



THE UNIVERSITY OF
SYDNEY

CONFIDENTIAL EXAM PAPER

This paper is not to be removed from the exam venue.

Computer Science

EXAMINATION

Semester 1- Practice, 2022

COMP2123 Data Structures and Algorithms

EXAM WRITING TIME: 2 hours

READING TIME: 10 minutes

EXAM CONDITIONS:

This is a OPEN book examination.

All submitted work must be **done individually** without consulting someone else's help, in accordance with the University's "Academic Dishonesty and Plagiarism" policies.

MATERIALS PERMITTED IN THE EXAM VENUE:

MATERIALS TO BE SUPPLIED TO STUDENTS:

INSTRUCTIONS TO STUDENTS:

Type your answers in your text editor (Latex, Word, etc.) and convert it into a pdf file.

Submit this pdf file via Canvas. No other file format will be accepted. Hand-written responses will **not** be accepted.

Start by typing you student ID at the top of the first page of your submission. Do **not** type your name.

Submit only your answers to the questions. Do **not** copy the questions.

Do **not** copy any text from the permitted materials. Always write your answers in your own words.

For examiner use only:

Problem	1	2	3	4	Total
Marks					
Out of	10	10	20	20	60

Problem 1.

a) Analyze the time complexity of this algorithm.

[3 marks]

```

1: def COMPUTE(A)
2:   result  $\leftarrow$  0
3:   for i = 0; i < n; i ++ do
4:     if A[i] > i then
5:       result  $\leftarrow$  result + A[i]
6:   return result

```

b) Solve the following recursion:

[3 marks]

$$T(n) = \begin{cases} T(n/2) + O(1) & \text{for } n > 1 \\ O(1) & \text{for } n = 1 \end{cases}$$

c) We are planning a board games event and we're using one of the shelves in my office to store the games. Unfortunately the shelf only has a certain amount of space S , so we need to carefully pick which games we want to bring. Every game takes some space s_i and has a fun factor f_i that indicates how much fun it is to play that game (for $1 \leq i \leq n$). [4 marks]

We want to maximize the amount of fun we'll have, so we want to maximize the sum of the fun factors of the games we pick (i.e., $\max \sum_{\text{picked game } i} f_i$), while

making sure that the games fit on my shelf, so the sum of the space the games we pick take should be at most S (i.e., $\sum_{\text{picked game } i} s_i \leq S$). For simplicity, you can

assume that all f_i , s_i , and S are all distinct positive integers.

The strategy of PICKLARGEST is to always pick the game with the highest fun factor until my shelf is full: it sorts the games by their fun factor f_i in decreasing order and adds a game when its required space is less than the remaining space on the shelf.

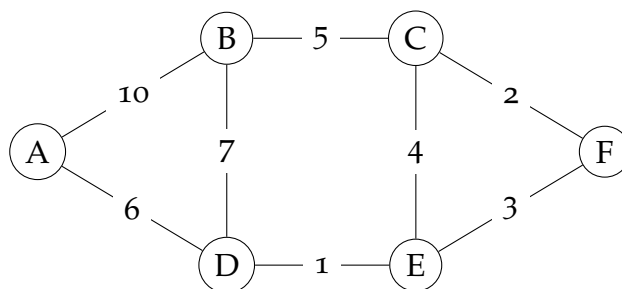
```

1: def PICKLARGEST(all  $f_i$  and  $s_i$ ,  $S$ )
2:   currentSpace  $\leftarrow$  0
3:   currentFun  $\leftarrow$  0
4:   Sort games by  $f_i$  and renumber such that  $f_1 \geq f_2 \geq \dots \geq f_n$ 
5:   for  $i \leftarrow 1$ ;  $i \leq n$ ;  $i++$  do
6:     if currentSpace +  $s_i \leq S$  then
7:       currentSpace  $\leftarrow$  currentSpace +  $s_i$            ▷ Pick the ith game
8:       currentFun  $\leftarrow$  currentFun +  $f_i$ 
9:   return currentFun

```

Show that PICKLARGEST doesn't always return the correct solution by giving a counterexample.

Problem 2. Consider the following edge weighted undirected graph G :



Your task is to:

a) Compute a minimum spanning tree T of G . List the edges in T .

[7 marks]

b) Indicate the order in which the edges are added by Kruskal's algorithm.

[3 marks]

(You do **not** have to explain your answer.)

Problem 3. Consider the *Dynamic Matrix* ADT for representing an matrix $A = \{a_{i,j}\}_{i,j=1}^n$ that supports the following operations:

- **CREATE()**: creates a 1×1 matrix where $a_{1,1} = 0$.
- **SET/GET(i, j)**: set or get the value of the entry $a_{i,j}$.
- **INCREASE-SIZE**: If the current size of the matrix is $n \times n$, increase it to $n + 1 \times n + 1$ such that the new entries are set of 0. In other words, A becomes A' such that $a'_{i,j} = a_{i,j}$ if $1 \leq i, j \leq n$, and $a'_{i,j} = 0$ otherwise.

Your task is to come up with a data structure implementation for the Dynamic Matrix ADT that uses $O(n^2)$ space, where n is the size of the matrix, and **CREATE**, **SET**, **GET** take $O(1)$ and **INCREASE-SIZE** takes $O(n)$ time. Remember to:

- a) Describe your data structure implementation in plain English. [8 marks]
- b) Prove the correctness of your data structure. [7 marks]
- c) Analyze the time and space complexity of your data structure. [5 marks]

Problem 4. Let G be a connected undirected graph on n vertices. We say that two distinct spanning trees T and S of G are *one swap away* from each other if $|T \cap S| = n - 2$; that is, T and S differ in only one edge.

For two distinct spanning trees T and S we say that R_1, R_2, \dots, R_k form a *swapping sequence* from T to S if:

1. $R_1 = T$,
2. $R_k = S$, and
3. for any $1 \leq i < k$, the trees R_i and R_{i+1} are one swap away from each other

Your task is to design a polynomial time algorithm that given G and two spanning trees T and S of G , constructs a minimum length swapping sequence. Remember to:

- a) Describe your algorithm in plain English. [8 marks]
- b) Prove the correctness of your algorithm. [7 marks]
- c) Analyze the time complexity of your algorithm. [5 marks]