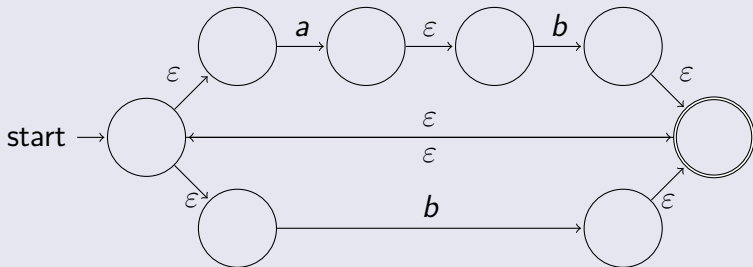
$ab+b$



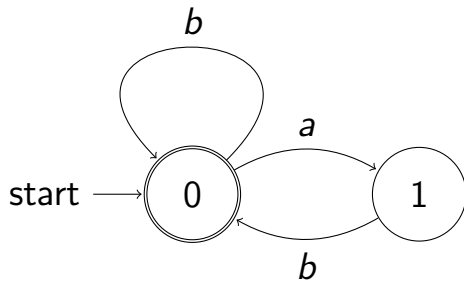
Example conversion of regex to ϵ NFA

$(ab+b)^*$

$(ab+b)^*$



Minimal NFA for $(ab+b)^*$



Describing Regular Languages

Regular Expression (regex)



Nondeterministic Finite Automata with ϵ moves
(ϵ NFA)



Nondeterministic Finite Automata (NFA)



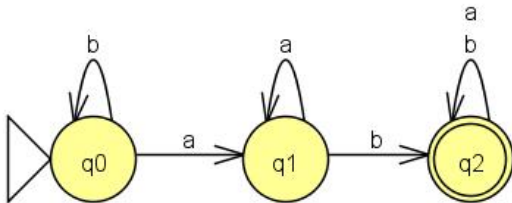
Deterministic Finite Automata (DFA)

Equivalence of all

To show that they are all essentially equivalent, we need to have DFA to regex conversion.

DFA \rightarrow RegEx Conversion

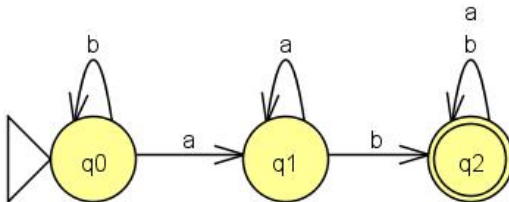
Using the state-elimination method



DFA \rightarrow RegEx Conversion

Using the state-elimination method

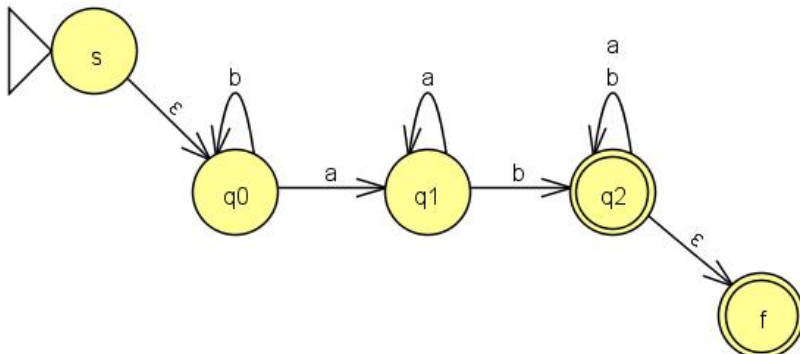
First step is to add new "super-state" (s) and (f), with ε -moves from (s) to the original start state, and from the original final states to (f)



DFA \rightarrow RegEx Conversion

Using the state-elimination method

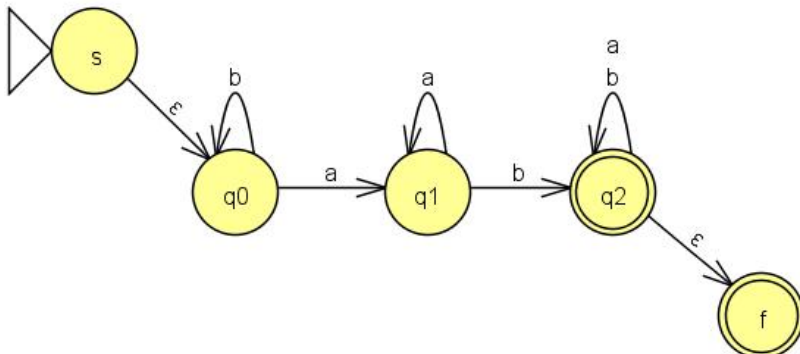
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DFA \rightarrow RegEx Conversion

Using the state-elimination method

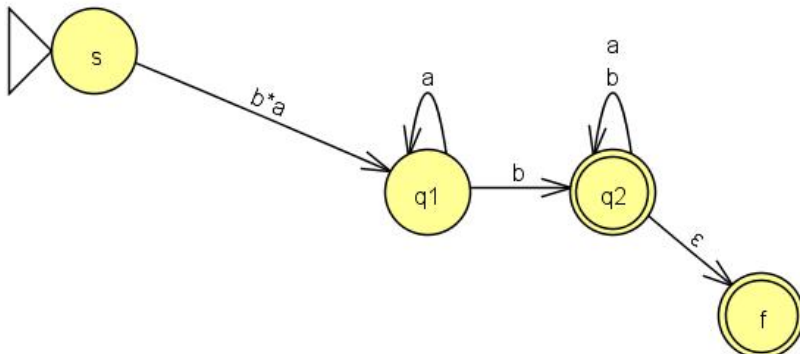
Then eliminate the original states one at a time, replacing transition labels with corresponding regular expressions



DFA \rightarrow RegEx Conversion

Using the state-elimination method

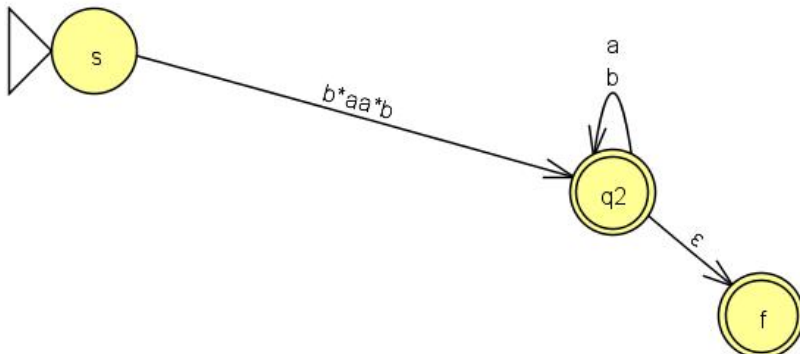
Then eliminate the original states one at a time, replacing transition labels with corresponding regular expressions



DFA \rightarrow RegEx Conversion

Using the state-elimination method

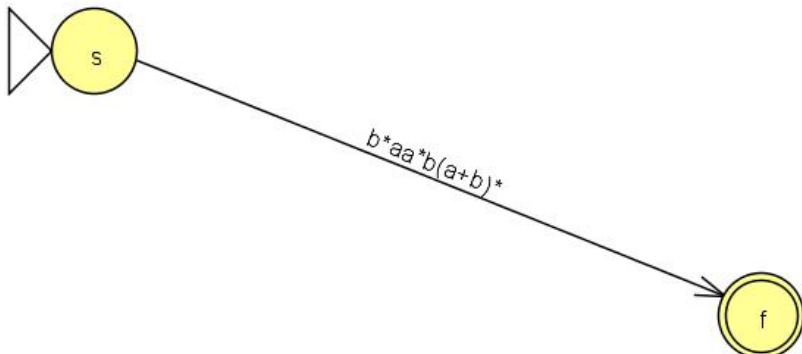
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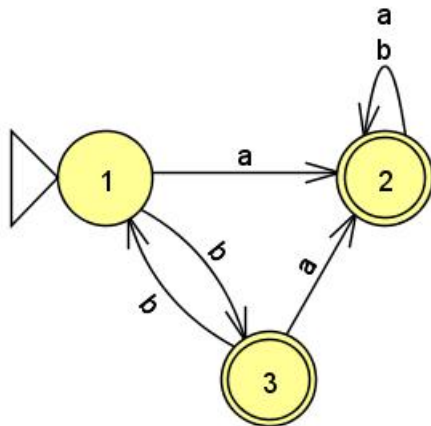
DFA \rightarrow RegEx Conversion

Using the state-elimination method

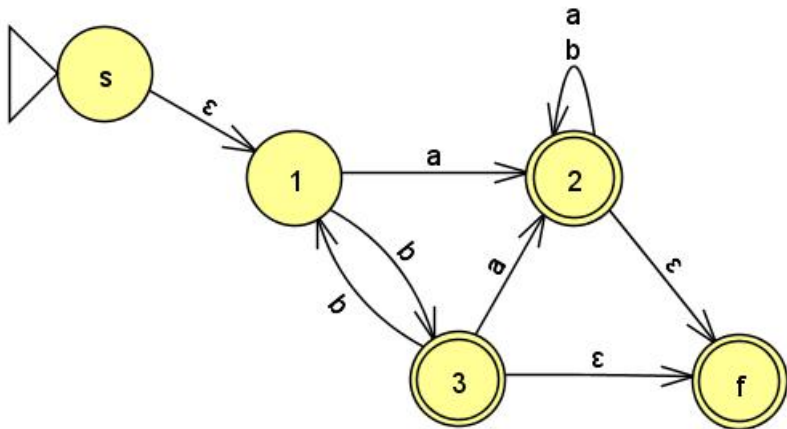
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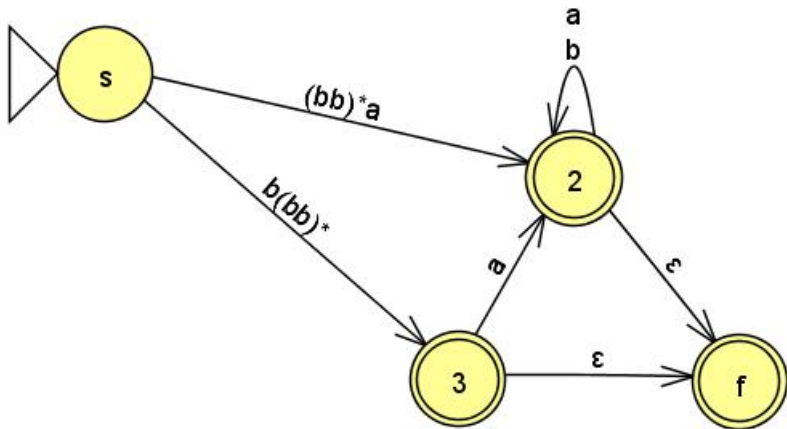
Another example



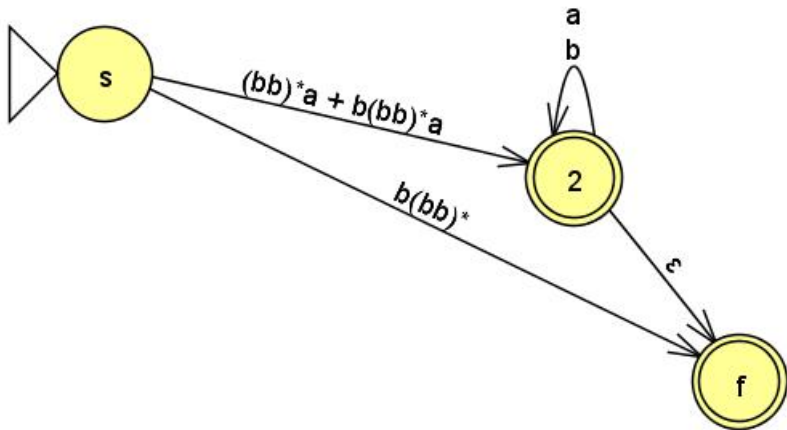
Another example



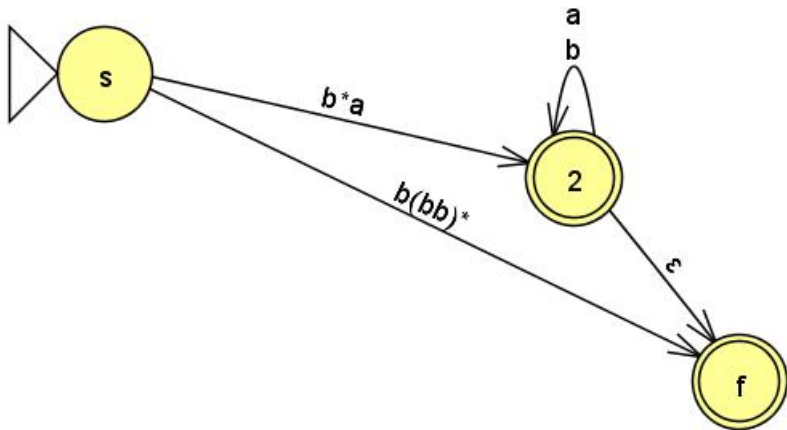
Another example



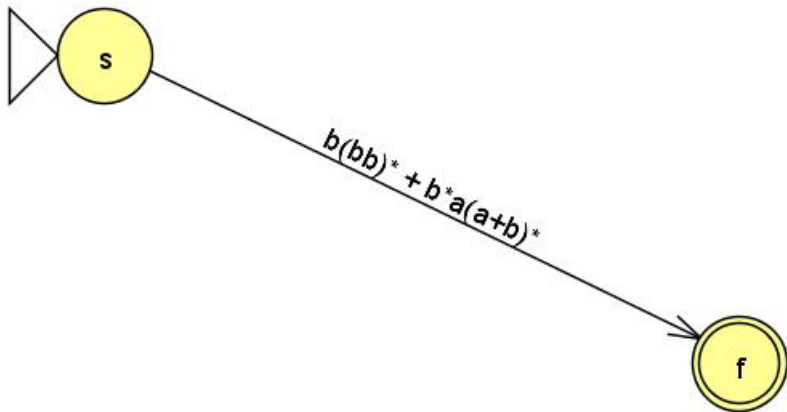
Another example



Another example

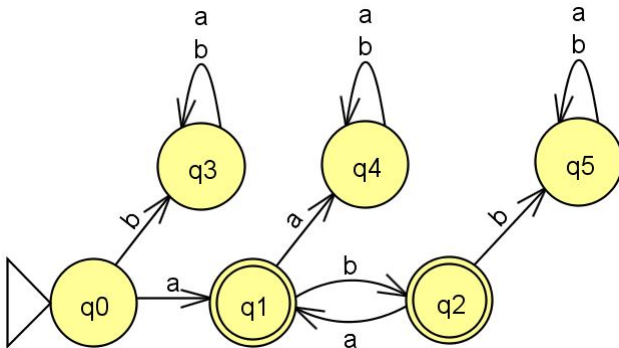


Another example



Exercise

- Does the sequence in which we eliminate states affect the resulting regular expression?
- Try finding a regular expression for the following DFA



References

- Previous slides on CMSC 141
- M. Sipser. Introduction to the Theory of Computation. Thomson, 2007.
- J.E. Hopcroft, R. Motwani and J.D. Ullman. Introduction to Automata Theory, Languages and Computation. 2nd ed, Addison-Wesley, 2001.
- E.A. Albacea. Automata, Formal Languages and Computations, UPLB Foundation, Inc. 2005
- JFLAP, www.jflap.org