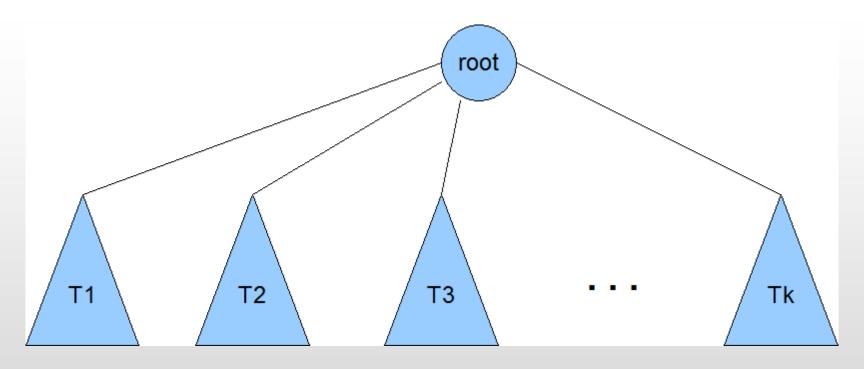
4. Trees

4.1 Basic Concepts and Terminology

- A tree is a collection of nodes.
- The collection can be empty, otherwise, a tree consists of a distinguished node r, called the root, and zero or more (sub)trees T₁, T₂, ..., T_k, each of whose roots are connected by a directed edge to r.
- no. of nodes = N, no. of edges = ?

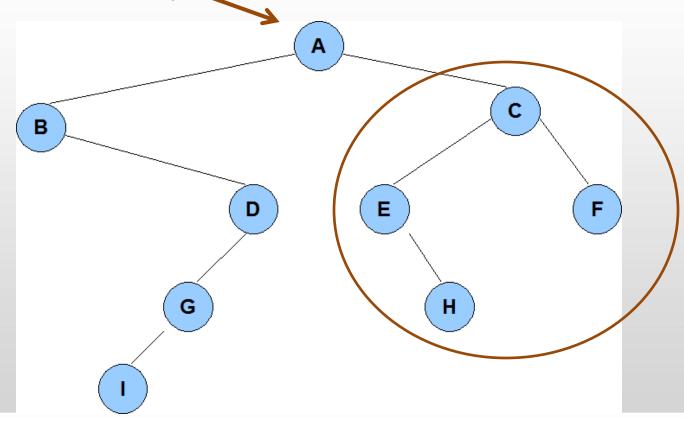


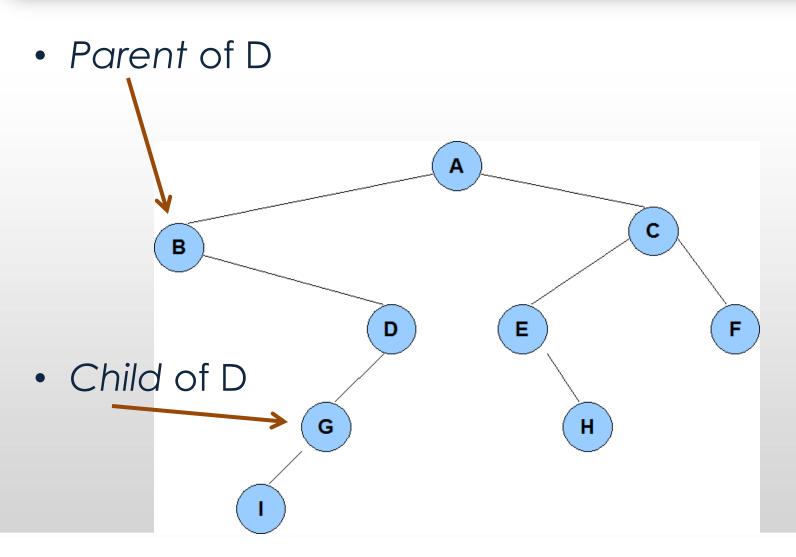
Generic Tree





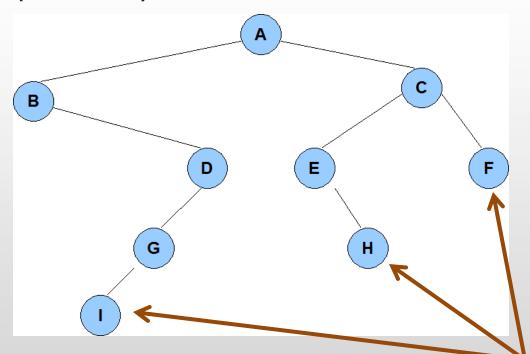
• The root of each subtree is said to be a child of r, and r is the parent of each subtree root.





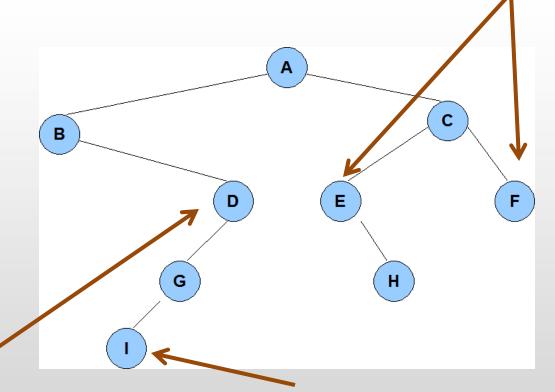


 Each node may have an arbitrary number of children, possibly zero.



Nodes with no children are known as leaves.

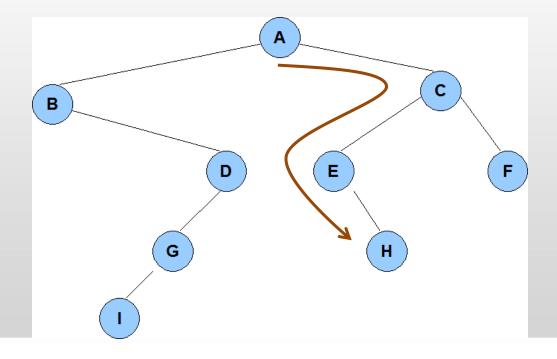
Nodes with the same parent are siblings.



Grandparent and grandchild



• A path from node n_1 to n_k is defined as a sequence of nodes n_1 , n_2 , ..., n_k such that n_i is the parent of n_{i+1} for $1 \le i \le k$.

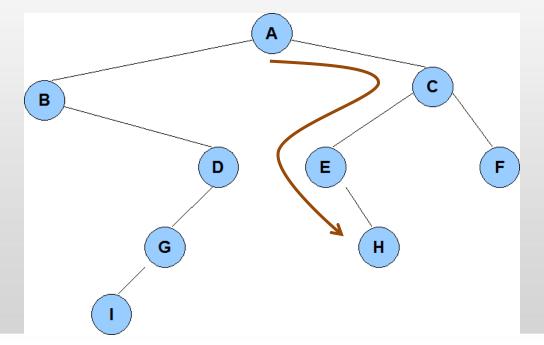




 The length of a path is the number of edges on the path.

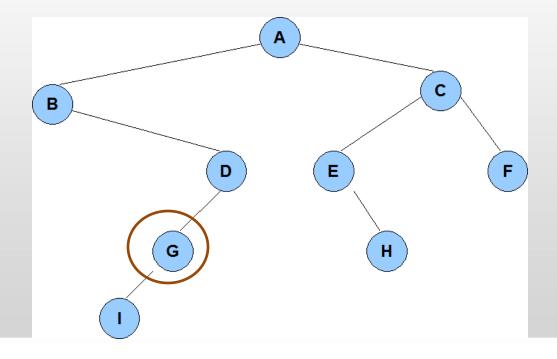
There is a path of length zero from every node

to itself.





- For any node n_i, the depth of n_i is the length of the unique path from the root to n_i.
- Thus the root is at depth 0.

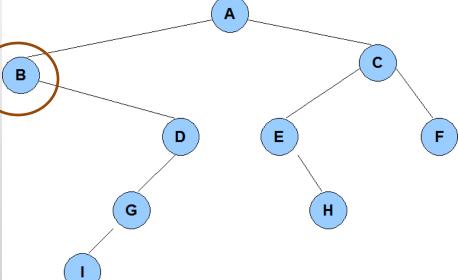




- The height of n_i is the longest path from n_i to a leaf.
- Thus all leaves are at height 0.

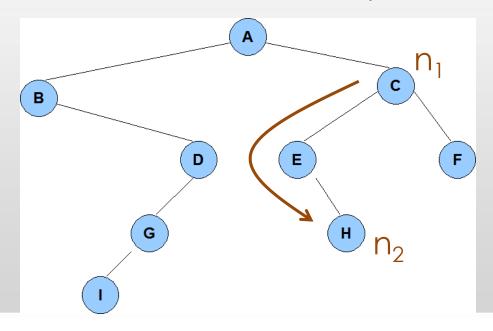
The height of a tree is equal to the height of the

root.



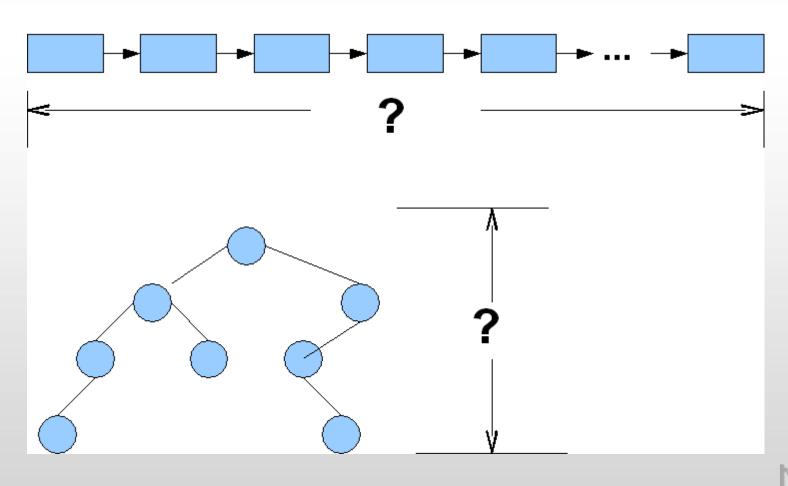


- If there is a path from n_1 to n_2 , then n_1 is an ancestor of n_2 and n_2 is a descendant of n_1 .
- If $n_1 \neq n_2$, then n_1 is a proper ancestor of n_2 and n_2 is a proper descendant of n_1 .





Tree vs. Linear List



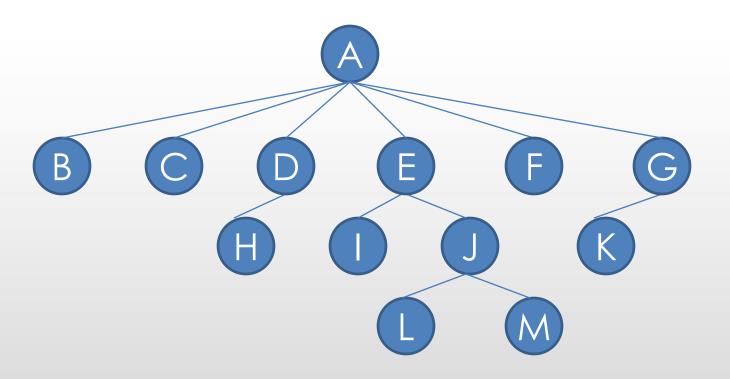


Implementation of Trees

- Linked representation
 - each node, besides its data has a pointer to each child of the node
- Left-child, Right-sibling representation
 - children of each node is kept in a linked list of tree nodes

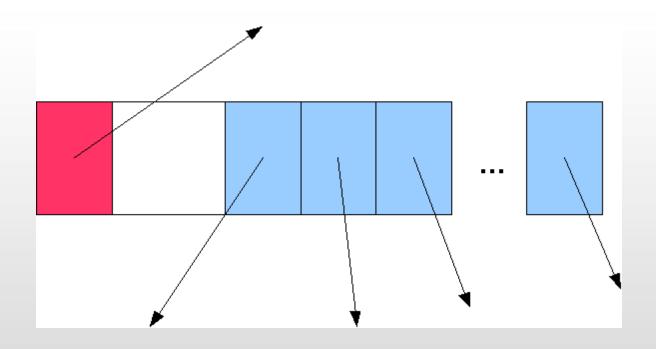


Linked Representation



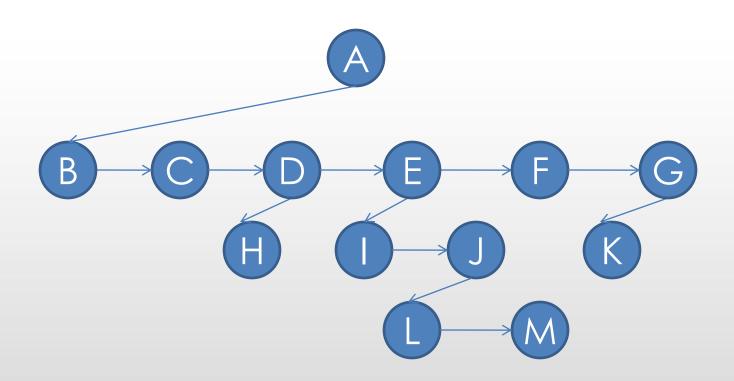


Linked Representation



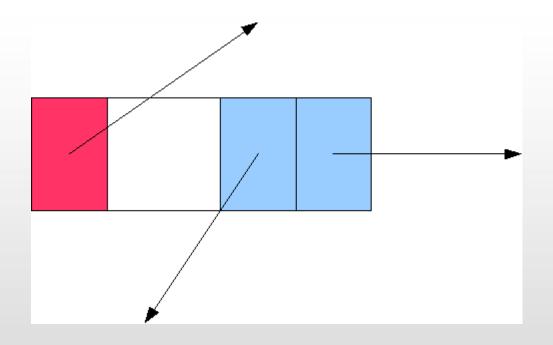


Left-child Right-sibling Representation





Left-child Right-sibling Representation



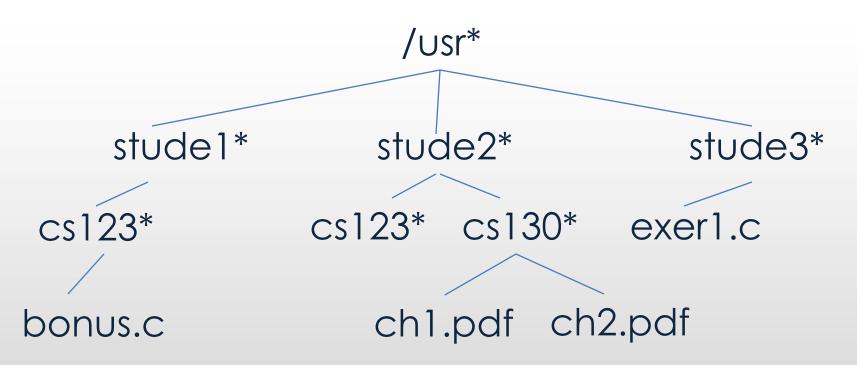


Applications

- Directory listing
- Getting the directory/file sizes
- Compiler design (e.g. Expression trees)



Directory Listing



/usr/stude1/cs123/bonus.c

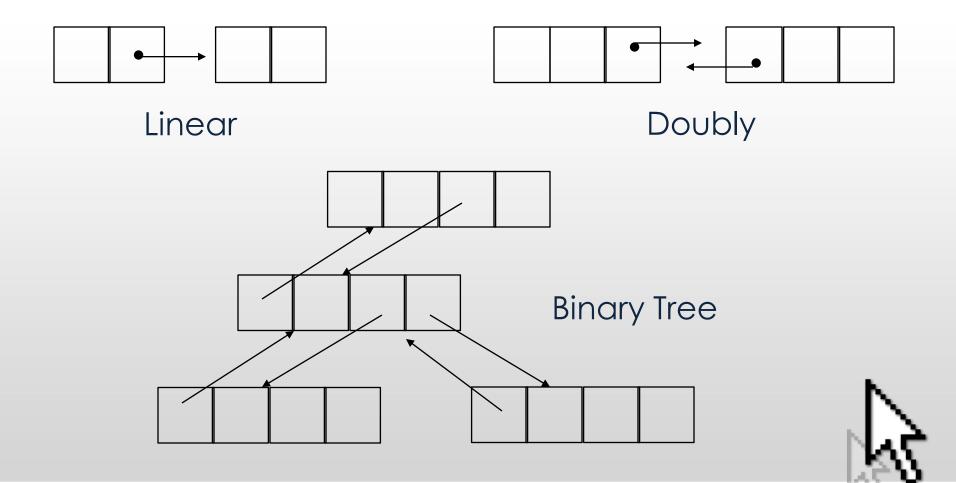


4. Trees

4.2 Binary trees and their implementations

- Linear linked list node has at most one pointer to another node
- Doubly linked list node has at most two pointers to two different nodes
- Binary tree node has at most three pointers to three different nodes with each node having a back pointer

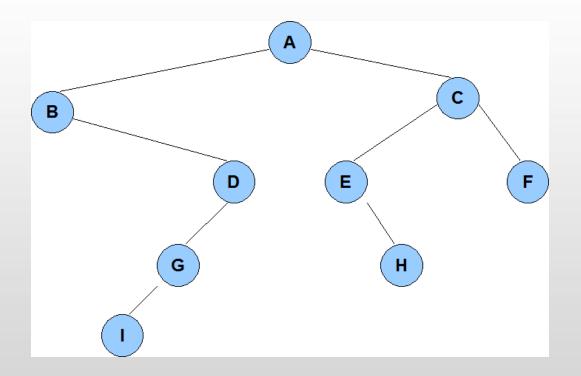




- A structure that is either empty
- Or one that contains a single node called the root
- Or a root with at most two descendants, each descendant being a root of another binary tree

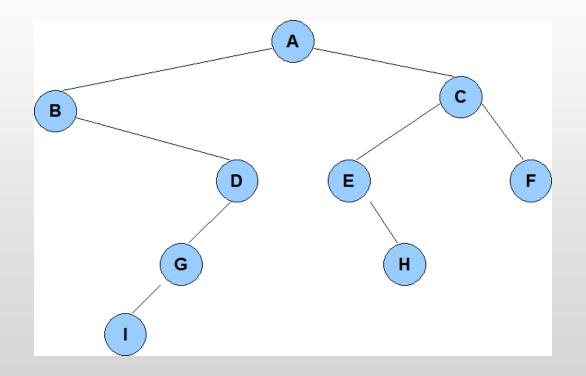


Internal nodes – nodes with at least one child



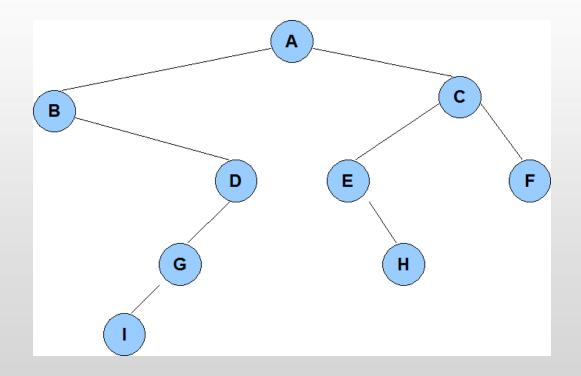


 Level of a node – 1 plus the level of its parent (root has level 0)



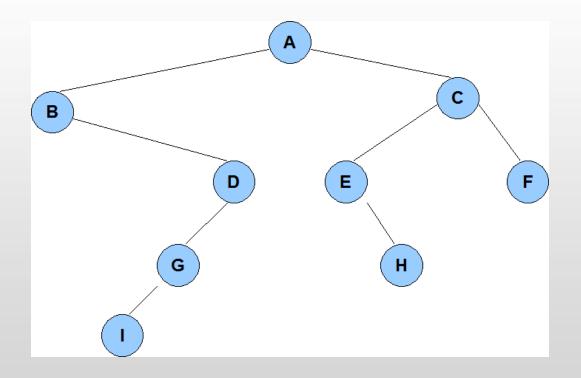


 Level i is full if there are exactly 2ⁱ nodes at this level



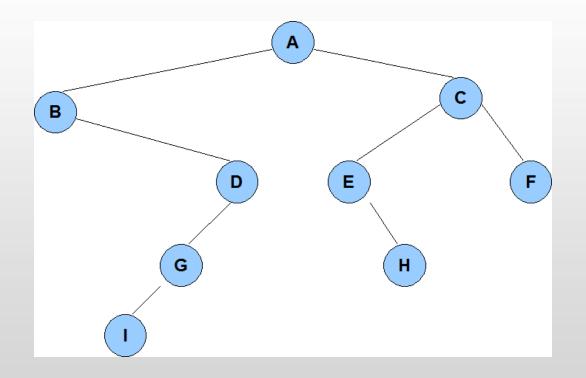


 A binary tree of height k is full if level k in this binary tree is full



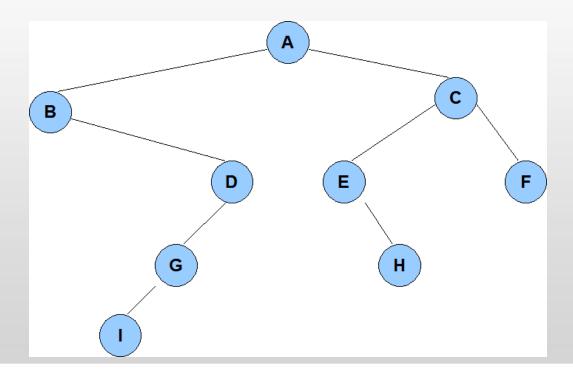


 A binary tree of height k is balanced if level k-1 in this binary tree is full



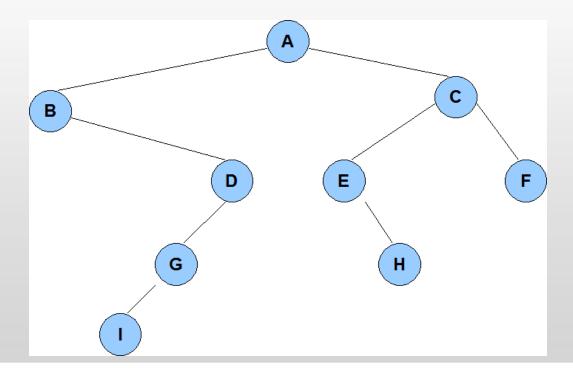


 The binary tree rooted at the left child of node v is the left subtree of v, and that rooted at the right child is the right subtree





 A binary tree is height balanced if, at every node in the tree, its left and right subtrees differ by no more than 1 in height





- Each node contains at least the value field, left child pointer, right child pointer, and a pointer to its parent
- If a parent or a child is not present, then the value of the pointer is NULL



```
typedef struct node{
  int value;
  struct node *left;
  struct node *right;
  struct node *parent;
}tree;
```



Binary Tree Traversal

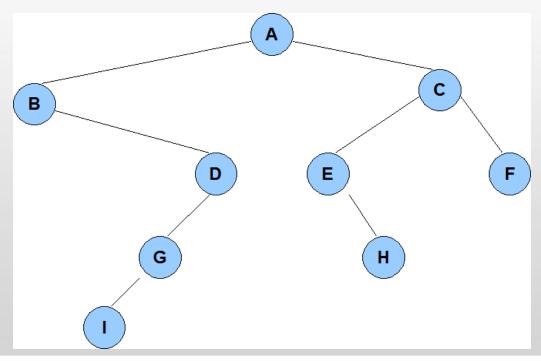
- Method of visiting each node of the tree at least once
- 3 ways:
 - Inorder
 - Preorder
 - Postorder



Binary Tree Traversal

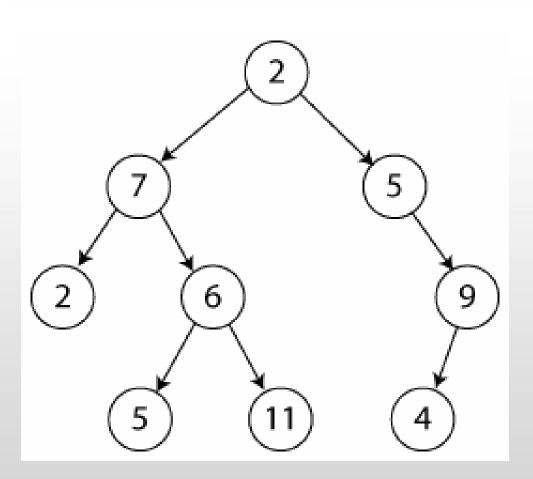
- Inorder
 - Left subtree is first visited, then the root node, and finally the right subtree
- Preorder
 - Root node is visited first, then left subtree, and finally the right subtree
- Postorder
 - Left subtree is visited first, then right subtree, and finally the root

- Inorder: BIGDAEHCF
- PreOrder: A B D G I C E H F
- PostOrder: I G D B H E F C A



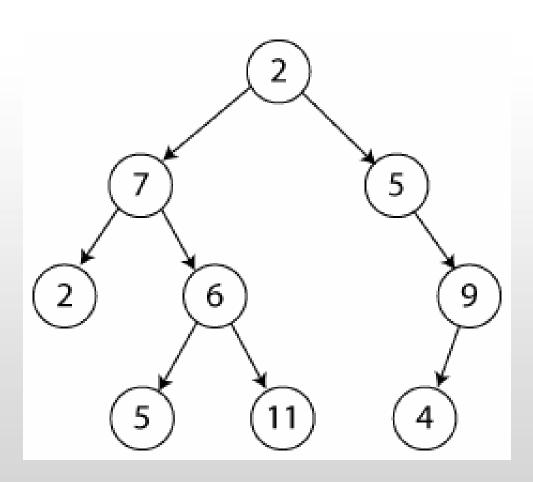


Recall



- root = 2
- parent of 6 = 7
- child of 9 = 4
- leaves = 2, 5, 11, 4
- sibling of 7 = 5
- path from 7 to 11 = 7, 6, 11
- length of path = 2

Recall



- depth of 4 = 3
- height of 7 = 2
- height of tree = 3
- level of 11 = 3
- Is level 2 full? No
- Is the tree full? No
- Is it balanced? No
- Is it height balanced? No

- A binary tree where the leaves hold the operands of the expression and the internal nodes hold the operators of the expression.
- The algorithm for creating an expression tree from the postfix form of the expression will require a stack of pointers to binary trees.



Algorithm

- 1. Read a symbol from the postfix form of the expression.
- 2. If the symbol is an operand, create a binary tree with one node from the operand. Then push into the stack a pointer to this one-node tree.



Algorithm

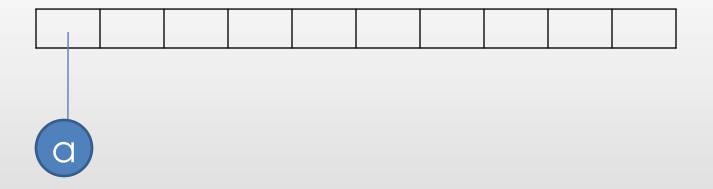
- 3. If the symbol is an operator, create a one-node tree out of the operator.
 - Pop the stack and make the right child of the operator node point to the binary tree just popped.
 - Pop the stack again and make the left child of the operator point to the binary tree just popped.
 - Push into the stack a pointer to the operator node.

- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-



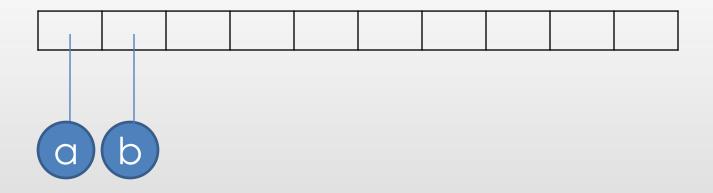


- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-



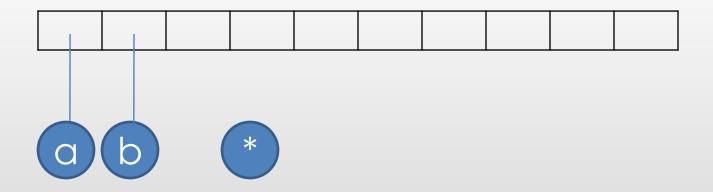


- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-



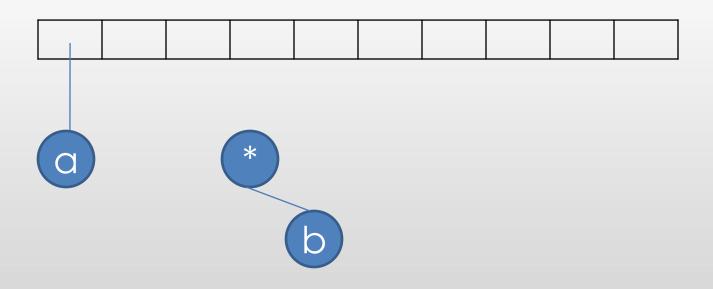


- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-





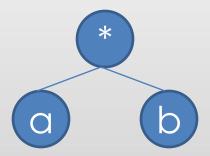
- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-





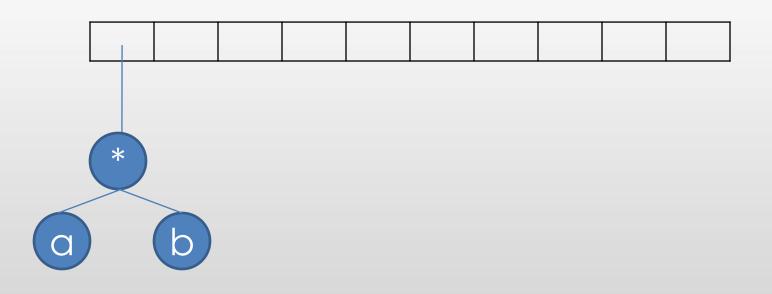
- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-





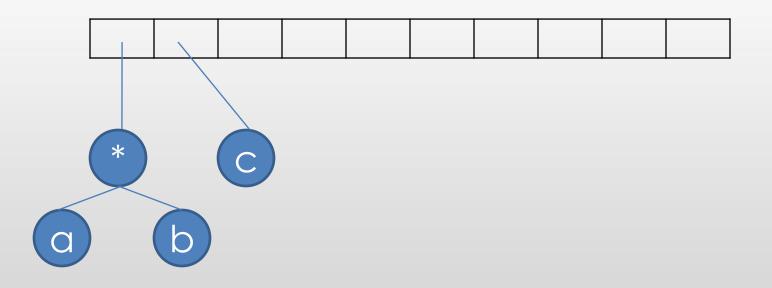


- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-



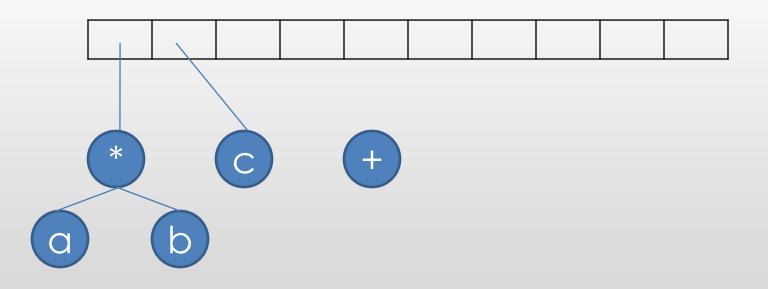


- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-



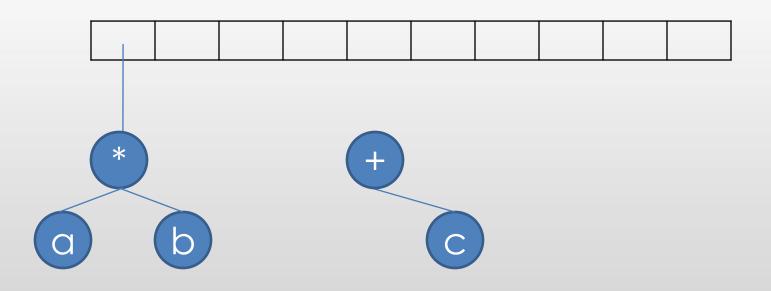


- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-





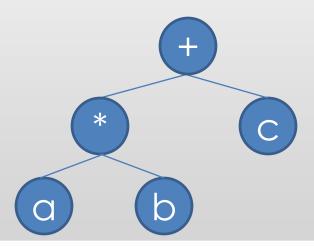
- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-





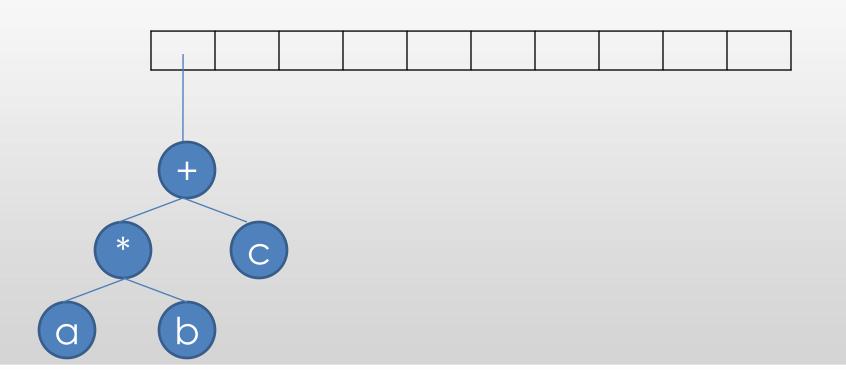
- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-





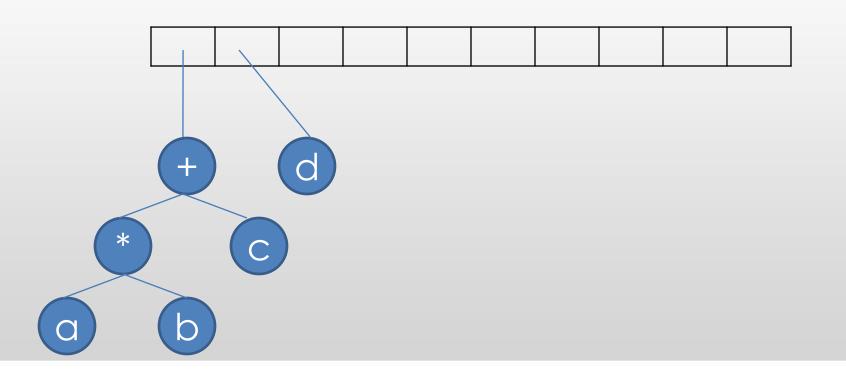


- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-



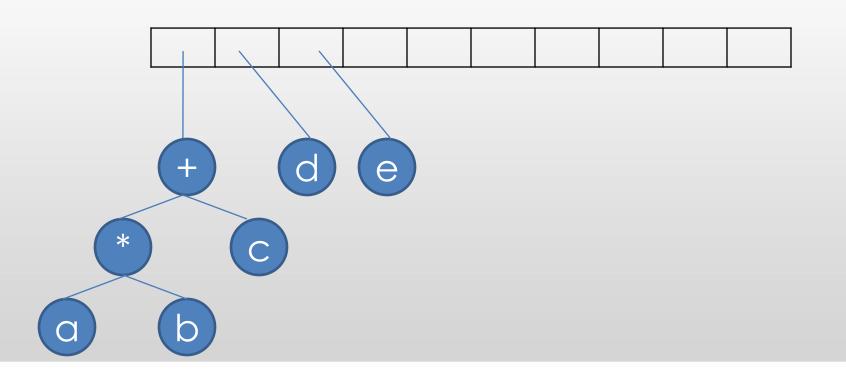


- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-



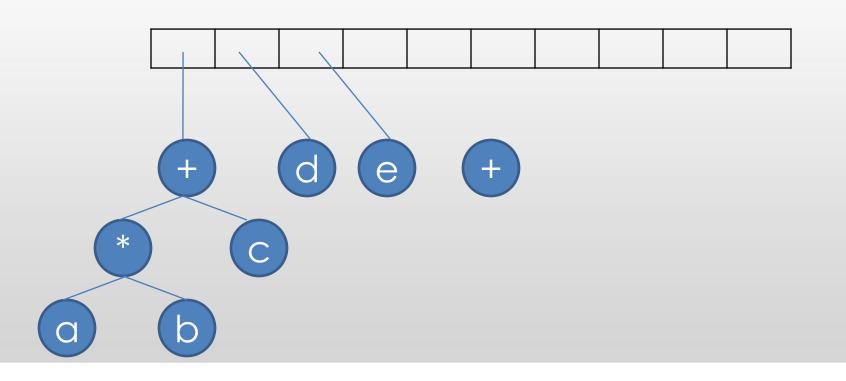


- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-



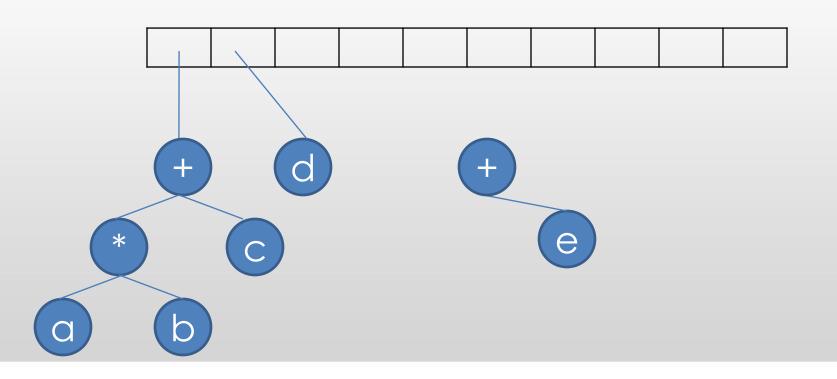


- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-



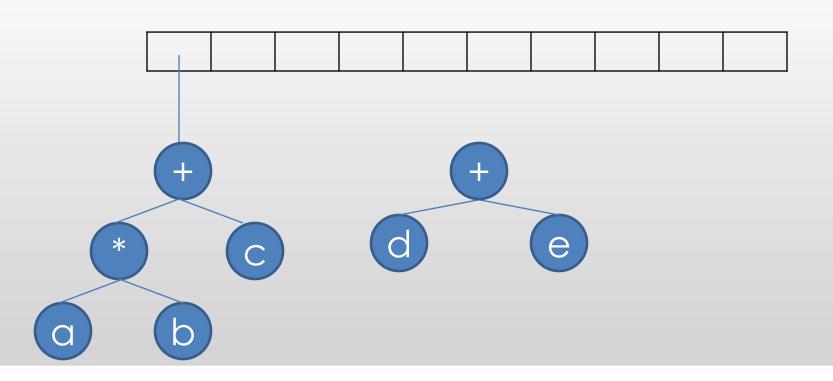


- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-



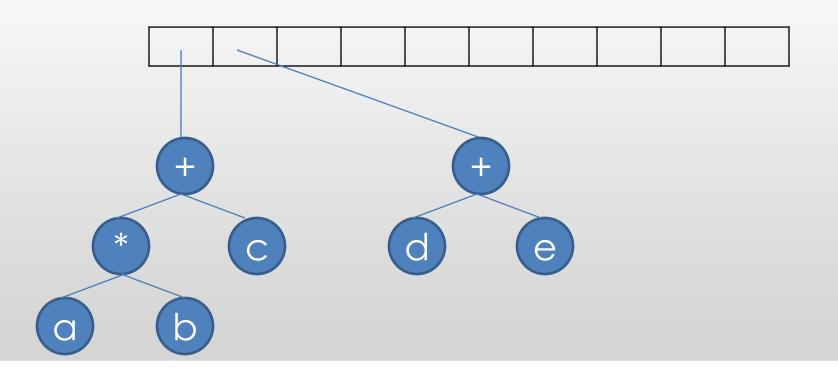


- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-



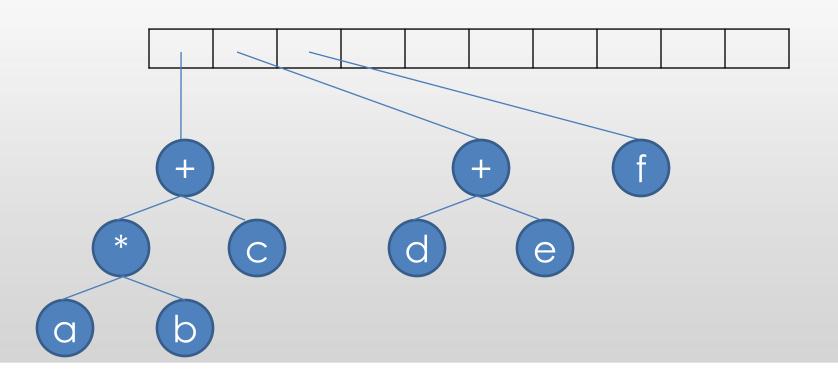


- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-



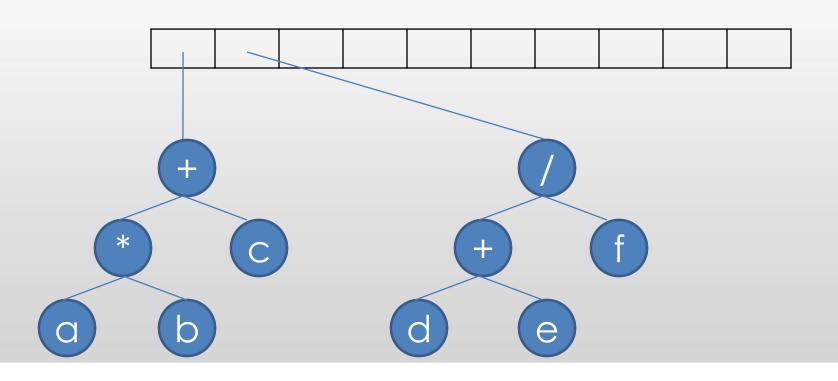


- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-



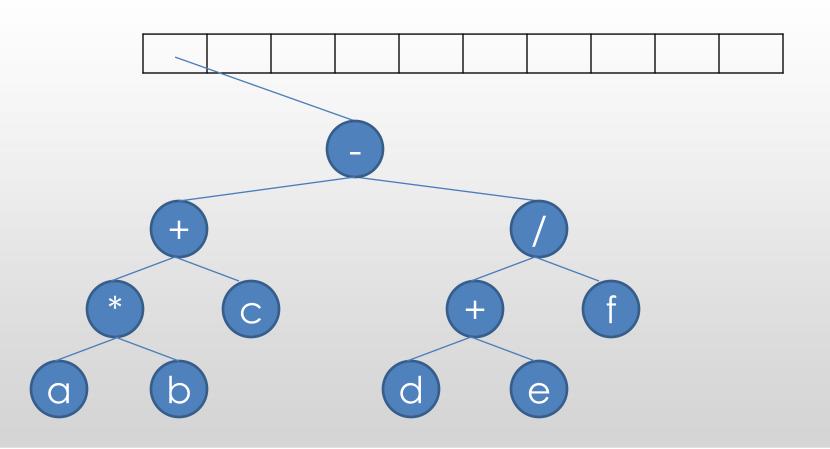


- Infix Expression: (a*b+c)-((d+e)/f)
- Postfix Form: ab*c+de+f/-





Postfix Form: ab*c+de+f/-





4. Trees

4.3 Binary Search Trees

Binary Search Tree (BST)

Definition: In a BST, the value of every node is greater than the values of the nodes in its left subtree and is less than the values of the nodes in its right subtree.

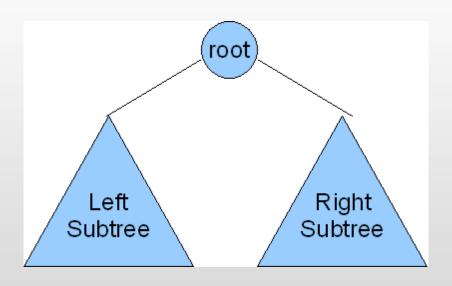
Operations:

- search, insert, delete
- minimum, maximum
- successor, predecessor (the value immediately greater and lesser than a node respectively)



BST

• Search order/tree property – i.e. Keys are organized in a special way, such that ...

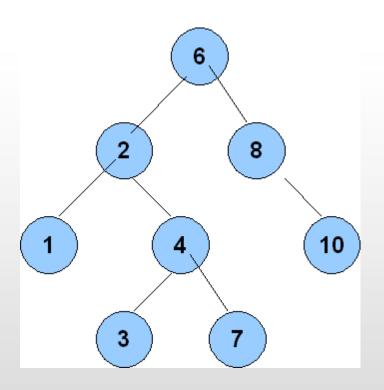


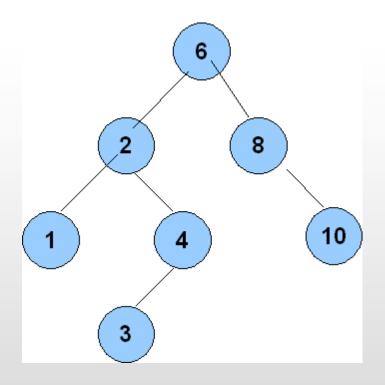
Two versions:

- a) keys are distinct
- b) duplicates are allowed



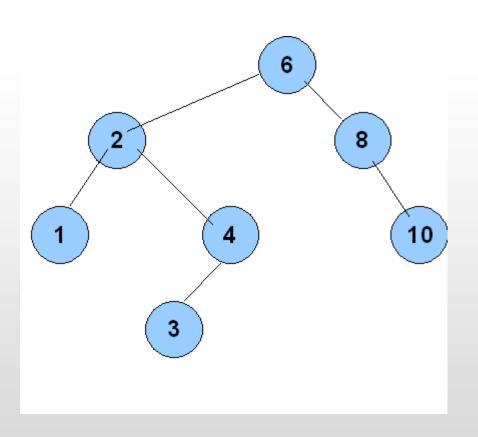
BST vs non-BST





Which is (not) a BST?





Two implementations:

- Non-recursive
- Recursive



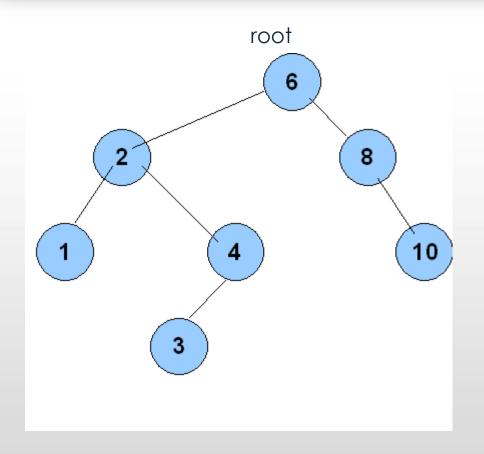
```
typedef struct node{
  int value;
  struct node *left;
  struct node *right;
  struct node *parent;
}BST;
```

```
BST *search(BST *root, int x) {
  BST *temp=root;

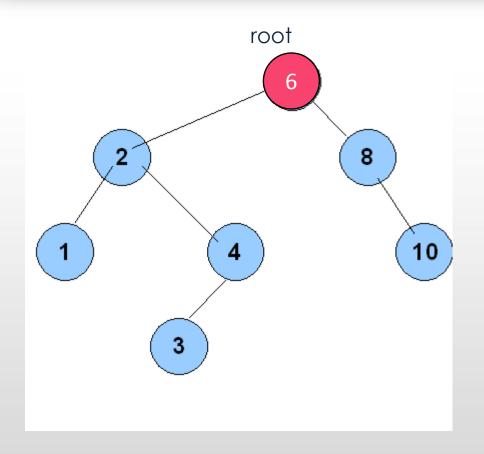
while((temp!=NULL)&&(temp->value!=x)) {
  if(x < temp->value)
    temp = temp->left;
  else
    temp = temp->right;
}

return temp;
}
```

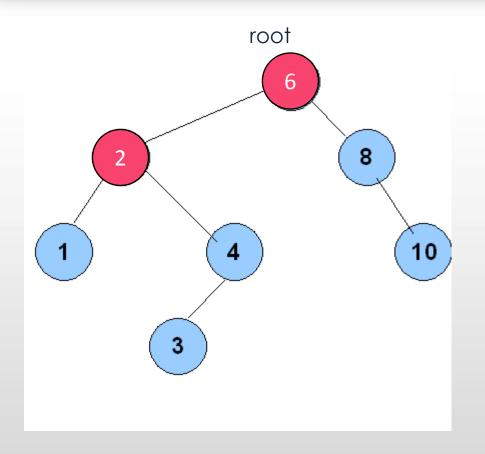




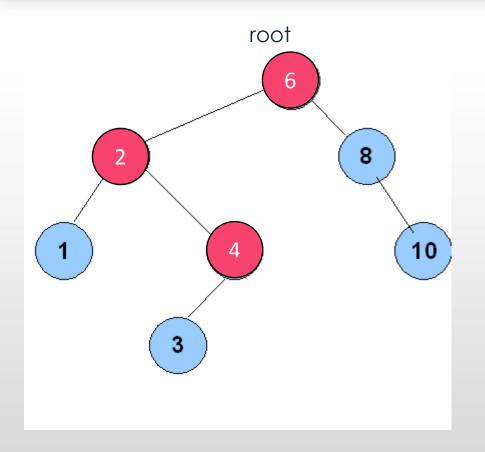






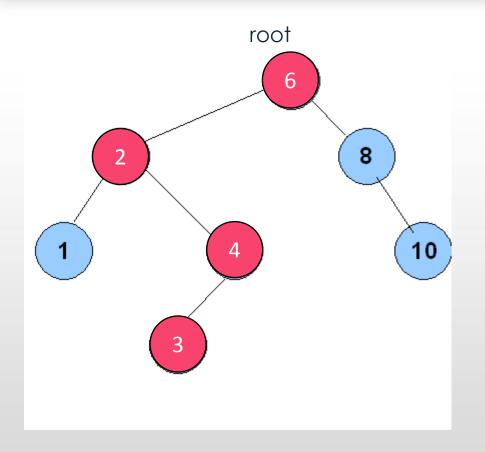








Search for a key



ptr = search(root, 3);



Search for a key

```
typedef struct node{
  int value;
  struct node *left;
  struct node *right;
  struct node *parent;
}BST;
```

```
BST *search(BST *root, int x){

if(root==NULL)
   return NULL;

if(x < root->value)
   return(search(root->left, x));

if(x > root->value)
   return(search(root->right, x));

else
   return root;
}
```

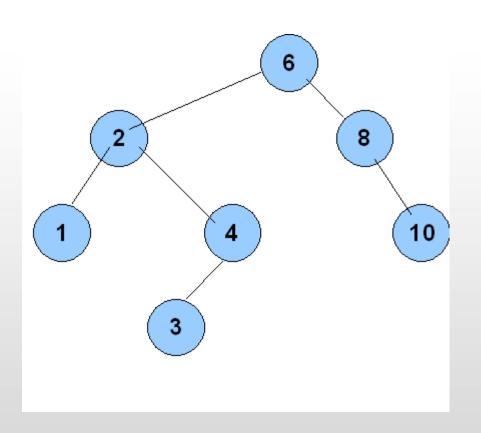


```
typedef struct node{
  int value;
  struct node *left;
  struct node *right;
  struct node *parent;
}BST;
```

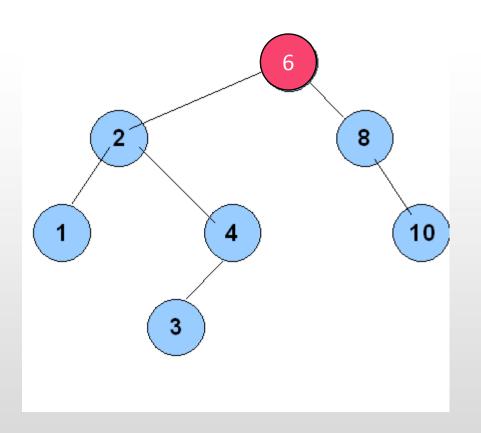
```
BST *find_min(BST *root) {
   BST *temp=root;

if(temp!=NULL) {
   while(temp->left!=NULL)
       temp = temp->left;
   return temp;
   }
   else
   return NULL;
}
```

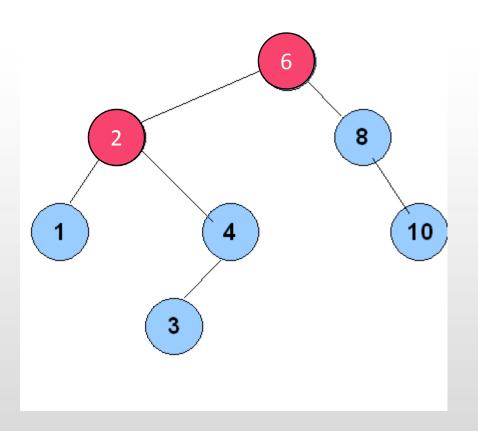




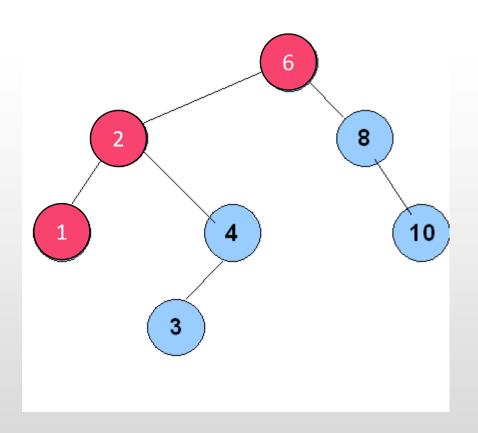














```
typedef struct node{
  int value;
  struct node *left;
  struct node *right;
  struct node *parent;
}BST;
```

```
BST *find_min(BST *root){

if(root==NULL)
   return NULL;
else if(root->left==NULL)
   return(root);
else
   return(find_min(root->left));
}
```



```
typedef struct node{
  int value;
  struct node *left;
  struct node *right;
  struct node *parent;
}BST;

int main(){
  BST *root=NULL;
}
```

```
void insert(BST *root, int x) {
  BST *temp;
  temp=(BST *) malloc(sizeof(BST));
  if(temp==NULL){
    printf("Insufficient Memory");
    exit(1);
  temp->value=x;
  temp->left=NULL;
  temp->right=NULL;
  temp->parent=NULL;
  insert2(root, temp);
```

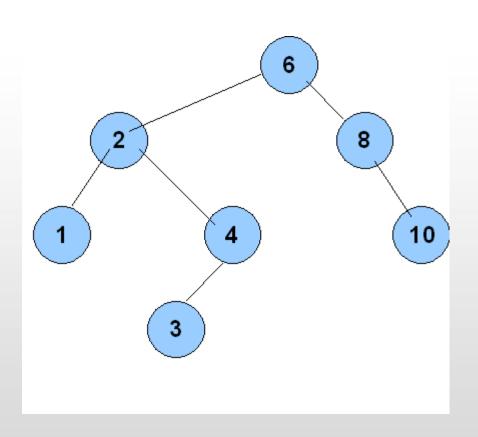
```
typedef struct node{
  int value;
  struct node *left;
  struct node *right;
  struct node *parent;
}BST;

int main(){
  BST *root=NULL;
}
```

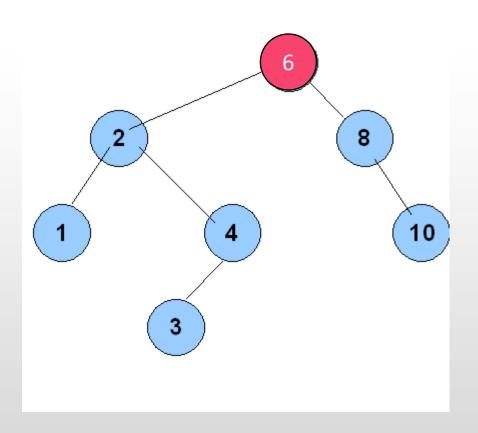
```
void insert2(BST *root, BST *temp){

if(root==NULL)
  root=temp;
else{
  temp->parent=root;
  if ((root)->value > temp->value)
    insert2(root->left, temp);
  else
    insert2(root->right, temp);
}
```

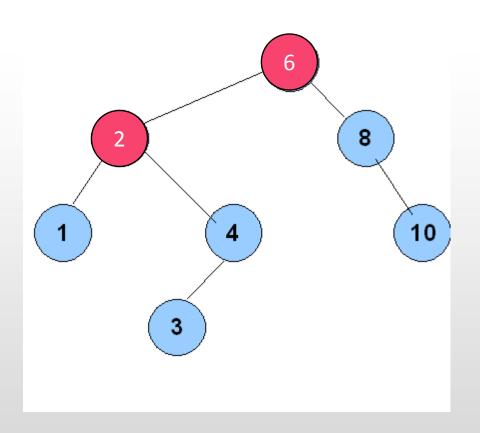




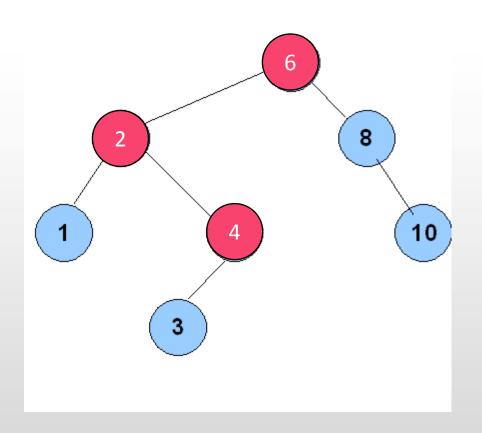




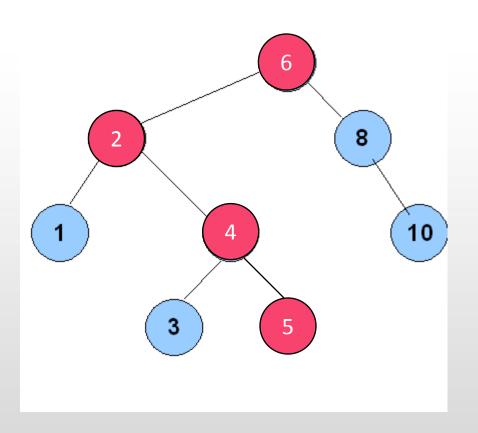












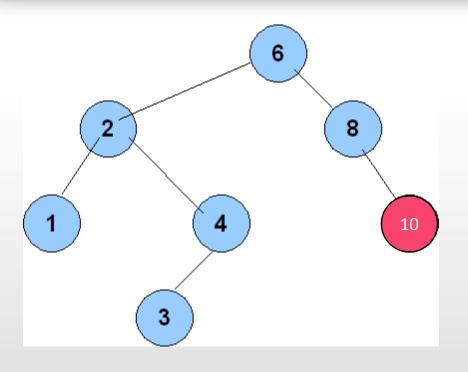


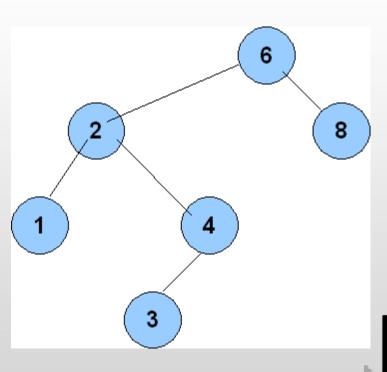
Deletion

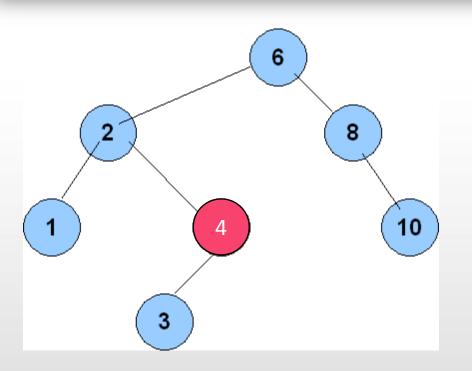
- 3 Cases
 - Node is a leaf
 - Node is non-leaf with 1 child
 - Node is non-leaf with 2 children

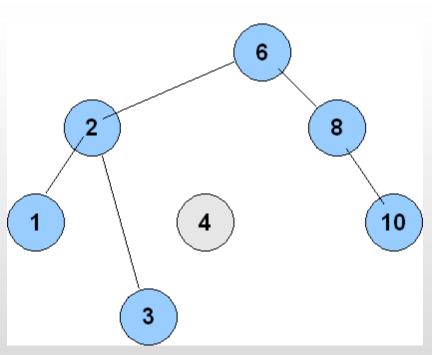


Deleting a leaf node

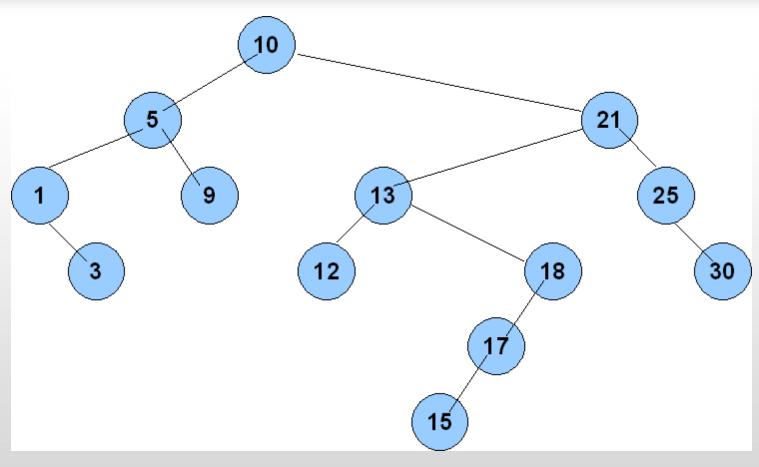






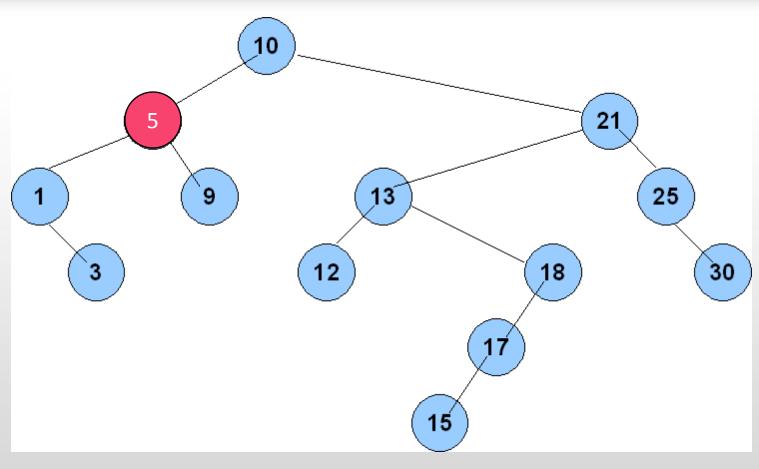






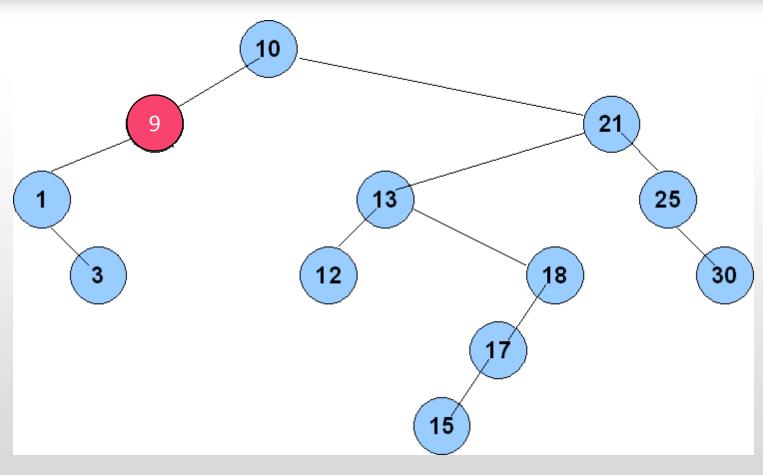
• Delete 5





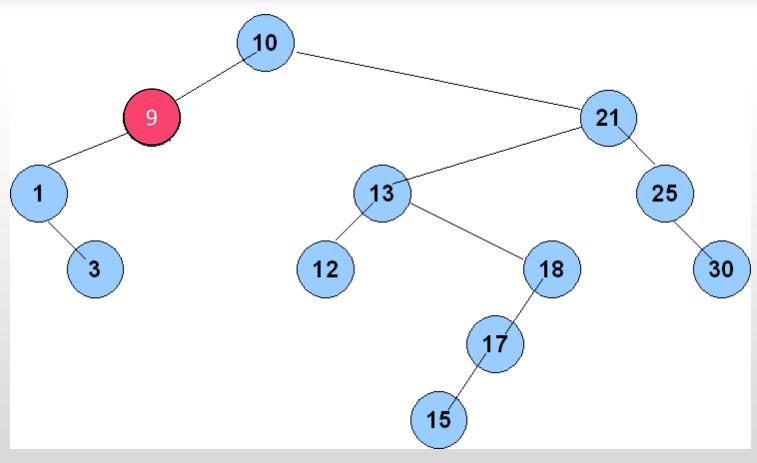
• Delete 5





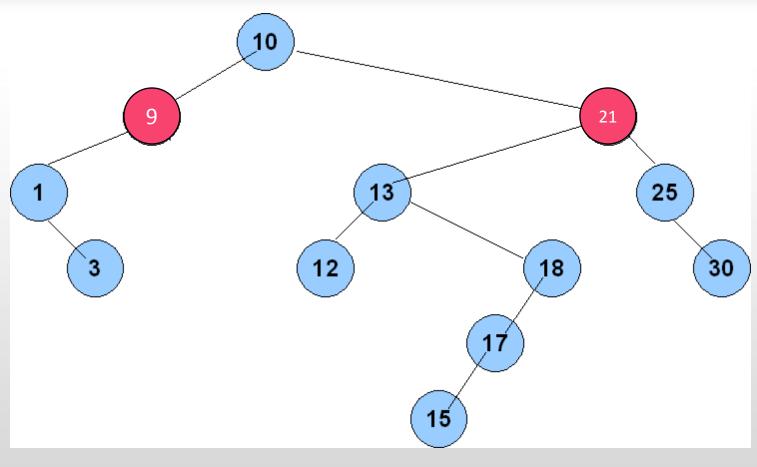
• Delete 5





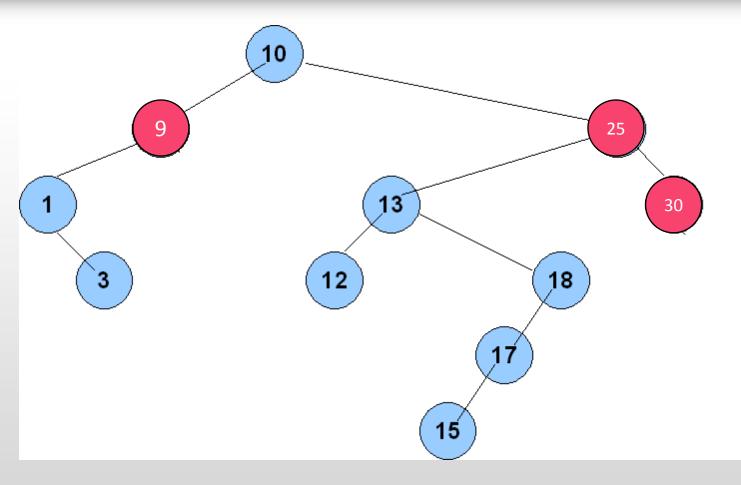
Delete 5, 21





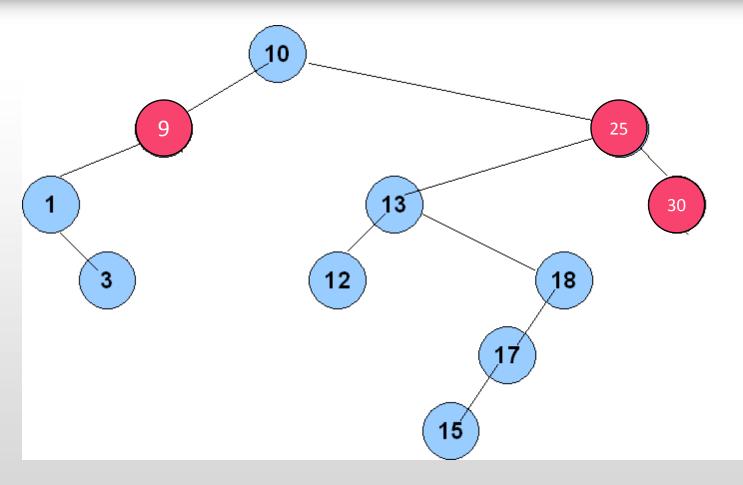
• Delete 5, 21





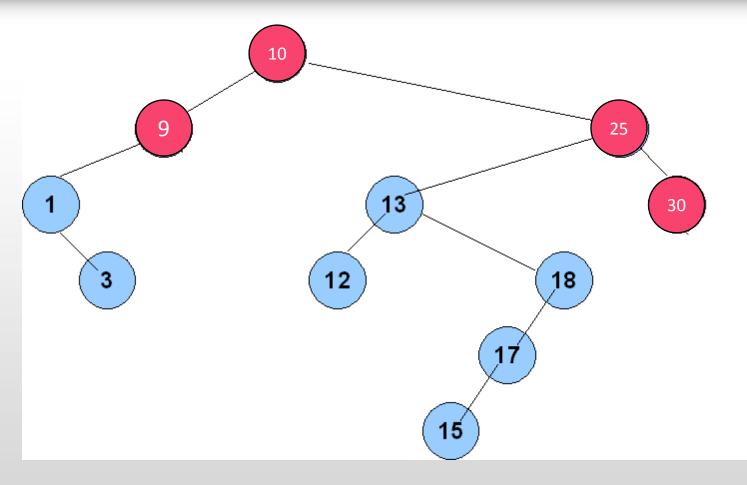
• Delete 5, 21





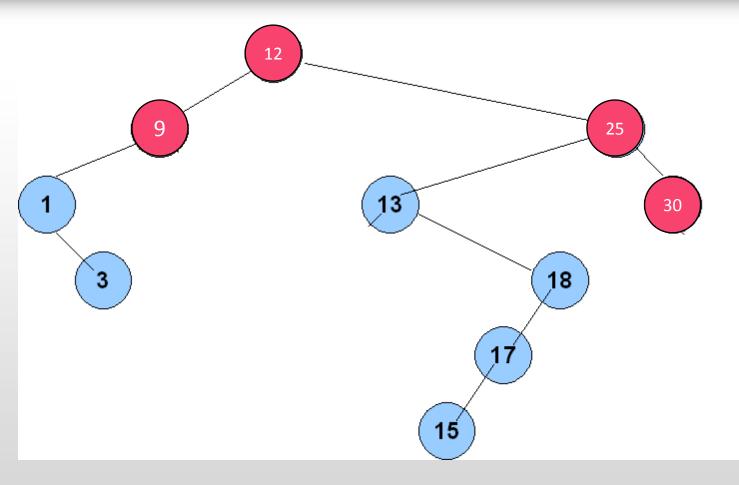
• Delete 5, 21, 10





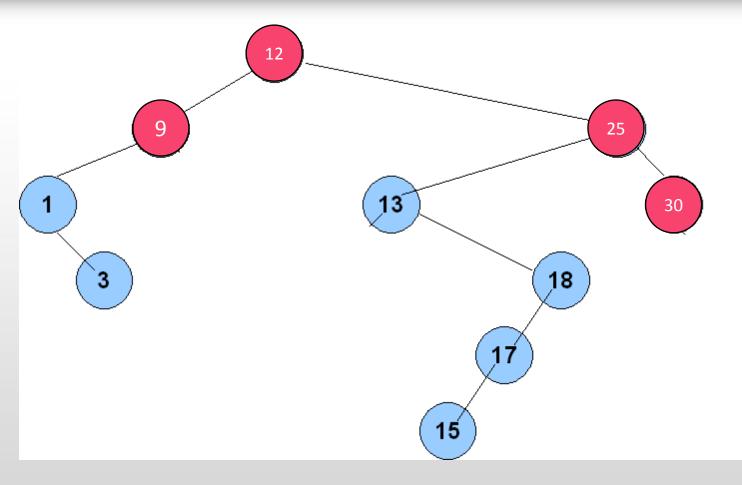
• Delete 5, 21, 10





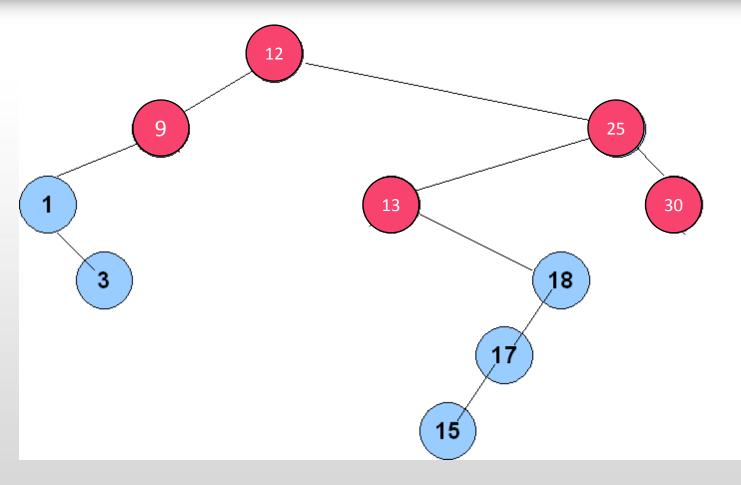
• Delete 5, 21, 10





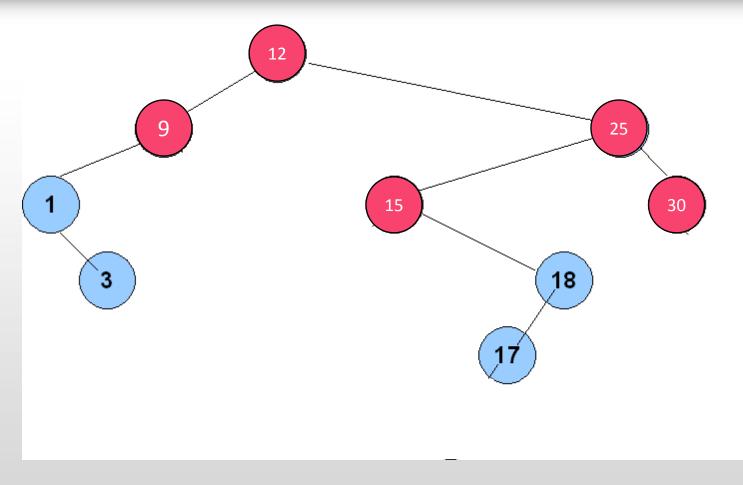
• Delete 5, 21, 10, 13





• Delete 5, 21, 10, 13





• Delete 5, 21, 10, 13



Other Operations

- Successor
 - find minimum element in the right subtree
- Predecessor
 - find maximum element in the left subtree



Print BST

```
typedef struct node{
  int value;
  struct node *left;
  struct node *right;
  struct node *parent;
}BST;

int main() {
  BST *root=NULL;
}
```

```
void print_sorted(BST *root) {
  if(root!=NULL) {
    print_sorted(root->left);
    printf("%d", root->value);
    print_sorted(root->right);
  }
}
```



BST - drawback

