

LINKED LISTS

OBJECTIVES

To learn what linked lists are

To insert and delete nodes in LL

What is a
LINKED LIST?

A data structure that consists of dynamic variables linked together to form a chain-like structure.

Used as an
ALTERNATIVE
to arrays

During execution, the linked list can
either...

GROW *or* **SHRINK**

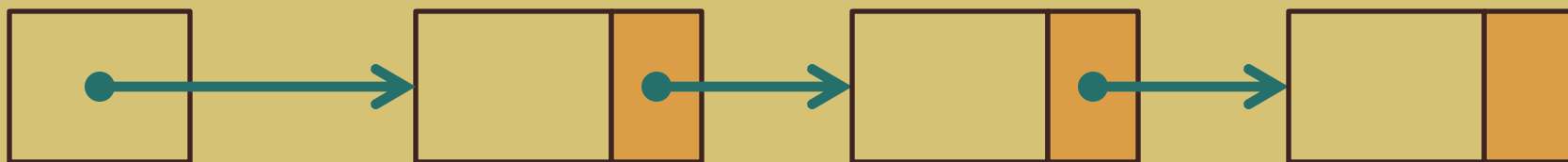
Used when the
DATA SIZE VARIES
during execution

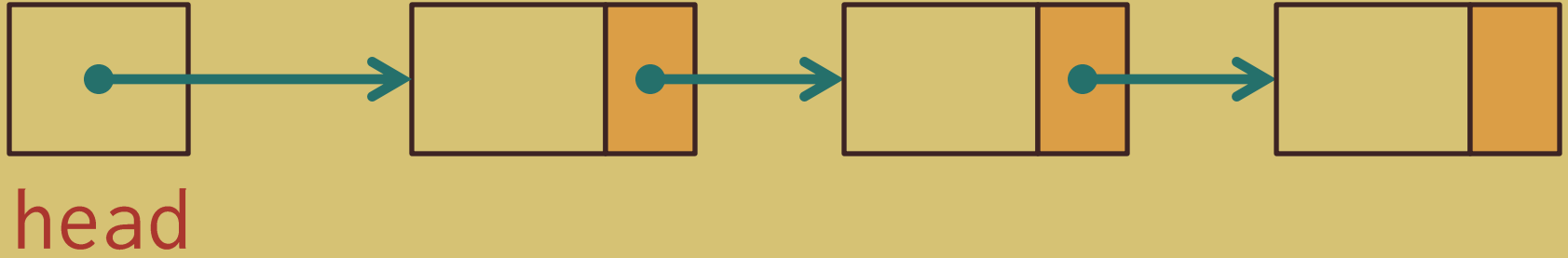
Unlike arrays, linked lists
SAVE MEMORY

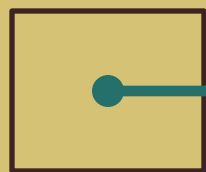
Allocated memory will
never exceed what is needed by
the program.

Dynamic arrays can handle the
change in the maximum size of
data, ...

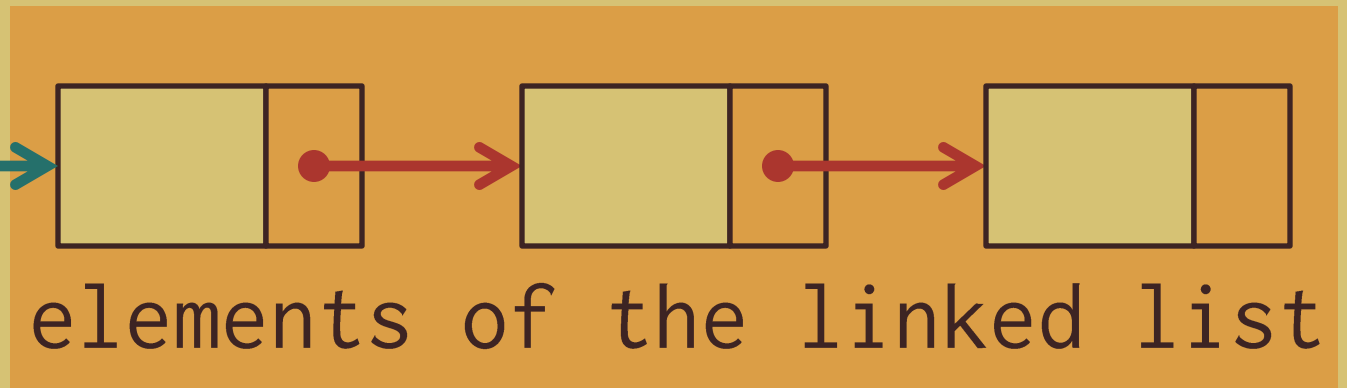
... but it is possible that the
allocated memory
will not be used.

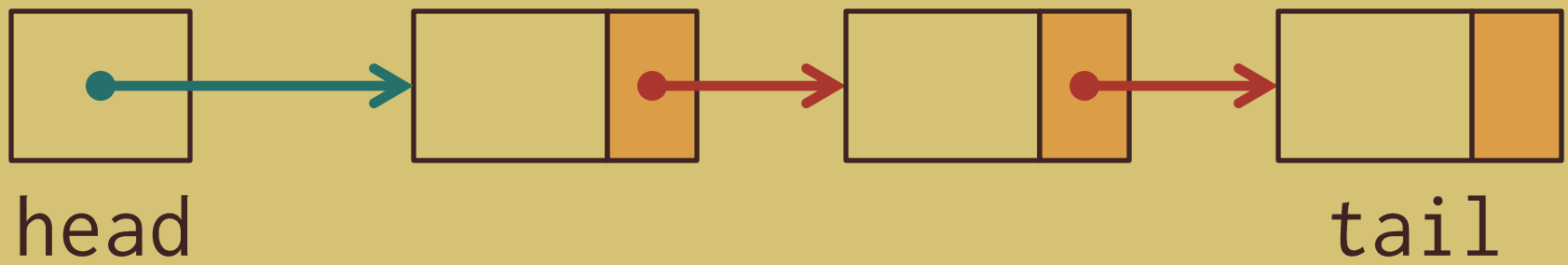


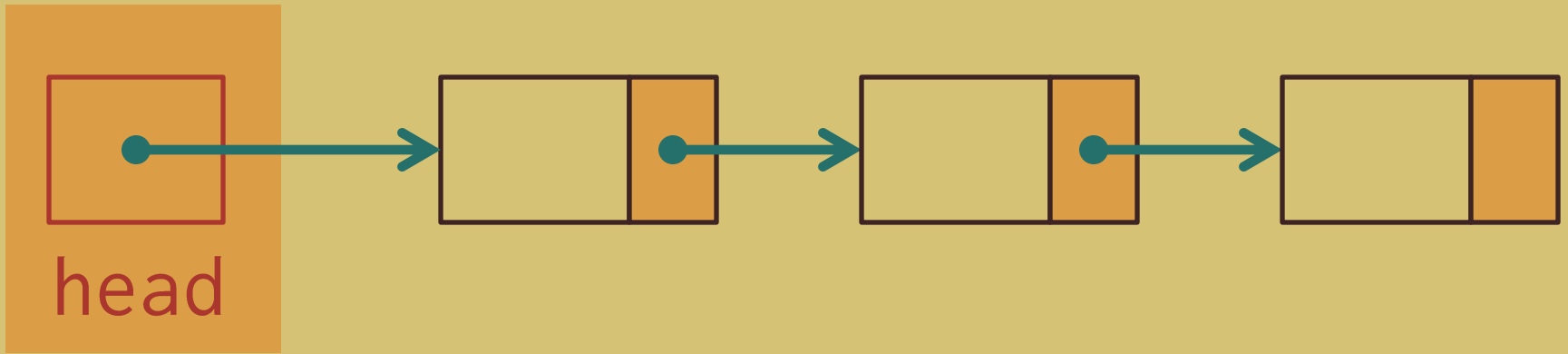




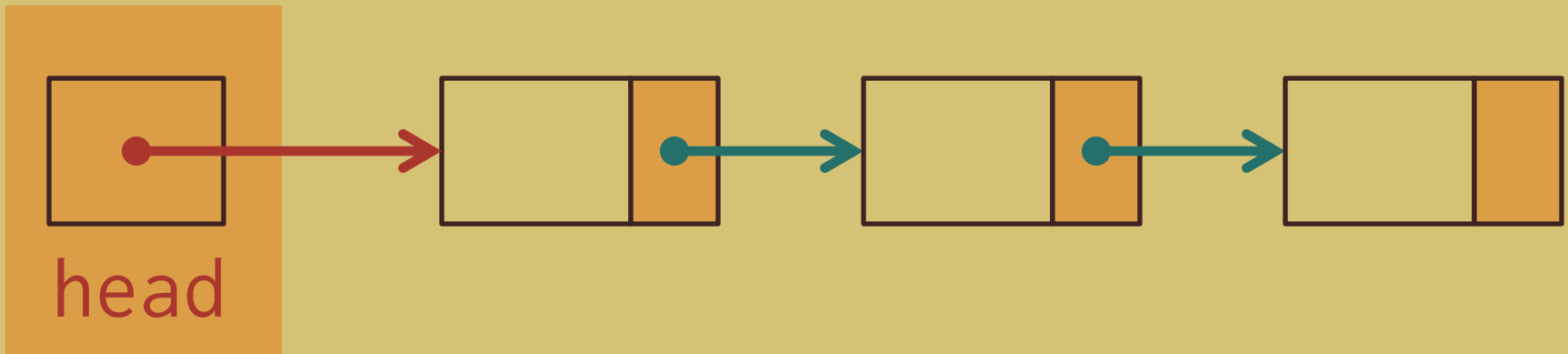
head



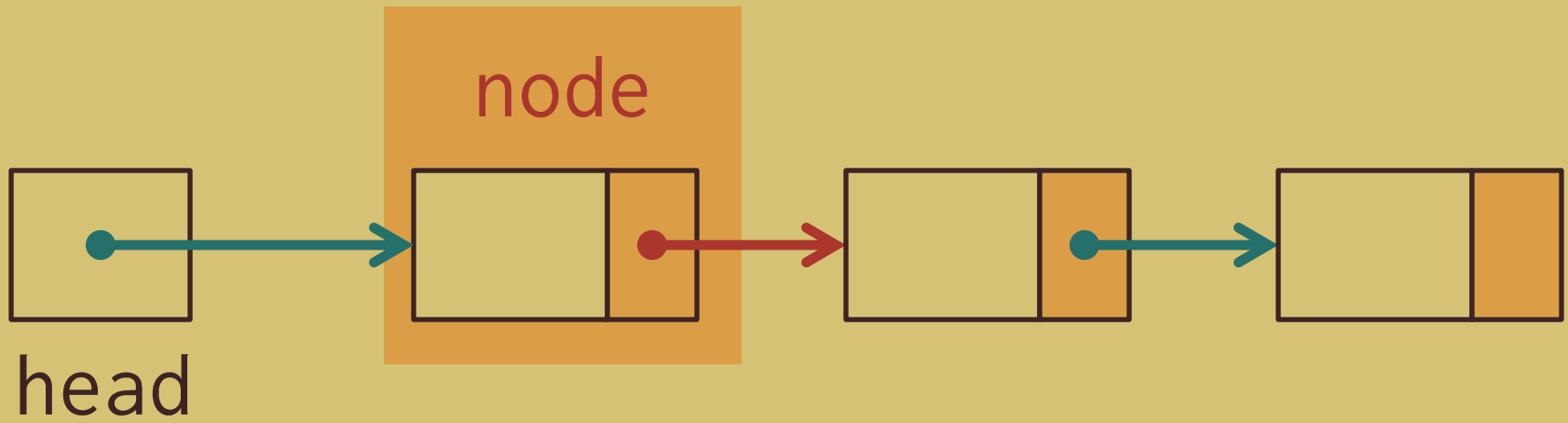


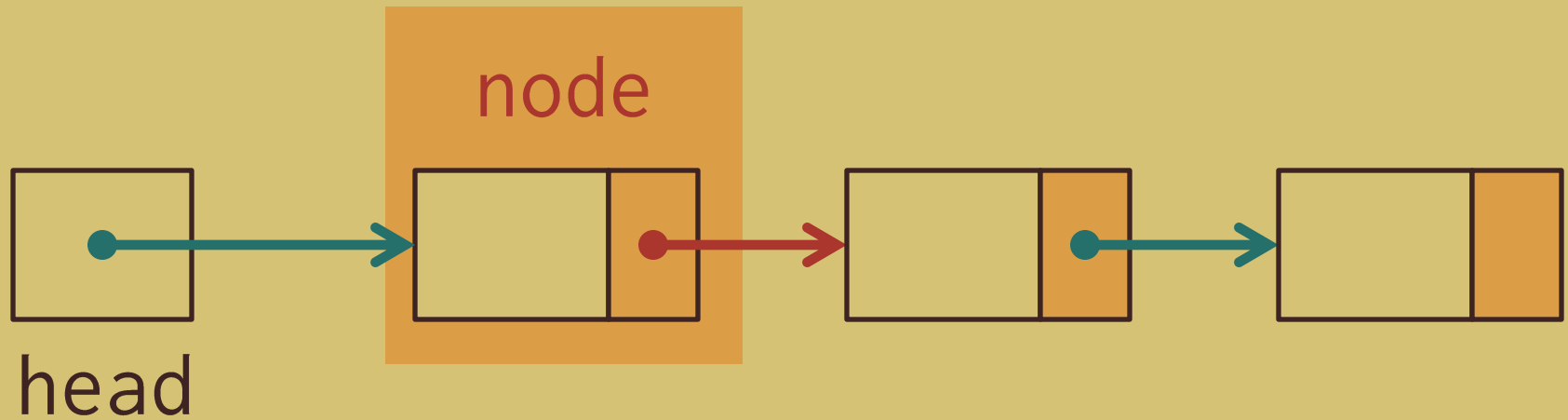


The head is a **pointer** (to a structure).

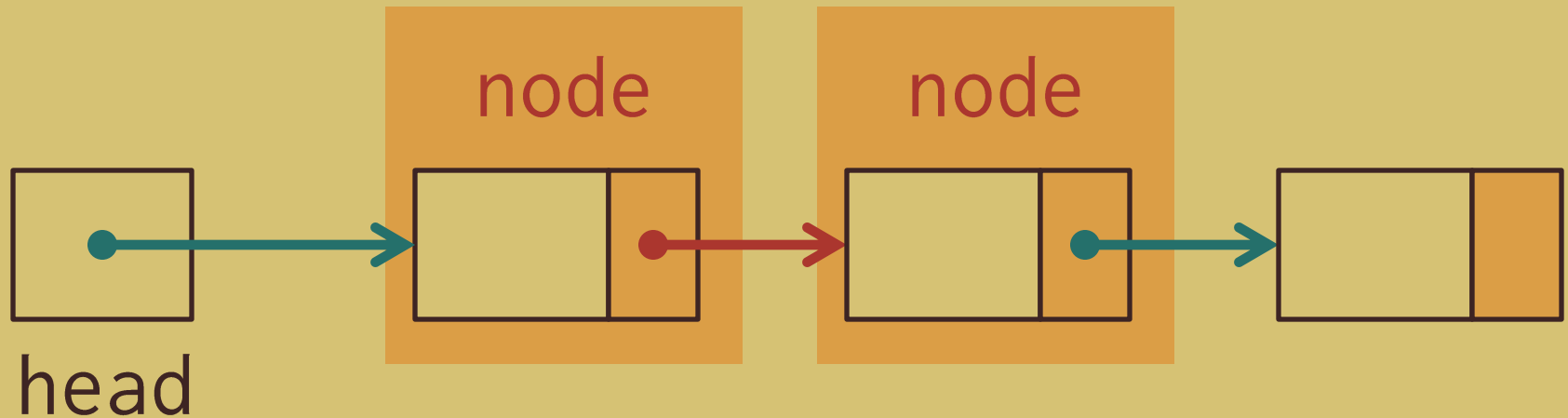


It holds the address of the first element of the linked list.



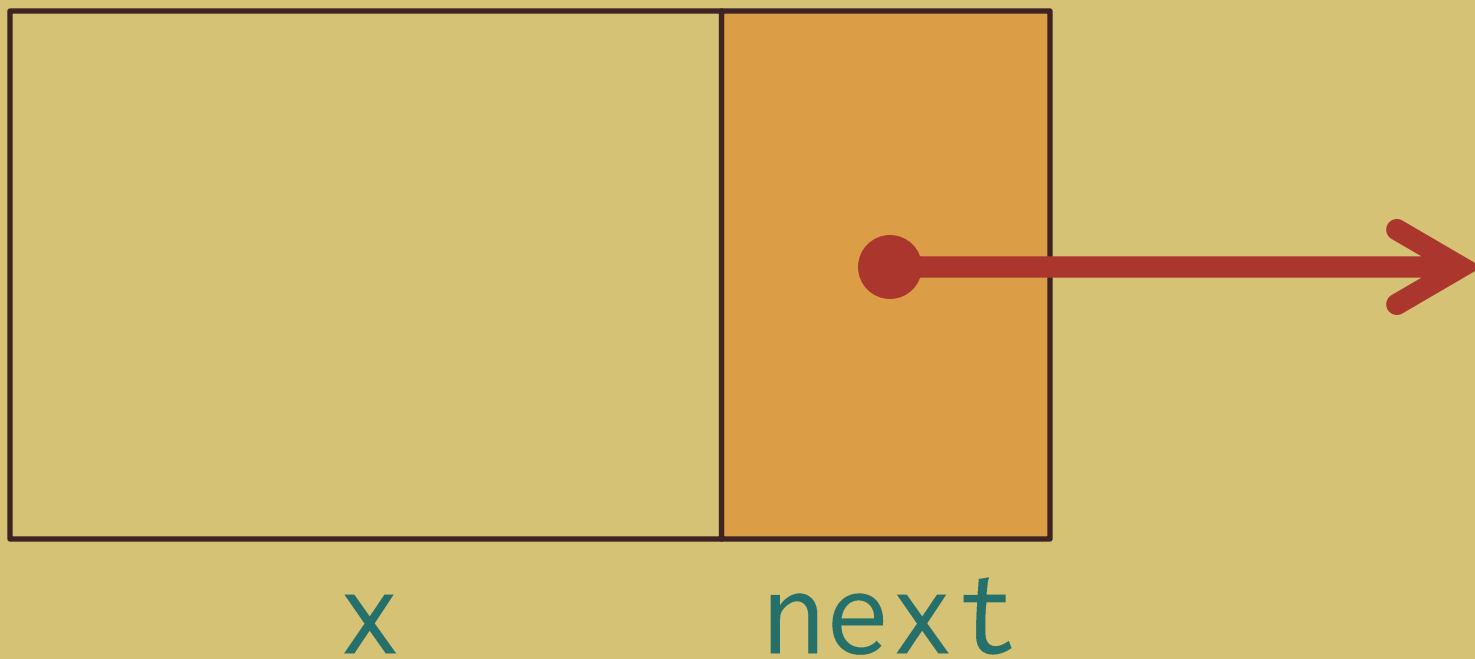


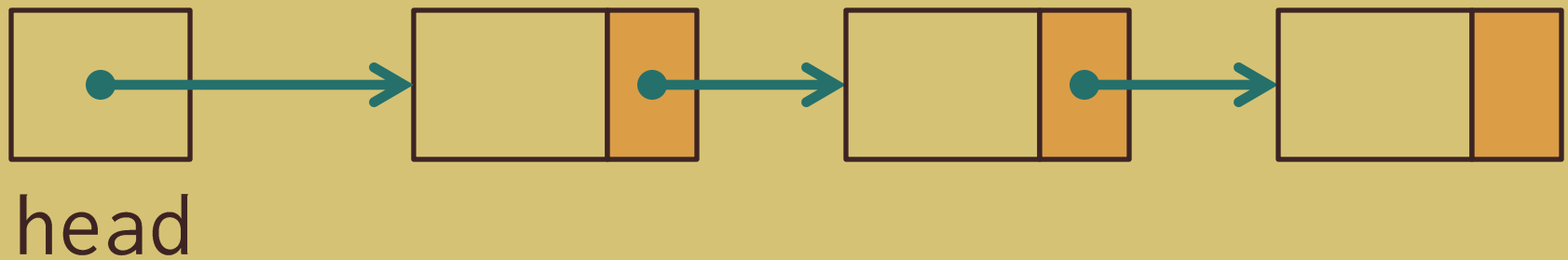
A node is a **self-referential structure**^{*}.



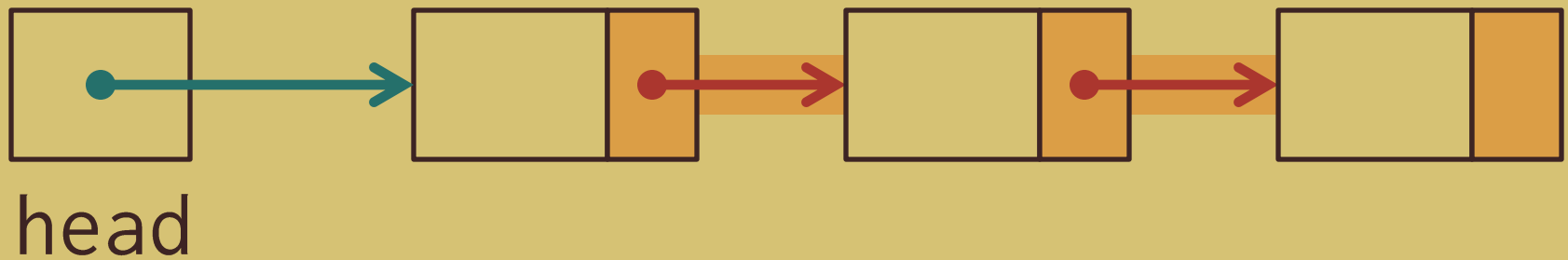
*A structure that has a pointer to an instance of itself as a field.

```
struct node {  
    int x;  
    //pointer to an instance of  
    //struct node  
    struct node *next;  
}
```

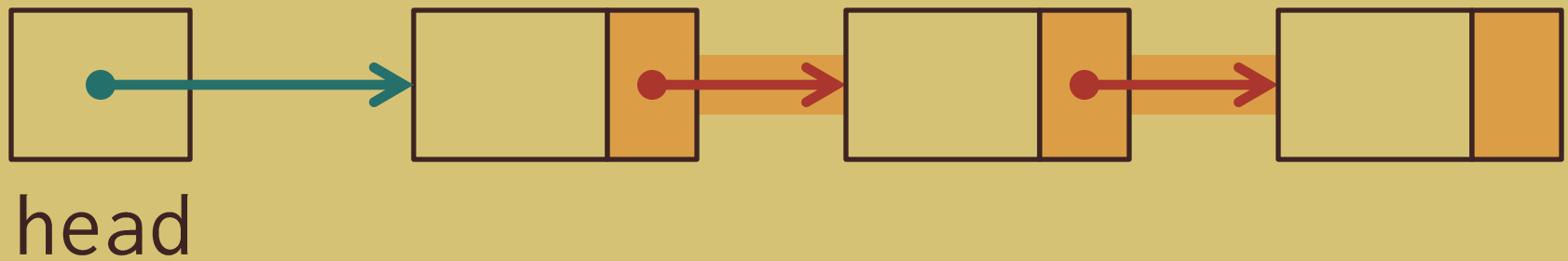




`malloc()` and `free()` are used to dynamically **grow** and **shrink** the linked list.



The pointer in each node **will point to the next** node in the list.



May have a **dummy node**.

Kinds of **LINKED LISTS:**

SINGLY LINKED LISTS

SINGLY LINKED LISTS

DOUBLY LINKED LISTS

CIRCULAR SINGLY LINKED LISTS

***CIRCULAR SINGLY
LINKED LISTS***

***CIRCULAR DOUBLY
LINKED LISTS***

A dummy node is a node in the linked list that does not contain data.

It is used to simplify some
linked lists operations.

The 4 basic
OPERATIONS ON
LINKED LISTS:

INSERT

DELETE

SEARCH

VIEW

INSERT

DELETE

INSERT

Deals with inserting values to the linked list.

INSERT

Different CASES of INSERT:

INSERT

AT HEAD

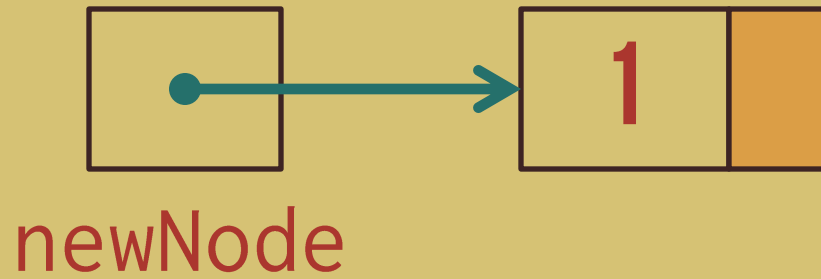
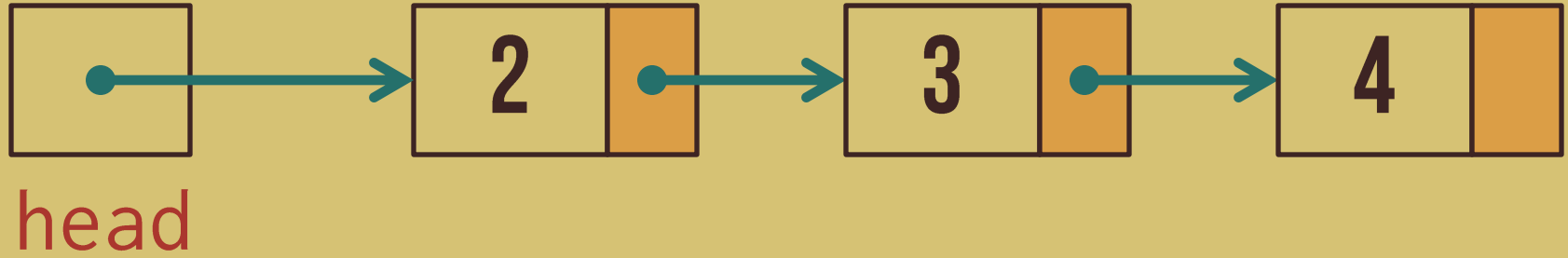
INSERT

AT HEAD

You have a new node that you
want to insert at the
beginning.

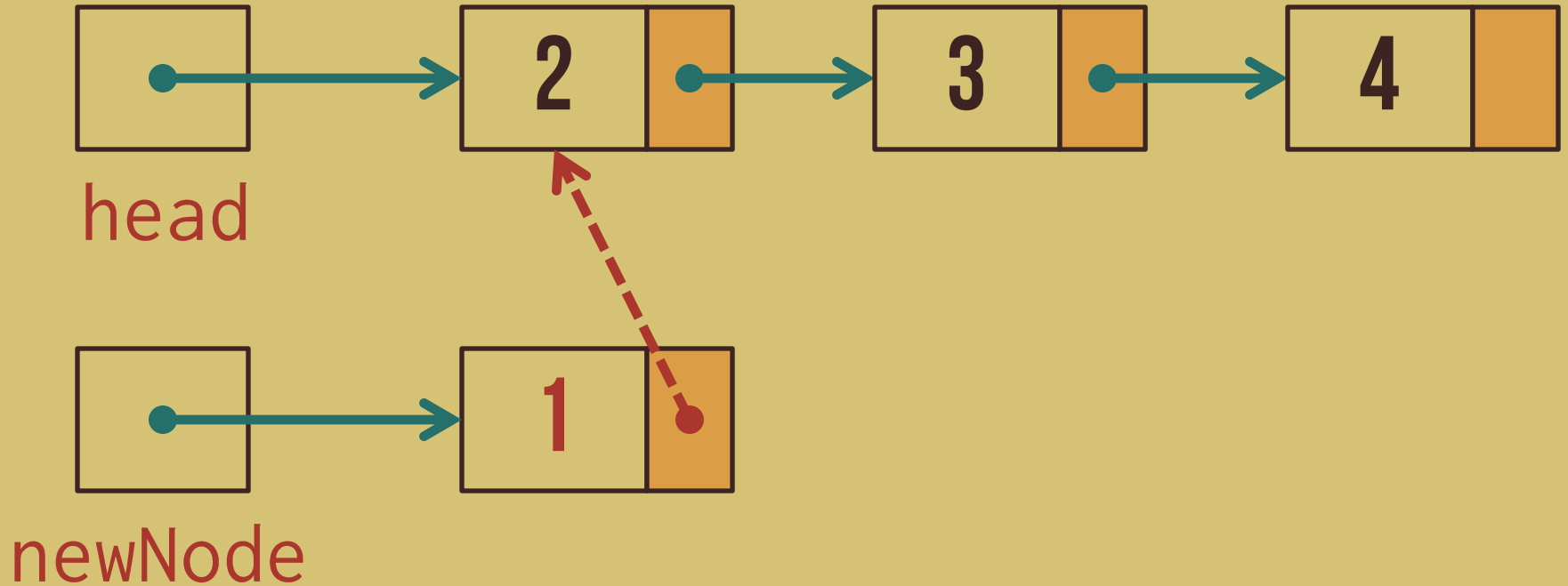
INSERT

AT HEAD



INSERT

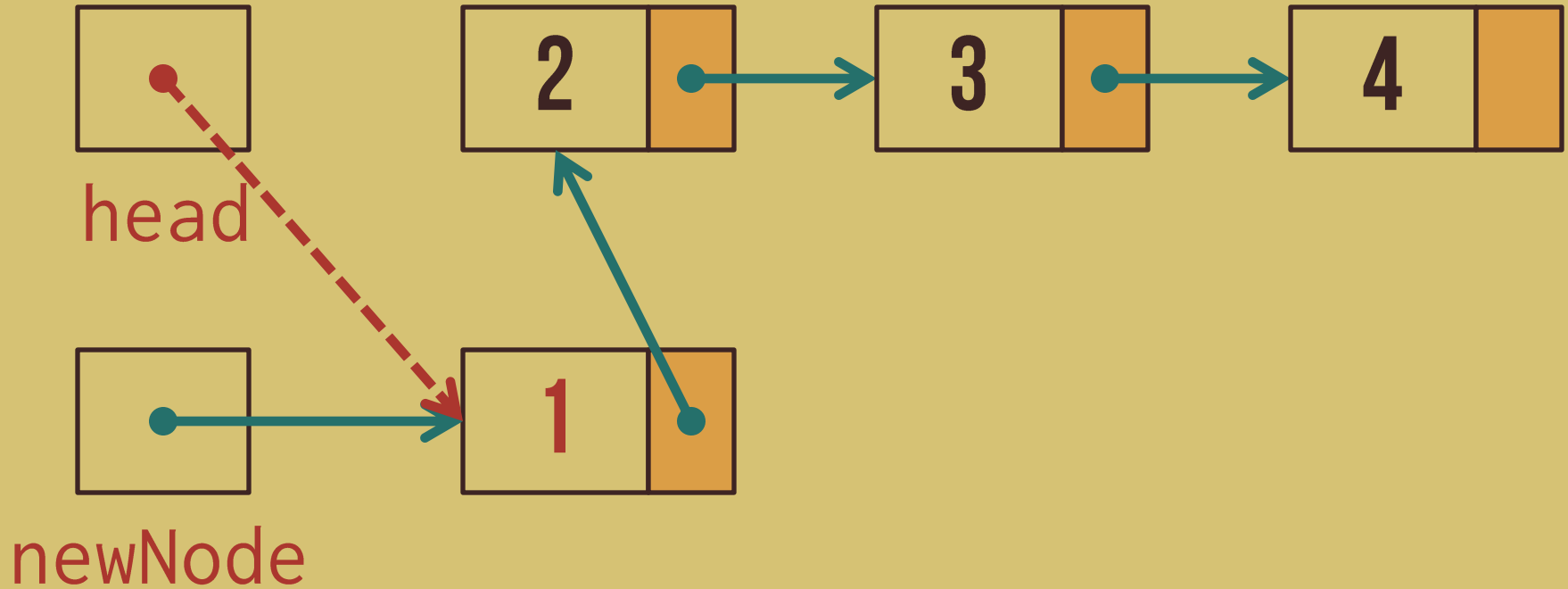
AT HEAD



1. Make the next pointer of the new node (**1**) point to the node pointed to by head (**2**).

INSERT

AT HEAD



2. Make the head pointer point to the new node (1).

INSERT

AT HEAD



newNode

INSERT

AT MIDDLE

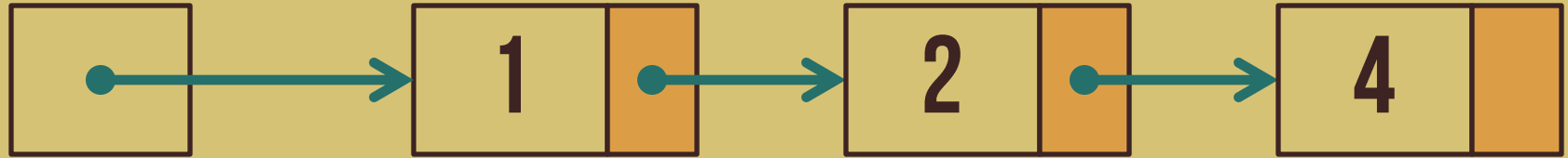
INSERT

AT MIDDLE

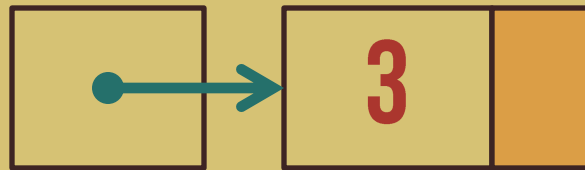
You have a new node that you want to insert somewhere in between the linked list.

INSERT

AT MIDDLE



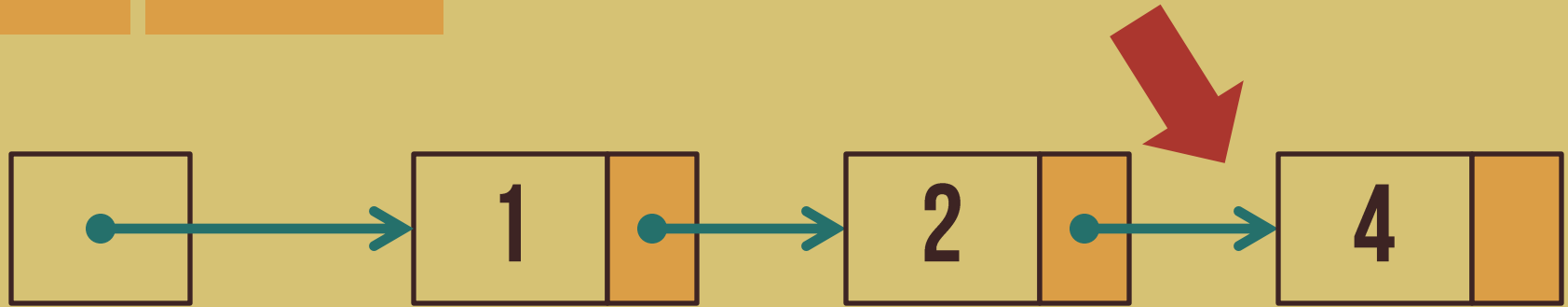
head



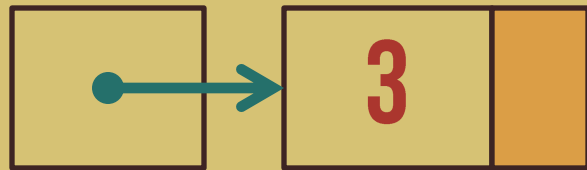
newNode

INSERT

AT MIDDLE



head

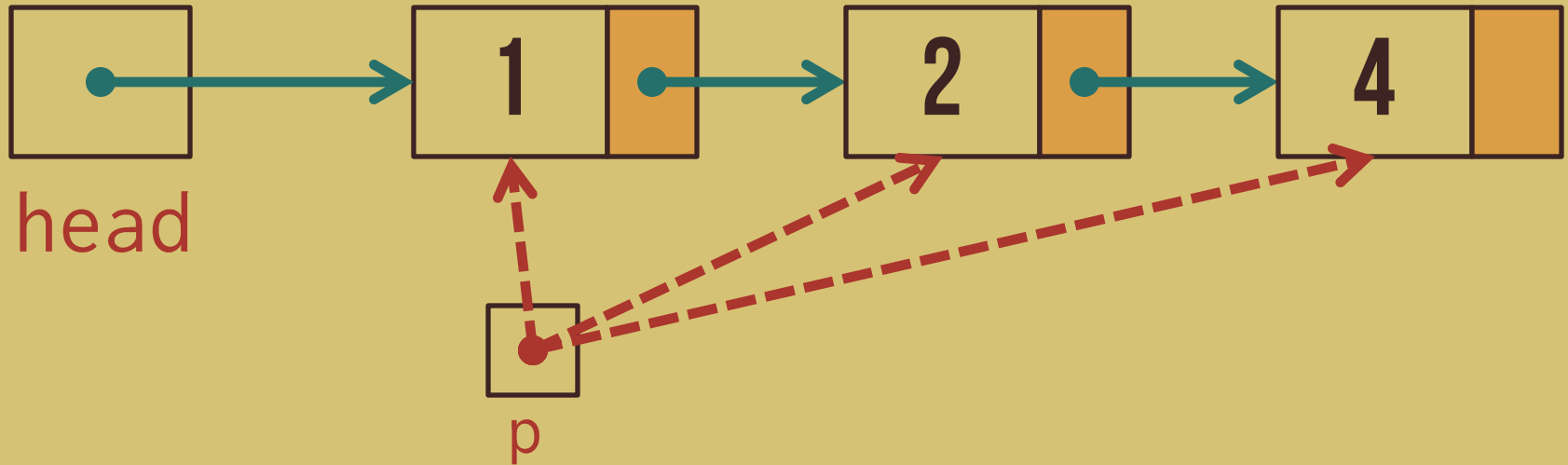


newNode

1a. Find the position where the node is to be inserted.

INSERT

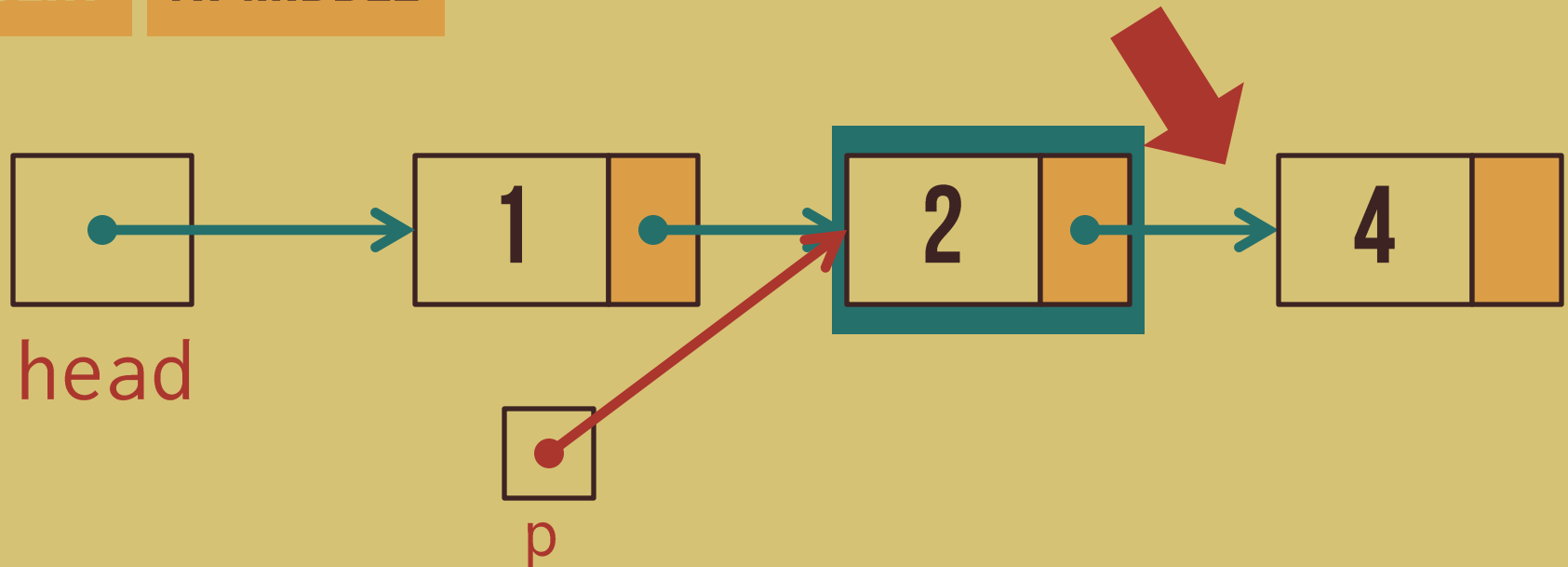
AT MIDDLE



We usually do this by traversing using a pointer (**p**).

INSERT

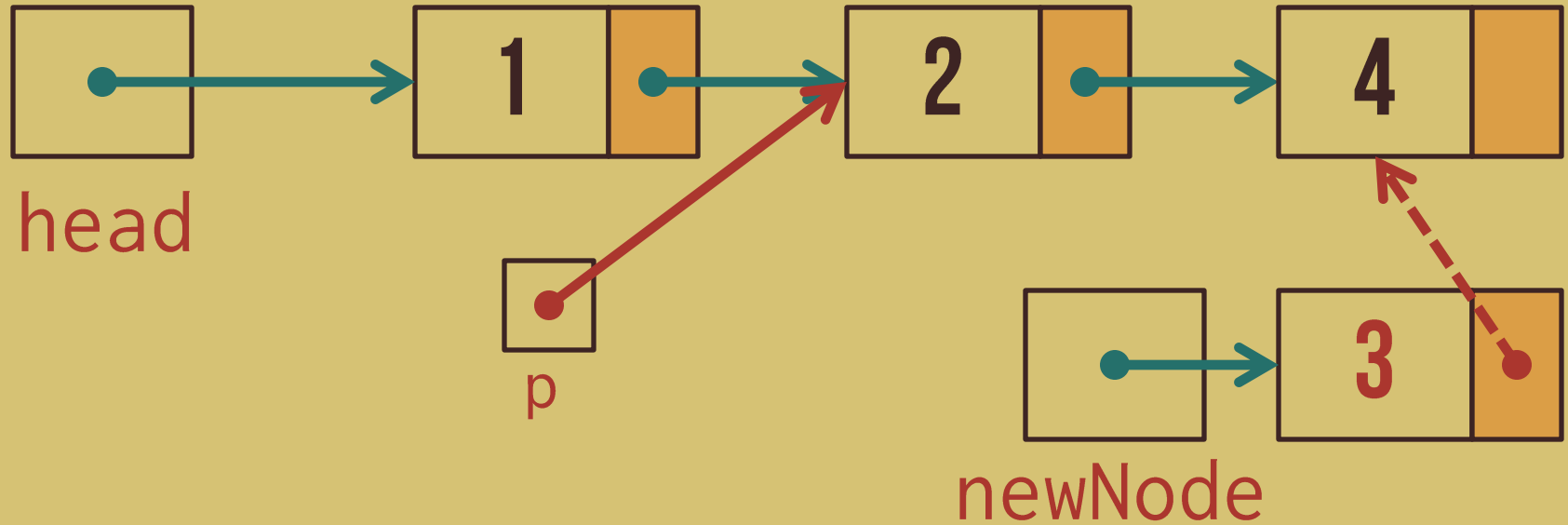
AT MIDDLE



1b. Let's **remember/select** the node *before* the position we want to insert to.

INSERT

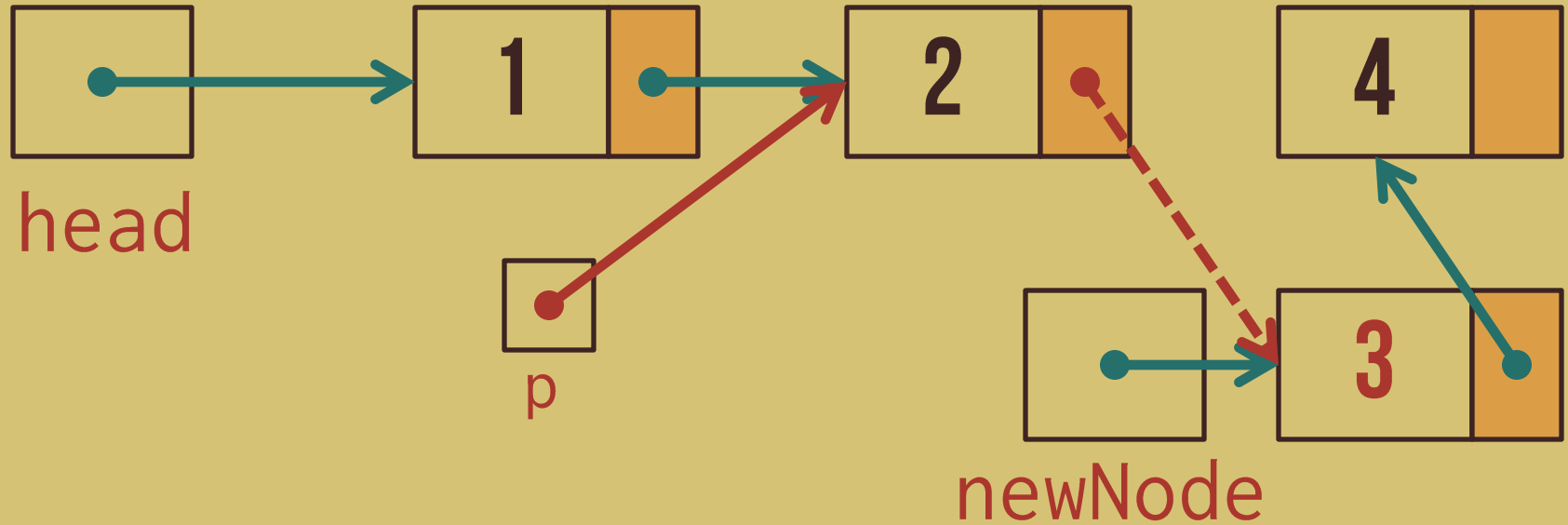
AT MIDDLE



2. Make the next pointer of the new node (**3**) point to the node next to the node we selected in step 1b (**2**).

INSERT

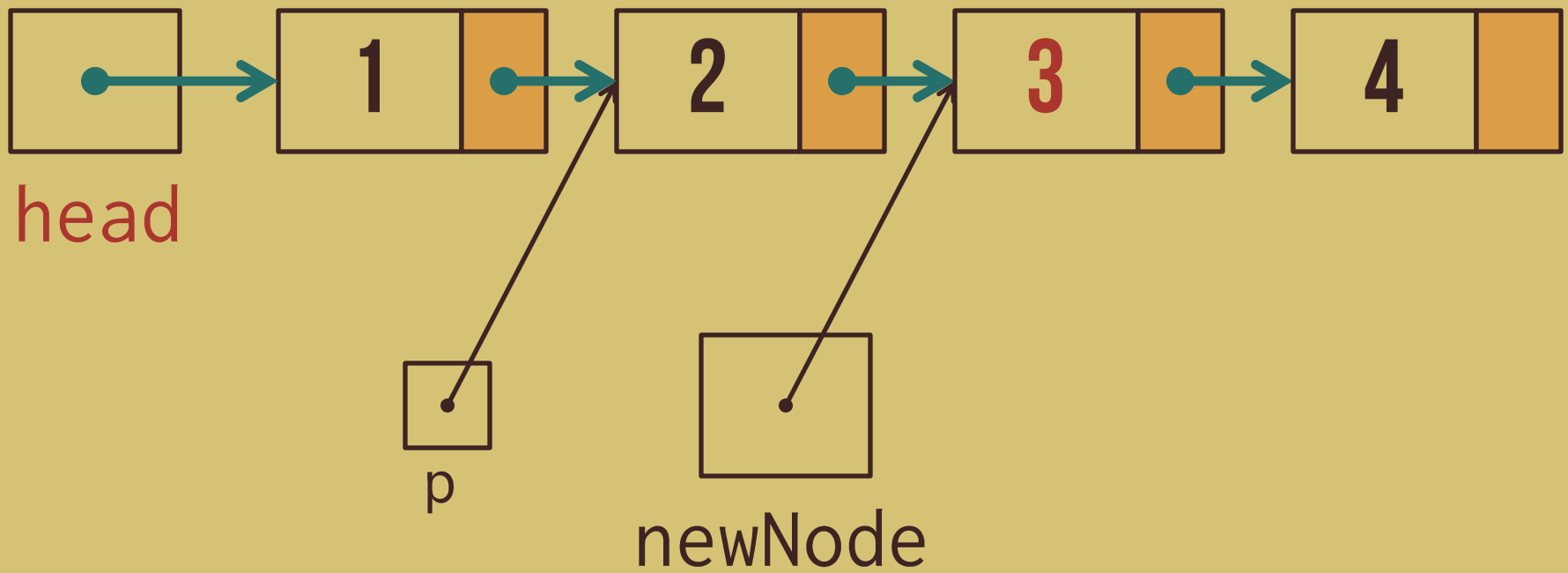
AT MIDDLE



3. Make the next pointer of the node selected in step 1b (**2**) point to the new node (**3**).

INSERT

AT MIDDLE



INSERT

AT MIDDLE

Usually used in conjunction
with insert at head.

INSERT

AT TAIL

INSERT

AT TAIL

Used when the position of the new value is at the end of the linked list.

INSERT

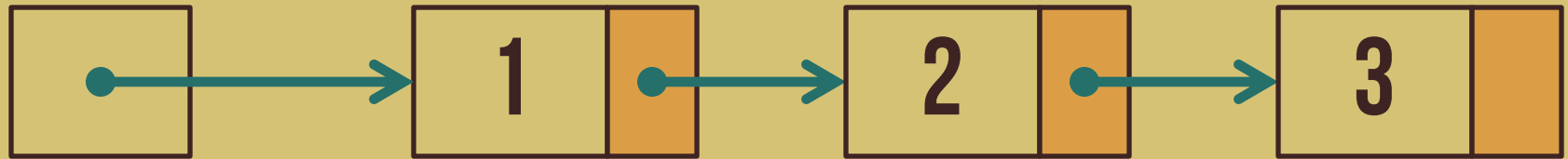
AT TAIL

Can be treated as a special
case of inserting at the
middle*.

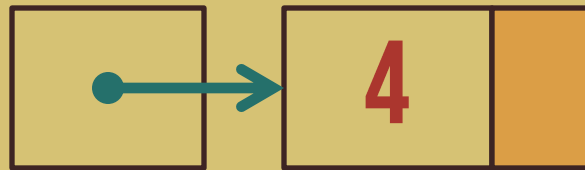
*There surely will be implementation differences.

INSERT

AT TAIL



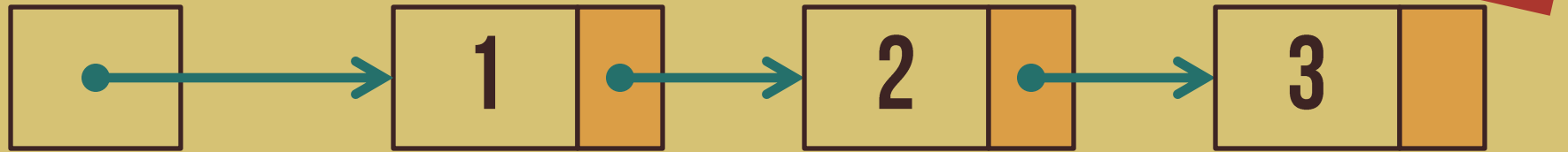
head



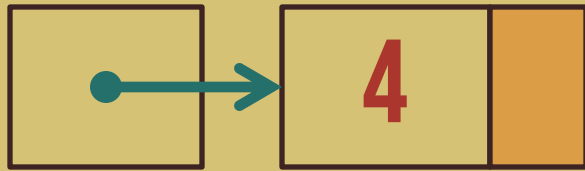
newNode

INSERT

AT TAIL



head

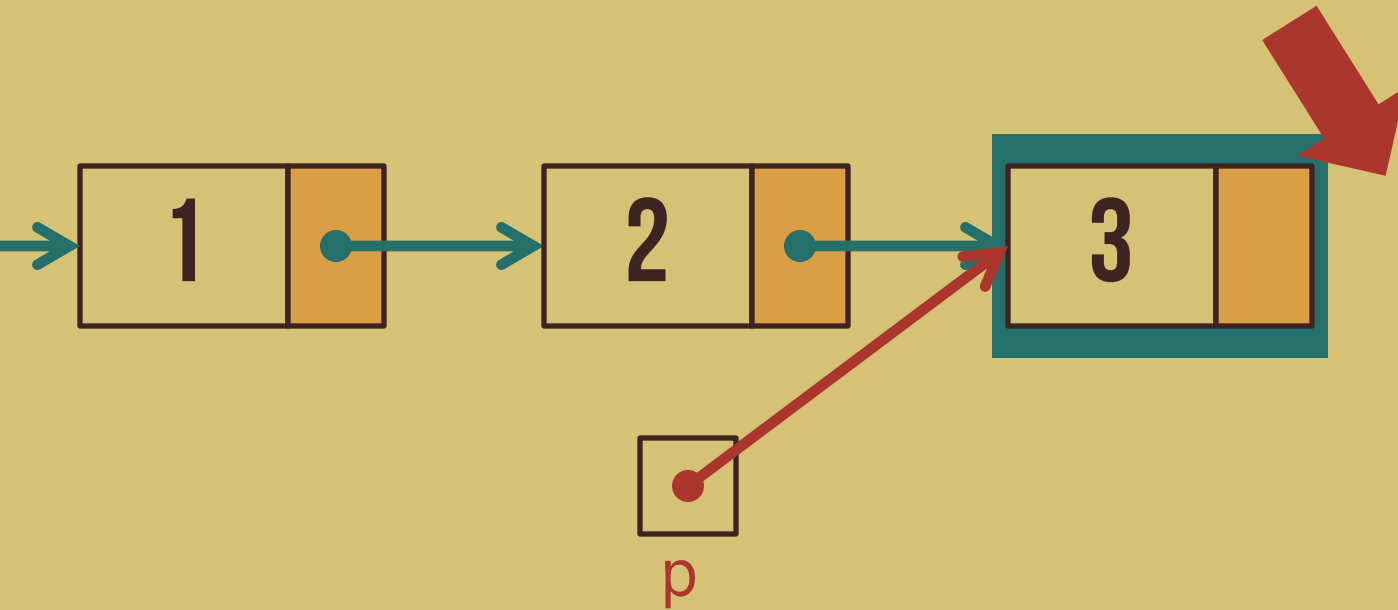


newNode

1a. Find the position where the node is to be inserted.

INSERT

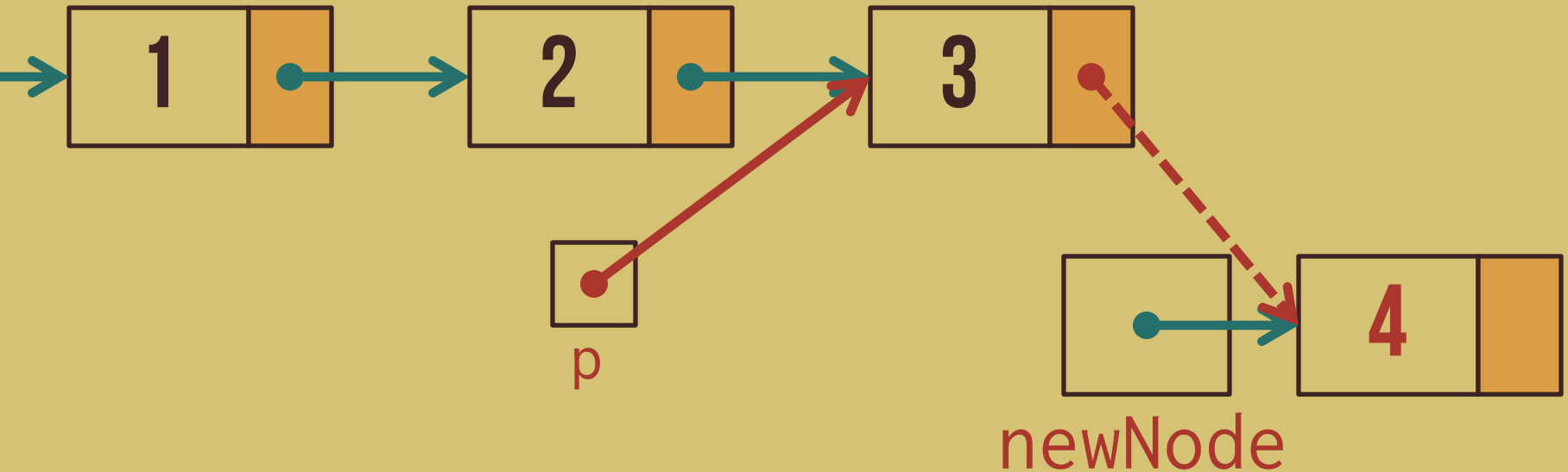
AT TAIL



1b. Let's **remember/select** the node *before* the position we want to insert to.

INSERT

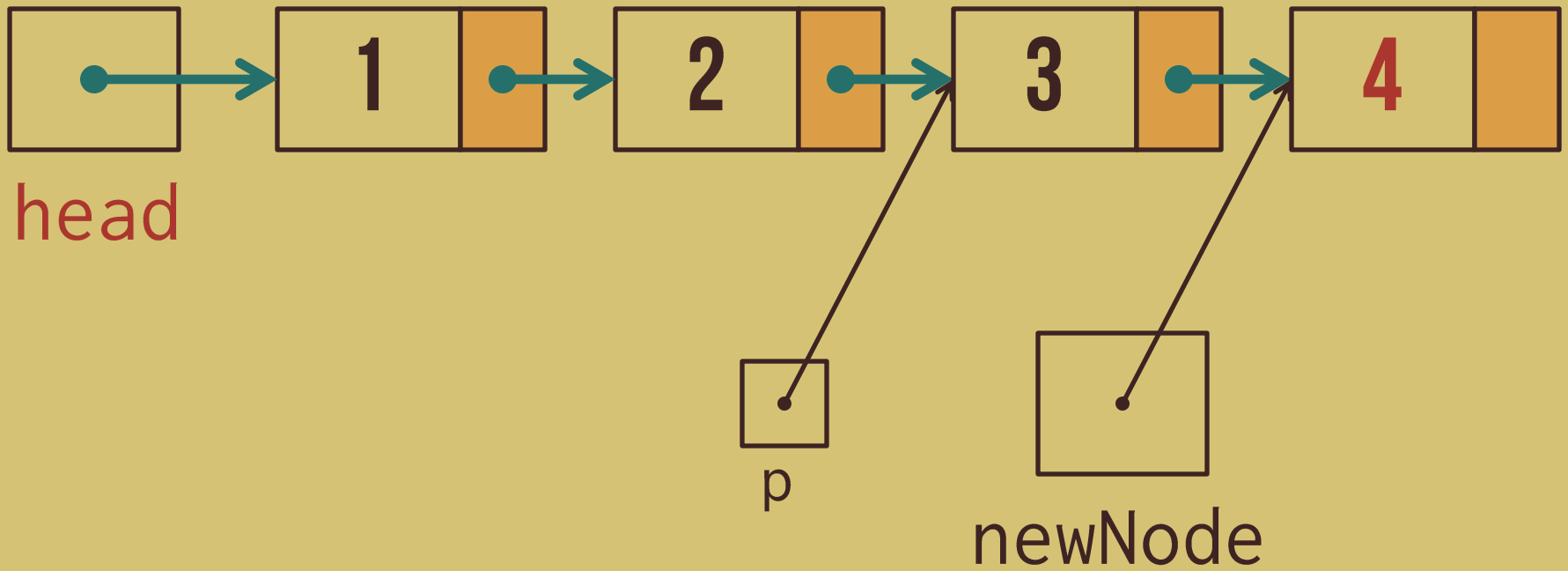
AT TAIL



2. Make the next pointer of the node selected in step 1b (**3**) point to the new node (**4**).

INSERT

AT TAIL



INSERT

SOME NOTES

To mark the end of the list, the next pointer of the last node is given the value **NULL**.

NULL is a const value defined
in `stdlib.h`.

In pointers, it's used to symbolize that a pointer is not pointing anywhere.

In our visualizations, we assume that a pointer field that does not have an outgoing arrow has a **NULL** value.

To prevent pointers having garbage values, ALWAYS initialize your pointers to NULL.

INSERT

DELETE

DELETE

Used to delete elements from
the linked list.

DELETE

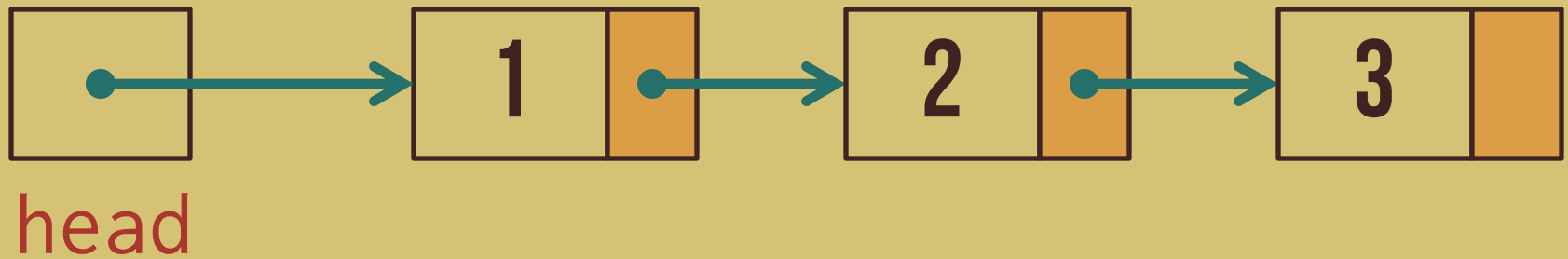
Different CASES of DELETE:

DELETE

AT HEAD

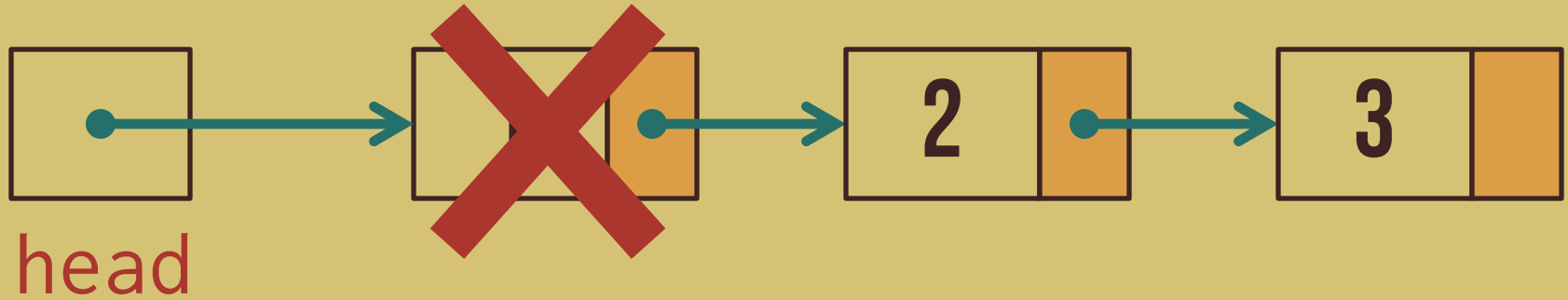
DELETE

AT HEAD



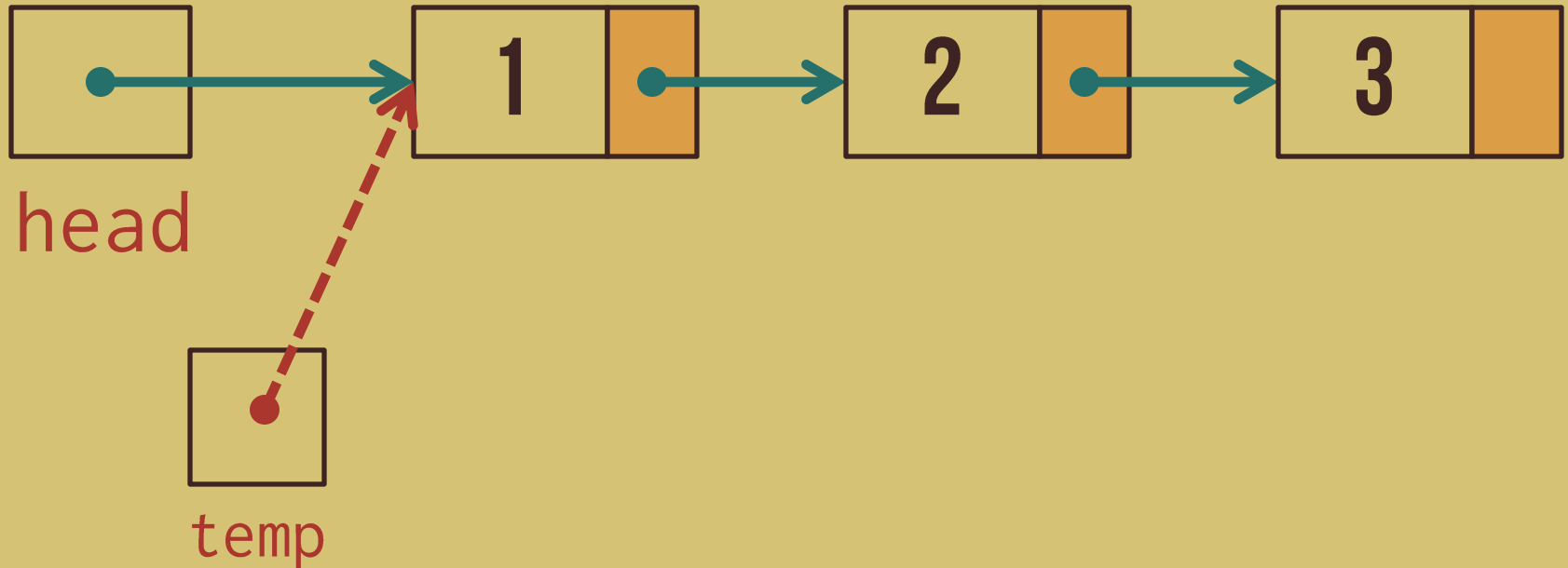
DELETE

AT HEAD



DELETE

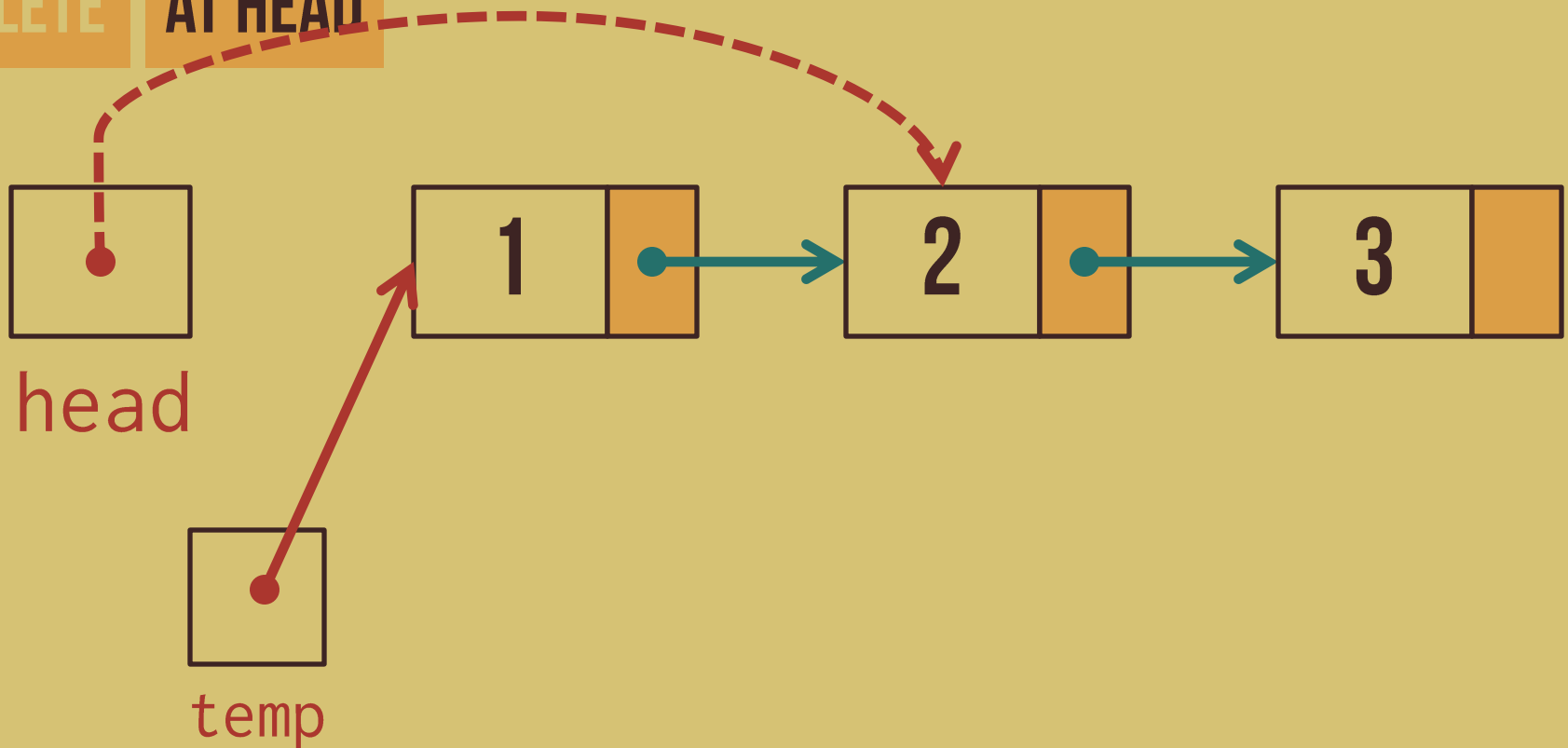
AT HEAD



1. Have **another pointer** (temp) point to the **first node**.

DELETE

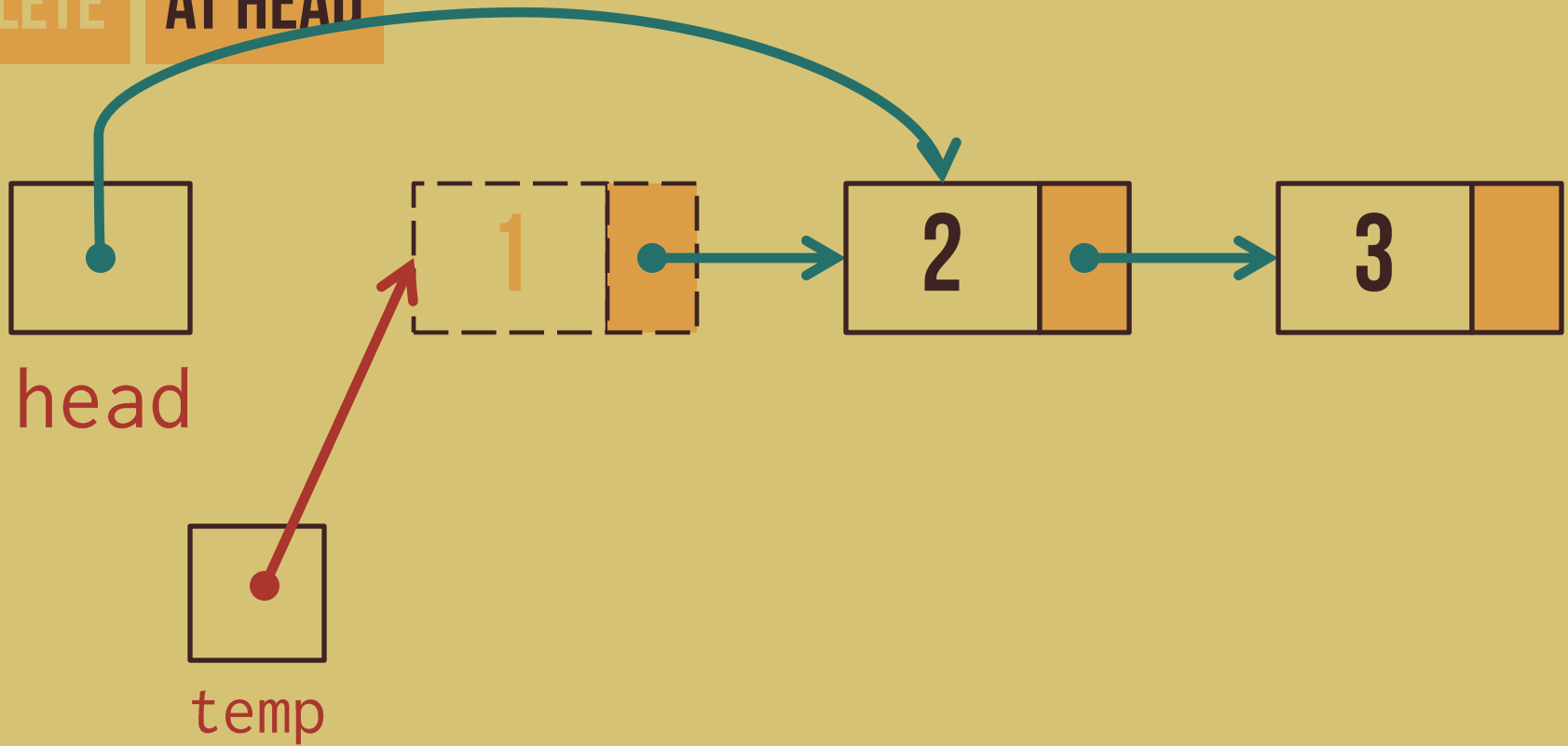
AT HEAD



2. Make **head** refer to the **second node**.

DELETE

AT HEAD



3. Free the **first node** using temp.

DELETE

AT HEAD

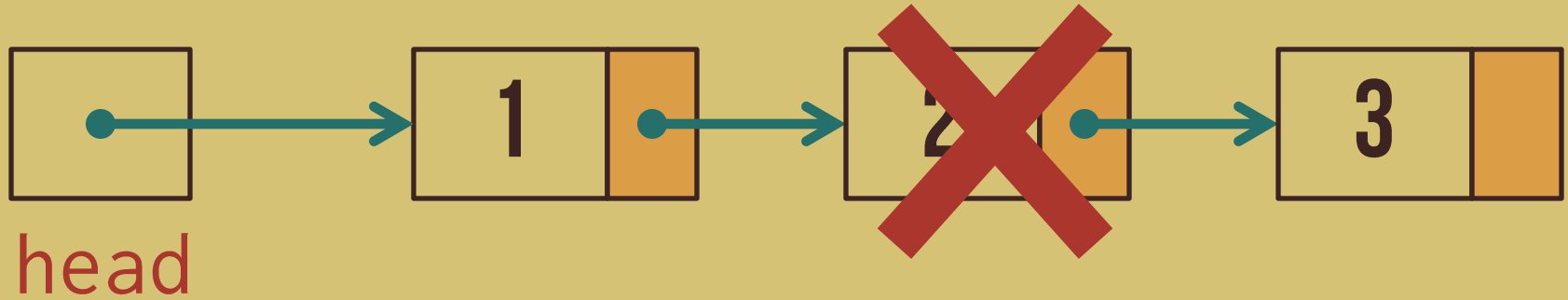


DELETE

AT MIDDLE

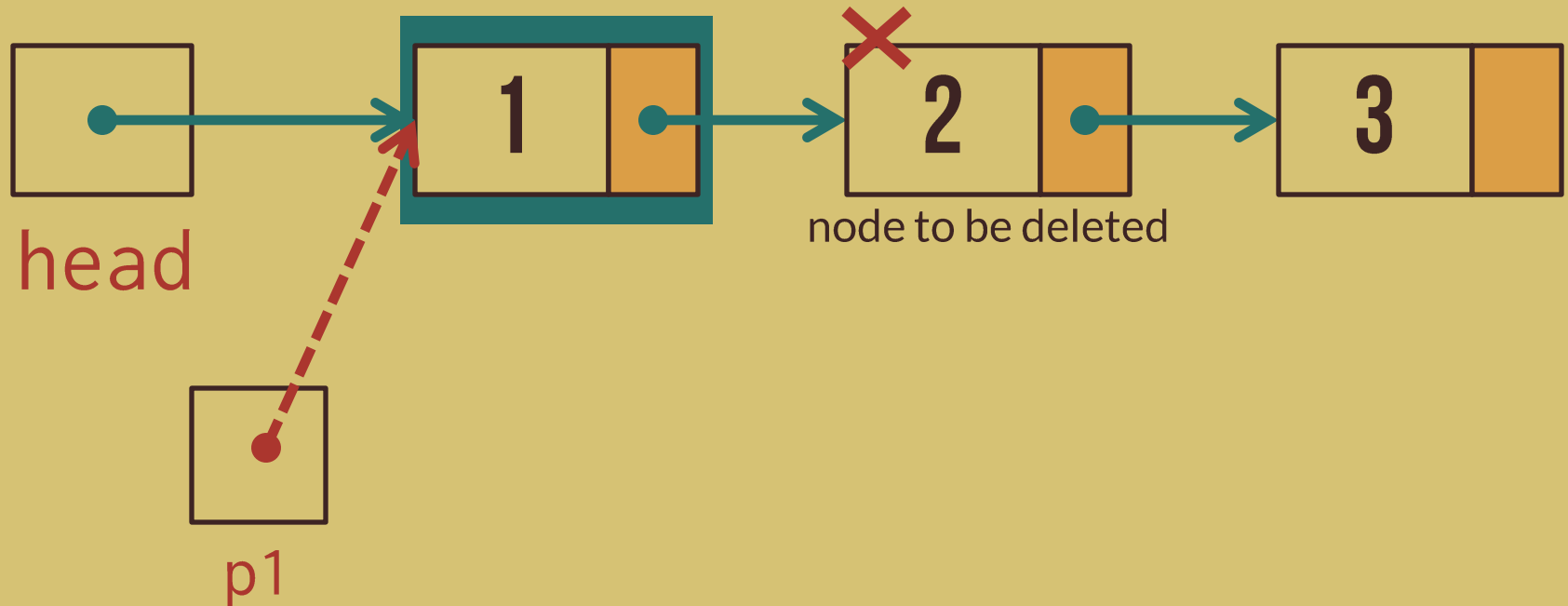
DELETE

AT MIDDLE



DELETE

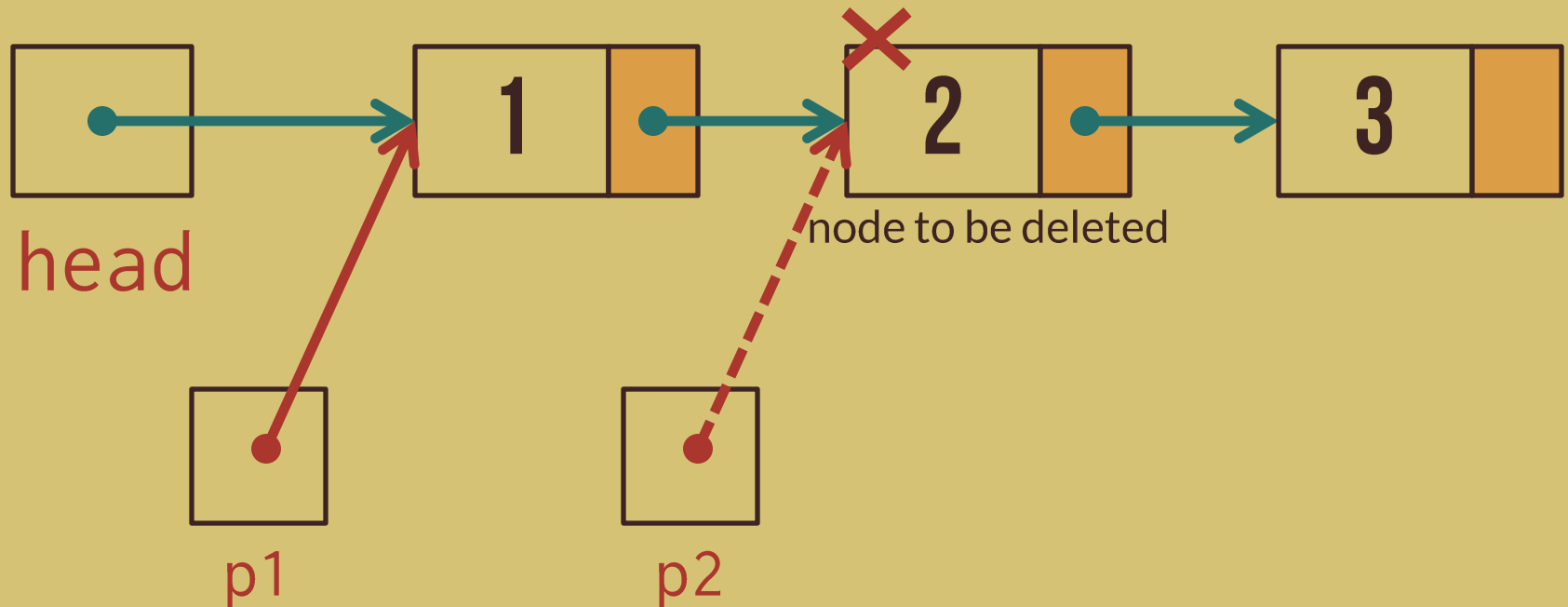
AT MIDDLE



1. Find the node *before* the node to be deleted and let a pointer (p1) point it.

DELETE

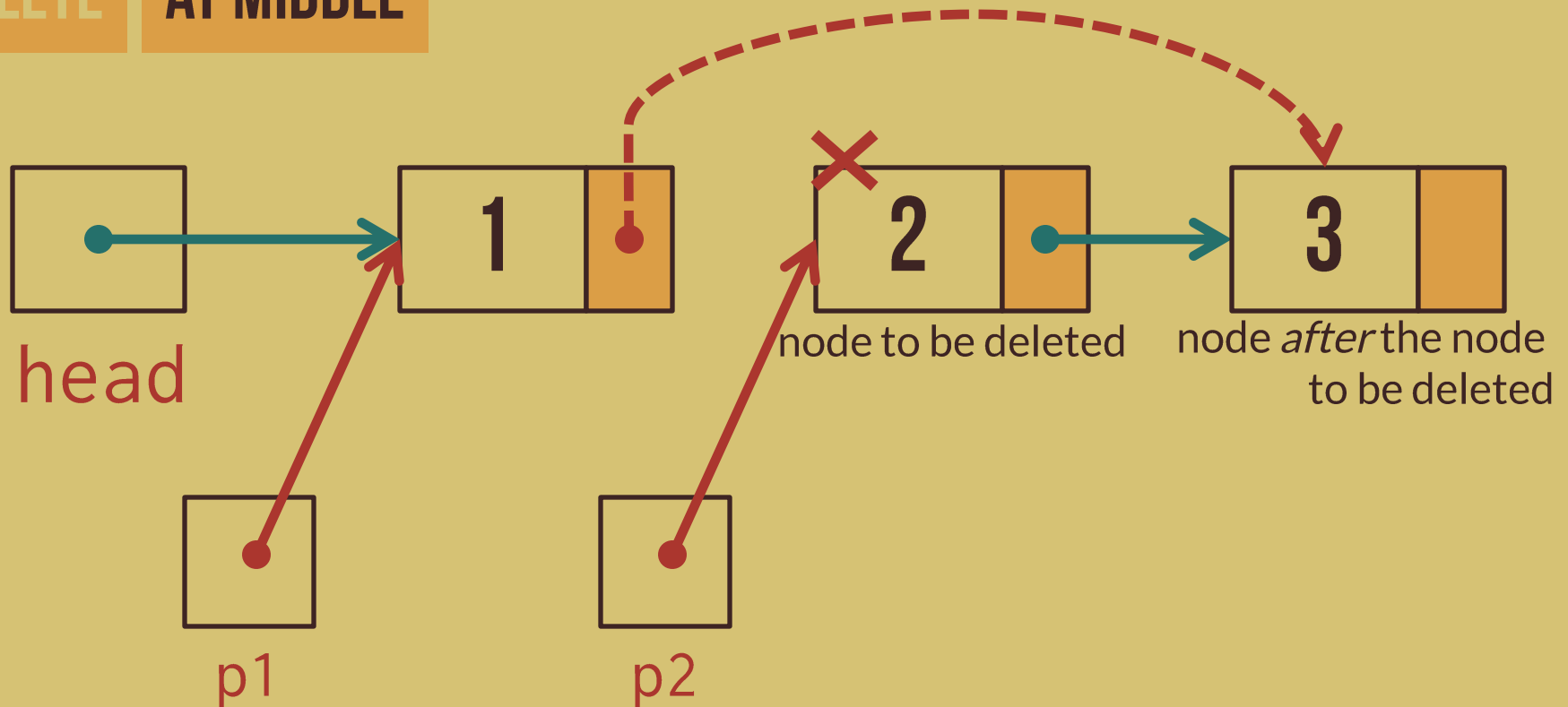
AT MIDDLE



2. Have another pointer (**p2**) refer to the node to be deleted.

DELETE

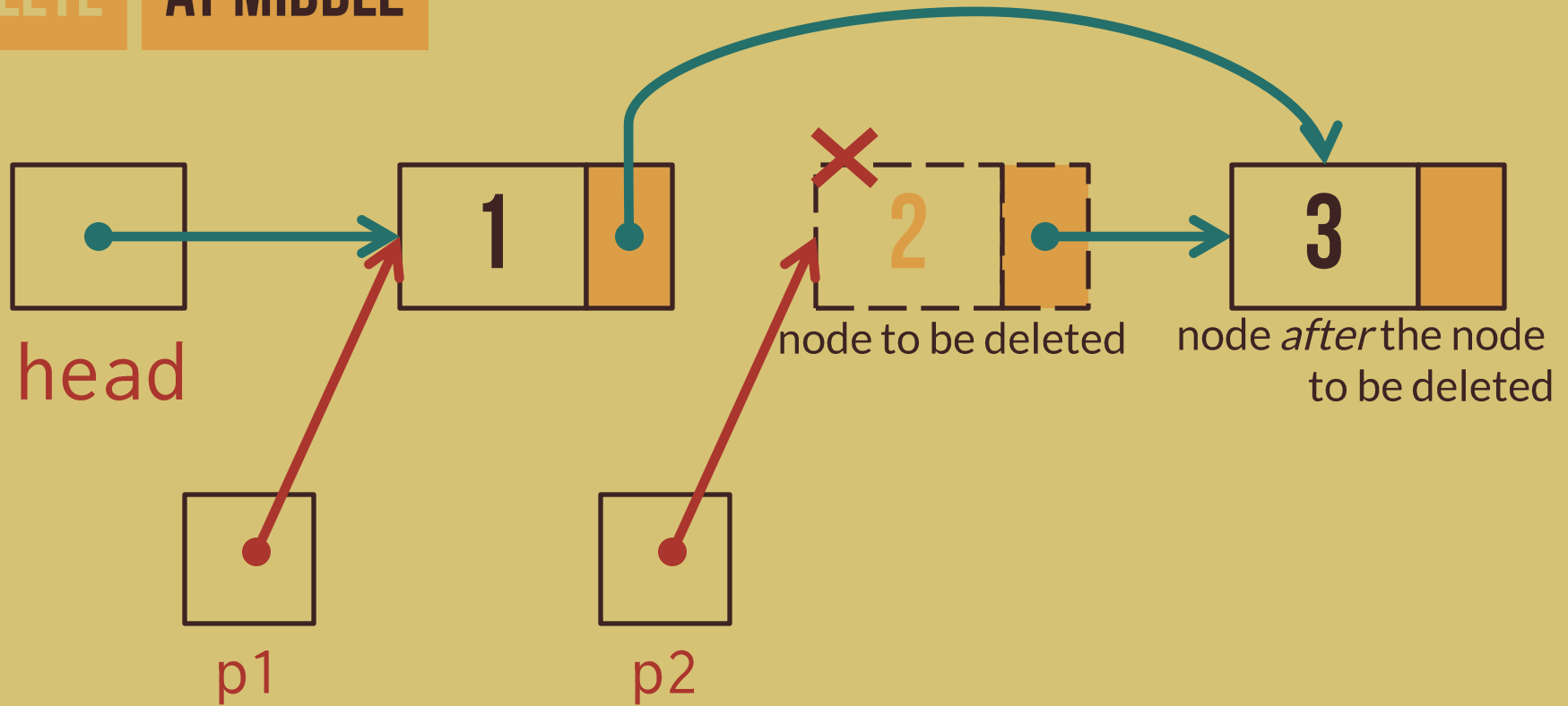
AT MIDDLE



3. Make the next pointer of the node pointed by **p1** refer to the node *after* the node to be deleted (using **p2**).

DELETE

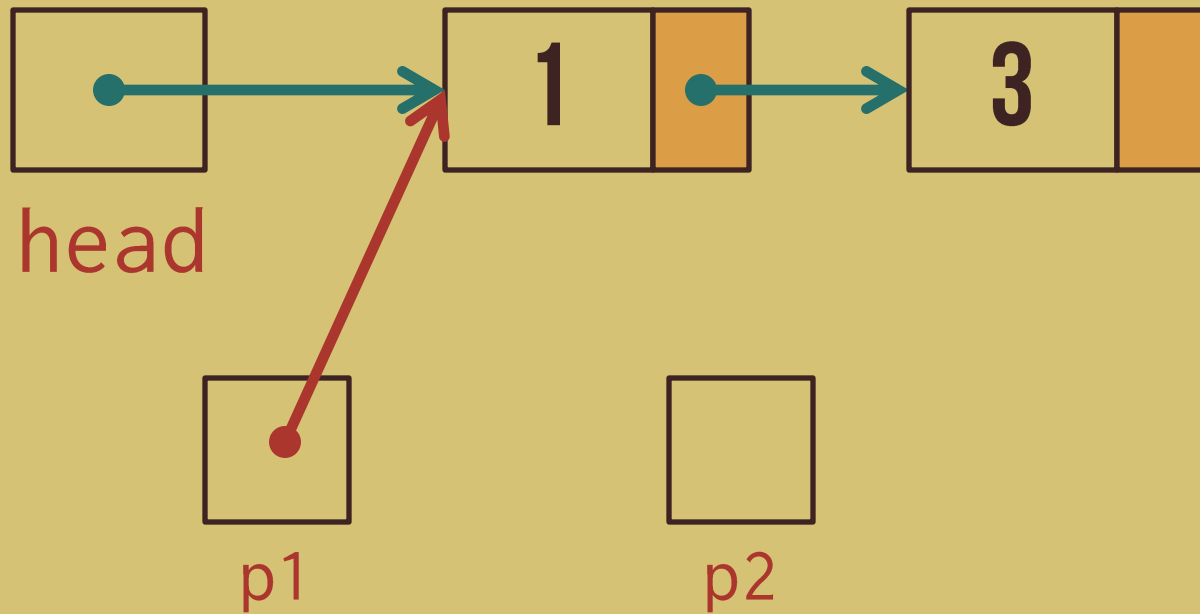
AT MIDDLE



4. Free the node to be deleted using **p2**.

DELETE

AT MIDDLE

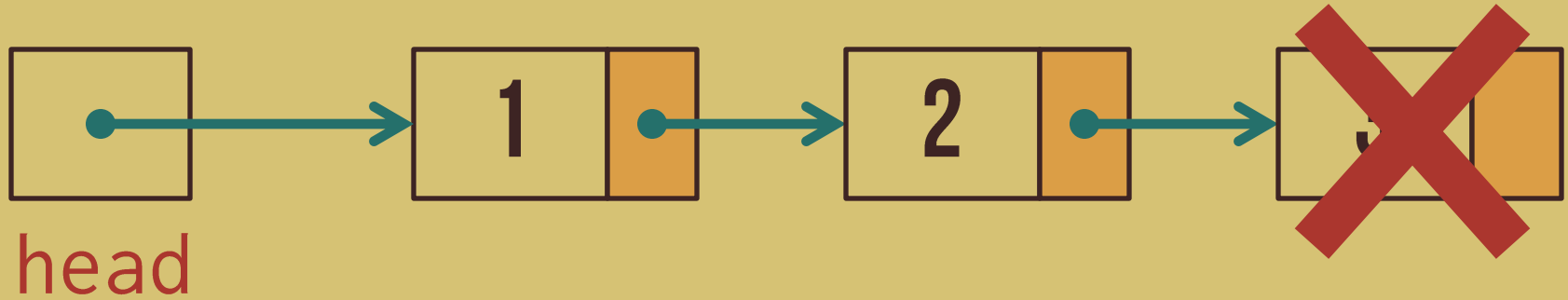


DELETE

AT TAIL

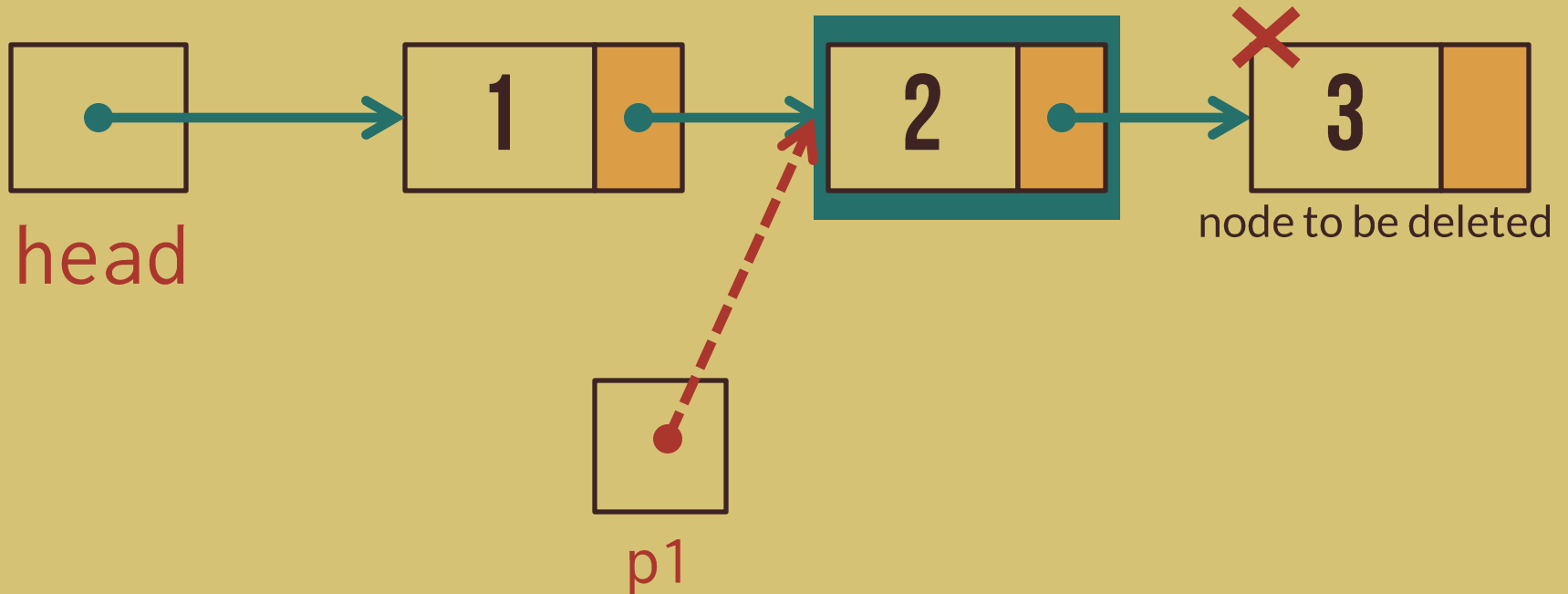
DELETE

AT TAIL



DELETE

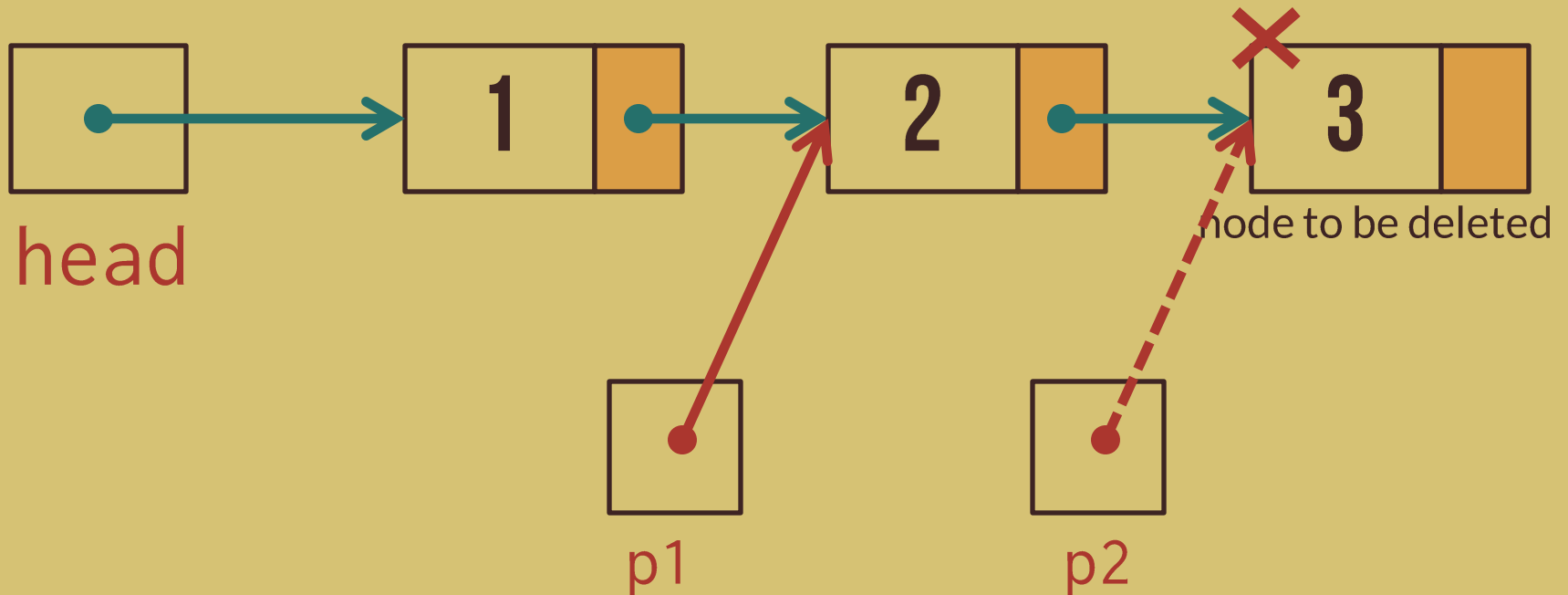
AT TAIL



1. Find the node *before* the tail using a pointer (p1).

DELETE

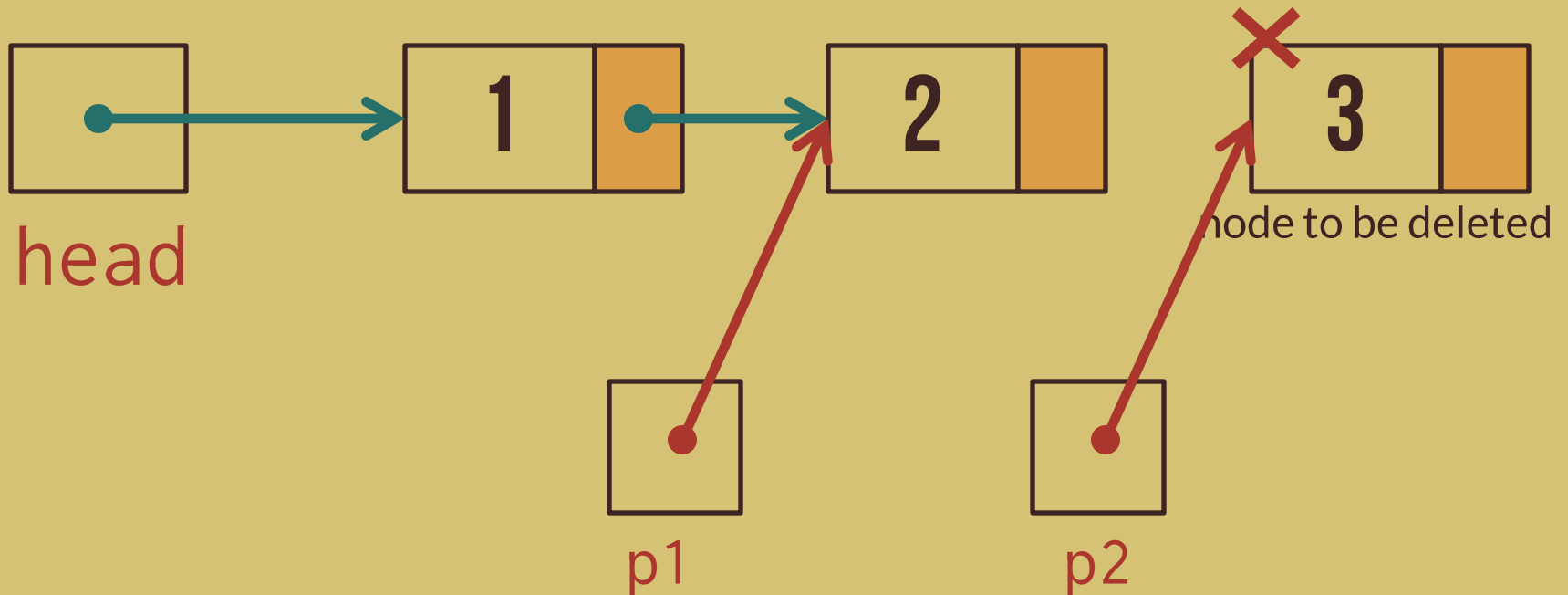
AT TAIL



2. Have another pointer (**p2**) refer to the last node.

DELETE

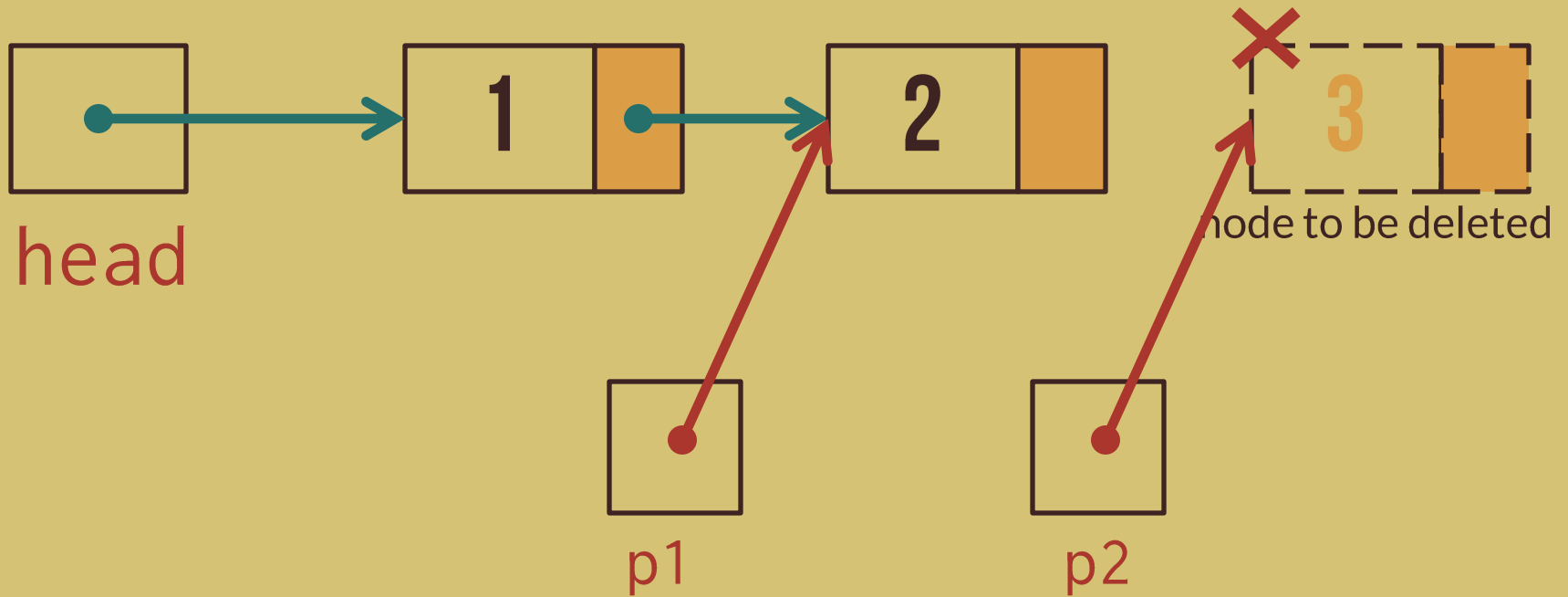
AT TAIL



3. Make the next pointer of the node referred to by **p1** to NULL.

DELETE

AT TAIL



4. Free the node pointed by **p2**.

DELETE

AT TAIL

