

Concurrency

Motivation

- Operating systems (and application programs) often need to be able to handle multiple things happening at the same time
 - Process execution, interrupts, background tasks, system maintenance
- Humans are not very good at keeping track of multiple things happening simultaneously
- Threads are an abstraction to help bridge this gap

Why Concurrency?

- Servers
 - Multiple connections handled simultaneously
- Parallel programs
 - To achieve better performance
- Programs with user interfaces
 - To achieve user responsiveness while doing computation
- Network and disk bound programs
 - To hide network/disk latency

Definitions

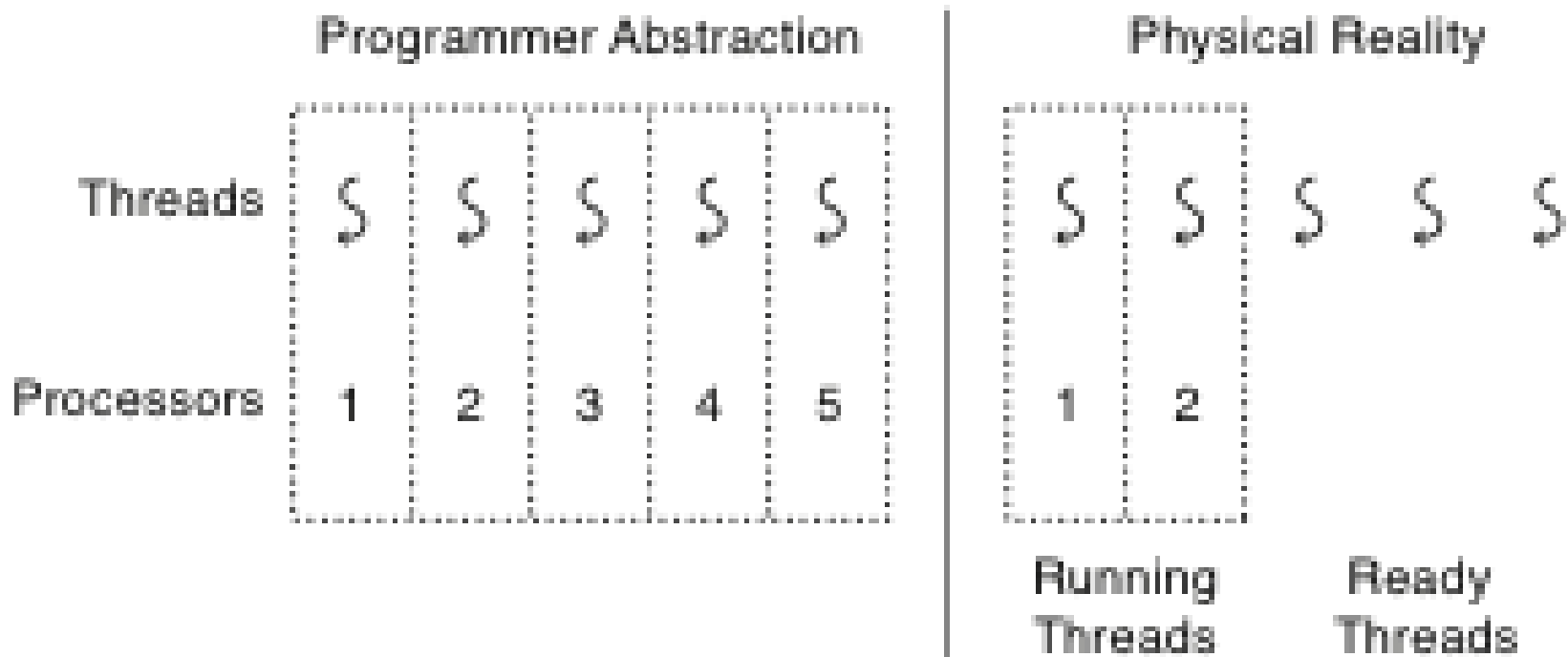
- A thread is a single execution sequence that represents a separately schedulable task
 - Single execution sequence: familiar programming model
 - Separately schedulable: OS can run or suspend a thread at any time
- Protection is an orthogonal concept
 - Can have one or many threads per protection domain

Threads in the Kernel and at User-Level

- Multi-threaded kernel
 - multiple threads, sharing kernel data structures, capable of using privileged instructions
- Multiprocess kernel
 - Multiple single-threaded processes
 - System calls access shared kernel data structures
- Multiple multi-threaded user processes
 - Each with multiple threads, sharing same data structures, isolated from other user processes

Thread Abstraction

- Infinite number of processors
- Threads execute with variable speed
 - Programs must be designed to work with any schedule



Programmer vs. Processor View

Programmer's View

.
.
.
x = x + 1;
y = y * x;
z = x + 5y;
.
.
.

Possible Execution #1

.
.
.
x = x + 1;
y = y * x;
z = x + 5y;
.
.
.

Possible Execution #2

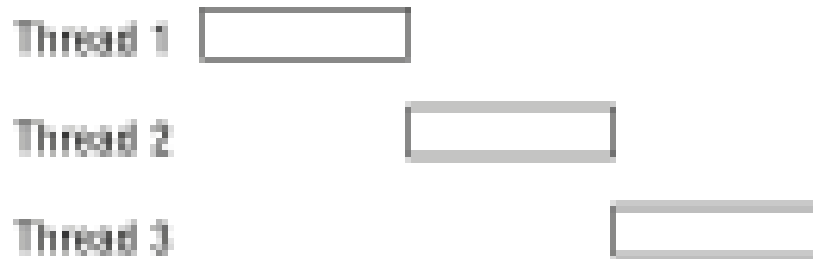
.
.
.
x = x + 1;
.....
Thread is suspended.
Other thread(s) run.
Thread is resumed.
.....
y = y * x;
z = x + 5y;

Possible Execution #3

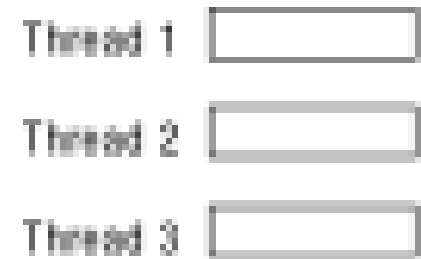
.
.
.
x = x + 1;
y = y * x;
.....
Thread is suspended.
Other thread(s) run.
Thread is resumed.
.....
z = x + 5y;

Possible Executions

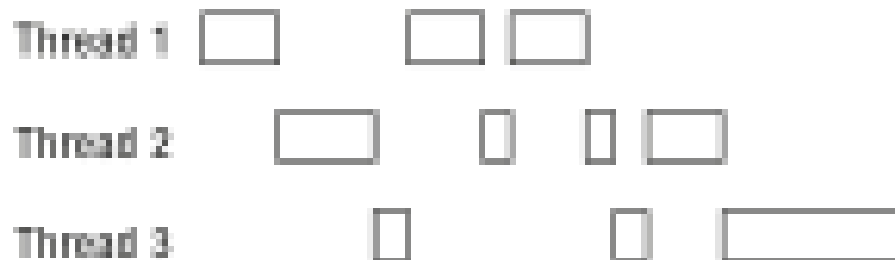
One Execution



Another Execution



Another Execution



Thread Operations

- `thread_create(thread, func, args)`
 - Create a new thread to run `func(args)`
- `thread_yield()`
 - Relinquish processor voluntarily
- `thread_join(thread)`
 - In parent, wait for forked thread to exit, then return
- `thread_exit`
 - Quit thread and clean up, wake up joiner if any

Example: threadHello

```
#define NTHREADS 10
thread_t threads[NTHREADS];
main() {
    for (i = 0; i < NTHREADS; i++)
        thread_create(&threads[i], &go, i);
    for (i = 0; i < NTHREADS; i++) {
        exitValue = thread_join(threads[i]);
        printf("Thread %d returned with %ld\n", i, exitValue);
    }
    printf("Main thread done.\n");
}

void go (int n) {
    printf("Hello from thread %d\n", n);
    thread_exit(100 + n);
    // REACHED?
}
```

threadHello: Example Output

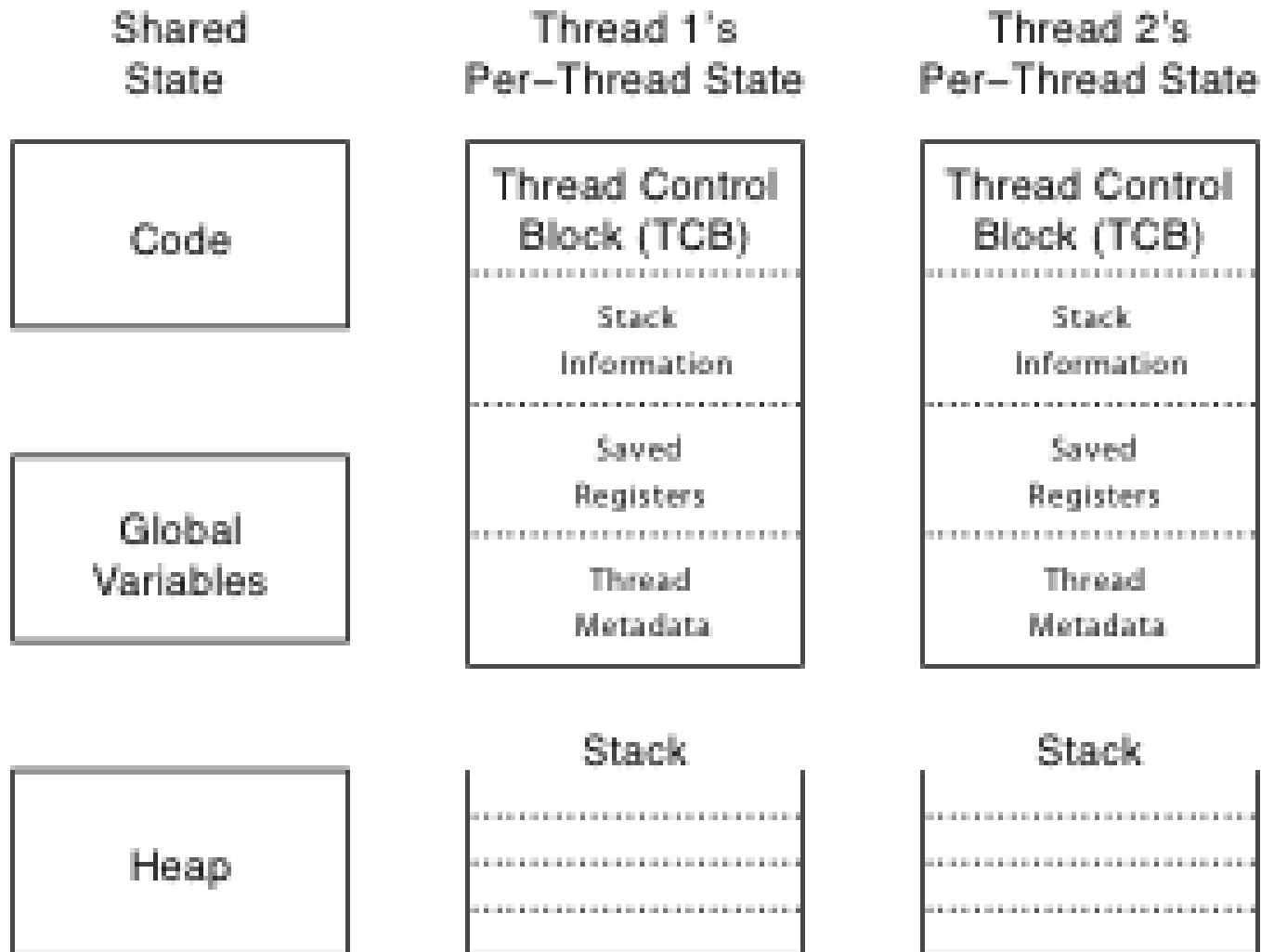
- Why must “thread returned” print in order?
- What is maximum # of threads running when thread 5 prints hello?
- Minimum?

```
bash-3.2$ ./threadHello
Hello from thread 0
Hello from thread 1
Thread 0 returned 100
Hello from thread 3
Hello from thread 4
Thread 1 returned 101
Hello from thread 5
Hello from thread 2
Hello from thread 6
Hello from thread 8
Hello from thread 7
Hello from thread 9
Thread 2 returned 102
Thread 3 returned 103
Thread 4 returned 104
Thread 5 returned 105
Thread 6 returned 106
Thread 7 returned 107
Thread 8 returned 108
Thread 9 returned 109
Main thread done.
```

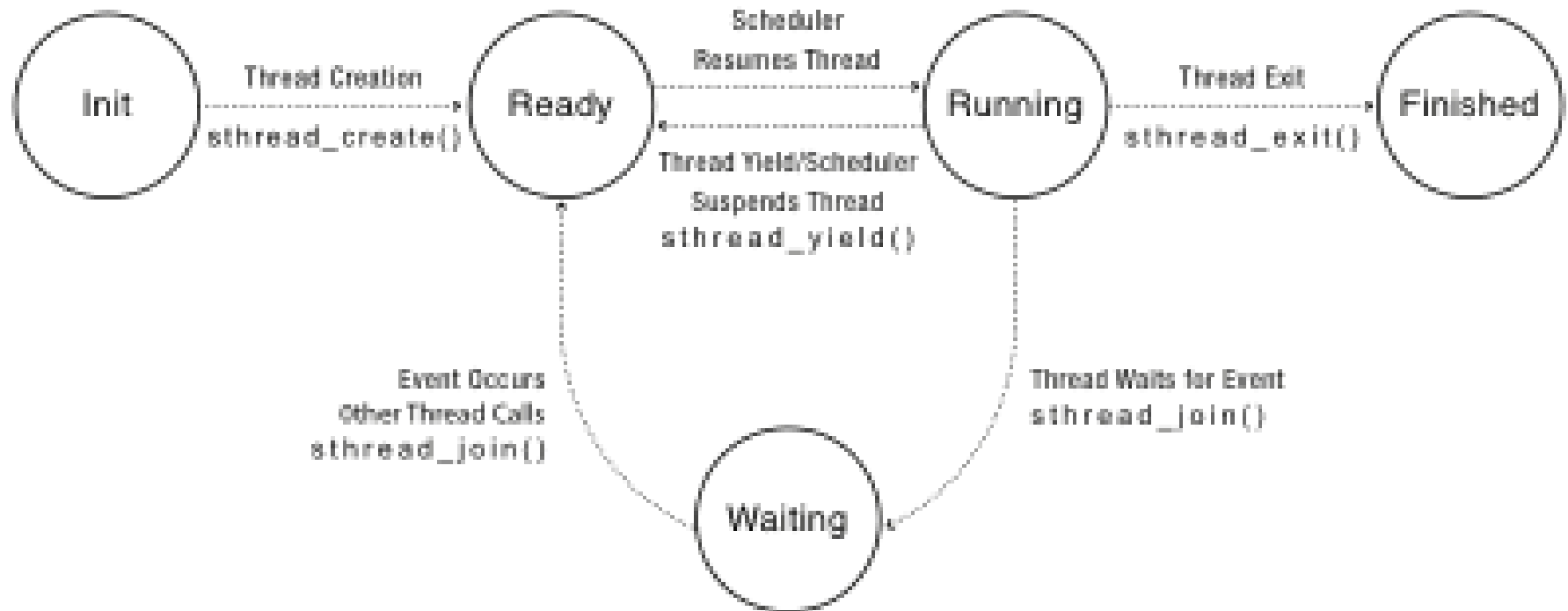
Fork/Join Concurrency

- Threads can create children, and wait for their completion
- Data only shared before fork/after join
- Examples:
 - Web server: fork a new thread for every new connection
 - As long as the threads are completely independent
 - Merge sort
 - Parallel memory copy

Thread Data Structures



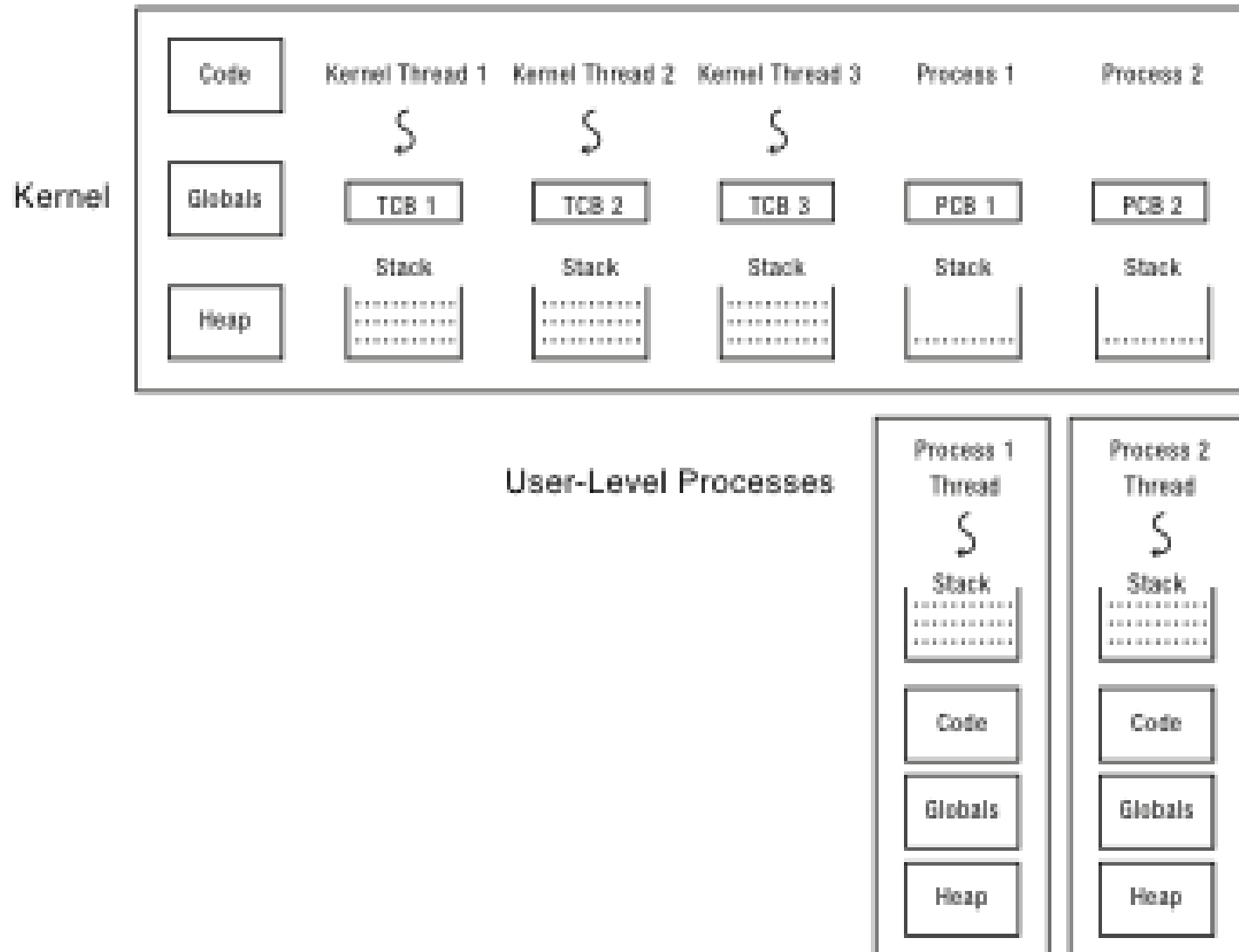
Thread Lifecycle



Implementing Threads: Roadmap

- Kernel threads
 - Thread abstraction only available to kernel
 - To the kernel, a kernel thread and a single threaded user process look quite similar
- Multithreaded processes using kernel threads (Linux, MacOS)
 - Kernel thread operations available via syscall
- User-level threads
 - Thread operations without system calls

Multithreaded OS Kernel



Implementing threads

- Thread_fork(func, args)
 - Allocate thread control block
 - Allocate stack
 - Build stack frame for base of stack (stub)
 - Put func, args on stack
 - Put thread on ready list
 - Will run sometime later (maybe right away!)
- stub(func, args): OS/161 mips_threadstart
 - Call (*func)(args)
 - If return, call thread_exit()

Thread Stack

- What if a thread puts too many procedures on its stack?
 - What happens in Java?
 - What happens in the Linux kernel?
 - What happens in OS/161?
 - What *should* happen?

Thread Context Switch

- Voluntary
 - Thread_yield
 - Thread_join (if child is not done yet)
- Involuntary
 - Interrupt or exception
 - Some other thread is higher priority

Voluntary thread context switch

- Save registers on old stack
- Switch to new stack, new thread
- Restore registers from new stack
- Return
- Exactly the same with kernel threads or user threads

OS/161 switchframe_switch

```
/* a0: old thread stack pointer          /* Get new stack pointer from new
 * a1: new thread stack pointer */      thread */
                                        lw    sp, 0(a1)
/* Allocate stack space for 10          nop                    /* delay slot for load
 registers. */                          */
addi sp, sp, -40

/* Save the registers */
sw    ra, 36(sp)
sw    gp, 32(sp)
sw    s8, 28(sp)
sw    s6, 24(sp)
sw    s5, 20(sp)
sw    s4, 16(sp)
sw    s3, 12(sp)
sw    s2, 8(sp)
sw    s1, 4(sp)
sw    s0, 0(sp)

/* Store old stack pointer in          /* Now, restore the registers */
old thread */
sw    sp, 0(a0)

                                        lw    s0, 0(sp)
                                        lw    s1, 4(sp)
                                        lw    s2, 8(sp)
                                        lw    s3, 12(sp)
                                        lw    s4, 16(sp)
                                        lw    s5, 20(sp)
                                        lw    s6, 24(sp)
                                        lw    s8, 28(sp)
                                        lw    gp, 32(sp)
                                        lw    ra, 36(sp)
                                        nop                    /* delay slot
                                        for load */

/* and return. */
j    ra
addi sp, sp, 40          /* in delay
slot */
```

x86 switch_threads

```
# Save caller's register
state
# NOTE: %eax, etc. are
ephemeral
pushl %ebx
pushl %ebp
pushl %esi
pushl %edi

# Get offsetof (struct
thread, stack)
mov thread_stack_ofs, %edx
# Save current stack pointer
to old thread's stack, if
any.
movl SWITCH_CUR(%esp), %eax
movl %esp, (%eax,%edx,1)

# Change stack pointer to
new thread's stack
# this also changes
currentThread
movl SWITCH_NEXT(%esp), %ecx
movl (%ecx,%edx,1), %esp

# Restore caller's register
state.
popl %edi
popl %esi
popl %ebp
popl %ebx
ret
```

A Subtlety

- Thread_create puts new thread on ready list
- When it first runs, some thread calls switchframe
 - Saves old thread state to stack
 - Restores new thread state from stack
- Set up new thread's stack as if it had saved its state in switchframe
 - “returns” to stub at base of stack to run func

Two Threads Call Yield

Thread 1's instructions

"return" from thread_switch
into stub
call go
call thread_yield
choose another thread
call thread_switch
save thread 1 state to TCB
load thread 2 state

return from thread_switch
return from thread_yield
call thread_yield
choose another thread
call thread_switch

Thread 2's instructions

"return" from thread_switch
into stub
call go
call thread_yield
choose another thread
call thread_switch
save thread 2 state to TCB
load thread 1 state

Processor's instructions

"return" from thread_switch
into stub
call go
call thread_yield
choose another thread
call thread_switch
save thread 1 state to TCB
load thread 2 state
"return" from thread_switch
into stub
call go
call thread_yield
choose another thread
call thread_switch
save thread 2 state to TCB
load thread 1 state
return from thread_switch
return from thread_yield
call thread_yield
choose another thread
call thread_switch

Involuntary Thread/Process Switch

- Timer or I/O interrupt
 - Tells OS some other thread should run
- Simple version (OS/161)
 - End of interrupt handler calls `switch()`
 - When resumed, return from handler resumes kernel thread or user process
 - Thus, processor context is saved/restored twice (once by interrupt handler, once by thread switch)

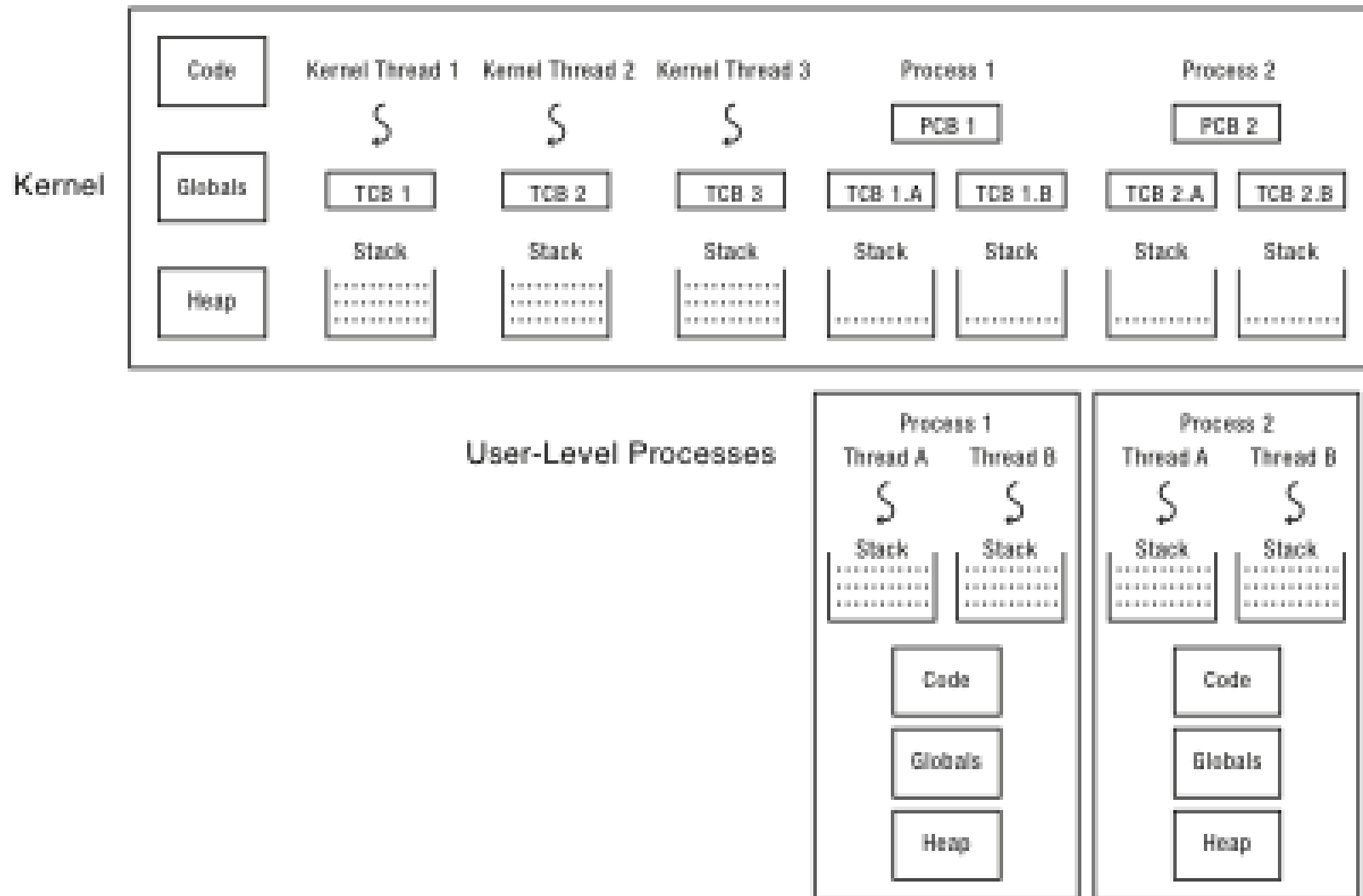
Faster Thread/Process Switch

- What happens on a timer (or other) interrupt?
 - Interrupt handler saves state of interrupted thread
 - Decides to run a new thread
 - Throw away current state of interrupt handler!
 - Instead, set saved stack pointer to trapframe
 - Restore state of new thread
 - On resume, pops trapframe to restore interrupted thread

Multithreaded User Processes (Take 1)

- User thread = kernel thread (Linux, MacOS)
 - System calls for thread fork, join, exit (and lock, unlock,...)
 - Kernel does context switch
 - Simple, but a lot of transitions between user and kernel mode

Multithreaded User Processes (Take 1)



Multithreaded User Processes (Take 2)

- Green threads (early Java)
 - User-level library, within a single-threaded process
 - Library does thread context switch
 - Preemption via upcall/UNIX signal on timer interrupt
 - Use multiple processes for parallelism
 - Shared memory region mapped into each process

Multithreaded User Processes (Take 3)

- Scheduler activations (Windows 8)
 - Kernel allocates processors to user-level library
 - Thread library implements context switch
 - Thread library decides what thread to run next
- Upcall whenever kernel needs a user-level scheduling decision
 - Process assigned a new processor
 - Processor removed from process
 - System call blocks in kernel