CSC282 HW2

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Question 6

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# Helper Function
function CCW(p1, p2, p3):
    #True if counter-clockwise turn
    return (p2.x - p1.x) * (p3.y - p1.y) - (p2.y - p1.y) * (p3.x - p1.x) > 0
function sortPointsLexicographically(points):
    # Sort the points by x-coordinate, breaking ties with y-coordinate
    # Time-complexity O(n log n)
    return sorted(points, key: (point.x, point.y))
# Main Function
function constructSimplePolygon(points):
    # Sort points lexicographically
    sortedPoints = sortPointsLexicographically(points)
    # Build the upper chain
    upperChain = []
    for point in sortedPoints:
        upperChain.append(point)
        while len(upperChain) >= 3 and not CCW_left(upperChain[len(upperChain) - 3],
                                                    upperChain[len(upperChain) - 2],
                                                    upperChain[len(upperChain) - 1]):
            # Remove the middle point to maintain the CCW property
            upperChain.pop(len(upperChain) - 2)
    # Build the lower chain
    lowerChain = []
    for point in reverse(sortedPoints):
        lowerChain.append(point)
        while len(lowerChain) >= 3 and not CCW_left(lowerChain[len(lowerChain) - 3],
                                                    lowerChain[len(lowerChain) - 2],
                                                    lowerChain[len(lowerChain) - 1]):
```

Remove the middle point to maintain the CCW property lowerChain.pop(len(lowerChain) - 2)

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# Combine the chains to form the simple polygon
upperChain.pop(len(upperChain) - 1)
lowerChain.pop(len(lowerChain) - 1)
polygon = upperChain + lowerChain
```

Return the polygon containing all points return polygon

Question 7

function findVisibleSegments(segments, p): # Collect all events (segment endpoints) and compute their angles from point p # Time Complexity: O(n) events = [] for segment in segments: angleStart = computeAngle(p, segment.start) angleEnd = computeAngle(p, segment.end) # Normalize angles to [0, 2pi) angleStart = normalizeAngle(angleStart) angleEnd = normalizeAngle(angleEnd) if angleStart <= angleEnd:</pre> events.append({'angle': angleStart, 'segment': segment, type': 'start'}) events.append({'angle': angleEnd, 'segment': segment, 'type': 'end'}) else: # Segment crosses the O angle events.append({'angle': angleStart, 'segment': segment, 'type': 'start'}) events.append({'angle': angleEnd + 2 * PI, 'segment': segment, 'type': 'end'}) # Sort all events by angle # Time Complexity: O(n log n) sortedEvents = sortByAngle(events) # The BST is ordered by distance from p along the current angle activeSegments = emptyBST() visibleSegments = emptySet() # Sweep through all sorted events to determine visibility # Time Complexity: O(n log n) for event in sortedEvents: angle = event['angle'] % (2 * PI) # Normalize angle within [0, 2pi) segment = event['segment'] if event['type'] == 'start': # Compute the distance from p to the segment along the current angle distance = computeDistanceToSegment(p, angle, segment)

```
# Insert the segment into the activeSegments BST with the computed distance insertSegment(activeSegments, segment, distance)
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Check if this segment is the closest one currently
closestSegment = activeSegments.findMin()
if closestSegment == segment:
 # If it's the closest add to wisibleSegments

If it's the closest, add to visibleSegments
visibleSegments.add(segment)

elif event['type'] == 'end':

Remove the segment from the activeSegments BST
removeSegment(activeSegments, segment)

After removal, check the new closest segment
closestSegment = activeSegments.findMin()
if closestSegment is not None:
 visibleSegments.add(closestSegment)

return visibleSegments

Question 8

The following code addressed the problem: counting intersections of axis-aligned segments in O(nlogn) time. Note that the codec is implemented in Python. The algorithm uses a sweep-line technique combined with a Binary Indexed Tree (BIT) to achieve this time complexity. Segments are represented with endpoints and classified as horizontal or vertical; y-coordinates are compressed into integer indices for efficient BIT operations. Events are created: horizontal segments generate add and remove events at their x-endpoints, while vertical segments generate query events at their x-coordinate. Events are sorted by x-coordinate, processing adds before queries and queries before removes when x-values are equal. As the sweep line advances, the BIT maintains the active set of horizontal segments' y-indices. During query events, the BIT efficiently counts active horizontal segments overlapping the vertical segment's y-range. Noticing that BIT costs $O(\log n)$ time updates and queries. Hence, the sweep line framework achieves the $O(n \log n)$ time complexity.

```
# Axis-Aligned Segment Intersection Counting Algorithm
# Time Complexity: O(n log n)
class Segment:
    def __init__(self, x1, y1, x2, y2):
        # Ensure (x1, y1) is the lower-left point and (x2, y2) is the upper-right point
        if x1 > x2 or y1 > y2:
            x1, x2 = min(x1, x2), max(x1, x2)
            y1, y2 = min(y1, y2), max(y1, y2)
        self.x1, self.y1 = x1, y1
        self.x2, self.y2 = x2, y2
        # Determine if the segment is horizontal or vertical
        self.is_horizontal = y1 == y2
        self.is_vertical = x1 == x2
def coordinate_compress(coordinates):
    unique_coords = sorted(set(coordinates))
    coord_dict = {coord: idx for idx, coord in enumerate(unique_coords)}
    return coord_dict
class BIT:
    def __init__(self, size):
        self.size = size + 2 # +2 to avoid indexing error
        self.tree = [0] * self.size
    def update(self, idx, val):
        idx += 1 # BIT uses 1-based indexing
        while idx < self.size:
            self.tree[idx] += val
```

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idx += idx & -idx
    def query(self, idx):
        idx += 1
        result = 0
        while idx > 0:
            result += self.tree[idx]
            idx -= idx & -idx
        return result
    def range_query(self, 1, r):
        return self.query(r) - self.query(l - 1)
def count_intersections(segments):
   horizontal_segments = []
   vertical_segments = []
    y_coords = []
    for seg in segments:
        if seg.is_horizontal:
            horizontal_segments.append(seg)
            y_coords.append(seg.y1)
        elif seg.is_vertical:
            vertical_segments.append(seg)
            y_coords.append(seg.y1)
            y_coords.append(seg.y2)
    # Coordinate compression for y-coordinates
   y_coord_map = coordinate_compress(y_coords)
   max_y_idx = len(y_coord_map)
    events = []
    # Create add and remove events for horizontal segments
    for seg in horizontal_segments:
        y_idx = y_coord_map[seg.y1]
        events.append((seg.x1, 0, y_idx)) # Add event
        events.append((seg.x2, 2, y_idx)) # Remove event
    # Create query events for vertical segments
    for seg in vertical_segments:
        y1_idx = y_coord_map[seg.y1]
        y2_idx = y_coord_map[seg.y2]
        events.append((seg.x1, 1, min(y1_idx, y2_idx), max(y1_idx, y2_idx)))
    # Sort events by x-coordinate and event type
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events.sort(key=lambda x: (x[0], x[1]))
bit = BIT(max_y_idx)
intersection_count = 0
for event in events:
    if event[1] == 0:
        # Add event: add horizontal segment's y-coordinate to BIT
        y_{idx} = event[2]
        bit.update(y_idx, 1)
    elif event[1] == 2:
        # Remove event: remove horizontal segment's y-coordinate from BIT
        y_idx = event[2]
        bit.update(y_idx, -1)
    else:
        # Query event: count overlapping horizontal segments
        y1_idx, y2_idx = event[2], event[3]
        count = bit.range_query(y1_idx, y2_idx)
        intersection_count += count
return intersection_count
```

Question 9

Picked Probelm: Simple Linear-Time Polygon Triangulation (https://topp.openproblem.net/p10AGR00)