

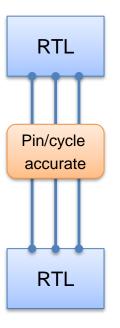
SystemC TLM

Stefano Spellini

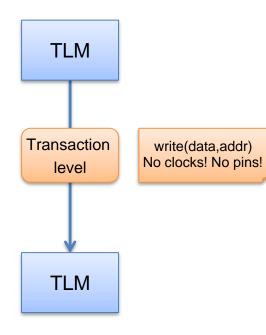


Transaction Level Modeling (TLM)

```
process(clock)
IF (clock'event and clock = '1')
THEN
CASE fsm_state IS:
           WHEN s0 \Rightarrow
             request_port <= '1';
            fsm state := s1;
           WHEN s1 =>
            IF (grant_port = '1')
              THEN
              fsm state := s2;
           WHEN s2 =>
            data port <= data;</pre>
             addr_port <= addr;</pre>
```



Simulate every event



100-10.000x faster simulation



Motivations for TLM

- Represents key architectural components of hardware platform
- Architectural exploration, performance modeling
- Software execution on virtual model of hardware platform
- Golden model for hardware functional verification
- Available before RTL
- Simulates much faster than RTL



MODELING AT TLM



TLM transaction

- Relies on the notion of *transaction*
 - Data transfer from a design module to another
 - Read/Write operation
 - Represented by a payload object
 - Exchanged with primitive calls
 - Contains both data and control information





TLM actors: initiators and targets

- Initiator
 - Initiates a transaction
 - Creates a transaction object
 - Calls target (or interconnect) methods to deliver it
- Target
 - Acts as final destination of a transaction
 - Elaborates the payload
- Master/Slave behavior typical of busses





TLM paths

- Forward path
 - Calling path from the initiator in the direction of a target
 - The Master calls the Slave
- Backward path
 - Calling path by which a target makes method call backs in the direction of the initiator
 - The Slave answers to the Master



ESD System Design Synthesis Verification Testing Power Communication

TLM Coding Styles

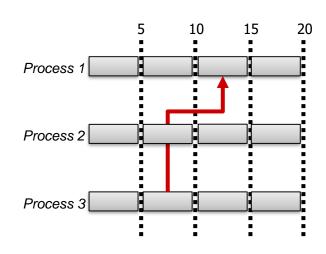
- Not levels of abstraction!
 - Defines how time and data are related together, according to designers' perception
 - More of a guideline, designers enjoy several degrees of freedom
- Specify how time and data are related together
 - Untimed
 - Blocking interface
 - Predefined synchronization points
 - Loosely timed
 - Blocking interface
 - Two synchronization points (invocation and return)
 - Temporal decoupling
 - Approximately timed
 - Non blocking transport interface
 - Timing annotation and multiple phases during the lifetime of a transaction
 - Beginning/end of request
 - Beginning/end of response

NOTION OF TIME



Temporal Decoupling

- Loosened Synchronization
 - Individual processes can run ahead for a certain time slice without advancing general simulation time
 - Delayed Synchronization with the system
 - Explicit Synchronization point
 - Very fast simulation





Generic Payload

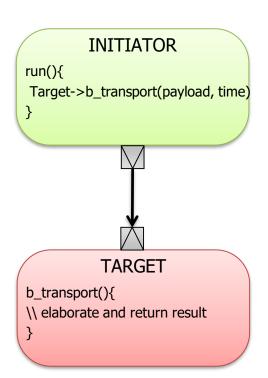
- Standard transaction objects
 - Modeled on memory-mapped bus
 - Command,
 - address,
 - data,
 - byte enables, ..
 - Transfered with function calls

Attribute	Туре
Command	Tlm_command
Address	Uint64
Data pointer	Unsigned char *
Data length	Unsigned int
Byte enable pointer	Unsigned char *
Byte enable length	Unsigned int
Streaming width	Unsigned int
DMI hint	Bool
Response status	Tlm_response_status
Extensions	Tlm_extension_base*



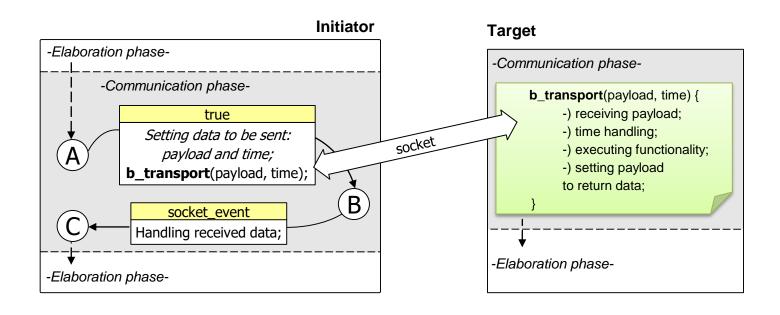
TLM blocking transport interface

- Support
 - Untimed
 - Loosely timed
- Initiator completes a transaction with the target in one function call
 - 2 synchronization points
 - Invocation
 - Return
- Uses only forward path





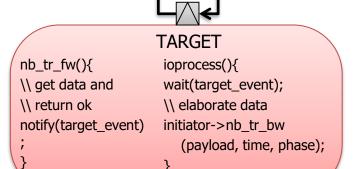
TLM blocking transport interface (Cont.d)



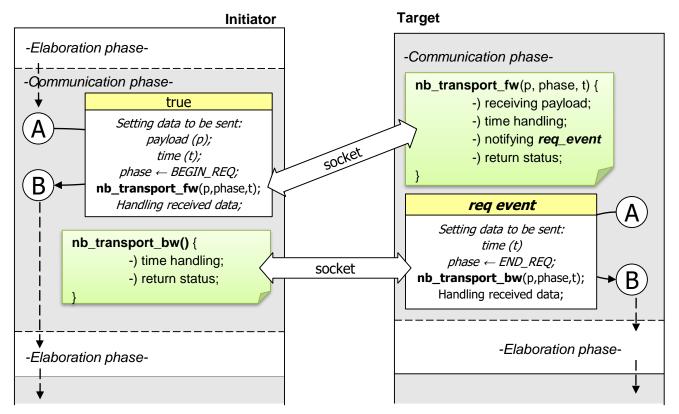


TLM non-blocking transport interface

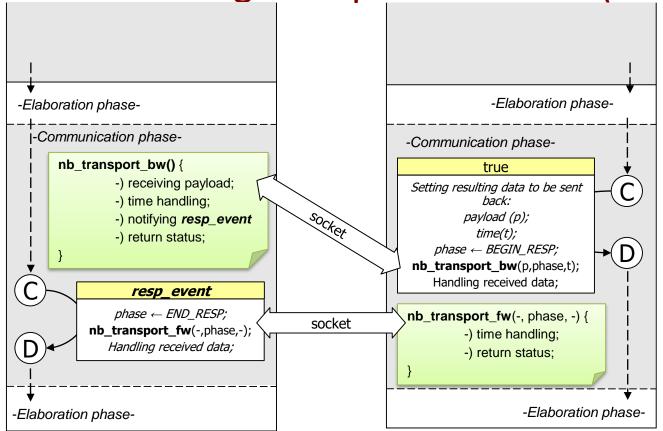
- Support
 - Approximately timed
- Detail a sequence of interactions between initiator and target
 - Phases
 - Begin/end request
 - Begin/end response
- Uses both forward path and backward path
- Classic protocol:
 - 4-phase handshaking protocol
 - AT4



TLM non-blocking transport interface (Cont.d)



TLM non-blocking transport interface (Cont.d)





SystemC Installation (ERRATA CORRIGE)

- Download SystemC
 - http://www.accellera.org/downloads/standards/systemc
 - From the elearning portal (SystemC 2.3.3 Source Code)
- Unzip the archive
 - \$> tar -xzvf systemc-2.3.3.tar.gz
- Configure the installation
 - \$> cd systemc-2.3.3
 - \$> mkdir build
 - \$> mkdir -p ~/pse_libraries/systemc
 - \$> cd build
 - \$> cmake -DCMAKE_INSTALL_PREFIX=/home/<your_user>/pse_libraries/systemc
 - - DBUILD_SHARED_LIBS=OFF ..
- Compile
 - \$> make -j
 - \$> make install



SystemC into action

- Uncompress the archive
 - \$ tar xzf 02_sources.tar.gz
 - \$ cd 02 sources/
 - \$ Is
 - dist div root root lesson root comparison
 - \$ cd root lesson
 - \$ Is
 - AT4 LT UT RTL
- Compile and run each TLM implementation, e.g.:
 - \$ cd UT/
 - \$ mkdir build && cd build
 - \$ cmake -DCMAKE_PREFIX_PATH=/home/<your_username>/pse_libraries/systemc ...
 - \$ make -j
 - \$./root_UT

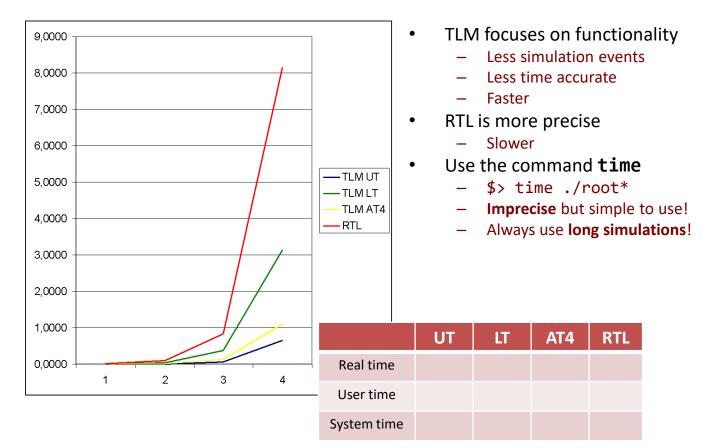


SystemC into action

- Go to the root_comparison folder
 - Root with 128000 interactions between target and initiator
- Compile all the TLM/RTL implementations, e.g.:
 - \$ cd RTL/
 - Ś mkdir build && cd build
 - \$ cmake -DCMAKE PREFIX PATH=/home/<your username>/pse libraries/systemc...
 - \$ make –j
 - \$./root_RTL
- Compare execution time with the time shell command
 - Real time
 - Elapsed time between invocation and termination
 - User time
 - Amount of time spent by the CPU performing actions for a program
 - System time
 - Amount of time spent by the CPU performing system calls on behalf of the program



SystemC into action





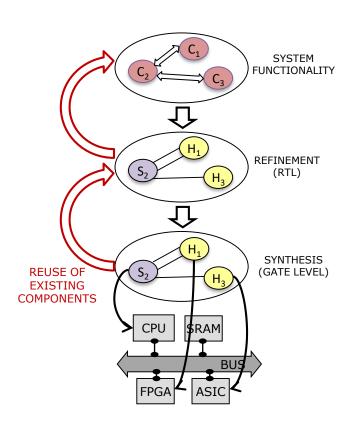
Connecting RTL and TLM

BASIC TRANSACTORS



Design

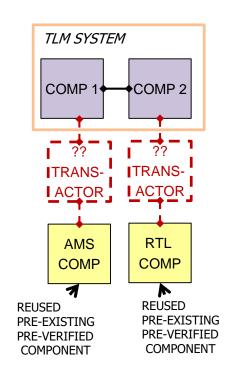
- Design flow
 - In theory: top-down approach
 - Choose and describe the functionality
 - Implement it with abstract communication (TLM)
 - Implement proper communication protocol (RTL)
 - Gate-level synthesis
 - In practice: reuse
 - Fundamental to meet time to market constraints and to save money (design+verification)





Mixed Modeling

- Mixed modeling
 - Implied by reuse
 - Integrate components developed at different levels of abstractions
 - Transactor
 - Translator in the communication between modules implemented at different levels of abstraction
 - Allows mixed modeling and testbench reuse

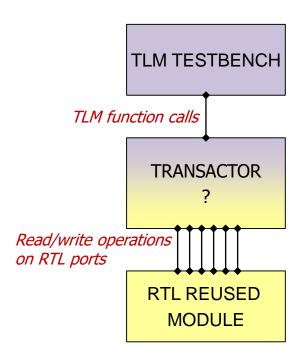




Transactor: main idea

Two APIs

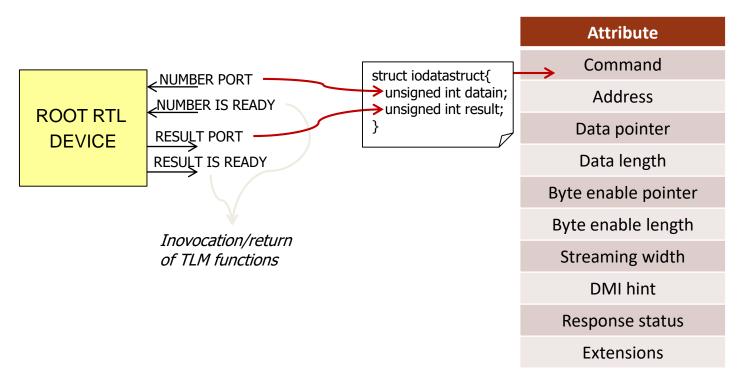
- Higher level of abstraction (TLM)
 - TLM standard functions
- Lower level of abstraction (RTL)
 - Sequence of read/write operations on RTL ports





Transactor: mapping

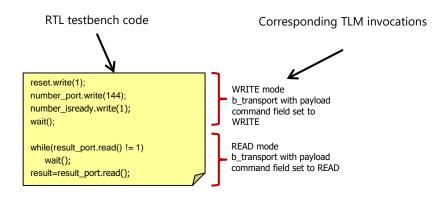
Define a mapping between RTL ports and TLM payload fields





Transactor: communication

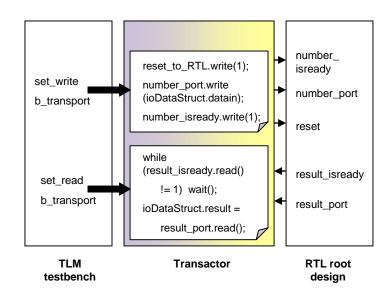
- Translation of TLM function calls to sequences of RTL signal operations
 - From blocking/non blocking primitives to cycle accurate temporization and handshaking
 - Root:
 - Two modes, read and write
 - Correspond to b_transport invocations with different command parameters





Transactor: structure

- Transactor structure
 - B_transport interface towards
 TLM
 - Depending on the command
 - Activates the WRITE process to set number_port and number isready port
 - Activates the READ process to get the value of the result_port and result_isready port





Transactor: Example

```
void root_RTL_transactor::WRITEPROCESS(){
                                                              while (true) {
                                                               wait(begin write);
void root RTL transactor::b transport
    (tlm::tlm generic payload& trans, sc time& t){
                                                               reset_to_RTL.write(1);
tlm::tlm command cmd = trans.get command();
                                                               p_Out_data.write(ioDataStruct.datain);
switch (cmd) {
                                                               p Out enable.write(1);
 case tlm::TLM WRITE COMMAND:
  ioDataStruct = *((iostruct*) trans.get data ptr());
                                                               end write.notify();
  begin write.notify();
                                                               wait();
  wait(end write);
  break;
  case tlm::TLM READ COMMAND:
  ioDataStruct = *((iostruct*) trans.get data ptr())
                                                             void root RTL transactor::READPROCESS(){
  begin read.notify();
                                                              while (true) {
  wait(end read);
                                                               wait(begin read);
  break;
 default:
                                                               while(p In enable.read() != 1)
  break:
                                                               wait();
                                                               ioDataStruct.result=p_In_data.read();
                                                               end read.notify();
                                                             }}
```



LECTURE ASSIGNMENTS



Fixed-Point Rational Binary Multiplication

- Algorithm for Binary Multiplication
- Requirements:
 - Operands: 16 bit Unsigned
 Fixed-Point rational value
 - 8 bit integer
 - 8 bit fraction
 - Forbidden use of sc fixed
 - Pay attention to the span of the **output!**
 - Max port width: 8 bit
 - No Limit on internal signals
 - Handshake protocol
- Well known algorithm, many examples on the web
 - Search for "Binary Multiplication" on Google!



Example with integers

6*9=54

```
Initialization:
a = 0110; // 6
b = 1001; // 9
temp = 00000110; // 6
product = 00000000;
Cicle 0:
b[0] is 1, thus:
product = 0 + 0110 = 00000110; // 6
temp = 00001100; // 12
Cicle 1:
b[1] is 0, thus:
temp = 00011000; // 24
Cicle 2:
b[2] = 0, thus:
temp = 00110000; // 48
Cicle 3:
b[3] = 1, thus:
product = 00000110 + 00110000 = 00110110; // 54
temp = 01100000; // 96
result = product; // 54
```



Assignments

- Implement the assigned design at RTL using SystemC
 - Remember: define the EFSM and FSMD before implementing the design!
- Implement the assigned design at TLM
 - Using the coding style of choice
 - Use the root examples as a guideline
- In the Report:
 - First page: FSMD and EFSM (structure)
 - Second page: Brief description of TLM implementation
 - Motivate the choice of the coding style
 - Compare it to the RTL in terms of simulation times



Assignments

- ONE directory compressed in ONE archive (tar.gz)
 - VRXXXXXX_Name_Surname
 - VRXXXXXX Name Surname.tar.gz
 - tar czf VRXXXXXX_Name_Surname.tar.gz VRXXXXXX_Name_Surname
- Within the directory: 2 sub-directory
 - Report:
 - Multiple PDFs: one PDF file per lecture
 - VRXXXXXX_Name_Surname_YY.pdf
 - YY = 01, 02, 03, ..., 10 etc... No 1, 2, 3
 - Solutions:
 - 1 Directory per report!



Example of tree

```
VR123456 Stefano Spellini/
      Reports/
            VR123456_Stefano_Spellini_01.pdf
            VR123456_Stefano_Spellini_02.pdf
            VR123456_Stefano_Spellini_03.pdf
      Solutions/
            01/
                   Empty
            02/
                   Mult
                          RTL/
                                  include, CMakeLists.txt, src
                          TLM *CodingStyle*/
                                                                  E.g. (TLM_UT)/
                                 include, CMakeLists.txt, src
            03/
```

Everything within VR123456_Stefano_Spellini.tar.gz



ASSIGNMENTS NOT PROPERLY DELIVERED WILL NOT BE CORRECTED, AND THE REPORT WILL NOT BE CONSIDERED!

Consegne che non rispettano le regole descritte sopra <u>non verranno corrette</u> e il report verrá considerato <u>non consegnato</u>!