

BLG 312E - Operating Systems

Homework 3

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1 1. Semaphore Fundamentals

1.1 (a)

Semaphore is counter. It show how many can go in. If value is 0, wait. If more than 0, pass.

1.2 (b)

- -1 is wrong. Not work. - 0 means block all. - 1 is mutex. - >1 means many thread can go.

1.3 (c)

We need queue to save who is waiting. Like ticket line.

1.4 (d)

Mutex is same like semaphore with 1 value. We use macros:

```
#define mutex_init(sem) sem_init(sem, 1)
#define mutex_lock(sem) sem_wait(sem)
#define mutex_unlock(sem) sem_post(sem)
#define mutex_destroy(sem) sem_destroy(sem)
```

2 2. Thread Scheduling and ucontext.h

2.1 (a)

I use `ucontext.h` to switch thread. When I do `swapcontext`, it go to another thread context.

2.2 (b)

Timer call signal, then call `thread_yield` or `thread_schedule`. Then it switch to next thread.

2.3 (c)

I make two queue: high and low. If high is not empty, pick from there. Else pick from low.

2.4 (d)

I check name. If name is "Player1" or "Player3", I put thread in high queue. So even if create later, they go first.

3 3. Synchronization

3.1 (a)

If two thread inside mutex and one use global var wrong, it still break. Mutex not fix logic bug.

3.2 (b)

Use less shared data. Also put all shared work in one mutex block. Use atomic if can.

4 4. File Descriptor and Communication

4.1 Why everything is file?

Because it is simple. Read, write, open all work same. We can use same code for socket, file, etc.

4.2 write vs dprintf

- `write` is low level. No format. - `dprintf` can format like `printf`.

4.3 send vs dprintf

- `send` is for socket. It can set flags. - `dprintf` can send string, but no flag.

4.4 recv vs read

- `recv` is like `read` but for socket. - `recv` can do more with flags.

4.5 Why no `dscanf`?

Because it hard to know when input finish. `scanf` need stream, socket is not same.

5 5. My Implementations

5.1 queue.c

queue_init:

```
void queue_init(Queue* q) {
    q->front = NULL;
    q->rear = NULL;
}
```

queue_enqueue:

```
void queue_enqueue(Queue* q, void* data) {
    QueueNode* node = malloc(sizeof(QueueNode));
    node->data = data;
    node->next = NULL;
    if (q->rear)
        q->rear->next = node;
    else
        q->front = node;
    q->rear = node;
}
```

queue_dequeue:

```
void* queue_dequeue(Queue* q) {
    if (!q->front) return NULL;
    QueueNode* temp = q->front;
    void* data = temp->data;
    q->front = temp->next;
    if (!q->front) q->rear = NULL;
    free(temp);
    return data;
}
```

queue_is_empty:

```
int queue_is_empty(Queue* q) {
    return q->front == NULL;
}
```

queue_remove:

```
void* queue_remove(Queue* q, void* data) {
    QueueNode *prev = NULL, *curr = q->front;
    while (curr) {
        if (curr->data == data) {
            if (prev) prev->next = curr->next;
            else q->front = curr->next;
            if (q->rear == curr) q->rear = prev;
            void* d = curr->data;
            free(curr);
            return d;
        }
        prev = curr;
        curr = curr->next;
    }
    return NULL;
}
```

queue_destroy:

```
void queue_destroy(Queue* q) {
    while (!queue_is_empty(q)) {
        queue_dequeue(q);
    }
}
```

5.2 semaphore.c

sem_init:

```
void sem_init(Semaphore* sem, int value) {
    sem->count = value;
    queue_init(&sem->waiting_queue);
}
```

sem_wait:

```
void sem_wait(Semaphore* sem) {
    sem->count--;
    if (sem->count < 0) {
        queue_enqueue(&sem->waiting_queue, current_thread);
        thread_block();
    }
}
```

sem_post:

```
void sem_post(Semaphore* sem) {
    sem->count++;
    if (sem->count <= 0) {
        Thread* t = queue_dequeue(&sem->waiting_queue);
        thread_unblock(t);
    }
}
```

sem_destroy:

```
void sem_destroy(Semaphore* sem) {
    queue_destroy(&sem->waiting_queue);
}
```

5.3 thread.c (priority)

I add 2 queues:

```
Queue high_priority_queue;
Queue low_priority_queue;
```

In thread_create:

```
if (strstr(thread->name, "Player4") || strstr(thread->name, "Player2"))
    queue_enqueue(&high_priority_queue, thread);
else
    queue_enqueue(&low_priority_queue, thread);
```

In thread_schedule:

```
if (!queue_is_empty(&high_priority_queue))
    next = queue_dequeue(&high_priority_queue);
else
    next = queue_dequeue(&low_priority_queue);
```

6 References

- <https://man7.org/linux/man-pages/man3/ucontext.3.html>
- <https://www.youtube.com/watch?v=mb6U5z4G0to>
- <https://www.geeksforgeeks.org/semaphores-in-process-synchronization/>
- My lecture notes from BLG 312E (ITU)
- <https://www.man7.org/linux/man-pages/>