

Quiz

Get 1/4

- 1-2. Enumerate the 2 influences on language design.**
- 3. What is the basis for imperative languages?**
- 4. Who is your lab instructor? What can you say about him/her?**

Chapter 3: Language Translation Issues

CMSC 124, 1st Semester, AY 2009-10



Chapter 3: Language Translation Issues

Topics Covered

- Stages in Translation
- Language Specification
 - ❖ Syntax
 - ❖ Semantics
- Syntax
 - ❖ Grammars and BNF
 - ❖ Syntax Diagrams
 - ❖ Finite State Automata
 - ❖ Pushdown Automata
- Semantics



Chapter 3: Language Translation Issues

Writing a Translator/Compiler: Span

- Programming Languages (**CMSC 124**)
- Machine Architecture (**CMSC 132**)
- Language Theory (**CMSC 141**)
- Algorithms (**CMSC 123**)
- Software Engineering (**CMSC 128**)
- Machine-Level Programming (**CMSC 131**)

Chapter 3: Language Translation Issues

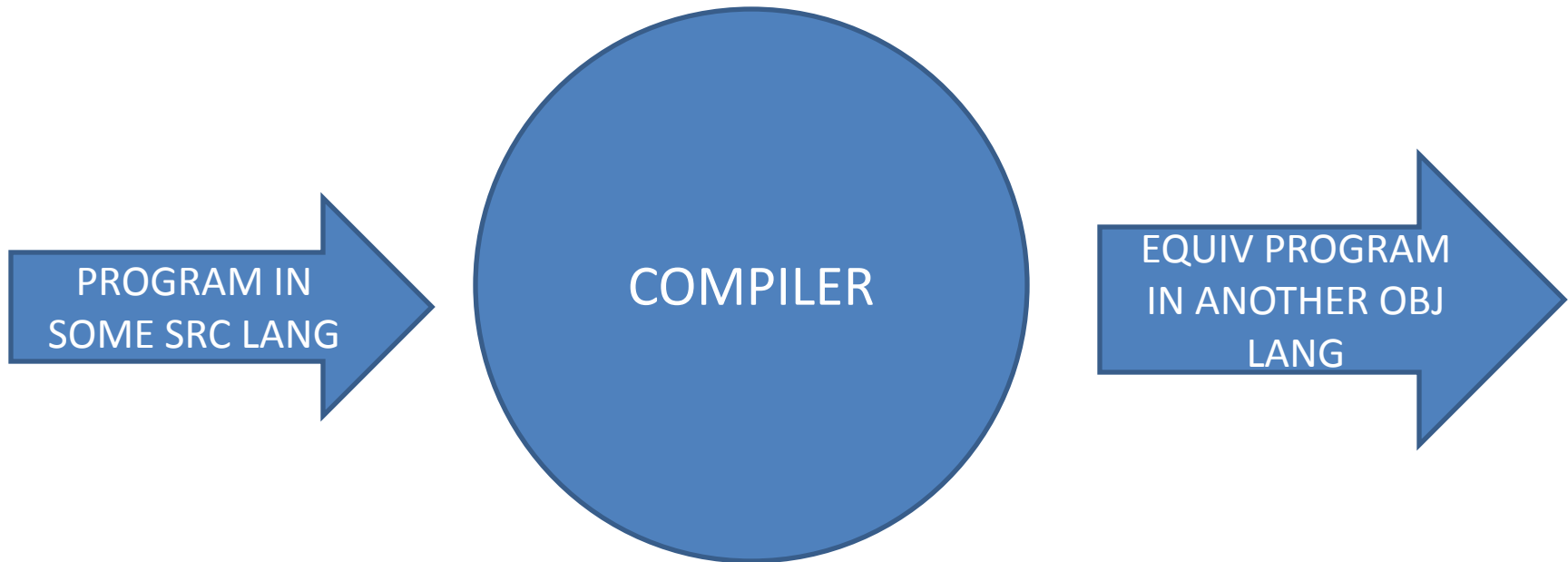
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Recall: Compilers



Chapter 3: Language Translation Issues

Analysis-Synthesis Model of Compilation

Two Parts of Compilation

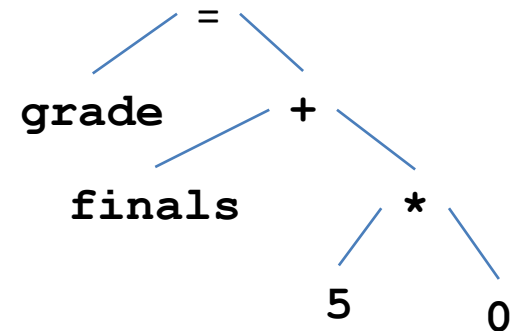
- **Analysis**

Breaks up the source program into constituent pieces and creates an intermediate represent'n of the source program.

- **Synthesis**

Constructs the desired target program from the intermediate representation

`grade = finals + 5 * 0`



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Analysis of the Source Program

- **Linear Analysis**

Stream of characters is read from left to right and group into tokens.

- **Tokens** – Sequences of chars. having a collective meaning.

- **Hierarchical Analysis**

Tokens are grouped hierarchically into nested collections with collective meaning.

- **Semantic Analysis**

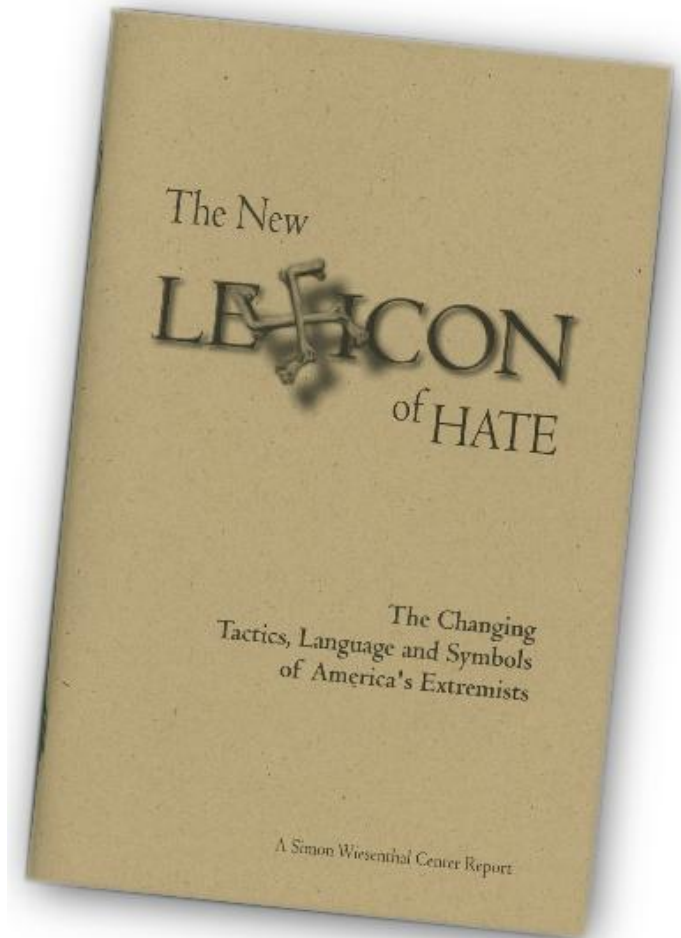
Performs certain checks to ensure that components of a program fit together meaningfully.



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Compilation Phase #1: Lexical Analysis

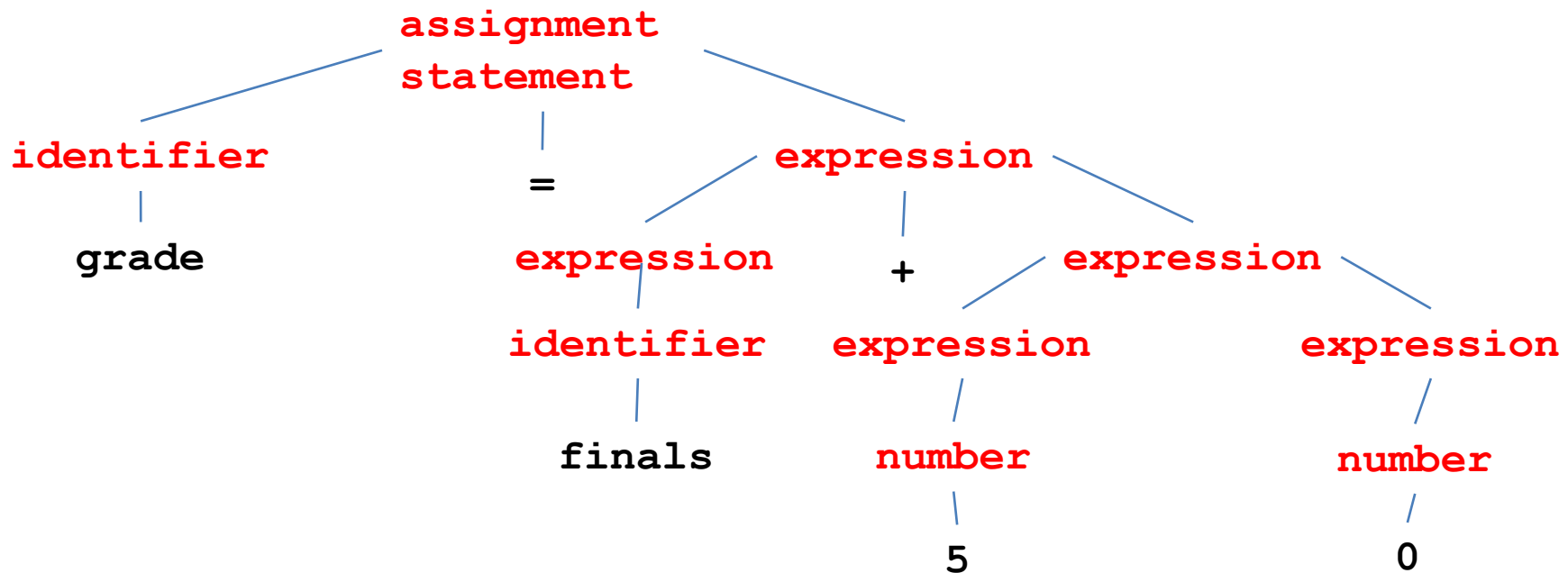
- AKA linear analysis or scanning.
- Consider the following:
grade = finals + 5 * 0
- Tokens:
 - identifier **grade**
 - assignment symbol **=**
 - identifier **finals**
 - plus sign **+**
 - number **5**
 - multiplication sign *****
 - number **0**
- Blanks are normally eliminated.



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Compilation Phase #2: Syntax Analysis

- AKA hierarchical analysis or parsing.
- Grouping tokens into grammatical phrases.
- The parse tree for `grade = finals + 5 * 0`



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Compilation Phase #2: Syntax Analysis

- **Expression Rules:**

- ✓ Any `identifier` is an expression.
- ✓ Any `number` is an expression.
- ✓ If `expression1` and `expression2` are expressions, then so are:
 - `expression1 + expression2`
 - `expression1 * expression2`
 - `(expression1)`



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Compilation Phase #2: Syntax Analysis

- **Statement Rules:**

- ✓ If **identifier** is an identifier **expression** and is an expression then:

- identifier = expression**

- is a statement.

- ✓ If **expression** is an expression and **statement** is a statement then:

- while (expression) do statement**

- if (expression) then statement**

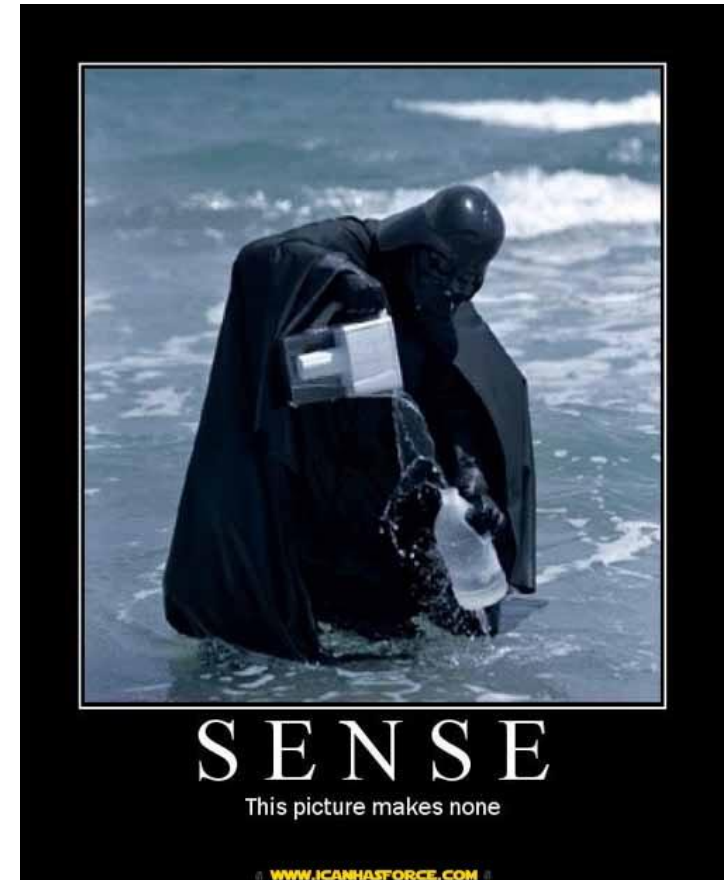
- are statements.

1. Statement Date
2. Cardholder's Name
3. Billing Cycle
4. Payment Due Date
5. Payment Due Amount
6. Total Amount Due
7. Minimum Payment Due
8. Interest Rate
9. Late Fee
10. Annual Fee
11. Cardholder's Address

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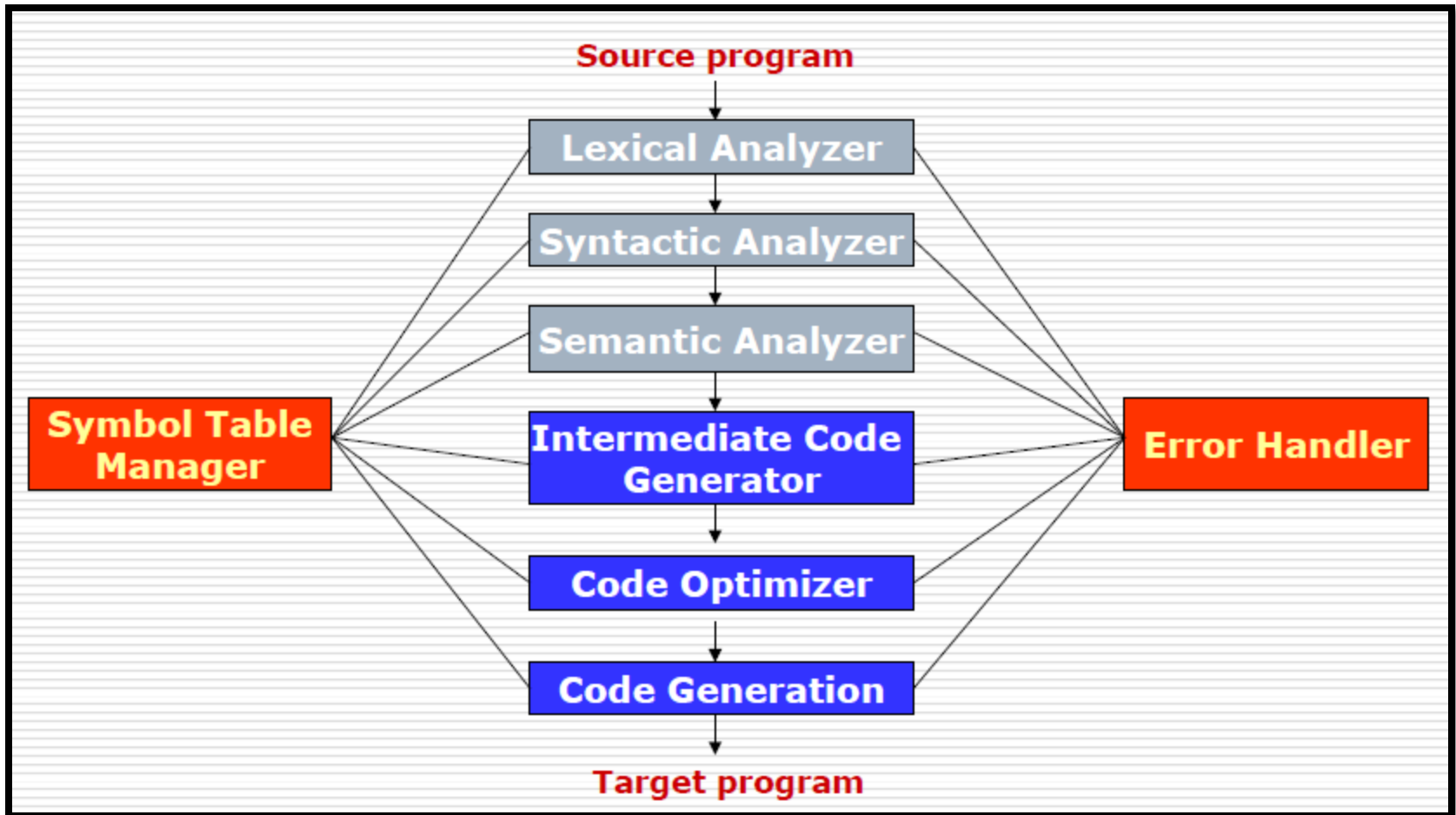
Compilation Phase #3: Semantic Analysis

- Checks for semantic errors.
- Performs type checking.
- Identifies the operators and operands of expressions and statements.
- Checks for valid operands.



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Compilation Phases



Chapter 3: Language Translation Issues

Two Other Activities in the Phases of Compilation

- Symbol-Table Maintenance
- Error Detection and Reporting



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Symbol-Table Maintenance

Symbol-Table

A data structure used by a language translator where each identifier in a program's source code is stored; Associated with information relating to its declaration or appearance in the source, such as its type, scope level and sometimes its location.

Global Symbol Table

NAME : string	TYPE : int	OFFSET : int	EXT_PTR : void *
a	int	0	NULL
b	long	4	NULL
c	int	12	NULL
f	function	N/A	
r	array	16	

Function Descriptor

Return_Type : int	Parameter_Type : vector<int>	SymbolTable : CSymbolTable *
int	int, int	

Function Symbol Table

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Size : int	Elementtype : int	EXT_PTR : void *
10	long	NULL

Array Descriptor

Size : int	Elementtype : int	EXT_PTR : void *
5	array	

Array Descriptor

Size : int	Elementtype : int	EXT_PTR : void *
5	int	NULL

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Symbol-Table Maintenance

Symbol-Table

A data structure that contains a record for each identifier.

- Type
- Scope
- For Procedures:
 - Parameter type
 - No. of parameters
 - Method of passing

Global Symbol Table

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Symbol-Table Maintenance

- It is essential to record identifiers.
- Identifiers are entered during the lexical analysis phase.

Global Symbol Table

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Error Detection and Reporting

- Each phase can encounter an error.
- To proceed, a phase must somehow deal with the error.
- A compiler that stops when it finds the first error is not as helpful.



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Error Detection and Reporting

- **Lexical Analysis Phase**

Characters do not form a valid token.

- **Syntax Analysis Phase**

Token streams violates the structure rules.

- **Semantic Analysis Phase**

Detect constructs that have the right syntax structure but no meaning to the operation involved.

- Eg. Add – an array and a procedure



Something to Ponder

The 3 phases of compilation discussed are indeed essential, important. Why?

Something to Ponder

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Review: Lexical Analysis (or Linear Analysis)

- A pattern matcher.
- It collects characters into logical groupings and assign internal groupings according to their structure.
- Character groupings are called **lexemes**.
- Internal codes are called **tokens**.

Eg:

```
result = oldsum - value / 100;
```

TOKEN	LEXEME
IDENT	result
ASSIGN_OP	=
IDENT	oldsum
SUBTRACT_OP	-
IDENT	value
DIVISION_OP	/
INT_LIT	100
SEMICOLON	;

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Review: Lexical Analysis (or Linear Analysis)

- **Goal:** Extract lexemes from a given input string and produce tokens.
- The process includes skipping comments and blanks.
- It also inserts lexemes for user-defined names into the symbol-table.
- The lexical analyzer is responsible for the initial construction of the symbol-table.

TOKEN	LEXEME
IDENT	result
ASSIGN_OP	=
IDENT	oldsum
SUBTRACT_OP	-
IDENT	value
DIVISION_OP	/
INT_LIT	100
SEMICOLON	;

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Review: Syntax Analysis (or Hierarchical Analysis)

- **Goals:**

1. The syntax analyzer must check the input program to determine whether it is syntactically correct.
 2. Produce a complete parse tree, or at least trace the structure of the complete parse tree for syntactically correct input.
- The parse tree is the basis for translation.

- **Eg.**

- Assume this is a statement in a certain PL: ($[x, y]$)

- These are the rules:

- $S ::= [A, S] \mid A$
- $A ::= C \mid (S)$
- $C ::= x \mid y \mid x$

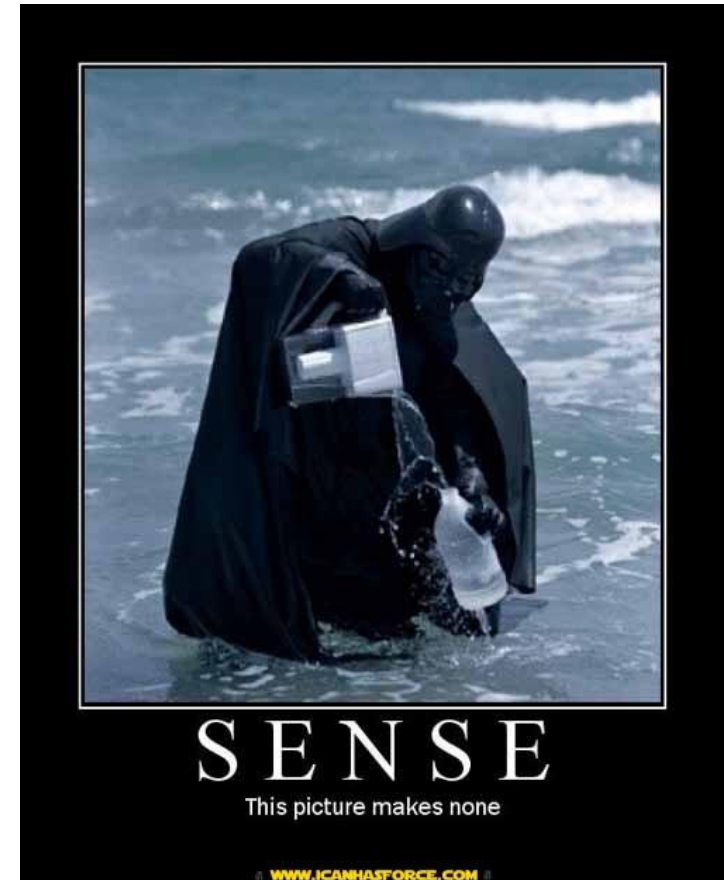
- **Generate** a parse tree?

- **Remember:** Grammars are used to describe syntax.

Chapter 3: Language Translation Issues

Review: Semantic Analysis

- Checks for semantic errors.
- Performs type checking.
- Identifies the operators and operands of expressions and statements.
- Checks for valid operands.



Chapter 3: Language Translation Issues

Compilation Phase #4: Intermediate Code Generation

Intermediate Code Generation

- Some compilers generate an explicit intermediate representation or a program for an abstract machine.
- Intermediate representation have many forms.
- **Eg: Three Address Code**
 - Similar to assembly language.
 - Every memo location can act as a register.
 - Consists of a sequence of instructions.
 - Each instruction has at most three operands.

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Compilation Phase #4: Intermediate Code Generation

Properties

- Each instruction has at most one operator plus assignment.
- Compiler must generate a temporary name to hold the value computed.
- Some instructions have fewer than three operands.

Eg:

```
position = initial + rate * 60
```

Remember: $id1 = id2 + id3 * 60$



```
temp1 = inttoreal(60)
```

```
temp2 = id3 * temp1
```

```
temp3 = id2 + temp2
```

```
id1 = temp3
```



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Compilation Phase #5: Code Optimization

- Improves the intermediate code to produce a faster running machine code.
- Compilers vary in the amount of time spent in this phase.

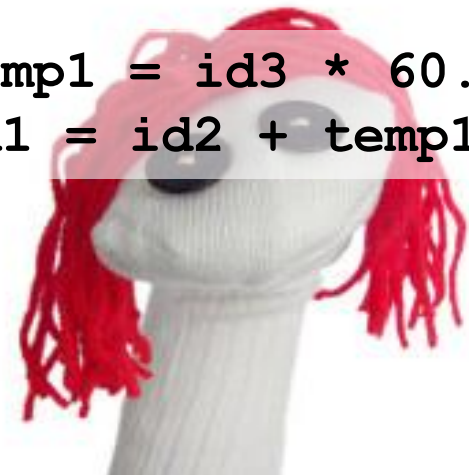
- **Eg:**

```
temp1 = inttoreal(60)
temp2 = id3 * temp1
temp3 = id2 + temp2
id1 = temp3
```



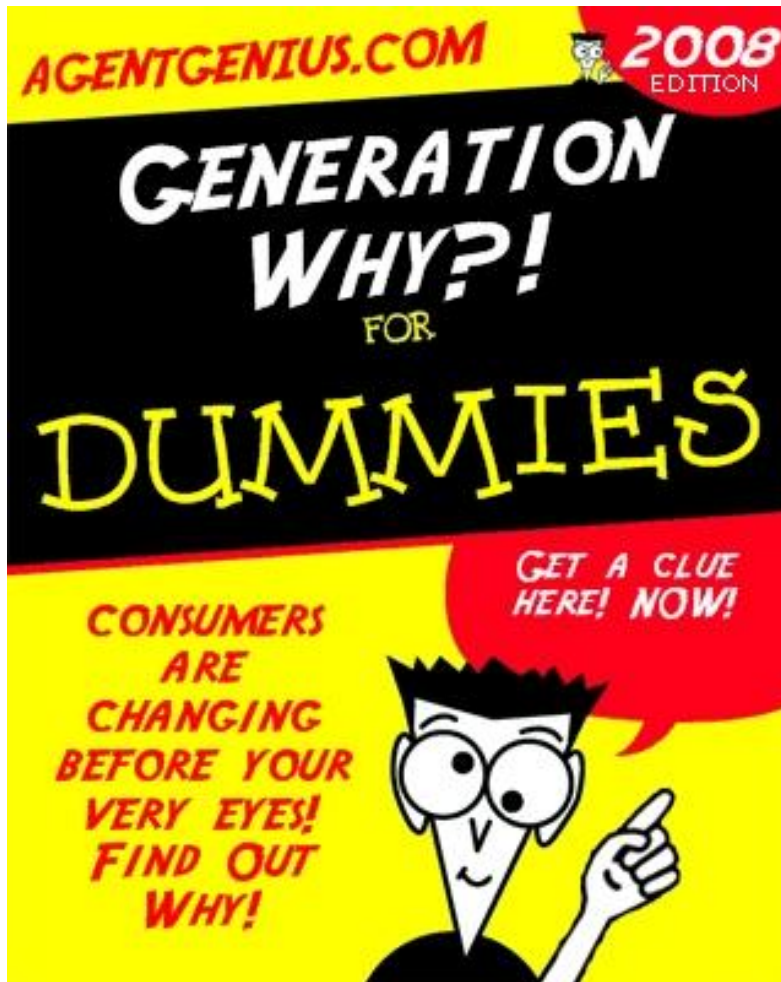
Optimize

```
temp1 = id3 * 60.0
id1 = id2 + temp1
```



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Compilation Phase #6: Code Generation



- Final phase.
- Generation of the target code from relocatable machine or assembly machine or assembly code.
 - Memory locations of variables are selected.
 - Each intermediate code are translated into an equiv. sequence of machine instructions.
- Assignment of variables to registers is **crucial!**

Chapter 3: Language Translation Issues

Translation of a Statement

Symbol Table

1	position	...
2	initial	...
3	rate	...
4		

Position = initial + rate * 60

Lexical Analyzer

id1 = id2 + id3 * 60

Syntactic Analyzer

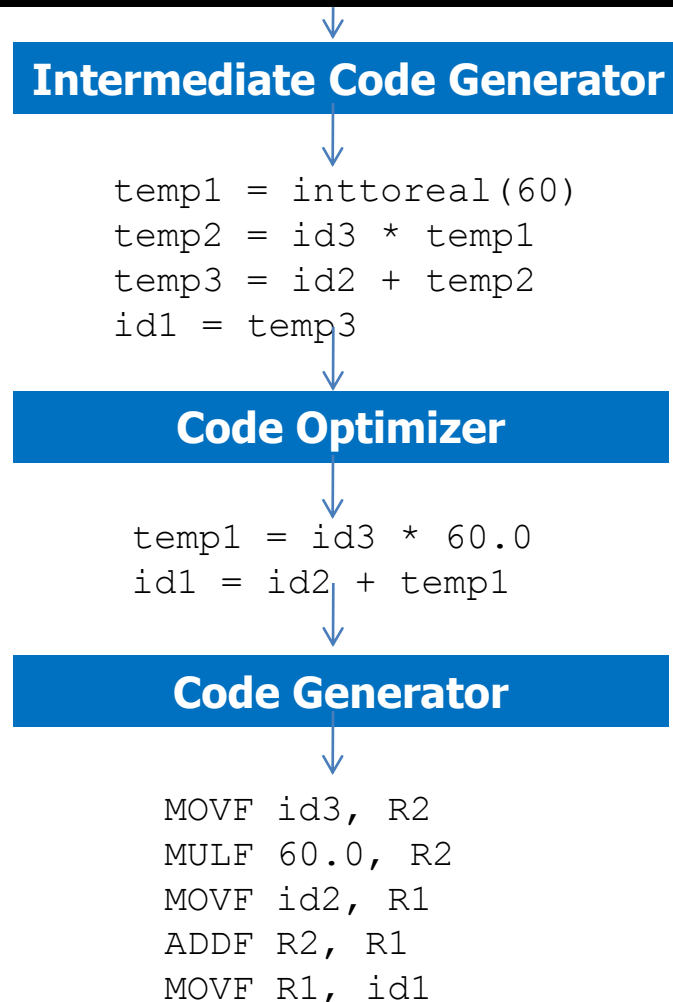
id1 = + id2 * 60

Semantic Analyzer

id1 = + id2 * inttoreal 60

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Translation of a Statement



Chapter 3: Language Translation Issues

Language Specification

Syntax

- “What its program **looks like**.”
- Refers to ways symbols may be combined to create well-formed sentences in the language.
- Formal methods include Context-Free Grammar or Backus Naur Form (BNF)

Semantics

- “What its program **means**.”
- Much more difficult to describe.
- Informal methods and suggestive examples are normally used.

