CSC279 HW5

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Nov/5/2023

Collaborator

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Question 15

PROBLEM 15—Closest point, farthest point.

- 1. Assume q is inside P. We want to find the closest point to q in $\{p_1, \ldots, p_n\}$.
- 2. Assume q is outside P. We want to find the closest point to q in $\{p_1, \ldots, p_n\}$.
- 3. Assume q is inside P. We want to find the farthest point to q in $\{p_1, \ldots, p_n\}$.
- 4. Assume q is outside P. We want to find the farthest point to q in $\{p_1, \ldots, p_n\}$.
- 5. Assume q is inside P. We want to find the closest point to q on P.
- 6. Assume q is outside P. We want to find the closest point to q on P.
- 7. Assume q is inside P. We want to find the farthest point to q on P.
- 8. Assume q is outside P. We want to find the farthest point to q on P.

Answer:

Question 1 - 4 Reasoning: The distance from q to the vertices p_i is unimodal function along the ordered sequence. This is true whether q is inside or outside P. Along the vertices, we can use ternary (binary) search on sequence of vertices to find MIN or MAX distance.

- 1. **Problem 1:** q inside P; find the closest point to q in $\{p_1, \ldots, p_n\}$.
 - Solution: $O(\log n)$ time.
 - **Reasoning:** The distance function is unimodal along the vertices, allowing ternary search.
- 2. **Problem 2:** q outside P; find the closest point to q in $\{p_1, \ldots, p_n\}$.

- Solution: $O(\log n)$ time.
- **Reasoning:** Similar to Problem 1, use ternary search due to the unimodal distance function.
- 3. **Problem 3:** q inside P; find the farthest point from q in $\{p_1, \ldots, p_n\}$.
 - Solution: $O(\log n)$ time.
 - **Reasoning:** Unimodal distance function allows ternary search for the maximum.
- 4. **Problem 4:** q outside P; find the farthest point from q in $\{p_1, \ldots, p_n\}$.
 - Solution: $O(\log n)$ time.
 - **Reasoning:** Same as Problem 3, apply ternary search on the unimodal distance function.

Question 5 - 6 Reasoning:

The distance from q to the boundary of the convex polygon P is a convex function along the perimeter, regardless of whether q is inside or outside P. (i.e. The distance decreases to a minimum point and then increases, forming a single through)

- 1. **Problem 5:** q inside P; find the closest point to q on P.
 - Solution: $O(\log n)$ time.
 - **Reasoning:** Perform binary search over edges to find the closest point where the minimum distance occurs.
- 2. **Problem 6:** q outside P; find the closest point to q on P.
 - Solution: $O(\log n)$ time.
 - **Reasoning:** Same as Problem 5, use binary search to find the closest point on edges.

Question 7 - 8 Reasoning:

The farthest point from q can lie anywhere along the perimeter of the convex polygon P, necessitating a check of all edges and vertices. The distance function for farthest points is not unimodal, preventing the use of efficient binary or ternary search methods. Hence, without preprocessing, we need $\Omega(n)$ time to do them.

- 1. **Problem 7:** q inside P; find the farthest point from q on P.
 - Solution: $\Omega(n)$ time.
 - **Reasoning:** Requires examining all edges using rotating calipers for antipodal points.
- 2. **Problem 8:** q outside P; find the farthest point from q on P.
 - Solution: $\Omega(n)$ time.
 - Reasoning: Similar to Problem 7, needs $\Omega(n)$ time to check all edges.

PROBLEM 16

Proof. For a n-polygon,

PROBLEM 17

```
import random
import matplotlib.pyplot as plt
from collections import deque
Input:
    - P = \{p1, p2, ..., pn\}: Set of n points in the plane
    - r: Radius of each circle Ci
    - 1: Radius of the target circle Q, where 1 > r
Output:
    - G: Set of "good" points
Algorithm FindGoodPoints(p_1, ..., p_n, r, 1):
1. Build the Delaunay triangulation (DT) of the points \{p_1, \ldots, p_n\}.
    - Time complexity: O(n log n)
2. Initialize an empty list GoodPoints.
3. For each point p_i in \{p_1, \ldots, p_n\}:
    - Time complexity: O(n)
    a. Initialize an empty list of intervals I_i.
    b. For each neighbor p_{-}j of p_{-}i in DT:
        - Time complexity: O(1).
        i. Compute d = distance between p_i and p_j.
        ii. If d <= 21:
            - Compute theta_j = angle between vector (p_j - p_i) and the x-axis.
            - Compute phi_j = arccos(d / (2 * (1 + r))).
                (Ensure the argument of arccos is within [-1, 1].)
            - If phi_j is real:
                - Add the interval [theta_j - phi_j, theta_j + phi_j] to I_i.
    c. Sort the intervals in I_i.
        - Time complexity: O(1), every Voronoi point have constant neighbor.
        - Merge overlapping intervals to get the union.
```

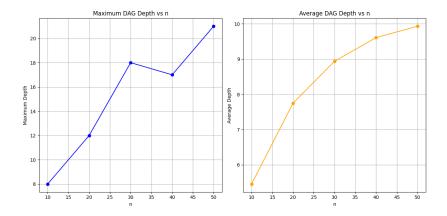
- d. If the union of intervals in I_i covers [0, 2pi]:
 - p_i is not good.
- e. Else:
 - p_i is good.
 - Add p_i to GoodPoints.
- 4. Return GoodPoints.

PROBLEM 18

Result:

$\overline{\textbf{Processing } n}$	Max Depth	Avg Depth
10	8	5.45273631840796
20	12	7.750312109862672
30	18	8.939478067740144
40	17	9.610121836925961
50	21	9.93361327734453

Table 1: Depth statistics for different values of n



Short Analysis:

The result I got is reasonable, with the average depths increasing in *logn* as expected, indicating that the algorithm effectively maintains a balanced DAG structure. Although the maximum depths are somewhat higher than theoretical predictions, they remain within an acceptable range considering the algorithm's randomness and the potential for local depth increases during edge legalization. **Implementation:**

The following code of randomized Delaunay triangulation algorithm with history DAG was implemented in Python:

```
# -----
# Data Structures Definitions
# -----
class Point:
   def __init__(self, x, y):
       self.x = x
       self.y = y
class SegmentWrapper:
   Helper class to uniquely identify a segment irrespective of point order.
   def __init__(self, p1, p2):
       # Ensure consistent ordering
       if (p1.x, p1.y) < (p2.x, p2.y):
           self.p1, self.p2 = p1, p2
       else:
           self.p1, self.p2 = p2, p1
   def __eq__(self, other):
       return (self.p1.x, self.p1.y) == (other.p1.x, other.p1.y) and \
              (self.p2.x, self.p2.y) == (other.p2.x, other.p2.y)
   def __hash__(self):
       return hash(((self.p1.x, self.p1.y), (self.p2.x, self.p2.y)))
   def __repr__(self):
       return f"Segment(({self.p1.x}, {self.p1.y}) - ({self.p2.x}, {self.p2.y}))"
class Triangle:
   def __init__(self, p1, p2, p3):
       self.vertices = [p1, p2, p3] # Points
       self.edges = [
           SegmentWrapper(p1, p2),
           SegmentWrapper(p2, p3),
           SegmentWrapper(p3, p1)
       self.neighbors = {} # edge -> adjacent triangle
       self.children = [] # For history DAG
   def contains_point(self, point):
       return point in self.vertices
```

```
def __repr__(self):
       verts = ', '.join([f"({p.x}, {p.y}))" for p in self.vertices])
       return f"Triangle({verts})"
class HistoryDAGNode:
   def __init__(self, triangle):
       self.triangle = triangle # Triangle object
       self.children = []
                                # List of HistoryDAGNode
   def add_child(self, child_node):
       self.children.append(child_node)
# -----
# Delaunay Triangulation Class
# -----
class DelaunayTriangulation:
   def __init__(self, points):
       self.points = points # List of Point objects
       self.segments = {} # SegmentWrapper -> list of adjacent Triangles
       self.triangles = [] # List of Triangle objects
       self.history_root = None
   def create_super_triangle(self):
       Create a super-triangle that encompasses all the points.
       min_x = min(p.x for p in self.points)
       max_x = max(p.x for p in self.points)
       min_y = min(p.y for p in self.points)
       max_y = max(p.y for p in self.points)
       dx = max_x - min_x
       dy = max_y - min_y
       delta_max = max(dx, dy) * 100 # Make it large enough
       # Create three points that form a super-triangle
       p1 = Point(min_x - delta_max, min_y - delta_max)
       p2 = Point(min_x + 2 * delta_max, min_y - delta_max)
       p3 = Point(min_x - delta_max, max_y + 2 * delta_max)
       super_triangle = Triangle(p1, p2, p3)
       self.triangles.append(super_triangle)
       self.history_root = HistoryDAGNode(super_triangle)
```

```
# Add segments of the super-triangle
    for edge in super_triangle.edges:
        self.segments.setdefault(edge, []).append(super_triangle)
def locate_containing_triangle(self, point):
    Traverse the history DAG to locate the triangle containing the point.
    node = self.history_root
    while node.children:
        found = False
        for child in node.children:
            if self.point_in_triangle(point, child.triangle):
                node = child
                found = True
                break
        if not found:
            break # Point not found in any child; fallback
    # Now, node.triangle should contain the point
    return node.triangle, node
@staticmethod
def point_in_triangle(p, triangle):
    Check if point p is inside the given triangle using barycentric coordinates.
    def sign(p1, p2, p3):
        return (p1.x - p3.x) * (p2.y - p3.y) - \
               (p2.x - p3.x) * (p1.y - p3.y)
   b1 = sign(p, triangle.vertices[0], triangle.vertices[1]) < 0.0
   b2 = sign(p, triangle.vertices[1], triangle.vertices[2]) < 0.0
   b3 = sign(p, triangle.vertices[2], triangle.vertices[0]) < 0.0
    return ((b1 == b2) and (b2 == b3))
def insert_point(self, point):
    Insert a single point into the triangulation.
    containing_triangle, containing_node = self.locate_containing_triangle(point)
    # Remove the containing triangle
    self.triangles.remove(containing_triangle)
    # Create new triangles by connecting the point
```

```
# to the vertices of the containing triangle
   t1 = Triangle(point, containing_triangle.vertices[0],
        containing_triangle.vertices[1])
   t2 = Triangle(point, containing_triangle.vertices[1],
        containing_triangle.vertices[2])
    t3 = Triangle(point, containing_triangle.vertices[2],
        containing_triangle.vertices[0])
   # Set neighbors
    self.set_neighbors(t1, containing_triangle)
    self.set_neighbors(t2, containing_triangle)
    self.set_neighbors(t3, containing_triangle)
    # Add new triangles
    self.triangles.extend([t1, t2, t3])
   # Update segments
    for t in [t1, t2, t3]:
        for edge in t.edges:
            self.segments.setdefault(edge, []).append(t)
    # Update history DAG
    child_nodes = [
        HistoryDAGNode(t1),
        HistoryDAGNode(t2),
       HistoryDAGNode(t3)
   for child in child_nodes:
        containing_node.add_child(child)
   # Edge Legalization
   for t in [t1, t2, t3]:
        for edge in t.edges:
            if point in [edge.p1, edge.p2]:
                self.legalize_edge(edge, t, point)
def set_neighbors(self, new_triangle, old_triangle):
    11 11 11
   Set the neighboring triangles for the new triangle.
   for edge in new_triangle.edges:
        if edge in self.segments:
            for neighbor in self.segments[edge]:
                if neighbor != old_triangle and neighbor != new_triangle:
                    new_triangle.neighbors[edge] = neighbor
                    neighbor.neighbors[edge] = new_triangle
```

```
def legalize_edge(self, edge, triangle, point):
    Legalize an edge to restore the Delaunay condition.
    if edge not in triangle.neighbors:
        return # Boundary edge
    neighbor = triangle.neighbors[edge]
    if neighbor is None:
        return
    # Find the opposite point in the neighbor triangle
    opposite_point = [v for v in neighbor.vertices
                        if v not in [edge.p1, edge.p2]][0]
    if self.in_circumcircle(opposite_point, triangle):
        # Perform edge flip
        new_triangles = self.edge_flip(triangle, neighbor, edge,
                                        point, opposite_point)
        for new_t in new_triangles:
           for e in new_t.edges:
                if point in [e.p1, e.p2]:
                    self.legalize_edge(e, new_t, point)
def edge_flip(self, t1, t2, edge, point, opposite_point):
    Flip the shared edge between two triangles.
    # Remove old triangles
    self.triangles.remove(t1)
    self.triangles.remove(t2)
    # Remove edge from segments
    self.segments[edge].remove(t1)
    self.segments[edge].remove(t2)
    if not self.segments[edge]:
        del self.segments[edge]
    # Create new edge
    new_edge = SegmentWrapper(point, opposite_point)
    # Create new triangles
   new_t1 = Triangle(point, edge.p1, opposite_point)
    new_t2 = Triangle(point, opposite_point, edge.p2)
```

```
# Update segments
    for t in [new_t1, new_t2]:
        for e in t.edges:
            self.segments.setdefault(e, []).append(t)
    # Update neighbors
    self.update_neighbors_after_flip(t1, t2, new_t1, new_t2, edge, new_edge)
    # Add new triangles
    self.triangles.extend([new_t1, new_t2])
    # Update history DAG
    parent_node = HistoryDAGNode(None)
    t1_node = self.find_history_node(self.history_root, t1)
    t2_node = self.find_history_node(self.history_root, t2)
    parent_node.add_child(t1_node)
   parent_node.add_child(t2_node)
    new_t1_node = HistoryDAGNode(new_t1)
    new_t2_node = HistoryDAGNode(new_t2)
    parent_node.add_child(new_t1_node)
    parent_node.add_child(new_t2_node)
    return [new_t1, new_t2]
def update_neighbors_after_flip(self, t1, t2, new_t1, new_t2, old_edge, new_edge):
    Update neighbor relationships after an edge flip.
    # Set neighbors for new_t1
    new_t1.neighbors[new_edge] = new_t2
   new_t2.neighbors[new_edge] = new_t1
    # Update other neighbors
    for e in new_t1.edges:
        if e != new_edge:
            for neighbor in self.segments[e]:
                if neighbor != new_t1:
                    new_t1.neighbors[e] = neighbor
                    neighbor.neighbors[e] = new_t1
    for e in new_t2.edges:
        if e != new_edge:
            for neighbor in self.segments[e]:
                if neighbor != new_t2:
                    new_t2.neighbors[e] = neighbor
                    neighbor.neighbors[e] = new_t2
```

```
def find_history_node(self, node, triangle):
    Find the history DAG node corresponding to the given triangle.
    if node.triangle == triangle:
        return node
    for child in node.children:
        result = self.find_history_node(child, triangle)
        if result is not None:
            return result
    return None
def in_circumcircle(self, point, triangle):
    Check if a point is inside the circumcircle of a triangle.
    ax, ay = triangle.vertices[0].x - point.x, triangle.vertices[0].y - point.y
    bx, by = triangle.vertices[1].x - point.x, triangle.vertices[1].y - point.y
    cx, cy = triangle.vertices[2].x - point.x, triangle.vertices[2].y - point.y
    det = (ax * (by * (cx**2 + cy**2) - cy * (bx**2 + by**2)) -
           ay * (bx * (cx**2 + cy**2) - cx * (bx**2 + by**2)) +
           (ax**2 + ay**2) * (bx * cy - cx * by))
    return det > 0
def build_triangulation(self):
   Build the triangulation by inserting all points.
    self.create_super_triangle()
    for point in self.points:
        self.insert_point(point)
    self.remove_super_triangle()
def remove_super_triangle(self):
    Remove any triangles that share a vertex with the super-triangle.
    # Super-triangle vertices
    super_vertices = set(self.history_root.triangle.vertices)
    self.triangles = [t for t in self.triangles
                        if not any(v in super_vertices for v in t.vertices)]
def calculate_dag_depths(self):
```

```
Calculate maximum and average depths of the history DAG.
    depths = []
    queue = deque([(self.history_root, 0)])
    while queue:
        node, depth = queue.popleft()
        if not node.children:
            depths.append(depth)
        else:
            for child in node.children:
                queue.append((child, depth + 1))
   max_depth = max(depths) if depths else 0
    avg_depth = sum(depths) / len(depths) if depths else 0
    return max_depth, avg_depth
def plot_triangulation(self):
    Plot the triangulation.
    plt.figure(figsize=(8, 8))
    for triangle in self.triangles:
        x_coords = [vertex.x for vertex in triangle.vertices + [triangle.vertices[0]]]
        y_coords = [vertex.y for vertex in triangle.vertices + [triangle.vertices[0]]]
        plt.plot(x_coords, y_coords, 'k-')
    plt.scatter([p.x for p in self.points],
                [p.y for p in self.points], color='red', s=10)
   plt.axis('equal')
   plt.title('Delaunay Triangulation')
   plt.show()
```