

操作系统研讨课

Course: B0911011Y

陈明宇、蒋德钧

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email: cmy@ict.ac.cn / jiangdejun@ict.ac.cn

office phone: 62601007



Lecture 3 Interactive OS and Process Management

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Schedule

- Project 3 assignment
- Project 2 due



Project 3 Interactive OS and Process Management

- Requirement I
 - Support interactive operation
 - Implement a simple shell terminal and support receiving and showing user commands
 - Implement a few shell commands, e.g. ps, clear



Project 3 Interactive OS and Process Management

- Requirement II
 - Support basic process management
 - Implement four system calls and the related shell commands
 - sys_spawn and its shell command: exec
 - sys_kill and its shell command: kill
 - sys_wait
 - sys_exit
 - Implement synchronization primitives: semaphores and barrier



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- Requirement III
 - Support inter-process communication
 - Implement synchronization primitives
 - semaphore
 - barrier
 - Implement inter-process communication mechanism
 - mailbox



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- Requirement IV
 - Support running processes on dual cores
 - Run dual cores
 - Support synchronization primitives、mailbox、binding core under dual cores



Project 3 Interactive OS and Process Management

- Simple terminal
 - Screen
 - Use the provided printf to show input command
 - Shell
 - A user level process: test_shell.c
 - Parse input command and invoke corresponding syscalls
 - Show the input command



Project 3 Interactive OS and Process Management

- Simple terminal

```
> [TASK] I am task with pid 2, I have acquired two mutex lock. (26)
> [TASK] I want to acquire a mute lock from task(pid=2).
> [TASK] I want to wait task (pid=2) to exit.

----- COMMAND -----
> root@UCAS_OS: ps
[PROCESS TABLE]
[0] PID : 0   STATUS : RUNNING
[1] PID : 1   STATUS : RUNNING
> root@UCAS_OS: spawn
> root@UCAS_OS:
```



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- Simple terminal
 - Note that shell runs immediately after the kernel starts and acts as PID 1
 - In this project, shell polls the serial port instead of using interrupt
 - Please read the start code (test_shell.c) we provide to you



Project 3 Interactive OS and Process Management

- Basic process management
 - Spawn
 - starts a new process with a new process ID
 - Run the program specified in `sys_spawn'` s arguments
 - Exit
 - Finish the running of a process in a normal way, and release **all its resources (e.g. lock, waiting processes)**



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- Process management
 - Wait
 - Waits on a process to complete its execution or to be killed
 - Kill
 - Sends signals to running processes to request the termination of the process, and **all its resources (e.g. lock, waiting processes)**



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- Implementing spawn
 - Shell command: `exec id`
 - Given an array of tasks
 - Spawn the process with the number *id* corresponding to the array subscript
 - `sys_spawn(task_info_t *info, void* arg, spawn_mode_t mode)`
 - A syscall to start a new process
 - Initialize the PCB and put the process into the ready queue



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- Implementing wait and exit
 - `sys_waitpid(pid_t)`
 - A syscall to wait on a process to terminate
 - Put the process into the corresponding wait queue
 - Note that, *pid_t* is hard coded in the task. If it is changed, pls. modify the task code
 - `sys_exit(void)`
 - Normally finish the running of the process
 - Reclaim all its resources



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- Implementing kill
 - Shell command: kill pid
 - Kill the process with the corresponding *pid*
 - sys_kill(pid_t)
 - A syscall to kill a process immediately no matter which queue it is in
 - Reclaim resources, such as PCB, stacks, and lock



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- Implement basic process operations

```
> [TASK] I am task with pid 2, I have acquired two mutex lock. (144)
> [TASK] I want to acquire a mute lock from task(pid=2).
> [TASK] I want to wait task (pid=2) to exit.
```

```
----- COMMAND -----
> root@UCAS_OS: exec 0
exec process[0].
> root@UCAS_OS: exec 1
exec process[1].
> root@UCAS_OS: exec 2
exec process[2].
```



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- You are encouraged to enrich your own shell to handle more user commands



Project 3 Interactive OS and Process Management

- Synchronization – semaphore
 - Control access to a shared resource
 - A value keeps track of the number of resources that are currently available
 - A queue holding waiting tasks
 - Main operations
 - Down: decrement value and block the process if the decremented value is less than zero
 - Up: increment value and unblock one waiting process



Project 3 Interactive OS and Process Management

- Synchronization – barriers
 - A barrier for a group of tasks is a location in code where any task must stop at this point and cannot proceed until all other tasks reach this barrier
 - Keep track of the number of tasks at barrier
 - Maintain queue holding waiting tasks
 - Main operations
 - Wait: block the task if not all the tasks have reached the barrier. Otherwise, unblock all



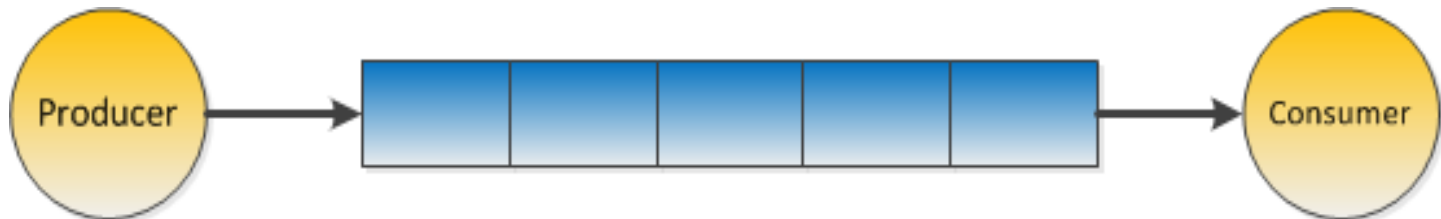
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- Synchronization
 - Note that
 - Pls. read the guide book and refer to the test case to see how these primitives are tested
 - These primitives are implemented in the kernel, and provide syscalls to user-level process
 - Pay attention to the impact of interrupt on implementing these primitives



Project 3 Interactive OS and Process Management

- IPC – Mailbox
 - Bounded buffer
 - Fixed size
 - FIFO
 - (Multiple) producers: put data into the buffer
 - (Multiple) consumers: remove data from the buffer



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- IPC – Mailbox
 - Producer-consumer problem
 - Two processes (producer and consumer) share a common fixed-size buffer used as a queue
 - The producer will not try to add data into the buffer if it is full
 - The consumer will not try to remove data from the buffer if it is empty



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- IPC – Mailbox
 - How to deal with producer-consumer problem?
 - Producer blocks if the buffer is full
 - Consumer blocks if the buffer is empty
 - How to notify the other part if the condition is satisfied?
 - Use your implemented synchronization primitives



Project 3 Interactive OS and Process Management

- Running processes on dual cores
 - The two cores share the same memory
 - Each core has its own set of registers
 - How to start the second core?
 - By default, both cores start to run initially
 - The second core continues to loop with bbl



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- Start dual cores
 - After setup, at bbl use loadbootm instruction to start both cores
 - a0 register or mhartid register holds core id
 - You can use core 0 as the main core
 - Use `sbi_send_ipi(const unsigned long *hart_mask)` to send inter-process interrupt to the second core
 - Use the mhartid register value to decide which code path to run in your kernel



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- Access shared variables in kernel under dual cores
 - Shared variables in kernel
 - e.g. ready queue, variables in synchronization primitives
 - Use a big lock to protect the whole kernel when entering the kernel
 - Use atomic operations to implement lock
 - Pls. refer to `arch/riscv/include/atomic.h`
 - You need to implement some kernel variables separately for each core, e.g. `current_running`



Project 3 Interactive OS and Process Management

- Test dual-core
 - Accelerating adding operations with dual cores
 - Run adding with single core
 - Run adding with dual cores
 - Acceleration is expected when running dual cores
 - Tips
 - Use a relatively large clock slice



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- Test dual-core
 - Support mailbox with dual core
 - Three processes can both send and receive mails from mailbox
 - A process blocks when it fails to send or receive mail
 - Pls. do not let deadlock occur



Project 3 Interactive OS and Process Management

- Test dual-core
 - Implement shell command: taskset
 - taskset cpumask threadID
 - taskset 1 0: run thread 0 on core 0
 - taskset -p cpumask threadID
 - taskset -p 1 4: run thread 4 on core 1
 - Please read test_affinity.c to know how to test
 - A process has 5 sub-tasks, each task by default run on the same core with the parent process



Project 3 Interactive OS and Process Management

- Step by step
 - Task 1 (S-Core)
 - Implement shell process to support user command *ps* and *clear*
 - At least, *ps* shows two process (PID 0 and 1)
 - Implement user command *exec* to invoke `sys_spawn` to run new process
 - Implement `sys_exit` and `sys_wait`
 - Implement user command *kill* to terminate a running process



Project 4 Synchronization Primitives and IPC

- Step by step
 - Task 2
 - Implement two primitives: semaphore and barrier. Verify them use the test cases.
 - Implement mailbox
 - S-Core: only need to implement barrier
 - A-Core: implement both primitives as well as mailbox



Project 4 Synchronization Primitives and IPC

- Step by step
 - Task 3 (A-Core)
 - Run dual-cores with the adding test case
 - Task 4 (C-Core)
 - Run mailbox under dual core
 - Task 5 (C-Core)
 - Support using taskset to bind process on specific core



Project 3 Interactive OS and Process Management

- Requirements for design review
 - Which commands can be supported or will be supported by your shell?
 - Show example code about spawn, kill, wait, and exit?
 - How do you handle the case when killing a task meanwhile it holds a lock?



Project 3 Interactive OS and Process Management

- Requirements for design review
 - How do you handle semaphores, and barrier? What to do if timer interrupt occurs?
 - Show the structure for mailbox. How do you protect concurrent accessing for mailbox?
 - How do you enable two CPU cores?



Project 3 Interactive OS and Process Management

- Requirements for design review
 - Any question about C-Core



Project 3 Interactive OS and Process Management

- P3 schedule
 - 25th Oct.: No class
 - P3 design review: 1st Nov.
 - P3 due: 8th Nov.

