



# **Embedded Operating System RTOS static priority servers**

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# Introduction

- In most real-time applications there are
  - Both periodic and aperiodic tasks
    - Typically periodic tasks are time-driven, hard real-time
    - Typically aperiodic tasks are event-driven, soft or hard RT
- Objectives:
  - Guarantee hard RT tasks
  - Provide good average response time for soft RT tasks

# Handling periodic and aperiodic tasks

Immediate service

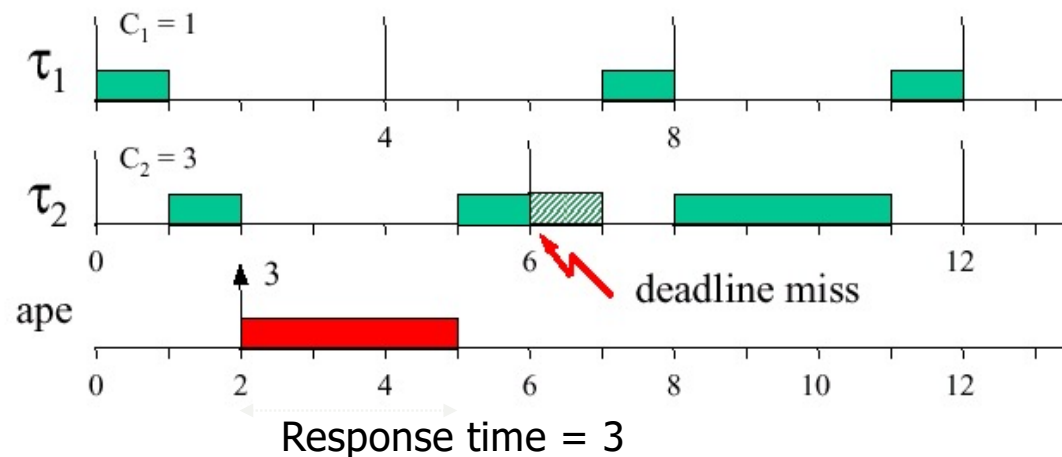
Background scheduling

Aperiodic servers

- Static priority servers
- Dynamic priority servers

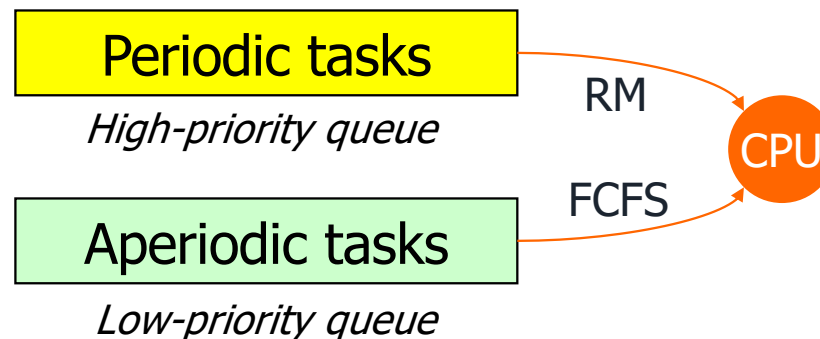
# Immediate service

- Aperiodic requests are served as soon as they arrive in the system
- Minimum response times for aperiodic requests
- Low guarantee of periodic tasks



# Background scheduling

- Soft aperiodic tasks in the background behind periodic tasks, that is, in the processor time left after scheduling all periodic tasks
- Aperiodic tasks get assigned a priority lower than any periodic one



# Background scheduling – Example

|          | $a_i$ | $C_i$ | $T_i$ |
|----------|-------|-------|-------|
| $\tau_1$ | 0     | 2     | 6     |
| $\tau_2$ | 0     | 4     | 10    |
| $ape_1$  | 2     | 1     | -     |
| $ape_2$  | 12    | 2     | -     |

RM for periodic tasks

# Background scheduling

Utilization factor under RM  $< 1$

➡ some processor time is left, it can be used for aperiodic tasks

High periodic load

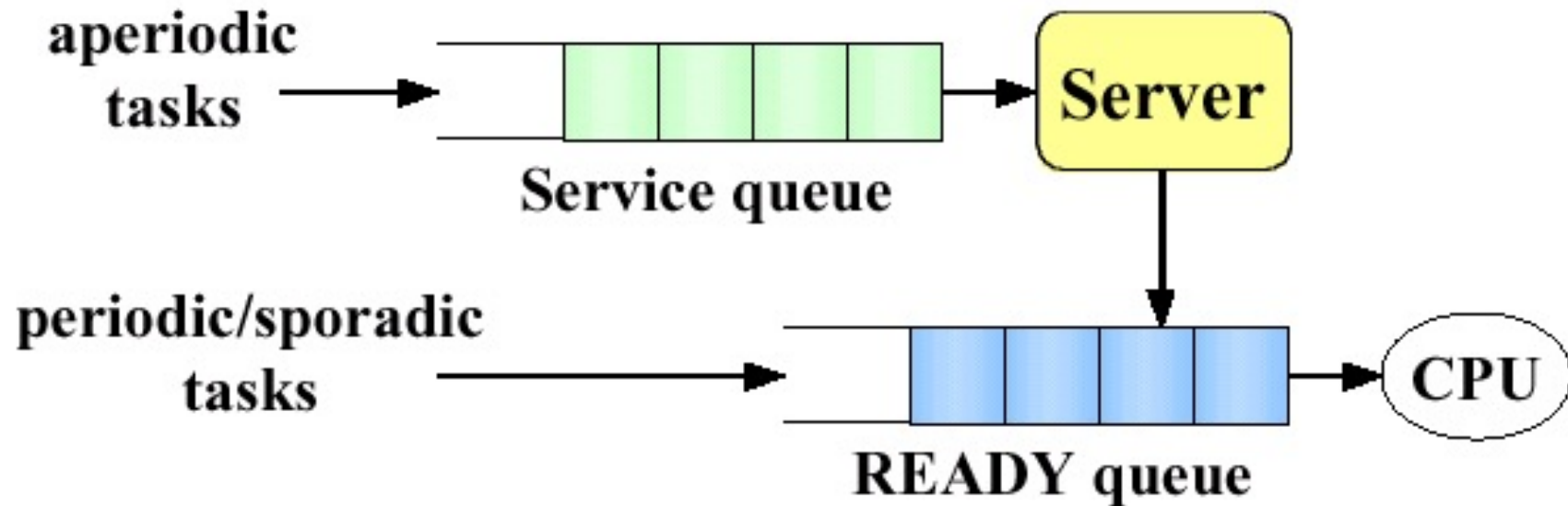
➡ bad response time for aperiodic tasks

Applicable only if no stringent timing requirements for aperiodic tasks

Major advantage: simplicity

# Priority servers

- To achieve more predictable aperiodic task handling
  - A specific periodic task (server) services aperiodic requests
  - The server is assigned
    - a period  $T_s$
    - a computation time  $C_s$  (capacity of the server)
  - The server is scheduled like any other periodic task





# Priority servers

Classified according to the priority scheme of the periodic scheduler

## Static priority servers

- Polling Server
- Deferrable server
- Priority exchange
- Sporadic server
- Slack stealing

## Dynamic priority servers

- Dynamic Polling Server
- Dynamic Deferrable Server
- Dynamic Sporadic Server
- Total Bandwidth Server
- Constant Bandwidth Server

# Polling Server (PS)

|                                 |   |   |
|---------------------------------|---|---|
| At the beginning of its period  | PS is (re)-charged at its full value $C_s$              |   |
|                                 | PS ready to serve any pending aperiodic requests        | within the limits of its capacity $C_s$   |
| If no aperiodic request pending | PS “suspends” itself until beginning of its next period | $C_s$ is discharged to 0  |
|                                 |   | Processor time is used for periodic tasks   |
|                                 |   | If aperiodic task arrives after suspension of PS, it is served in the next period |
| If aperiodic requests pending   | PS serves them until $C_s > 0$                          |   |

# Polling server - example

|          | $C_i$ | $T_i$ |
|----------|-------|-------|
| $\tau_1$ | 1     | 4     |
| $\tau_2$ | 2     | 6     |

| Server    |
|-----------|
| $C_s = 2$ |
| $T_s = 5$ |

|      | $C_i$ | $a_i$ |
|------|-------|-------|
| ape1 | 2     | 2     |
| ape2 | 1     | 8     |
| ape3 | 2     | 12    |
| ape4 | 1     | 19    |

# Polling server analysis

- In the worst-case, PS behaves as a periodic task with  $U_s = C_s/T_s$

$$U_p + U_s \leq U_{lub}(n+1) \quad \sum_{i=1}^n \frac{C_i}{T_i} + \frac{C_s}{T_s} \leq (n+1)[2^{1/(n+1)} - 1]$$

- More precise schedulability analysis

$$U_{lub}^{RM+PS}(n) = U_s + n \left[ \left( \frac{2}{U_s + 1} \right)^{1/n} - 1 \right] \Rightarrow U_p \leq n \left[ \left( \frac{2}{U_s + 1} \right)^{1/n} - 1 \right]$$

# Deferrable Server

Like PS but

DS preserves its capacity if no requests are pending at invocation of the server

Capacity is maintained until server period

Aperiodic requests arriving at any time are served as long as the capacity has not been exhausted

At the beginning of any server period

Capacity is replenished at its full value

But no cumulation

# Deferrable server – Example

|          | $C_i$ | $T_i$ |
|----------|-------|-------|
| $\tau_1$ | 1     | 4     |
| $\tau_2$ | 2     | 6     |

| Server    |
|-----------|
| $C_s = 2$ |
| $T_s = 5$ |

|      | $C_i$ | $a_i$ |
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| ape1 | 2     | 2     |
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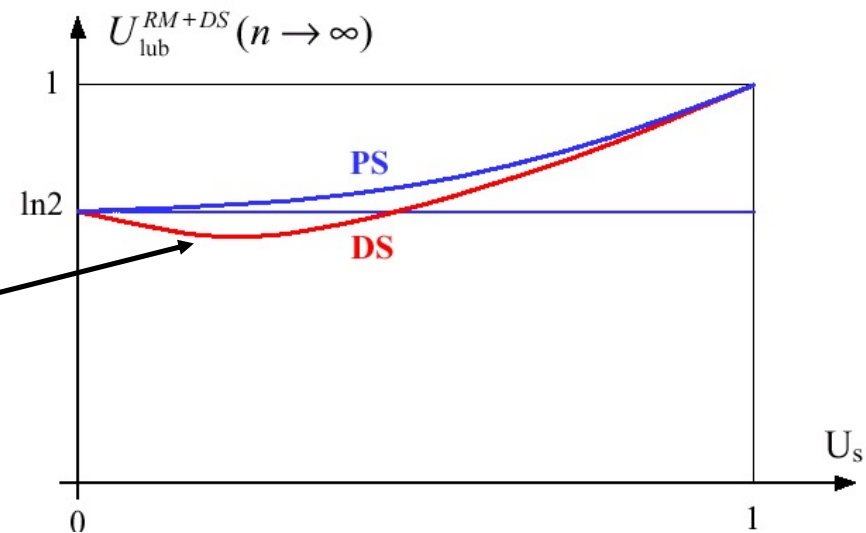
# Priority server – PS vs. DS

- Utilization

$$U_{\text{lub}}^{RM+DS}(n) = U_s + n \left[ \left( \frac{U_s + 2}{2U_s + 1} \right)^{\frac{1}{n}} - 1 \right]$$

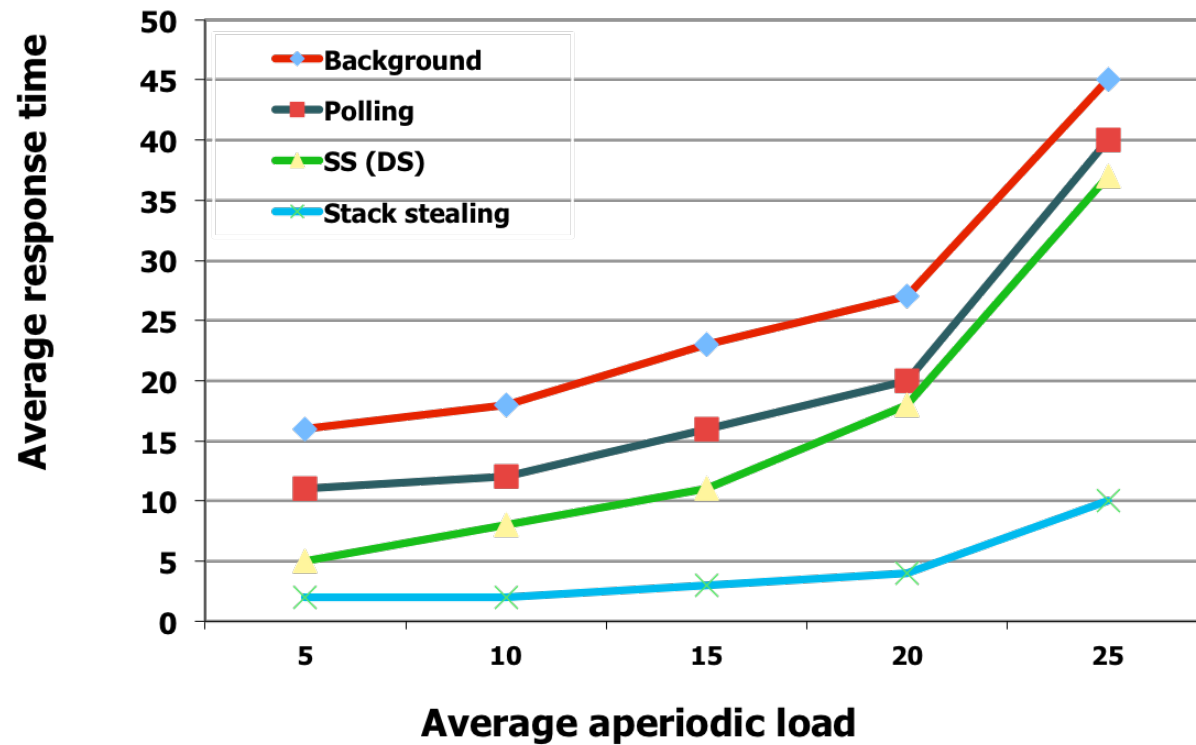
- Comparing PS and DS

Keeping the budget improves responsiveness, but decreases the utilization bound.



# Comparison of fixed priority servers (1)

( $U_p=0.69$ )





## Comparison of fixed priority servers (2)

|                   | Performance | Computational complexity | Memory requirement | Implementation complexity |
|-------------------|-------------|--------------------------|--------------------|---------------------------|
| Background server | <b>C</b>    | <b>A</b>                 | <b>A</b>           | <b>A</b>                  |
| Polling Server    | <b>C</b>    | <b>A</b>                 | <b>A</b>           | <b>A</b>                  |
| Deferrable Server | <b>B</b>    | <b>A</b>                 | <b>A</b>           | <b>A</b>                  |
| Priority Exchange | <b>B</b>    | <b>B</b>                 | <b>B</b>           | <b>B</b>                  |
| Slack Stealing    | <b>A</b>    | <b>C</b>                 | <b>C</b>           | <b>C</b>                  |

A=excellent B=good C=poor