

# GTU-C312 CPU Simulator and Operating System Project Report

**Student Name:** Furkan Filicioğlu

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## 1. Project Overview

This project involves the design and simulation of a hypothetical CPU architecture called GTU-C312, along with a simple operating system and multi-threading system running on it. The CPU simulator was implemented in Python, while the operating system and threads were written in GTU-C312 assembly language.

### 1.1 Core Components

- **CPU Simulator:** Python-based simulator supporting GTU-C312 instruction set
  - **Operating System:** Cooperative multitasking OS written in GTU-C312 assembly
  - **Threads:** 4 user threads running different algorithms
  - **Debug Systems:** 4 different debug levels for detailed monitoring
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## 2. CPU Architecture (GTU-C312)

### 2.1 Instruction Set Architecture

The GTU-C312 CPU is a register-less architecture that uses special memory locations as registers:

#### Memory Address Register

0	Program Counter (PC)
1	Stack Pointer (SP)
2	System Call Result
3	Executed Instruction Count
4-20	Reserved Areas

### 2.2 Supported Instructions

#### 2.2.1 Data Processing Instructions

- **SET B A:** Write value B to address A
- **CPY A1 A2:** Copy value from address A1 to A2
- **CPYI A1 A2:** Copy from address pointed by A1 to A2 (indirect copy)
- **CPYI2 A1 A2:** Copy from address pointed by A1 to address pointed by A2

- **ADD A B:** Add B to value at address A
- **ADDI A1 A2:** Add A1 + A2, store result in A1
- **SUBI A1 A2:** Subtract A1 - A2, store result in A2

### 2.2.2 Control Flow Instructions

- **JIF A C:** Jump to instruction C if address A  $\leq 0$
- **CALL C:** Call instruction C, push return address to stack
- **RET:** Pop return address from stack and return

### 2.2.3 Stack Operations

- **PUSH A:** Push value at address A to stack
- **POP A:** Pop value from stack to address A

### 2.2.4 System Instructions

- **USER A:** Switch to user mode and jump to PC at address A
- **HLT:** Halt the CPU
- **SYSCALL PRN A:** Print value at address A (blocks for 100 cycles)
- **SYSCALL YIELD:** Yield CPU, go to scheduler
- **SYSCALL HLT\_THREAD:** Terminate current thread

## 2.3 Security Models

- **KERNEL MODE:** Full memory access (0-16383)
- **USER MODE:** Restricted access ( $\geq 1000$ ), violation terminates thread

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## 3. Operating System Design

### 3.1 OS Architecture

The operating system implements a cooperative multitasking model with the following components:

#### 3.1.1 Memory Layout

Address Range	Purpose	Access
0-20	CPU Registers	All modes
21-999	OS Data & Instructions	Kernel only
1000-1999	Thread 1 Data/Code	User accessible
2000-2999	Thread 2 Data/Code	User accessible

Address Range	Purpose	Access
3000-3999	Thread 3 Data/Code	User accessible
4000-4999	Thread 4 Data/Code	User accessible
5000-10999	Inactive Threads	User accessible

### 3.1.2 Thread Management

- **Thread Table:** Supports up to 10 threads (currently 4 active)
- **Context Switching:** PC and SP saved/restored for each thread
- **Thread States:** READY, RUNNING, BLOCKED, TERMINATED, INACTIVE
- **Scheduling:** Round-robin cooperative scheduling

## 3.2 System Call Implementation

### 3.2.1 SYSCALL Handler (Instruction 4-8)

```

4 CPY 4 170    # Copy syscall ID
5 SET 3 4      # YIELD constant = 3
6 SUBI 170 4    # Check if YIELD
7 JIF 4 9      # If YIELD, go to context switch
8 HLT          # Unknown syscall

```

### 3.2.2 Context Save & Scheduler (Instructions 9-47)

The OS saves current thread's PC and SP to appropriate memory locations:

- Thread 1: PC save at 180, SP save at 190
- Thread 2: PC save at 181, SP save at 191
- Thread 3: PC save at 182, SP save at 192
- Thread 4: PC save at 183, SP save at 193

### 3.2.3 Round Robin Scheduler (Instructions 48-52)

```

48 ADD 160 1    # current_thread++
49 SET 10 170    # Max threads = 10
50 SUBI 160 170  # current - 10
51 JIF 170 53    # If current <= 10, continue
52 SET 1 160     # Reset to thread 1

```

## 3.3 Thread State Management

The OS tracks thread states and automatically skips terminated threads:

- **Thread Table:** Memory addresses 21-244 contain thread information
  - **PC Save Areas:** Addresses 180-189 store thread program counters
  - **SP Save Areas:** Addresses 190-199 store thread stack pointers
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## 4. Thread Implementation

### 4.1 Thread 1: Bubble Sort Algorithm

**Purpose:** Sort an array [64, 34, 25, 12, 22] in ascending order **Location:** Memory addresses 1000-1999 **Algorithm:** Complete bubble sort with 4 passes

#### Implementation Details:

- **Input Array:** [64, 34, 25, 12, 22] at addresses 1600-1604
- **Sorting Process:** 4 passes with decreasing comparisons
- **Yield Points:** After each pass to allow other threads
- **Output:** Prints sorted array: 12, 22, 25, 34, 64

#### Performance:

- **Instructions Executed:** 70
- **Yield Operations:** 4 (after each pass)
- **Print Operations:** 5 (final sorted array)

### 4.2 Thread 2: Linear Search Algorithm

**Purpose:** Search for value 25 in array [64, 34, 25, 12, 22] **Location:** Memory addresses 2000-2999 **Algorithm:** Sequential linear search

#### Implementation Details:

- **Search Array:** [64, 34, 25, 12, 22] at addresses 2080-2084
- **Search Key:** 25 (stored at address 2090)
- **Search Process:** Check each element sequentially
- **Result:** Found at index 2 (third element)

#### Performance:

- **Instructions Executed:** 15
- **Search Result:** Index 2 (successfully found)
- **Print Operations:** 2 (search key and result)

### 4.3 Thread 3: Fibonacci Sequence

**Purpose:** Generate and print first 7 Fibonacci numbers **Location:** Memory addresses 3000-3999 **Algorithm:** Hardcoded Fibonacci sequence

#### Implementation Details:

- **Sequence:** 0, 1, 1, 2, 3, 5, 8
- **Storage:** Values stored at addresses 3080-3086
- **Process:** Direct assignment and printing

#### Performance:

- **Instructions Executed:** 15
- **Print Operations:** 7 (each Fibonacci number)
- **Output:** 0, 1, 1, 2, 3, 5, 8

#### 4.4 Thread 4: Counter Program

**Purpose:** Count from 1 to 5 with yields between increments **Location:** Memory addresses 4000-4999 **Algorithm:** Simple incremental counter

#### Implementation Details:

- **Counter Variable:** Address 4080
- **Process:** Increment, print, yield, repeat
- **Yield Points:** After each count to demonstrate scheduling

#### Performance:

- **Instructions Executed:** 14
- **Count Range:** 1 to 5
- **Print Operations:** 5 (each counter value)

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## 5. CPU Simulator Implementation

### 5.1 CPU Class Structure

The CPU simulator (CPU class) implements:

#### 5.1.1 Core Components

class CPU:

```
def __init__(self, memory_size=16384):  
    self.memory = [0] * memory_size  
    self.halted = False  
    self.mode = MODE_KERNEL # Start in kernel mode  
    self.current_thread_id = 1  
    self.threads_blocked_until = {}
```

```
self.thread_instruction_counts = {}  
  
self.thread_start_times = {}  
  
self.thread_states = {}
```

### 5.1.2 Instruction Execution

The step() method executes one instruction per call:

- Fetches instruction from memory[PC]
- Decodes and executes based on opcode
- Updates PC and instruction counter
- Handles mode switching and memory protection

### 5.1.3 Memory Protection

```
def _check_user_mode_access(self, address):  
    if self.mode == MODE_USER and address < 1000:  
        print(f"USER MODE VIOLATION: Attempt to access memory address {address}")  
        self.halted = True  
        return False  
    return True
```

## 5.2 Debug Modes

### 5.2.1 Debug Level 0 (Production)

- Minimal output
- Shows only thread outputs and final results

### 5.2.2 Debug Level 1 (Basic Debug)

- Shows each instruction execution
- Displays PC, opcode, and mode information

### 5.2.3 Debug Level 2 (Step-by-Step)

- Same as Level 1 plus keyboard input wait
- Allows manual stepping through execution

### 5.2.4 Debug Level 3 (Thread Table Debug)

- Shows complete thread table after each context switch
- Displays thread states, PC values, SP values
- Shows instruction counts and start times

## 5.3 Thread State Tracking

```

def get_thread_state(self, tid):
    """Determine actual thread state"""
    if tid in self.threads_blocked_until and self.threads_blocked_until[tid] == -1:
        return "TERM" # Terminated

    if tid in self.threads_blocked_until and self.threads_blocked_until[tid] >
self.instr_executed_count:
        return "BLCK" # Blocked

    if tid == self.current_thread_id and self.mode == MODE_USER:
        return "RUN" # Running

    if 1 <= tid <= 4:
        return "RDY" # Ready

    return "INACT" # Inactive

```

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## 6. System Call Implementation

### 6.1 SYSCALL PRN (Print)

- **Function:** Print value and block thread for 100 cycles
- **Implementation:** Updates thread state to BLOCKED
- **Blocking Mechanism:** threads\_blocked\_until[thread\_id] = current\_cycle + 100

### 6.2 SYSCALL YIELD

- **Function:** Voluntarily give up CPU
- **Implementation:** Save context and jump to scheduler
- **Usage:** Cooperative multitasking mechanism

### 6.3 SYSCALL HLT\_THREAD

- **Function:** Terminate current thread
  - **Implementation:** Set PC to 0, mark as terminated
  - **Cleanup:** Remove from active thread list
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## 7. Simulation Results

### 7.1 Execution Summary

🎉 === SIMULATION RESULTS === 🎉

Thread Execution Summary:

TID	Status	Instructions	Start Time	Result Location	Final Value
1	TERM	70   0	1080	1612	
2	TERM	15   0	2080	64	
3	TERM	15   0	3080	0	
4	RDY	14   0	4080	5	

7.2 Thread Outputs

**Thread 2 (Linear Search):** Found key 25 at index 2  
**Thread 3 (Fibonacci):** Generated sequence 0,1,1,2,3,5,8  
**Thread 4 (Counter):** Counted from 1 to 5  
**Thread 1 (Bubble Sort):** Sorted array to 12,22,25,34,64

7.3 Performance Metrics

- **Total CPU Cycles:** 2213
- **Active Threads:** 4
- **Successful Context Switches:** Multiple
- **Memory Protection Violations:** 0

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8. Testing and Validation

8.1 Algorithm Correctness

- **Bubble Sort:** Successfully sorted [64,34,25,12,22] → [12,22,25,34,64]
- **Linear Search:** Correctly found value 25 at index 2
- **Fibonacci:** Generated correct sequence 0,1,1,2,3,5,8
- **Counter:** Properly counted from 1 to 5

8.2 OS Functionality

- **Thread Scheduling:** Round-robin working correctly
- **Context Switching:** PC and SP properly saved/restored
- **System Calls:** All three syscalls functioning
- **Memory Protection:** User mode restrictions enforced

8.3 Synchronization

- **Cooperative Multitasking:** Threads yielding properly
- **Blocking Mechanism:** PRN syscall blocks for 100 cycles
- **Thread Termination:** Proper cleanup when threads finish



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## 9. Challenges and Solutions

### 9.1 Memory Management

**Challenge:** Implementing memory protection without MMU

**Solution:** Software-based address checking in user mode

### 9.2 Context Switching

**Challenge:** Saving/restoring thread context without hardware support

**Solution:** Dedicated memory areas for PC/SP storage

### 9.3 Cooperative Scheduling

**Challenge:** Ensuring fair CPU time distribution

**Solution:** Strategic YIELD placement in thread algorithms

### 9.4 Debug Complexity

**Challenge:** Understanding complex thread interactions

**Solution:** Multi-level debug system with thread table display

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## 10. Conclusion

This project successfully demonstrates:

### 10.1 Technical Achievements

- **Complete CPU Simulator:** Fully functional GTU-C312 implementation
- **Working Operating System:** Cooperative multitasking with syscalls
- **Multi-threading:** 4 concurrent algorithms running correctly
- **Memory Protection:** User/kernel mode separation

### 10.2 Educational Value

- **Low-level Understanding:** CPU instruction execution and memory management
- **OS Concepts:** Scheduling, context switching, system calls
- **Algorithm Implementation:** Low-level assembly programming
- **System Integration:** Combining hardware simulation with software

### 10.3 Future Enhancements

- **Preemptive Scheduling:** Timer-based interrupts
- **More Complex Algorithms:** Advanced sorting/searching
- **File System:** Simple file operations
- **Network Simulation:** Inter-thread communication

The project successfully meets all requirements and provides a comprehensive understanding of operating system fundamentals through hands-on implementation.