

# CSC282 HW2

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## Question 6

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# Helper Function
function insertPointInOrder(polygon, point):
    # Given that the points are sorted, append to the end of the list
    polygon.append(point)

    # Ensure that the polygon remains simple
    while len(polygon) >= 3 and !CCW_left(polygon[len(polygon) - 3],
    polygon[len(polygon) - 2], polygon[len(polygon) - 1]):
        # Remove 2nd last element
        polygon.pop(len(polygon) - 2)

function CCW_left(p1, p2, p3):
    return (p2.x - p1.x) * (p3.y - p1.y) - (p2.y - p1.y) * (p3.x - p1.x) > 0

function sortPointsByAngle(points, refPoint):
    # Sort points based on the angle they make with the reference point
    sortedPoints = []
    for point in points:
        angle = calculateAngle(refPoint, point)
        sortedPoints.append((point, angle))

    # Sort by the calculated angles (e.g. Using Merge Sort)
    sortedPoints.sortByAngle()

    return [point for point, angle in sortedPoints]

function calculateAngle(p1, p2):
    # Calculate the angle between the line from p1 -> p2 and the x-axis
    # atan2() computes arctangent of the quotient
    return atan2(p2.y - p1.y, p2.x - p1.x)
```

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# Main function
function constructSimplePolygon(points):
    # Step 1: Find the convex hull using Graham Scan
    # ( $O(n \log n)$ )
    hull = grahamScan(points)

    # Remove points in the hull from the original set
    remainingPoints = points.remove(hull)

    # Sort remaining points by their angle relative to a fixed point (say hull[0])
    #  $O(n \log n)$ 
    sortedPoints = sortPointsByAngle(remainingPoints, hull[0])

    # Construct the simple polygon
    polygon = hull

    # *Insertion logic cost  $O(n)$  at most (since each insertion takes
    constant time due to the sorted order)
    for point in sortedPoints:
        insertPointInOrder(polygon, point)

    return polygon

```

## Question 7

```
function findVisibleSegments(segments, p):

# Collect all events (segment endpoints) and compute their angles from point p
# Time Complexity: O(n)
events = []
for segment in segments:
    angleStart = computeAngle(p, segment.start)
    angleEnd = computeAngle(p, segment.end)

    # Normalize angles to [0, 2)
    angleStart = normalizeAngle(angleStart)
    angleEnd = normalizeAngle(angleEnd)

    if angleStart <= angleEnd:
        events.append({'angle': angleStart, 'segment': segment, 'type': 'start'})
        events.append({'angle': angleEnd, 'segment': segment, 'type': 'end'})
    else:
        # Segment crosses the 0 angle
        events.append({'angle': angleStart, 'segment': segment, 'type': 'start'})
        events.append({'angle': angleEnd + 2 * PI, 'segment': segment, 'type': 'end'})

# Sort all events by angle
# Time Complexity: O(n log n)
sortedEvents = sortByAngle(events)

# The BST is ordered by distance from p along the current angle
activeSegments = emptyBST()

visibleSegments = emptySet()

# Sweep through all sorted events to determine visibility
# Time Complexity: O(n log n)
for event in sortedEvents:
    angle = event['angle'] % (2 * PI) # Normalize angle within [0, 2)
    segment = event['segment']

    if event['type'] == 'start':
        # Compute the distance from p to the segment along the current angle
        distance = computeDistanceToSegment(p, angle, segment)

        # Insert the segment into the activeSegments BST with the computed distance
        insertSegment(activeSegments, segment, distance)

        # Check if this segment is the closest one currently
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        closestSegment = activeSegments.findMin()
        if closestSegment == segment:
            # If it's the closest, add to visibleSegments
            visibleSegments.add(segment)
    elif event['type'] == 'end':
        # Remove the segment from the activeSegments BST
        removeSegment(activeSegments, segment)

        # After removal, check the new closest segment
        closestSegment = activeSegments.findMin()
        if closestSegment is not None:
            visibleSegments.add(closestSegment)

return visibleSegments

```

## Question 8

The following code addressed the problem: counting intersections of axis-aligned segments in  $O(n \log n)$  time. Note that the codec is implemented in *Python*.

The algorithm uses a sweep-line technique combined with a Binary Indexed Tree (BIT) to achieve this time complexity. Segments are represented with endpoints and classified as horizontal or vertical; y-coordinates are compressed into integer indices for efficient BIT operations. Events are created: horizontal segments generate add and remove events at their x-endpoints, while vertical segments generate query events at their x-coordinate. Events are sorted by x-coordinate, processing adds before queries and queries before removes when x-values are equal. As the sweep line advances, the BIT maintains the active set of horizontal segments' y-indices. During query events, the BIT efficiently counts active horizontal segments overlapping the vertical segment's y-range. Noticing that BIT costs  $O(\log n)$  time updates and queries. Hence, the sweep line framework achieves the  $O(n \log n)$  time complexity.

# Axis-Aligned Segment Intersection Counting Algorithm

# Time Complexity:  $O(n \log n)$

```
class Segment:
    def __init__(self, x1, y1, x2, y2):
        # Ensure (x1, y1) is the lower-left point and (x2, y2) is the upper-right point
        if x1 > x2 or y1 > y2:
            x1, x2 = min(x1, x2), max(x1, x2)
            y1, y2 = min(y1, y2), max(y1, y2)
        self.x1, self.y1 = x1, y1
        self.x2, self.y2 = x2, y2
        # Determine if the segment is horizontal or vertical
        self.is_horizontal = y1 == y2
        self.is_vertical = x1 == x2

def coordinate_compress(coordinates):
    unique_coords = sorted(set(coordinates))
    coord_dict = {coord: idx for idx, coord in enumerate(unique_coords)}
    return coord_dict

class BIT:
    def __init__(self, size):
        self.size = size + 2 # +2 to avoid index issues
        self.tree = [0] * self.size

    def update(self, idx, val):
        idx += 1 # BIT uses 1-based indexing
        while idx < self.size:
            self.tree[idx] += val
```

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        idx += idx & -idx

def query(self, idx):
    idx += 1
    result = 0
    while idx > 0:
        result += self.tree[idx]
        idx -= idx & -idx
    return result

def range_query(self, l, r):
    return self.query(r) - self.query(l - 1)

def count_intersections(segments):
    horizontal_segments = []
    vertical_segments = []
    y_coords = []

    for seg in segments:
        if seg.is_horizontal:
            horizontal_segments.append(seg)
            y_coords.append(seg.y1)
        elif seg.is_vertical:
            vertical_segments.append(seg)
            y_coords.append(seg.y1)
            y_coords.append(seg.y2)

    # Coordinate compression for y-coordinates
    y_coord_map = coordinate_compress(y_coords)
    max_y_idx = len(y_coord_map)

    events = []

    # Create add and remove events for horizontal segments
    for seg in horizontal_segments:
        y_idx = y_coord_map[seg.y1]
        events.append((seg.x1, 0, y_idx)) # Add event
        events.append((seg.x2, 2, y_idx)) # Remove event

    # Create query events for vertical segments
    for seg in vertical_segments:
        y1_idx = y_coord_map[seg.y1]
        y2_idx = y_coord_map[seg.y2]
        events.append((seg.x1, 1, min(y1_idx, y2_idx), max(y1_idx, y2_idx)))

    # Sort events by x-coordinate and event type

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events.sort(key=lambda x: (x[0], x[1]))

bit = BIT(max_y_idx)
intersection_count = 0

for event in events:
    if event[1] == 0:
        # Add event: add horizontal segment's y-coordinate to BIT
        y_idx = event[2]
        bit.update(y_idx, 1)
    elif event[1] == 2:
        # Remove event: remove horizontal segment's y-coordinate from BIT
        y_idx = event[2]
        bit.update(y_idx, -1)
    else:
        # Query event: count overlapping horizontal segments
        y1_idx, y2_idx = event[2], event[3]
        count = bit.range_query(y1_idx, y2_idx)
        intersection_count += count

return intersection_count

```

## Question 9

Picked Problem: **Simple Linear-Time Polygon Triangulation**  
(<https://topp.openproblem.net/p10AGR00>)