



Applied Electronics

C5 – Bus Protocols

- Addressing
- Arbitrage
- Performance parameters
- Source sync transfers
- DDR transfers



Lecture C5: bus protocols

- Addressing techniques (single-Master systems)
- Channel allocation (multi-Master systems)
 - ◆ Basic allocation mechanisms and techniques
 - ◆ Priority, starvation, fairness
- Analysis and improvement of the performance of a bus
 - ◆ Source-Synchronous Transfers
 - ◆ Double Data Rate (DDR) Transfers
- References:
 - ◆ D. Del Corso: Telecommunication Electronics: Ch. 5.4



Services provided by the cycle level

- Transfer of information units (byte, word, ...), respecting the time specifications of the destination register
 - ◆ Independent of the information TYPE (address, data, ...)
- There is **one source** of information and **one destination**
 - ◆ Extended protocols support asynchronous 1-to-N and N-to-M transfers (N-partner protocols)
- Cycle level operations can be **read** or **write**
- They need **power** and **time** to provide the services.

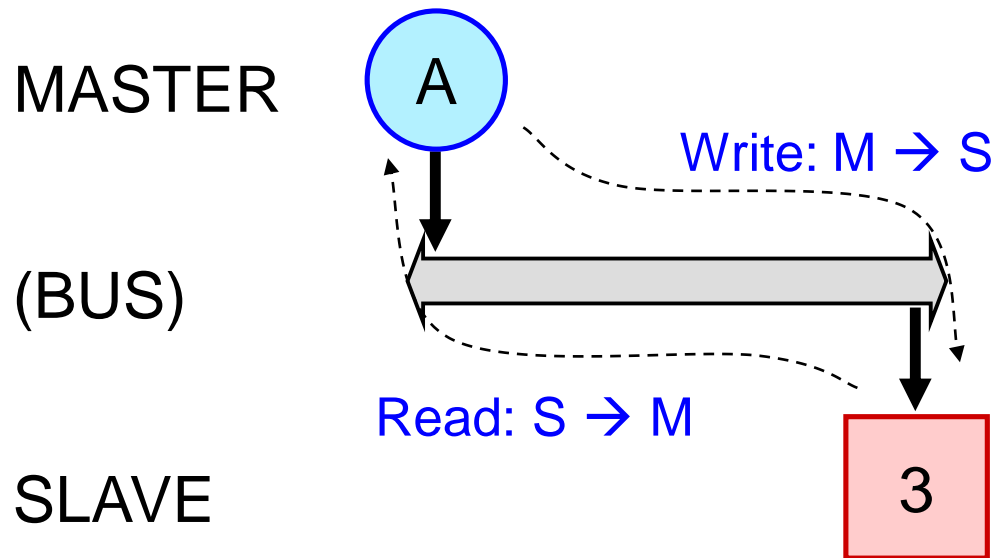
Definition of transaction

- Sequence of **one or more cycles**, which transfer elements of information with a **definite meaning** (addresses, data, priority vector, ...)
- Transaction-level items
 - ◆ **MASTER:** activates the operations
 - ◆ **SLAVE:** responds to Master commands



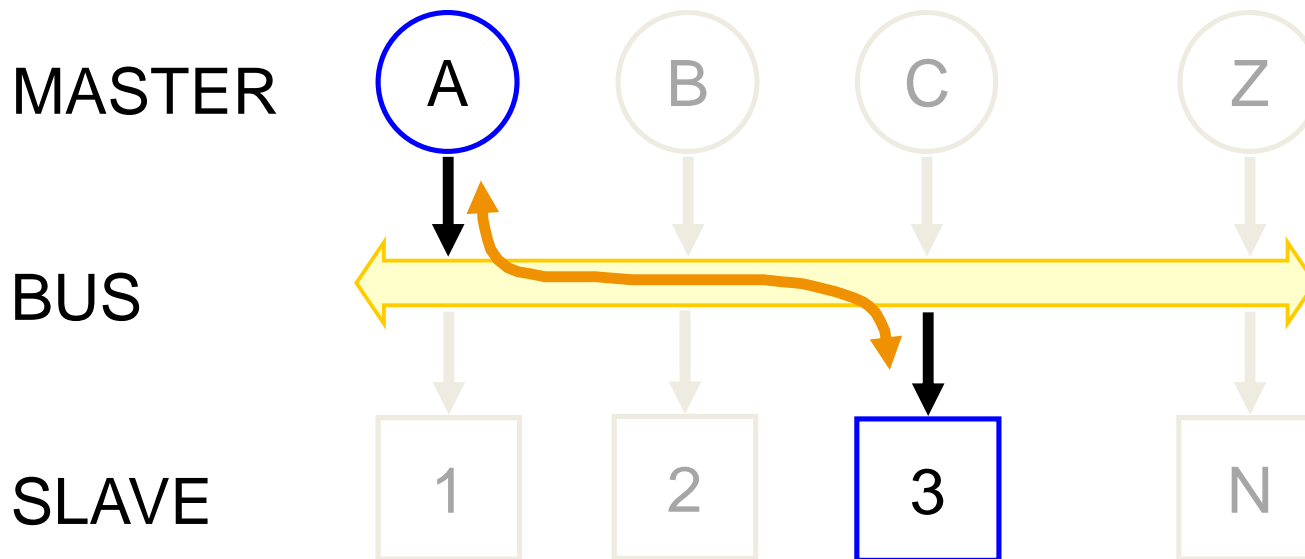
Point-to-point systems

- Transfer **point-to-point**
 - ◆ The two elements between which the information is transferred are defined beforehand



Multi-point systems (bus)

- More units use the same physical channel:
the BUS
 - ◆ Step 1: Select the units participating in the transfer
 - ◆ Step 2: perform the transfer





Bus systems

- Systems with **bus**: one can change the configuration by removing/adding modules
 - ◆ **MODULAR OPEN** system
- All units must follow the same rules for transferring information
 - ◆ **PROTOCOL**
- Parameters of a bus interface
 - ◆ Physical-electrical structure (connector, formats, levels, ..)
 - ◆ Transaction type, speed, and width (protocol)
 - ◆



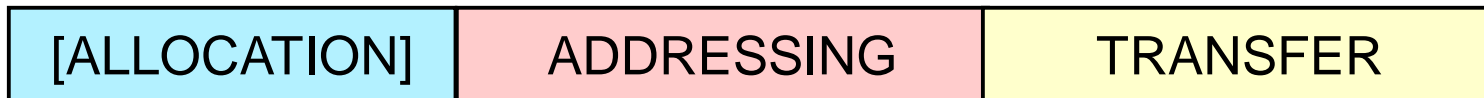
Comparison of bus and point-to-point

- **Point-to-point** connections
 - ◆ Multiple connection ports required on each module
 - ◆ Independent physical channels; well-defined op. conditions
 - ◆ Requires routing for communication between various unit pairs
 - ◆ Can reach higher speeds
- **Multi-point / bus** connections
 - ◆ One port on each unit, one physical channel
 - ◆ Variable electrical conditions (depending on board number)
 - ◆ Requires choice of who controls the channel (arbitrage)
 - ◆ Requires partner selection mechanism (addressing)
 - ◆ Maximum modularity



Transaction protocols

- In multipoint (bus) systems, specific cycles are required to select the participating units
 - ♦ Master selection → **Allocation** (of the channel)
 - ♦ Slave selection → **Addressing**
- Only **after** the transfer can be performed



cycle 1

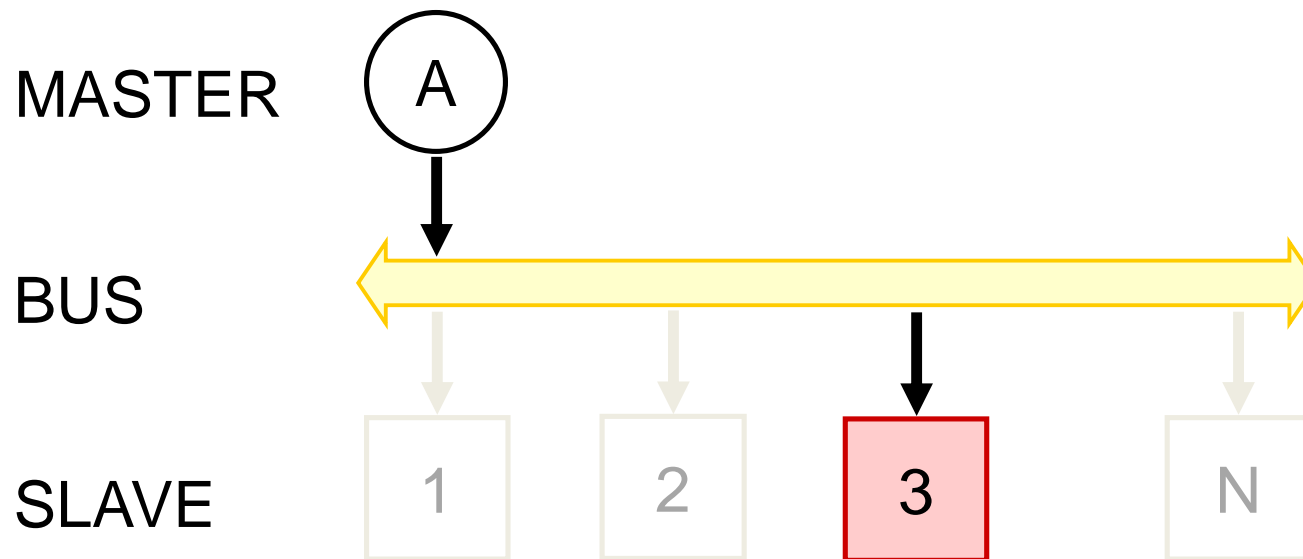
cycle 2

cycle 3



Single master - multislave systems

- The master is predefined (unique)
- The slave is selected with an operation of **ADDRESSING**



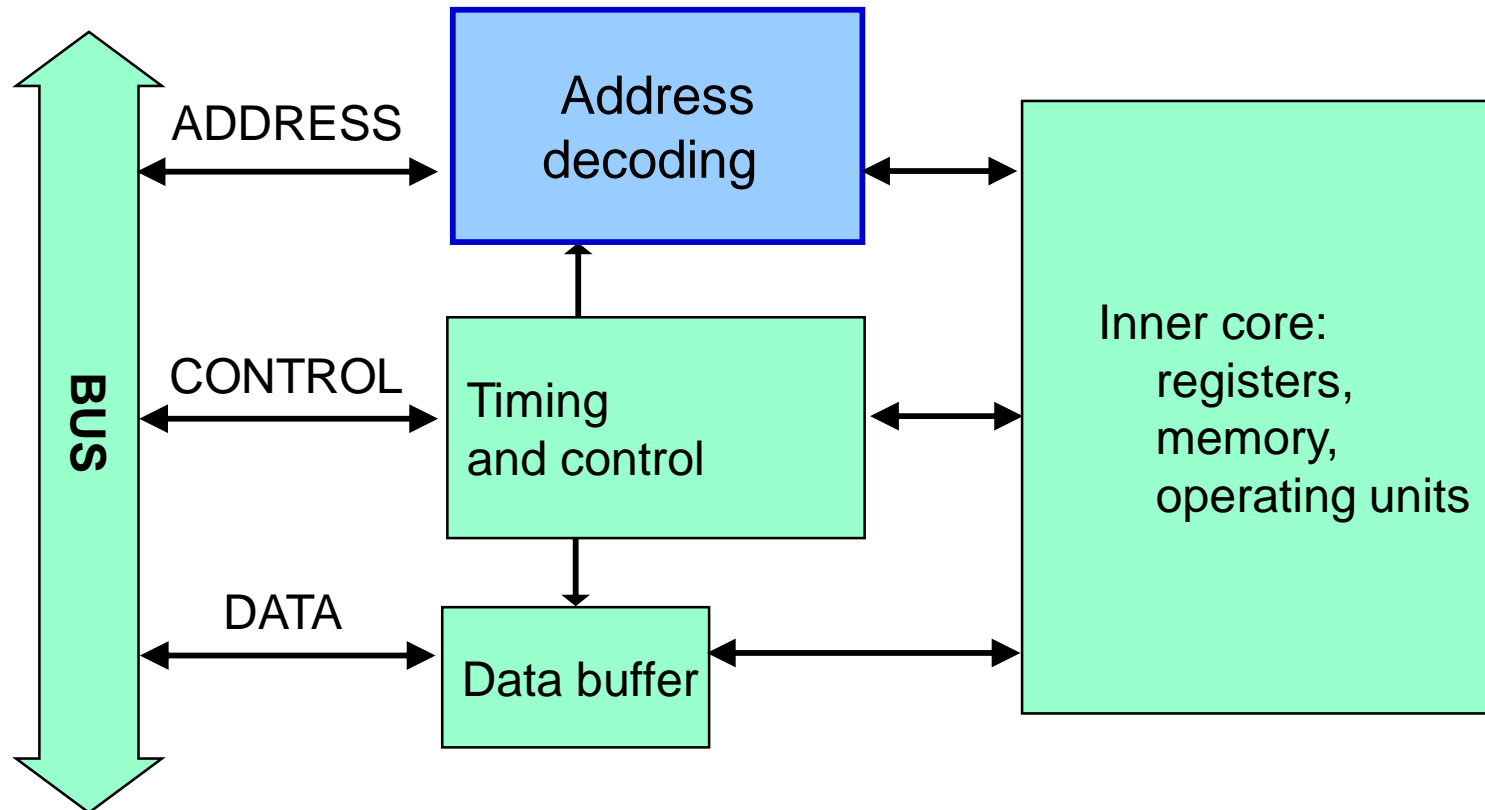


Structure of the Slave units

- Address decoding
 - ◆ Recognizes the address issued by the Master
 - ◆ Uses comparators, decoders, gate logic,
- Timing and control
 - ◆ Generates internal command signals
 - ◆ Manages synchronization (handshake) and data buffers
- Inner core
 - ◆ Registers, memory, operating units, ...
- Data buffer
 - ◆ Adequate for driver and receiver, monolithic transfer to the bus



Block diagram of the Slave units





Addressing techniques

- Coded / Decoded Selection
 - ◆ Encoded: N bits select 1 element out of 2^N
 - Selection of memory and peripheral cells (with address)
 - ◆ Decoded: M bits select 1 (or more) of M elements
 - Direct selection of registers
 - Selection of memory cells after the address decoder
- Logical / Geographic Addressing (position)
 - ◆ Logical: Selection depends on the "name" of the Slave (logical)
 - Memory, registers, ...
 - ◆ Geographic: Selection depends on location (connector)
 - Used to determine configuration and identify cards at system startup

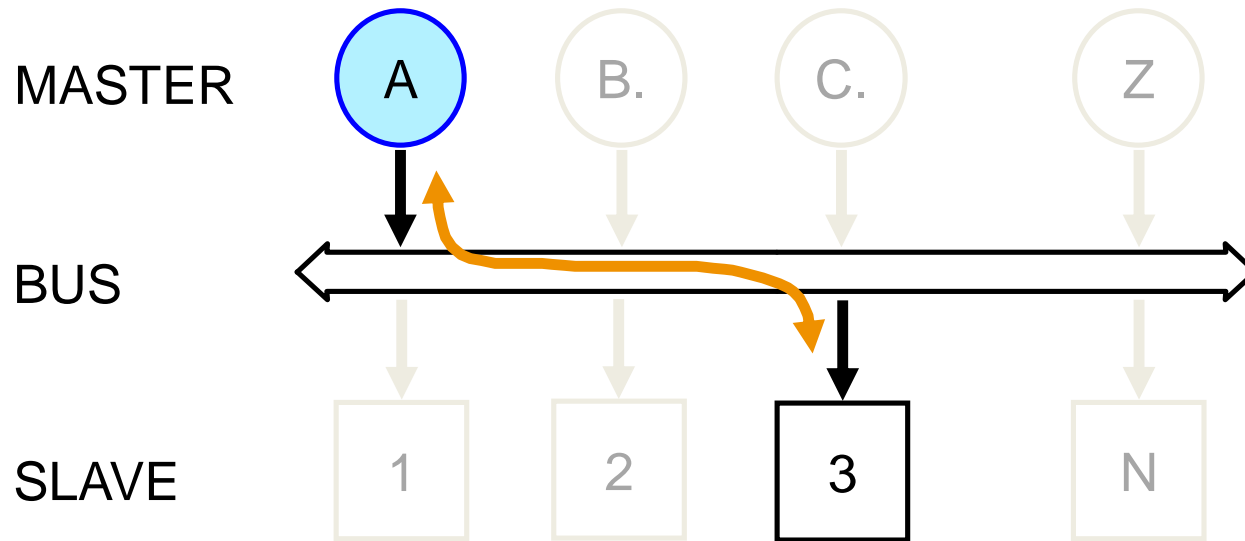


Lecture C5: bus protocols

- Addressing techniques for Single-Master systems
- Channel allocation for Multi-Master systems
 - ◆ Basic allocation mechanisms and techniques
 - ◆ Priority, starvation, fairness
- Performance analysis of a bus
- Techniques to improve the performance
 - ◆ Source-Synchronous transfers
 - ◆ Double Data Rate (DDR) transfers

Multimaster systems: who uses the bus?

- Multimaster / multislave systems
 - ◆ Must **identify** the couple **Master / Slave** for transfers
 - ◆ Operation of **ALLOCATION (ARBITRATION)**





Allocation techniques

- **Collision**: simultaneous access of two users
 - ◆ It must be avoided → **allocation** of the channel before use
- **Allocation mechanisms**
 - ◆ **Token passing**
 - Single token, moved between masters (without evaluating requests)
 - ◆ Collision detector for **collision detection**
 - GRANT to all those who ask for access → **collision**
 - Collision detection mechanisms and retry
 - ◆ **Arbitration**
 - Evaluation of REQUESTS, issue a single GRANT
 - No collision
 - Deterministic access, controllable parameters

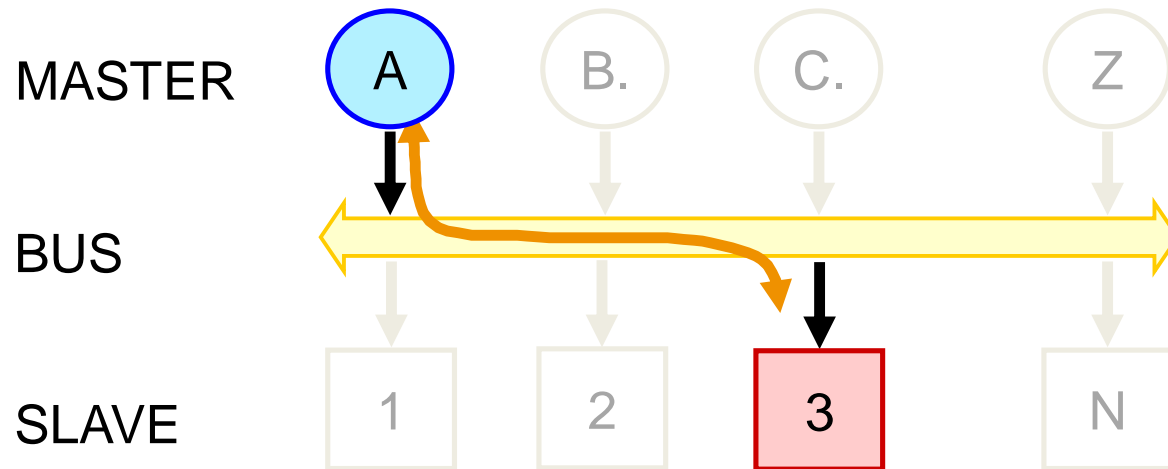


Arbitration policies

- **Time-based** priority: FCFS (First-Come, First-Served)
 - ◆ Limited time resolution
 - ◆ May not order all events → need another mechanism
- **Logic** priority
 - ◆ Define a "rank" of the requests (priority level)
 - ◆ High-priority request sequences can block low priority ones
 - Starvation
 - ◆ Must guarantee servicing in a limited time
 - Fairness
- Next slides: examples of requested / grant sequences:
 - ◆ FCFS referee and priority
 - ◆ Starvation and fairness

Complete sequence summary

- Op. 1: Master selection (arbitration)
- Op. 2: Slave selection (addressing)
 - ◆ Master and Slave form a Point-to-Point “virtual channel”
- Op. 3: M-S pair transfers the information as point-to-point





Lecture C5: bus protocols

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 - ◆ DDR transfers



Performance of a bus

- The amount of information transferred in a unit of time is **THROUGHPUT (T)**
- $T = W \times S$
 - ◆ W: bus width
 - ◆ S: bus speed (cycles/second)
if the cycle has duration t_C , $S = 1 / t_C$
- The duration of the cycle t_C depends on
 - ◆ Electrical parameters: t_{TX}, t_K
 - ◆ Module parameters: $t_{SU}, t_H, t_{WR}, t_{EN}, \dots$
 - ◆ Protocol (number of transitions/cycle, type of protocol, ...)



Performance of protocols

- **Synchronous** protocol (READ/WRITE)
 - ◆ **Fixed speed**, limited by the slowest device
 - ◆ Write cycle time: depends on t_K
 - ◆ Reading cycle time: depends on t_{TX} (without pipeline)
- **Asynchronous** protocol
 - ◆ **Variable speed**, adaptable to the involved devices
 - ◆ Cycle time: depends on t_{TX} , for both READ and WRITE
- Best protocol: **Source Synchronous**
 - ◆ Control signals managed by both units
 - Only the necessary delays

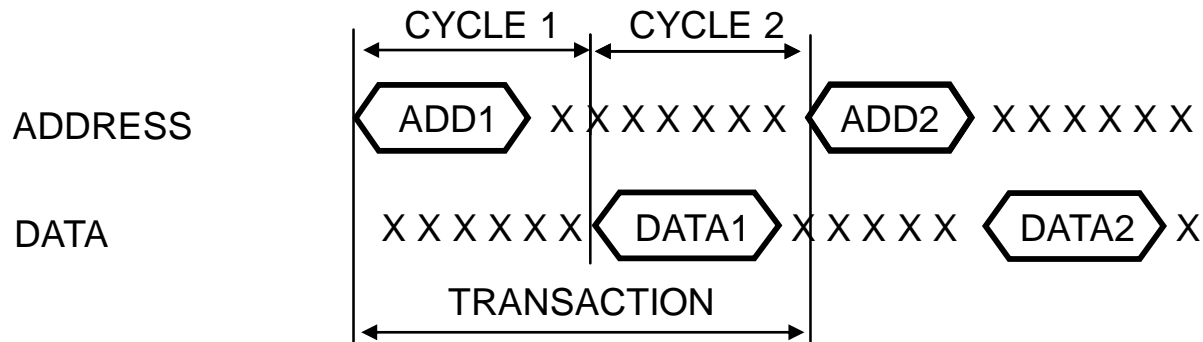


Multiplexed buses

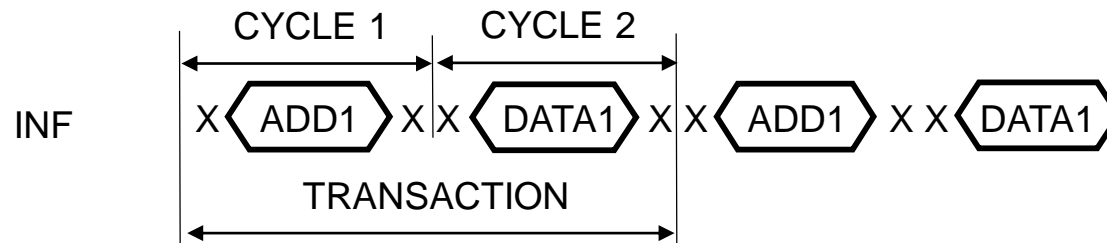
- Basic structure of parallel buses
 - ◆ Addresses and data use different conductors/tracks
 - ◆ Driver/receiver, tracks on the backplane, connector pins are **expensive resources** (require space, use energy, ...)
- Objective
 - ◆ Reduce the number of (physical) conductors
- Solution
 - ◆ Use the same conductor for multiple signals
- Multiple, **MULTIPLEXED**, bus
 - ◆ Addresses, data, other information use the **same conductors**, but at **different times**

Comparison of parallel/multiplexed buses

- Parallel: **separate conductors** for addresses and data



- Multiplexed** (multiplexed): addresses, data, other information use **same conductors**



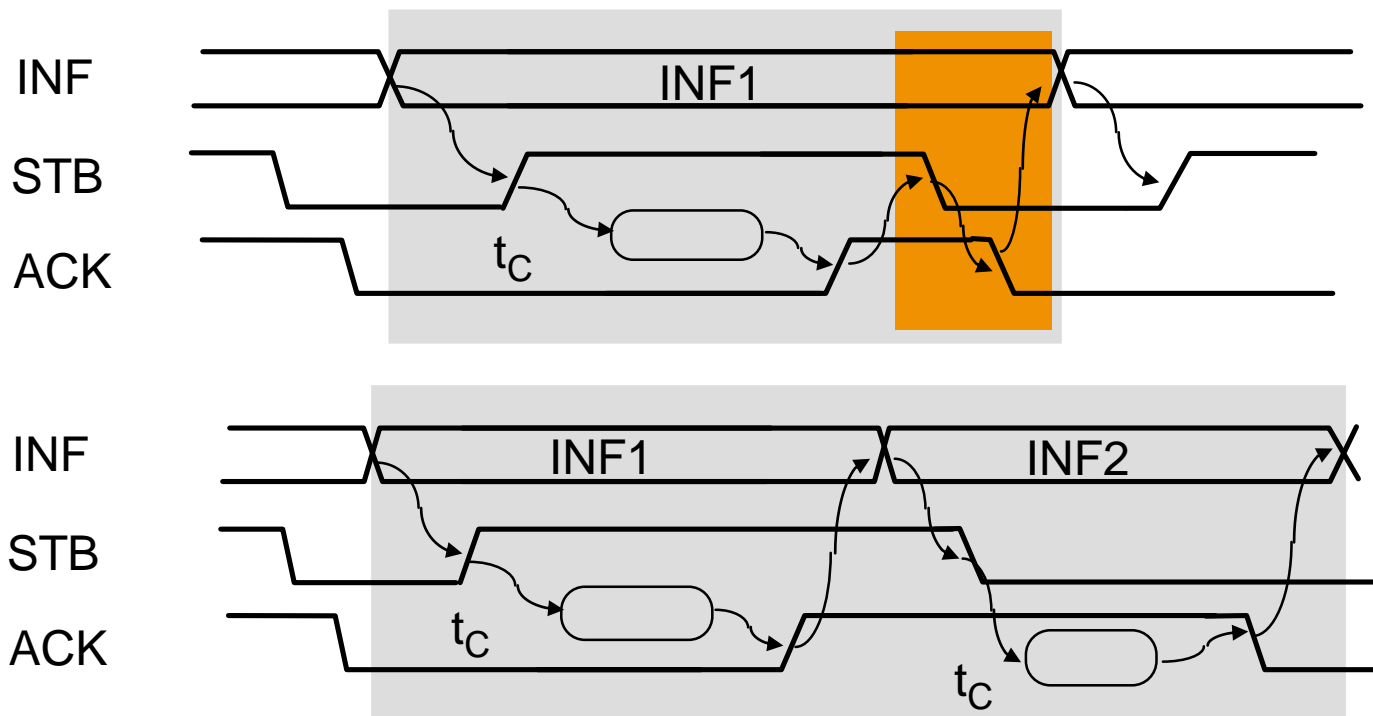


Performance improvement

- Improving performance for the same power consumption?
- Standard Read/Write cycle
 - ◆ Two transitions/cycle for control signals (but *one* used for timing)
 - ◆ Bandwidth and consumption are linked to number of transitions
 - ◆ Data bandwidth lower than control signal bandwidth
 - One transition/cycle compared to two
- DDR, **Double Data Rate**, cycle
 - ◆ Uses **both transitions** of the control signals STB/ACK
 - ◆ Lower consumption
 - ◆ Higher speed
 - ◆ Higher complexity

Two-edge cycles (Dual-edge): DDR

- Removes the unused transitions
- Reduces the consumption / increases the speed



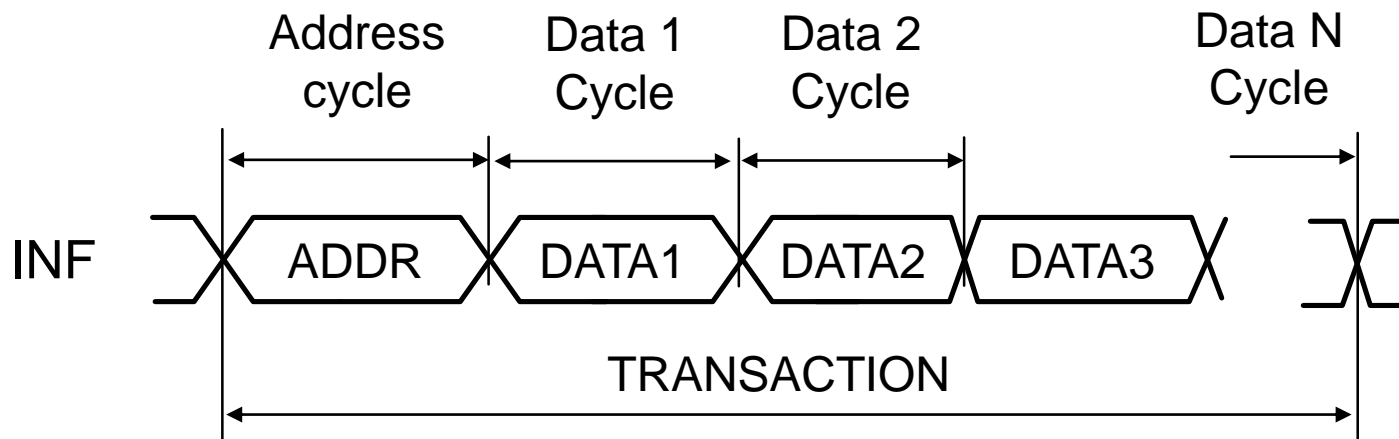


Source Synchronous cycles

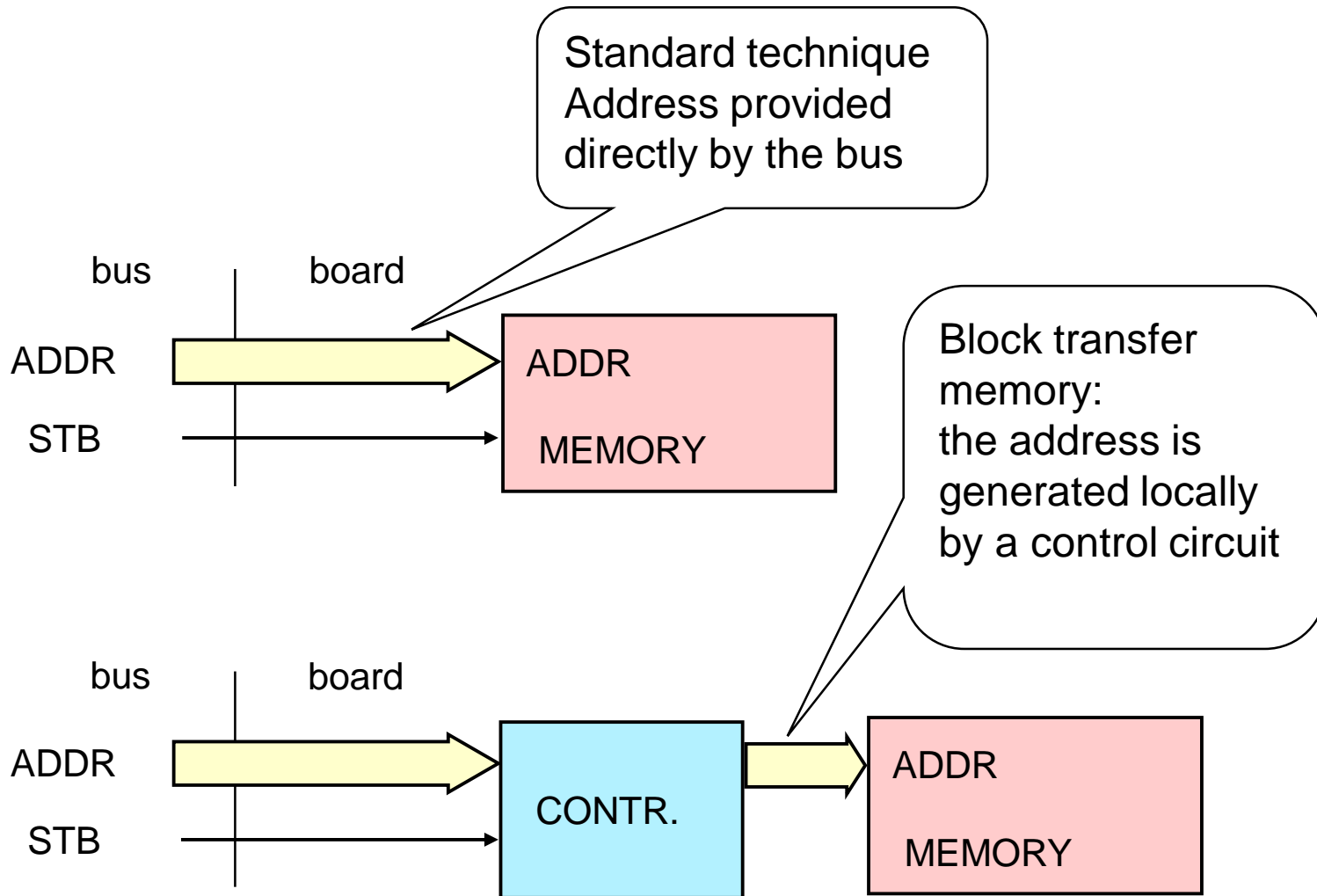
- In the writing cycles, INF and STB move in the same direction, Master \rightarrow Slave
- Two timing parameters
 - ◆ Cycle **duration**, t_C , or cadence in cycles/s f_C ($t_C = 1 / f_C$)
 - Depends on skew t_K , **not** on t_{TX}
 - ◆ Data **latency**, t_L , how long to wait to get the data
 - Depends on t_{TX} **and** on skew t_K
- A new transfer can start **before** the previous one is over
 - ◆ Pipeline technique applied to transfers
 - ◆ Pipeline example: **Source Synchronous** protocols
 - Used for high-speed RAM (**SSTL-2** and later)

Block Transfers (Burst)

- Most transfers in program execution involve sequences of successive memory cells
 - ◆ Send only the first address
 - ◆ **Double the INF amount** for long packages transferred per time
 - Skips the address transfer

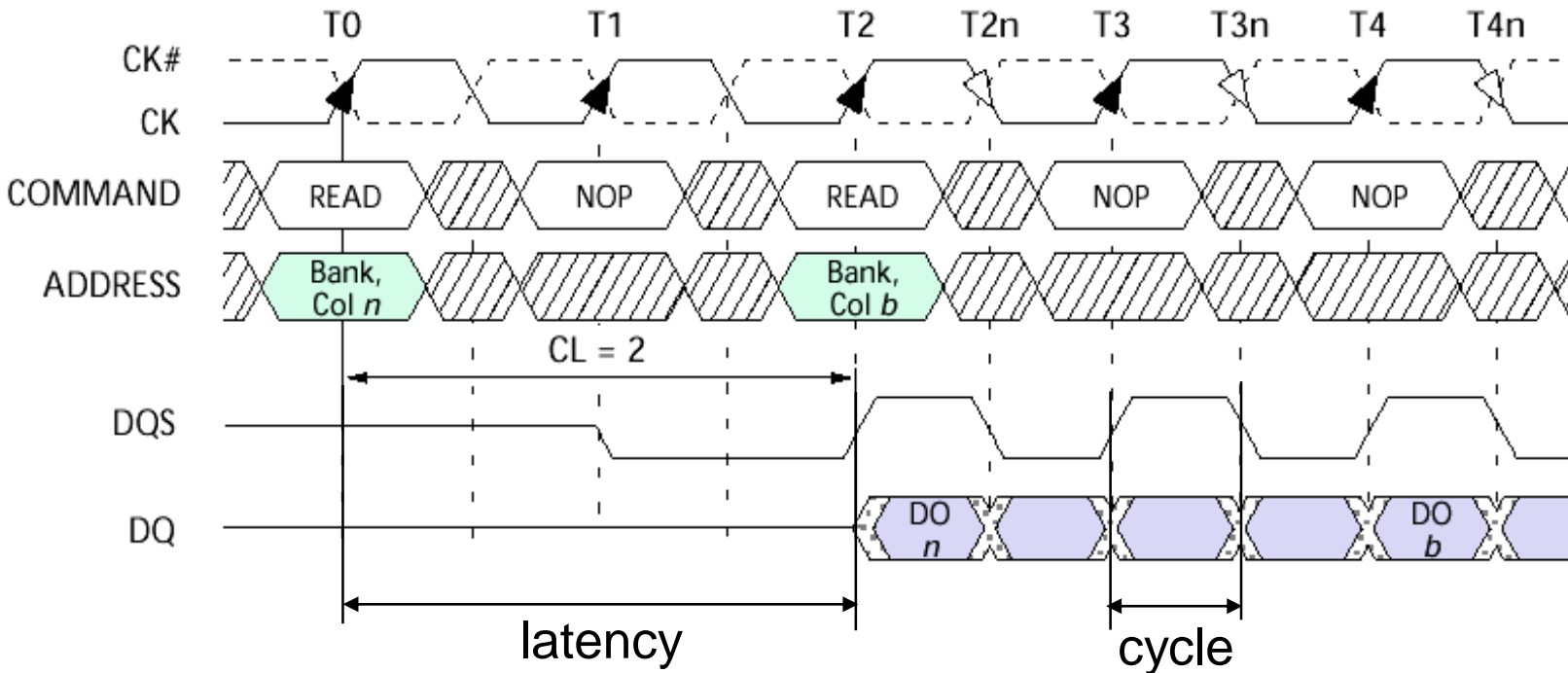


Block transfer interface



Burst transfers: DDR memories

- **Source Synchronous DDR Block transfer**
 - ◆ DQS (Data Strobe) driven by the memory (double edge)
 - ◆ Latency of two clock periods





Beyond Source Synchronous

- With a Source Synchronous (SS) protocol
 - ◆ An SS read requires Slave-controlled STBs
 - ◆ Access time delays are still present
 - ◆ Pipeline for even better performance
- Skew t_K is the key parameter for speed
 - ◆ t_K due to INF and STB using different physical channels (wires)
 - ◆ How to eliminate the skew? Use only one wire!
- Modulations with embedded clock or self-synchronizing
 - ◆ Timing and INF on the same physical signal
 - ◆ Convert parallel transfers into multiple serial channels



Lecture C5 – final test

- When and why must there be an addressing operation?
- In which systems is an arbitration operation necessary?
- Describe the operation of a bus arbiter.
- What is meant by “performance” of a bus?
- What parameters determine the performance of a bus?
- What does the minimum cycle time depend on?
- Compare min cycle times for synch. and asynch. protocol.
- Describe the benefits of DDR protocols.
- What are the advantages of multiplexed buses?
- Describe the benefits of block transfer transactions.