

B561 Advanced Database Concepts

Assignment 2

Fall 2022

Assignment solved by: Nikhil Vemula
Collaborated with: Shreeja Deshpande, Ami Sharma,
Atharva Pore

September 23, 2022

1. Translating SQL queries with query predicates and subquery expressions into safe TRC

1. • Consider the query ‘Find the bno and title of each book that was bought by exactly one student.’ This query can be expressed as the SQL query

Solution:

$$\{(b.bno, b.title) \mid Book(b) \wedge \exists s \in Student (Buys(s.sid, b.bno)) \wedge \forall s1 \in Student \forall s2 \in Student (Buys(s1.sid, b.bno) \wedge Buys(s2.sid, b.bno) \rightarrow (s1 = s2))\}$$

2. • Consider the query ‘Find each pair (m, b) where m is a major and b is the bno of a book bought by a student who has major m and such that the price of b is the lowest among the set of books bought by students with major m .’

Solution:

$$\{(m.major, b.bno) \mid Major(m) \wedge Book(b) \wedge \exists t \in Buys (b.bno = t.bno \wedge \exists hm \in hasMajor (t.sid = hm.sid \wedge hm.major = m.major)) \wedge \neg \exists t \in Buys \neg \exists b1 \in Books (t.bno = b1.bno \wedge hasMajor(t.sid, m.major) \wedge b1.price < b.price)\}$$

2. Expressing queries in (Extended) Safe TRC and Pure SQL with and without subquery expressions

3. Consider the query ‘Find the *bno* and title of each book that is bought by a student who is (strictly) younger than each student who majors in Chemistry and who also bought that book.’
- (a) • Express this query in Safe TRC (i.e, with quantifiers ‘*exists*’ or ‘ \forall ’).

Solution:

$$\{(b.bno, b.title) \mid Book(b) \wedge \exists t1 \in Buys(b.bno = t1.bno \wedge \exists s1 \in Student \exists s2 \in Student(s1.sid = t1.sid \wedge s1.birthyear > s2.birthyear \wedge \exists t2 \in Buys(t2.sid = s2.sid \wedge \exists hm \in hasMajor(hm.sid = t2.sid \wedge hm.major = 'Chemistry'))))\}$$

- (b) • Express this query in Extended Safe TRC (i.e, with subquery expressions).

Solution:

$$\{(b.bno, b.title) \mid Book(b) \wedge exists(1 \mid Buys(t1) \wedge t1.bno = b.bno \wedge exists(1 \mid Student(s1) \wedge Student(s2) \wedge s1.sid = t1.sid \wedge s1.birthyear > s2.birthyear \wedge exists(1 \mid Buys(t2) \wedge t2.sid = s2.sid \wedge exists(1 \mid hasMajor(hm) \wedge hm.sid = t2.sid \wedge hm.major = 'Chemistry'))))\}$$

4. Consider the query ‘Find each student-book pair (*s*, *b*) where *s* is the sid of a student who majors in CS and who bought each book that costs no more than book *b*.’
- (a) • Express this query in Extended Safe TRC (i.e, with subquery expressions).

Solution:

$$\{(s.sid, b.bno) \mid Student(s) \wedge Book(b) \wedge exists(1 \mid hasMajor(hm) \wedge hm.major = 'CS' \wedge hm.sid = s.sid \wedge notexists(1 \mid Book(b1) \wedge$$

$$b1.price \leq b.price \wedge \text{notexists}(1 \mid Buys(t) \wedge t1.bno = b1.bno \wedge t1.sid = s.sid)))\}$$

5. Consider the query ‘Find the sid and name of each student who bought all-but-one book that cost strictly more than \$30.’

- (a) • Express this query in Safe TRC (i.e, with quantifiers ‘exists’ or ‘ \forall ’).

Solution:

$$\{(s.sname, s.sname \mid Student(s) \wedge \exists b1 \in Book(\neg \exists t1 \in Buys(t1.sid = s.sid \wedge b1.bno = t1.bno) \wedge \neg \exists b2 \in Book(b2.bno \neq b1.bno \wedge b2.price \geq 30 \wedge \neg \exists t2 \in Buys((t2.sid = s.sid) \rightarrow (b2.bno = t2.bno))) \wedge b1.price \geq 30))\}$$

3. Expressing queries in Relational Algebra and RA SQL

10. Reconsider the query in Problem 1 ‘Find the bno and title of each book that was bought by exactly one student.’

- (a) • Express this query in Relational Algebra in standard notation.

Solution:

$$\{ \pi_{B.bno, B.title}(B) - \pi_{B1.bno, B1.title}(B1 \bowtie_{B1.bno=T1.bno} (T1 \bowtie_{T1.bno=T2.bno \wedge T1.sid \neq T2.sid} T2)) \}$$

11. Reconsider the query in Problem 2 ‘Find each pair (m, b) where m is a major and b is the bno of a book bought by a student who has major m and such that the price of b is the lowest among the set of books bought by students with major m.’

- (a) • Express this query in Relational Algebra in standard notation.

Solution:

$$\{ \pi_{hM.major, T.bno} (hM \bowtie_{T.sid=hM.sid} (T \bowtie_{T.bno=B.bno} B)) - \pi_{hM.major, T.bno} (hM \bowtie_{T.sid=hM.sid} (T \bowtie_{T.bno=B.bno} (B \bowtie_{hM.major=hM_1.major} (hM_1 \bowtie_{T_1.sid=hM_1.sid} (T_1 \bowtie_{T_1.bno=B_1.bno \wedge B.price > B_1.price} (B_1)))))) \}$$

12. Reconsider the query in Problem 3 ‘Find the bno and title of each book that is bought by a student who is (strictly) younger than each student who majors in Chemistry and who also bought that book.’

- (a) • Express this query in Relational Algebra in standard notation.

Solution:

$$\pi_{bno, title} (Book \bowtie_{b.bno=t1.bno} Buys_1 \bowtie_{t1.sid=s1.sid} Student_1 \bowtie_{s1.birthyear > s2.birthyear} Student_2 \bowtie_{s2.sid=t2.sid} Buys_2 \bowtie_{hm.sid=t2.sid \wedge \sigma_{major='Chemistry'}} hasMajor)$$

13. Reconsider the query in Problem 4 ‘Find each student-book pair (s, b) where s is the sid of a student who majors in CS and who bought each book that costs no more than book b.’

- (a) • Express this query in Relational Algebra in standard notation.

Solution:

$$\pi_{sid, bno} (Student \bowtie_{hm.sid=s.sid \wedge hm.major='CS'} hasMajor) - \pi_{sid, bno} (Student \bowtie_{hm.sid=s.sid \wedge hm.major='CS'} hasMajor \cap \pi_{sid, bno} (Student \bowtie_{hm.sid=s.sid \wedge hm.major='CS'} hasMajor \bowtie_{Book_1.price \leq b.price} Book_1 \bowtie_{t1.sid=s.bno \wedge t1.bno=b1.bno} Buys))$$

14. Reconsider the query in Problem 5 ‘Find the sid and name of each student who bought all-but-one book that cost strictly more than \$30.’

- (a) • Express this query in Relational Algebra in standard notation.

Solution:

$$1 - \pi_{sid} (Student) \bowtie_{t1.sid=s.sid} Buys \bowtie_{t.bno \neq b.bno \wedge b.price > 30} Books$$

4. Expressing constraints using Relational Algebra

19. • Among the books that cite a book, there are books that cite the same set of other books.

Solution:

$$\begin{aligned}
 & \{ \pi_{c_1.bno1}(\pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2}(\sigma_{c_1.bno2=c_2.bno2}(C_1 \times C_2))) \\
 & \cap \pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2}(\sigma_{c_1.bno1 \neq c_2.bno1}(C_1 \times C_2)) \\
 & \cap \pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2}(\pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2}(C_1 \times C_2) - \\
 & \pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2} \\
 & (\pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2,c_3.bno1,c_3.bno2} \\
 & (\sigma_{c_1.bno1=c_3.bno1}(C_1 \times C_2 \times C_3)) \cap \\
 & \pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2,c_3.bno1,c_3.bno2} \\
 & (\pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2,c_3.bno1,c_3.bno2}(C_1 \times C_2 \times C_3) - \\
 & \pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2,c_3.bno1,c_3.bno2} \\
 & (\pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2,c_3.bno1,c_3.bno2,c_4.bno1,c_4.bno2} \\
 & (\sigma_{c_2.bno1=c_4.bno1}(C_1 \times C_2 \times C_3 \times C_4)) \cap \\
 & \pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2,c_3.bno1,c_3.bno2,c_4.bno1,c_4.bno2}(\sigma_{c_3.bno2=c_4.bno2}(C_1 \times \\
 & C_2 \times C_3 \times C_4)))))) \cap \\
 & \pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2}(\pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2}(C_1 \times C_2) - \\
 & \pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2}(\pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2,c_3.bno1,c_3.bno2} \\
 & (\sigma_{c_2.bno1=c_3.bno1}(C_1 \times C_2 \times C_3)) \cap \\
 & \pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2,c_3.bno1,c_3.bno2} \\
 & (\pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2,c_3.bno1,c_3.bno2}(C_1 \times C_2 \times C_3) - \\
 & \pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2,c_3.bno1,c_3.bno2} \\
 & (\pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2,c_3.bno1,c_3.bno2,c_4.bno1,c_4.bno2} \\
 & (\sigma_{c_1.bno1=c_4.bno1}(C_1 \times C_2 \times C_3 \times C_4)) \cap \\
 & \pi_{c_1.bno1,c_1.bno2,c_2.bno1,c_2.bno2,c_3.bno1,c_3.bno2,c_4.bno1,c_4.bno2}(\sigma_{c_3.bno2=c_4.bno2}(C_1 \times \\
 & C_2 \times C_3 \times C_4)))))) \neq \emptyset \}
 \end{aligned}$$

20. • Some student who majors only bought books that were bought by students who major in Math.

Solution:

$$\pi_{h.sid}(\pi_{h.sid,h.major}(hasMajor) -$$

$$\pi_{h.sid, h.major}(\pi_{h.sid, h.major, b.sid, b.bno}(\sigma_{b.sid=h.sid}(hasMajor \times buys))) \cap$$

$$\pi_{h.sid, h.major, b.sid, b.bno}(\pi_{h.sid, h.major, b.sid, b.bno}(hasMajor \times buys) - \pi_{h.sid, h.major, b.sid, b.bno}(\sigma_{h_1.sid < h_2.sid \wedge b_1.sid = h_1.sid \wedge h_1.major = Math}(hasMajor \times buys \times hasMajor_1 \times buys_1)))) \neq \emptyset$$

21. • There are pairs of majors that have no common students who have those majors.

Solution:

$$\pi_{m_1.major, m_2.major}(\pi_{m_1.major, m_2.major}(\sigma_{m_1.major \neq m_2.major}(major_1 \times major_2))) \cap$$

$$\pi_{m_1.major, m_2.major}(\pi_{m_1.major, m_2.major}(major_1 \times major_2) - \pi_{m_1.major, m_2.major}(\pi_{m_1.major, m_2.major, h_1.major}(\sigma_{h_1.major = m_1.major}(major_1 \times major_2 \times hasMajor_1))) \cap$$

$$\pi_{m_1.major, m_2.major, h_1.major}(\sigma_{h_2.major = m_2.major \wedge h_1.sid = h_2.sid}(h_1.major = m_1.major(major_1 \times major_2 \times hasMajor_1 \times hasMajor_2)))) \neq \emptyset$$

22. • Attribute 'sid' in the relation hasMajor is a foreign key referencing the primary key 'sid' in the relation Student

Solution:

$$\pi_{h.sid}(\pi_{h.sid}(hasMajor) - \pi_{h.sid}(\sigma_{h.sid=s.sid}(hasMajor \times student))) = \emptyset$$